Unified and Optimized Linear Collision Attacks and Their Application in a Non-Profiled Setting

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Power Analysis Attacks

Divide & Conquer

- Differential Power Analysis
- Correlation Power Analysis
- Template Attack
- ▶ ...

Alternatives

- Algebraic side-channel attacks
- Side-channel collision attacks



Motivations

- Getting rid of the leakage model
- Main idea

same output \Rightarrow same leakage.

leakage model choice \rightarrow similarity metric choice

- [Schramm et al. '03] collision in the f function of DES
- [Bogdanov '07] collision between S-box computations
 - ► Software: table implementation of an S-boxes
 - \blacktriangleright Hardware: high area constraints \rightarrow S-box reuse



This work

Target: linear collision attacks

- 1. Enhancing collision attacks
 - Handling errors generically
 - Exploiting non-colliding events
- 2. Collision-attack relevance
 - Comparison with unprofiled attacks
 - Software context



Overview

Linear Collision Attacks

Linear Collision and Coding Theory

Experiments



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Linear Collision Attacks Principle





Linear Collision Attacks Principle



$X_2 \oplus K_2 = X_{11} \oplus K_{11} \quad \Longrightarrow \quad K_2 \oplus K_{11} = X_2 \oplus X_{11}$



Linear Collision Attacks Principle





$$\begin{cases}
K_1 \oplus K_5 &= \Delta K_{1,5} \\
K_1 \oplus K_2 &= \Delta K_{1,2} \\
K_2 \oplus K_8 &= \Delta K_{2,8} \\
K_3 \oplus K_4 &= \Delta K_{3,4} \\
K_6 \oplus K_7 &= \Delta K_{6,7}
\end{cases}$$



$$\begin{cases} K_1 \oplus K_5 &= \Delta K_{1,5} \\ K_1 \oplus K_2 &= \Delta K_{1,2} \\ K_2 \oplus K_8 &= \Delta K_{2,8} \\ K_3 \oplus K_4 &= \Delta K_{3,4} \\ K_6 \oplus K_7 &= \Delta K_{6,7} \end{cases}$$

$$\begin{array}{c} 2^{24} \text{ values for } (K_1, K_3, K_6) \\ & \downarrow \\ 2^{24} \text{ keys } (K_1, K_2, K_3, K_4, K_5, K_6, K_7, K_8) \\ & \text{ instead of } 2^{64} \end{array}$$



Limitations

- 1. Information available only if a collision occurs
 - Non-colliding event also brings information
- 2. Errors
 - Inconsistency in the system
 - Undetectable erroneous system
- Techniques to enhance the attack
 - [Bogdanov '08] binary and ternary voting
 - [Moradi et al. '10] determining $\Delta K_{a,b}$ using correlation



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Decoding Problem



Coding: adding redundancy to the message



$$\begin{array}{c} \mathsf{key} \\ (K_1, \dots, K_{16}) \end{array}$$



$$(K_1, \ldots, K_{16})$$

 $(\Delta K_{a,b})_{1 < a < b < 16}$

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 \Rightarrow why not using a decoding algorithm?



Contributions

- 1. General framework
- 2. Enhancement of current tools
 - LDPC soft decoding
 - Bayesian extensions
- 3. Experiments on software implementations































$$\Delta K \stackrel{\text{def}}{=} (\Delta K_{1,2}, \dots, \Delta K_{15,16})$$

- 120 $\Delta K's \rightarrow$ dimension-15 subspace
- Constraints involving 3 positions

$$\underbrace{\Delta K_{a,b}}_{\mathcal{K}_{a} \oplus K_{b}} \oplus \underbrace{\Delta K_{a,c}}_{\mathcal{K}_{a} \oplus K_{c}} = \underbrace{\Delta K_{b,c}}_{K_{b} \oplus K_{c}}$$



$$\Delta K \stackrel{\text{def}}{=} (\Delta K_{1,2}, \dots, \Delta K_{15,16})$$

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Sparse constraints



$$\Delta K \stackrel{\text{def}}{=} (\Delta K_{1,2}, \dots, \Delta K_{15,16})$$

- 120 $\Delta K's \rightarrow$ dimension-15 subspace
- Constraints involving 3 positions



$$\Delta K \stackrel{\text{def}}{=} (\Delta K_{1,2}, \dots, \Delta K_{15,16})$$

- 120 $\Delta K's \rightarrow$ dimension-15 subspace
- Constraints involving 3 positions





Bayesian extensions

$$\left. \begin{array}{c} \text{scores} \ (s_1, \dots, s_n) \\ \Pr\left[S|\text{coll}\right] \\ \Pr\left[S|\text{non-coll}\right] \end{array} \right\} \xrightarrow{Bayes} \Pr\left[\Delta K_{a,b} = \delta\right]$$

Unprofiled setting:

- Theoretical models
- On-line parameter estimation



Bayesian extensions

1. Euclidean distance (normalized distance)

$$\operatorname{NED}(T, T') = \sum_{j} \frac{\left(T_{j} - T'_{j}\right)^{2}}{2\sigma_{j}^{2}}$$

2. Correlation-enhanced (Fisher transform)

$$\operatorname{Fisher}(c) = \operatorname{arctanh}(c)$$

Overview

Experiments

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Reference implementation





Furious implementation





Attacking the Reference Implementation





Attacking the Reference Implementation





Attacking the Reference Implementation













Attacking the Furious Implementation





















Theoretical context

Non-linear leakage





Conclusions

Collision attacks as a decoding problem

- General framework
- Improvements of former attacks
 - Soft decoding algorithm
 - Bayesian extensions

Experiments performed

- Usually less efficient than stochastic DPA
- May be useful in challenging implementation contexts



Perspectives

- List decoding
- Application to masked implementations
- Application to non-linear collisions

