

# Raccoon Attack: Finding and Exploiting Most-Significant-Bit-Oracles in TLS-DH(E)

Robert Merget,<sup>1</sup> Marcus Brinkmann,<sup>1</sup> Nimrod Aviram,<sup>2</sup>  
Juraj Somorovsky,<sup>3</sup> Johannes Mittmann,<sup>4</sup> Jörg Schwenk<sup>1</sup>

## Real World Crypto 2021

<sup>1</sup> Ruhr University Bochum

<sup>2</sup> Department of Electrical Engineering, Tel Aviv University

<sup>3</sup> Paderborn University

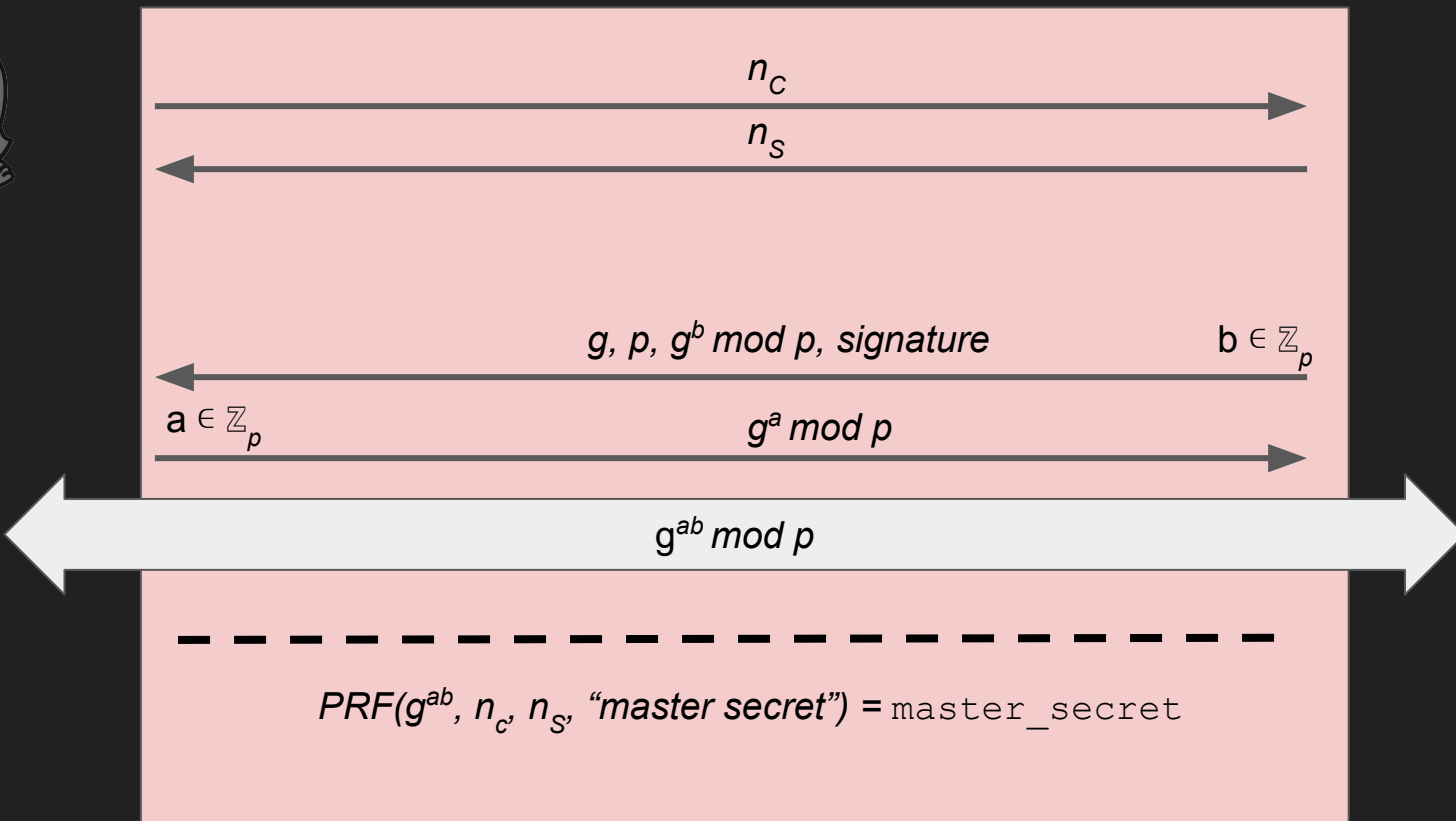
<sup>4</sup> Bundesamt für Sicherheit in der Informationstechnik (BSI), Germany



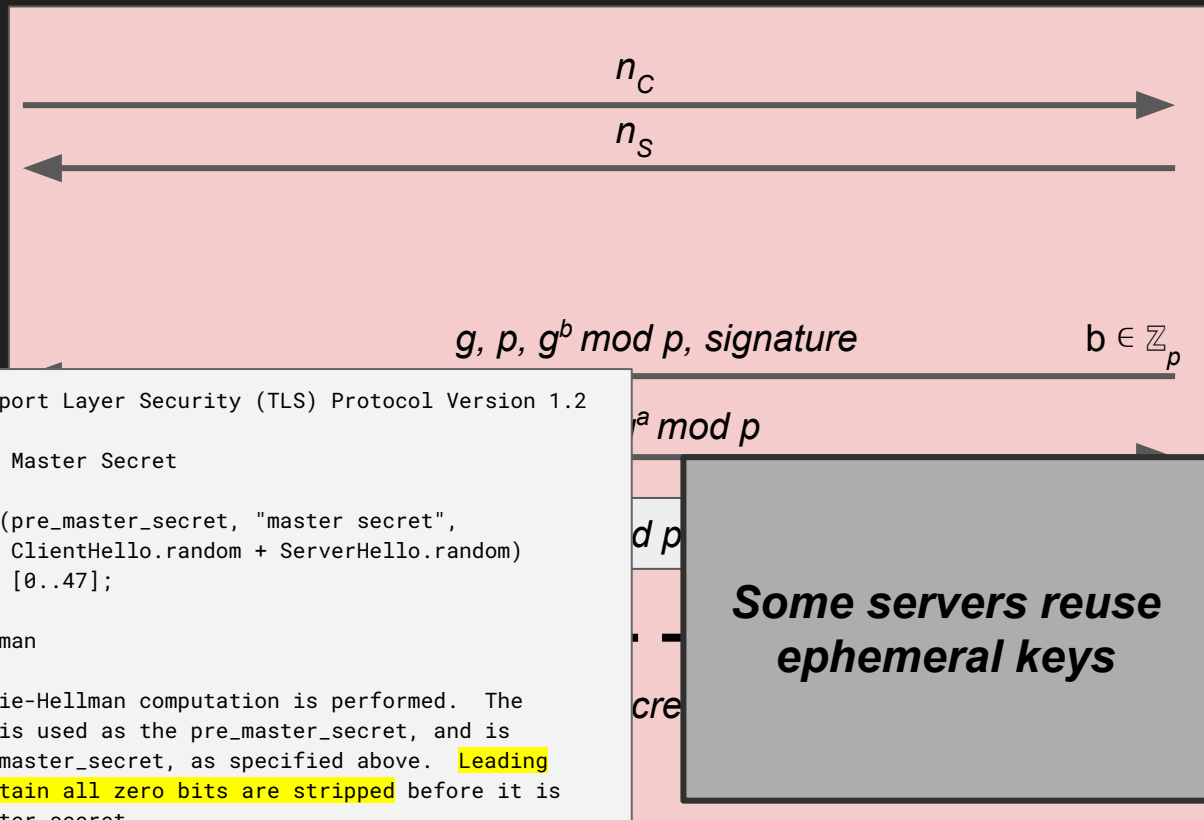
# Transport Layer Security



# TLS-DH(E)



# TLS-DH(E)



RFC 5246: The Transport Layer Security (TLS) Protocol Version 1.2

## 8.1. Computing the Master Secret

```
master_secret = PRF(pre_master_secret, "master secret",
                    ClientHello.random + ServerHello.random)
                    [0..47];
```

### 8.1.2. Diffie-Hellman

A conventional Diffie-Hellman computation is performed. The negotiated key ( $Z$ ) is used as the `pre_master_secret`, and is converted into the `master_secret`, as specified above. **Leading bytes of  $Z$  that contain all zero bits are stripped** before it is used as the `pre_master_secret`.

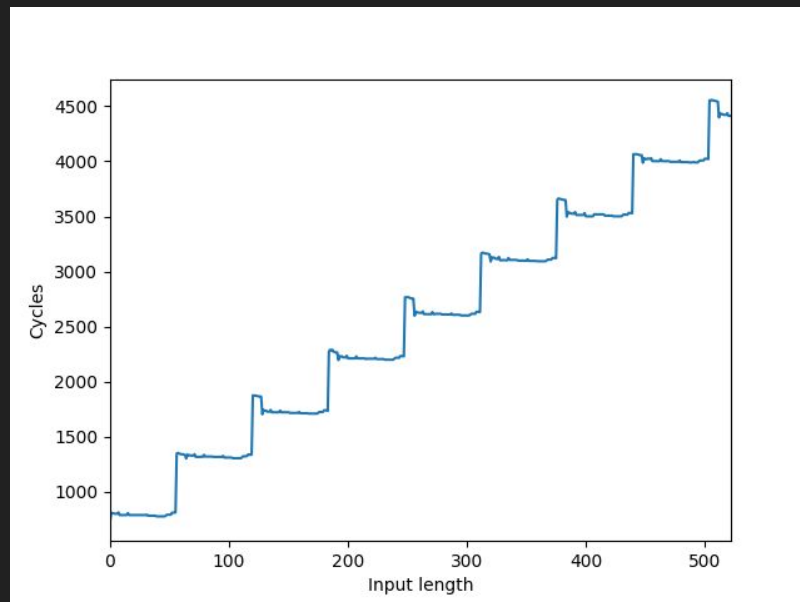
# Key derivation and Constant Time Execution

TLS key derivation is based on hash functions

Hash functions operate in  $O(n)$  not  $O(1)$

This creates various side-channels:

- Compression function invocation
- Hash function invocation
- Key padding
- Direct side-channel



Example: SHA-256 in OpenSSL

# From Side-channel to Exploit



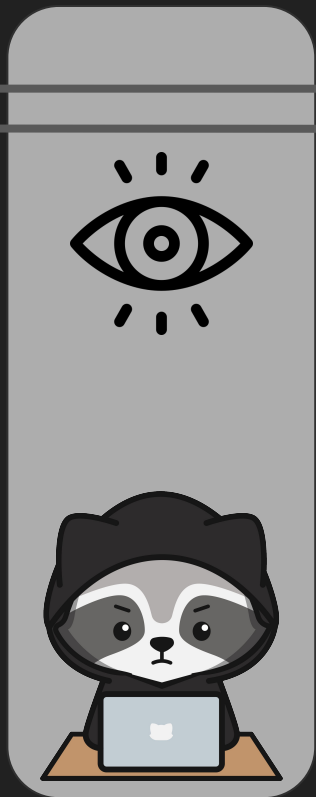
$$\text{MSB}_k(g^{ab}) = 0?$$

$\text{CKE}(g^a), \text{CCS}, \text{FIN}_{\text{Inv}}$

ALERT



# Attack Overview



$g^b$

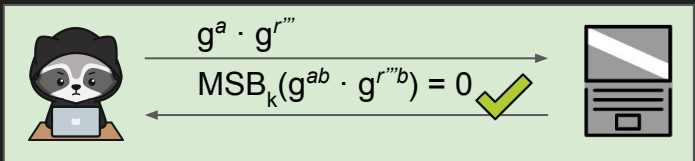
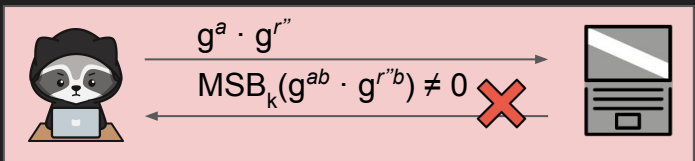
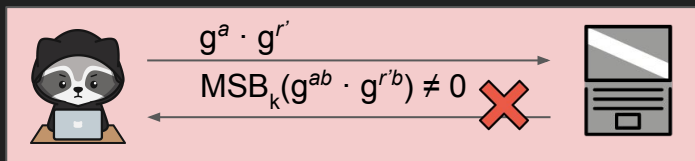
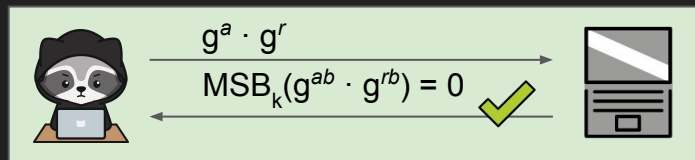
$g^a$

$g^a \cdot g^r$



$PMS = (g^a \cdot g^r)^b$   
 $= g^{ab} \cdot g^{rb}$

# Retrieving the PMS



...

Constructing Instance of Hidden Number Problem:  
 $\alpha = g^{ab}, t_i = g^{r_i b}, 0 < y_i < 2^{n-k}$

$$\alpha \cdot t_1 = y_1 \pmod{p},$$

$$\alpha \cdot t_2 = y_2 \pmod{p},$$

$$\alpha \cdot t_3 = y_3 \pmod{p}, \dots$$

HNP  
Solver

$$\alpha = g^{ab}$$



# Performance

DH Group	Bit length	k=24	k=20	k=16	k=12	k=8
<b>RFC 5114</b>	<b>1024</b>	d=50 T=6s	d=60 T=10s	d=80 T=26s	d=100 T=111s	d=200 T=9295s
<b>LibTomCrypt</b>	<b>1036</b>	d=50 T=6s	d=60 T=10s	d=80 T=28s	d=100 T=52s	d=180 T=5613s
<b>SKIP</b>	<b>2048</b>	d=100 T=112s	d=120 T=207s	d=160 T=977s	Unsolved	Unsolved

d = Number of equations required  
T = Time required to solve HNP  
k = Leading Zero bits leaked

# Impact

## Scan of Alexa Top 100k:

- 32% of the scanned servers supported DHE cipher suites
- 10.9% of those servers reused their ephemeral keys

Firefox was the last to drop support in September 2020

No major browser supports DHE anymore



# Countermeasure

## Generally:

- Do not leak partial information about secret values
- Make secrets constant size

## For TLS:

- Clients should avoid DH(E)
- Servers should not reuse ephemeral keys
- Servers and clients should not use DH



# Raccoon and ECDH(E)

TLS does not strip leading zero bytes of shared ECDH secrets

Requires implementation specific side-channels

Further research required, currently not exploitable



# Raccoon and TLS 1.3

TLS 1.3 does not strip leading zero bytes of shared secrets

Foresight by David Benjamin in Draft-13 proved useful

Ephemeral key reuse is uncommon



# Conclusion

- No need to panic, exploitation is difficult
- The Raccoon attack is not TLS specific
- First time HNP is used to attack DH

More info: <https://raccoon-attack.com>

Tool to scan your servers:

<https://github.com/tls-attacker/TLS-Scanner>



@ic0nz1



[robert.merget@rub.de](mailto:robert.merget@rub.de)

