

# Beyond Birthday Bound Secure Fresh Rekeying: Application to Authenticated Encryption

Bart Mennink

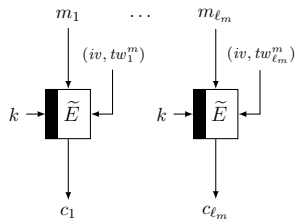
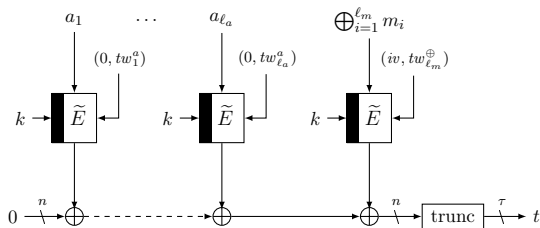
Radboud University (The Netherlands)



ASIACRYPT 2020  
December 7–11, 2020

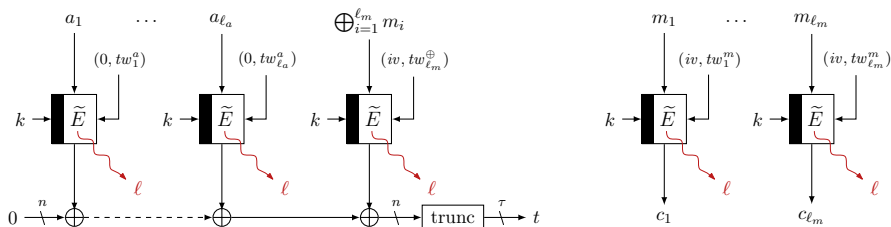
# Introduction

- Block cipher based constructions invoke the secret key many times
- Typical examples:
  - CTR and CBC encryption
  - OCBx, and its generalization  $\Theta$ CB [KR11] (depicted for integral data)



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Repeated evaluation of the key  $\longrightarrow$  repeated leakage of the key

# Countermeasures Against Leakage

## Implementation Protection



- Protection on top
- Masking or hiding

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## Rekeying



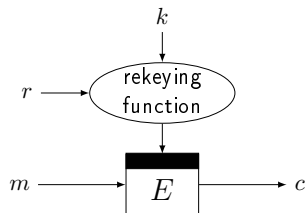
- Method-in-the-middle
- Leveled implementation

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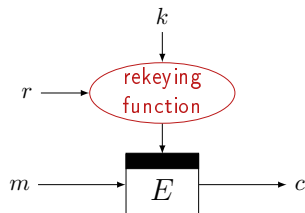
# Fresh Parallel Rekeying



## Idea

- Make scarce use of key material
- Strong protection only needed for cryptographically light building blocks

# Fresh Parallel Rekeying

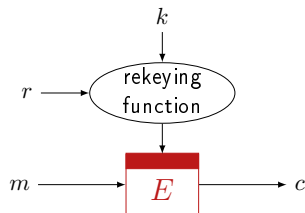


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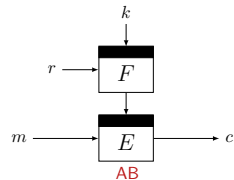


## Idea

- Make scarce use of key material
- Strong protection only needed for cryptographically light building blocks
- **Rekeying**: strong protection (e.g., against DPA), but not necessarily cryptographically strong
- **Core**: must be cryptographically strong, but only needs lighter protection (e.g., against SPA)

# Rekeying

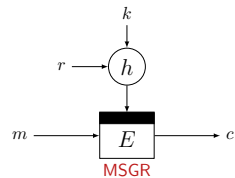
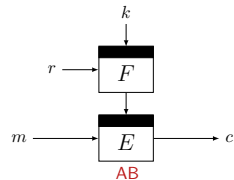
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# Rekeying

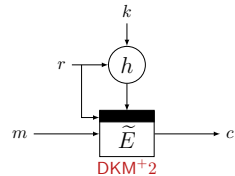
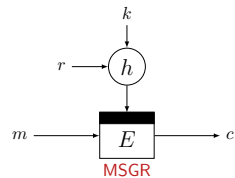
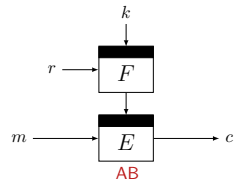
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formalized idea of rekeying

2010 • MSGR — Medwed et al.  
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# Rekeying

- 2000 • **AB** — Abdalla and Bellare formalized idea of rekeying
- 2010 • **MSGR** — Medwed et al. minimalized subkeying
- 2014 • **DKM<sup>+</sup>1/DKM<sup>+</sup>2** — Dobraunig et al. attack and 2 remedies for **MSGR**



## Rekeying Versus Tweakable Block Ciphers



- The idea of rekeying reminds a bit of **tweakable block ciphers**
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- Known as **tweak-rekeyable tweakable block ciphers**

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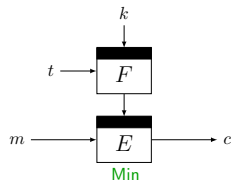
Minematsu — Min  
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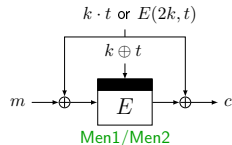
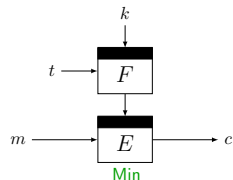
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**Mennink** — **Men1/Men2**  
beyond birthday bound security

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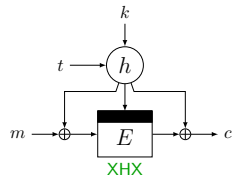
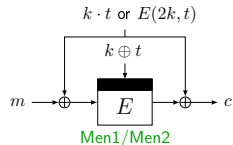
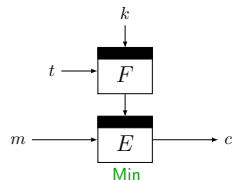
Mennink — Men1/Men2  
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Wang et al. — **WGZ<sup>+</sup><sub>i</sub>**  
generalization of Men<sub>i</sub>: 32 schemes

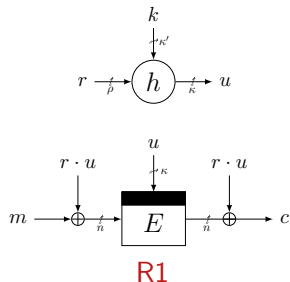
Naito — **XKX**  
targeting AE

Jha et al. — **XHX**  
generalized construction

# Tweak-Rekeyable TBC

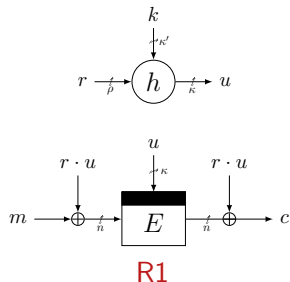


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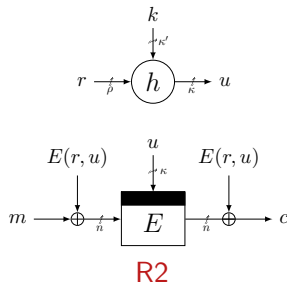


- Derivative of Men1
- $\kappa = \rho = n$
- $\approx 2n/3$ -bit security

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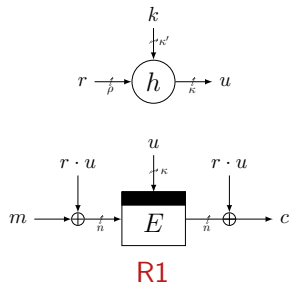


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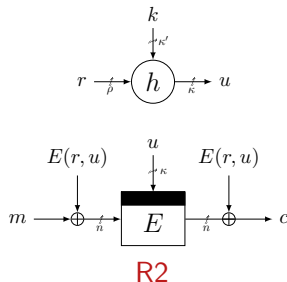


- Inspired by **WGZ<sup>+</sup>12**
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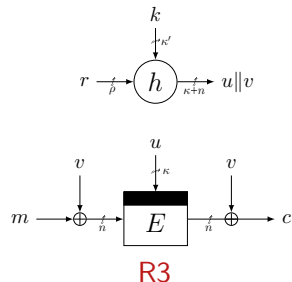
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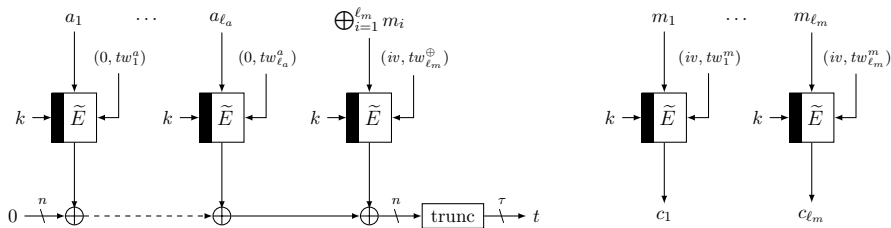
- Adaptation of **XHX**
- $\approx (\kappa + n)/2$ -bit security
- Covers  $p$ -based rekeying

## Cost Comparison

- Keep  $\kappa = \rho = n$  for simplicity
- $F$  is RF,  $\tilde{E}$  is TBC,  $E$  is BC
- For sake of counting: consider  $n$ -bit finite field multiplication for  $h$

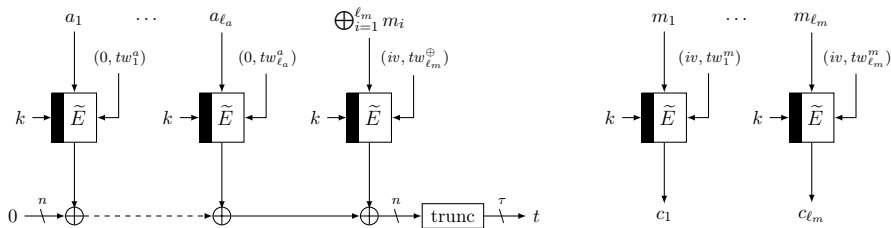
scheme	subkey		core			keysize	security
	$F$	$\otimes/h$	$\tilde{E}$	$E$	$\otimes$		
AB (2000)	1	0	0	1	0	$n$	$2^n$ (as PRF)
MSGR (2010)	0	1	0	1	0	$n$	$2^{n/2}$
DKM <sup>+</sup> 1 (2014)	0	1	0	2	0	$n$	$2^{n/2}$
DKM <sup>+</sup> 2 (2014)	0	1	1	0	0	$n$	$2^n$
<b>R1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>1</b>	<b><math>n</math></b>	<b><math>2^{2n/3}</math></b>
<b>R2</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>2</b>	<b>0</b>	<b><math>n</math></b>	<b><math>2^n</math></b>
<b>R3</b>	<b>0</b>	<b>2</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b><math>2n</math></b>	<b><math>2^n</math></b>

# Application to Authenticated Encryption



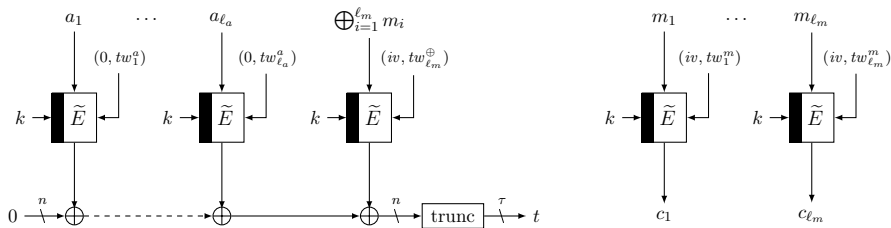
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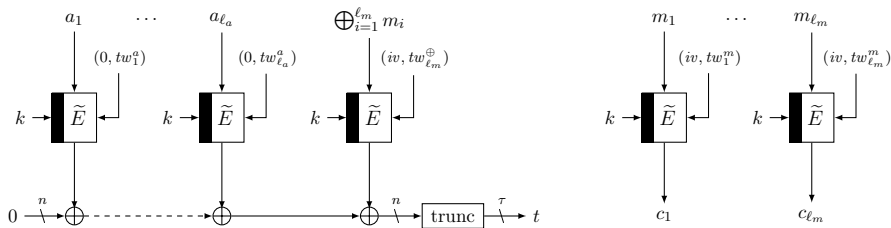
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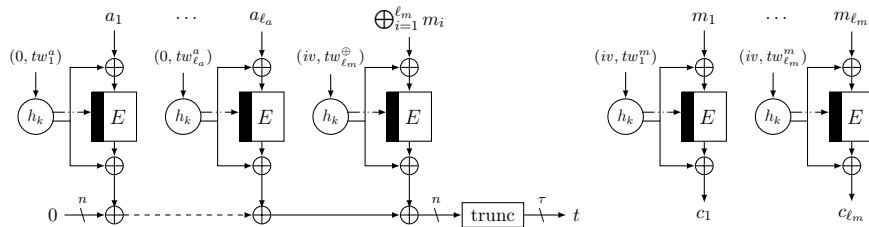


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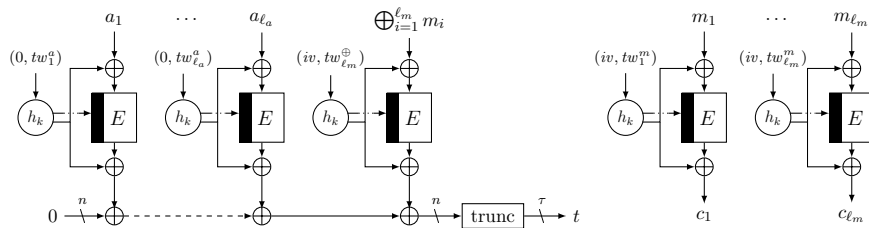
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  - Deoxys-I [JNPS16]  $\equiv \Theta$ CB instantiated with **Deoxys-BC** design
  - One can just as well instantiate it with **a rekeying scheme**

## Application to Authenticated Encryption: $\Theta$ CB-R3



- Instantiation of  $\Theta$ CB with **R3**

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- Instantiation of  $\Theta$ CB with **R3**
- Features:
  - $n$ -bit security (in ideal model)
  - $\ell_a + \ell_m + 1$  lightly protected  $E$  calls
  - $2(\ell_a + \ell_m + 1)$  strongly protected  $h$  calls
  - By design easier to protect against side-channel attacks

# Application to Authenticated Encryption: Comparison

## $\Theta$ CB-R3

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## Comparison to OCB3 [KR11]

- $n/2$ -bit security (but in standard model)
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## Comparison to DTE [BKP+17]

- Different goal: nonce-misuse resistance
- $\ell_a + \ell_m + 1$  unprotected  $E$  calls (approx., for hashing)
- $2\ell_m$  lightly protected  $E$  calls
- 2 strongly protected  $E$  calls

# Conclusion

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**Thank you for your attention!**