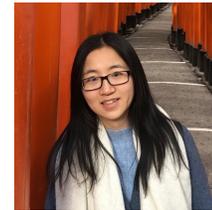


# Private Set Intersection in the Internet Setting From Lightweight Oblivious PRF



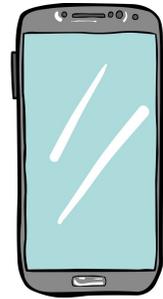
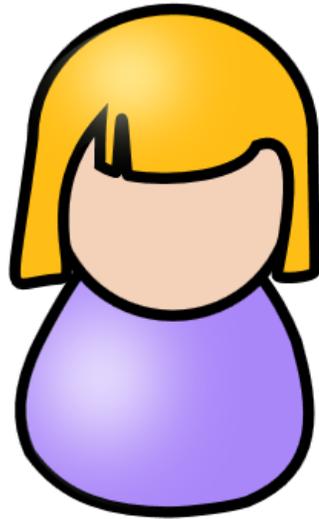
Melissa Chase



Peihan Miao

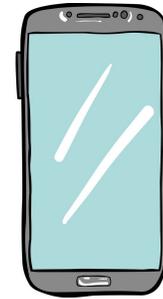


What is Private Set Intersection?

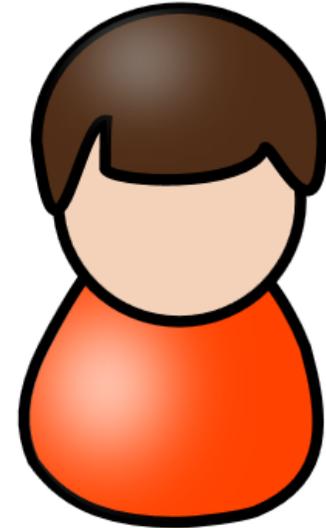


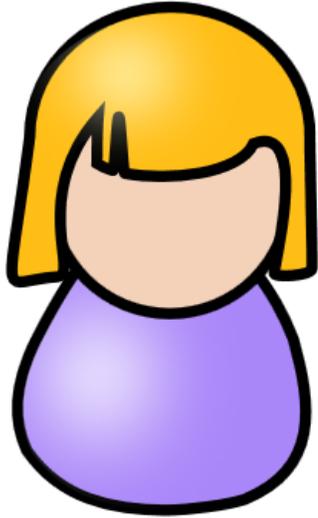
Charlie
Peihan
David
Eve
Melissa

Common friends?

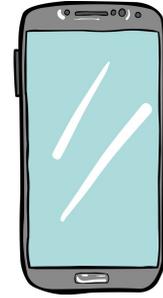


Eric
Frank
Melissa
Peihan
Grace

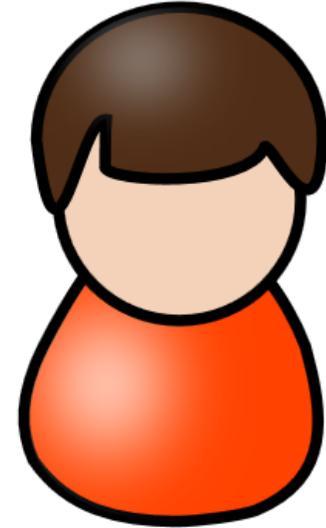


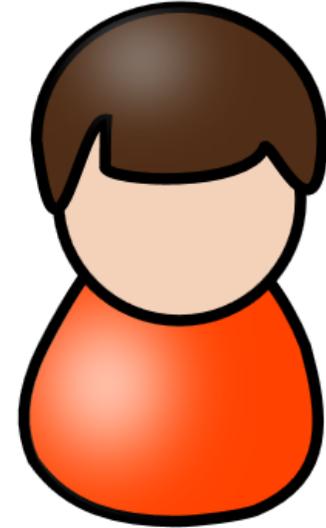
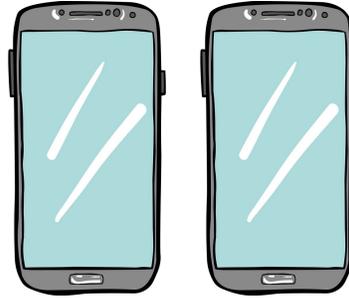
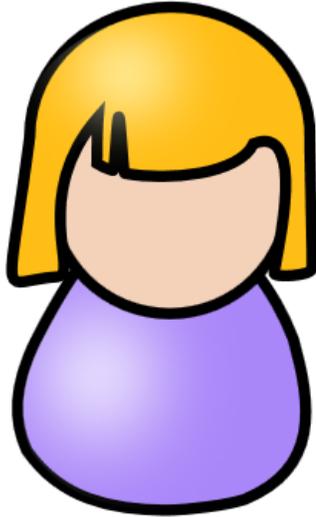


Charlie
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Eve
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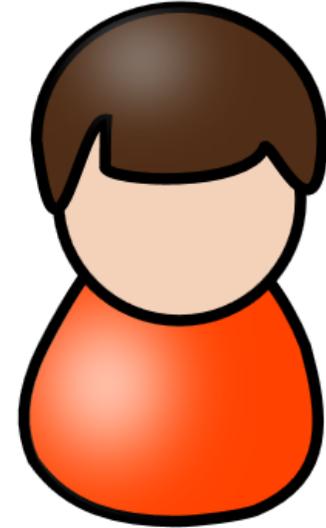
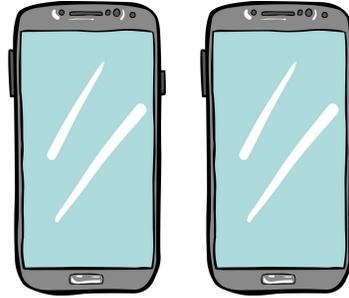
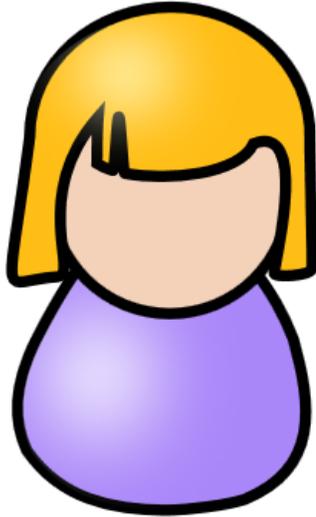


Eric
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Melissa
Peihan
Grace

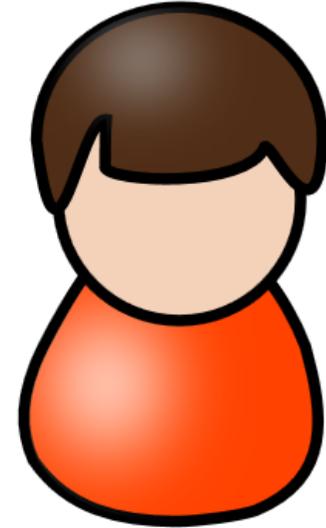
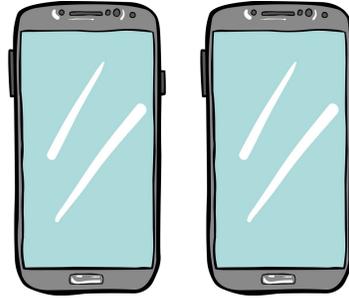
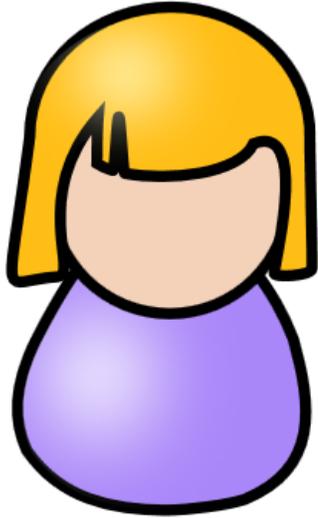




Charlie	Eric
Peihan	Frank
David	Melissa
Eve	Peihan
Melissa	Grace



Charlie	Eric
Peihan	Frank
David	Melissa
Eve	Peihan
Melissa	Grace



Charlie	Eric
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David	Melissa
Eve	Peihan
Melissa	Grace

Eurocrypt 2020

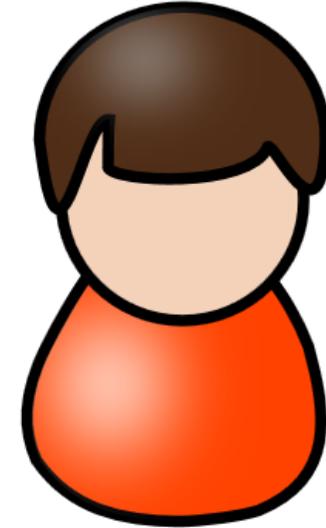
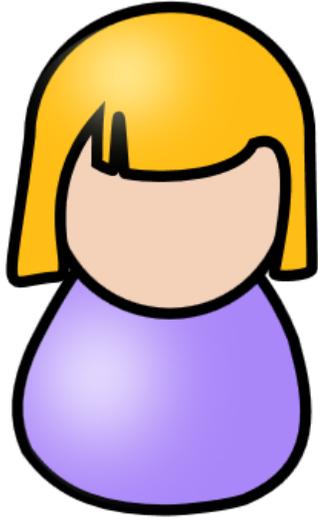


May 11-15 2020  
Virtual conference

Crypto 2020



August 17-21 2020  
Virtual



Charlie	Eric
Peihan	Frank
David	Melissa
Eve	Peihan
Melissa	Grace

PKC 2020

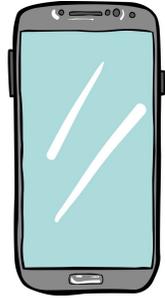
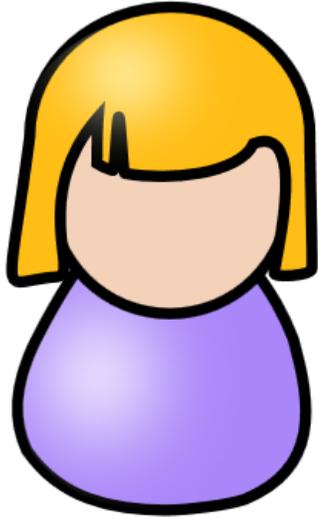


June 1-4 2020  
Virtual

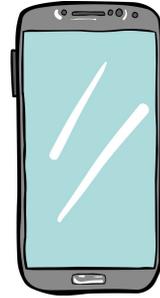
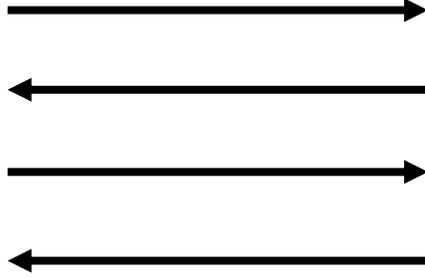
CHES 2020



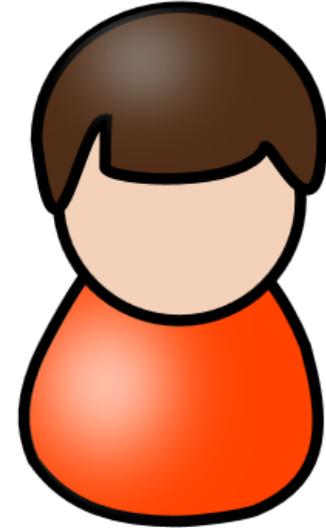
14-18 September 2020  
Virtual Conference

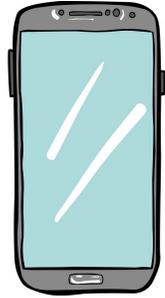
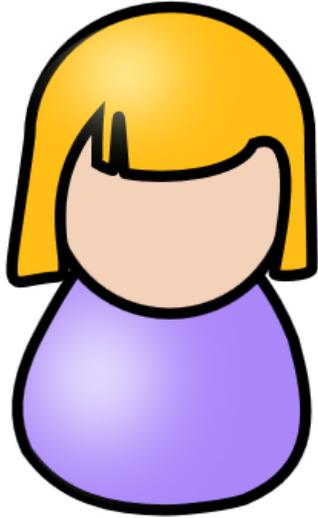


Charlie
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Eve
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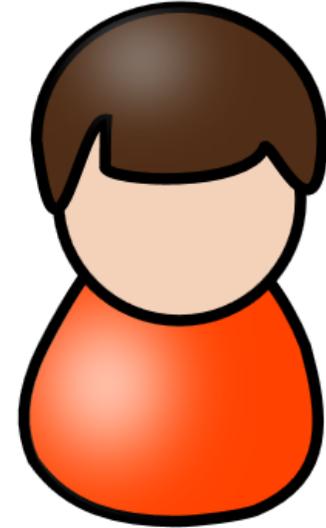
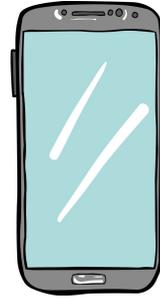
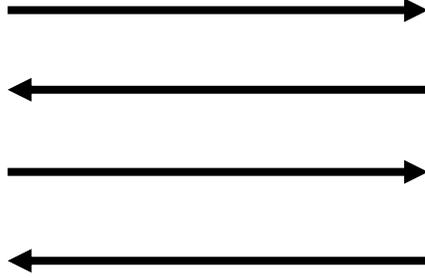


Eric
Frank
Melissa
Peihan
Grace

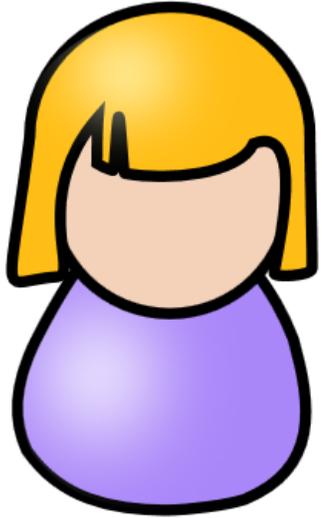




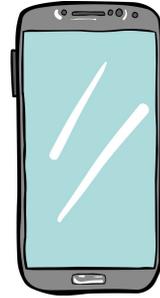
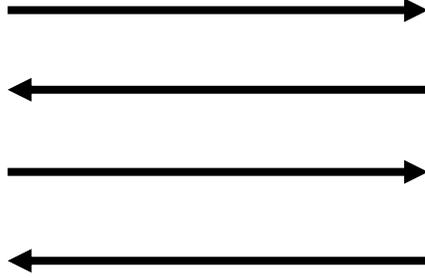
Charlie
Peihan
David
Eve
Melissa



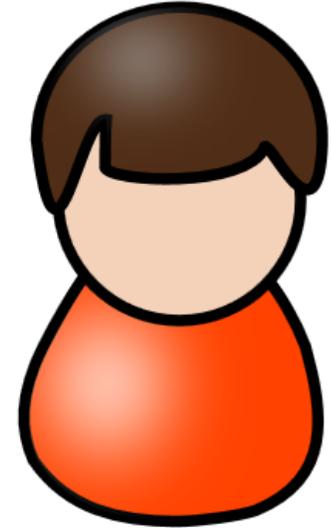
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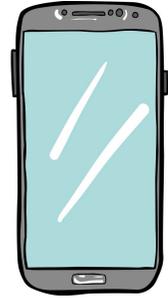
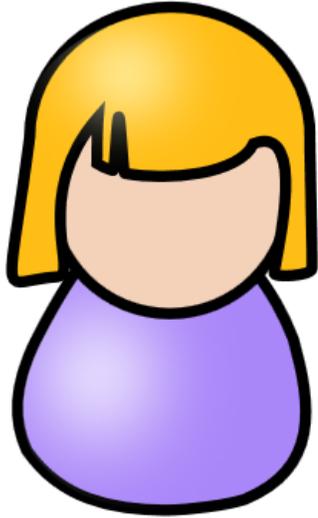


Charlie
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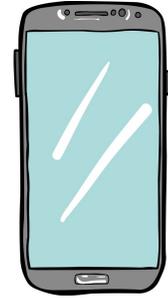
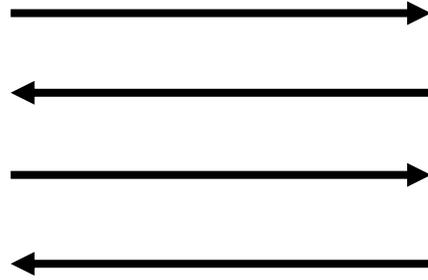


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Melissa
Peihan
???

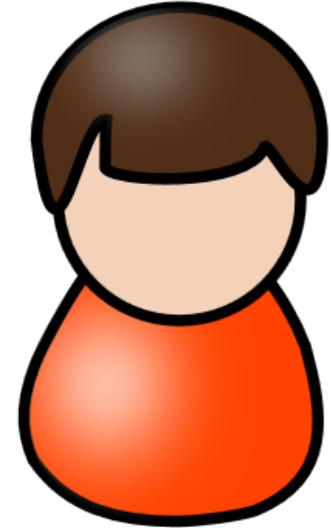




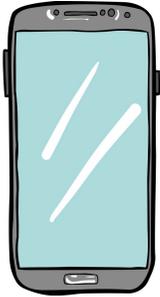
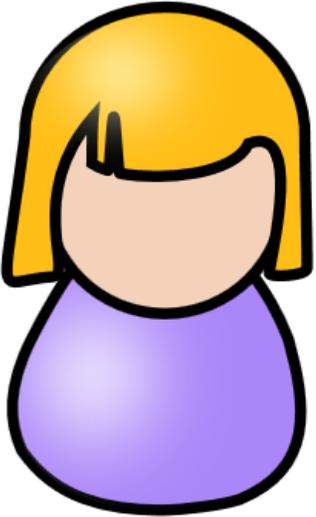
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Peihan
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???
Melissa



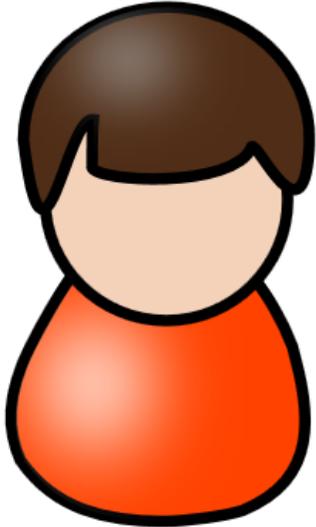
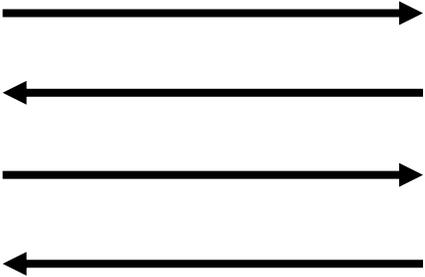
Eric
Frank
Melissa
Peihan
Grace



# Private Set Intersection (PSI)



Charlie
Peihan
David
Eve
Melissa



Eric
Frank
Melissa
Peihan
Grace

# Google open-sources cryptographic tool to keep data sets private

by RAVIE LAKSHMANAN — Jun 20, 2019 in SECURITY



79 SHARES



<https://tnw.to/C0PP1>

Poorly secured databases are a top privacy and security concern — and Google now wants to plug that leak.

The internet giant has said it's open sourcing [Private Join and Compute](#), a new secure multi-party computation (MPC) tool designed to help organizations work together with confidential data sets.

Ads Conversion Measurement

## Google Security Blog

The latest news and insights from Google on security and safety on the Internet

### Protect your accounts from data breaches with Password Checkup

February 5, 2019

Posted by Jennifer Pullman, Kurt Thomas, and Elie Bursztein, Security and Anti-abuse research

*Update (Feb 6):* We have updated the post to clarify a protocol used in the design is centered around private set intersection.

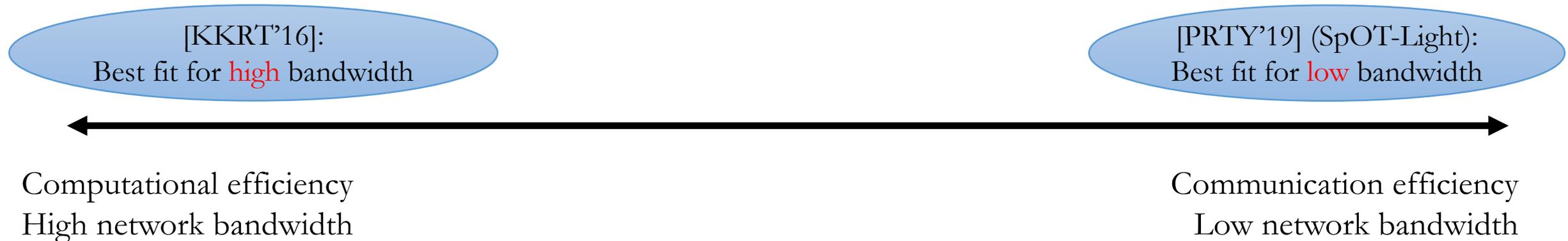
Google helps keep your account safe from hijacking with a defense in depth strategy that spans [prevention, detection, and mitigation](#). As part of this, we regularly reset the passwords of Google accounts affected by [third-party data breaches](#) in the event of password reuse. This strategy has helped us protect over 110 million users in the last two years alone. Without these safety measures, users would be at [ten times the risk](#) of account hijacking.

We want to help you stay safe not just on Google, but elsewhere on the web as well. This is where the new [Password Checkup Chrome extension](#) can help. Whenever you sign in to a site, Password Checkup will trigger a warning if the username and password you use is one of over 4 billion credentials that Google knows to be unsafe.

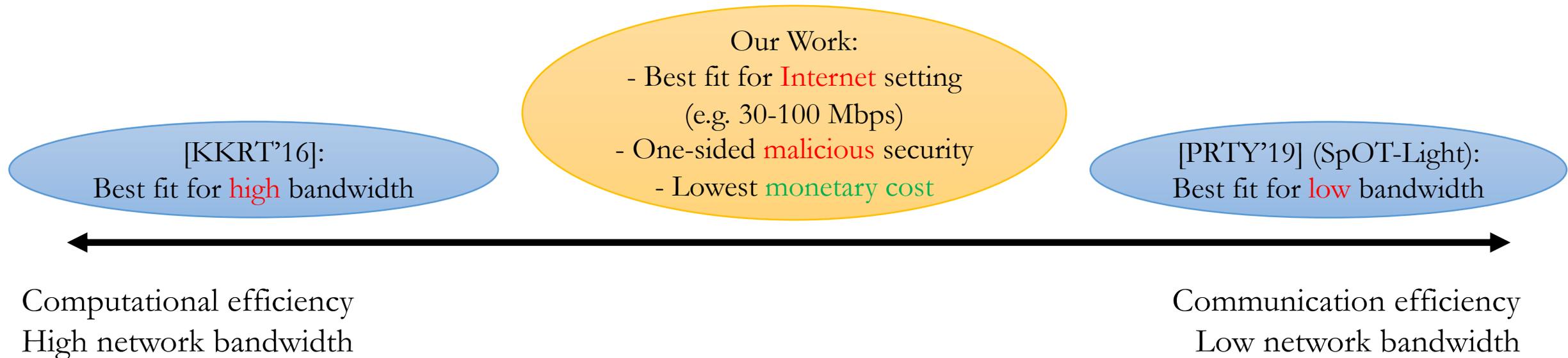
Password Breach Alert

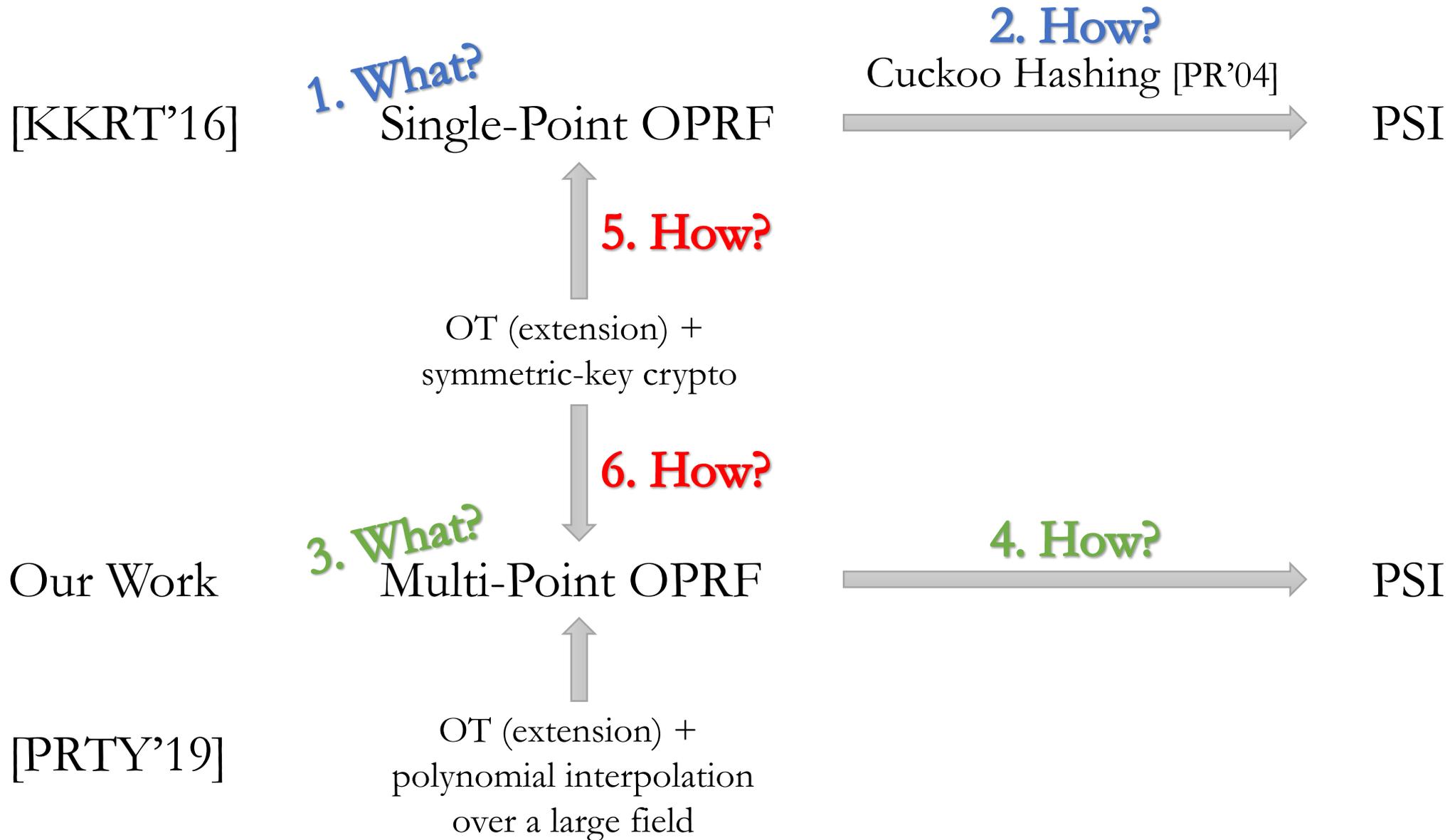
# State-of-the-art PSI

- Prior work [FNP'04, DSMRY'09, DCKT'10, ADCT'11, HEK'12, DCW13, PSZ14, PSSZ15, PSZ'14, KKRT'16, RR17a, RR17b, CLR17, DRRT18, PSWW18, GN'19, PRTY'19, PRTY'20, ...]
- This work: **semi-honest** security, sets of the **same size**
- Which protocol should we adopt?
  - Most efficient one!
  - Tradeoff between Computation & Communication
  - **Monetary cost** [PRTY'19]

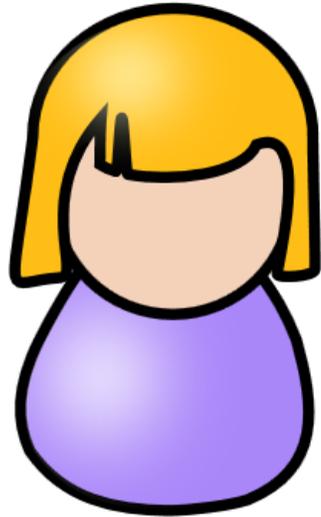


# Can we achieve a better balance?



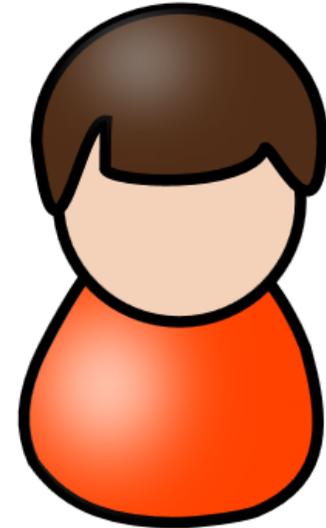
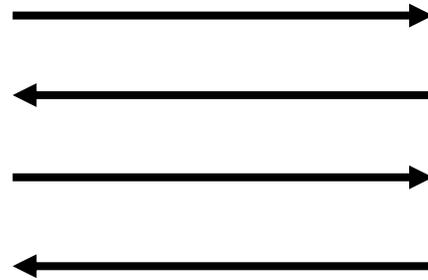


# Single-Point Oblivious Pseudorandom Function (OPRF)



Input:  $\perp$

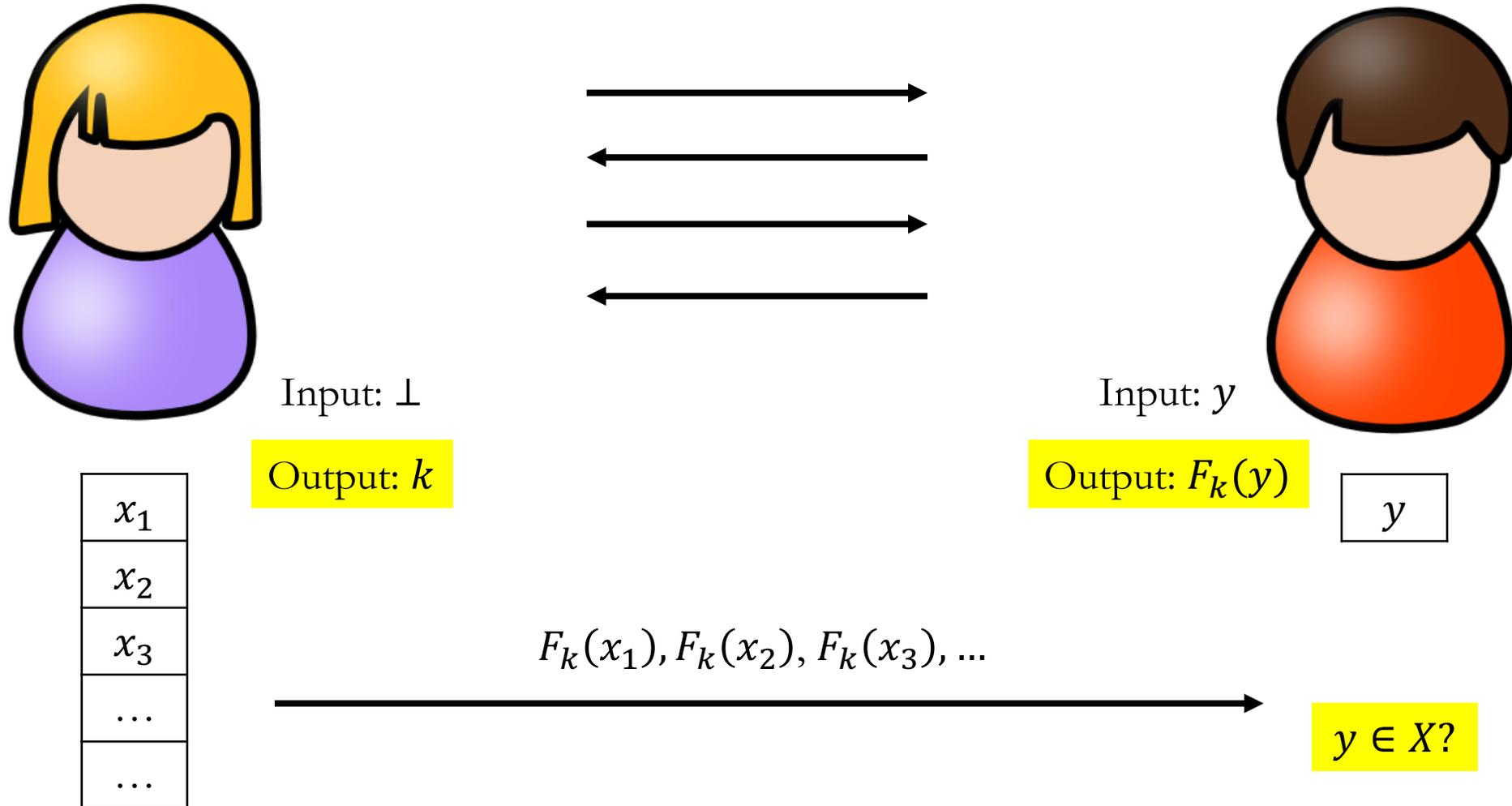
Output:  $k$



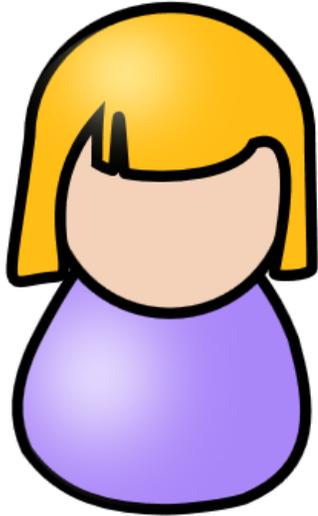
Input:  $y$

Output:  $F_k(y)$

# Single-Point PSI from Single-Point OPRF

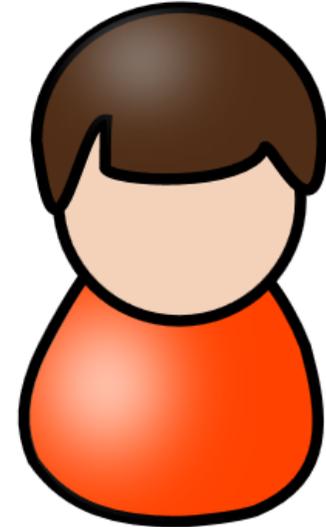


# PSI from Single-Point OPRF?



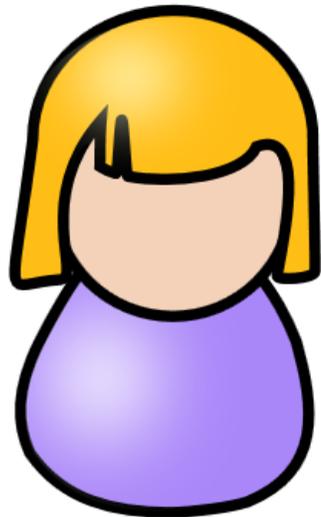
$x_1$
$x_2$
$x_3$
...
...

Computation & Communication:  $O(n^2)$

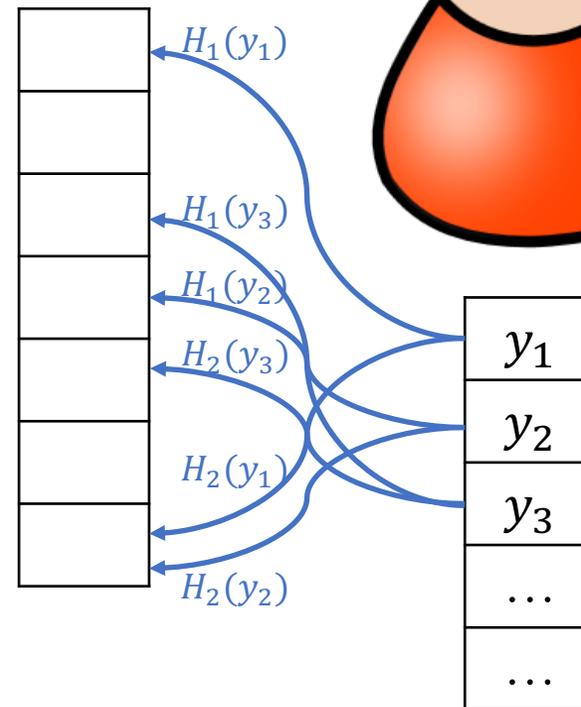
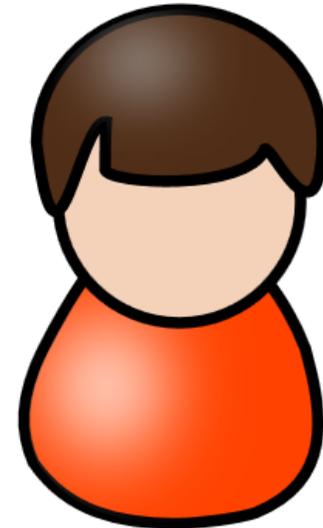


$y_1$	$y_1 \in X?$
$y_2$	$y_2 \in X?$
$y_3$	$y_3 \in X?$
...	...
...	...

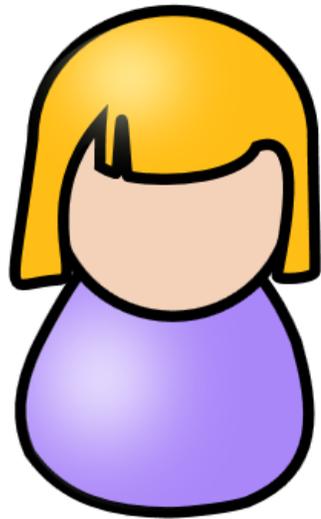
# PSI from Single-Point OPRF using Cuckoo Hashing [KKRT'16]



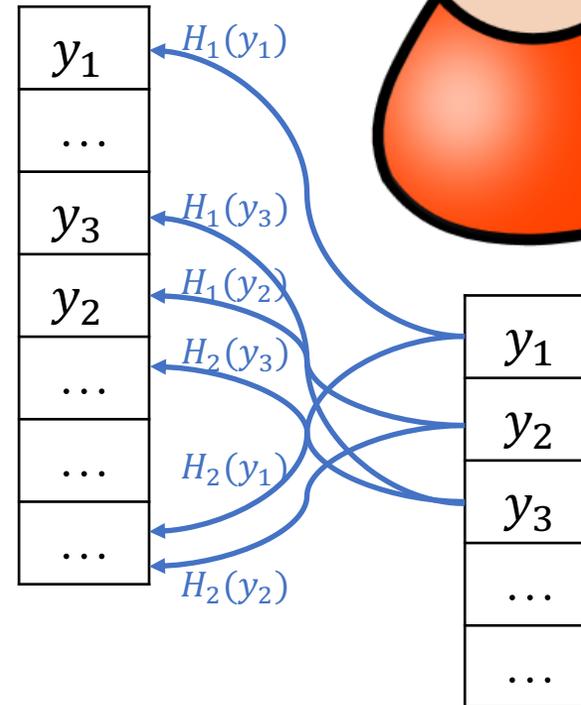
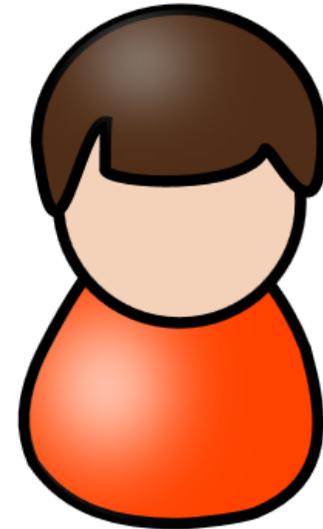
$x_1$
$x_2$
$x_3$
...
...



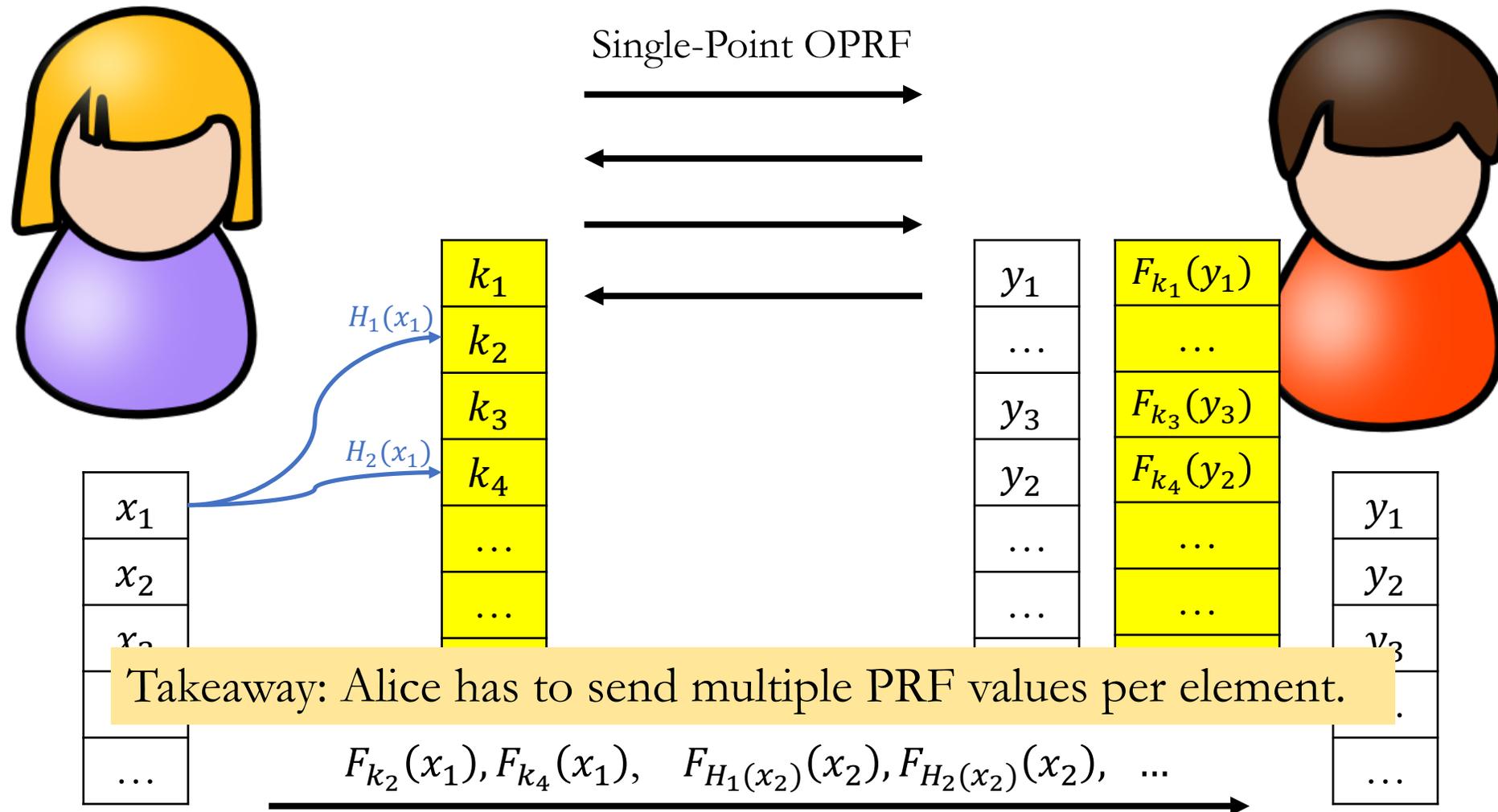
# PSI from Single-Point OPRF using Cuckoo Hashing [KKRT'16]



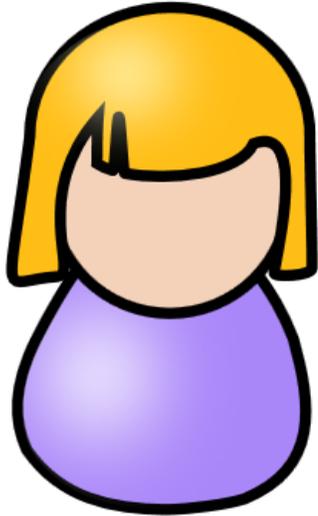
$x_1$
$x_2$
$x_3$
...
...



# PSI from Single-Point OPRF using Cuckoo Hashing [KKRT'16]

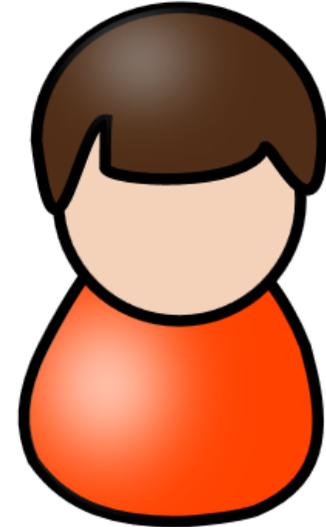
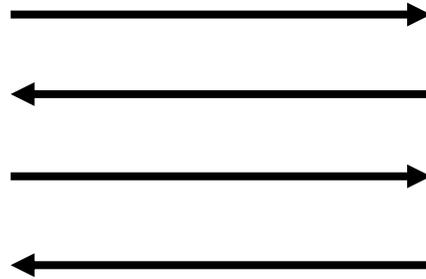


# Single-Point OPRF



Input:  $\perp$

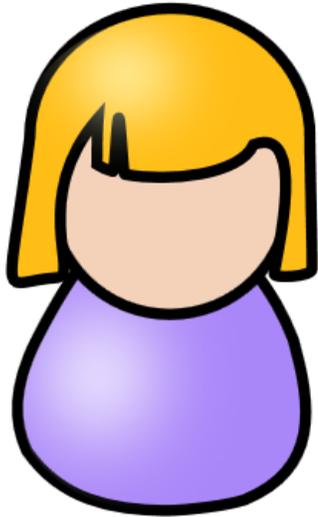
Output:  $k$



Input:  $y$

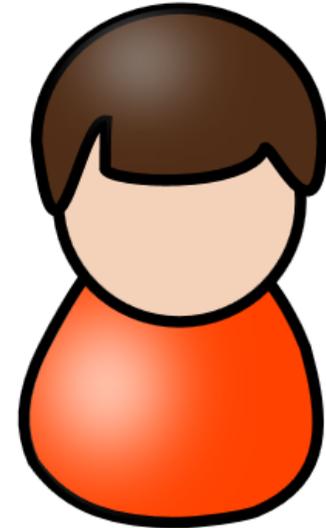
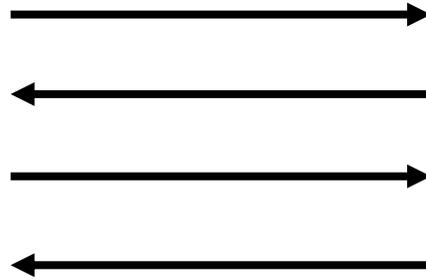
Output:  $F_k(y)$

# Multi-Point OPRF



Input:  $\perp$

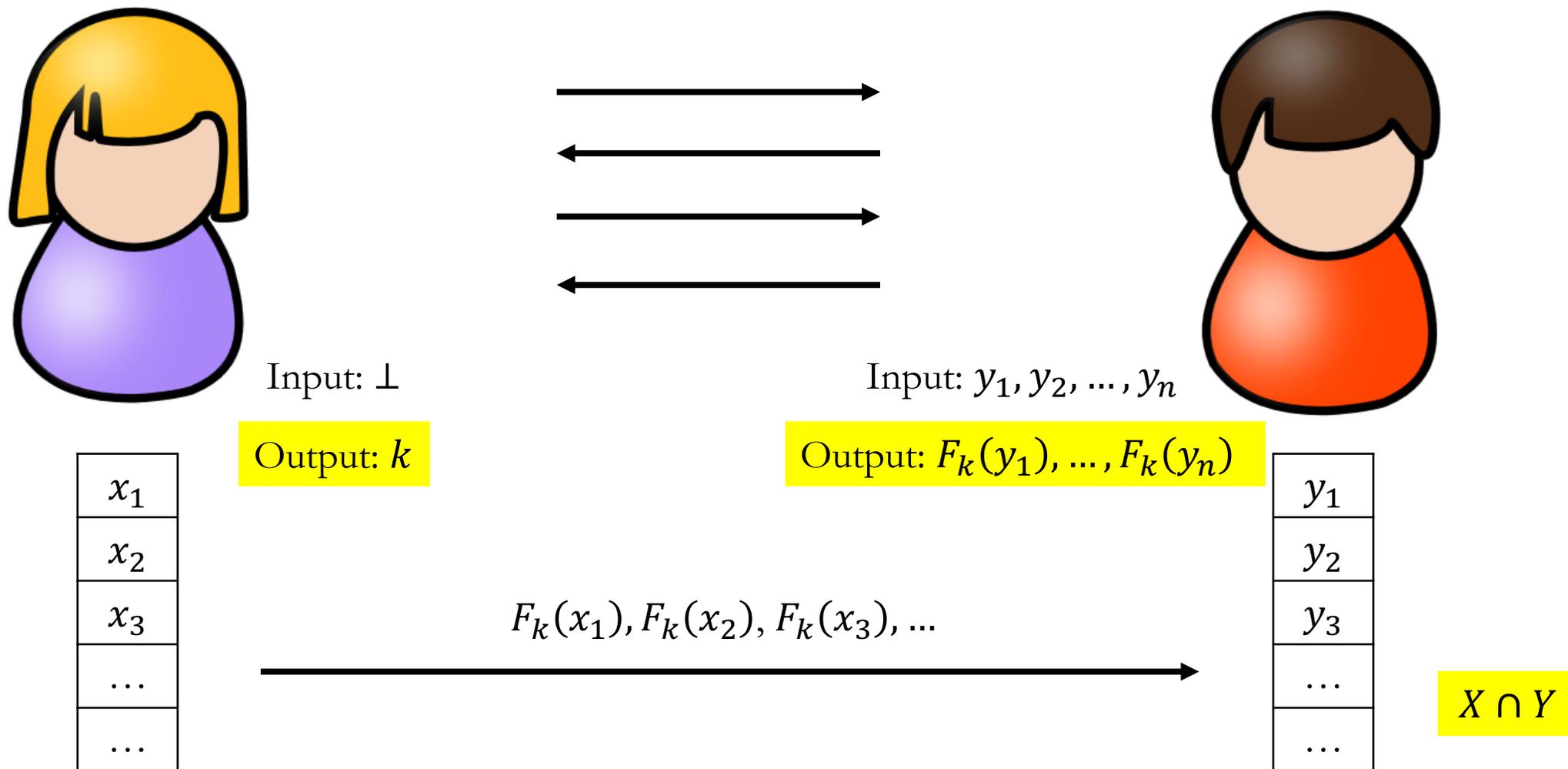
Output:  $k$

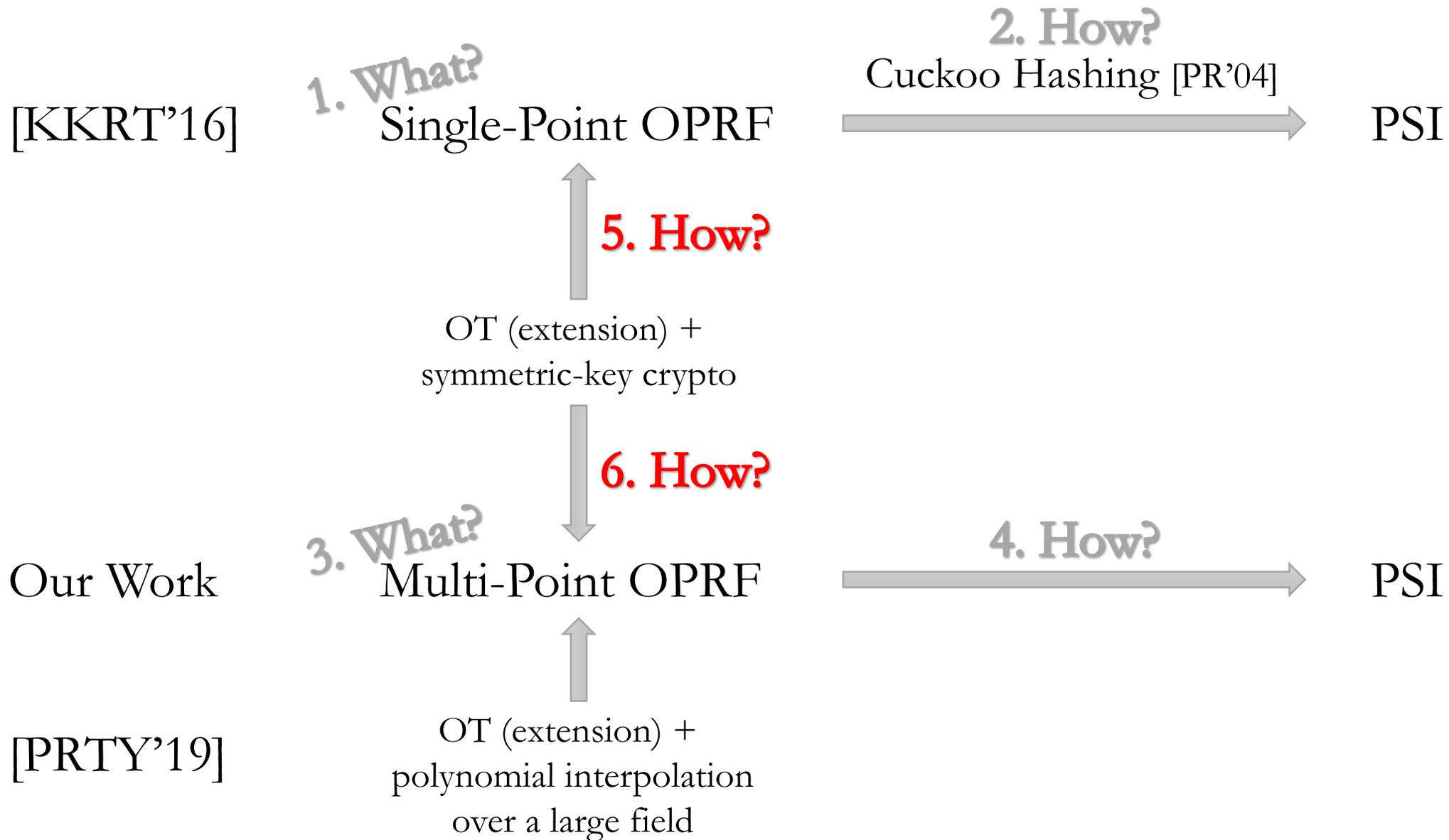


Input:  $y_1, y_2, \dots, y_n$

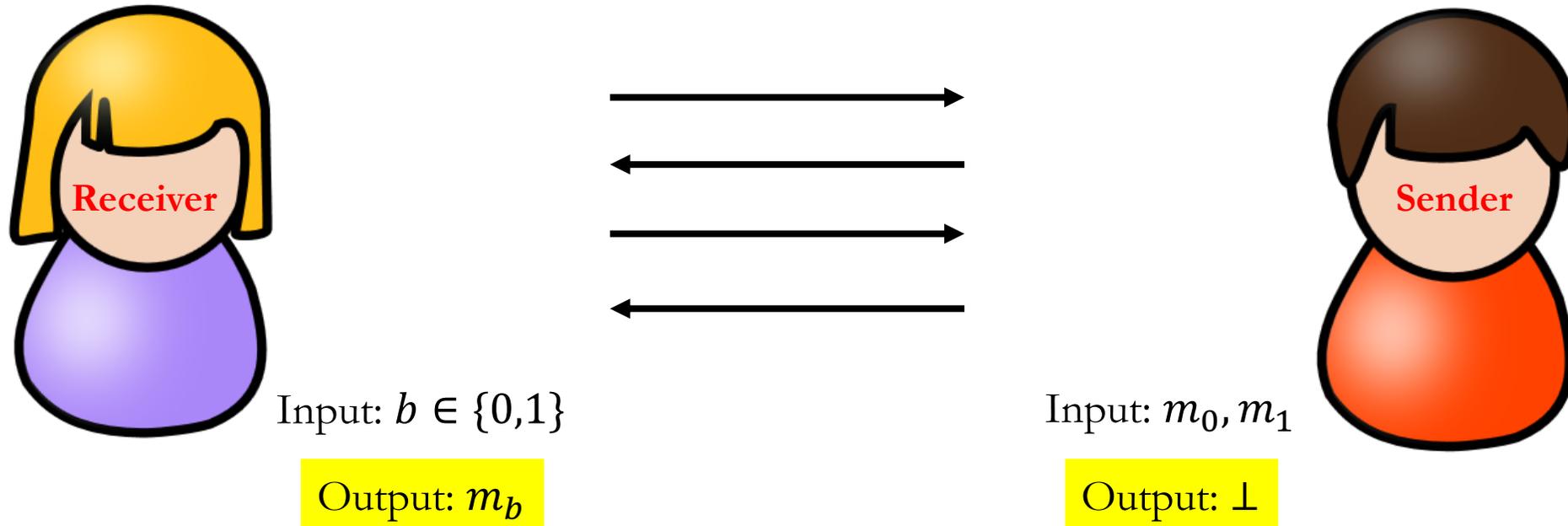
Output:  $F_k(y_1), \dots, F_k(y_n)$

# PSI from Multi-Point OPRF



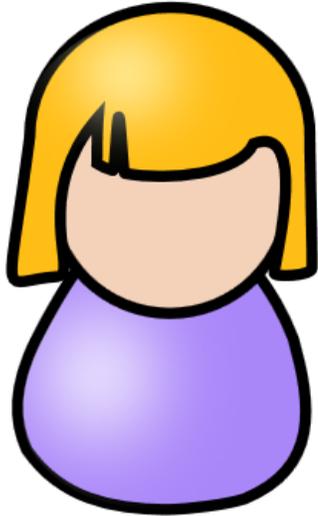


# Oblivious Transfer (OT) [Rabin'05]

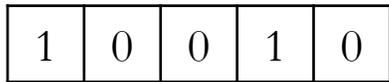


Note:  $N$  OTs can be done by  $O(\lambda)$  public-key operations and  $O(N)$  symmetric-key operations using OT extension [IKNP'03].

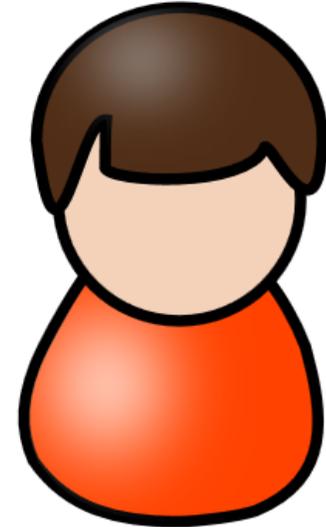
# Single-Point OPRF [KKRT'16]



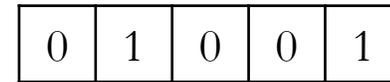
Input:  $\perp$



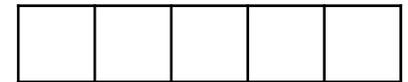
$s \leftarrow \{0,1\}^\lambda$



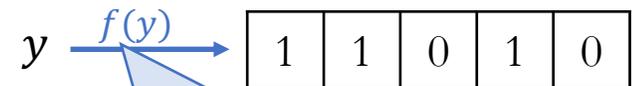
Input:  $y$



$r_0 \leftarrow \{0,1\}^\lambda$

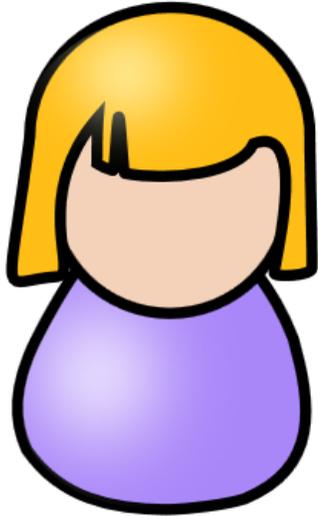


$r_1 \in \{0,1\}^\lambda$



pseudorandom, deterministic,  
e.g.  $PRF_k(y)$  with public  $k$

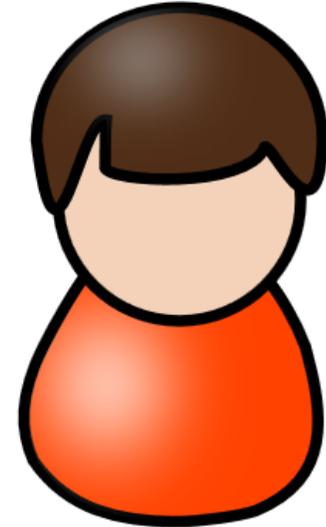
# Single-Point OPRF [KKRT'16]



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$



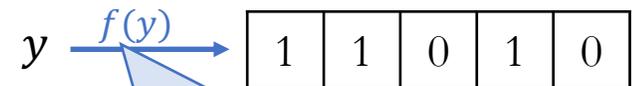
Input:  $y$

0	1	0	0	1
---	---	---	---	---

$r_0 \leftarrow \{0,1\}^\lambda$

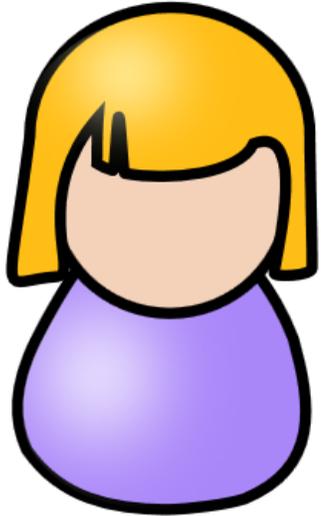
1	0	0	1	1
---	---	---	---	---

$r_1 = r_0 \oplus f(y)$



pseudorandom, deterministic,  
e.g.  $PRF_k(y)$  with public  $k$

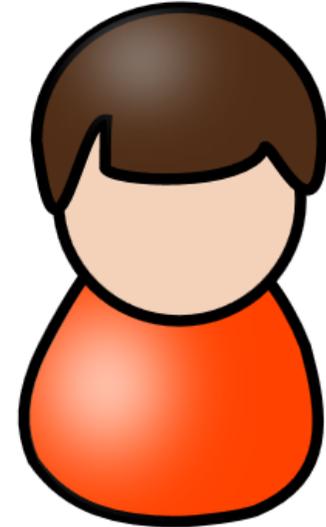
# Single-Point OPRF [KKRT'16]



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$



Input:  $y$

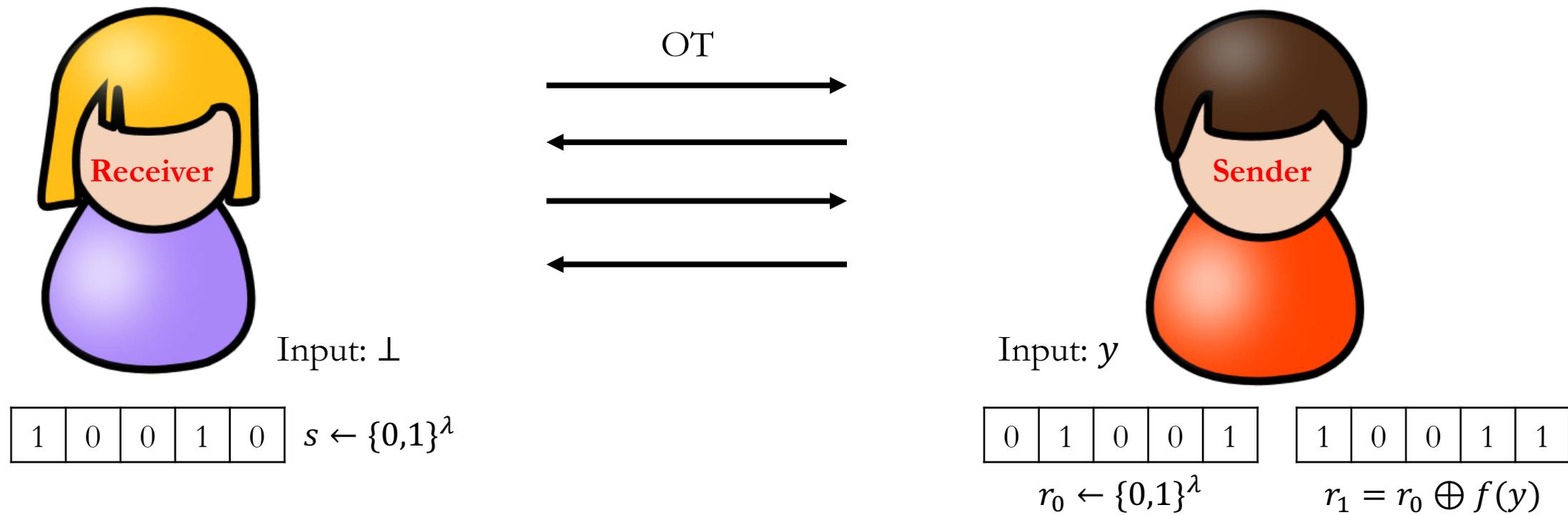
0	1	0	0	1
---	---	---	---	---

$r_0 \leftarrow \{0,1\}^\lambda$

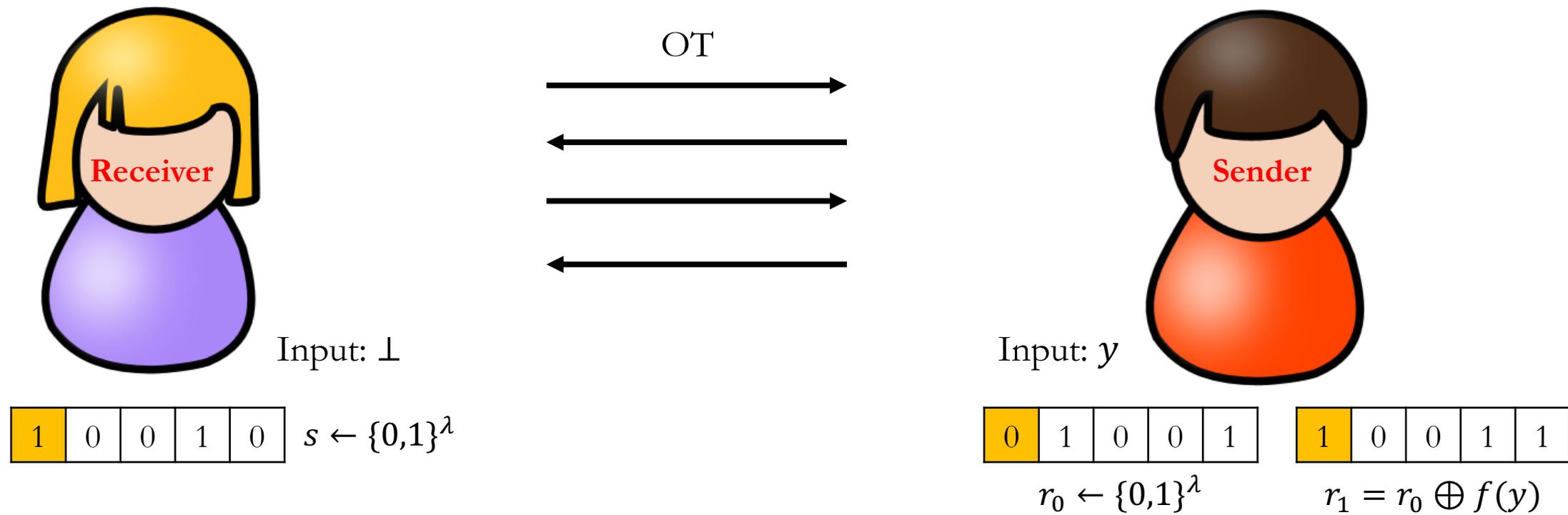
1	0	0	1	1
---	---	---	---	---

$r_1 = r_0 \oplus f(y)$

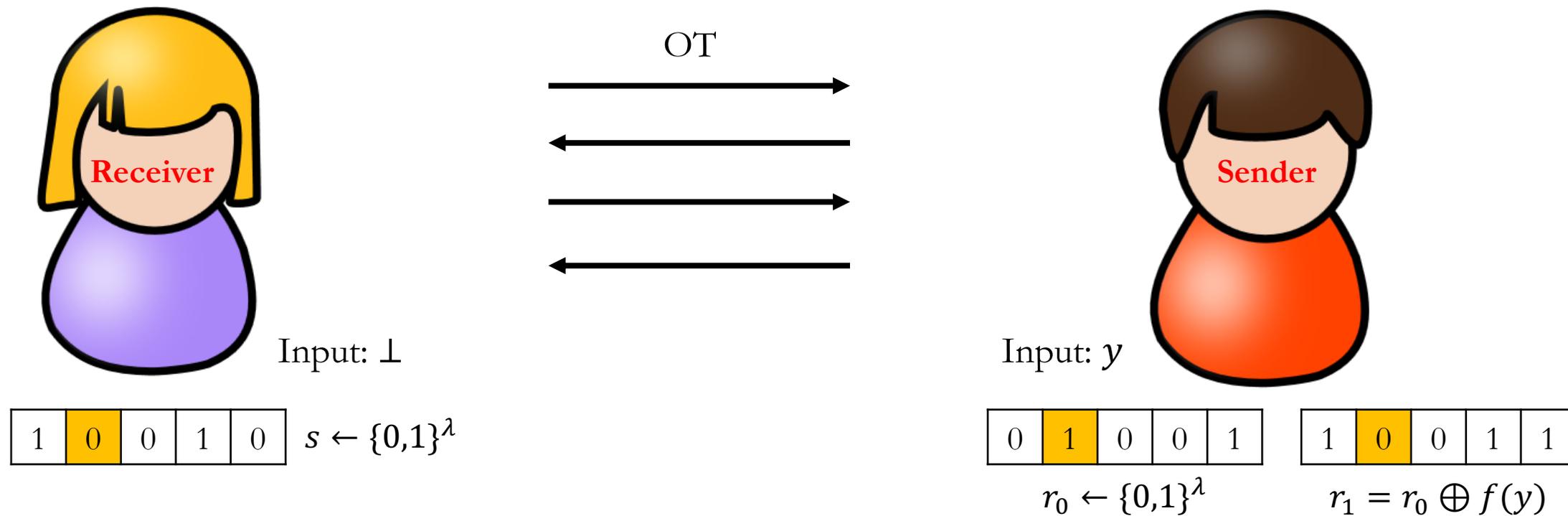
# Single-Point OPRF [KKRT'16]



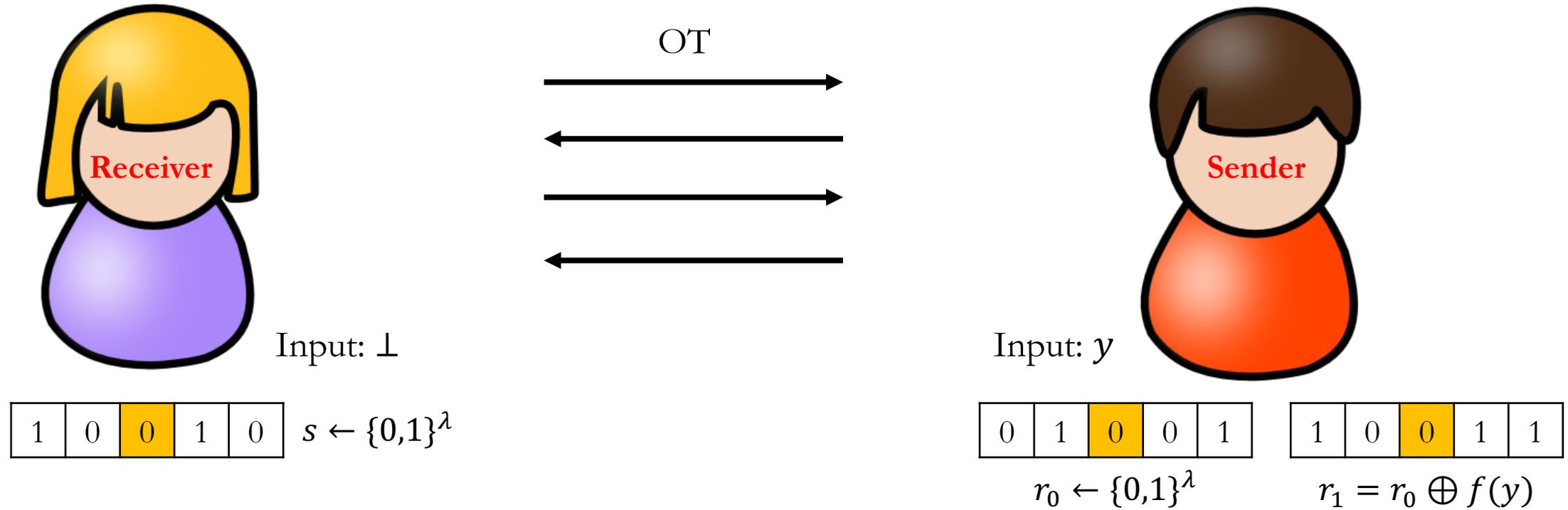
# Single-Point OPRF [KKRT'16]



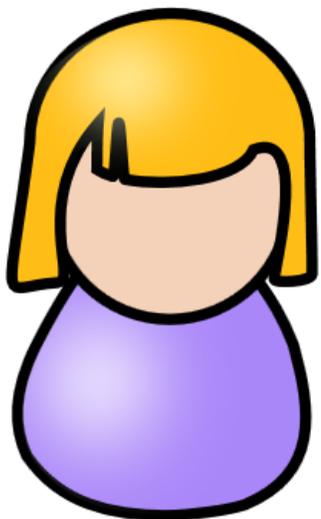
# Single-Point OPRF [KKRT'16]



# Single-Point OPRF [KKRT'16]



# Single-Point OPRF [KKRT'16]



Input:  $x$

1	0	0	1	0
---	---	---	---	---

$$s \leftarrow \{0,1\}^\lambda$$

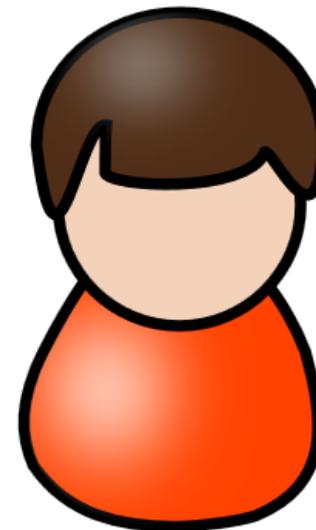
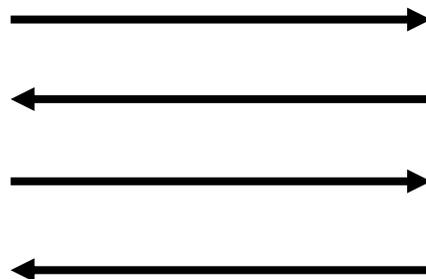
1	1	0	1	1
---	---	---	---	---

$$q = r_0 \oplus [s \cdot f(y)]$$

$x \xrightarrow{f(x)}$

0	1	0	0	1
---	---	---	---	---

OT



Input:  $y$

0	1	0	0	1
---	---	---	---	---

$$r_0 \leftarrow \{0,1\}^\lambda$$

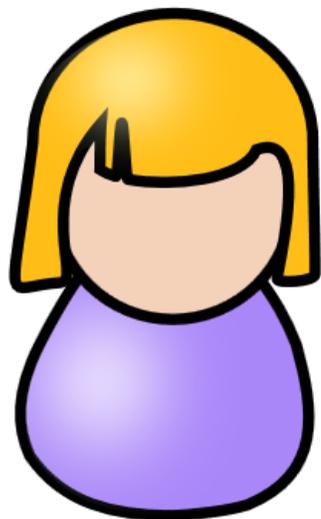
1	0	0	1	1
---	---	---	---	---

$$r_1 = r_0 \oplus f(y)$$

Output:  $k = (s, q)$ ,  $F_k(x) = H(q \oplus [s \cdot f(x)])$

Output:  $F_k(y) = H(r_0)$

# Single-Point OPRF [KKRT'16]



Input:  $x$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$

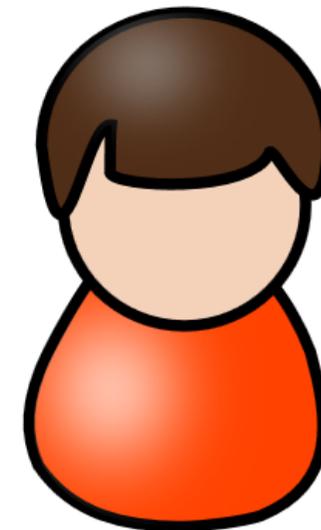
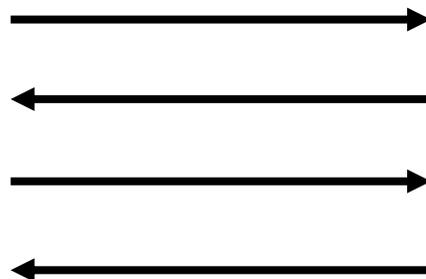
1	1	0	1	1
---	---	---	---	---

$q = r_0 \oplus [s \cdot f(y)]$

$x \xrightarrow{f(x)}$

0	1	0	0	1
---	---	---	---	---

OT



Input:  $y$

0	1	0	0	1
---	---	---	---	---

$r_0 \leftarrow \{0,1\}^\lambda$

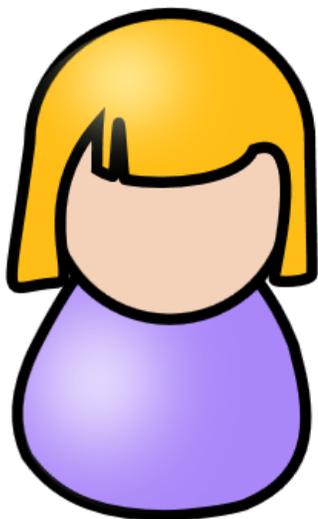
1	0	0	1	1
---	---	---	---	---

$r_1 = r_0 \oplus f(y)$

Output:  $k = (s, q)$ ,  $F_k(x) = H(q \oplus [s \cdot f(x)])$

Output:  $F_k(y) = H(r_0)$

# Single-Point OPRF [KKRT'16]



Input:  $x$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$

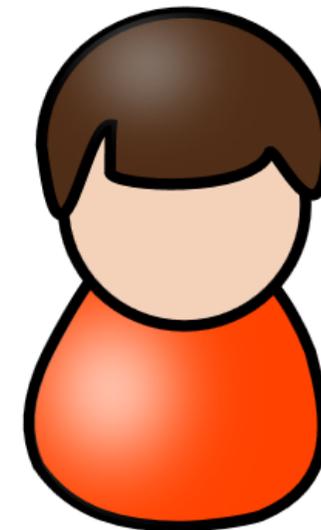
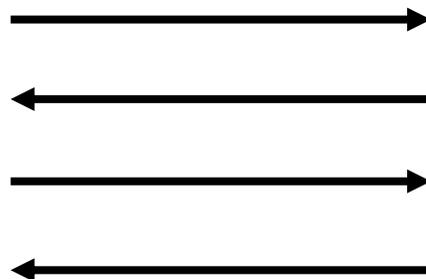
1	1	0	1	1
---	---	---	---	---

$q = r_0 \oplus [s \cdot f(y)]$

$x \xrightarrow{f(x)}$

0	1	0	0	1
---	---	---	---	---

OT



Input:  $y$

0	1	0	0	1
---	---	---	---	---

$r_0 \leftarrow \{0,1\}^\lambda$

1	0	0	1	1
---	---	---	---	---

$r_1 = r_0 \oplus f(y)$

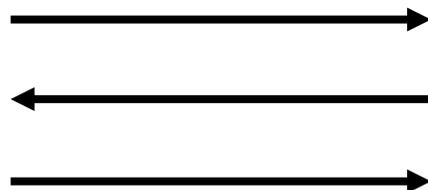
Output:  $k = (s, q)$ ,  $F_k(x) = H(q \oplus [s \cdot f(x)])$

Output:  $F_k(y) = H(r_0)$

# Single-Point OPRF [KKRT'16]



OT



Takeaways:

- (a) If  $x = y$ , then  $F_k(x) = H(r_0)$  no matter what  $s$  is chosen.
- (b) If  $x \neq y$ , then  $F_k(x)$  is hard to guess.

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$

1	1	0	1	1
---	---	---	---	---

$q = r_0 \oplus [s \cdot f(y)]$

$x \xrightarrow{f(x)}$

0	1	0	0	1
---	---	---	---	---

0	1	0	0	1
---	---	---	---	---

$r_0 \leftarrow \{0,1\}^\lambda$

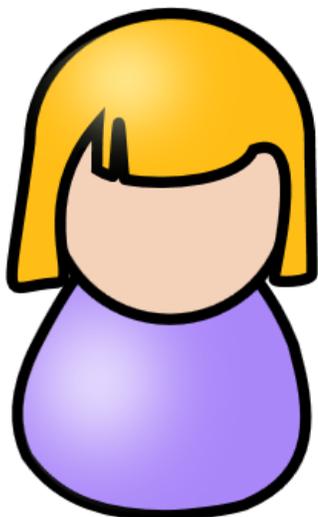
1	0	0	1	1
---	---	---	---	---

$r_1 = r_0 \oplus f(y)$

Output:  $k = (s, q)$ ,  $F_k(x) = H(q \oplus [s \cdot f(x)])$   
 $= H(r_0 \oplus [s \cdot (f(x) \oplus f(y))])$

Output:  $F_k(y) = H(r_0)$

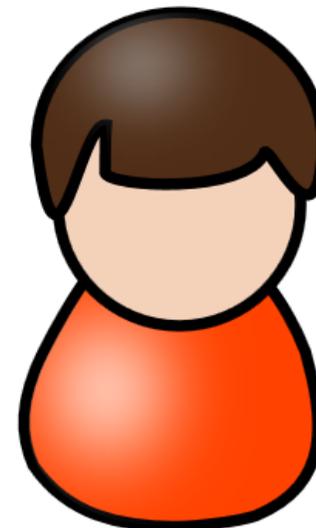
# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

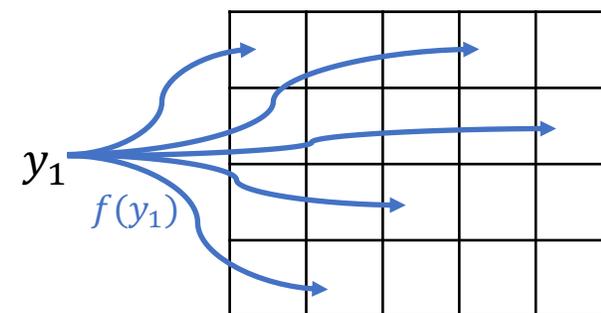
$s \leftarrow \{0,1\}^\lambda$



Input:  $y_1, y_2, \dots, y_n$

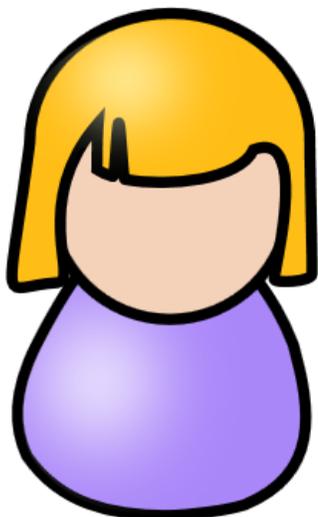
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$



$R_1 \in \{0,1\}^{n \times \lambda}$

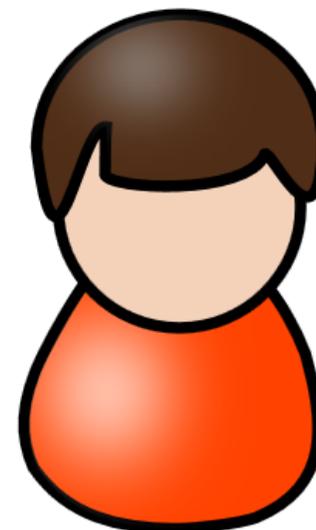
# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$



Input:  $y_1, y_2, \dots, y_n$

0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

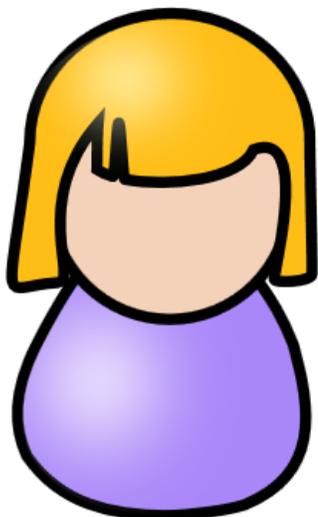

$R_1 \in \{0,1\}^{n \times \lambda}$

$y_1$

$f(y_1)$

pseudorandom, deterministic,  
e.g.  $PRF_k(y)$  with public  $k$

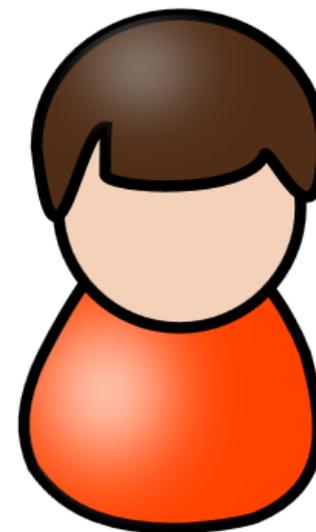
# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$



Input:  $y_1, y_2, \dots, y_n$

0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

	0		1	
				0
			0	
		0		

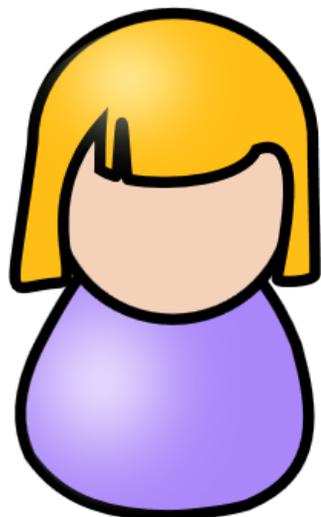
$R_1 \in \{0,1\}^{n \times \lambda}$

$y_1$

$f(y_1)$

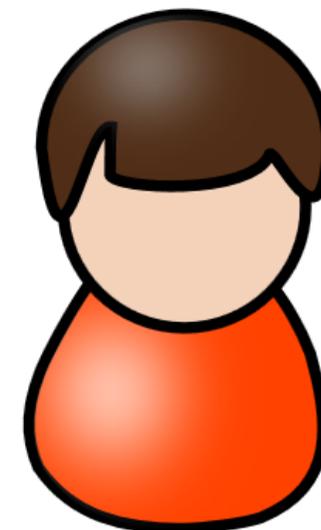
pseudorandom, deterministic,  
e.g.  $PRF_k(y)$  with public  $k$

# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

 $s \leftarrow \{0,1\}^\lambda$ 

Input:  $y_1, y_2, \dots, y_n$

0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

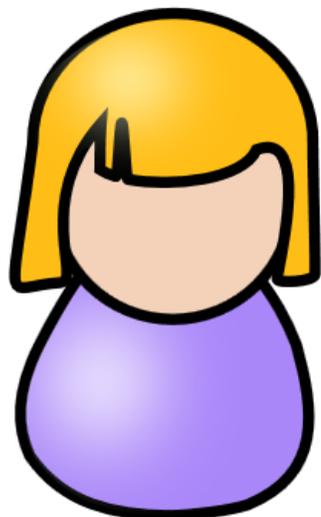
$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

	0			1	
					0
$y_2$				0	
				0	

$f(y_2)$

$R_1 \in \{0,1\}^{n \times \lambda}$

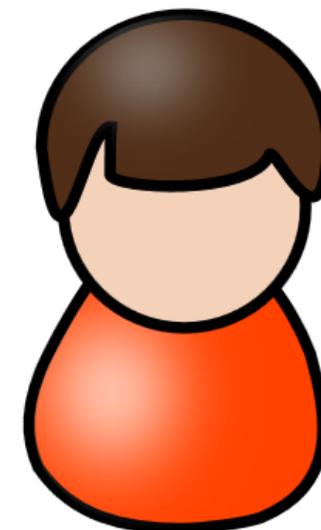
# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$



Input:  $y_1, y_2, \dots, y_n$

0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

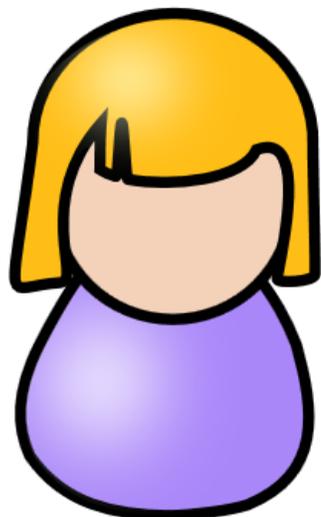
	0			1	
					0
			0		

$y_2$  (with arrows pointing to the first three rows)

$f(y_2)$  (with arrows pointing to the yellow cells)

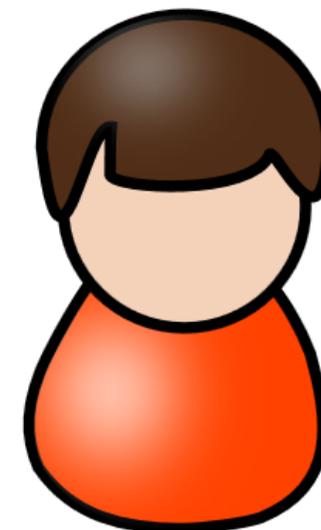
$R_1 \in \{0,1\}^{n \times \lambda}$

# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

 $s \leftarrow \{0,1\}^\lambda$ 

Input:  $y_1, y_2, \dots, y_n$

0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

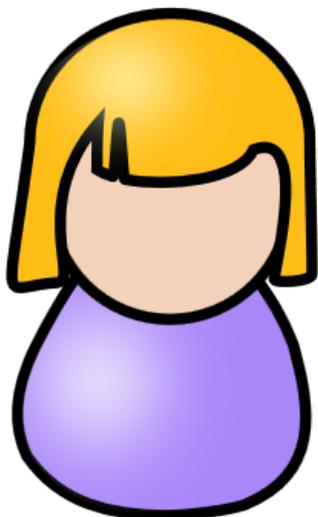
$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

	0	1		1	
					0
			0		
$y_2$					
	0	0		1	

$f(y_2)$

$R_1 \in \{0,1\}^{n \times \lambda}$

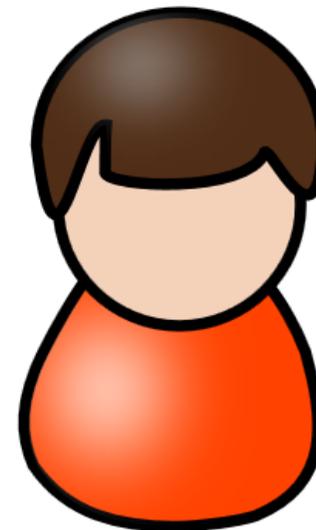
# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

$s \leftarrow \{0,1\}^\lambda$



Input:  $y_1, y_2, \dots, y_n$

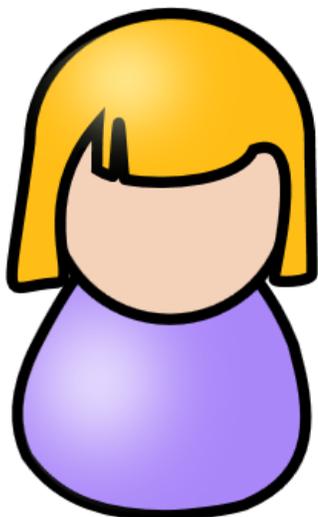
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

0	1		1	
				0
		0		
0	0		1	

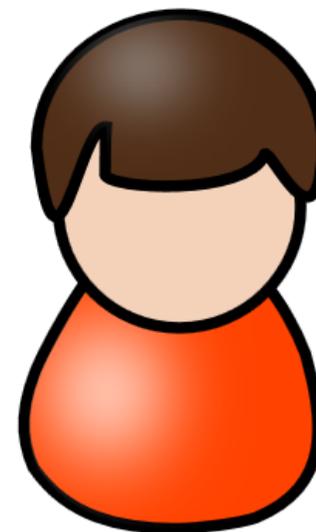
$R_1 \in \{0,1\}^{n \times \lambda}$

# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

 $s \leftarrow \{0,1\}^\lambda$ 

Input:  $y_1, y_2, \dots, y_n$

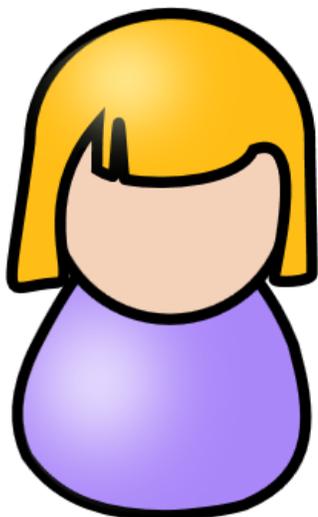
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

0	1	0	1	
1			1	0
		0		1
0	0		1	

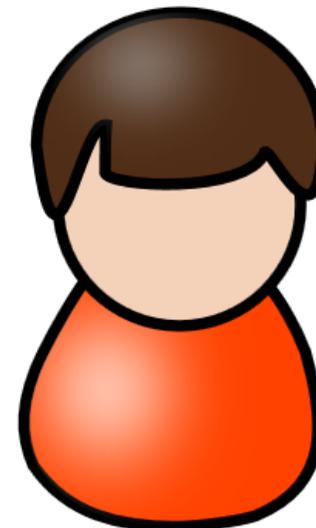
$R_1 \in \{0,1\}^{n \times \lambda}$

# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

 $s \leftarrow \{0,1\}^\lambda$ 

Input:  $y_1, y_2, \dots, y_n$

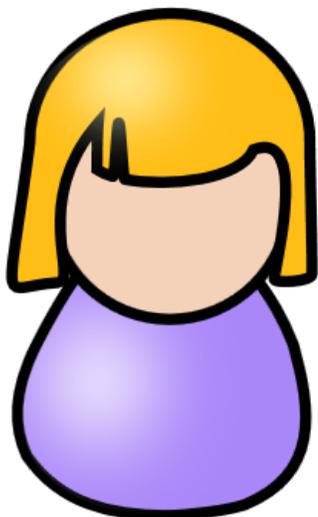
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

0	1	0	1	
1			1	0
		0		1
0	0		1	

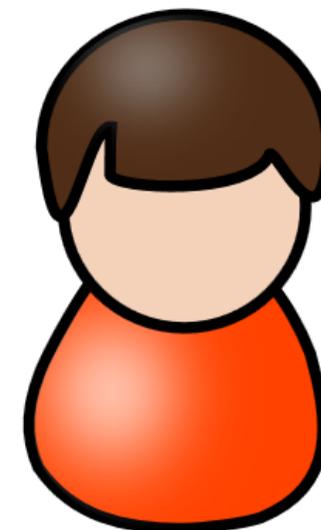
$R_1 \in \{0,1\}^{n \times \lambda}$

# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

 $s \leftarrow \{0,1\}^\lambda$ 

Input:  $y_1, y_2, \dots, y_n$

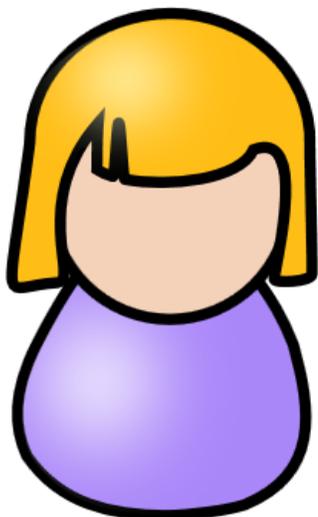
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

0	1	0	1	0
1	1	1	1	0
0	0	0	1	1
0	0	0	1	0

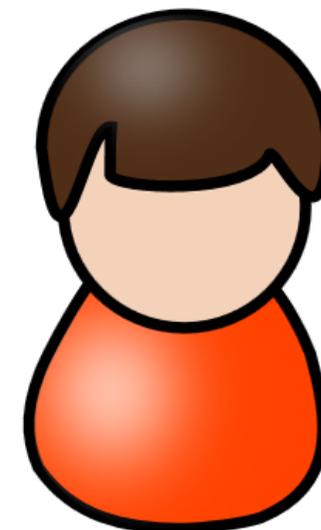
$R_1 \in \{0,1\}^{n \times \lambda}$

# Multi-Point OPRF



Input:  $\perp$

1	0	0	1	0
---	---	---	---	---

 $s \leftarrow \{0,1\}^\lambda$ 

Input:  $y_1, y_2, \dots, y_n$

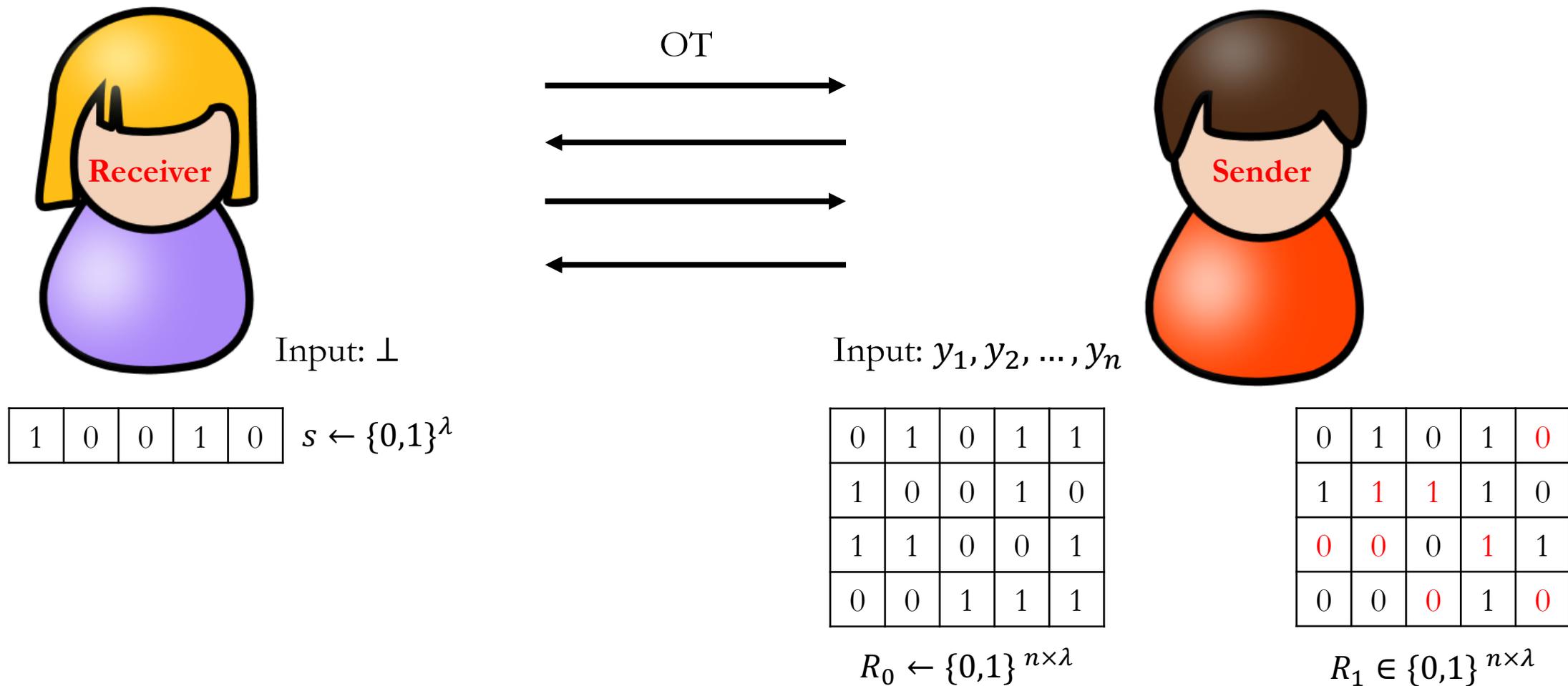
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

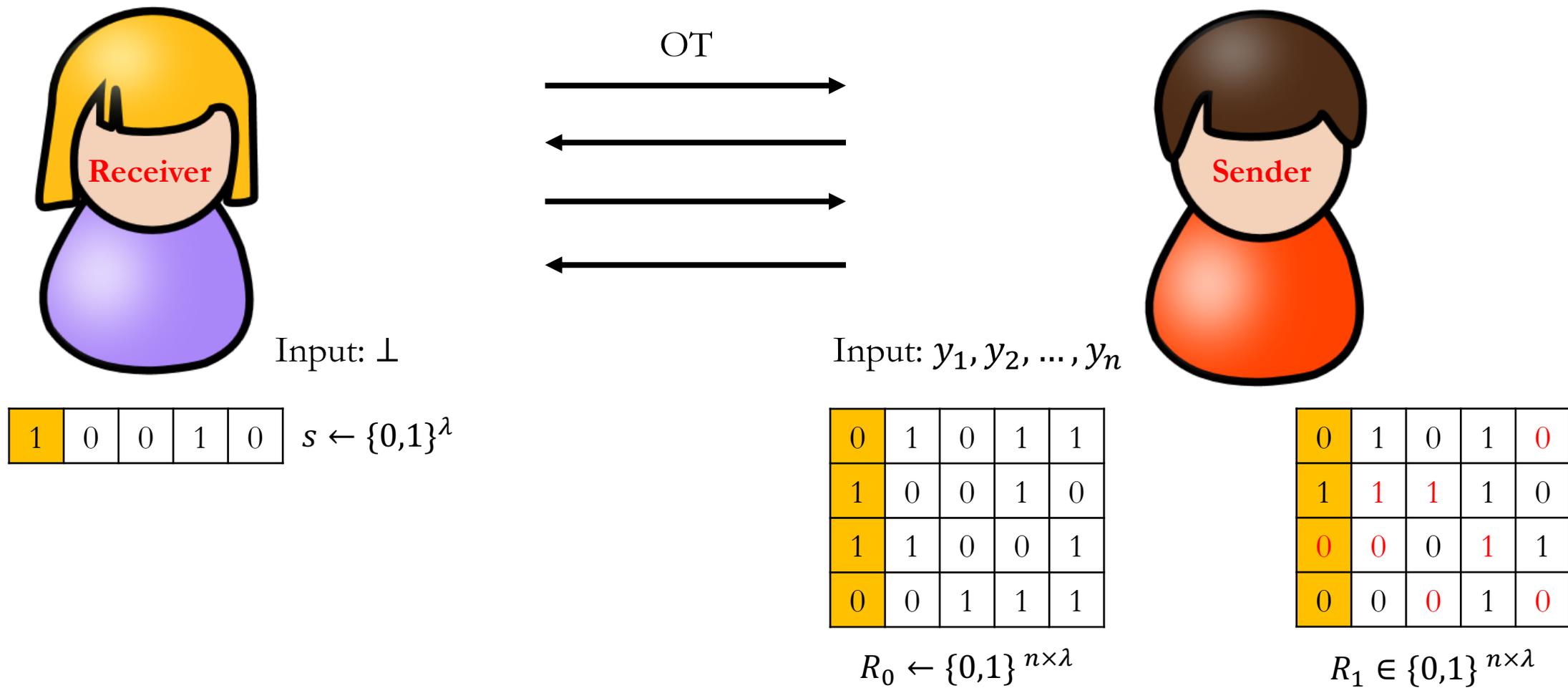
0	1	0	1	0
1	1	1	1	0
0	0	0	1	1
0	0	0	1	0

$R_1 \in \{0,1\}^{n \times \lambda}$

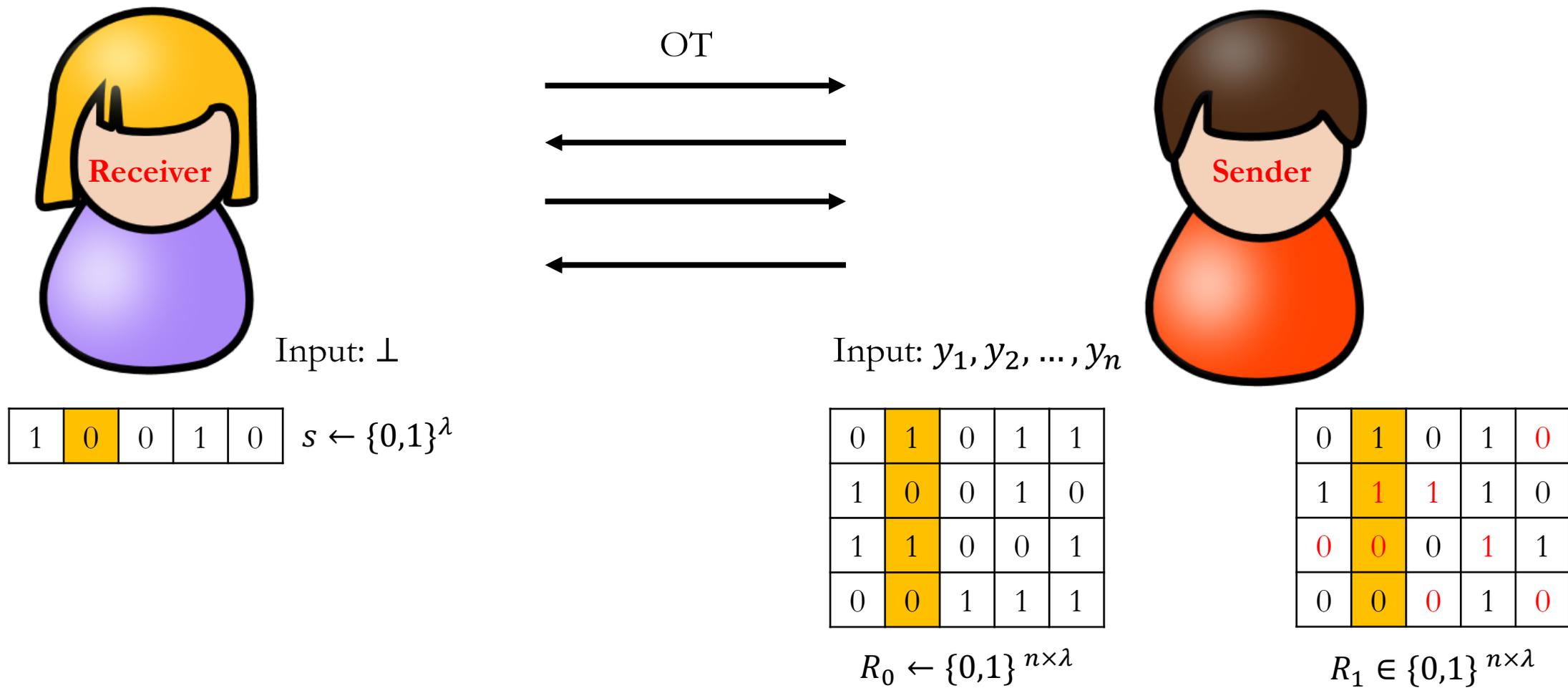
# Multi-Point OPRF



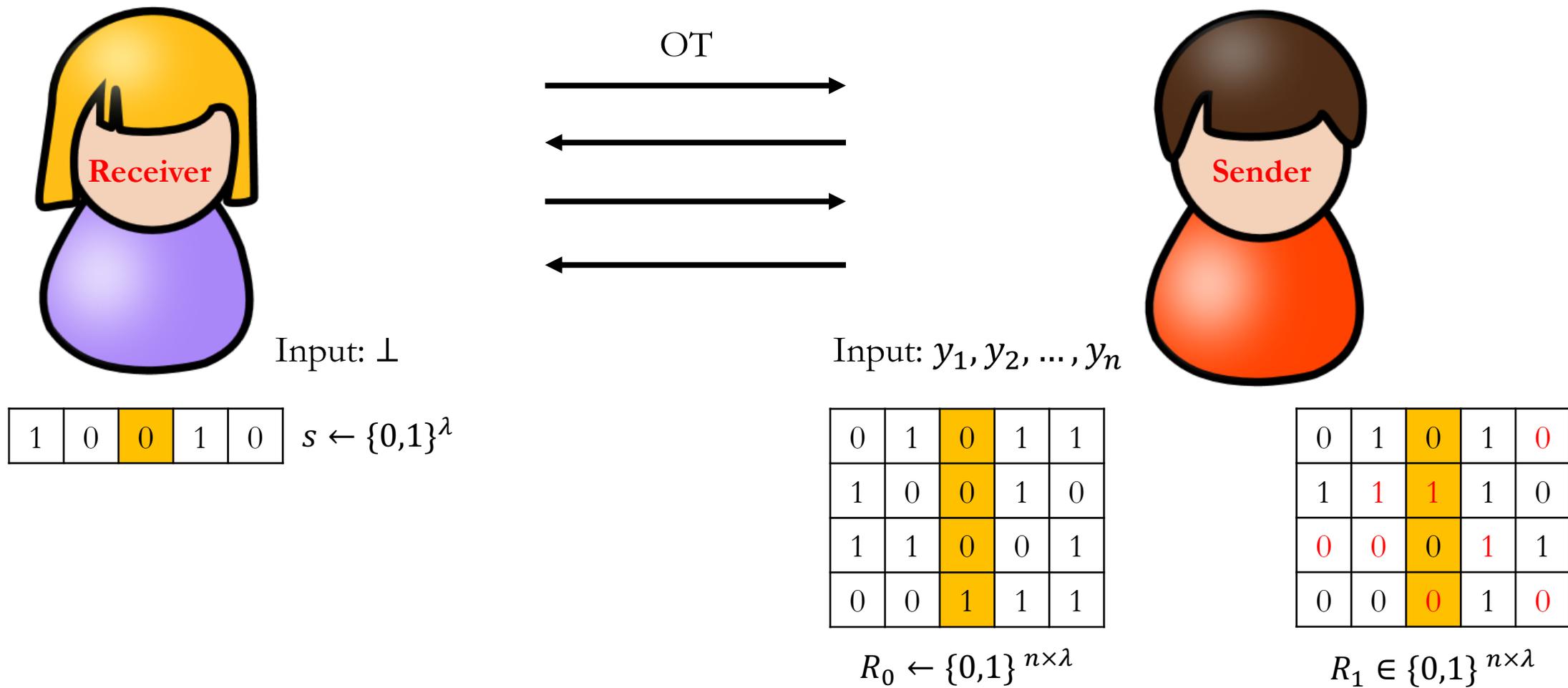
# Multi-Point OPRF



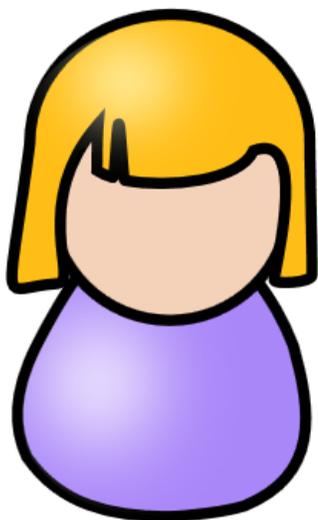
# Multi-Point OPRF



# Multi-Point OPRF

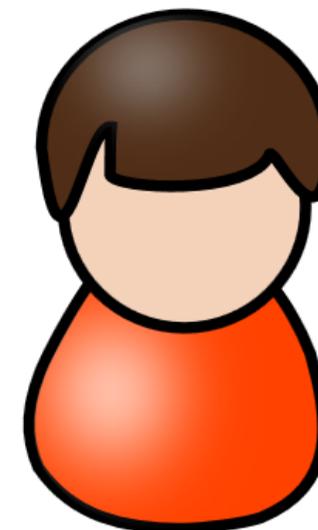
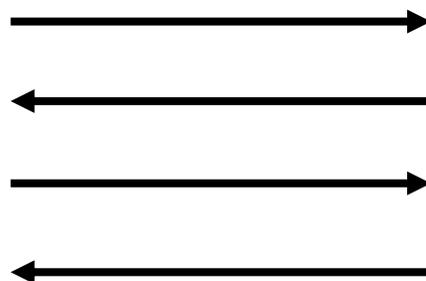


# Multi-Point OPRF



Input:  $\perp$

OT



Input:  $y_1, y_2, \dots, y_n$

1	0	0	1	0
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$s \leftarrow \{0,1\}^\lambda$

$F_k(x)$ :

$x$

$f(x)$

0	0	0	1	1
1	1	0	1	0
0	1	0	1	1
0	0	1	1	1

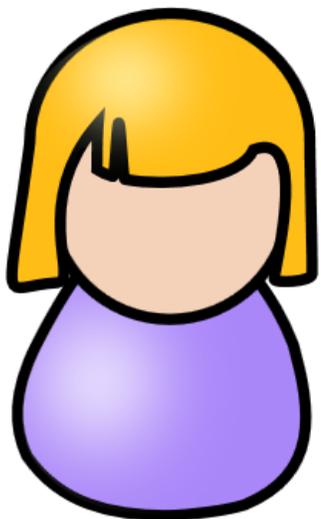
0	1	0	1	1
1	0	0	1	0
1	1	0	0	1
0	0	1	1	1

$R_0 \leftarrow \{0,1\}^{n \times \lambda}$

0	1	0	1	0
1	1	1	1	0
0	0	0	1	1
0	0	0	1	0

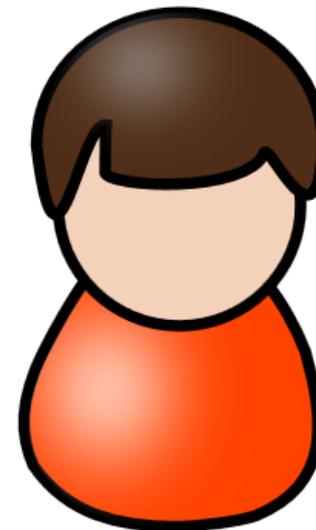
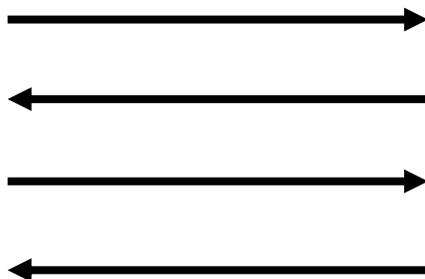
$R_1 \in \{0,1\}^{n \times \lambda}$

# Multi-Point OPRF

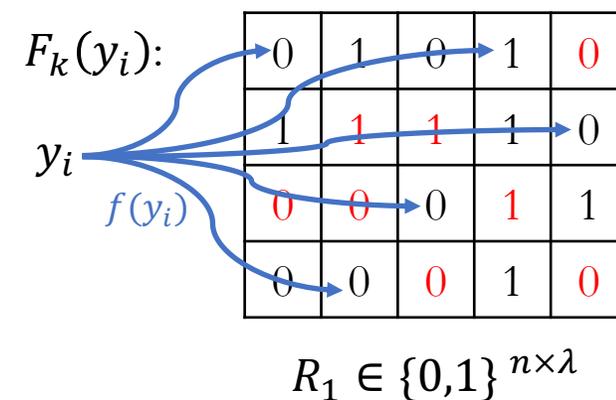
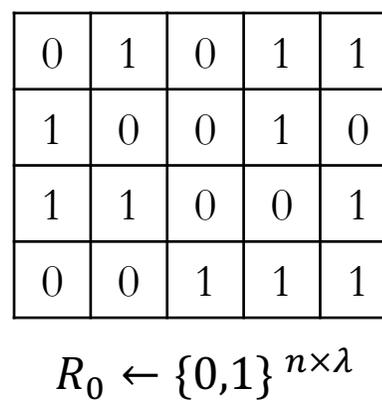
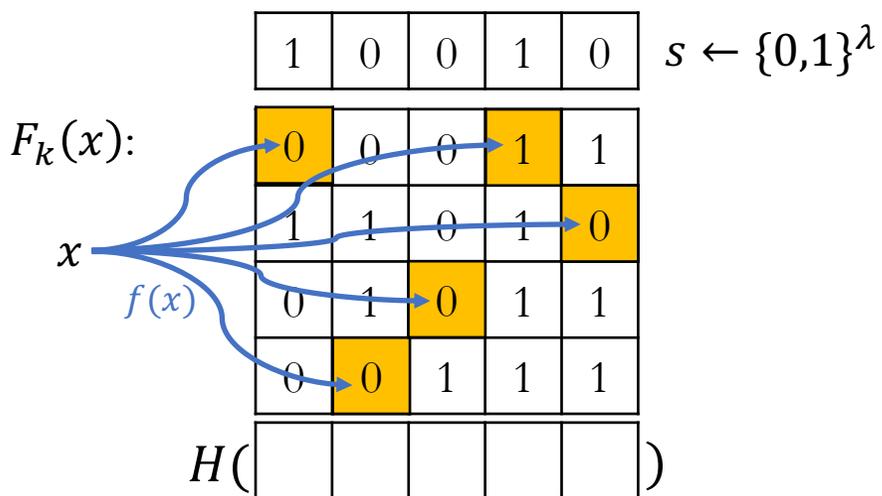


Input:  $\perp$

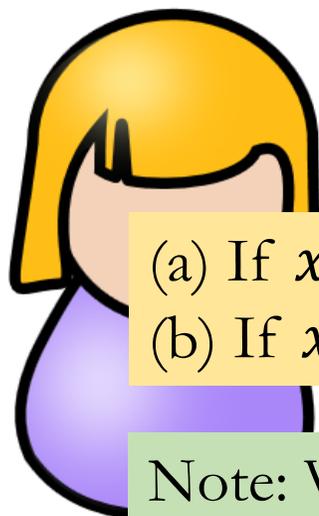
OT



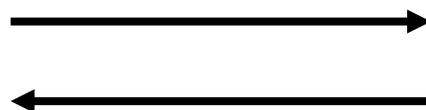
Input:  $y_1, y_2, \dots, y_n$



# Multi-Point OPRF

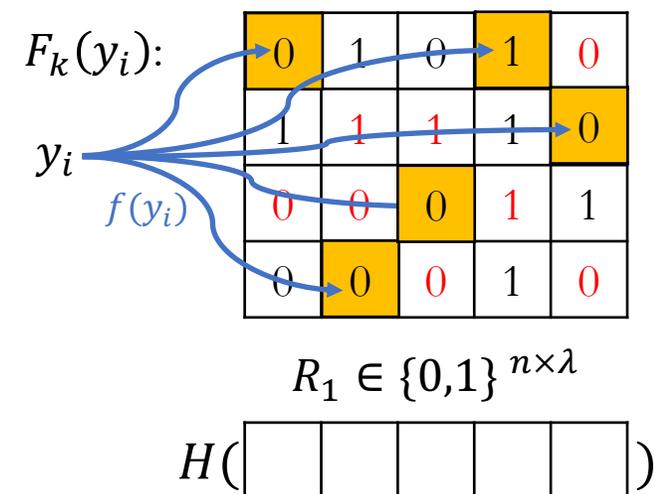
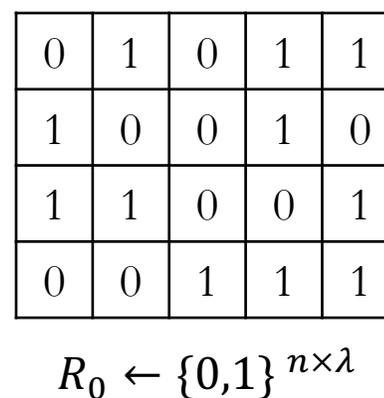
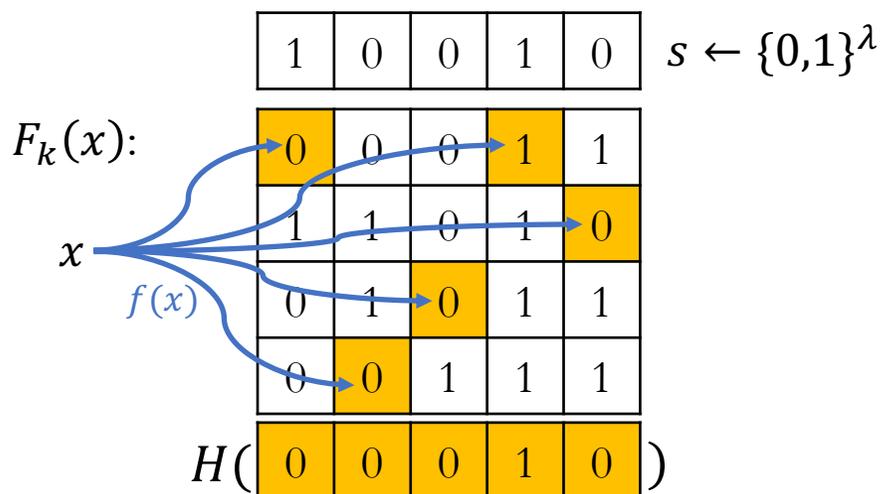


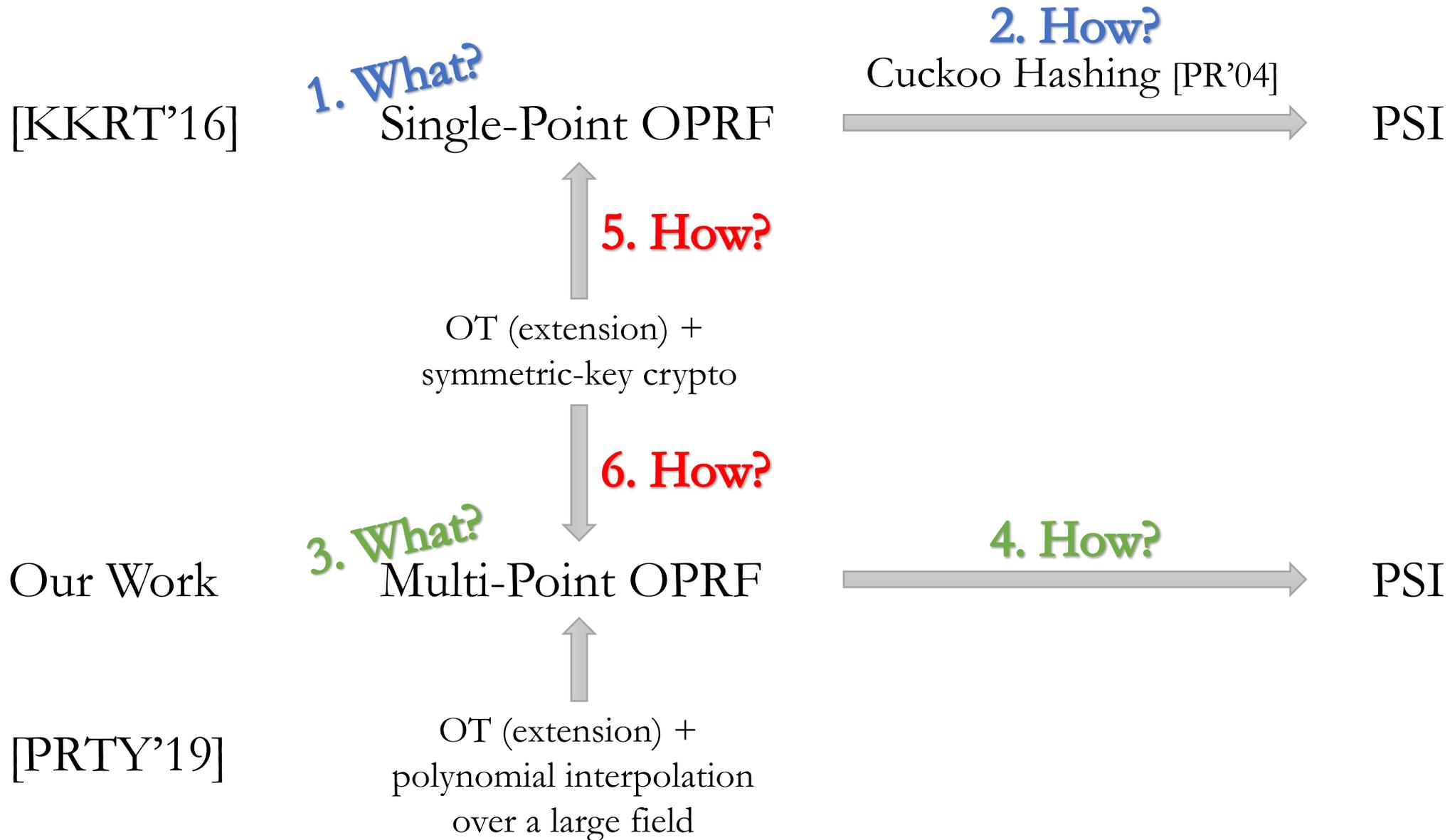
OT



- (a) If  $x = y_i$ , then  $F_k(x) = F_k(y_i)$  no matter what  $s$  is chosen.
- (b) If  $x \notin Y$ , then  $F_k(x)$  is hard to guess.

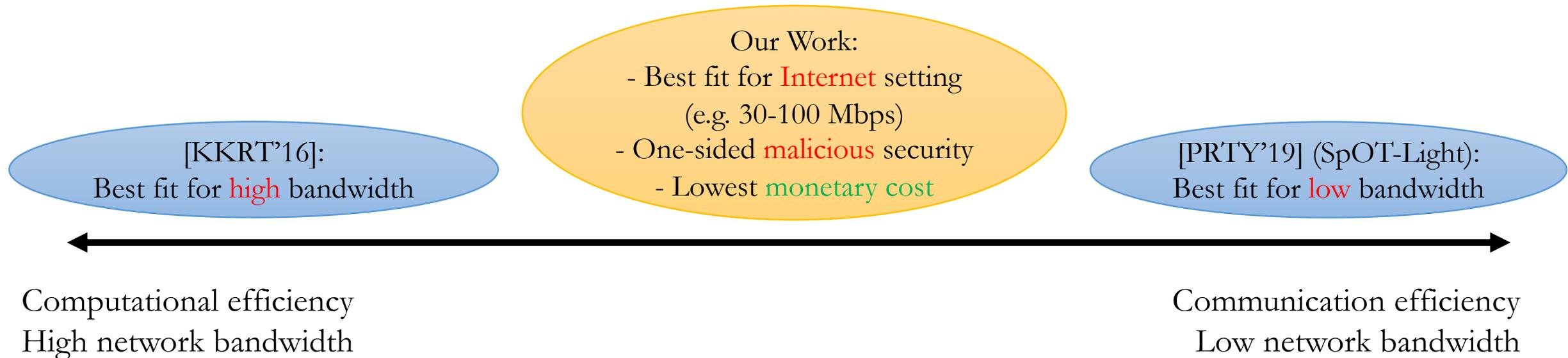
Note: We can prove security against malicious Alice.





# Open Problems

- Best computation & communication?
- Malicious security against Bob?



Thank you!