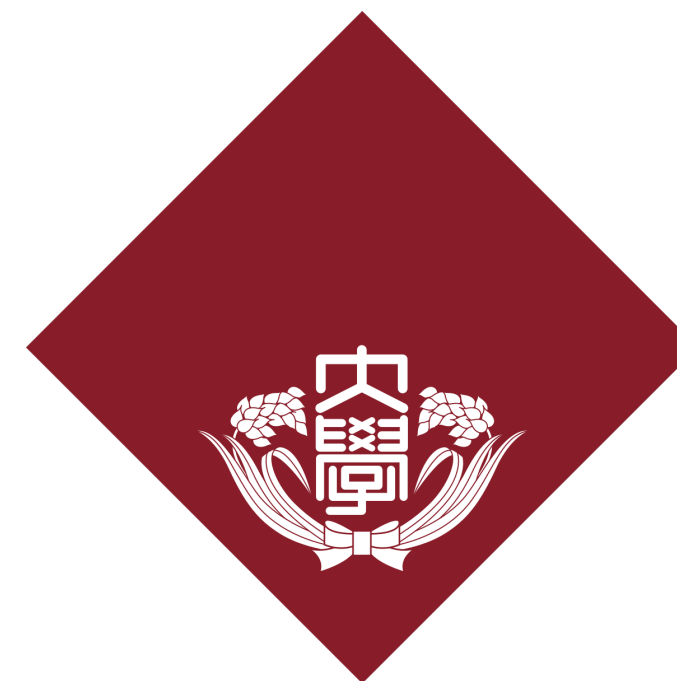


# Compact FE for Unbounded Attribute-Weighted Sums for Logspace from SXDH

**Pratish Datta**  
NTT Research

**Tapas Pal**  
NTT SIL

**Katsuyuki Takashima**  
Waseda University



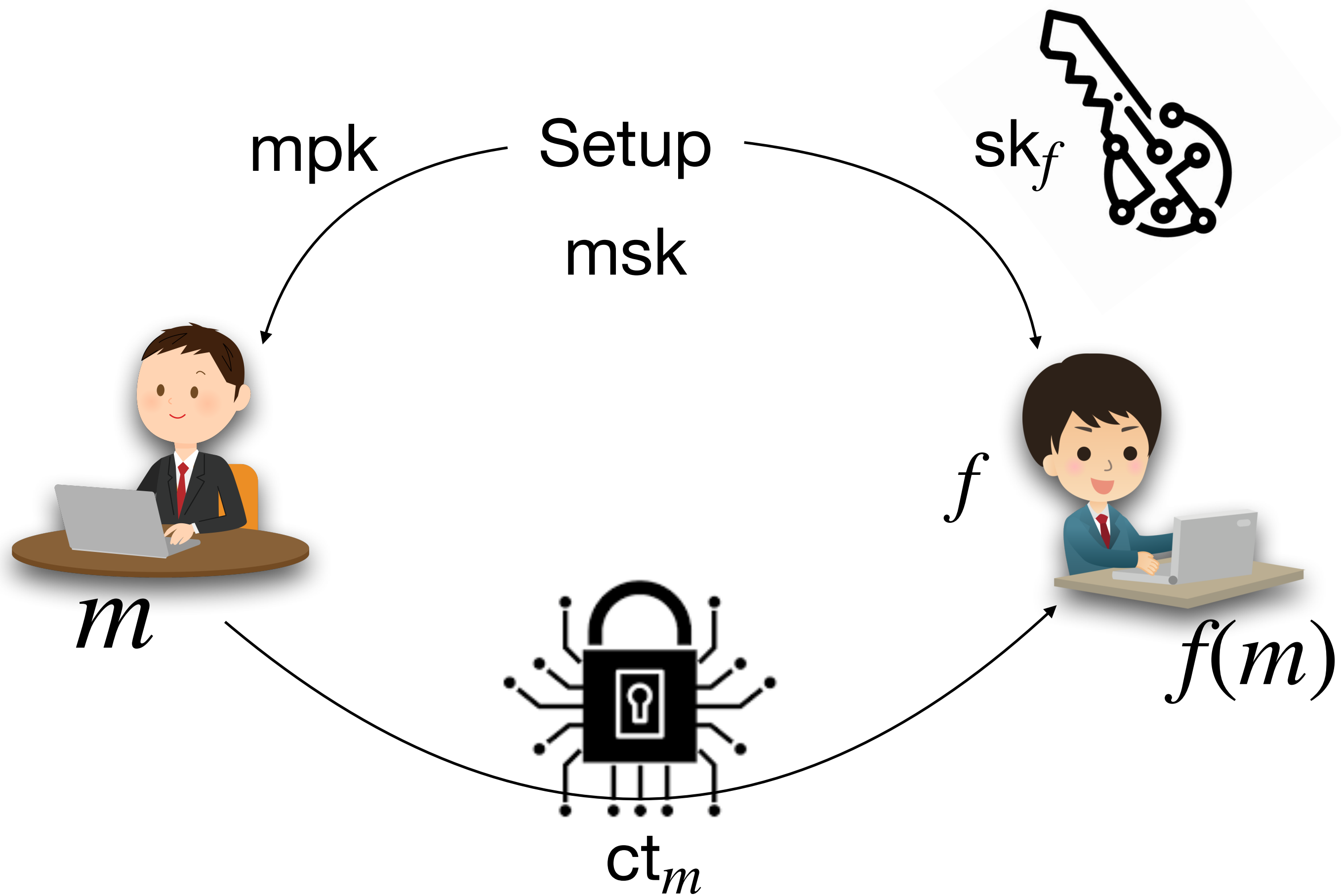
# Functional Encryption [BSW11]

$$\text{Setup}(1^\lambda) \rightarrow (\text{mpk}, \text{msk})$$

$$\text{KeyGen}(\text{msk}, f) \rightarrow \text{sk}_f$$

$$\text{Enc}(\text{mpk}, m) \rightarrow \text{ct}_m$$

$$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$$



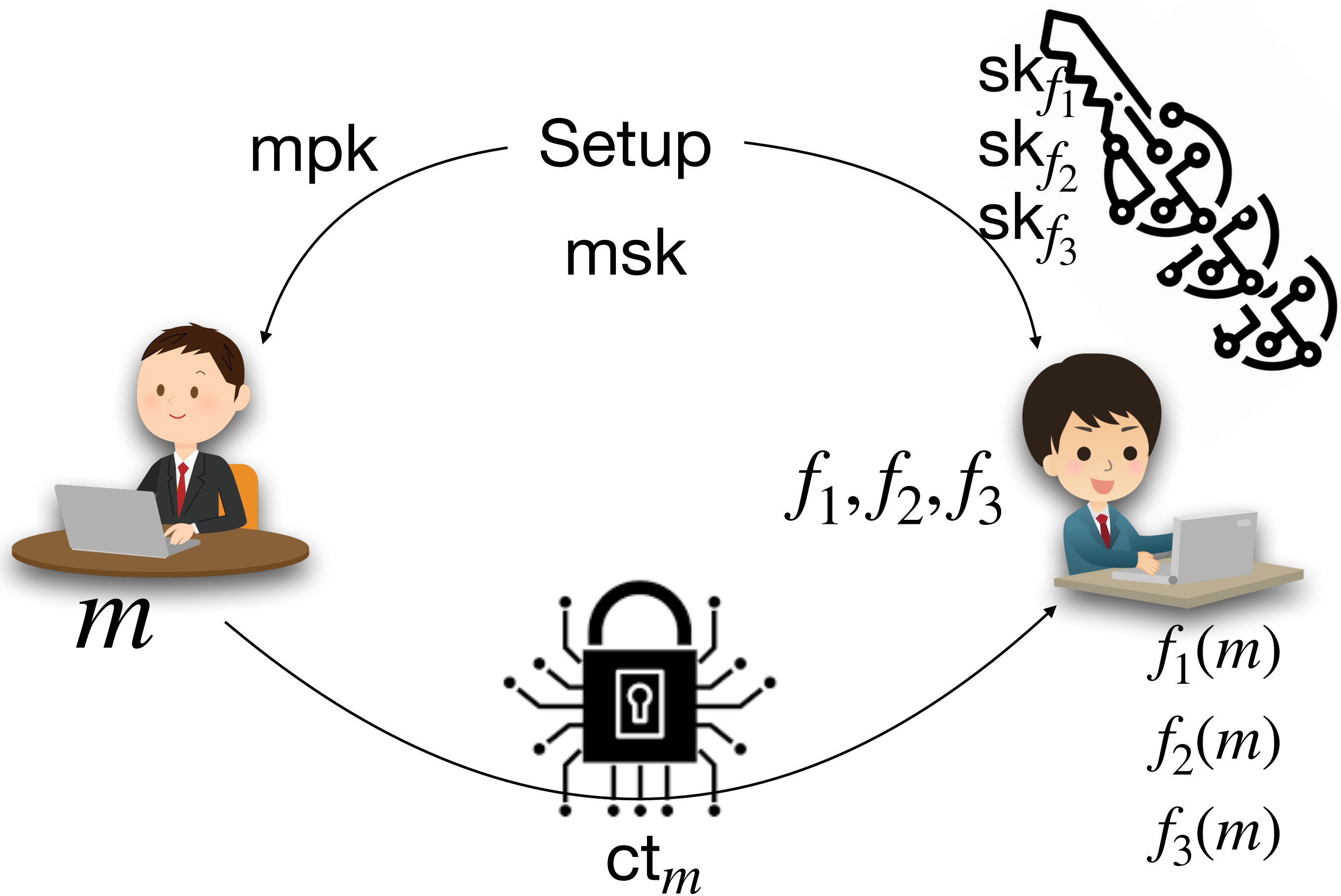
# Simulation Security of FE [O’Niell11,BSW11]

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# Adaptive Simulation Security

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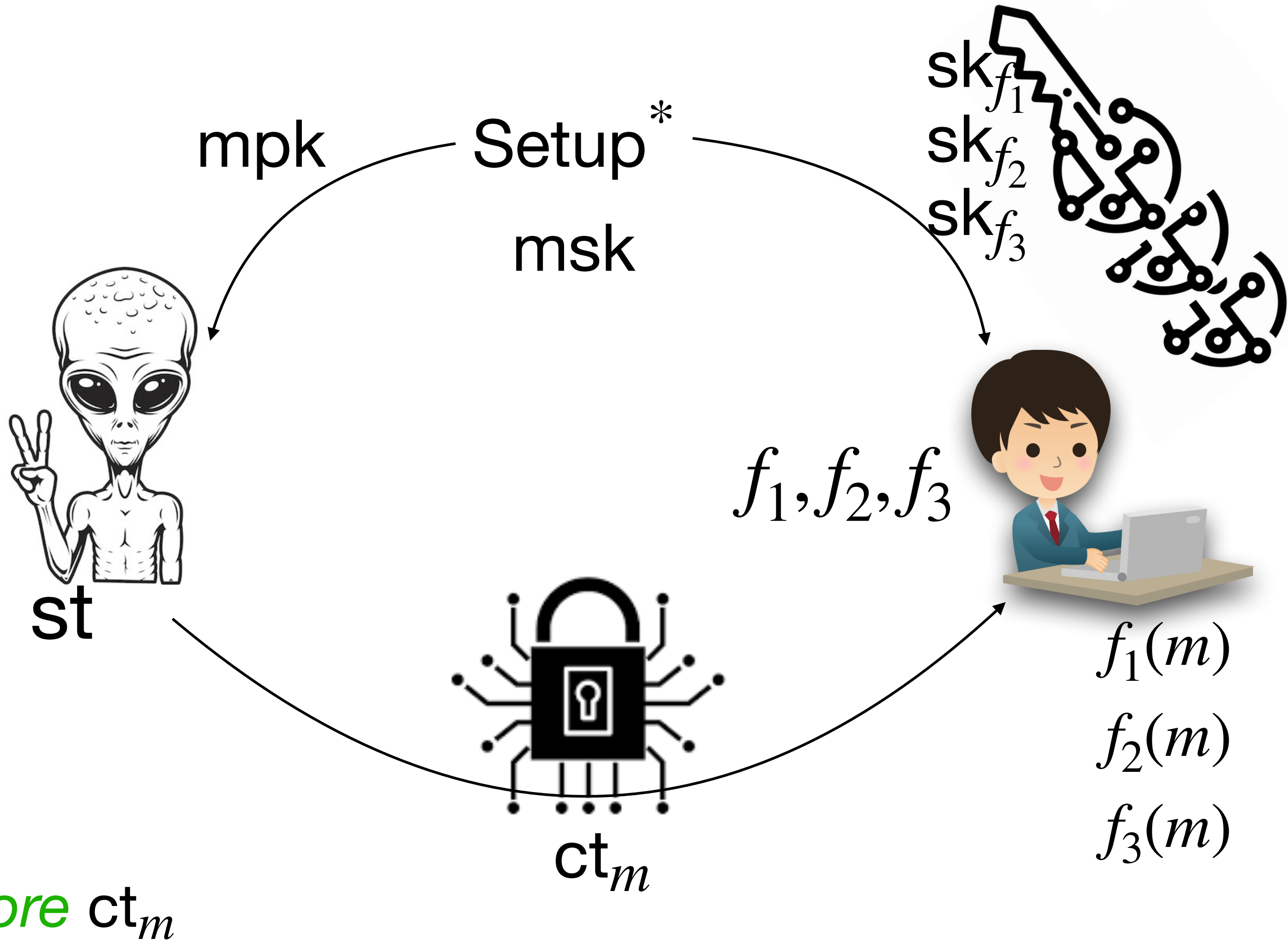
$$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$$

---


$$\text{Setup}^*(1^\lambda) \rightarrow (\text{mpk}, \text{msk}^*)$$

$$\text{KeyGen}_0^*(\text{msk}^*, f) \rightarrow \text{sk}_f$$

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*before*  $ct_m$

$st$  contains  $\{f_i(m)\}_i$  for already queried  $\{f_i\}_i$

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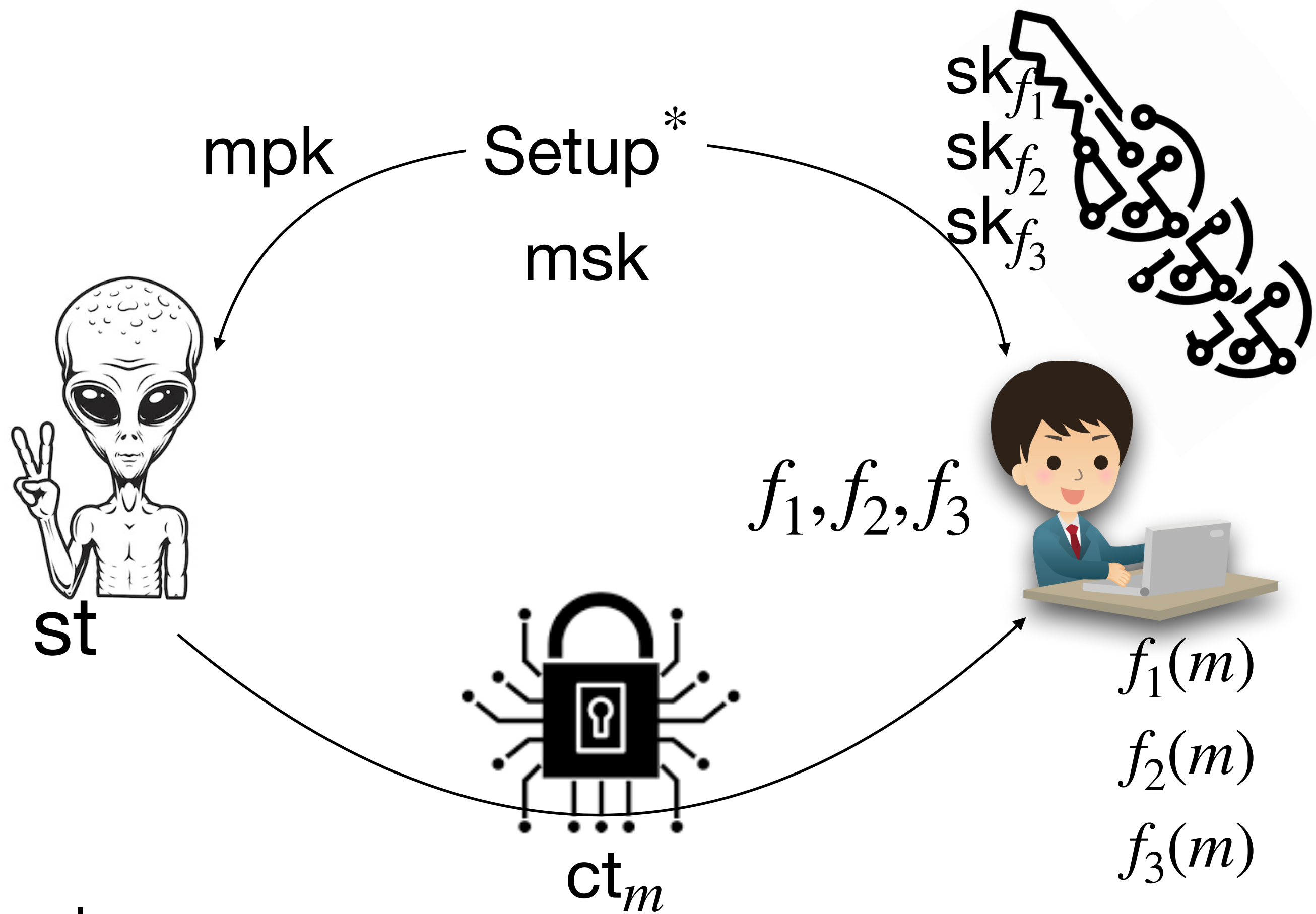
$$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$$

---


$$\text{Setup}^*(1^\lambda) \rightarrow (\text{mpk}, \text{msk}^*)$$

$$\text{KeyGen}_1^*(\text{msk}^*, f, f(m)) \rightarrow \text{sk}_f$$

$$\text{Enc}^*(\text{mpk}, \text{msk}^*, \text{st}) \rightarrow \text{ct}_m$$



after  $\text{ct}_m$

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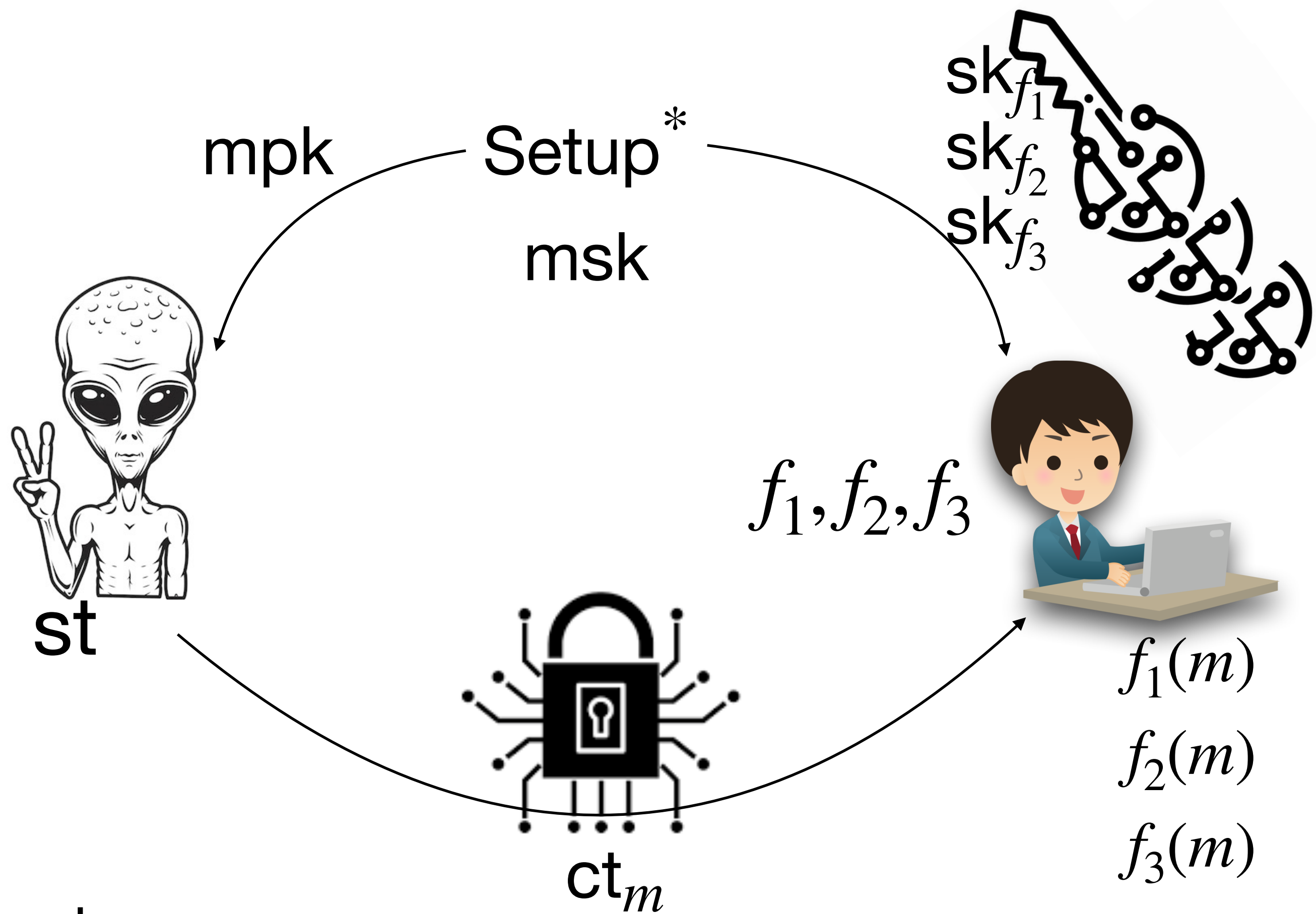
$$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$$

$\approx_c$

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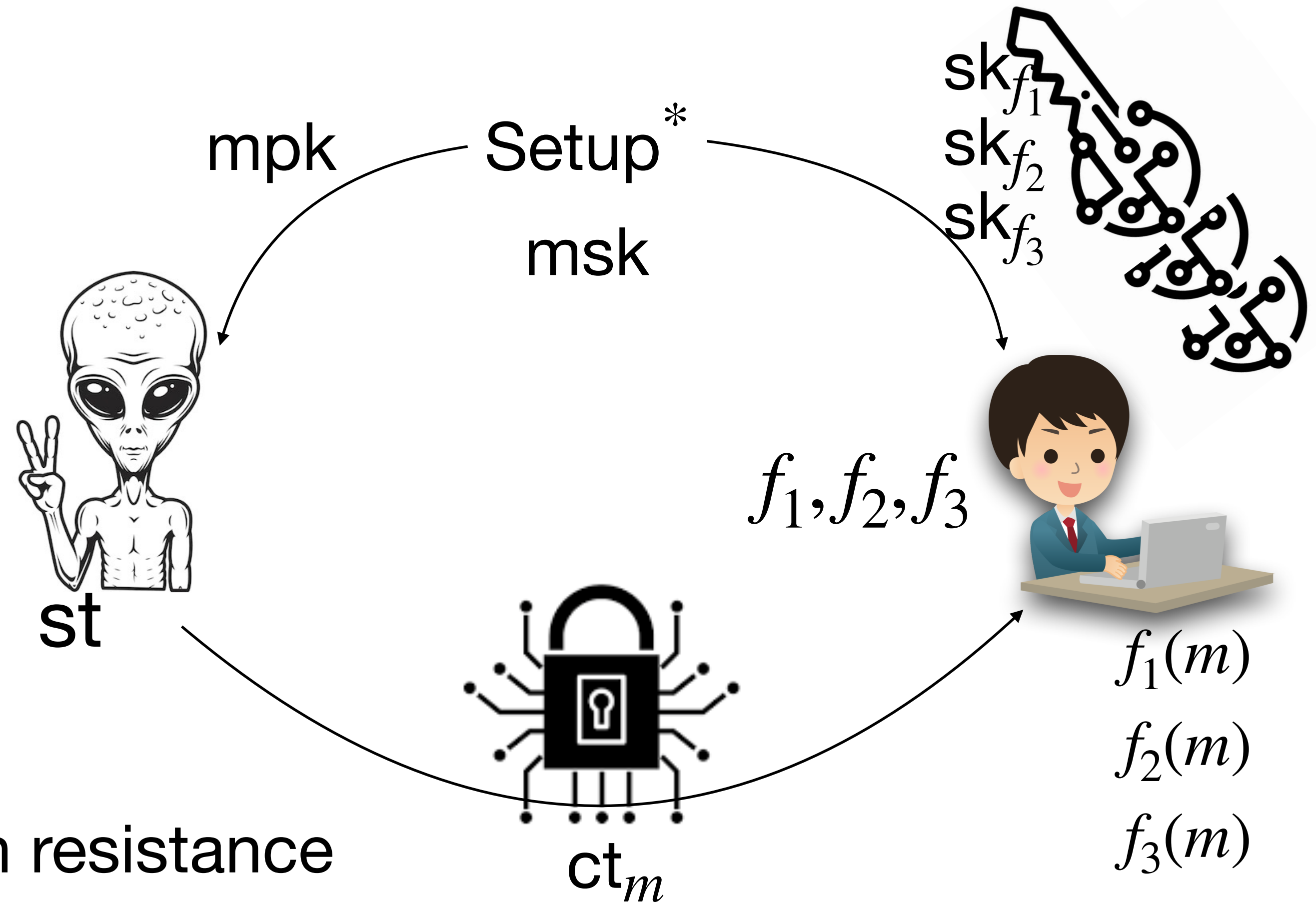
# Adaptive Simulation Security: Bounded or Full Collusion resistance

$\text{Setup}(1^\lambda) \rightarrow (\text{mpk}, \text{msk})$

$\text{KeyGen}(\text{msk}, f) \rightarrow \text{sk}_f$

$\text{Enc}(\text{mpk}, m) \rightarrow \text{ct}_m$

$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$



- **bounded**  $\#\text{sk}_{f_i} \rightarrow$  **bounded** collusion resistance
- **any**  $\#\text{sk}_{f_i} \rightarrow$  **full** collusion resistance

# FE for Various Function Classes

---

General Class:  
TMs or All Circuits

---

[GKP+13,GGG+14,BGJS15,AJ15,  
BKS16,AR17,AM18,AMVY21,JLS22....]

- **Inefficient** and **complex**
- **bounded** collusion-resistant
- Assumptions: **IO** or **SubExp** LWE
- **Not** yet practical



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## Specific Class: Linear, Quadratic or its variants

---

[ABDP15,ALS16,BCFG17,TT18,G20,  
GQ20,Wee20,ACGU20,AGW20,DP21,MKMS22....]

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- **Full** collusion-resistant
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- 
- **This work advances FE for a specific class**

# Functional Encryption for Attribute-Weighted Sums[AGW20]

$\text{Setup}(1^\lambda, 1^n, 1^m) \rightarrow (\text{mpk}, \text{msk})$

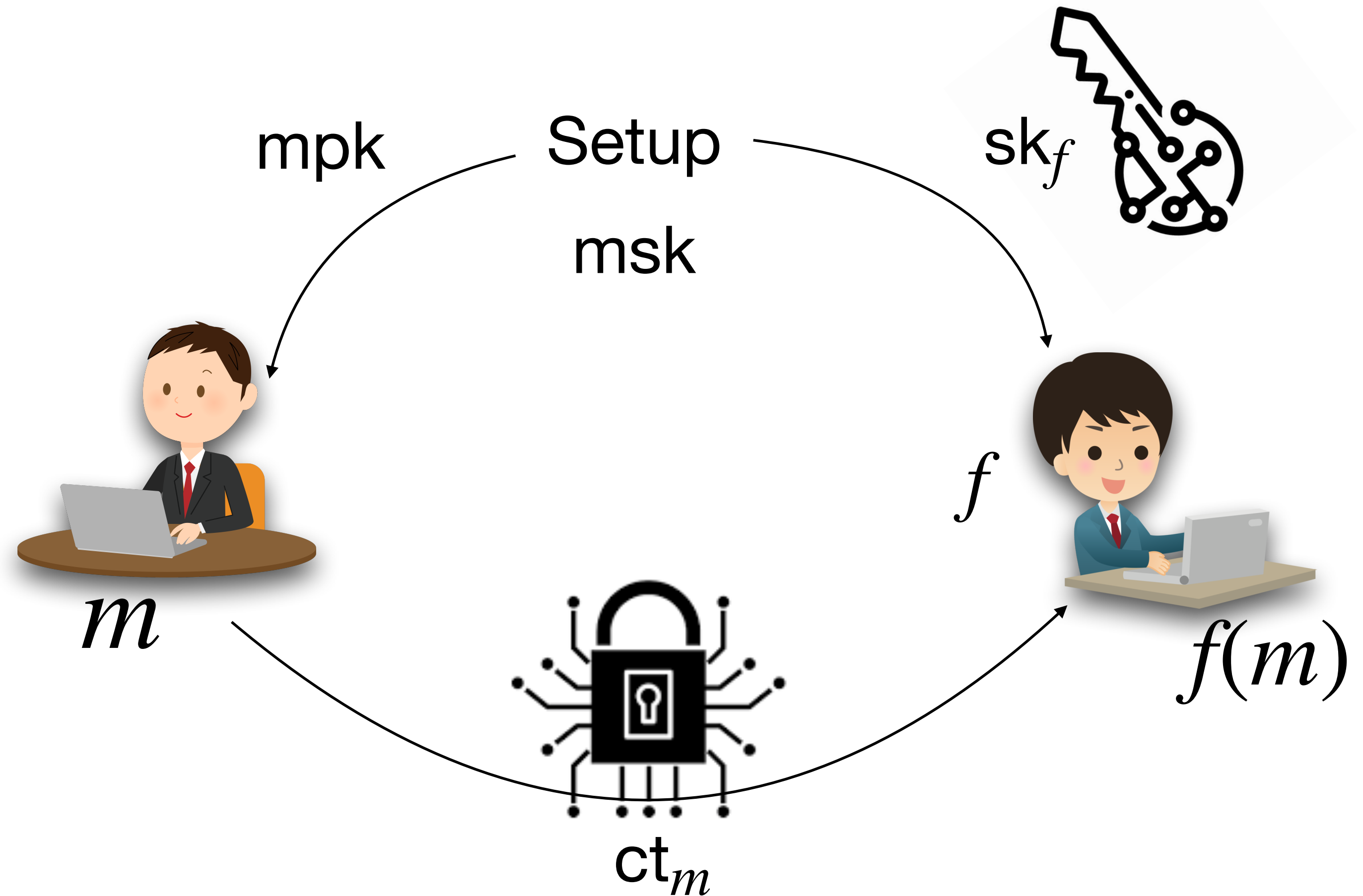
$\text{KeyGen}(\text{msk}, f) \rightarrow \text{sk}_f$

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$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$

- Function:  $f : \mathbb{Z}_p^m \rightarrow \mathbb{Z}_p^n$
- Message:  $m = (\mathbf{x}, \mathbf{z}) \in \mathbb{Z}_p^{m \times n}$
- Output:  $f(m) = f(\mathbf{x}) \cdot \mathbf{z}$

$\mathbf{x}$  is public,  $\mathbf{z}$  is private



# Functional Encryption for Attribute-Weighted Sums (AWS)

---

$$f(m) = f(\mathbf{x}) \cdot \mathbf{z}, \quad f \in \text{ABP}$$

## Prior Works

- Function class = ABP
- Setup  
 $|\text{mpk}| = O(|\mathbf{x}|, |\mathbf{z}|)$
- $|\text{ct}_m| = O(|\mathbf{z}|)$
- AD-SIM [DP21]  
 $|\text{ct}_m| = O(|\mathbf{x}|, |\mathbf{z}|)$

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## Applications

- IPFE:  $f(\mathbf{x}) = \mathbf{y}$   
 $f(m) = \mathbf{y} \cdot \mathbf{z}$
- ABE:  $f(\mathbf{x}) = 1/0$ ,  $\mathbf{z} = M$   
 $f(m) = M$  if  $f$  satisfies  $\mathbf{x}$
- AB-IPFE:  $f(\mathbf{x}) = \mathbf{y}g(\mathbf{x})$   
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## Limitations

- non-uniform, non-dynamic
- bounded Setup  
 $|\text{mpk}| \neq O(\lambda)$
- $|\text{ct}_m| \neq \text{input-specific}$
- bounded FE:  
IPFE, ABE,...

# Functional Encryption for Attribute-Weighted Sums (AWS)

$$f(m) = f(\mathbf{x}) \cdot \mathbf{z}, f \in \text{TM}$$

## Applications

- Function class = TM
- Setup  
 $|\text{mpk}| = O(|\mathbf{x}|, |\mathbf{z}|)$
- $|\text{ct}_m| = O(|\mathbf{z}|)$
- AD-SIM [DP21]  
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- Setup  
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## Limitations

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## Limitations

- uniform, dynamic
- unbounded Setup  
 $|\text{mpk}| = O(\lambda)$
- $|\text{ct}_m| = \text{input-specific}$
- bounded FE:  
IPFE, ABE,...

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$$f(m) = f(\mathbf{x}) \cdot \mathbf{z}, f \in \text{TM}$$

## This Work

- Function class = TM
- Setup  
 $|\text{mpk}| = O(\lambda)$
- $|\text{ct}_m| = O(|\mathbf{z}|, |\mathbf{x}|)$
- AD-SIM

## Applications

- UIPFE:  $f(\mathbf{x}) = \mathbf{y}$   
 $f(m) = \mathbf{y} \cdot \mathbf{z}$
- UABE:  $f(\mathbf{x}) = 1/0$ ,  $\mathbf{z} = M$   
 $f(m) = M$  if  $f$  satisfies  $\mathbf{x}$
- UAB-IPFE:  $f(\mathbf{x}) = \mathbf{y}g(\mathbf{x})$   
 $f(m) = \mathbf{y} \cdot \mathbf{z}$  if  $g$  satisfies  $\mathbf{x}$

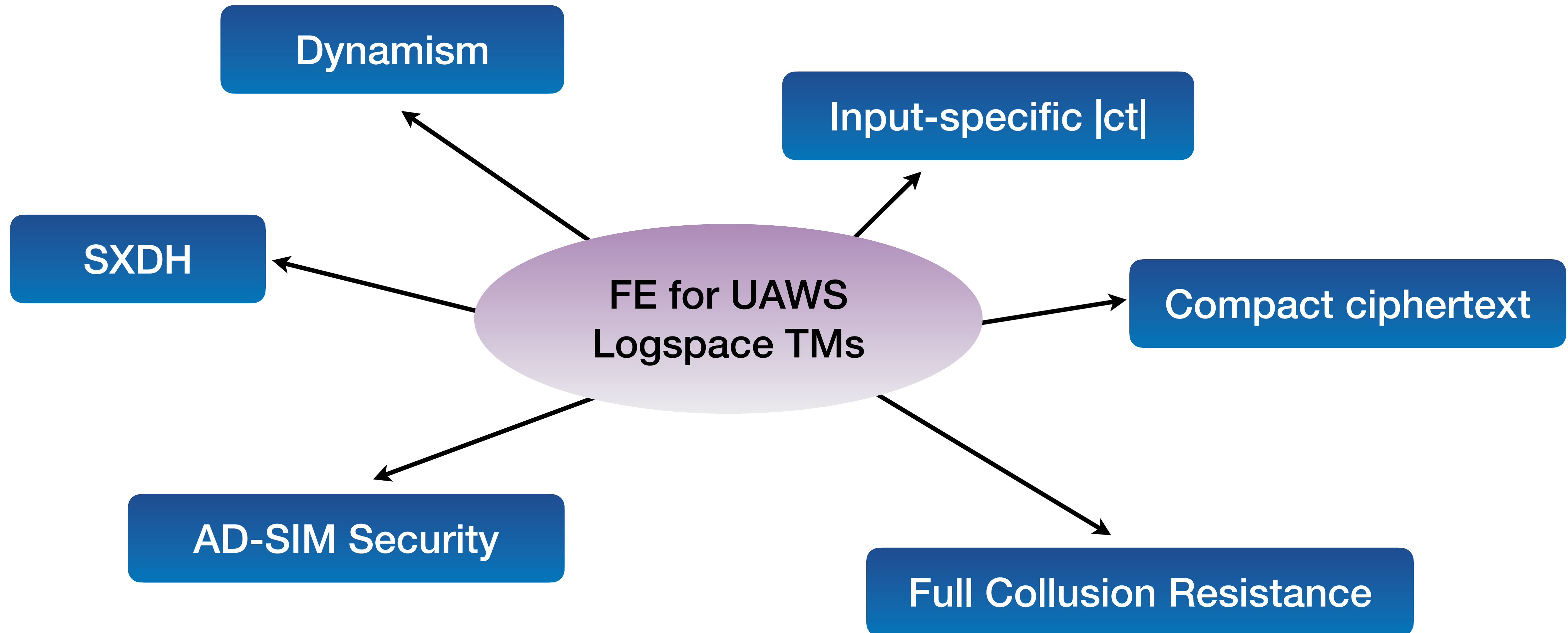
## Limitations

- uniform, dynamic
- unbounded Setup  
 $|\text{mpk}| = O(\lambda)$
- $|\text{ct}_m| = \text{input-specific}$
- Unbounded FE:  
UIPFE, UABE

# Summary of our Results

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- Define FE for Unbounded Attribute-Weighted Sums (**UAWS**)



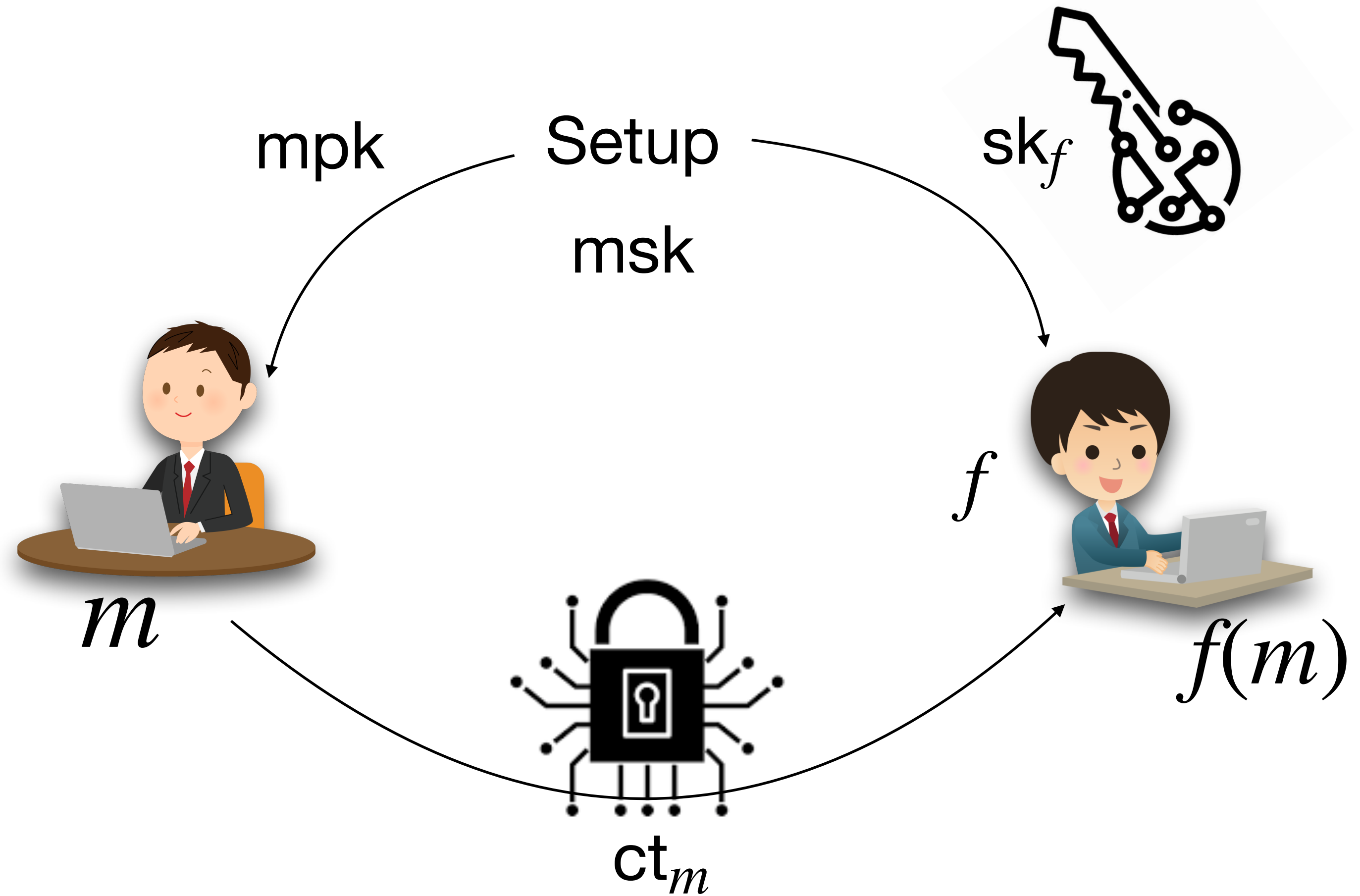
# How to Define: FE for Unbounded AWS (UAWs)

$\text{Setup}(1^\lambda) \rightarrow (\text{mpk}, \text{msk})$

$\text{KeyGen}(\text{msk}, f) \rightarrow \text{sk}_f$

$\text{Enc}(\text{mpk}, m) \rightarrow \text{ct}_m$

$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$



- Function:  $f = (M_k)_{k \in I}$  s.t.  $M_k \in \text{TM}$
- Message:  $m = (\mathbf{x}, \mathbf{z}) \in \{0,1\}^* \times \mathbb{Z}_p^n$
- Output:  $f(m) = \sum_{k \in I} M_k(\mathbf{x})z_k$  iff  $I \subseteq [n]$

$\mathbf{x}$  is public,  $\mathbf{z}$  is private

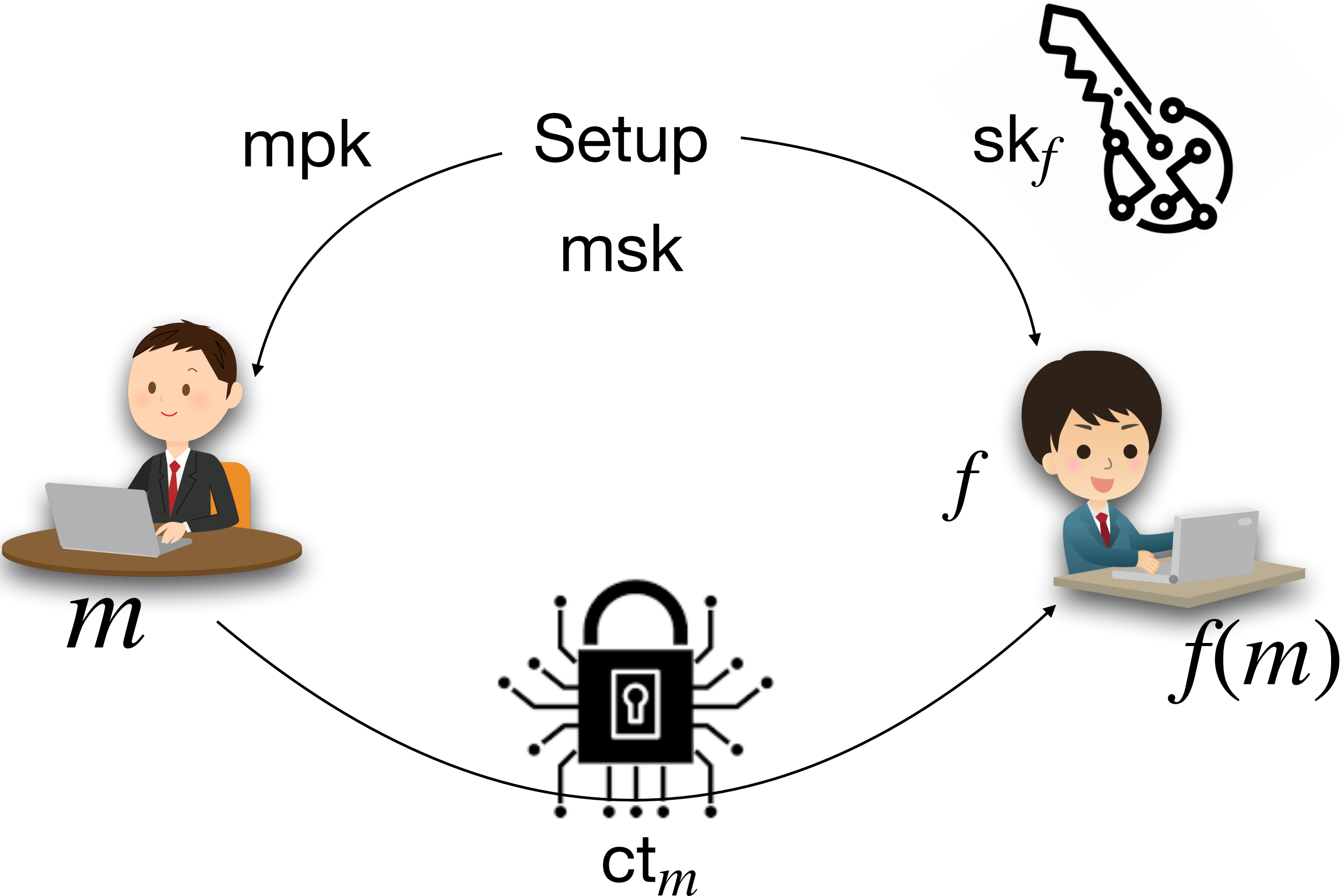
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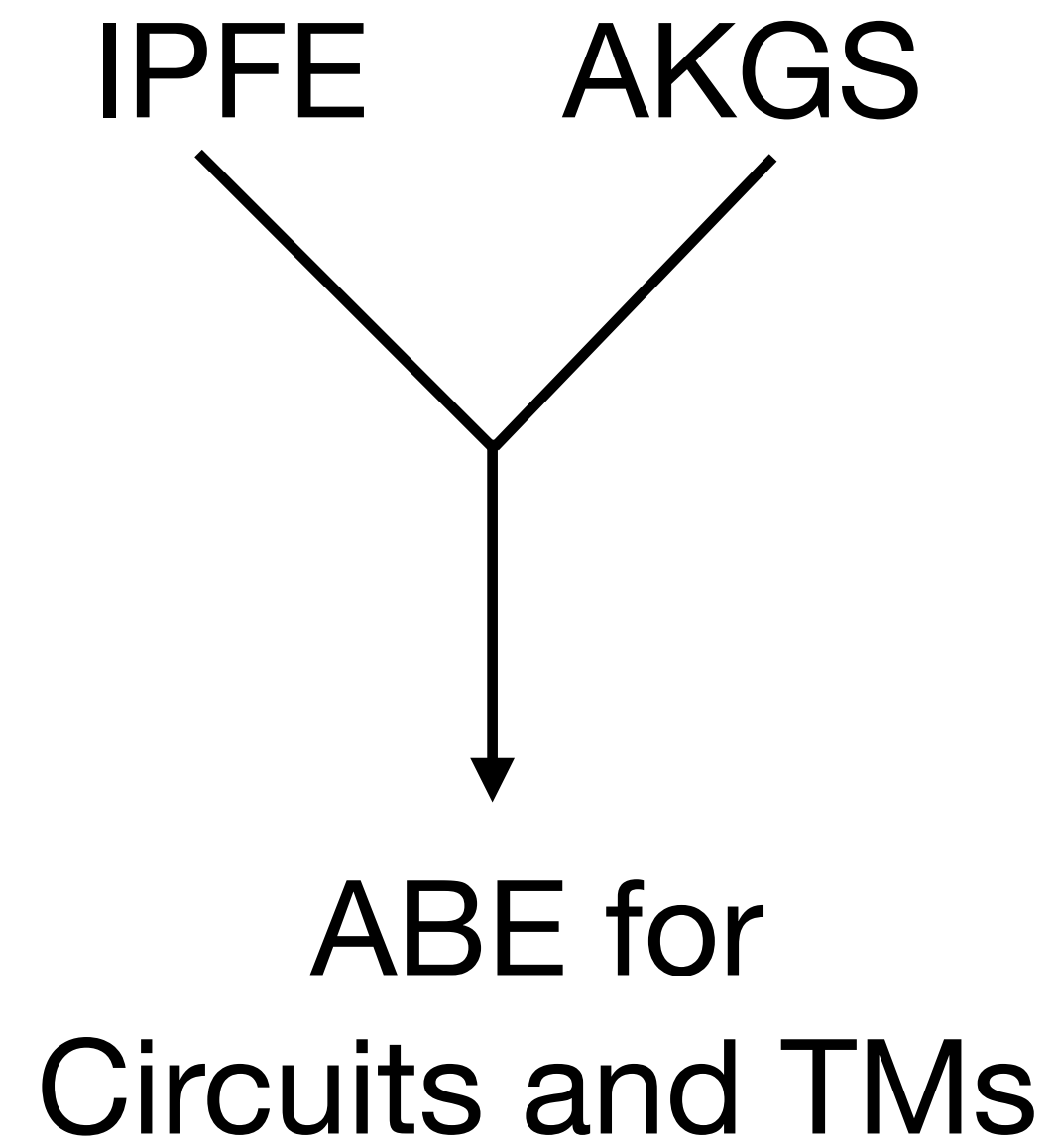
$\mathbf{x}$  is public,  $\mathbf{z}$  is private

$$f = (M_1, M_2, M_4), m = (\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3, z_4, z_5))$$
$$f(m) = M_1(\mathbf{x})z_1 + M_2(\mathbf{x})z_2 + M_4(\mathbf{x})z_4$$

# Roadmap towards FE for UAWS

---

[LL20]

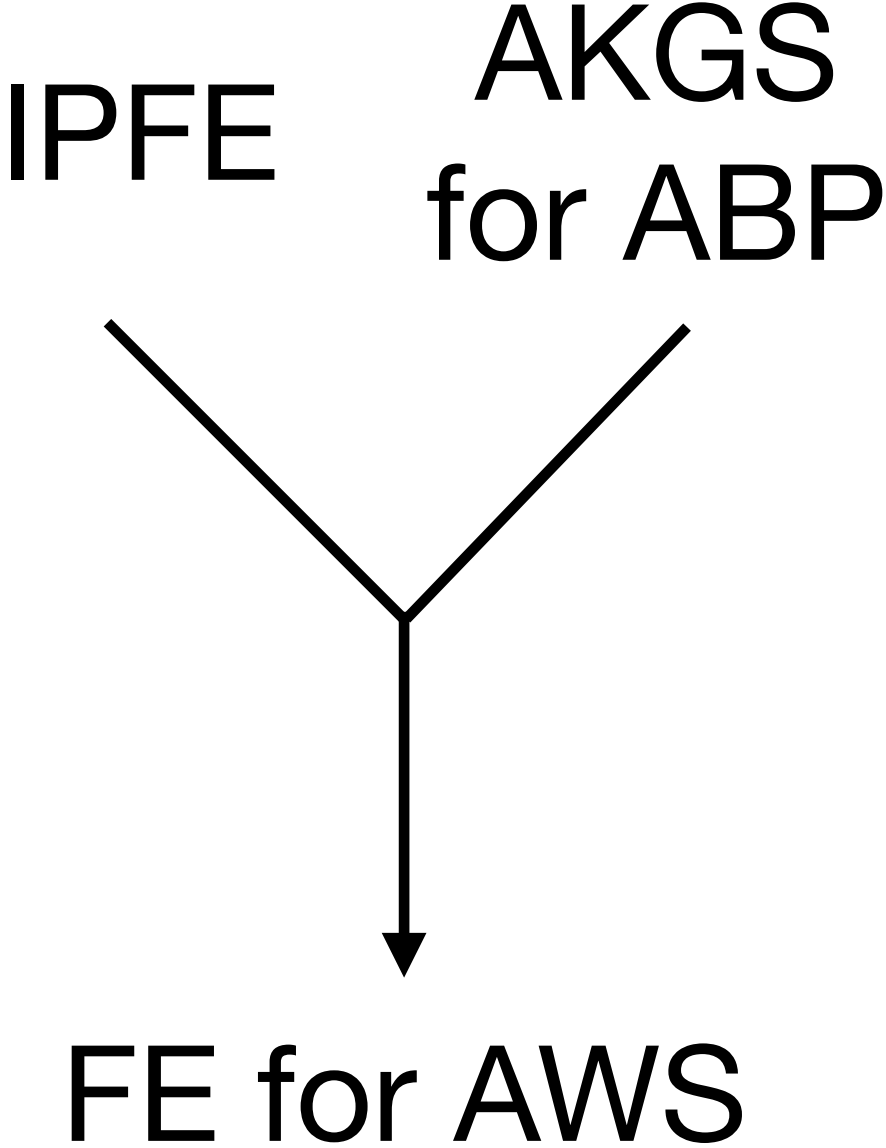
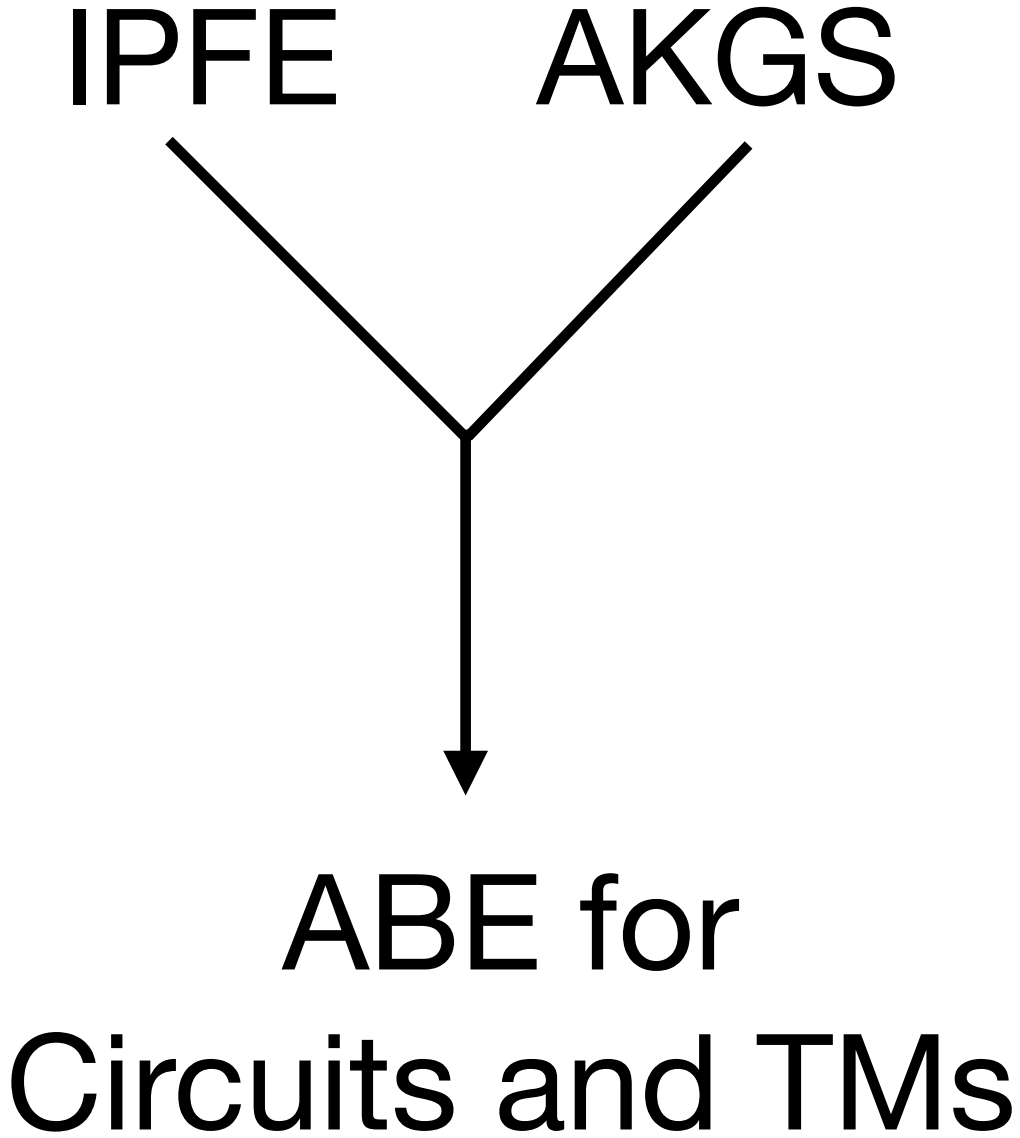


Payload-hiding

# Roadmap towards FE for UAWS

[LL20]

[DP21]



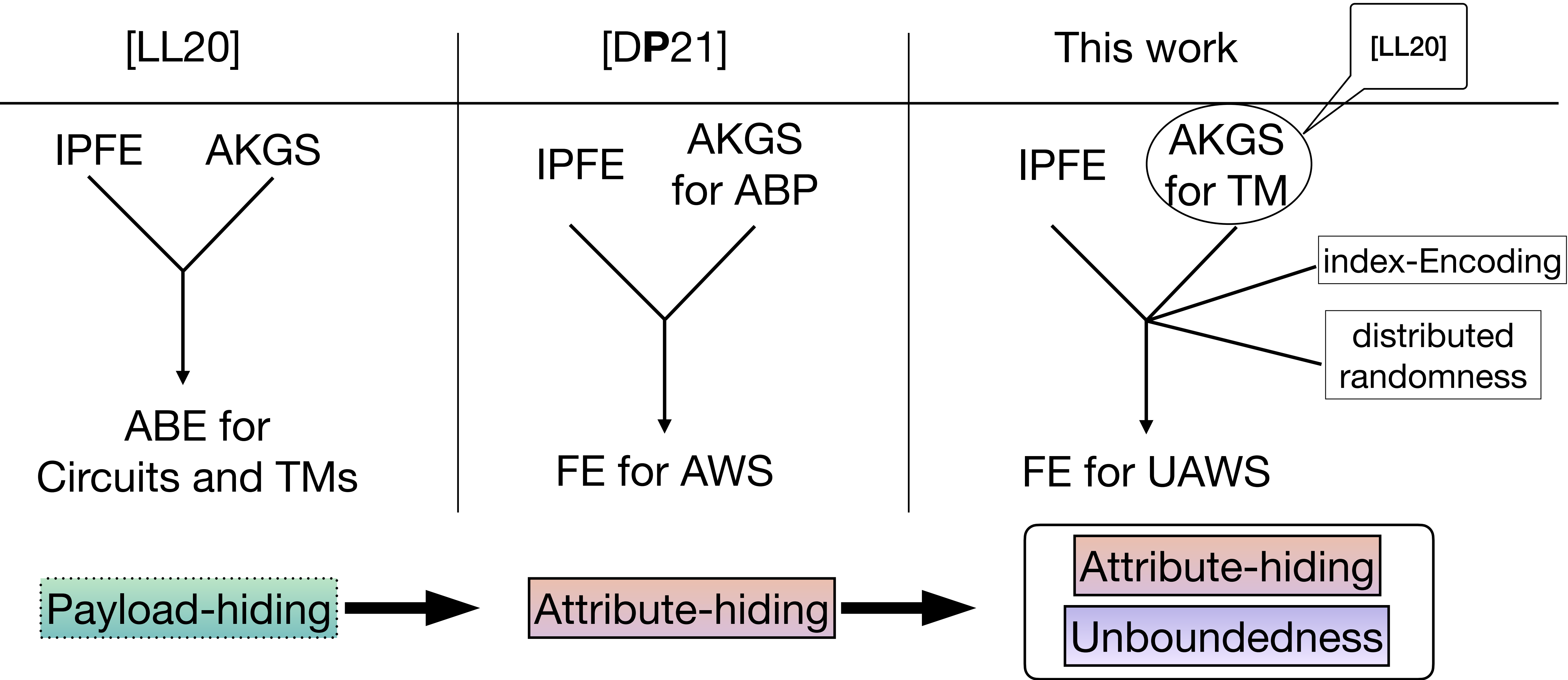
Payload-hiding



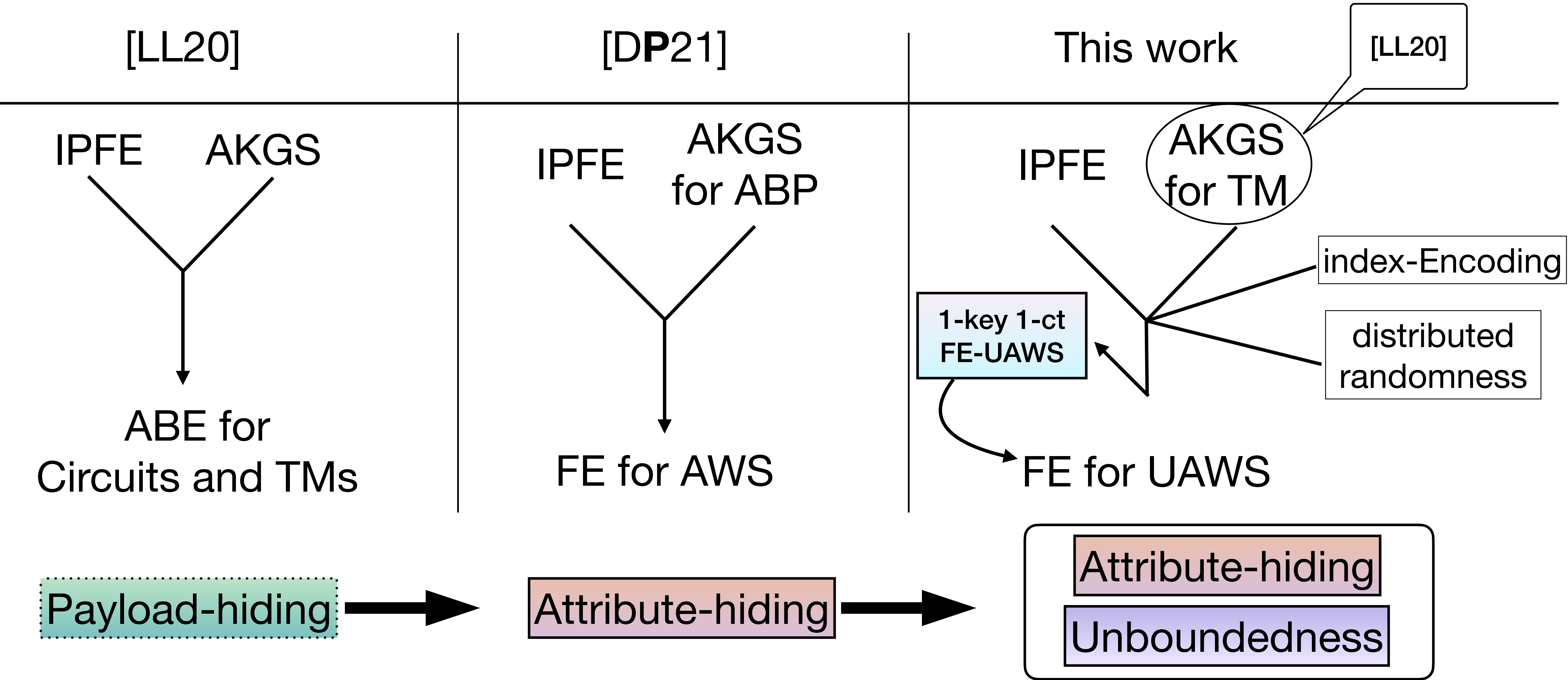
Attribute-hiding



# Roadmap towards FE for UAWS



# Roadmap towards FE for UAWS



# Inner Product Functional Encryption (IPFE)

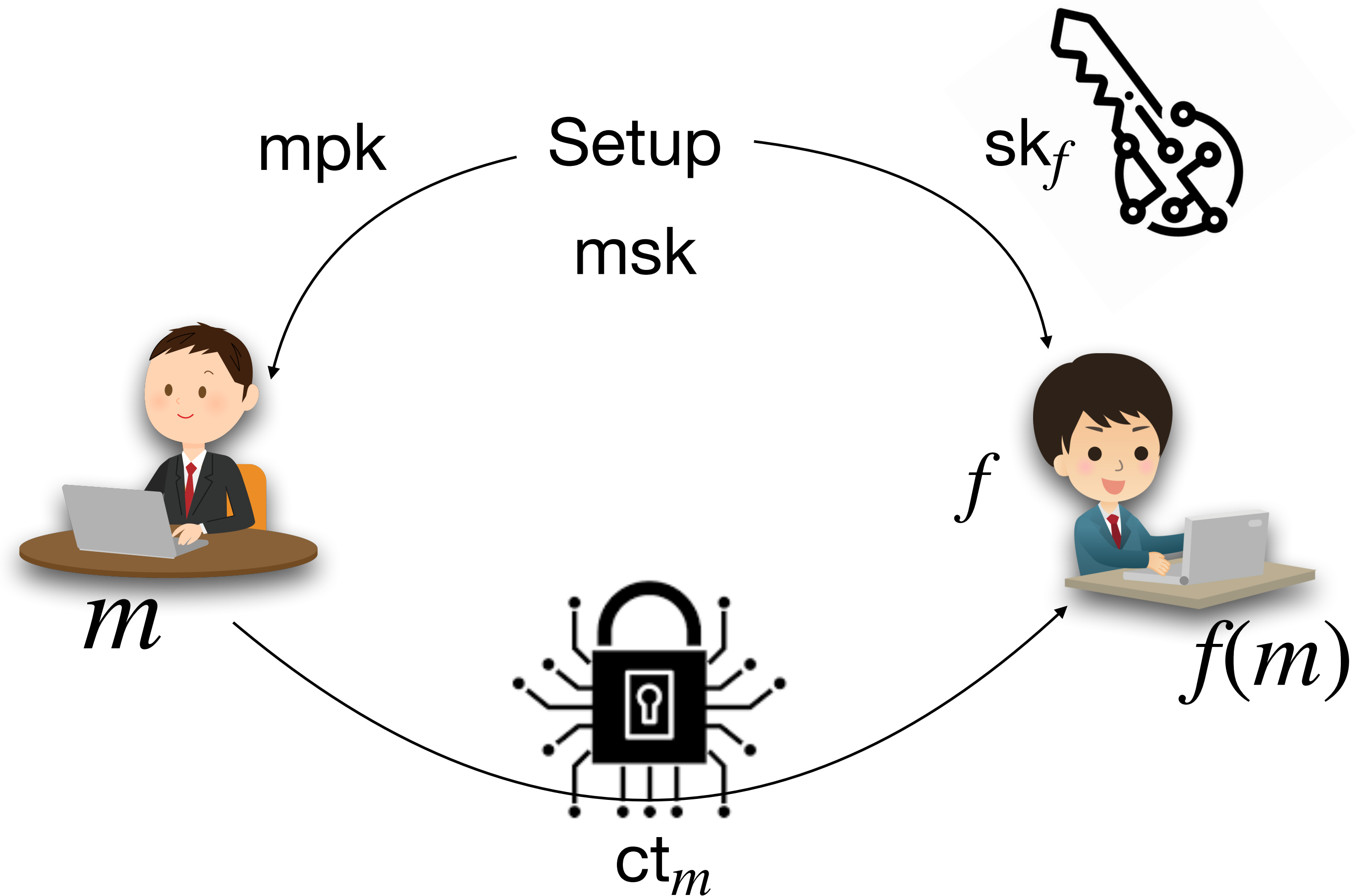
$\text{Setup}(1^\lambda, 1^n) \rightarrow \text{msk}$

$\text{KeyGen}(\text{msk}, f) \rightarrow \text{sk}_f$

$\text{Enc}(\text{msk}, m) \rightarrow \text{ct}_m$

$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$

- Functions:  $f = \mathbf{y} \in \mathbb{Z}_p^n$
- Message:  $m = \mathbf{x} \in \mathbb{Z}_p^n$
- Output:  $f(m) = \mathbf{x} \cdot \mathbf{y}$



# Inner Product Functional Encryption (IPFE)

$\text{Setup}(1^\lambda, 1^n) \rightarrow \text{msk}$

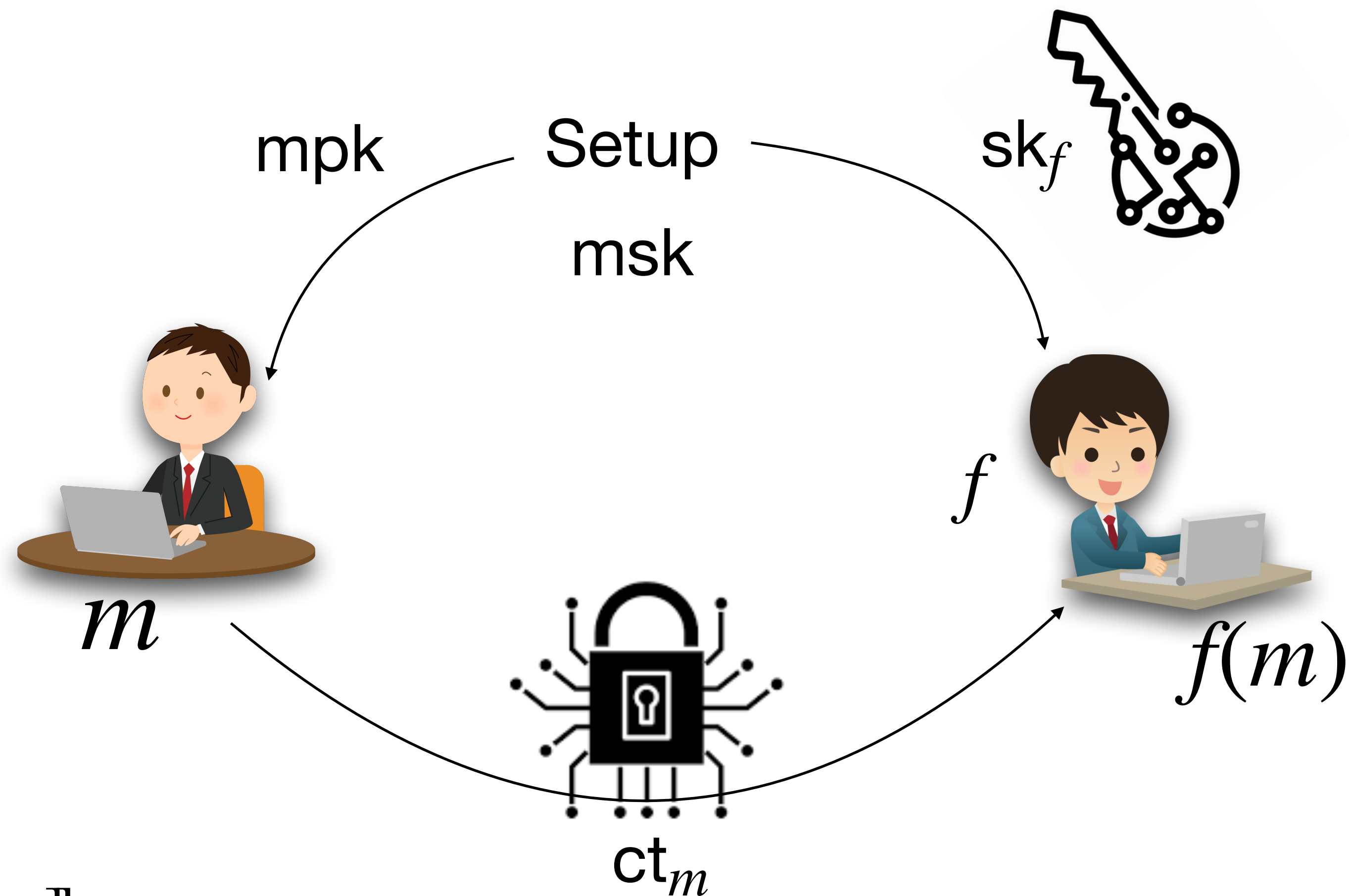
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$\text{Enc}(\text{msk}, m) \rightarrow \text{ct}_m$

$\text{Dec}(\text{sk}_f, \text{ct}_m) \rightarrow f(m)$

$e : \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{G}_T$

- Functions:  $f = \llbracket \mathbf{y} \rrbracket_2 \in \mathbb{G}_2^n$
- Message:  $m = \llbracket \mathbf{x} \rrbracket_1 \in \mathbb{G}_1^n$
- Output:  $f(m) = e(\llbracket \mathbf{x} \rrbracket_1, \llbracket \mathbf{y} \rrbracket_2) = \llbracket \mathbf{x} \cdot \mathbf{y} \rrbracket_T$



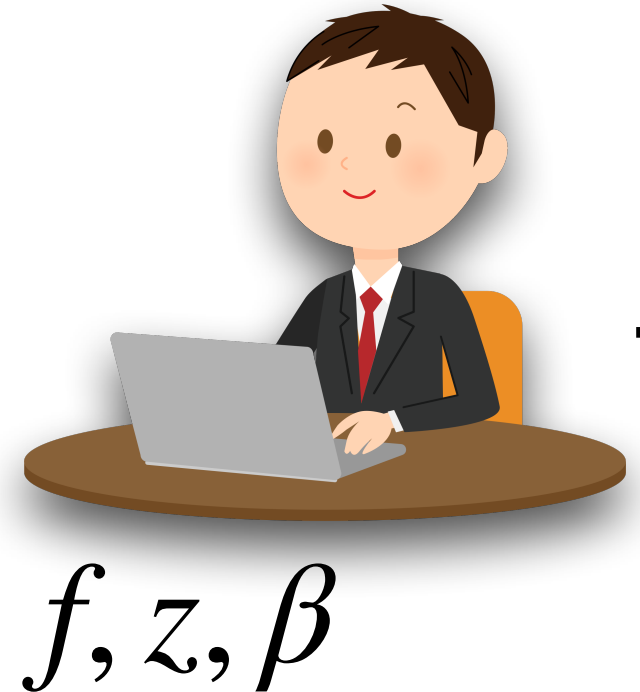
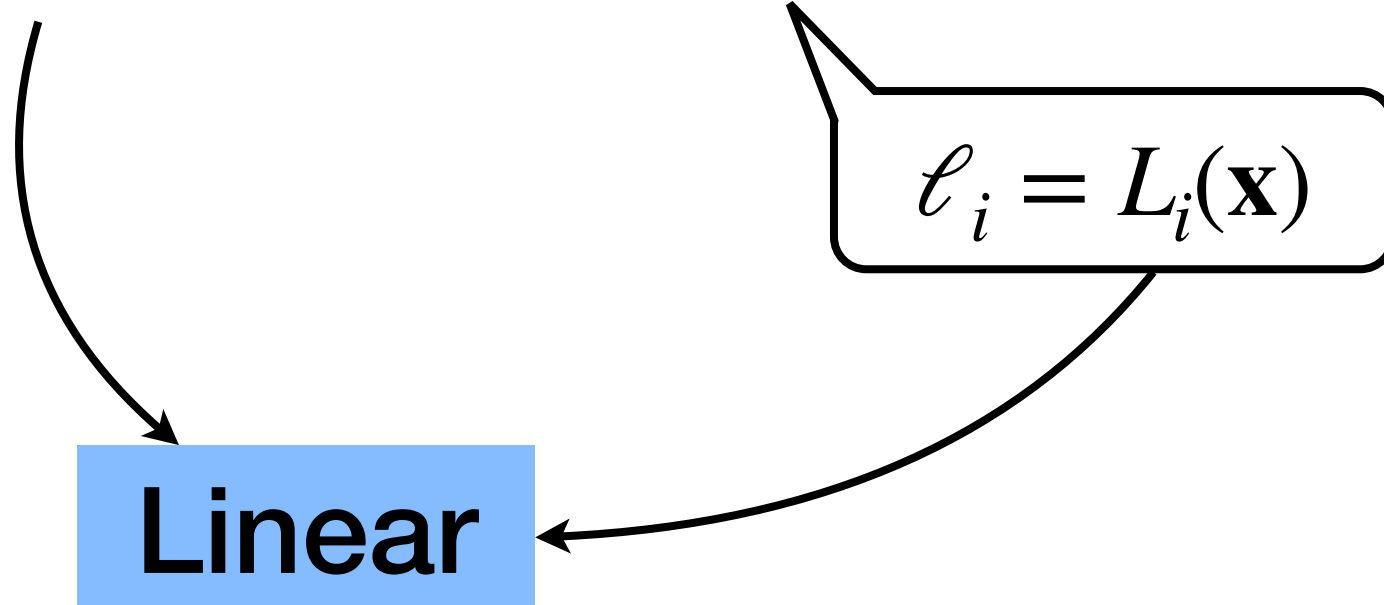
FH-IPFE:  $\{\text{sk}_{\mathbf{y}_0}, \text{ct}_{\mathbf{x}_0}\} \approx_c \{\text{sk}_{\mathbf{y}_1}, \text{ct}_{\mathbf{x}_1}\}$  if  $\llbracket \mathbf{x}_0 \cdot \mathbf{y}_0 \rrbracket_T = \llbracket \mathbf{x}_1 \cdot \mathbf{y}_1 \rrbracket_T$

# Arithmetic Key Garbling Scheme over $\mathbb{Z}_p$ [IW14]

$$\text{Garble}(f, z, \beta) \rightarrow (L_1, \dots, L_t), \quad f: \mathbb{Z}_p^m \rightarrow \mathbb{Z}_p$$

$$\text{Eval}(f, \mathbf{x}, \ell_1, \dots, \ell_t) \rightarrow zf(\mathbf{x}) + \beta$$

$f, \mathbf{x}$  are **public**,  $z, \beta$  are **private**



$f, \mathbf{x}, (L_1(\mathbf{x}), \dots, L_t(\mathbf{x}))$   
 $f, \mathbf{x}, (\ell_1, \dots, \ell_t)$



$zf(\mathbf{x}) + \beta$

$\text{Sim}(f, \mathbf{x}, zf(\mathbf{x}) + \beta)$

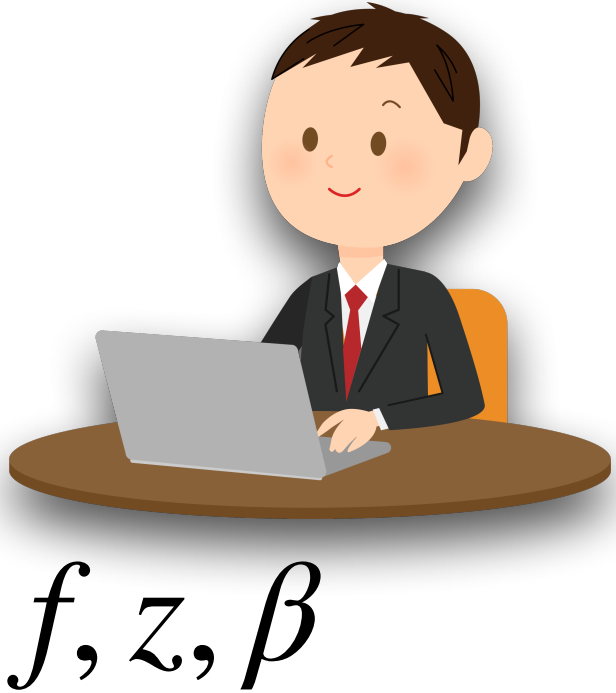
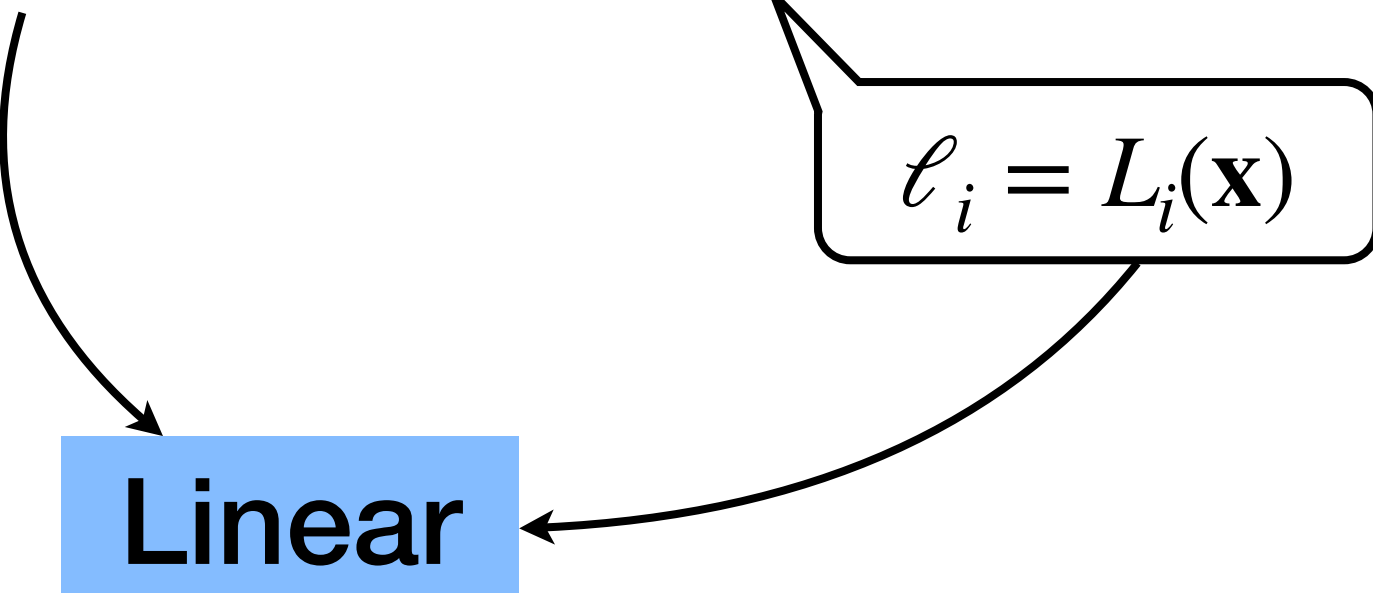
$$\text{Sim}(f, \mathbf{x}, zf(\mathbf{x}) + \beta) \rightarrow (\ell_1, \dots, \ell_t)$$

# Arithmetic Key Garbling Scheme: Piecewise Security [LL20]

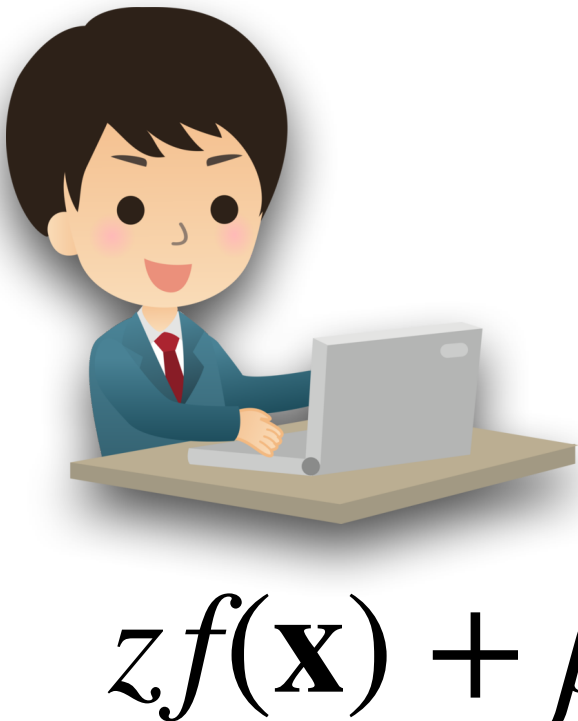
$$\text{Garble}(f, z, \beta) \rightarrow (L_1, \dots, L_t), \quad f: \mathbb{Z}_p^m \rightarrow \mathbb{Z}_p$$

$$\text{Eval}(f, \mathbf{x}, \ell_1, \dots, \ell_t) \rightarrow zf(\mathbf{x}) + \beta$$

$f, \mathbf{x}$  are **public**,  $z, \beta$  are **private**



$f, \mathbf{x}, (L_1(\mathbf{x}), \dots, L_t(\mathbf{x}))$   
 $f, \mathbf{x}, (\ell_1, \ell_2, \dots, \ell_t)$



$\text{RevSamp}(f, \mathbf{x}, zf(\mathbf{x}) + \beta, \ell_2, \dots, \ell_t)$

$$\text{Sim}(f, \mathbf{x}, zf(\mathbf{x}) + \beta) \rightarrow (\ell_1, \dots, \ell_t)$$

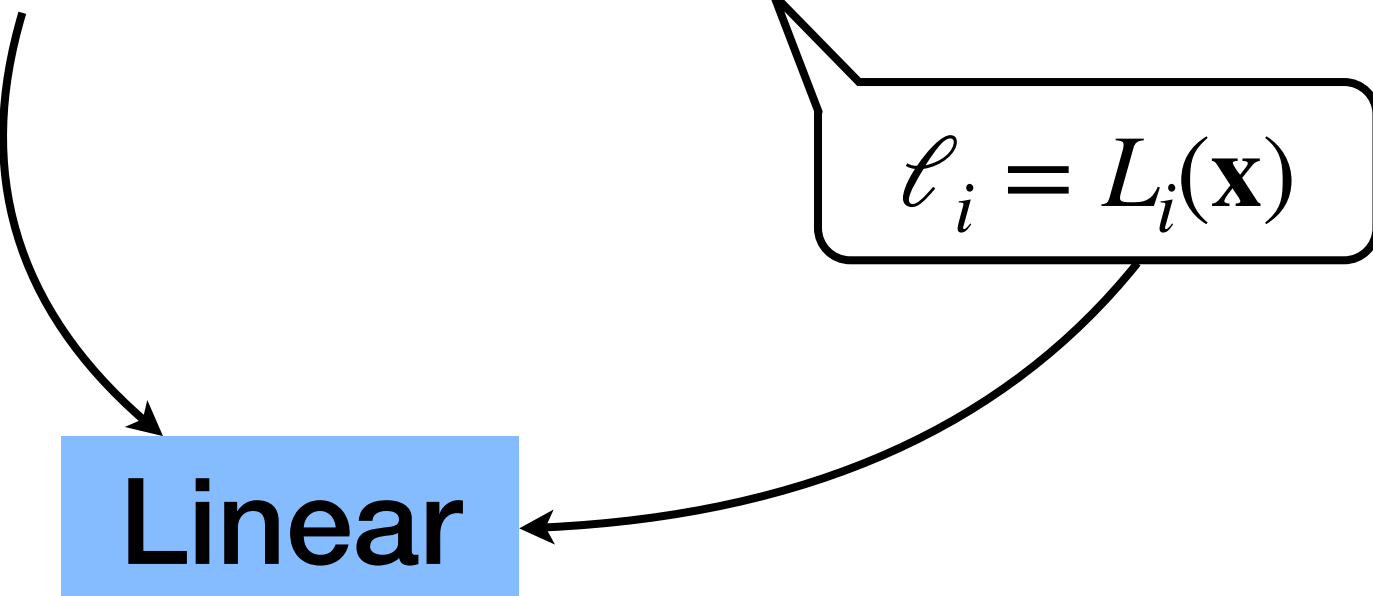
$$\text{RevSamp}(f, \mathbf{x}, zf(\mathbf{x}) + \beta, \ell_2, \dots, \ell_t) \rightarrow \ell_1$$

# Arithmetic Key Garbling Scheme: Piecewise Security [LL20]

$$\text{Garble}(f, z, \beta) \rightarrow (L_1, \dots, L_t), \quad f: \mathbb{Z}_p^m \rightarrow \mathbb{Z}_p$$

$$\text{Eval}(f, \mathbf{x}, \ell_1, \dots, \ell_t) \rightarrow zf(\mathbf{x}) + \beta$$

$f, \mathbf{x}$  are public,  $z, \beta$  are private



$f, \mathbf{x}, (L_1(\mathbf{x}), \dots, L_t(\mathbf{x}))$   
 $f, \mathbf{x}, (\ell_1, r_2, \dots, \ell_t)$



$zf(\mathbf{x}) + \beta$

$\ell_2 \leftarrow \$ \text{ Given } \ell_3, \dots, \ell_t$

$$\text{Sim}(f, \mathbf{x}, zf(\mathbf{x}) + \beta) \rightarrow (\ell_1, \dots, \ell_t)$$

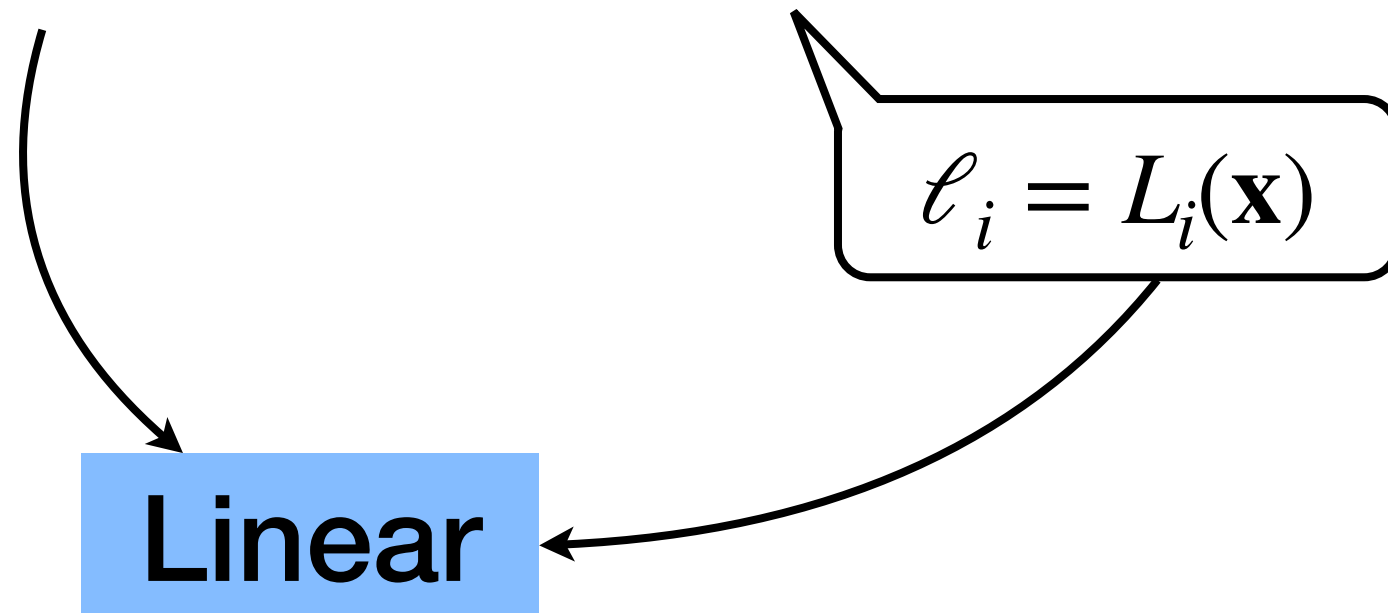
$$\text{RevSamp}(f, \mathbf{x}, zf(\mathbf{x}) + \beta, \ell_2, \dots, \ell_t) \rightarrow \ell_1$$

# Arithmetic Key Garbling Scheme: Piecewise Security [LL20]

$$\text{Garble}(f, z, \beta) \rightarrow (L_1, \dots, L_t), \quad f: \mathbb{Z}_p^m \rightarrow \mathbb{Z}_p$$

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$f, \mathbf{x}$  are **public**,  $z, \beta$  are **private**



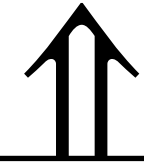
$$\frac{f, \mathbf{x}, (L_1(\mathbf{x}), \dots, L_t(\mathbf{x}))}{f, \mathbf{x}, (\ell_1, r_2, \dots, r_t)}$$



$f, z, \beta$

$zf(\mathbf{x}) + \beta$

$$\text{Sim}(f, \mathbf{x}, zf(\mathbf{x}) + \beta) \rightarrow (\ell_1, \dots, \ell_t)$$



$$\text{RevSamp}(f, \mathbf{x}, zf(\mathbf{x}) + \beta, \ell_2, \dots, \ell_t) \rightarrow \ell_1$$

Marginal Randomness:  $\ell_{j>1} \leftarrow \$$ , Given  $\ell_{j+1}, \dots, \ell_t$

**Piecewise Security**



# Core Idea of [DP21] for FE-AWS for ABPs

---

function

$$f = (f_1, f_2) \text{ s.t. } f_k : \mathbb{Z}_p^n \rightarrow \mathbb{Z}_p$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2))$$

output

$$z_1 f_1(\mathbf{x}) + z_2 f_2(\mathbf{x})$$

---

# Core Idea of [DP21] for FE-AWS for ABPs

function

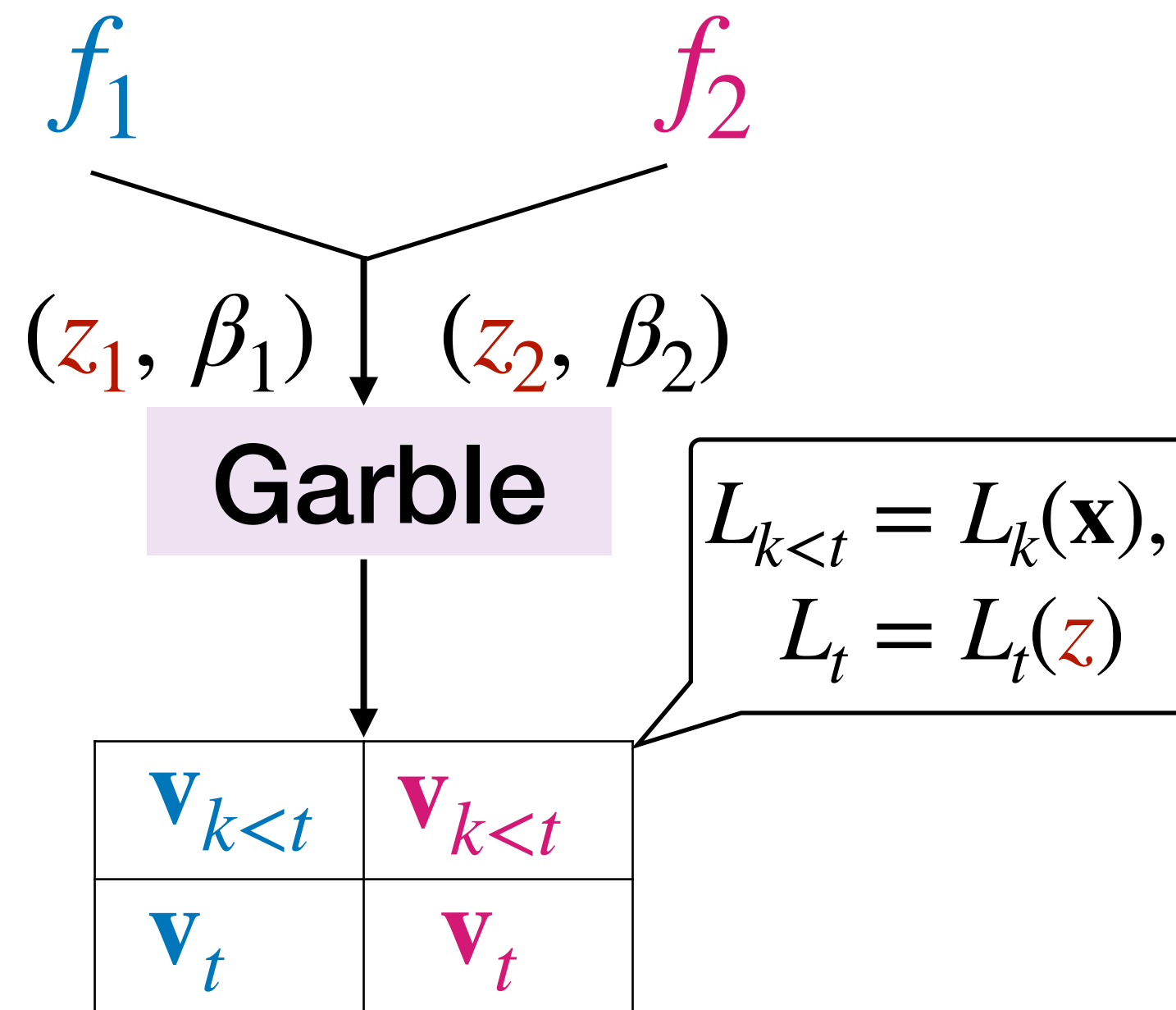
$$f = (f_1, f_2) \text{ s.t. } f_k : \mathbb{Z}_p^n \rightarrow \mathbb{Z}_p$$

input

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# Core Idea of [DP21] for FE-AWS for ABPs

function

input

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$$(\mathbf{x}, \mathbf{z} = (z_1, z_2))$$

$$z_1 f_1(\mathbf{x}) + z_2 f_2(\mathbf{x})$$

$f_1$   $f_2$

$(z_1, \beta_1)$   $(z_2, \beta_2)$

Garble

$$L_{k < t} = L_k(\mathbf{x}),$$

$$L_t = L_t(\mathbf{z})$$

$\mathbf{v}_{k < t}$	$\mathbf{v}_{k < t}$
$\mathbf{v}_t$	$\mathbf{v}_t$

IPFE, IPFE<sub>t</sub>

$sk_{k < t}, sk_{k < t}, sk_t, sk_t$

# Core Idea of [DP21] for FE-AWS for ABPs

function

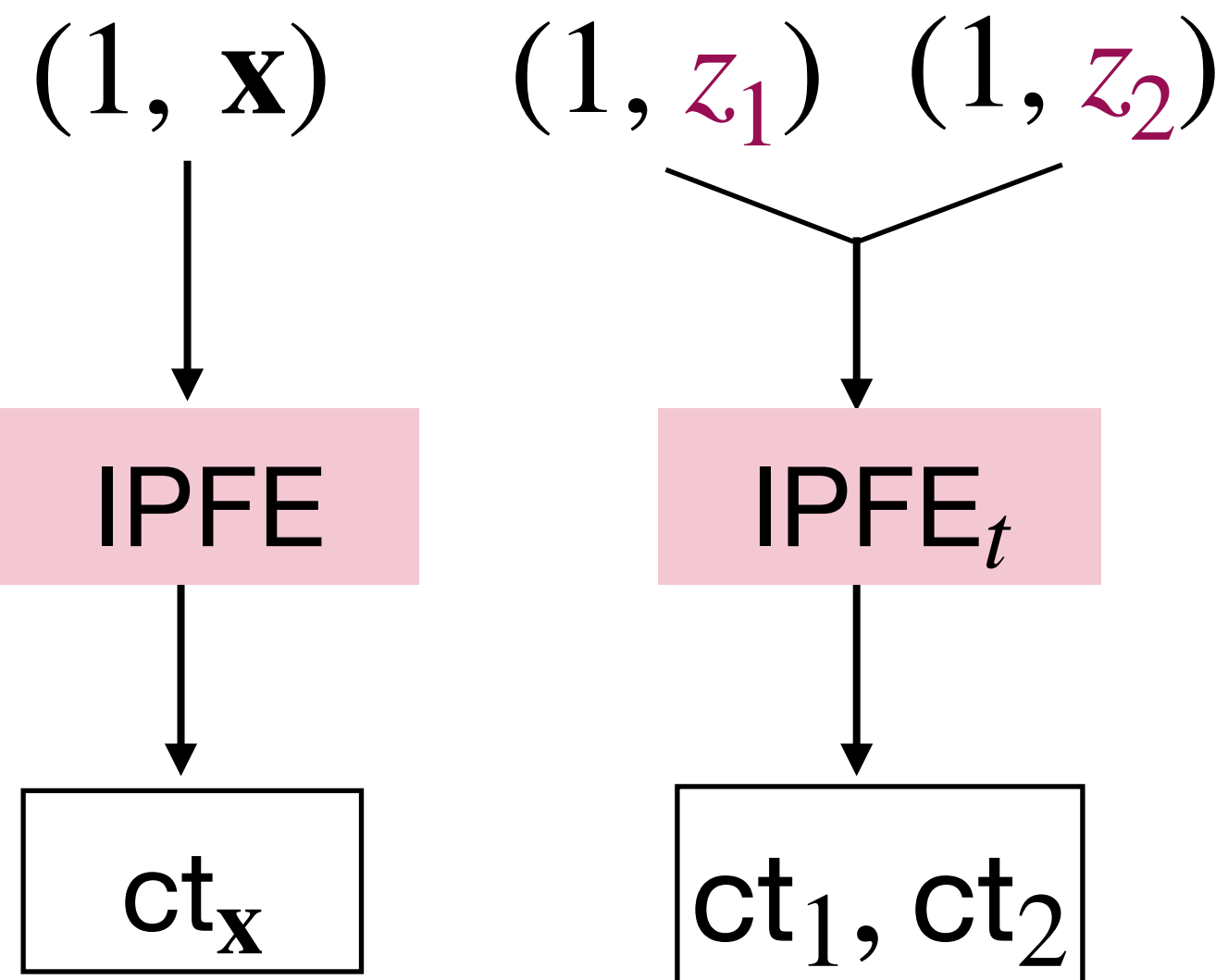
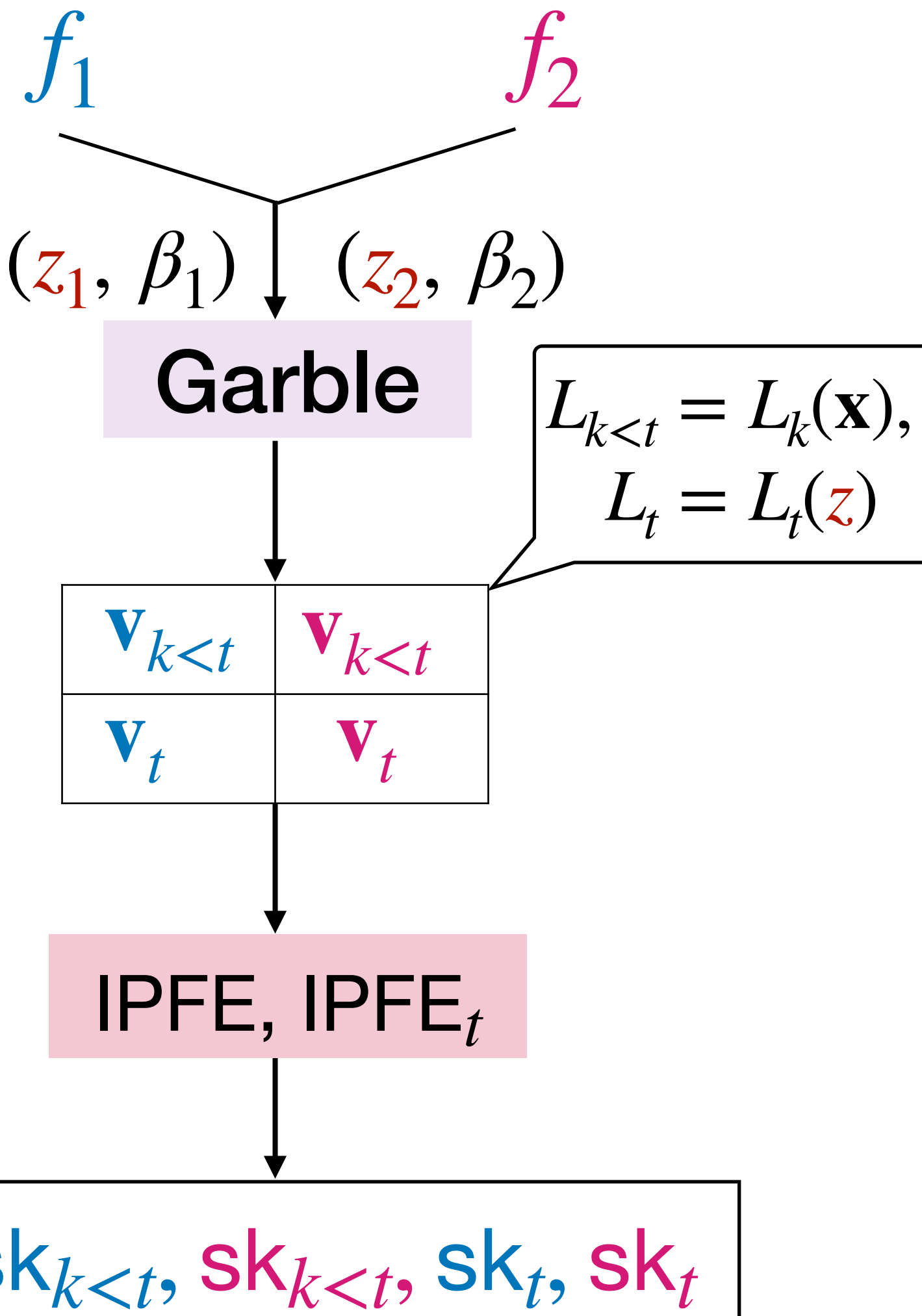
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input

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output

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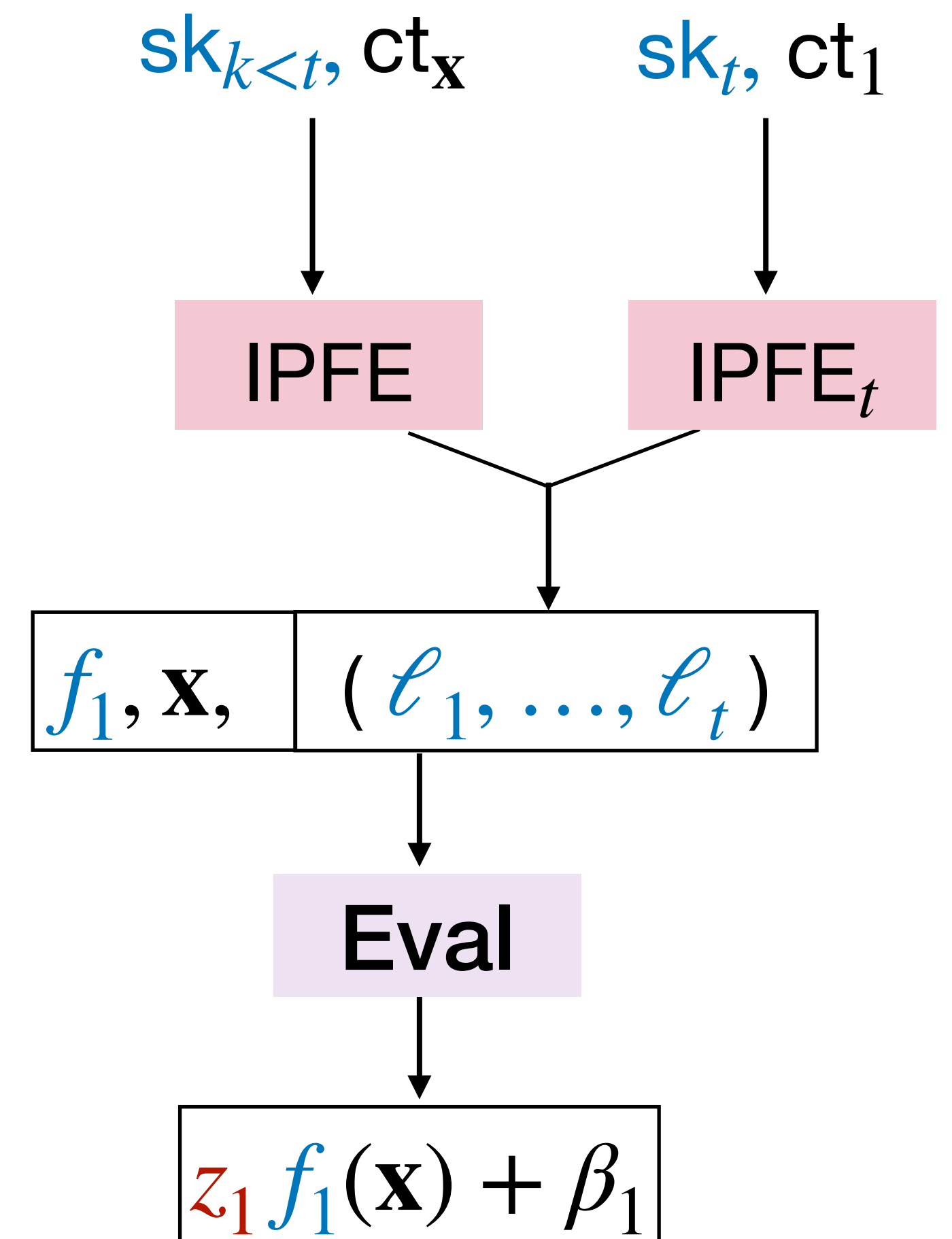
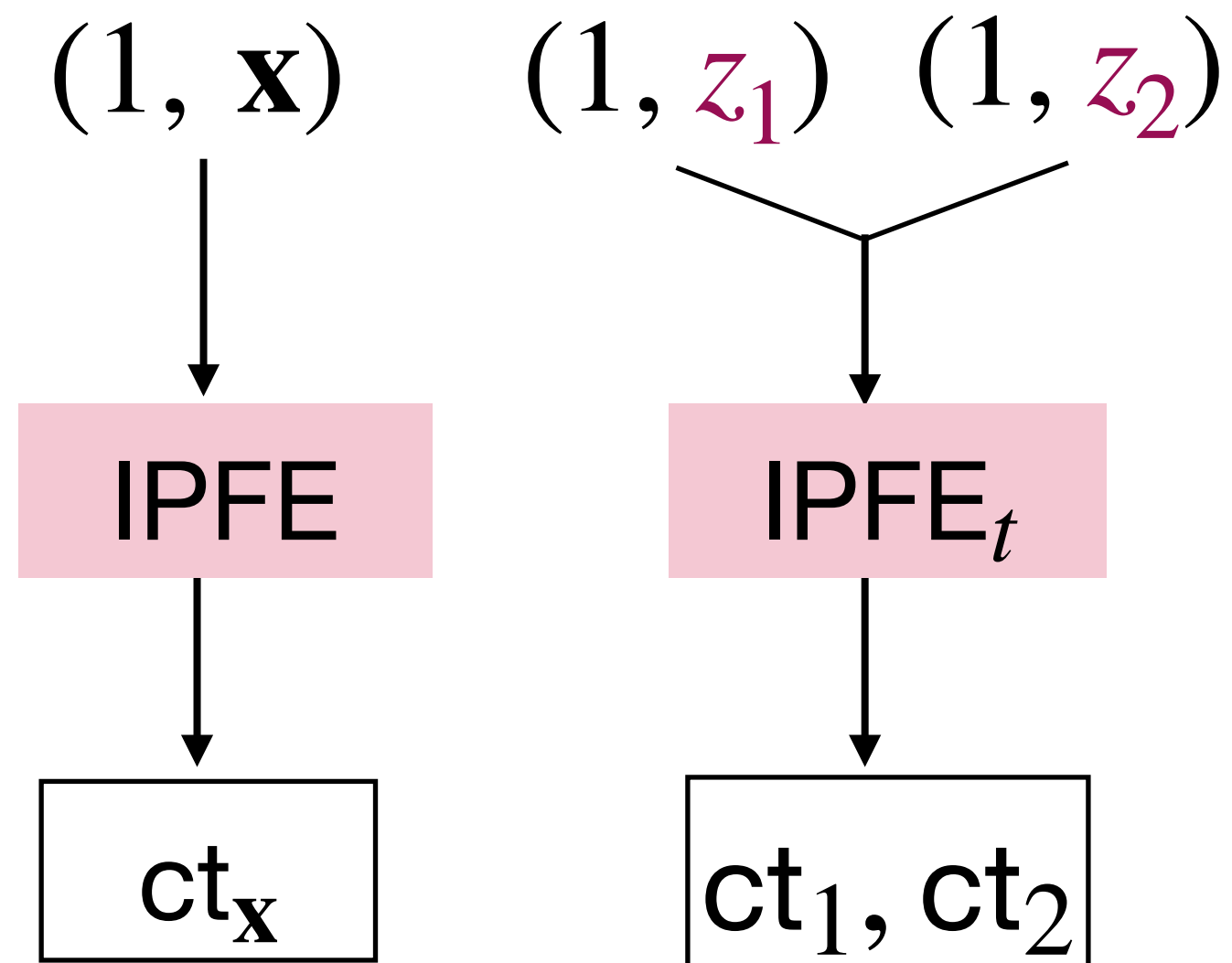
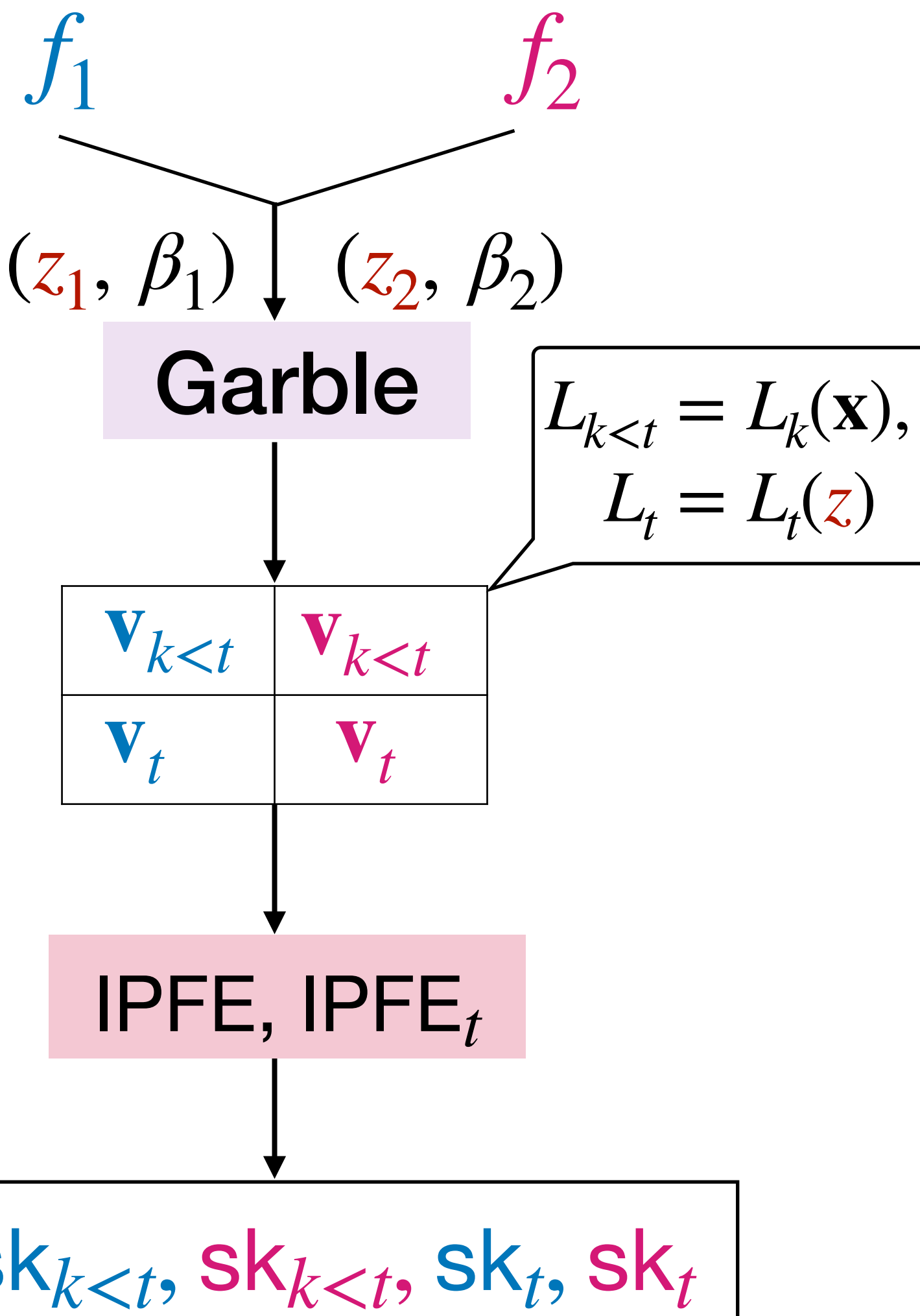
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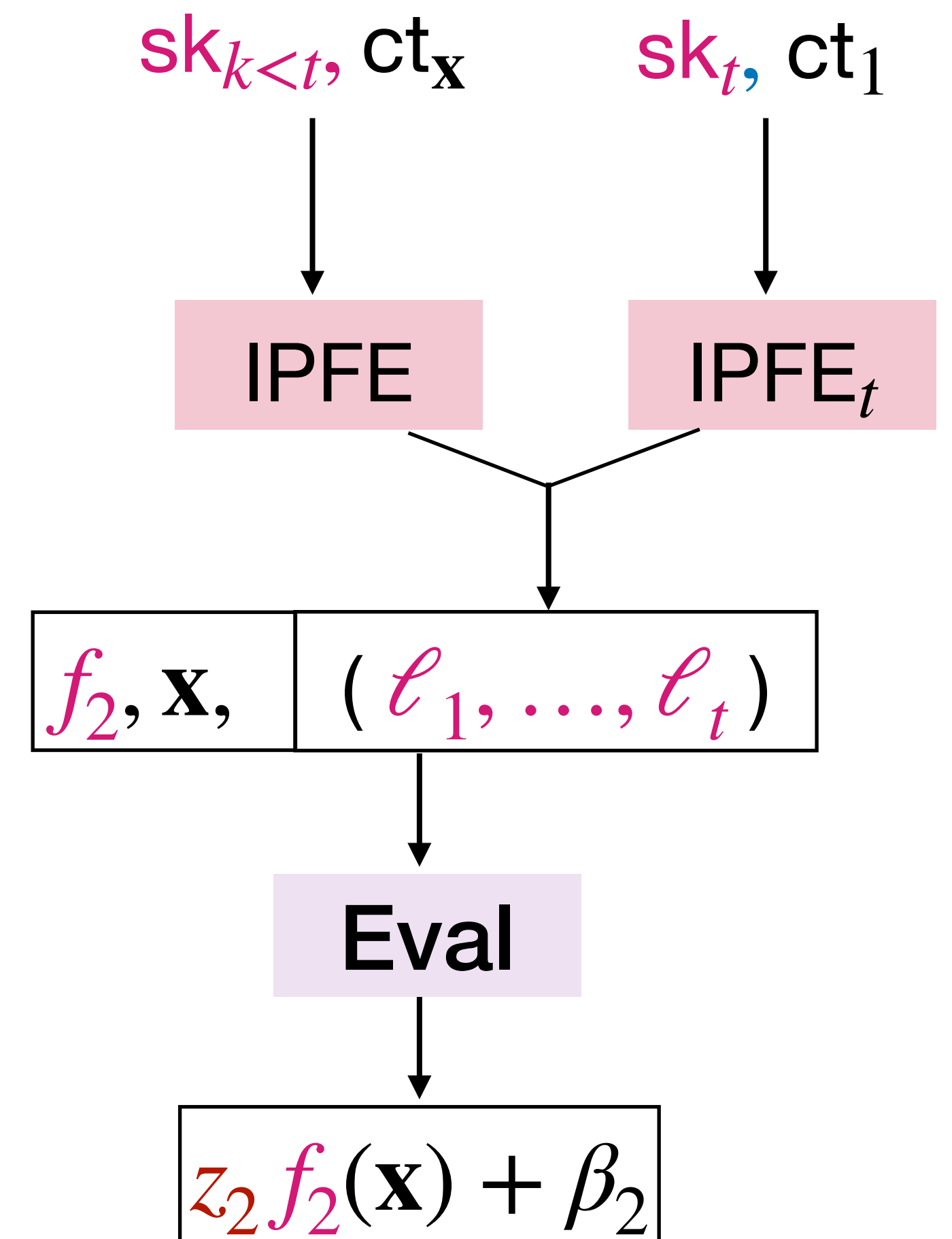
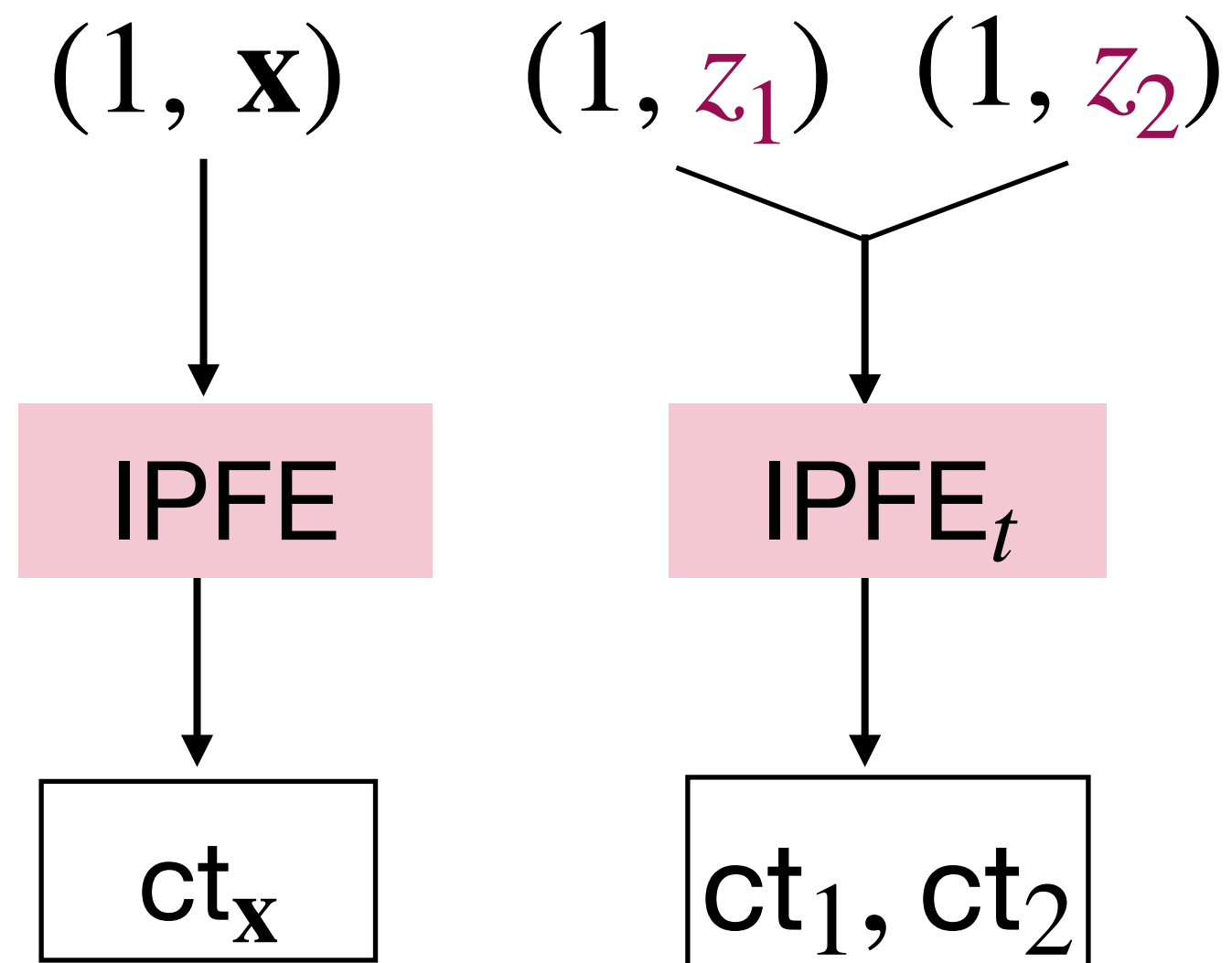
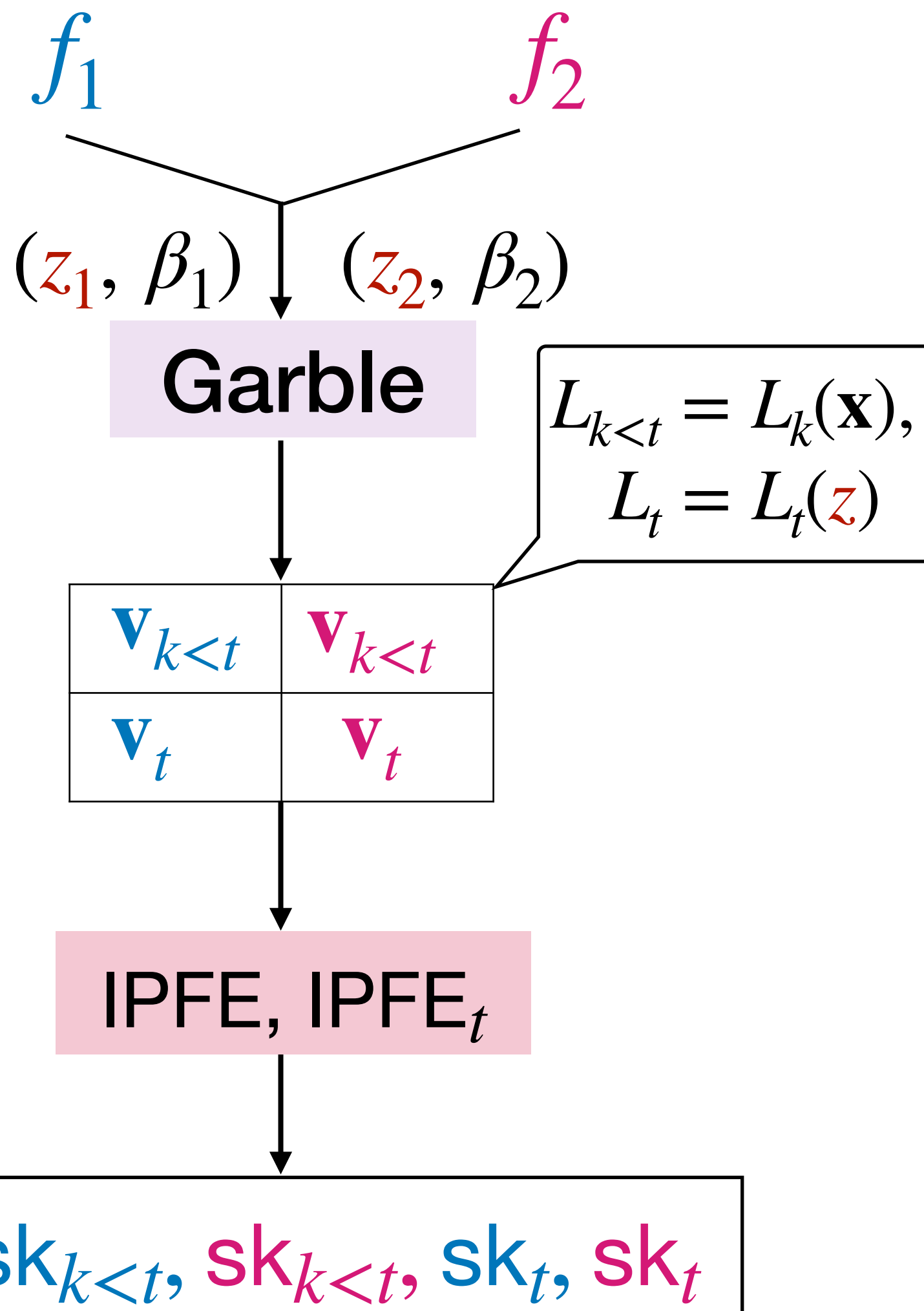
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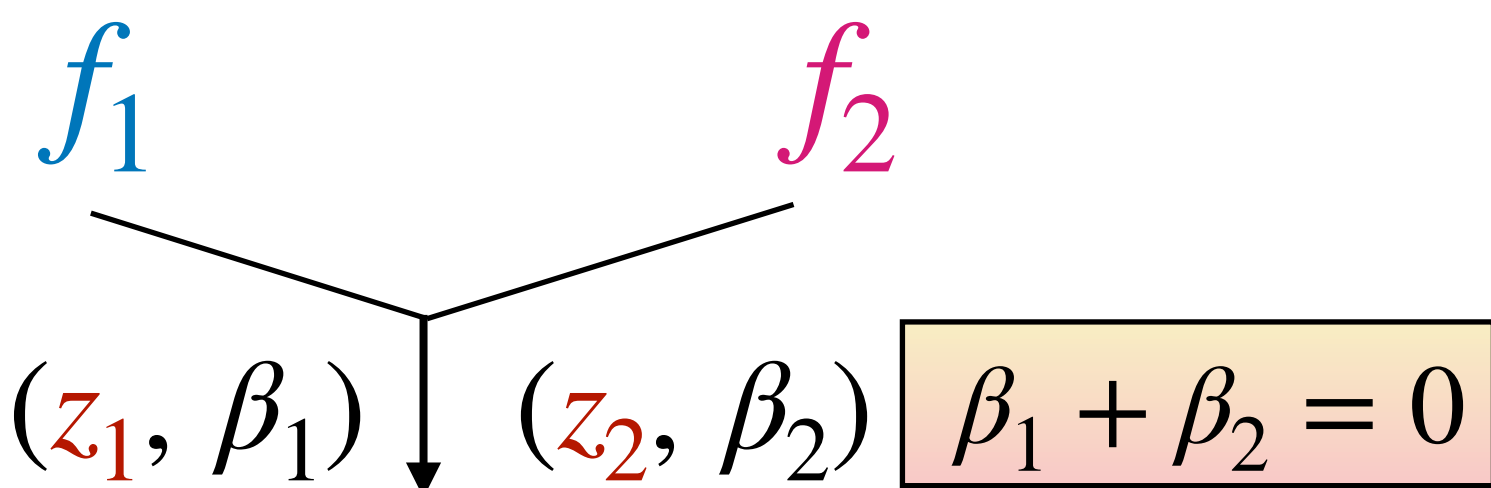
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input

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output

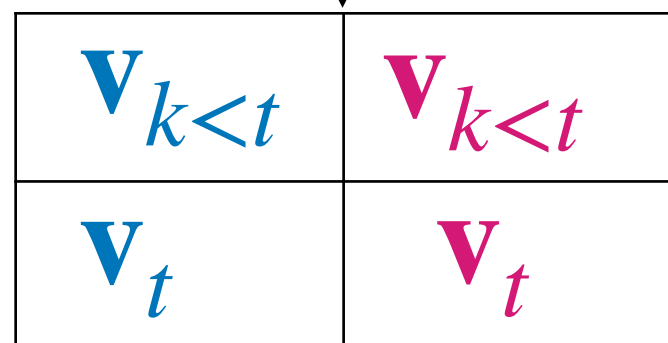
$$z_1 f_1(\mathbf{x}) + z_2 f_2(\mathbf{x})$$



Garble

$$L_{k < t} = L_k(\mathbf{x}),$$

$$L_t = L_t(\mathbf{z})$$



IPFE, IPFE<sub>t</sub>

$$sk_{k < t}, sk_{k < t}, sk_t, sk_t$$

$$(1, \mathbf{x})$$

IPFE

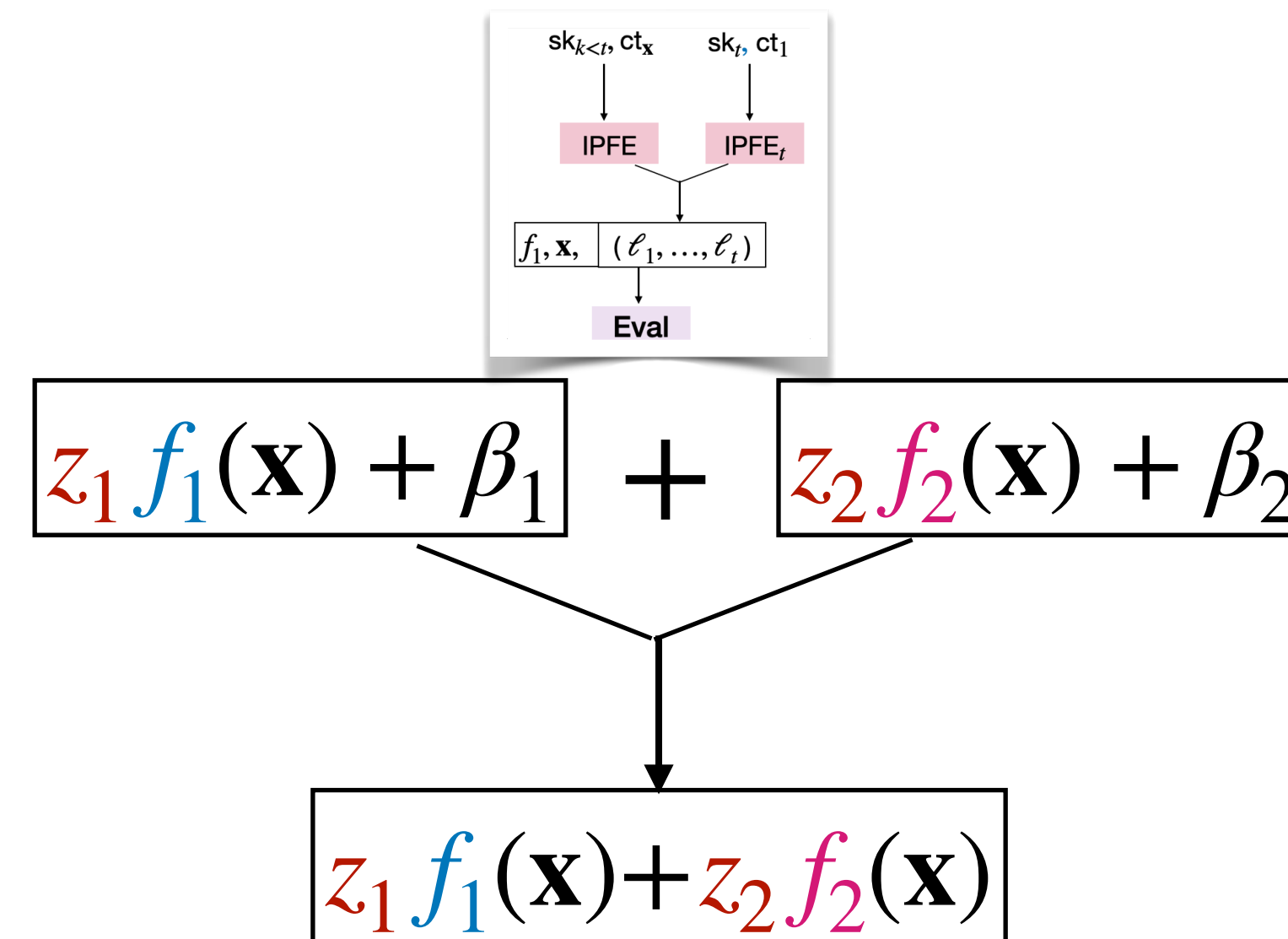
$$ct_{\mathbf{x}}$$

$$(1, z_1)$$

IPFE<sub>t</sub>

$$ct_1, ct_2$$

$$(1, z_2)$$



# Our Idea for FE-UAWS for TMs

---

function

input

output

$M = (M_k)_{k \in I}$  s.t.  $M_k : \{0,1\}^* \rightarrow \{0,1\}$   $(\mathbf{x}, \mathbf{z} = (z_1, \dots, z_n))$

$\sum_{k \in I} M_k(\mathbf{x}) z_k$

---



# Our Idea for FE-UAWS for TMs

---

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

---

# Our Idea for FE-UAWS for TMs

function

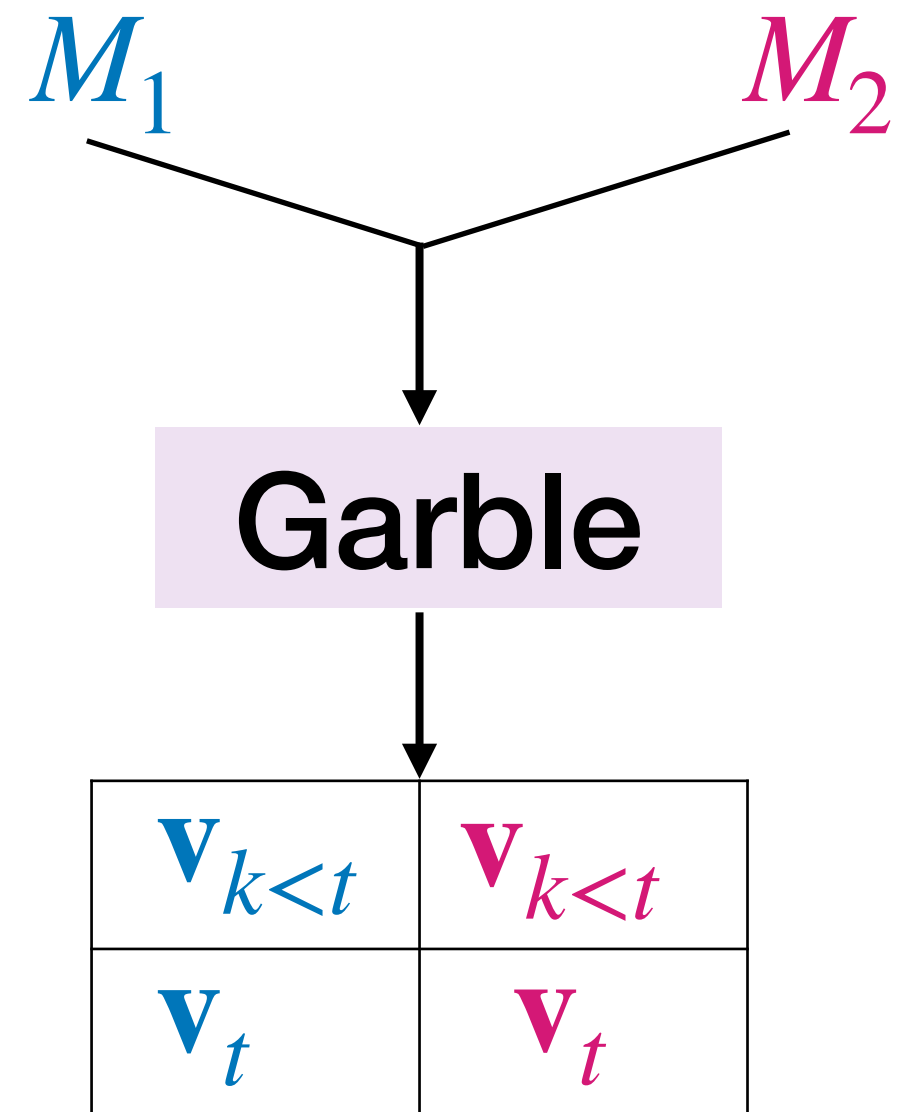
$$M = (M_1, M_2)$$

input

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# Our Idea for FE-UAWS for TMs

function

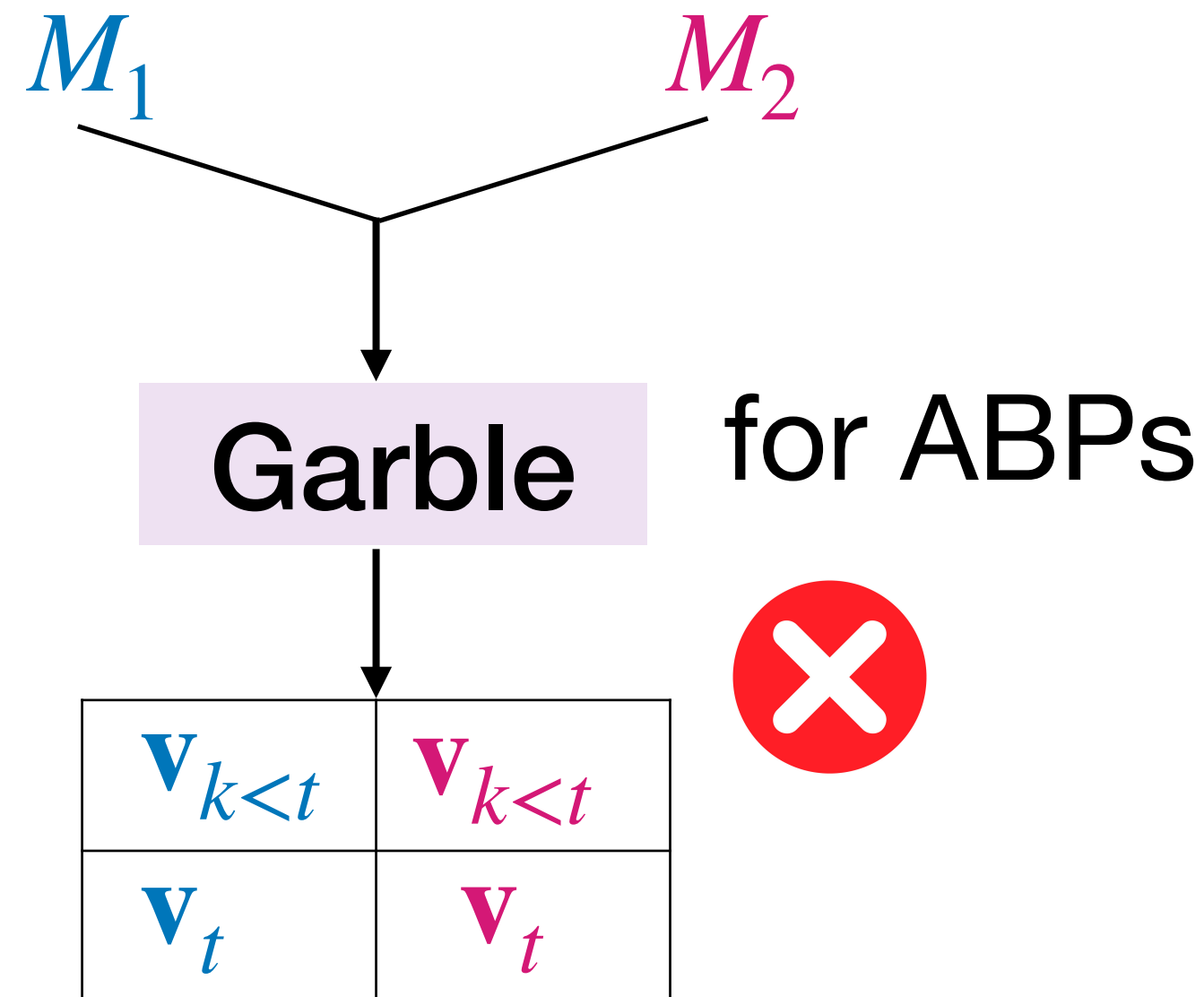
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



AKGS for TM computation [LL20]

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

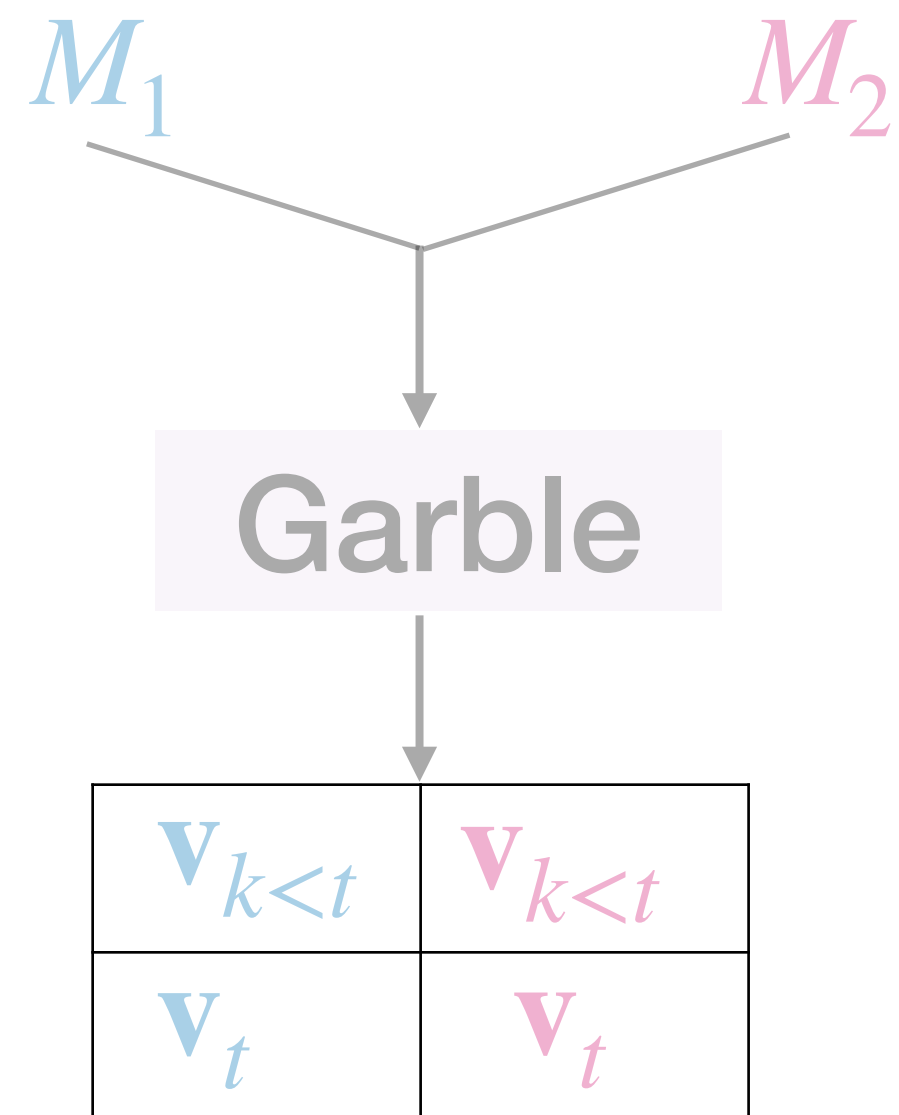
input

output

$$M = (M_1, M_2)$$

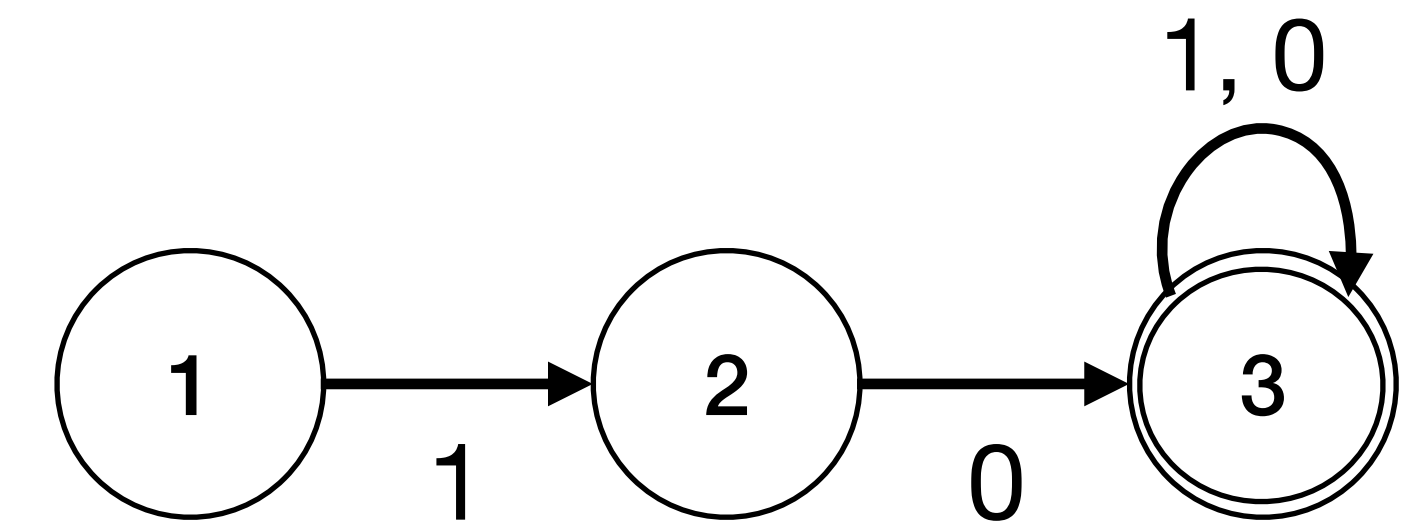
$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## Deterministic Finite Automaton (DFA)

- ◆ states  $\{1, 2, \dots, Q\}$ 
  - initial state 1
  - accepting state  $q_{acc}$
- ◆ transition function  $\delta$



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

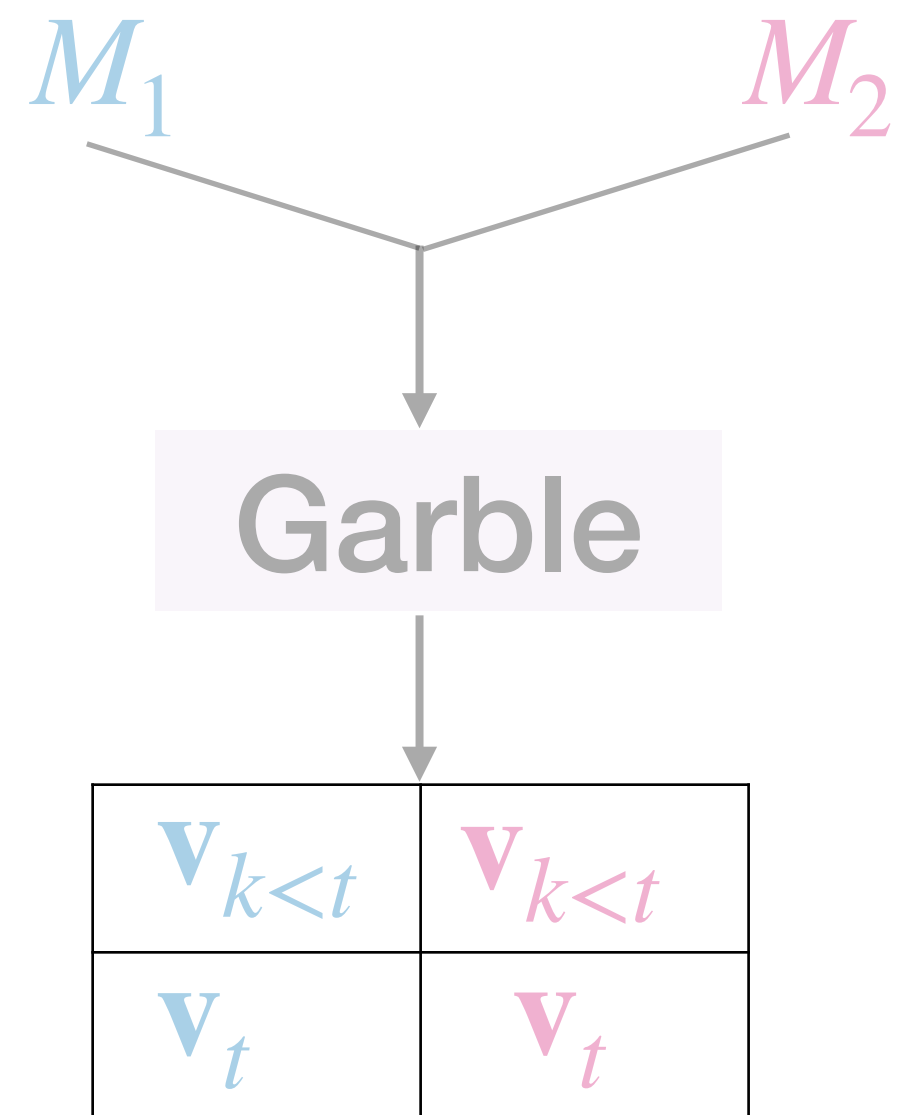
input

output

$$M = (M_1, M_2)$$

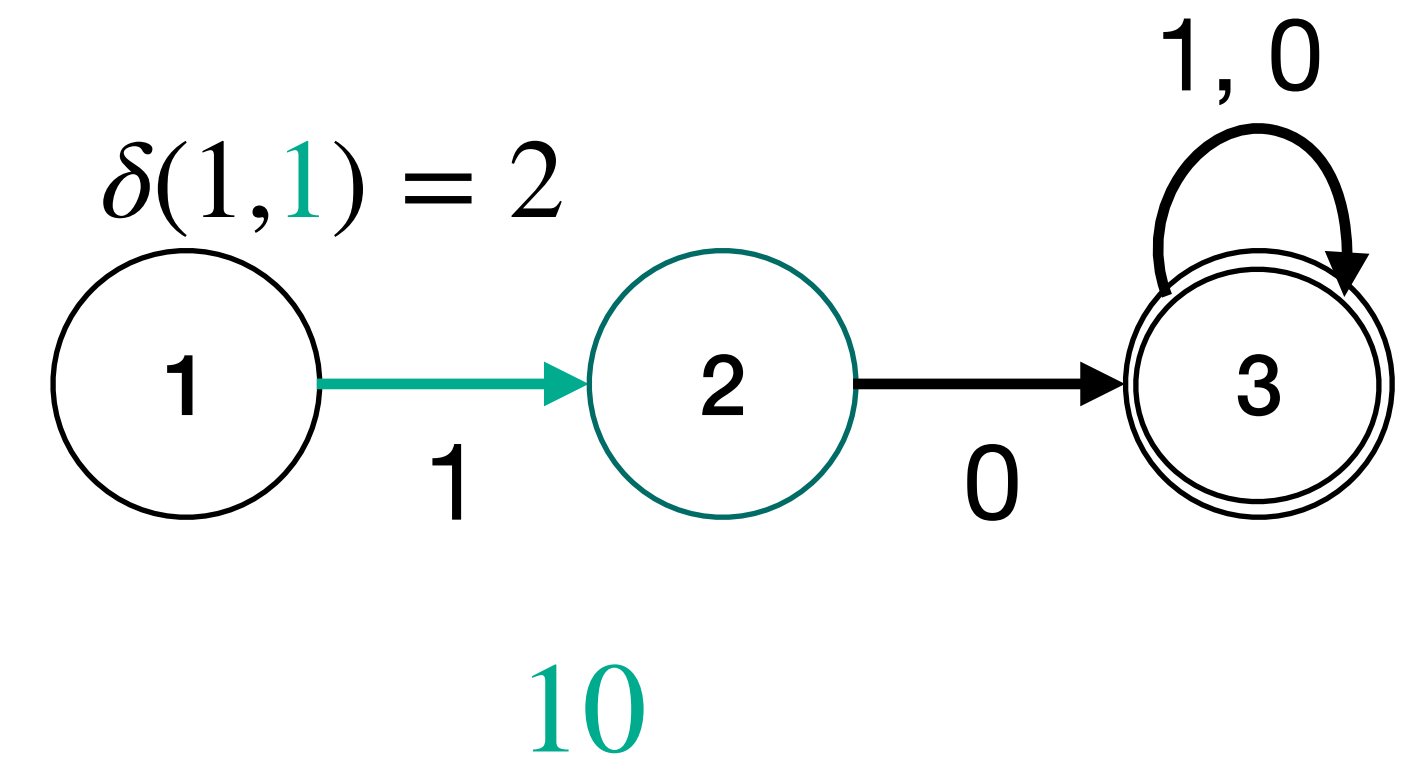
$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

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function

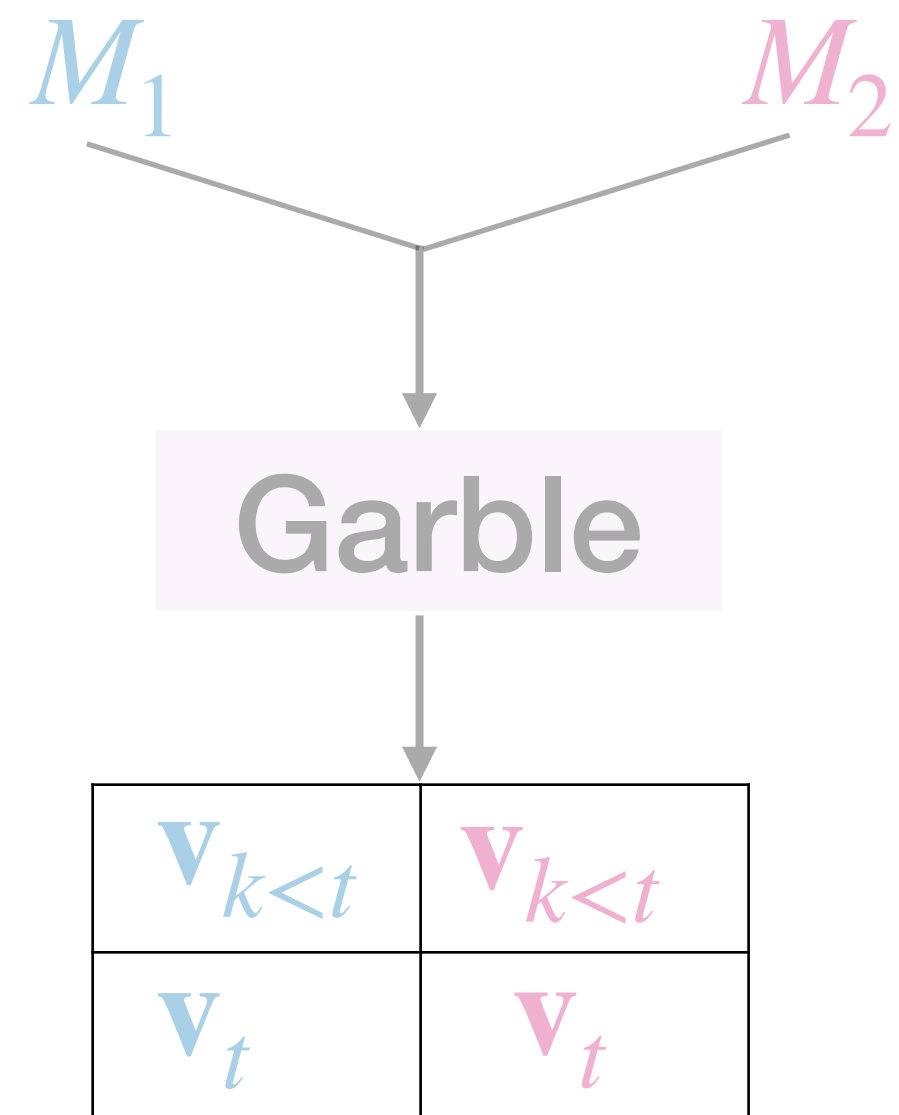
input

output

$$M = (M_1, M_2)$$

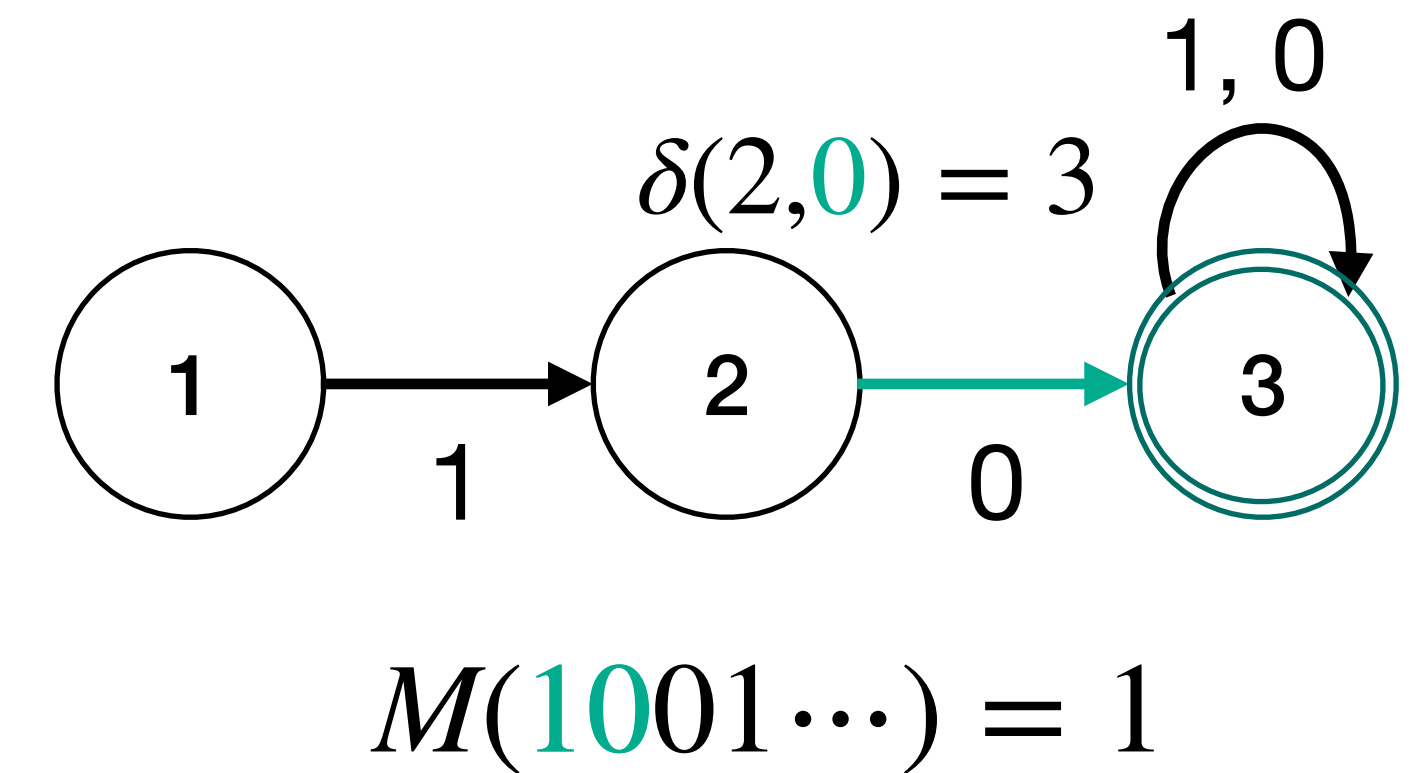
$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

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## Deterministic Finite Automaton (DFA)

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  - initial state 1
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# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

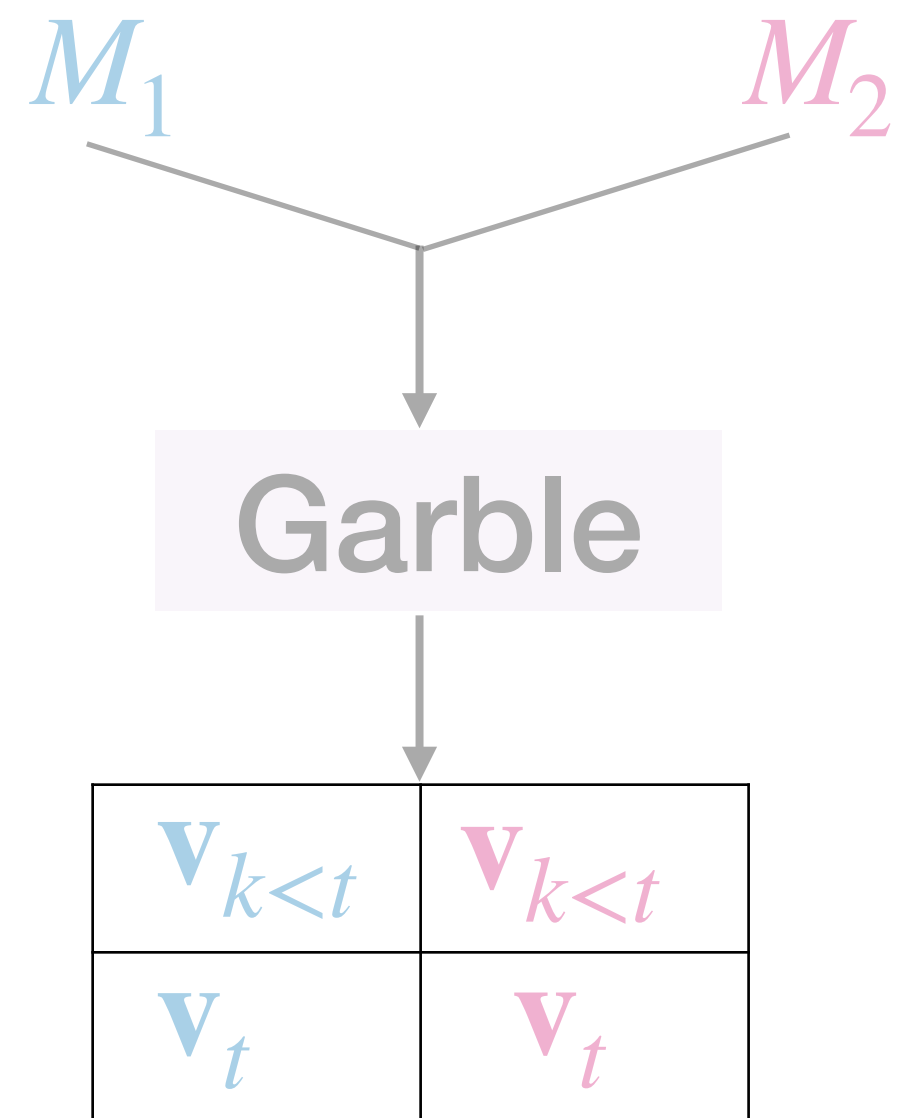
input

output

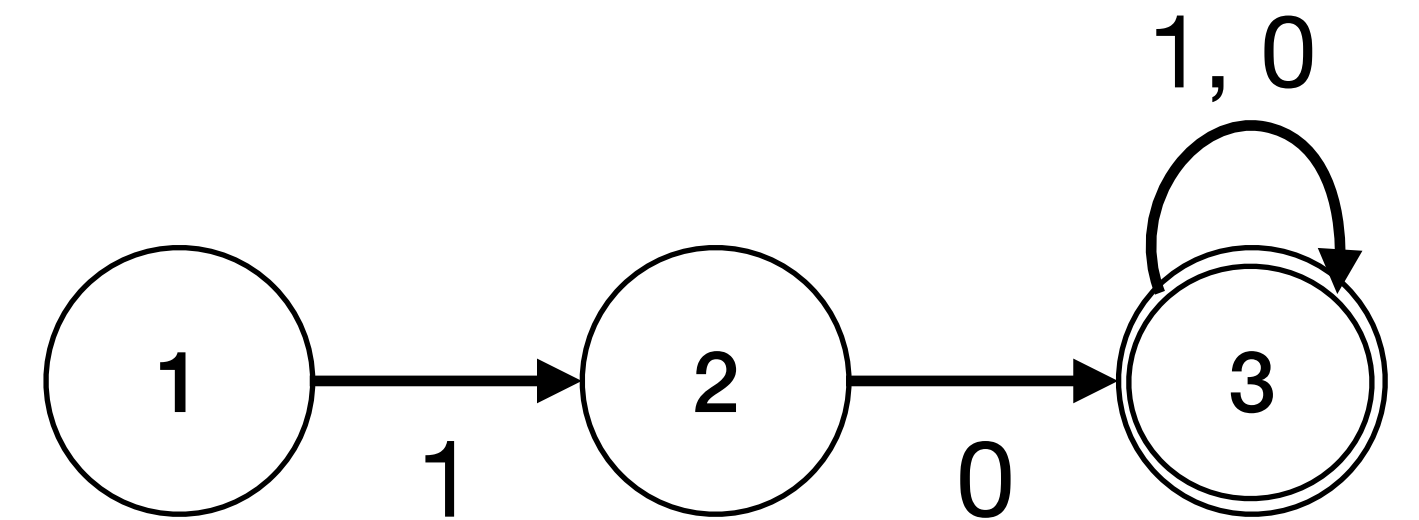
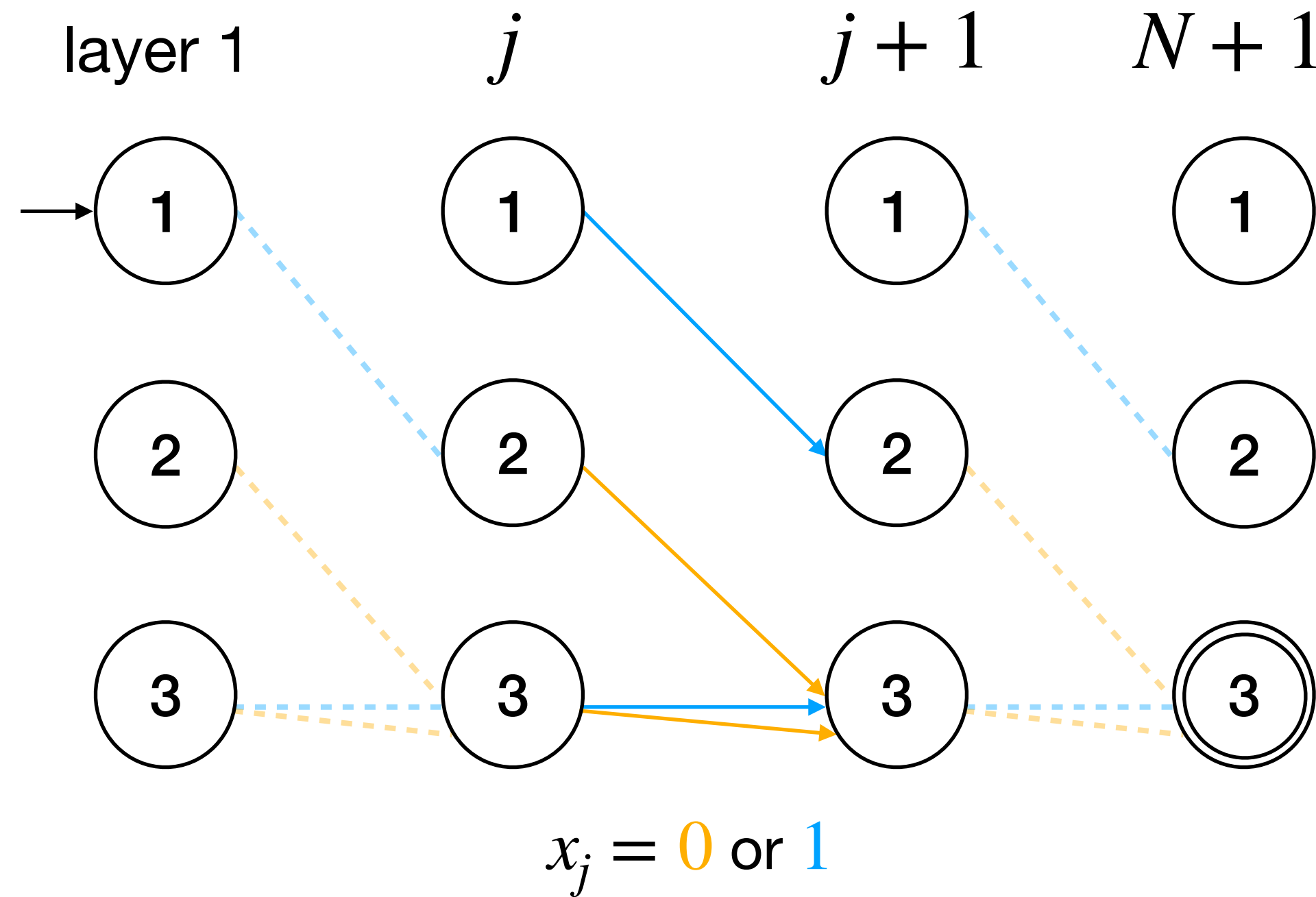
$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## DFA as Branching Program



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

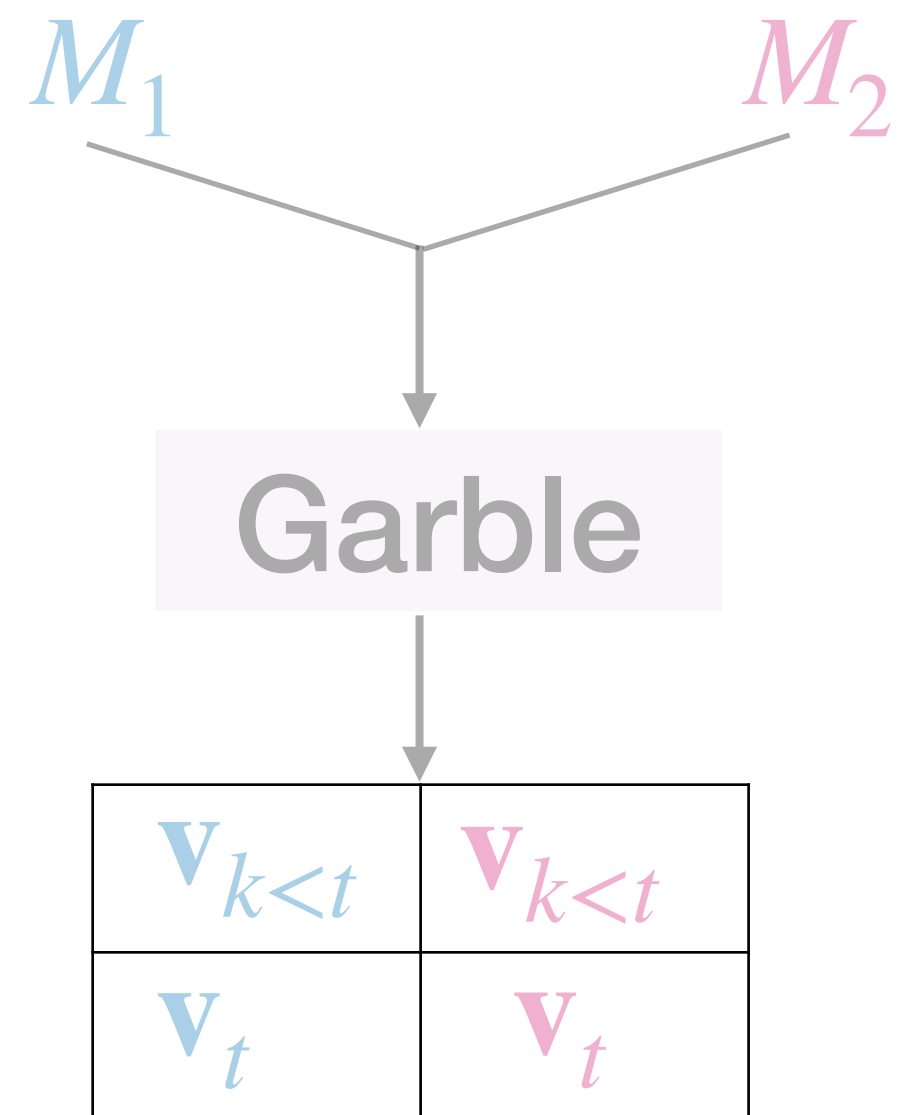
input

output

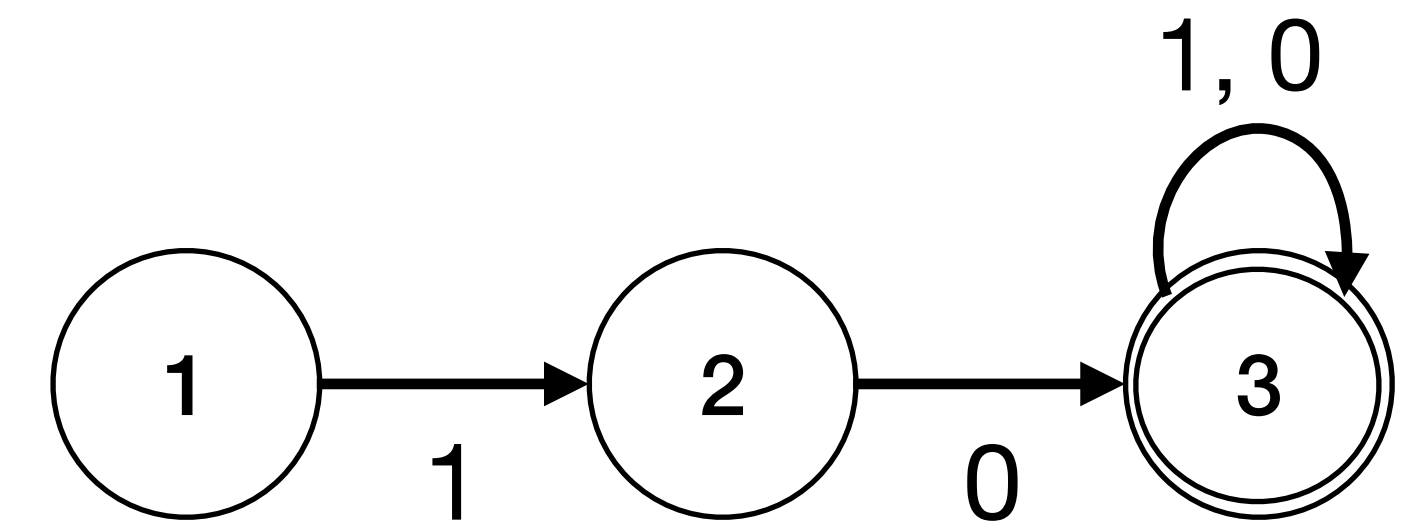
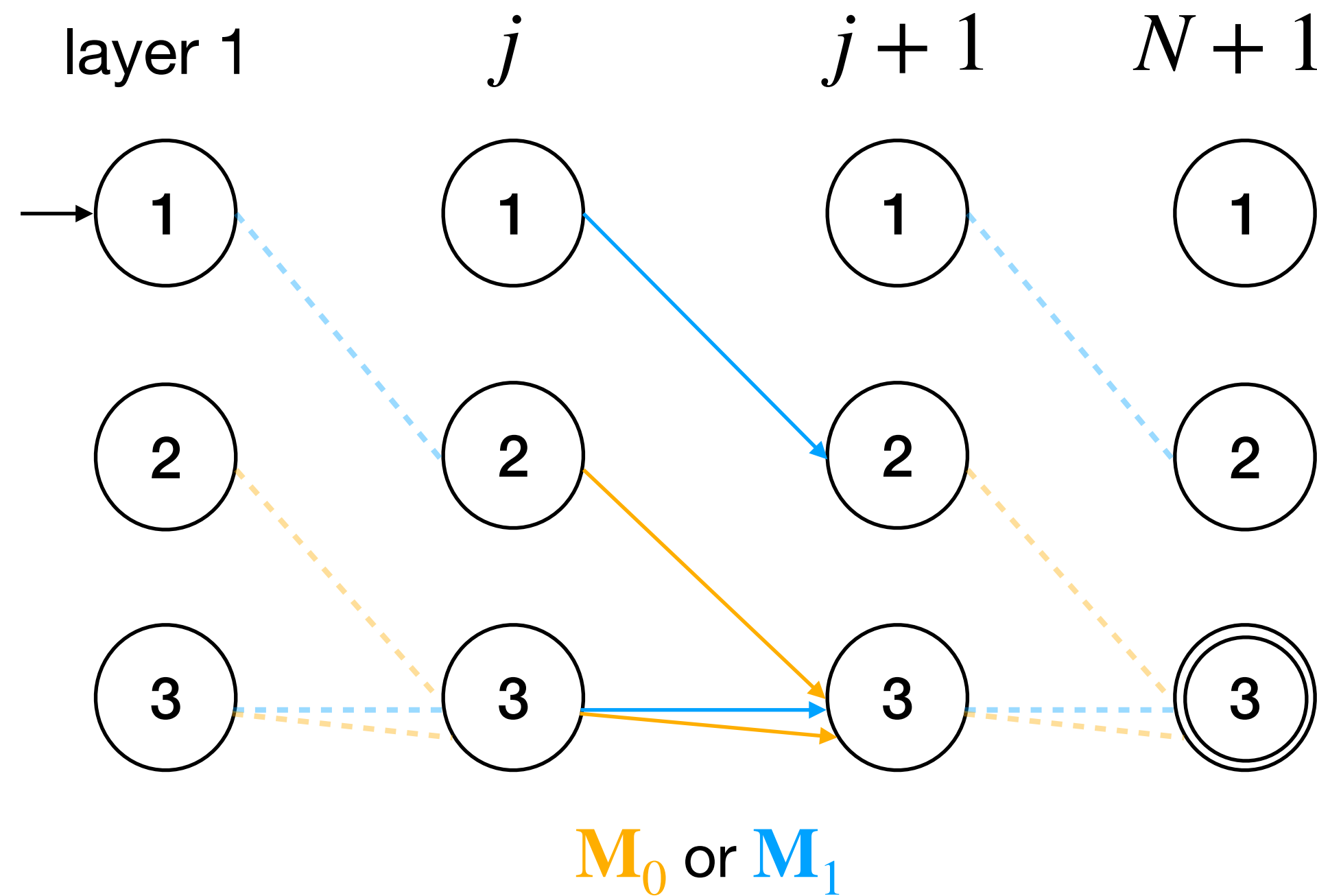
$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## DFA as Matrix Multiplication



$$\mathbf{e}_1^T \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N}$$



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

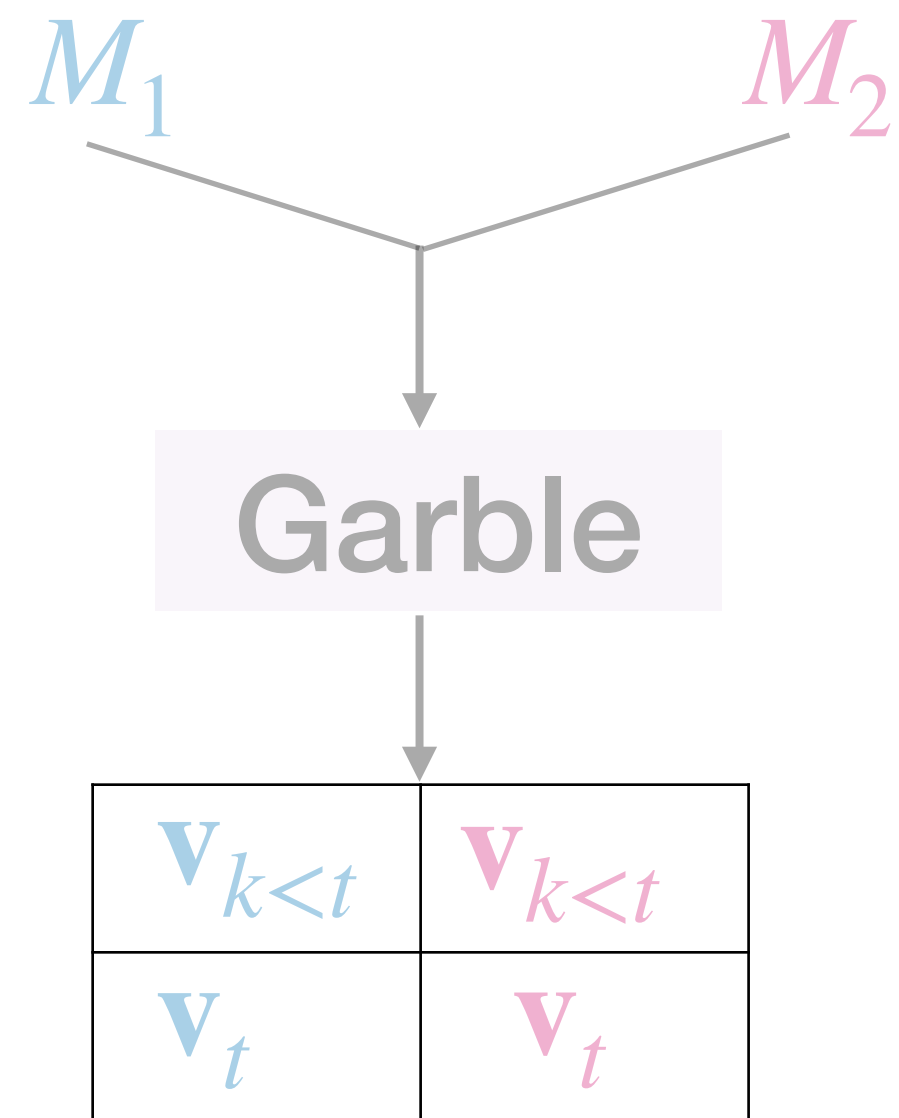
input

output

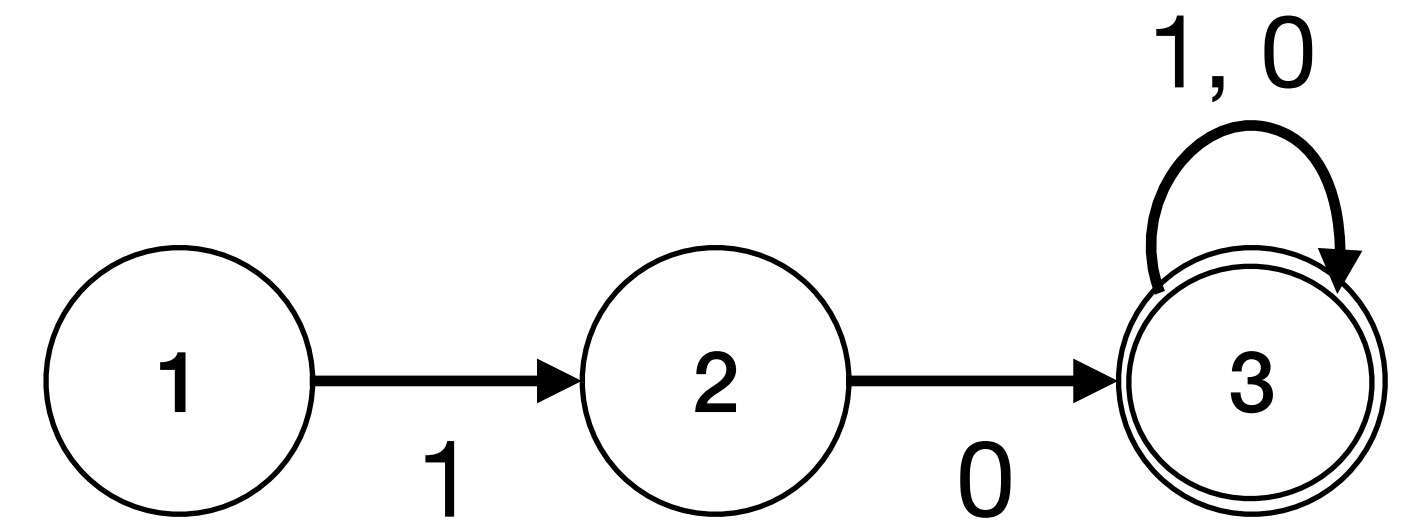
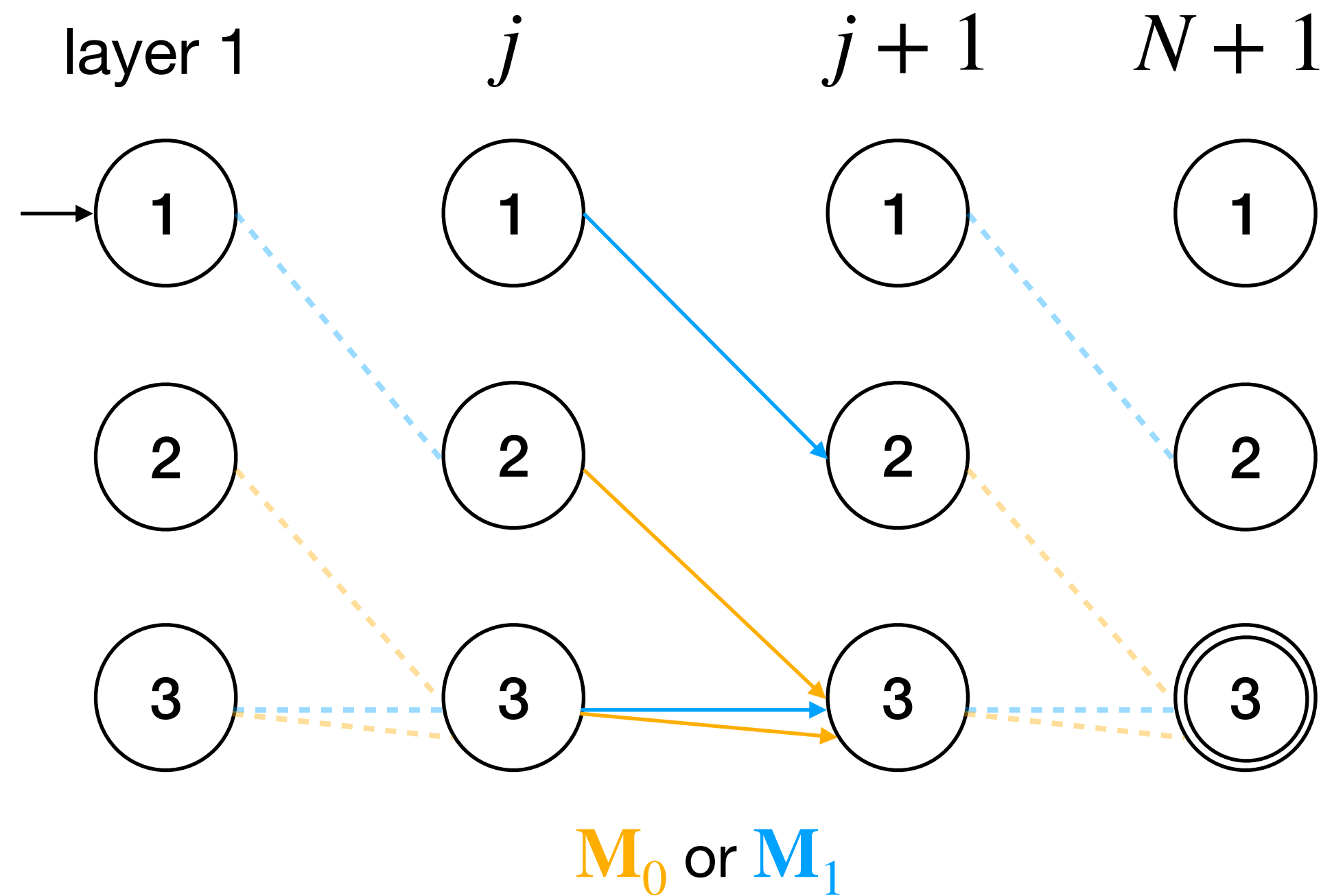
$$M = (M_1, M_2)$$

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$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## DFA as Matrix Multiplication



$$\mathbf{e}_1^T \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} = 1 \text{ or } 0$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

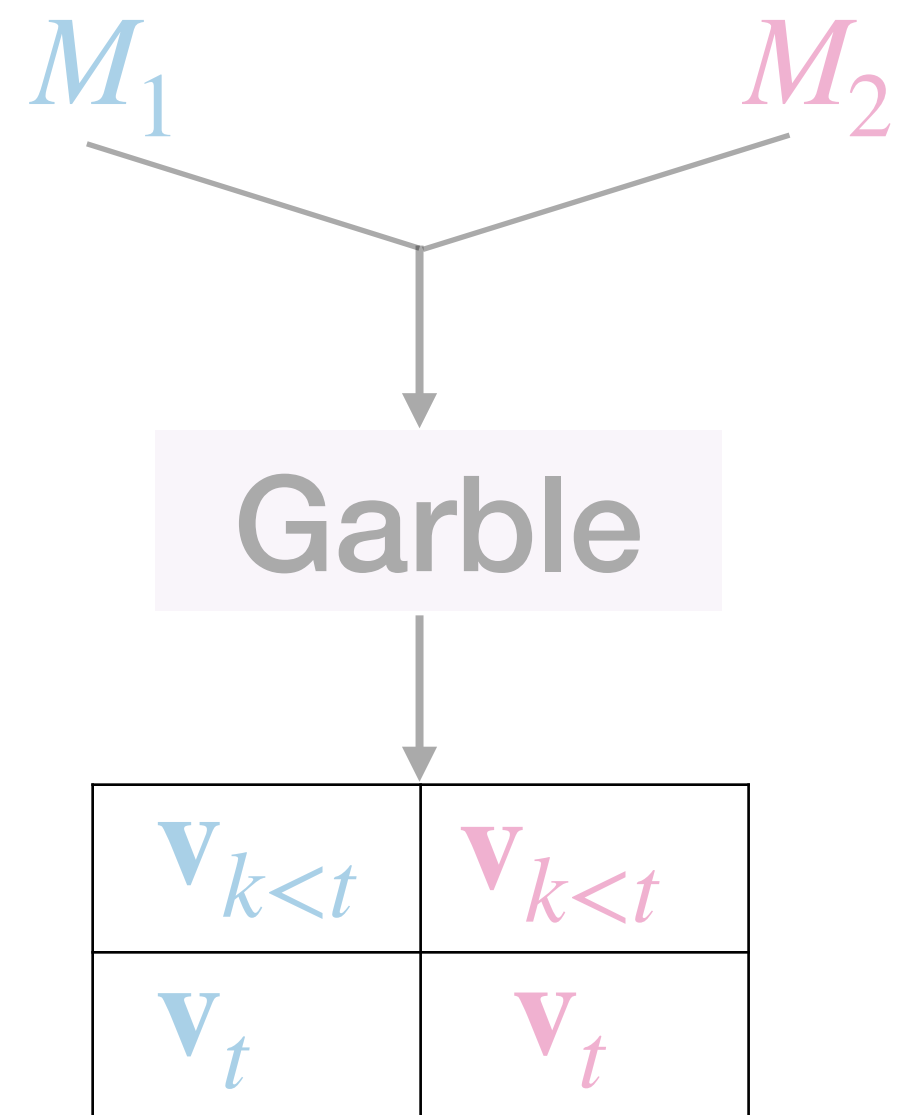
input

output

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$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } \mathbf{z} \cdot \underbrace{\mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}}}_{M(\mathbf{x})} + \beta$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

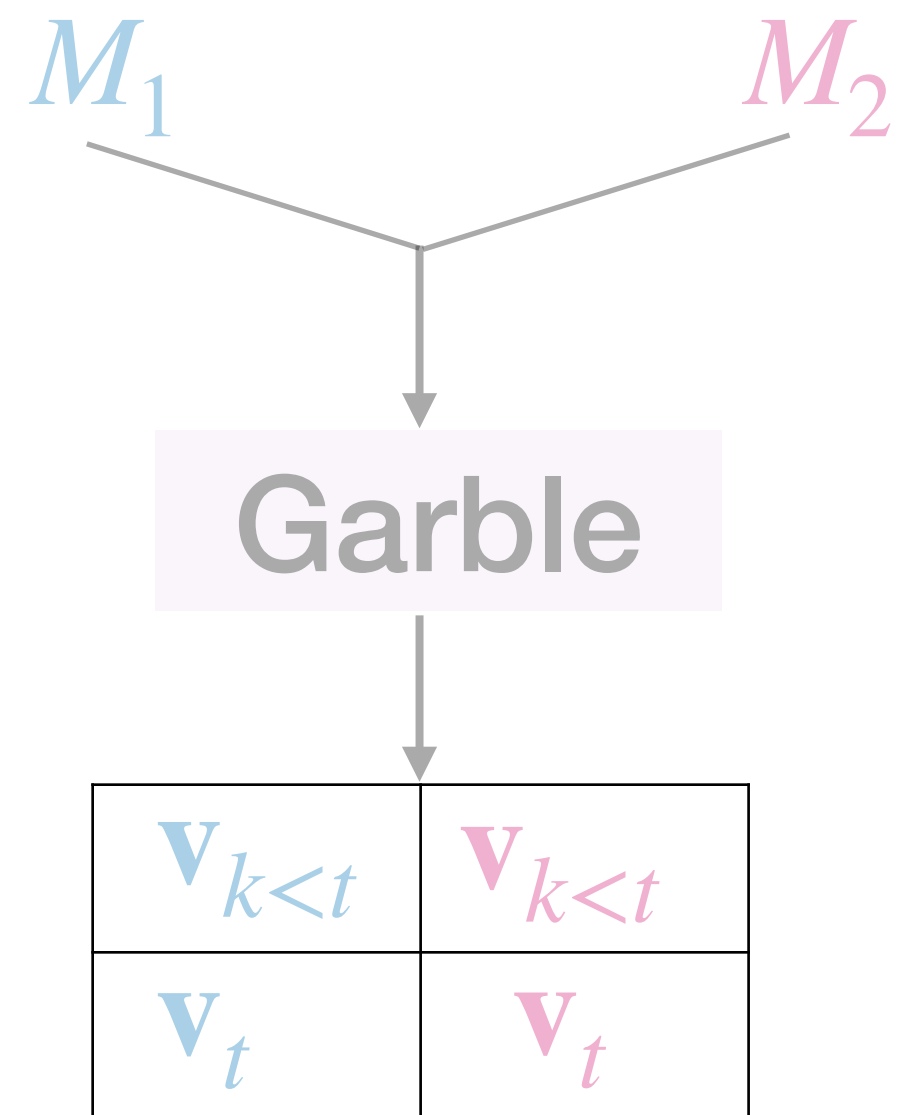
input

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$$z \mathbf{e}_1^T \mathbf{M}_{x_1} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$= \beta + \mathbf{e}_1^T \mathbf{r}_0 + \mathbf{e}_1^T (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1) + \mathbf{e}_1^T \mathbf{M}_{x_1} (-\mathbf{r}_1 + z \mathbf{e}_{q_{\text{acc}}})$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

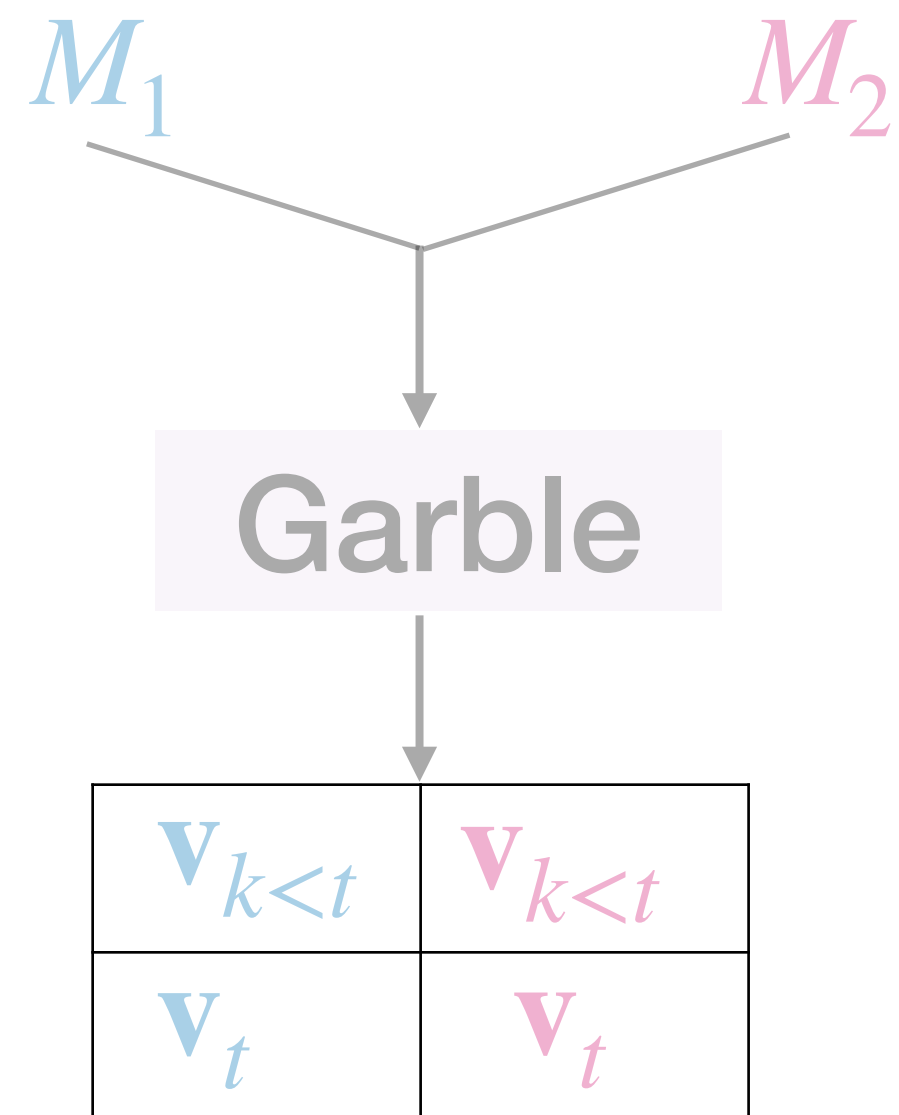
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$$\mathbf{z} \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$= \underbrace{\beta + \mathbf{e}_1^\top \mathbf{r}_0}_{L_0(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1)}_{L_1(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top \mathbf{M}_{x_1} (-\mathbf{r}_1 + \mathbf{z} \mathbf{e}_{q_{\text{acc}}})}_{L_2(\mathbf{z})}$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

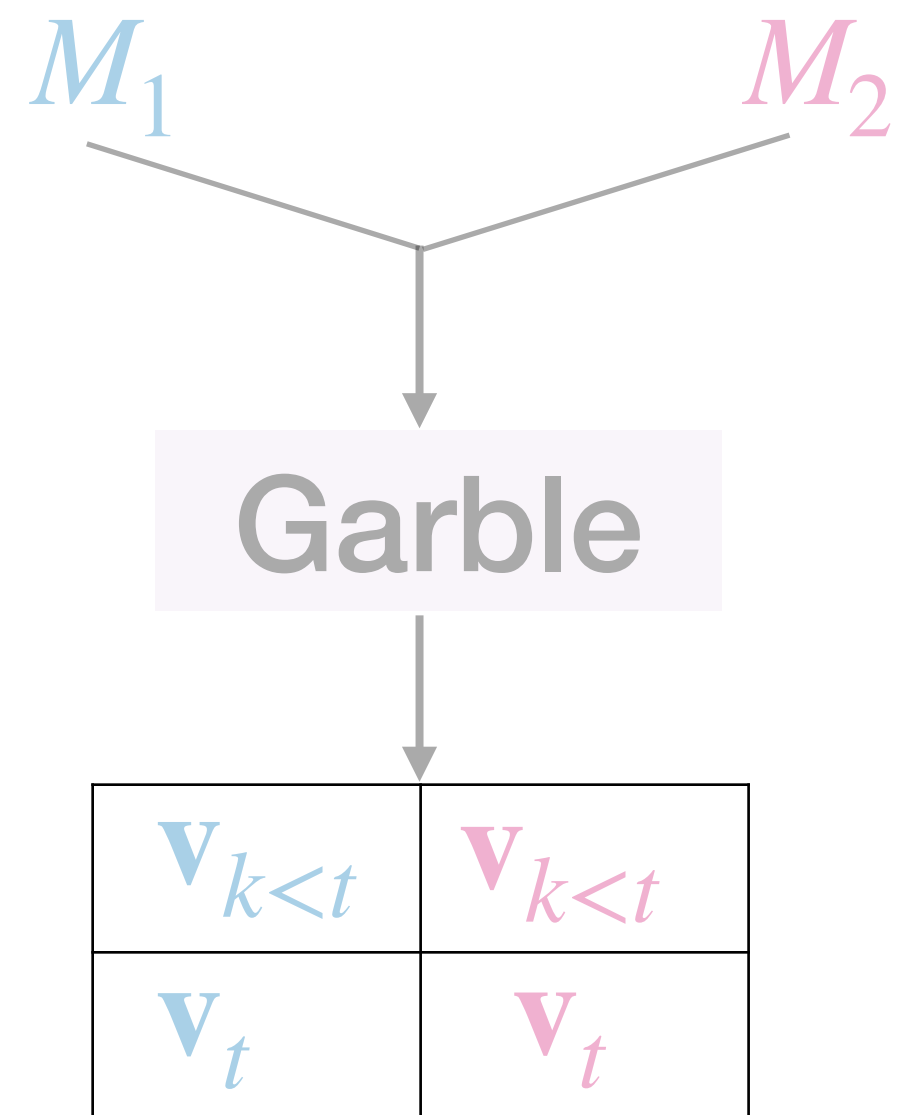
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$$z \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$= \underbrace{\beta}_{L_0(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top \mathbf{r}_0}_{L_1(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1)}_{L_1(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_1 + \mathbf{M}_{x_2} \mathbf{r}_2)}_{L_2(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top \mathbf{M}_{x_2} (-\mathbf{r}_2 + z \mathbf{e}_{q_{\text{acc}}})}_{L_3(z)}$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

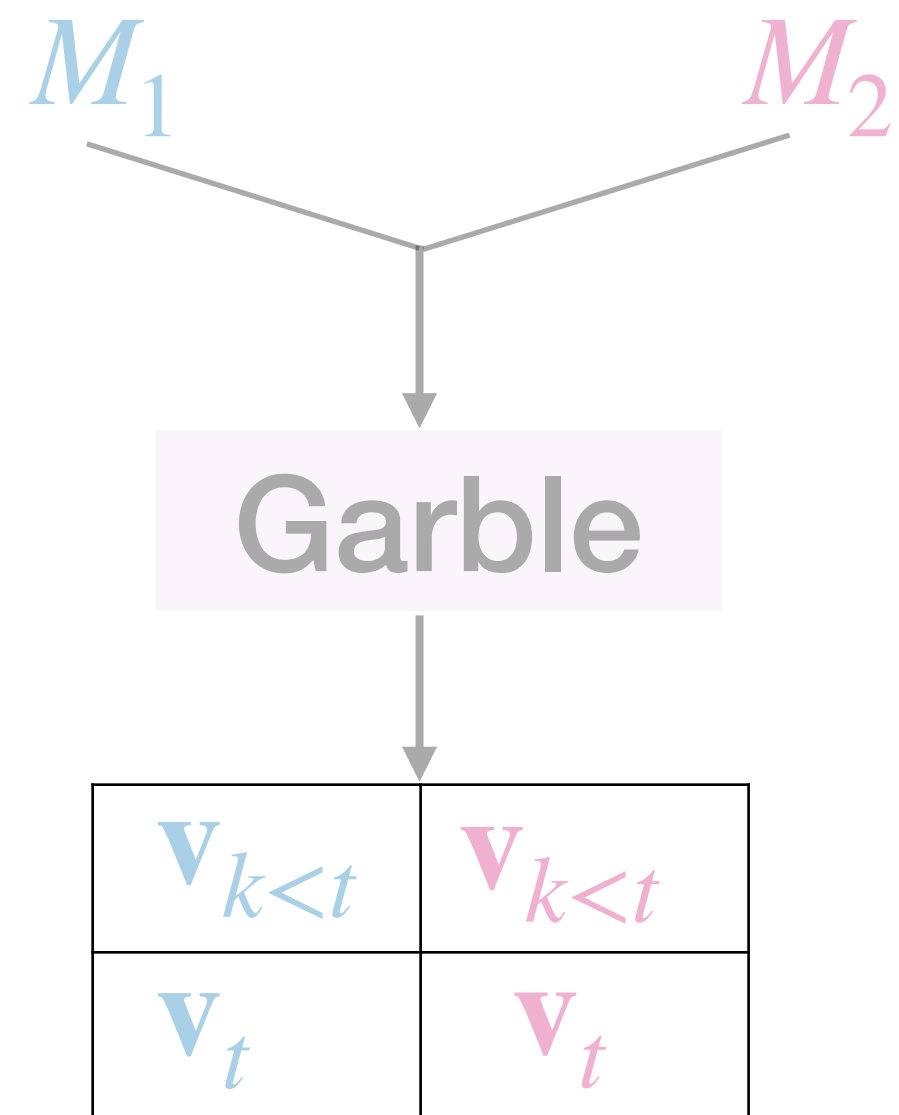
input

output

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$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$\mathbf{z} \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$= \underbrace{\beta + \mathbf{e}_1^\top \mathbf{r}_0}_{L_0(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1)}_{L_1(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_1 + \mathbf{M}_{x_2} \mathbf{r}_2)}_{L_2(\mathbf{x})} + \cdots + \underbrace{\mathbf{e}_1^\top \mathbf{M}_{x_N} (-\mathbf{r}_N + \mathbf{z} \mathbf{e}_{q_{\text{acc}}})}_{L_{N+1}(z)}$$

$$\text{total size} = NQ + 1$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

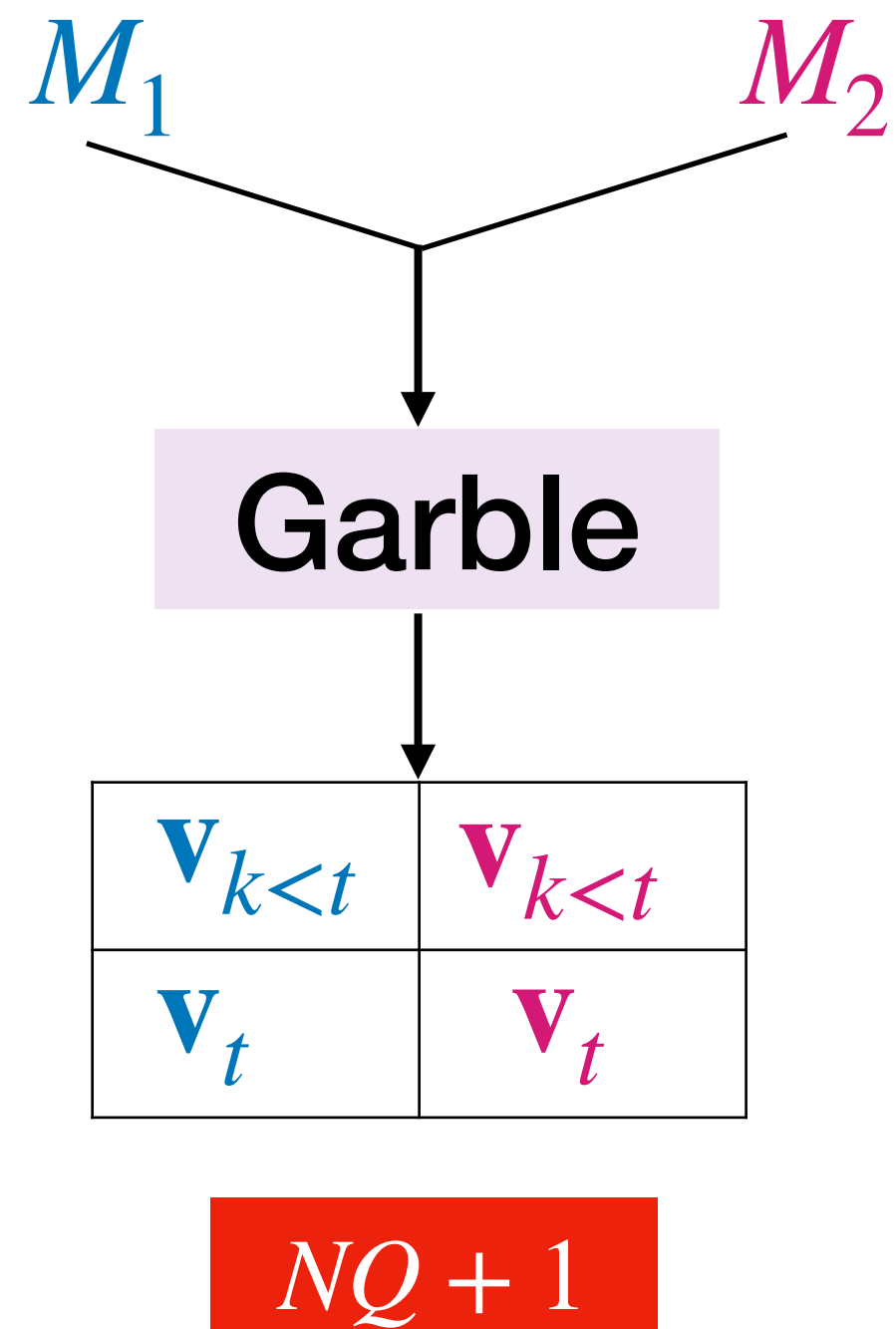
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } z \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$z \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$= \underbrace{\beta + \mathbf{e}_1^\top \mathbf{r}_0}_{L_0(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1)}_{L_1(\mathbf{x})} + \underbrace{\mathbf{e}_1^\top (-\mathbf{r}_1 + \mathbf{M}_{x_2} \mathbf{r}_2)}_{L_2(\mathbf{x})} + \cdots + \underbrace{\mathbf{e}_1^\top \mathbf{M}_{x_N} (-\mathbf{r}_N + z \mathbf{e}_{q_{\text{acc}}})}_{L_{N+1}(z)}$$

$$\text{total size} = NQ + 1$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

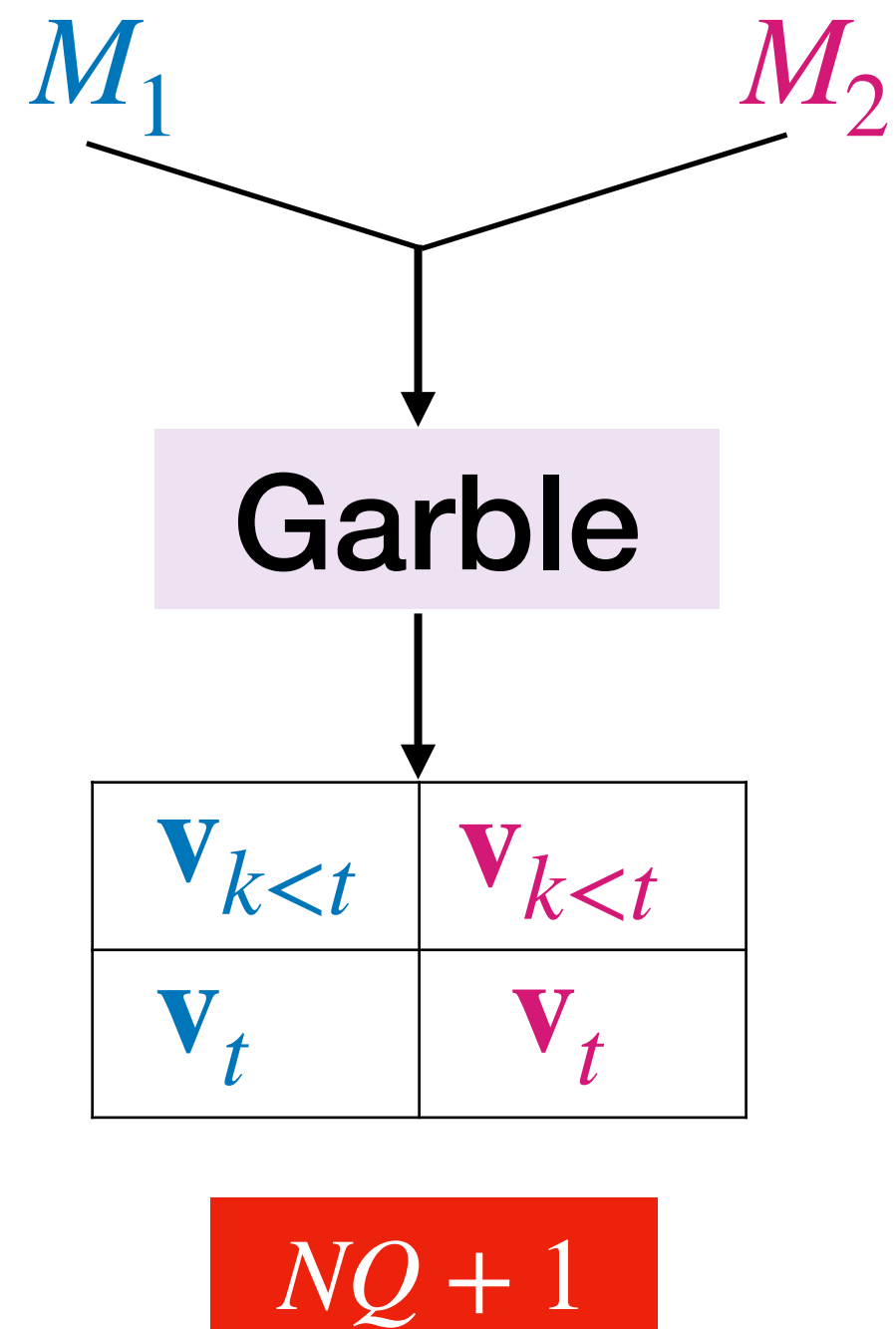
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$L_{q,j}(\mathbf{x}) = -\mathbf{r}_{j-1}[q] + (\mathbf{M}_{x_j} \mathbf{r}_j)[q]$$

$$\beta + \mathbf{e}_1^\top \mathbf{r}_0 + \mathbf{e}_1^\top (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1) + \mathbf{e}_1^\top (-\mathbf{r}_1 + \mathbf{M}_{x_2} \mathbf{r}_2) + \cdots + \mathbf{e}_1^\top \mathbf{M}_{x_N} (-\mathbf{r}_N + z \mathbf{e}_{q_{\text{acc}}})$$



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

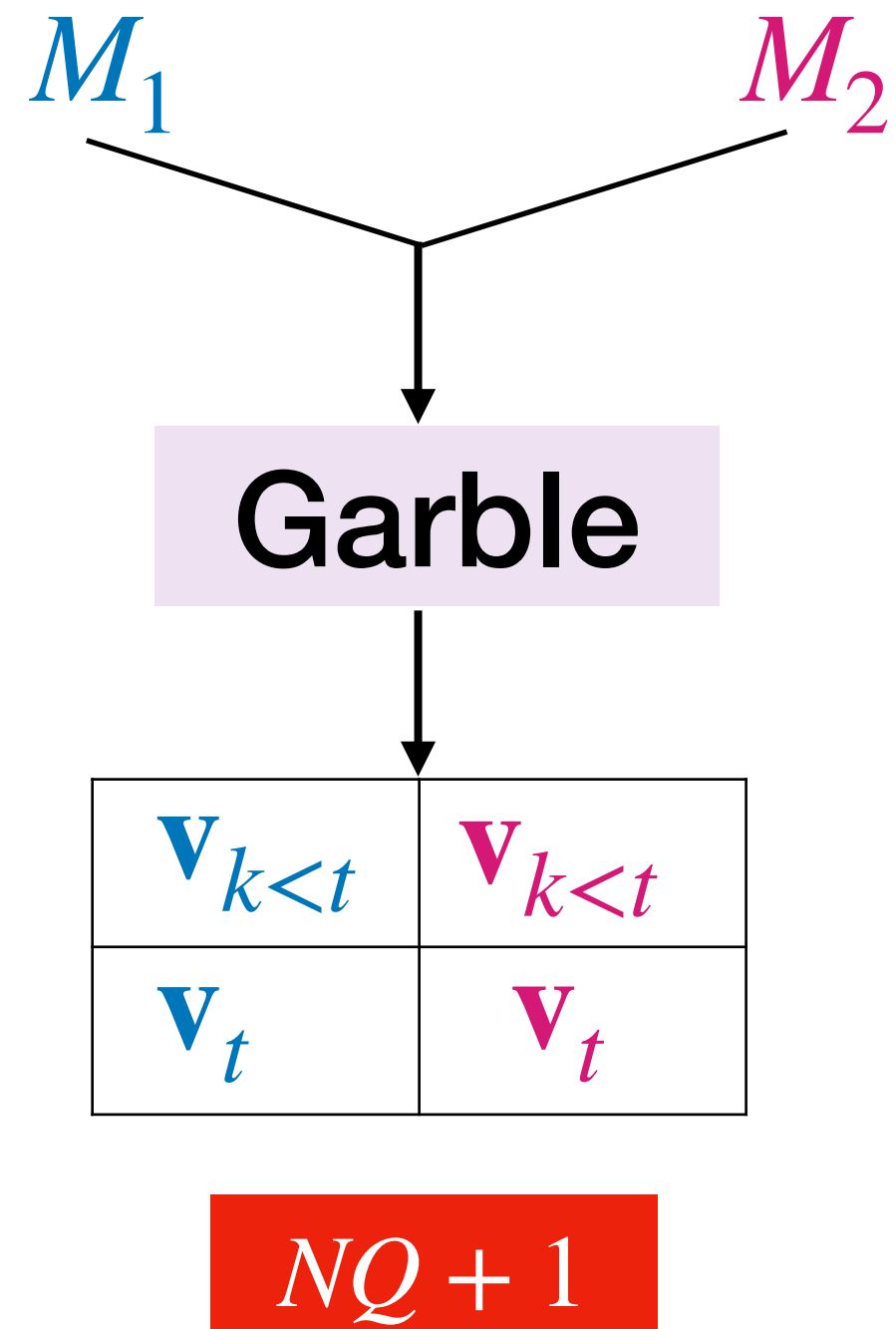
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

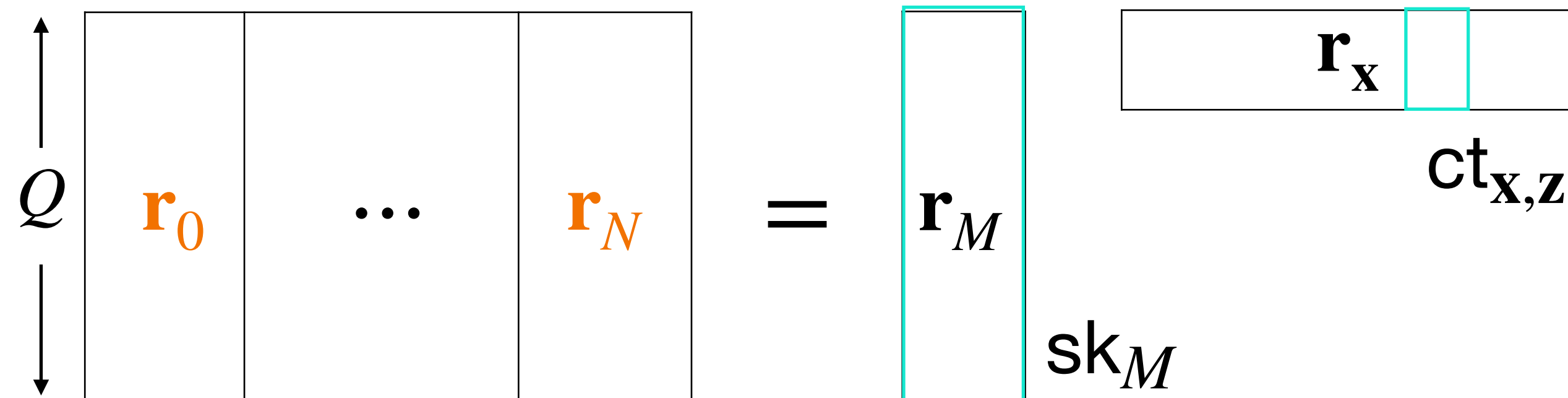


AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

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$$\mathbf{r}_j = \mathbf{r}_x[j] \cdot \mathbf{r}_M$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

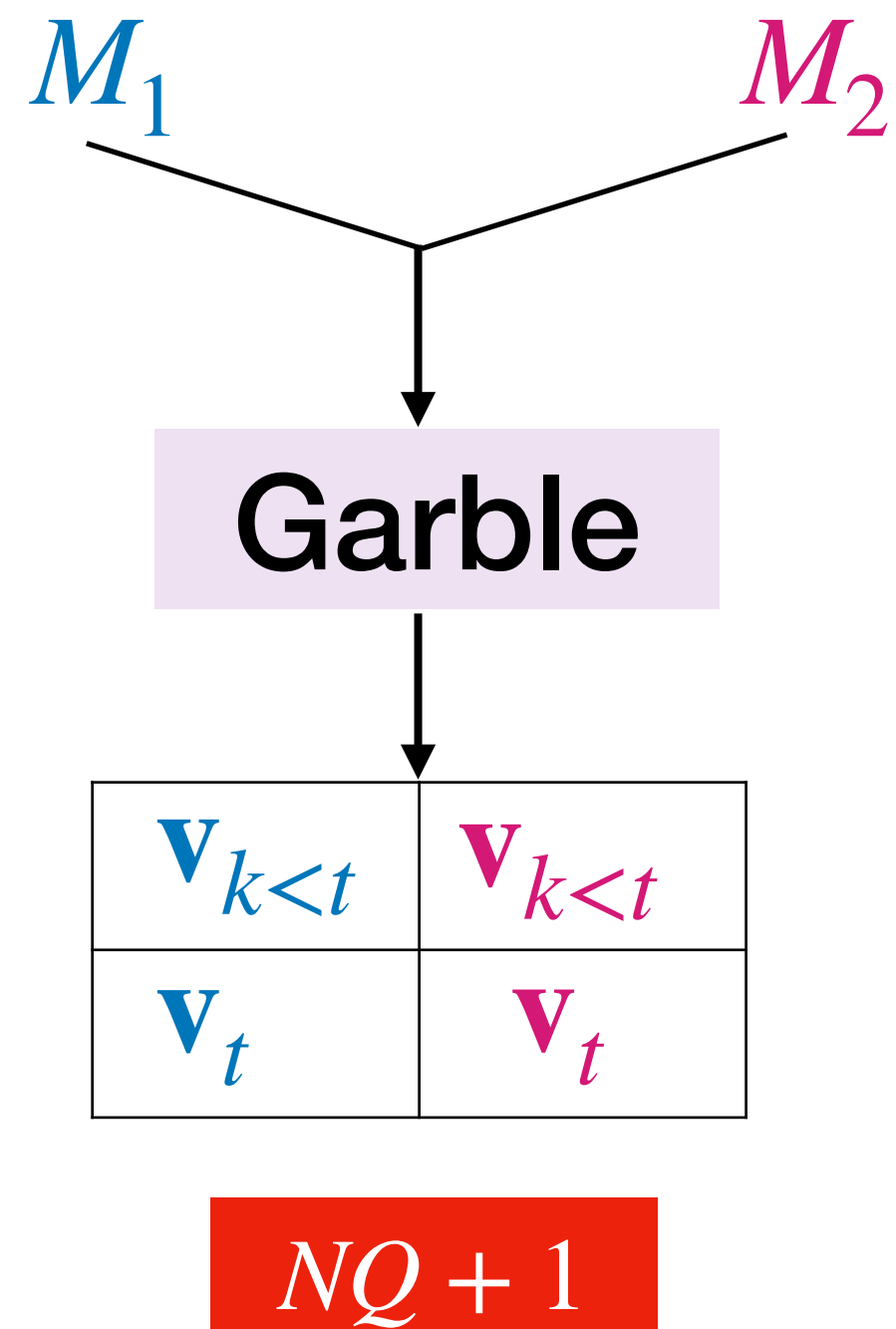
input

output

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$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



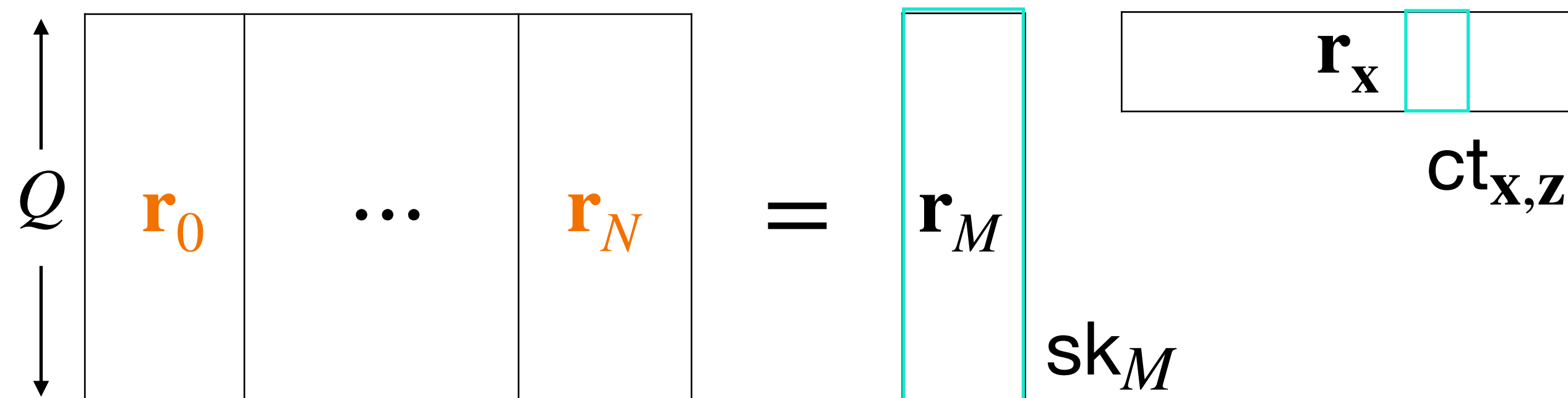
AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$L_{q,j}(\mathbf{x}) = -\mathbf{r}_{j-1}[q] + (\mathbf{M}_{x_j} \mathbf{r}_j)[q]$$

$$= (\mathbf{r}_x[j-1], \mathbf{r}_x[j]) \cdot (-\mathbf{r}_M[q], (\mathbf{M}_{x_j} \mathbf{r}_M)[q])$$

$$\beta + \mathbf{e}_1^\top \mathbf{r}_0 + \mathbf{e}_1^\top (-\mathbf{r}_0 + \mathbf{M}_{x_1} \mathbf{r}_1) + \mathbf{e}_1^\top (-\mathbf{r}_1 + \mathbf{M}_{x_2} \mathbf{r}_2) + \cdots + \mathbf{e}_1^\top \mathbf{M}_{x_N} (-\mathbf{r}_N + z \mathbf{e}_{q_{\text{acc}}})$$



$$\mathbf{r}_j = \mathbf{r}_x[j] \cdot \mathbf{r}_M$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

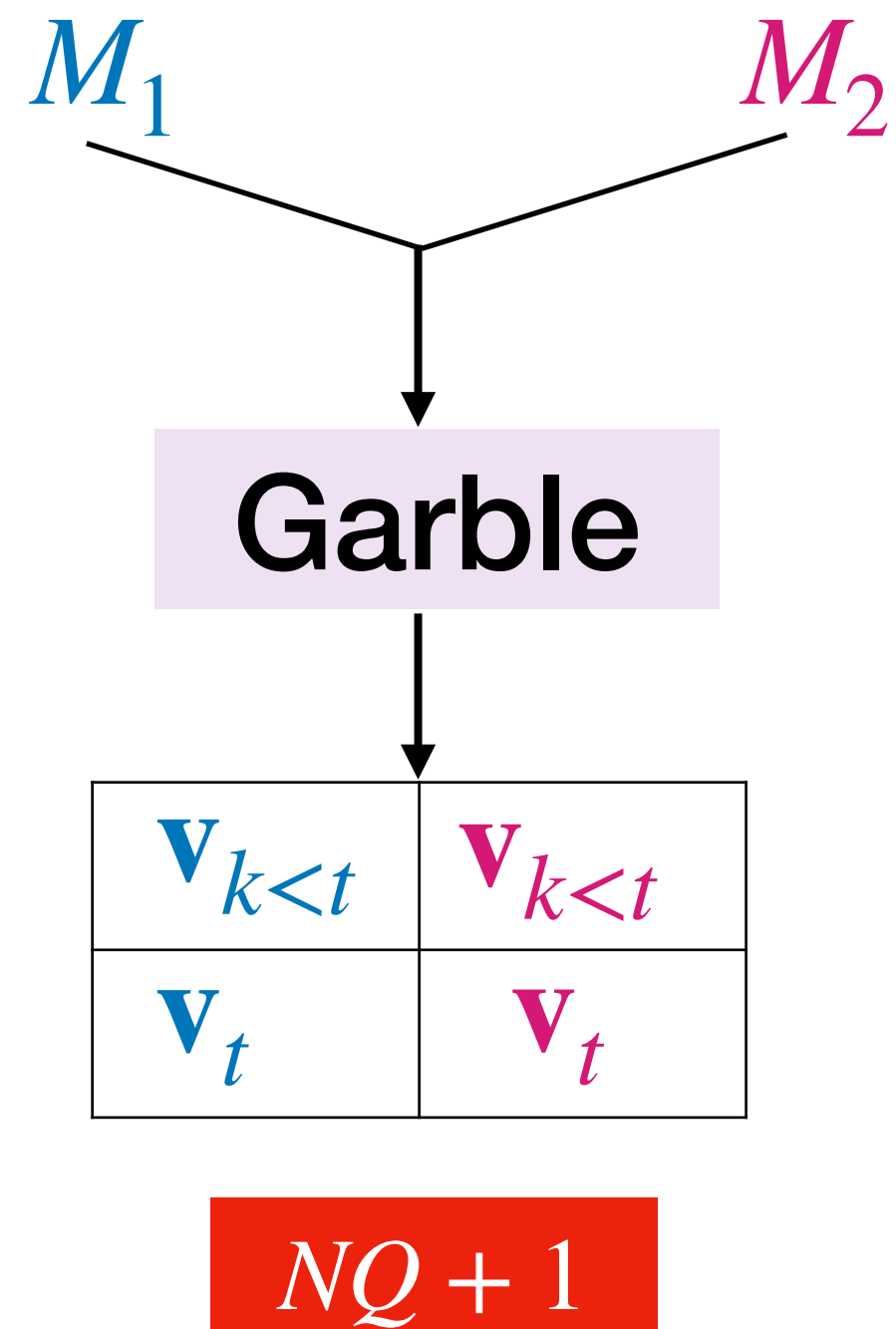
input

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$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

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AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$L_{q,j}(\mathbf{x}) = -\mathbf{r}_{j-1}[q] + (\mathbf{M}_{x_j} \mathbf{r}_j)[q]$$

$$= (\mathbf{r}_x[j-1], \mathbf{r}_x[j]) \cdot (-\mathbf{r}_M[q], (\mathbf{M}_{x_j} \mathbf{r}_M)[q])$$

Known to the **Encrypter**

Known to the **Key Generator**

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

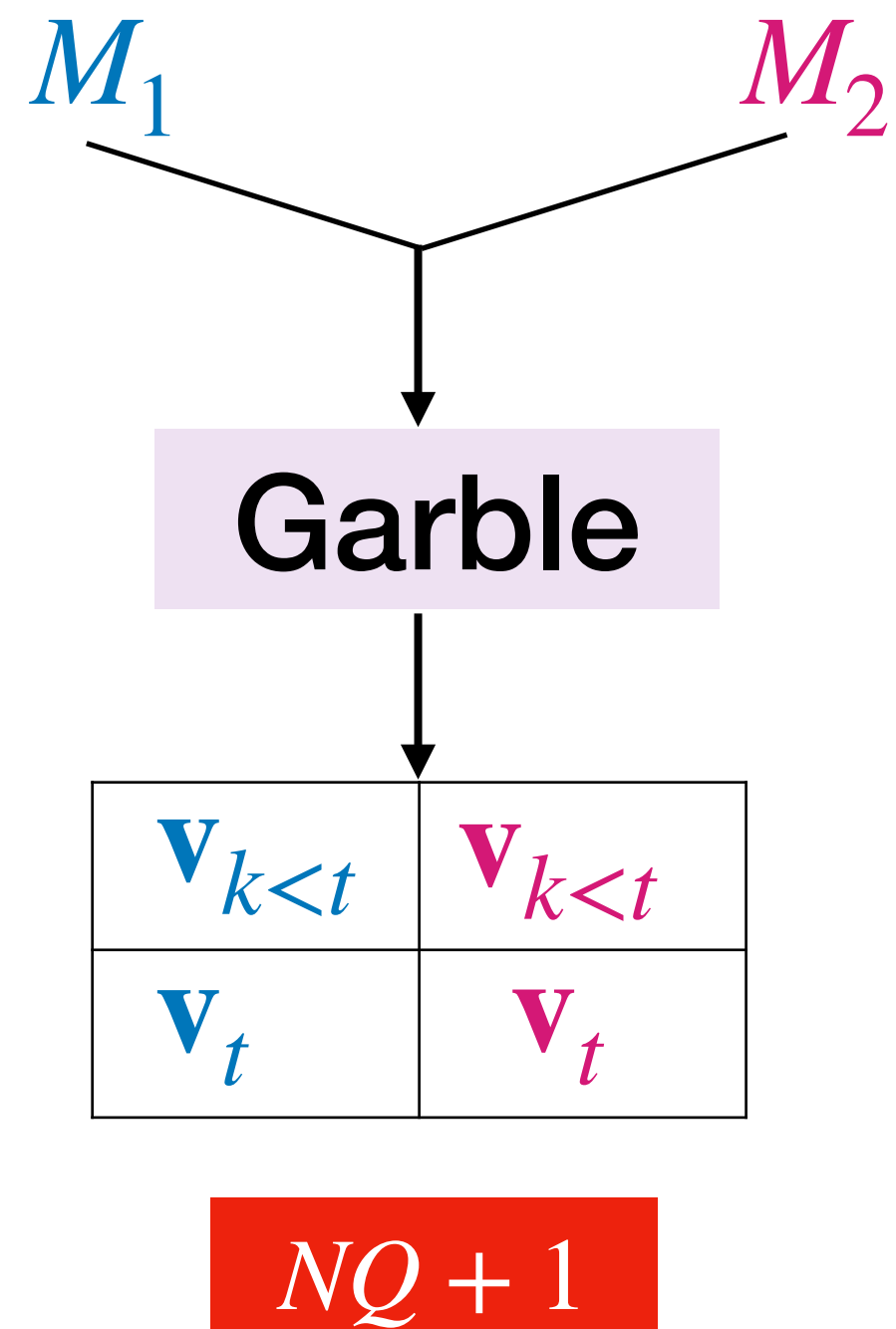
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



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$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$L_{q,j}(\mathbf{x}) = -\mathbf{r}_{j-1}[q] + (\mathbf{M}_{x_j} \mathbf{r}_j)[q] = \mathbf{u}_j \cdot \mathbf{v}_q$$

$$= (\mathbf{r}_x[j-1], \mathbf{r}_x[j]) \cdot (-\mathbf{r}_M[q], (\mathbf{M}_{x_j} \mathbf{r}_M)[q])$$

Known to the **Encrypter**

Known to the **Key Generator**

$$\mathbf{v}_q = ( -\mathbf{r}_M[q], (\mathbf{M}_{x_j} \mathbf{r}_M)[q] )$$

$$\mathbf{u}_j = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

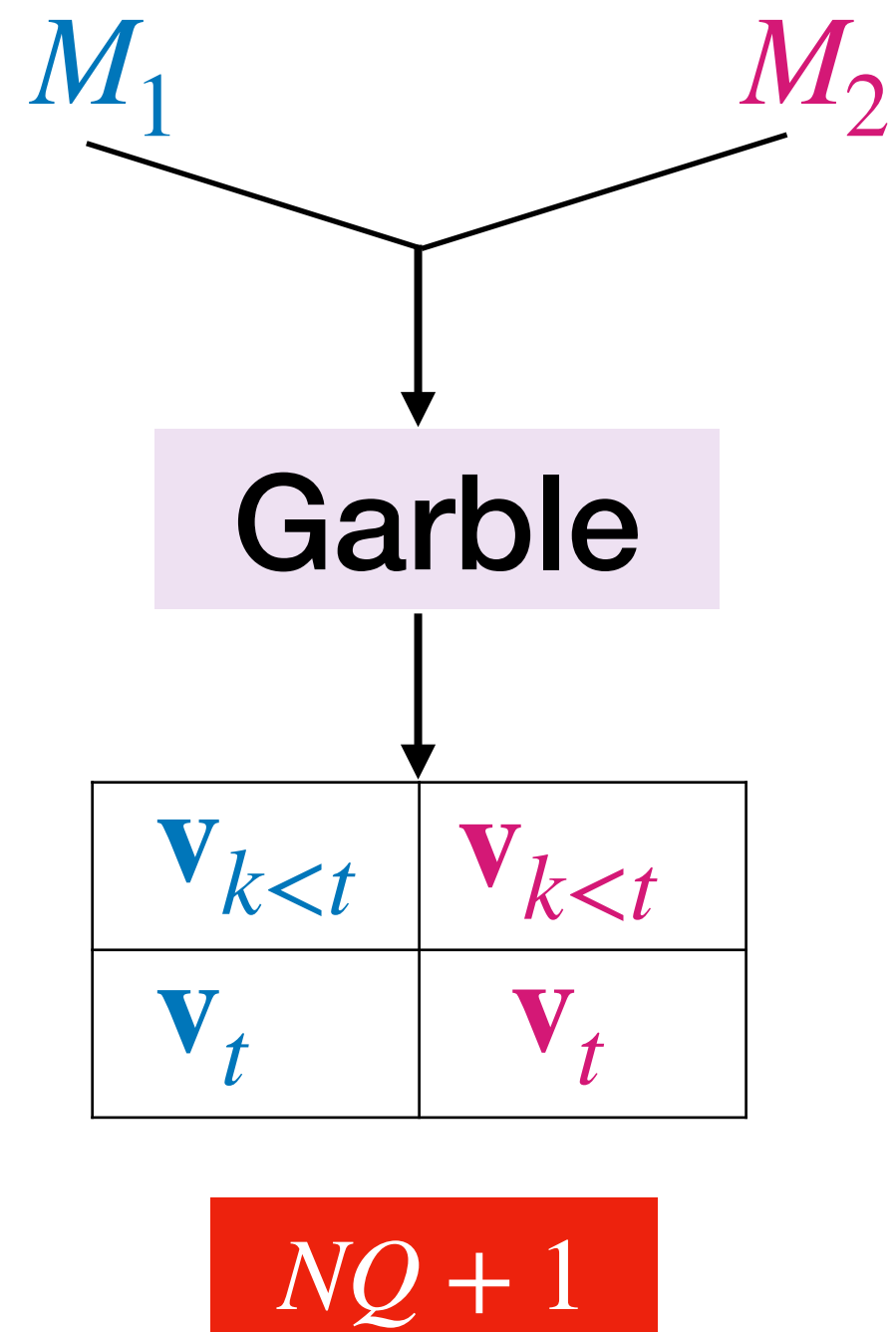
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } \mathbf{z} \cdot \mathbf{e}_1^\top \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$L_0(\mathbf{x}) = \beta + \mathbf{e}_1^\top \mathbf{r}_0 = \mathbf{u}_0 \cdot \mathbf{v}_0$$

$$= (\mathbf{r}_x[0], 1) \cdot (-\mathbf{r}_M[1], \beta)$$

$$\mathbf{v}_0 = (-\mathbf{r}_M[1], \beta)$$

$$\mathbf{v}_q = (-\mathbf{r}_M[q], (\mathbf{M}_{x_j} \mathbf{r}_M)[q])$$

$$\mathbf{u}_0 = (\mathbf{r}_x[0], 1)$$

$$\mathbf{u}_j = (\mathbf{r}_x[j-1], \mathbf{r}_x[j])$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

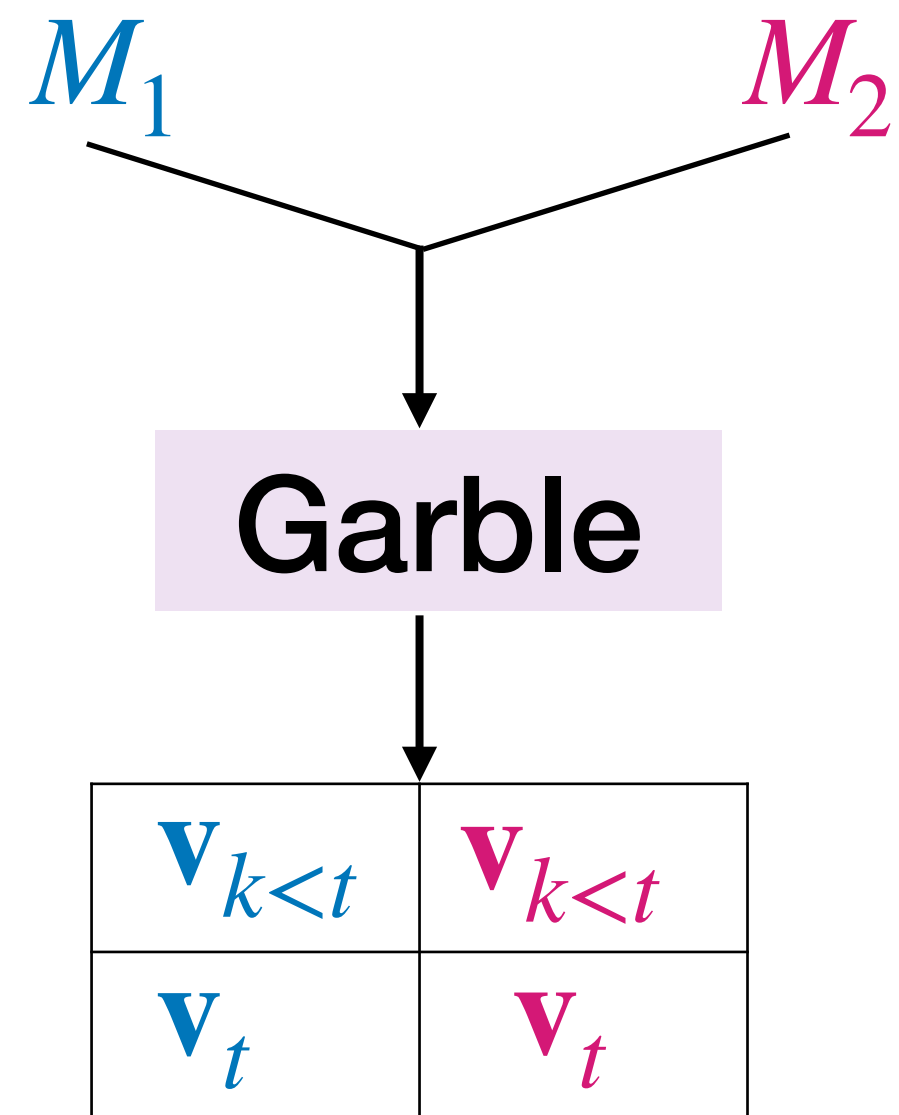
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



$$NQ + 1$$

AKGS for DFA  $\equiv$  AKGS for Matrix multiplication [LL20]

$$\text{Garble: } z \cdot \mathbf{e}_1^T \mathbf{M}_{x_1} \mathbf{M}_{x_2} \cdots \mathbf{M}_{x_N} \mathbf{e}_{q_{\text{acc}}} + \beta$$

$$L_{q,j}(z) = -\mathbf{r}_N[q] + z \mathbf{e}_{q_{\text{acc}}}[q] = \tilde{\mathbf{u}}_j \cdot \tilde{\mathbf{v}}_q$$

$$= (\mathbf{r}_x[N], z) \cdot (-\mathbf{r}_M[q], \mathbf{e}_{q_{\text{acc}}}[q])$$

process of garbling is distributed between **key generator** and **encrypter**

$$\mathbf{v}_0 = (-\mathbf{r}_M[1], \beta)$$

$$\mathbf{v}_q = (-\mathbf{r}_M[q], (\mathbf{M}_{x_j} \mathbf{r}_M)[q])$$

$$\tilde{\mathbf{v}}_q = (-\mathbf{r}_M[q], \mathbf{e}_{q_{\text{acc}}}[q])$$

$$\mathbf{u}_0 = (\mathbf{r}_x[0], 1)$$

$$\mathbf{u}_j = (\mathbf{r}_x[j-1], \mathbf{r}_x[j])$$

$$\tilde{\mathbf{u}}_j = (\mathbf{r}_x[N], z)$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$M_1$   $M_2$

$\mathbf{v}_0, \mathbf{v}_q$	$\mathbf{v}_0, \mathbf{v}_q$
$\tilde{\mathbf{v}}_q$	$\tilde{\mathbf{v}}_q$

IPFE, IPFE

$$sk_0, sk_q, sk_q, \tilde{sk}_0, \tilde{sk}_q, \tilde{sk}_q$$

$$\mathbf{v}_{0,k} = ( -\mathbf{r}_{M_k}[1], \beta_k )$$

$$\mathbf{v}_{q,k} = ( -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( -\mathbf{r}_{M_k}[q], \mathbf{e}_{q_{acc,k}}[q] )$$

$$\mathbf{u}_0 = ( \mathbf{r}_x[0], 1 )$$

$$\mathbf{u}_j = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

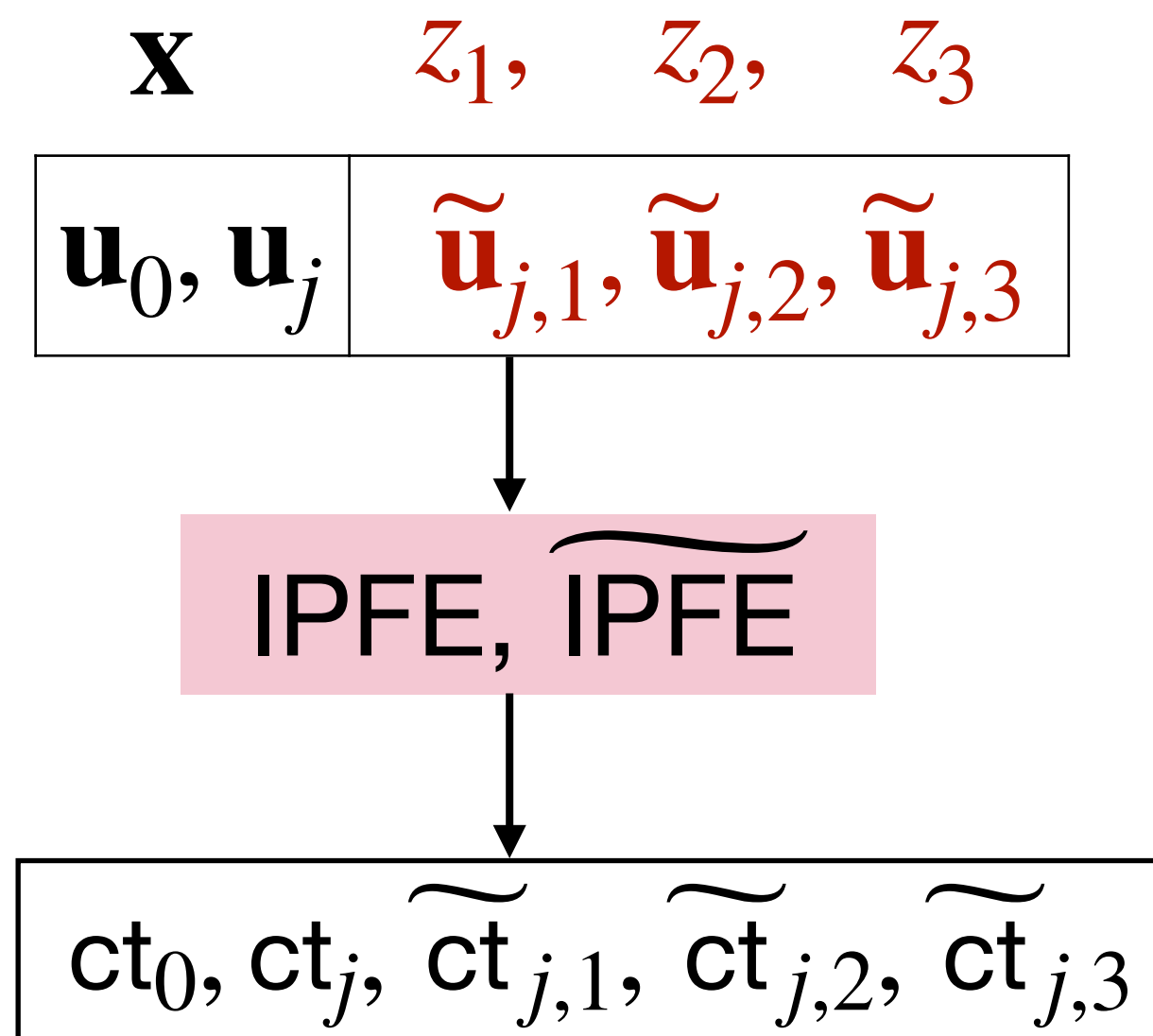
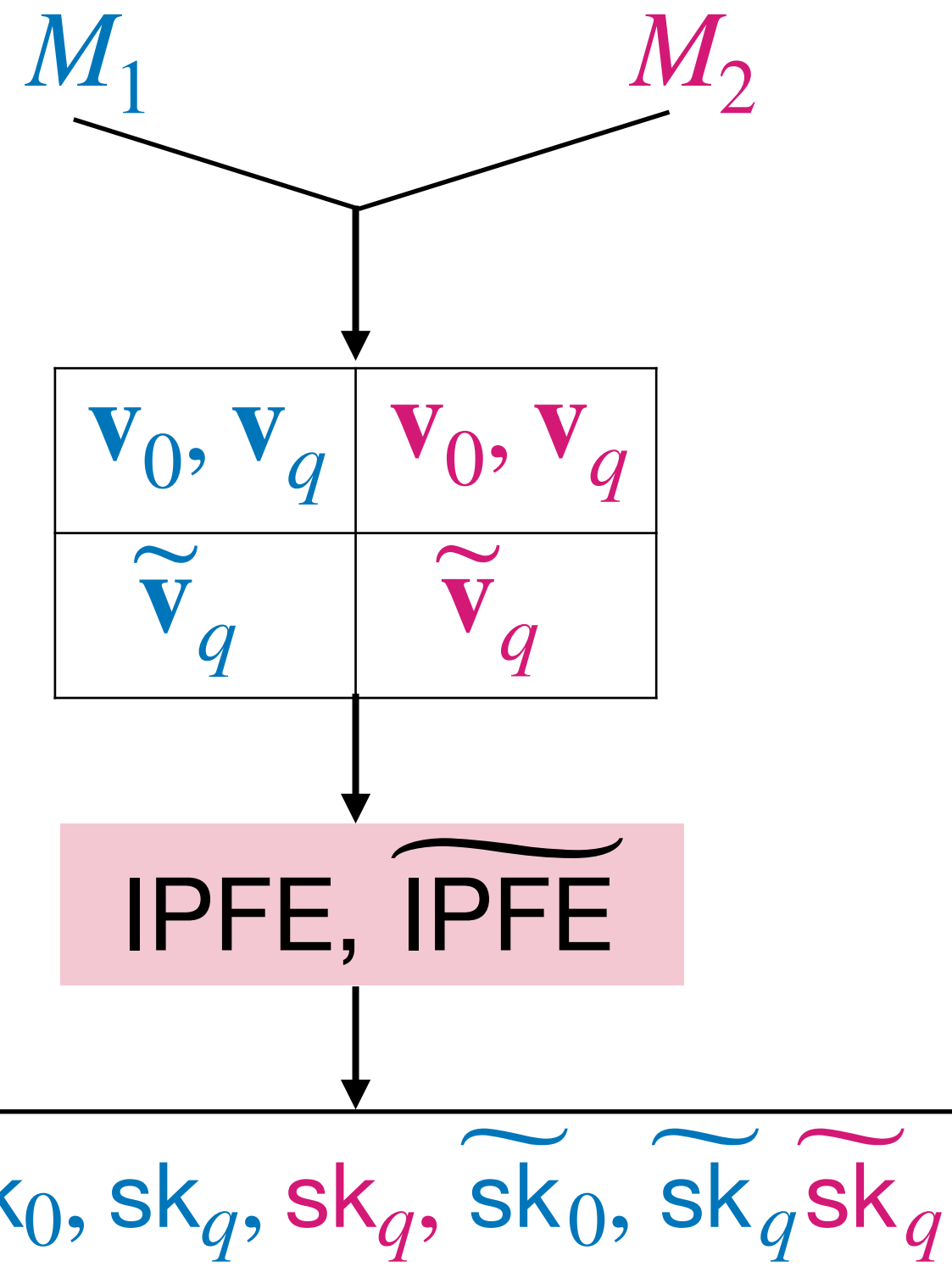
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



$$\mathbf{v}_{0,k} = ( -\mathbf{r}_{M_k}[1], \beta_k )$$

$$\mathbf{v}_{q,k} = ( -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( -\mathbf{r}_{M_k}[q], \mathbf{e}_{q_{acc,k}}[q] )$$

$$\mathbf{u}_0 = ( \mathbf{r}_x[0], 1 )$$

$$\mathbf{u}_j = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

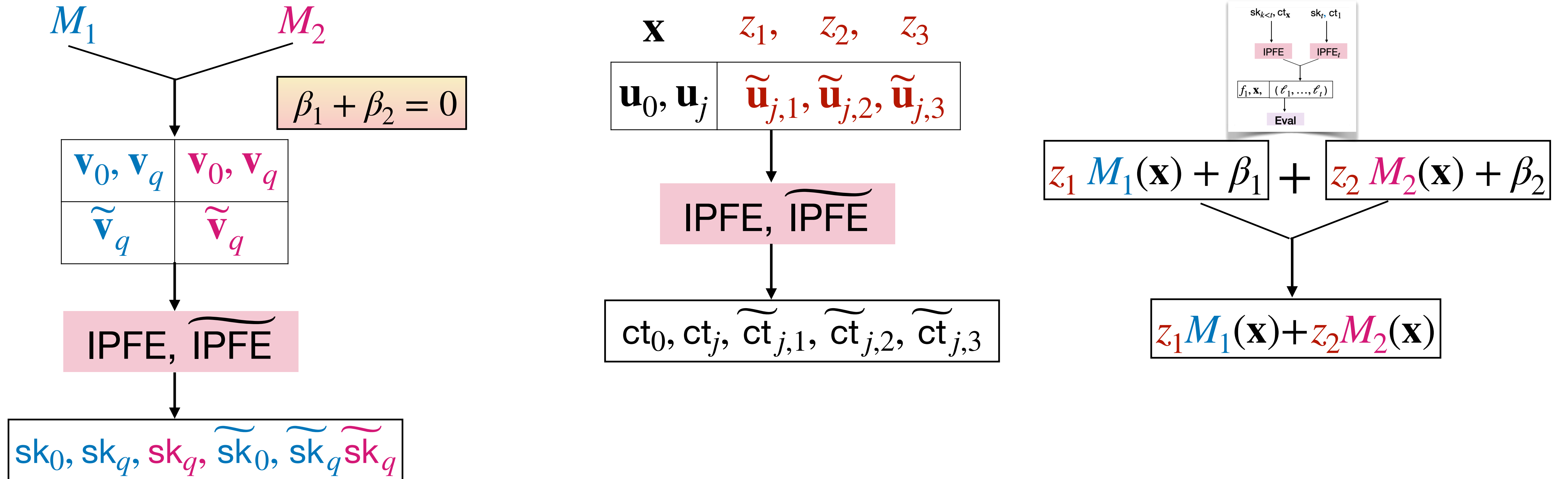
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



$$\mathbf{v}_{0,k} = ( -\mathbf{r}_{M_k}[1], \beta_k )$$

$$\mathbf{v}_{q,k} = ( -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( -\mathbf{r}_{M_k}[q], \mathbf{e}_{q_{acc,k}}[q] )$$

$$\mathbf{u}_0 = ( \mathbf{r}_x[0], 1 )$$

$$\mathbf{u}_j = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

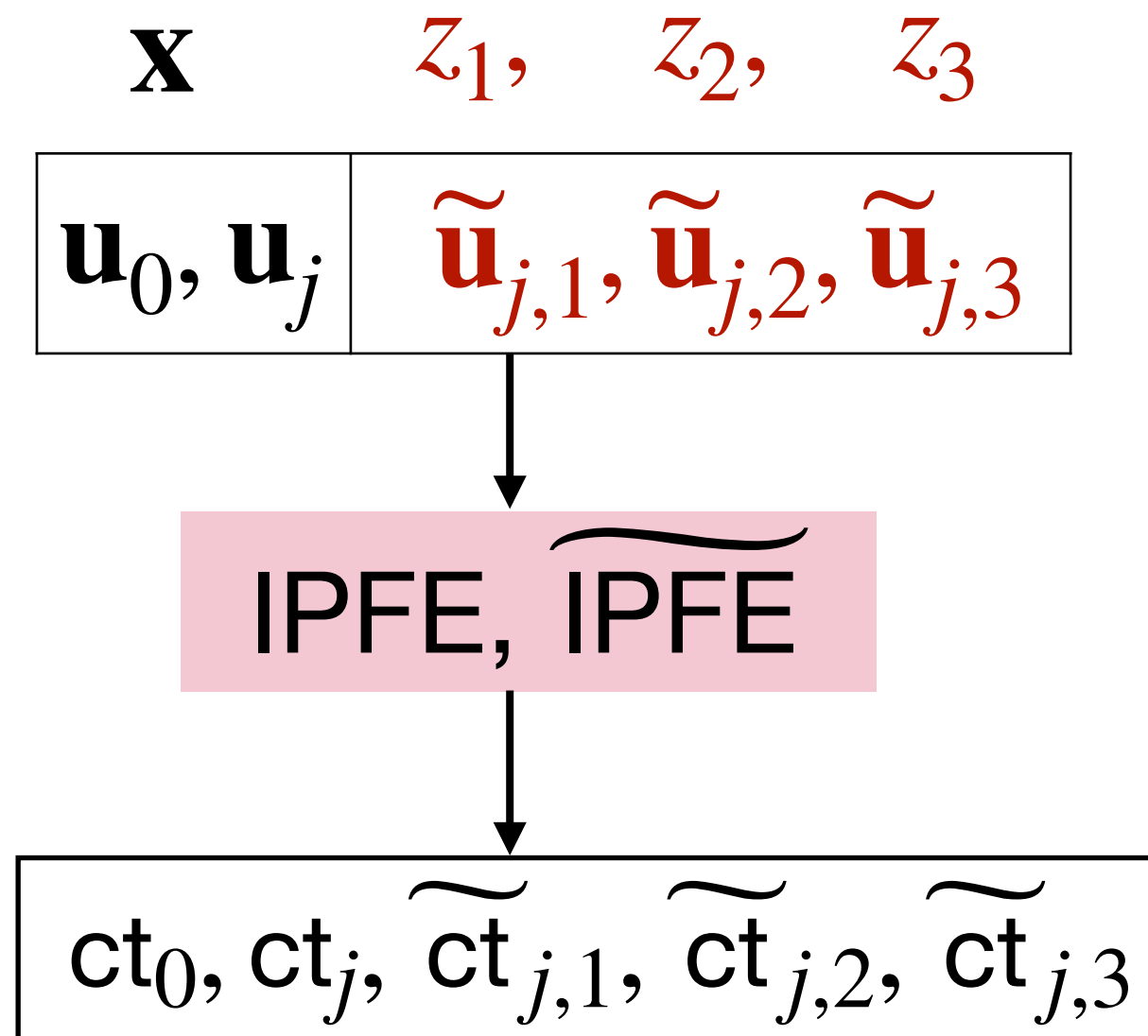
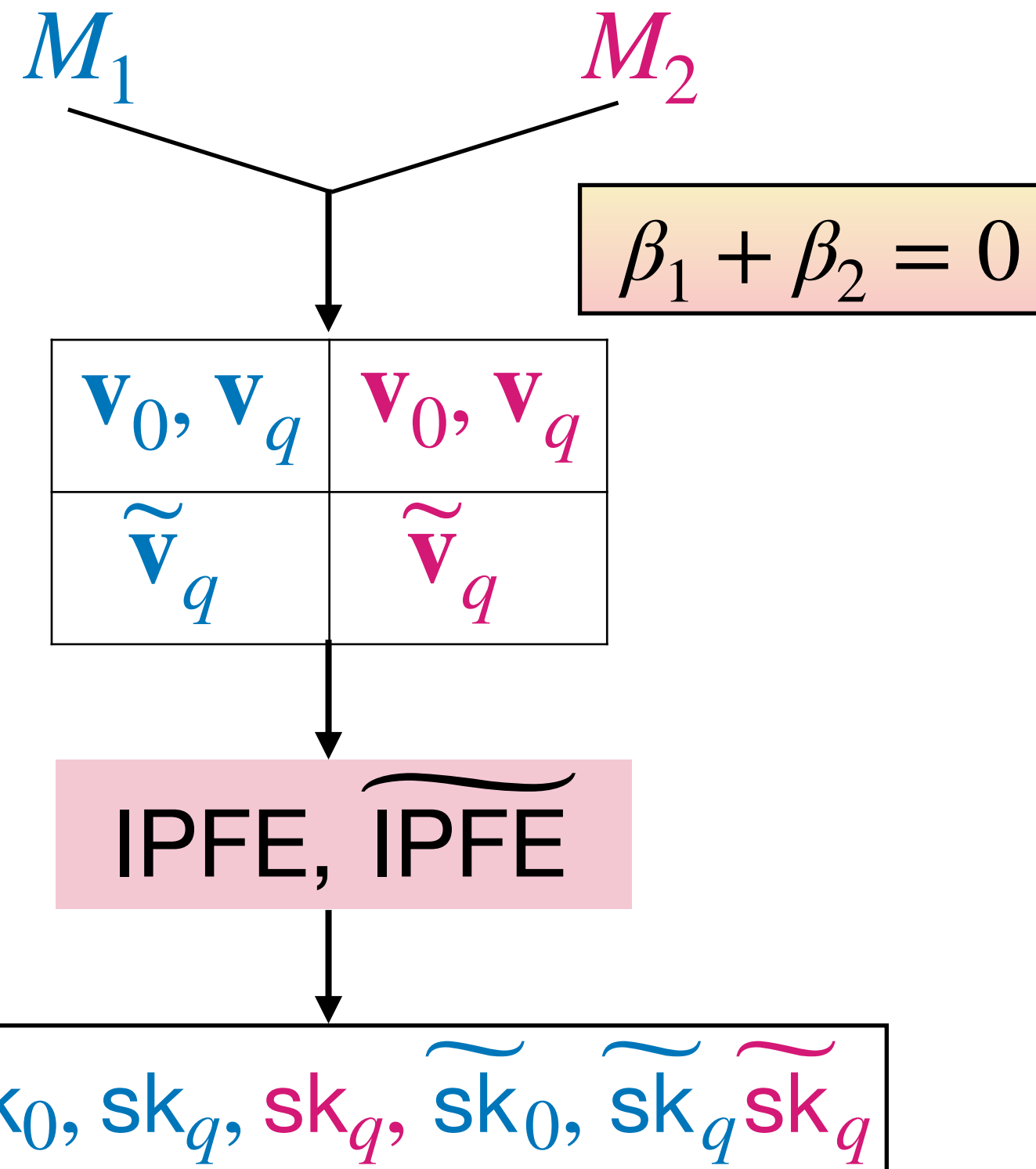
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

$$\mathbf{v}_{0,k} = ( -\mathbf{r}_{M_k}[1], \beta_k )$$

$$\mathbf{v}_{q,k} = ( -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j, k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( -\mathbf{r}_{M_k}[q], \mathbf{e}_{q_{\text{acc}, k}}[q] )$$

$$\mathbf{u}_0 = ( \mathbf{r}_x[0], 1 )$$

$$\mathbf{u}_j = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

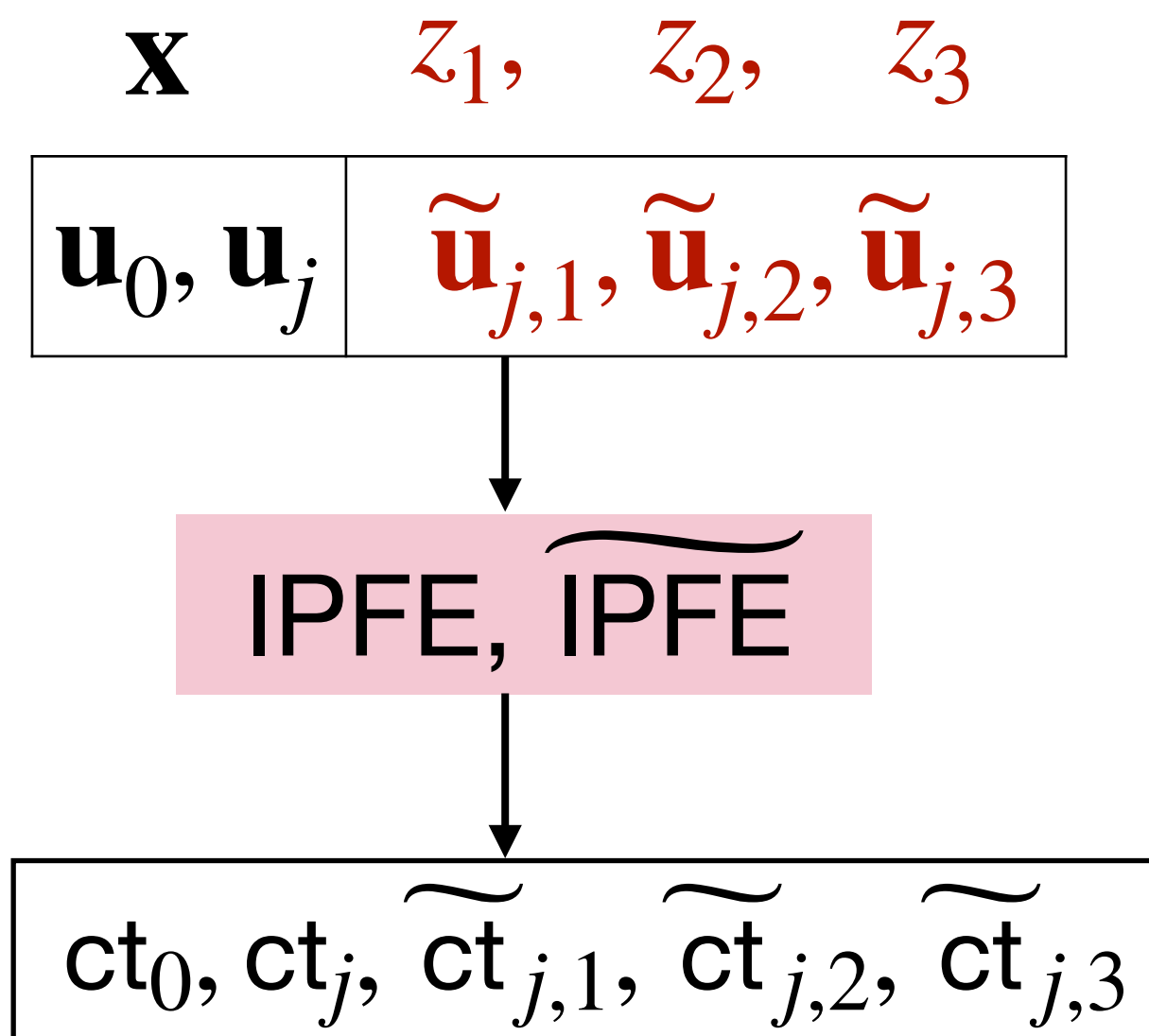
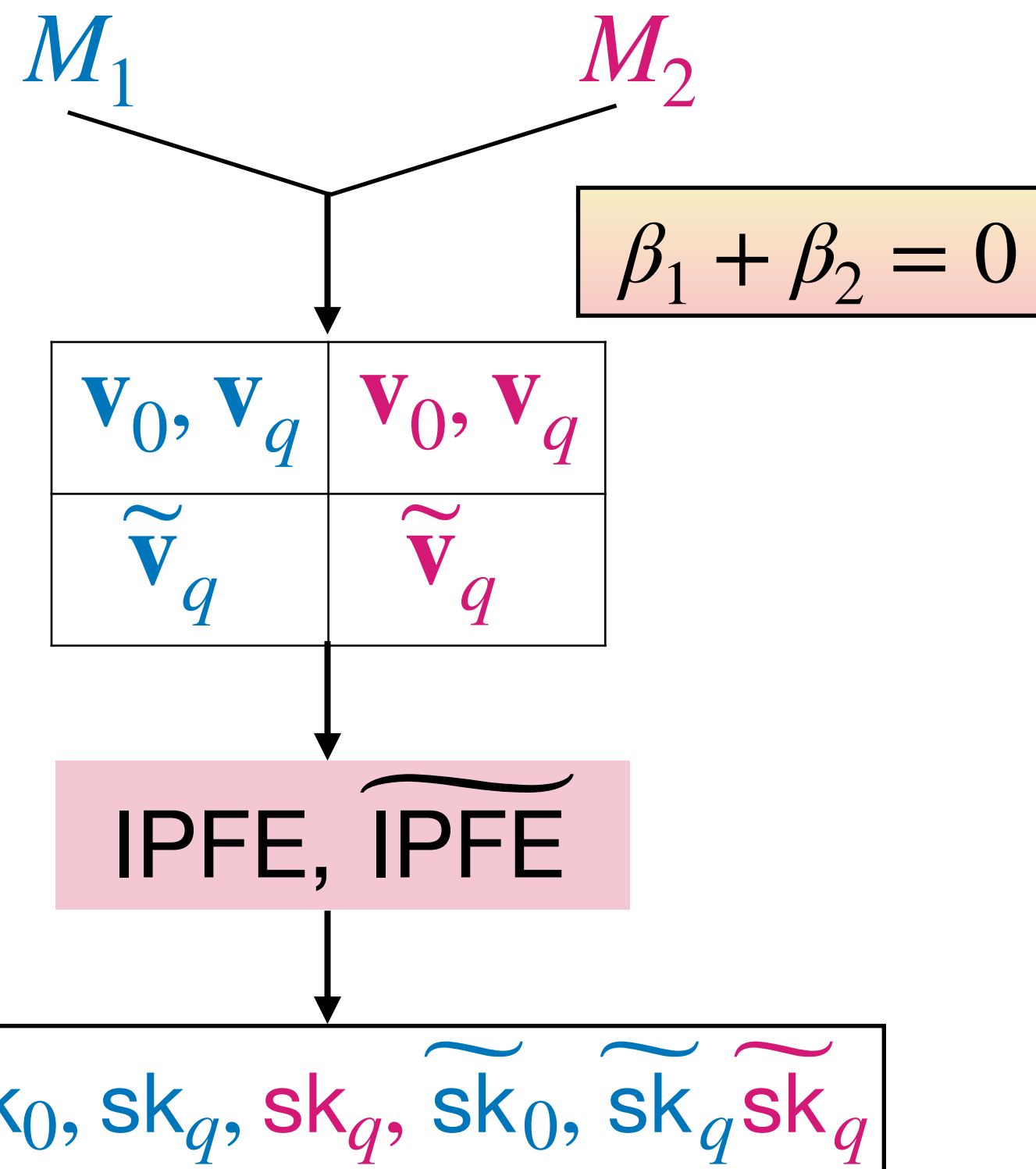
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## Steps: Simulation Security

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$$\tilde{\mathbf{v}}_{q,k} = ( -\mathbf{r}_{M_k}[q], \mathbf{e}_{q_{\text{acc},k}}[q] )$$

$$\mathbf{u}_0 = ( \mathbf{r}_x[0], 1 )$$

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$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$

not enough space

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

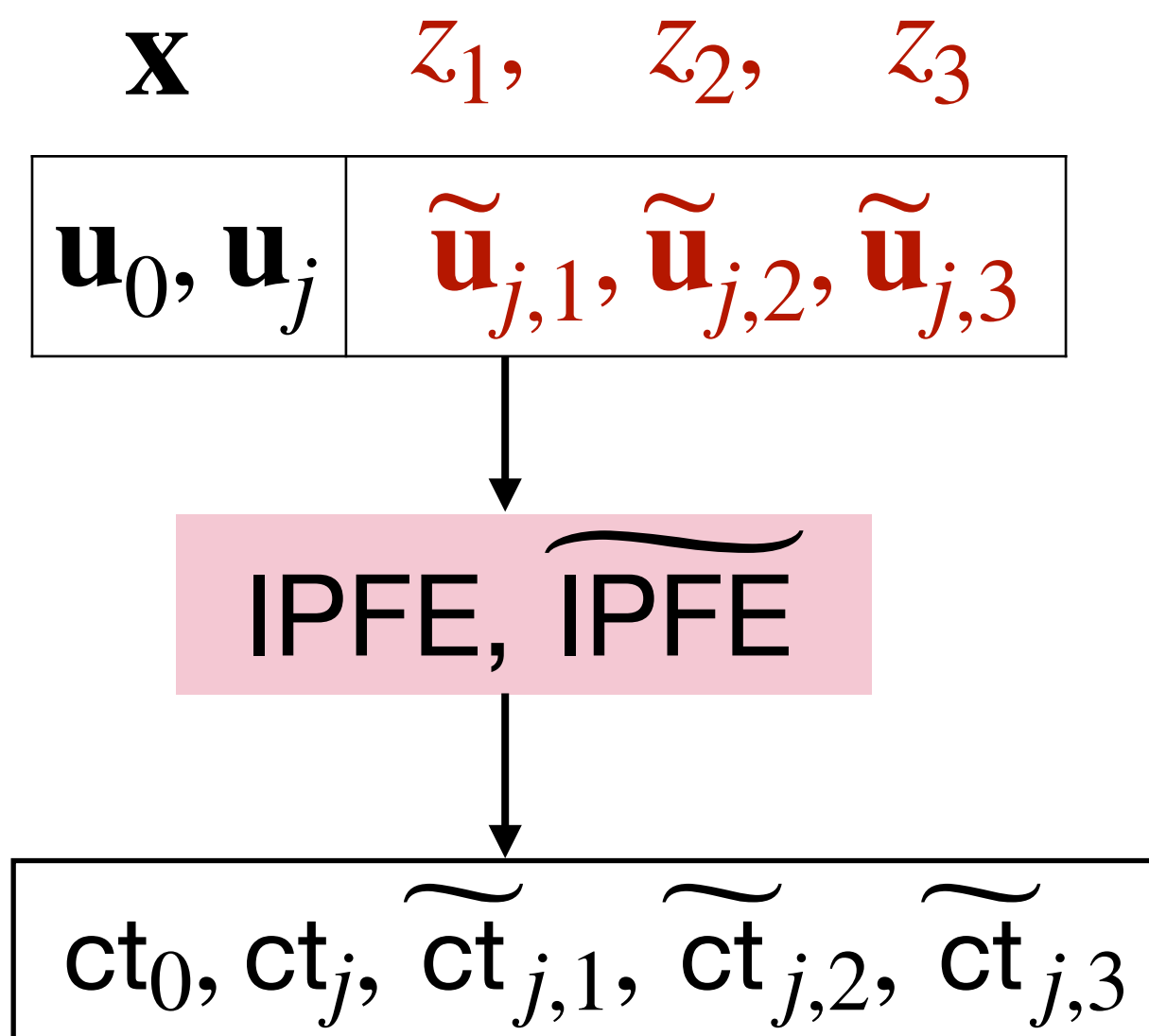
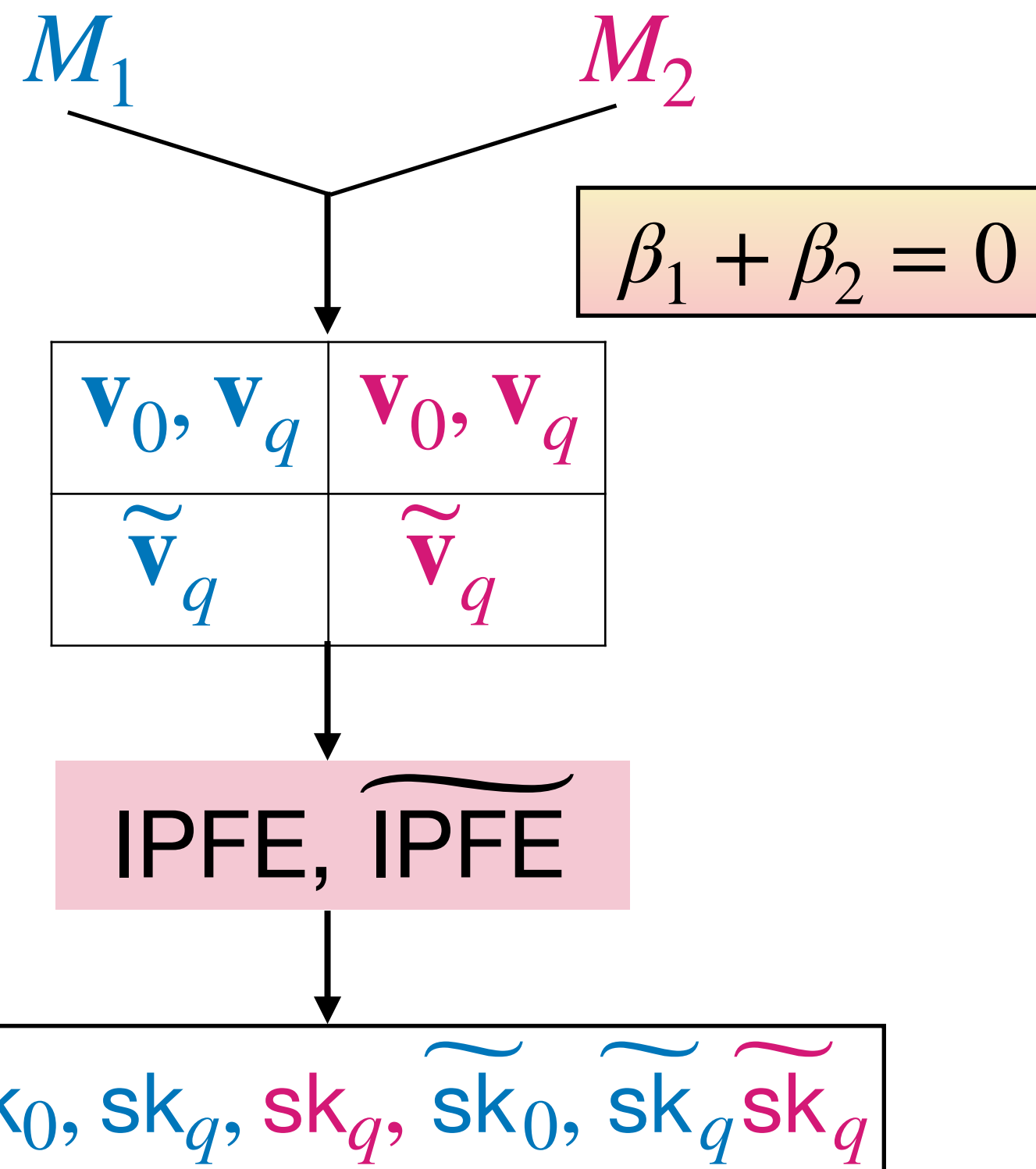
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( -\mathbf{r}_{M_k}[1], \beta_k )$$

$$\mathbf{v}_{q,k} = ( -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j, k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( -\mathbf{r}_{M_k}[q], \mathbf{e}_{q_{\text{acc}, k}}[q] )$$

$$\mathbf{u}_0 = ( \mathbf{r}_x[0], 1 )$$

$$\mathbf{u}_j = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$

not enough space

not enough space

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

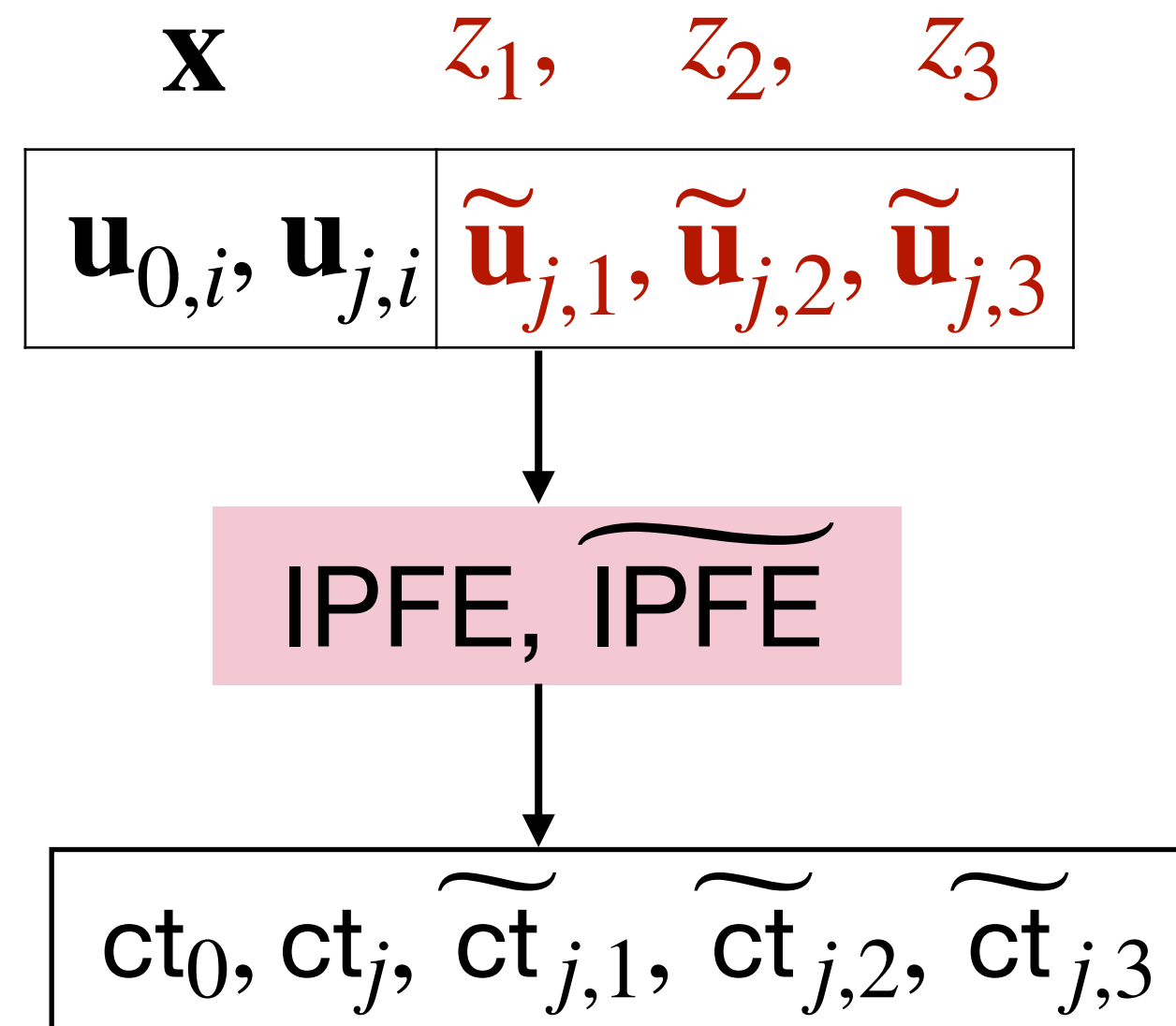
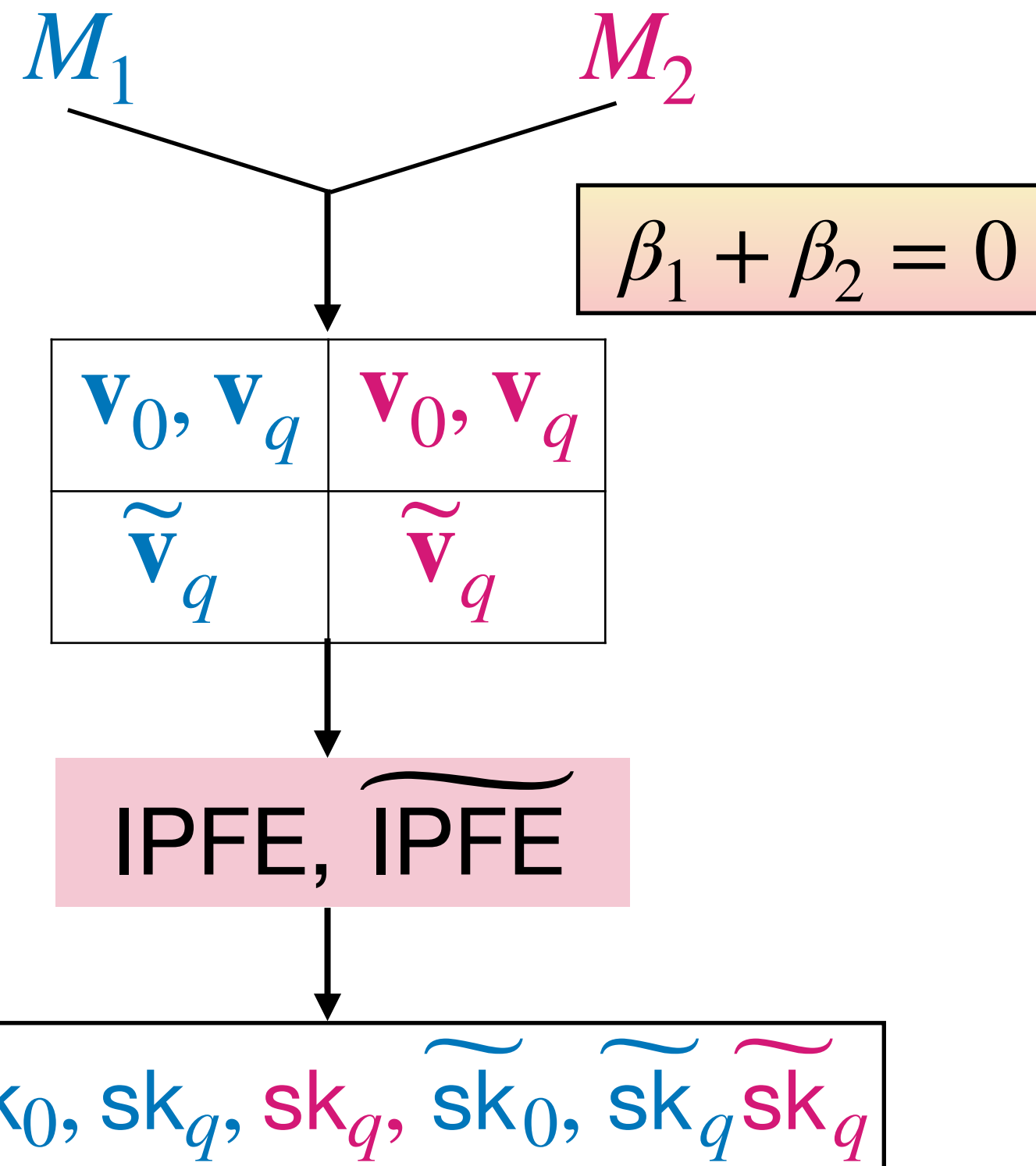
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

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$$\mathbf{u}_{0,i} = ( \mathbf{r}_x[0], 1 )$$

$$\mathbf{u}_{j,i} = ( \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \mathbf{r}_x[N], z_i )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

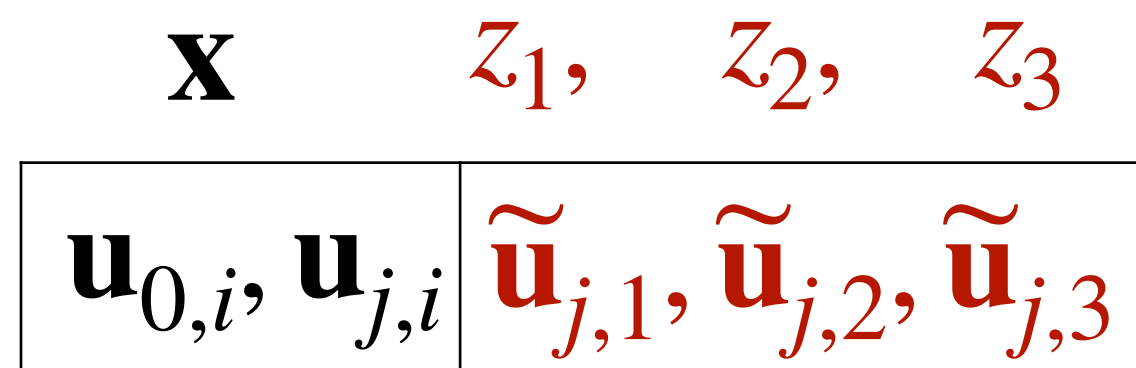
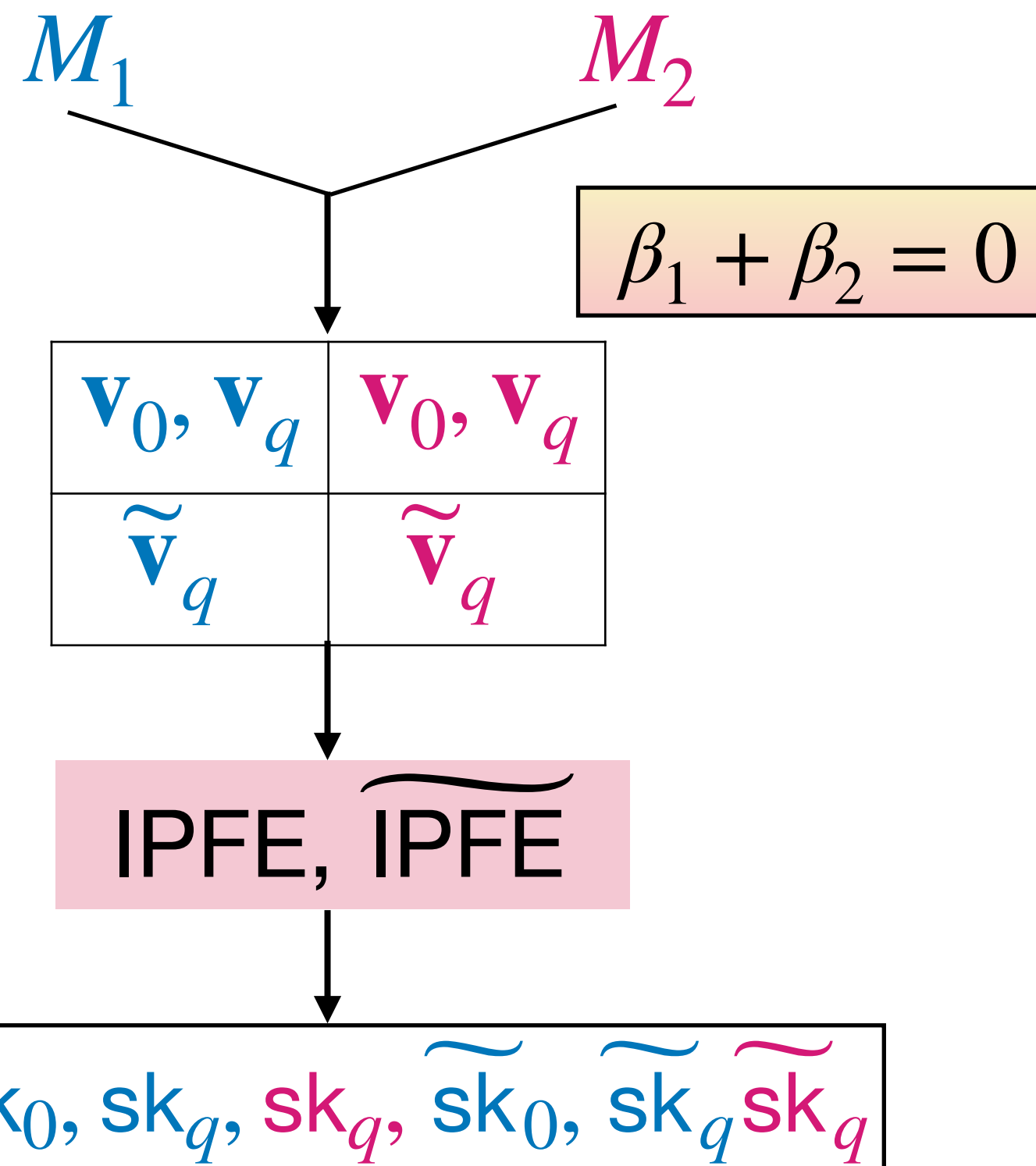
$$M = (M_1, M_2)$$

input

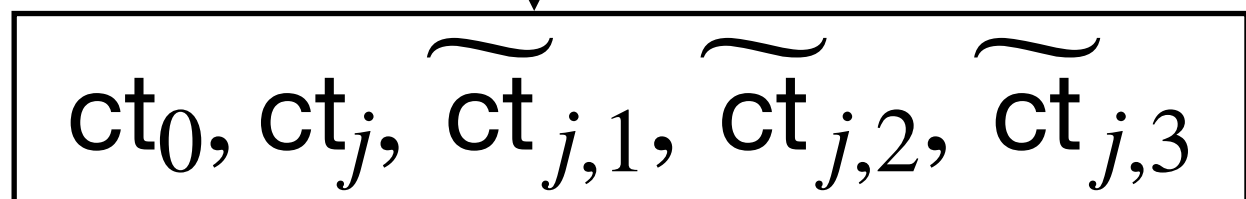
$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



IPFE, IPFE



Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,1} = ( -\mathbf{r}_{M_1}[1], \beta_1 )$$

$$\mathbf{u}_{0,2} = ( \mathbf{r}_{\mathbf{x}}[0], 1 )$$

$$\mathbf{v}_{q,1} = ( -\mathbf{r}_{M_1}[q], (\mathbf{M}_{x_j,1} \mathbf{r}_{M_1})[q] )$$

$$\mathbf{u}_{j,2} = ( \mathbf{r}_{\mathbf{x}}[j-1], \mathbf{r}_{\mathbf{x}}[j] )$$

$$\tilde{\mathbf{v}}_{q,1} = ( -\mathbf{r}_{M_1}[q], \mathbf{e}_{q_{\text{acc},1}}[q] )$$

$$\tilde{\mathbf{u}}_{j,2} = ( \mathbf{r}_{\mathbf{x}}[N], z_2 )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

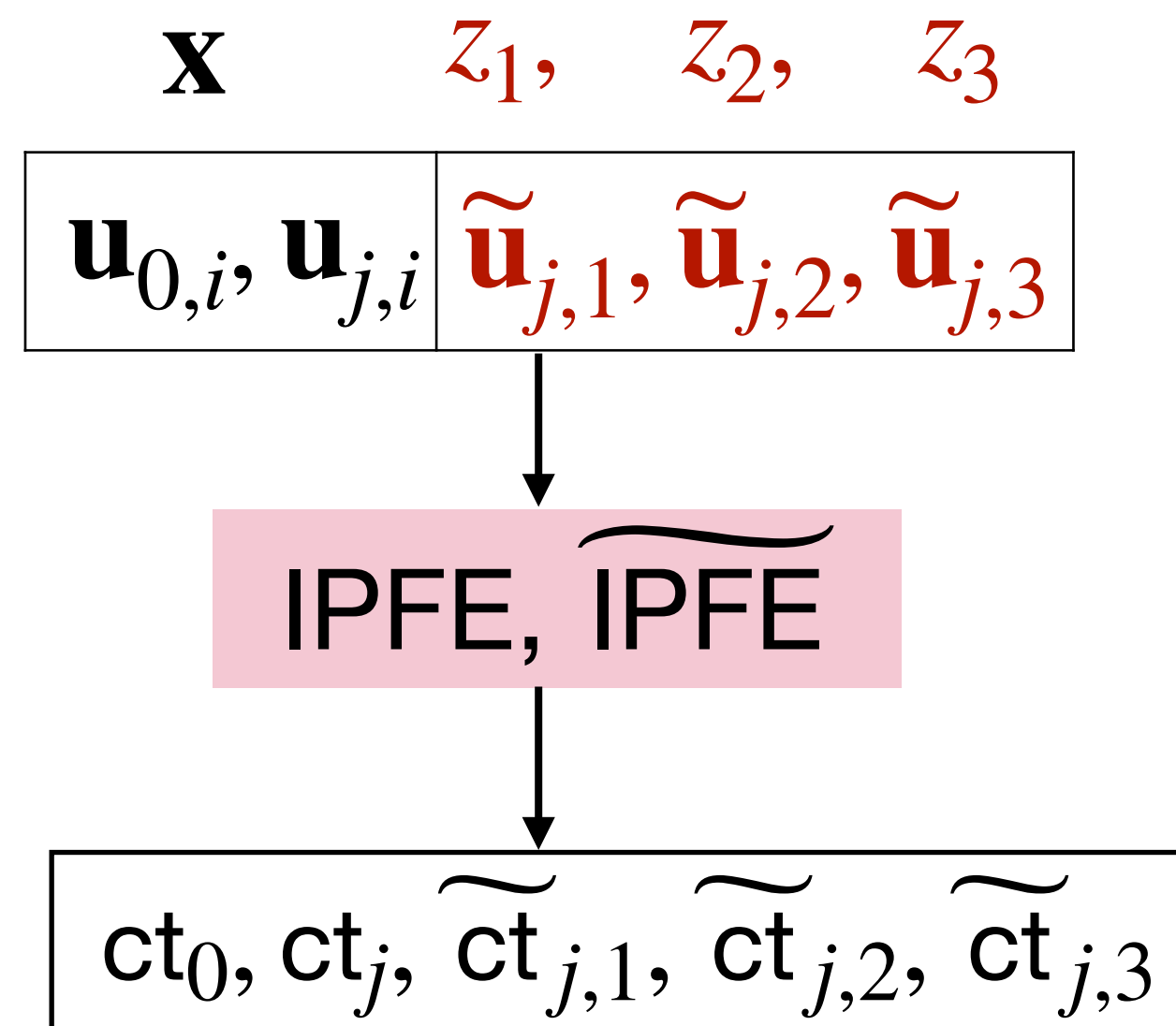
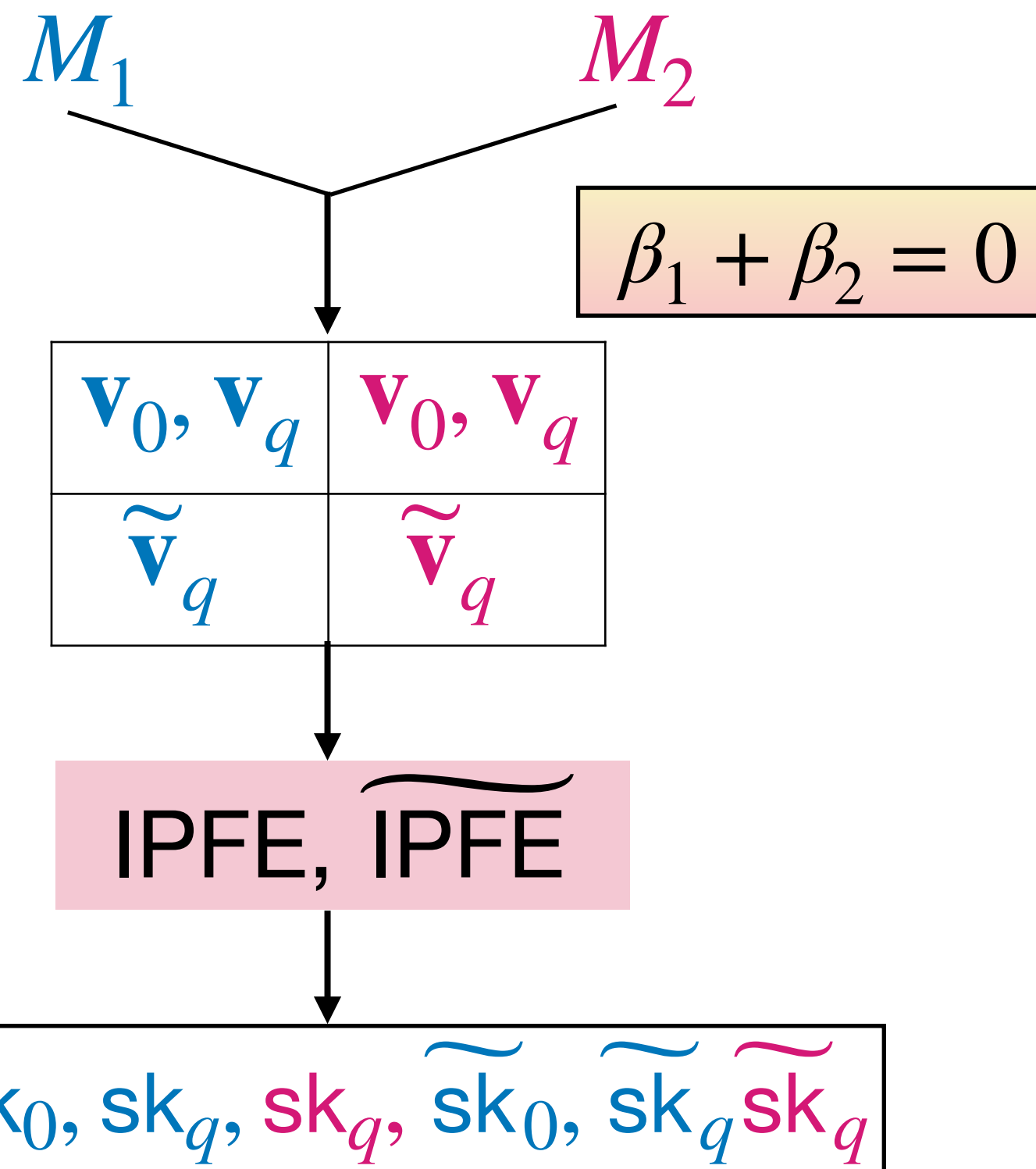
input

output

$$M = (M_1, M_2)$$

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,2} = ( -\mathbf{r}_{M_2}[1], \beta_2 )$$

$$\mathbf{u}_{0,1} = ( \mathbf{r}_{\mathbf{x}}[0], 1 )$$

$$\mathbf{v}_{q,2} = ( -\mathbf{r}_{M_2}[q], (\mathbf{M}_{x,2} \mathbf{r}_{M_2})[q] )$$

$$\mathbf{u}_{j,1} = ( \mathbf{r}_{\mathbf{x}}[j-1], \mathbf{r}_{\mathbf{x}}[j] )$$

$$\tilde{\mathbf{v}}_{q,2} = ( -\mathbf{r}_{M_2}[q], \mathbf{e}_{q_{\text{acc},2}[q]} )$$

$$\tilde{\mathbf{u}}_{j,1} = ( \mathbf{r}_{\mathbf{x}}[N], z_1 )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

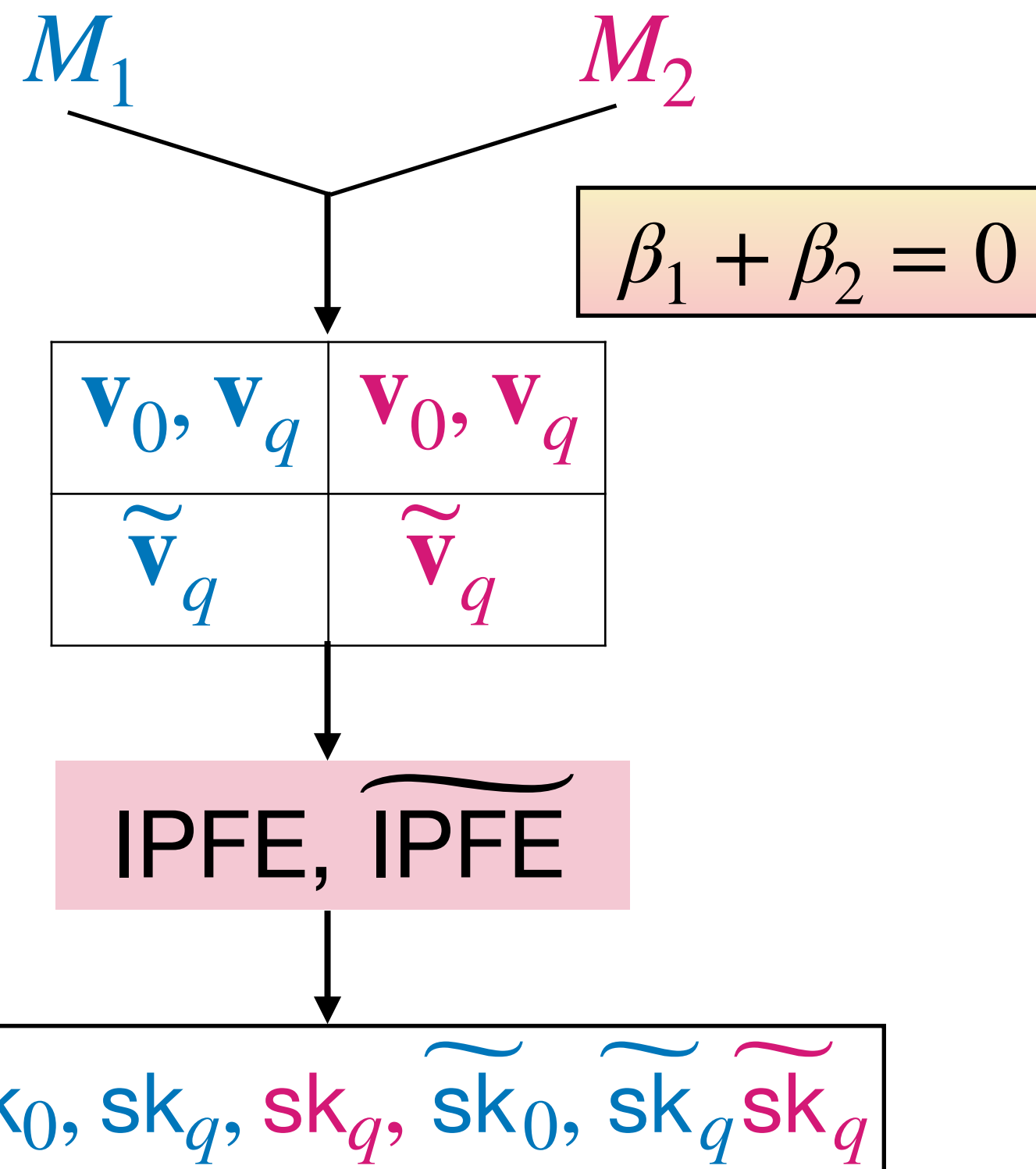
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_2(\mathbf{x}) + z_2 M_1(\mathbf{x})$$



$$\mathbf{x} \quad z_1, z_2, z_3$$

$$\mathbf{u}_{0,i}, \mathbf{u}_{j,i} \quad \tilde{\mathbf{u}}_{j,1}, \tilde{\mathbf{u}}_{j,2}, \tilde{\mathbf{u}}_{j,3}$$

IPFE, IPFE

$$ct_0, ct_j, \tilde{ct}_{j,1}, \tilde{ct}_{j,2}, \tilde{ct}_{j,3}$$

Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,2} = ( -\mathbf{r}_{M_2}[1], \beta_2 )$$

$$\mathbf{u}_{0,1} = ( \mathbf{r}_{\mathbf{x}}[0], 1 )$$

$$\mathbf{v}_{q,2} = ( -\mathbf{r}_{M_2}[q], (\mathbf{M}_{x,2} \mathbf{r}_{M_2})[q] )$$

$$\mathbf{u}_{j,1} = ( \mathbf{r}_{\mathbf{x}}[j-1], \mathbf{r}_{\mathbf{x}}[j] )$$

$$\tilde{\mathbf{v}}_{q,2} = ( -\mathbf{r}_{M_2}[q], \mathbf{e}_{q_{\text{acc},2}}[q] )$$

$$\tilde{\mathbf{u}}_{j,1} = ( \mathbf{r}_{\mathbf{x}}[N], z_1 )$$



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

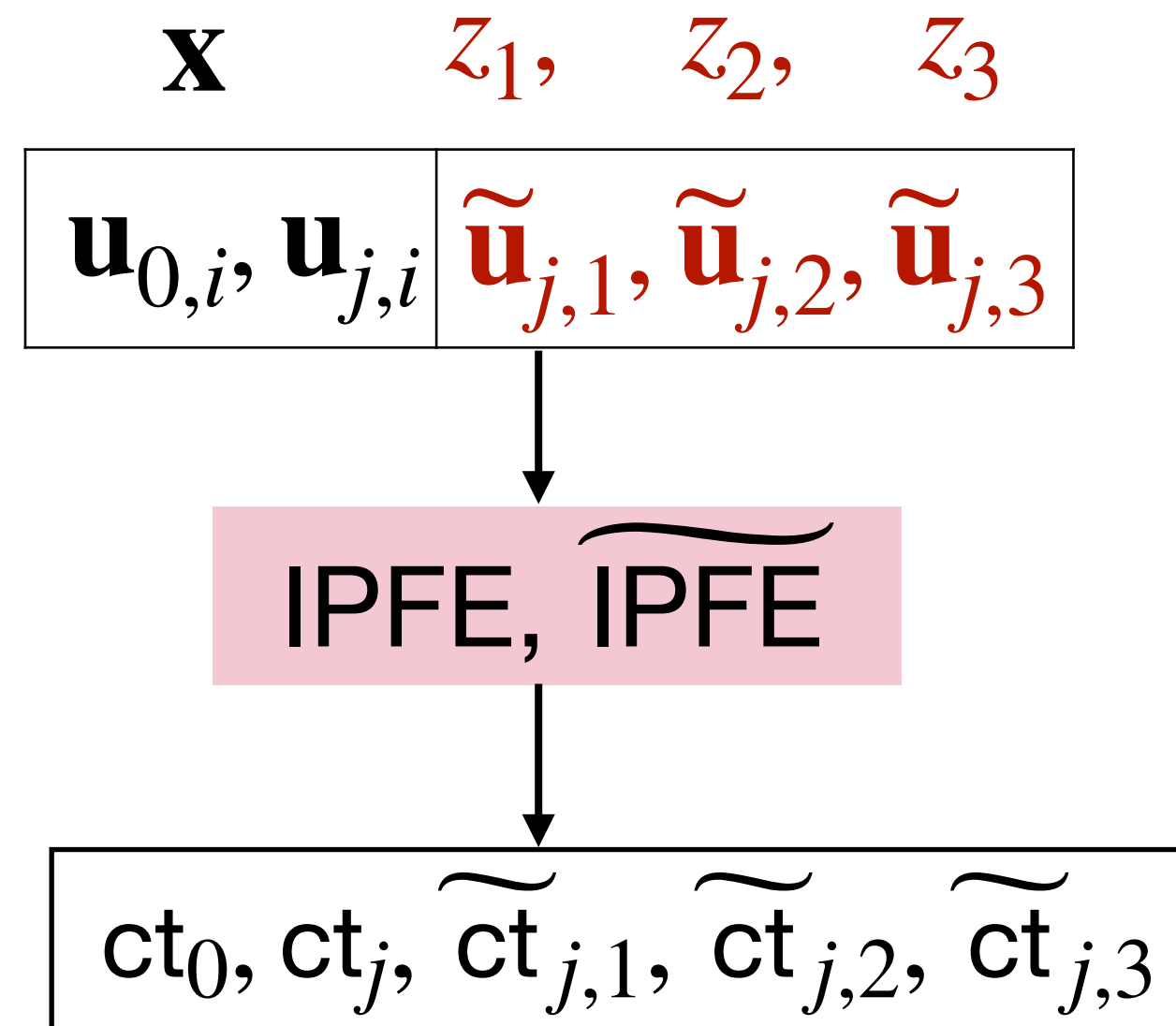
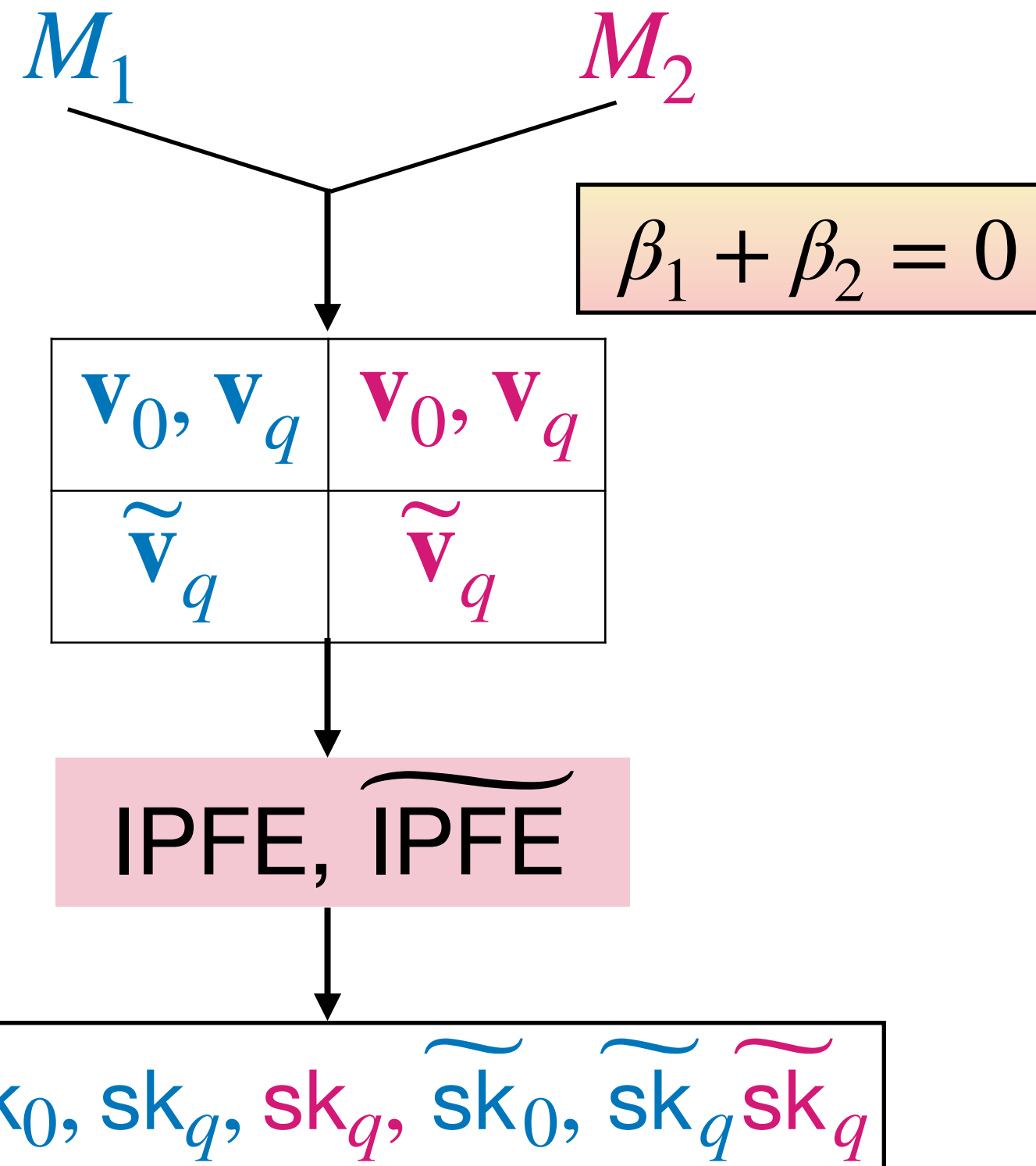
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_2(\mathbf{x}) + z_2 M_1(\mathbf{x})$$



Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

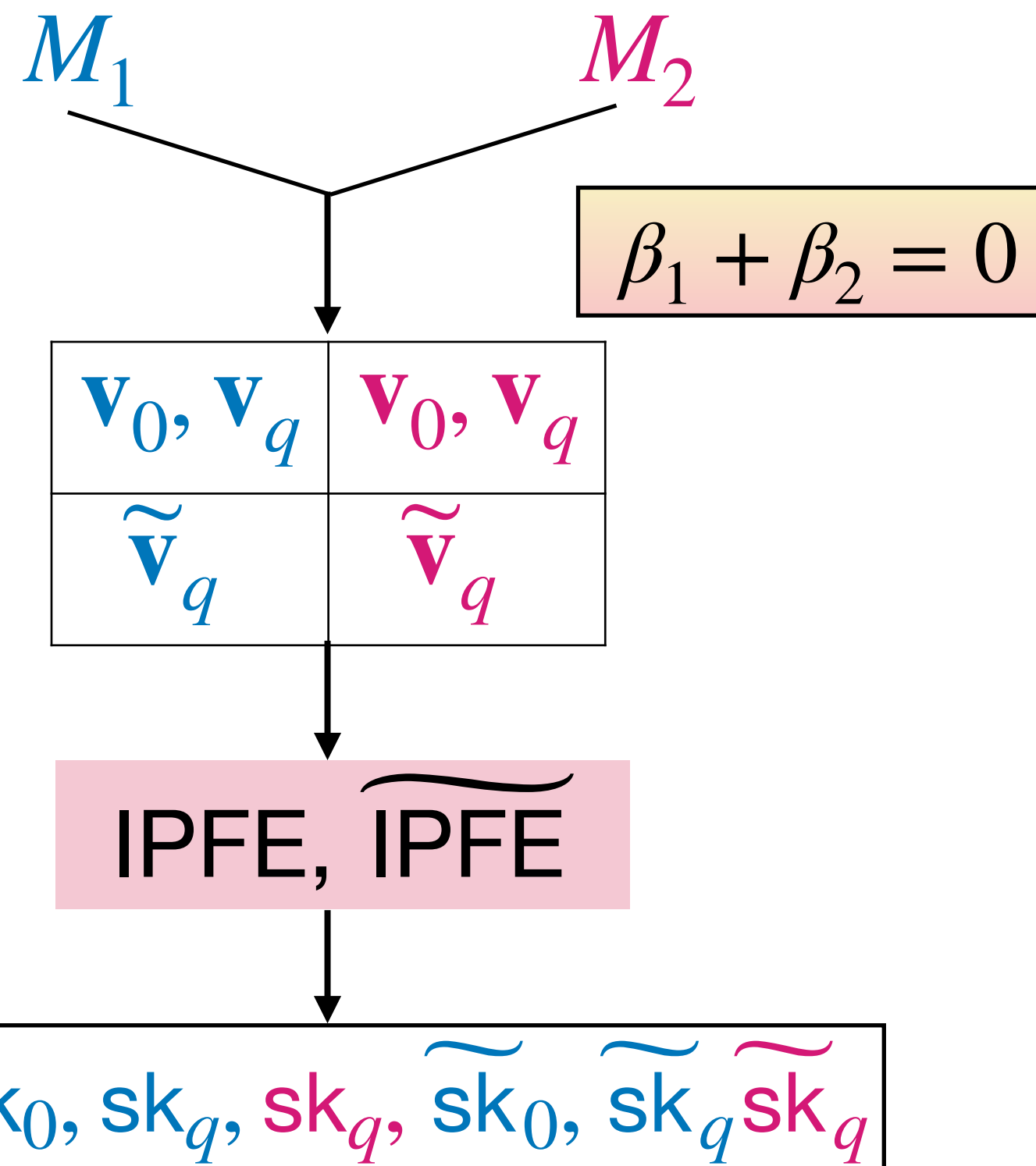
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_2(\mathbf{x}) + z_2 M_1(\mathbf{x})$$



$$\mathbf{x} \quad z_1, z_2, z_3$$

$$\mathbf{u}_{0,i}, \mathbf{u}_{j,i} \quad \tilde{\mathbf{u}}_{j,1}, \tilde{\mathbf{u}}_{j,2}, \tilde{\mathbf{u}}_{j,3}$$

**IPFE, IPFE**

$$ct_0, ct_j, \tilde{ct}_{j,1}, \tilde{ct}_{j,2}, \tilde{ct}_{j,3}$$

Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,1} = (\pi_1(1,1), \dots)$$

$$\mathbf{u}_{0,2} = (\rho_2(-1,2), \dots)$$

$$\mathbf{v}_{q,1} = (\pi_1(1,1), -\mathbf{r}_{M_1}[q], (\mathbf{M}_{x_j,1} \mathbf{r}_{M_1})[q])$$

$$\mathbf{u}_{j,2} = (\rho_2(-1,2), \mathbf{r}_x[j-1], \mathbf{r}_x[j])$$

$$\tilde{\mathbf{v}}_{q,1} = (\pi_1(1,1), \dots)$$

$$\tilde{\mathbf{u}}_{j,2} = (\rho_2(-1,2), \dots)$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

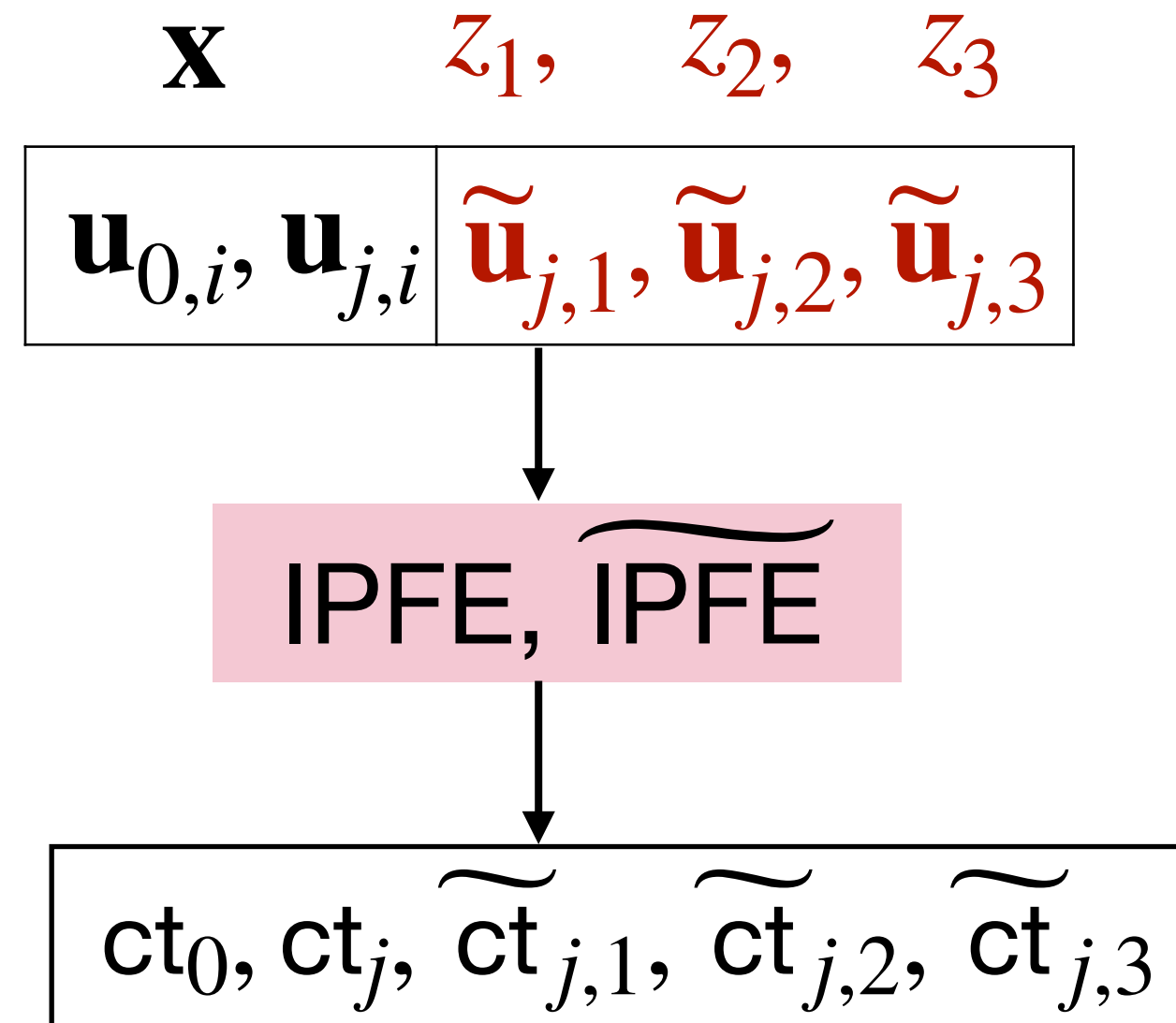
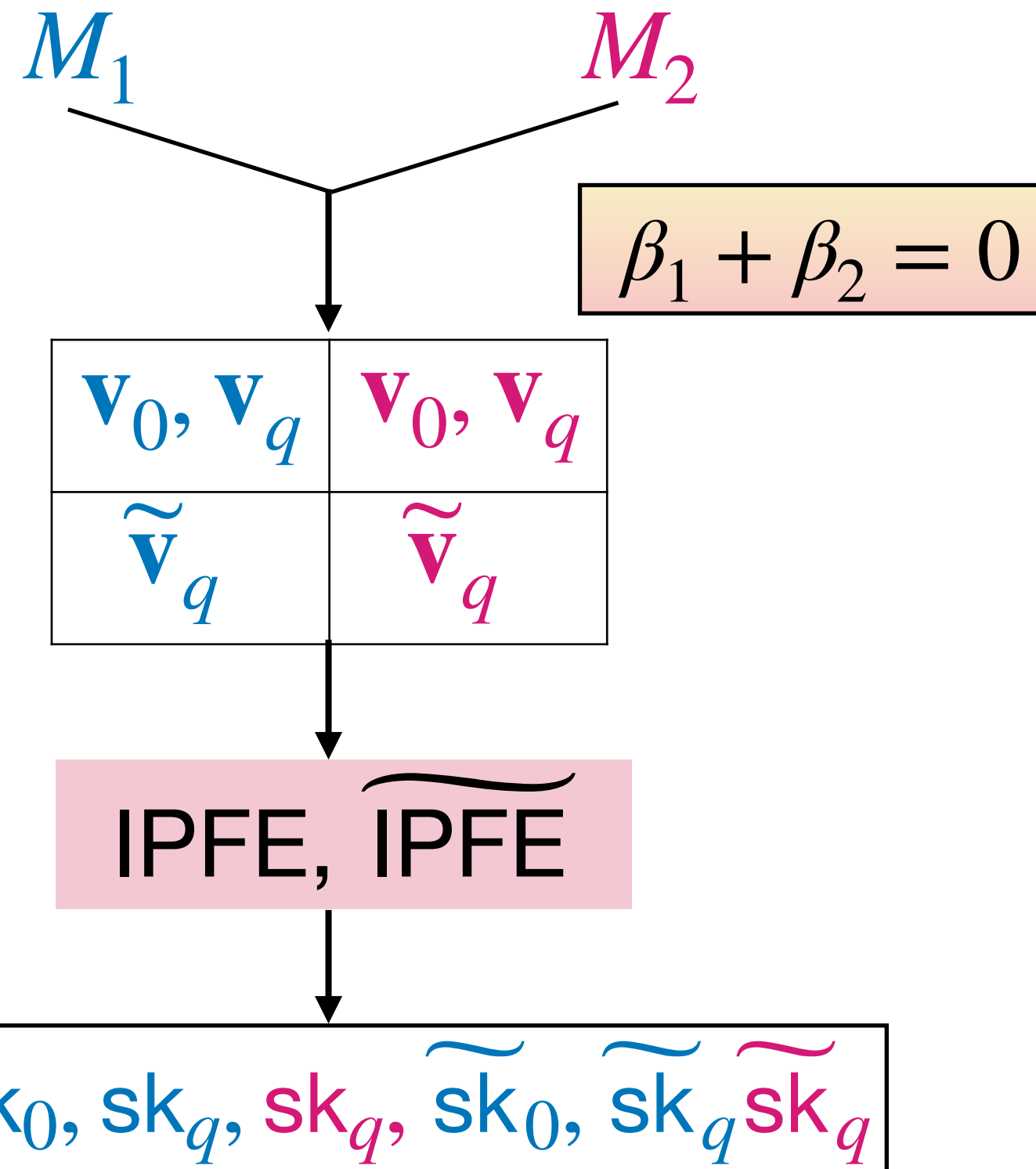
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = (\pi_k(k,1), \dots)$$

$$\mathbf{v}_{q,k} = (\pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q])$$

$$\tilde{\mathbf{v}}_{q,k} = (\pi_k(k,1), \dots)$$

$$\mathbf{u}_{0,i} = (\rho_i(-1,i), \dots)$$

$$\mathbf{u}_{j,i} = (\rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j])$$

$$\tilde{\mathbf{u}}_{j,i} = (\rho_i(-1,i), \dots)$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \quad -\mathbf{r}_{M_2}[1], \quad \beta_k \quad )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \quad \mathbf{r}_{\mathbf{x}}[0], \quad 1 \quad )$$

**Steps: Simulation Security**

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \quad \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), \quad -\mathbf{r}_{M_k}[q], \quad (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] \quad )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \quad \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \quad \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \quad \mathbf{r}_{\mathbf{x}}[j-1], \quad \mathbf{r}_{\mathbf{x}}[j] \quad )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \quad \dots )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \boxed{1}, 0 )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \boxed{\ell_{0,i}}, \boxed{1} )$$

FH-IPFE

**Steps: Simulation Security**

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2, M_4)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \boxed{1}, 0 )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \boxed{\ell_{0,i}}, \boxed{1} )$$

$$\boxed{\rho_i(-1,i) \cdot \pi_4(4,1) \neq 0}$$

## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\boxed{\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)}$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\boxed{\ell_{j,i} \leftarrow \$ \text{ for all } j > 1}$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2, M_4)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \boxed{1}, 0 )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \boxed{\ell_{0,i}}, \boxed{1} )$$

$$\rho_i(-1,i) \cdot \pi_4(4,1) \neq 0$$

additional entropy from *index encoding*

## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots ) \quad \mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] ) \quad \tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots ) \quad \mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] ) \quad \tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2, M_4)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \boxed{1}, 0 )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \boxed{\ell_{0,i}}, \boxed{1} )$$

$$\boxed{\rho_i(-1,i) \cdot \pi_4(4,1) \neq 0} \Rightarrow \boxed{\beta_4 \leftarrow \$} \Rightarrow \boxed{\beta_1 + \beta_2 + \beta_4 \neq 0}$$

additional entropy from *index encoding*

along with **FH-IPFE**

## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\boxed{\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)}$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\boxed{\ell_{j,i} \leftarrow \$ \text{ for all } j > 1}$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots )$$



# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2, M_4)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \boxed{1}, 0 )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \boxed{\ell_{0,i}}, \boxed{1} )$$

$$\rho_i(-1,i) \cdot \pi_4(4,1) \neq 0$$

$\Rightarrow$

$$\beta_4 \leftarrow \$$$

$\Rightarrow$

$$\beta_1 + \beta_2 + \beta_4 \neq 0$$

additional entropy from *index encoding*

along with **FH-IPFE**

## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

$$\tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots )$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\begin{array}{l} \mathbf{v}_{q,k} = ( \pi_k(k,1), \quad -\mathbf{r}_{M_k}[q], \quad (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] ) \\ \mathbf{u}_{j,i} = ( \rho_i(-1,i), \quad \mathbf{r}_x[j-1], \quad \mathbf{r}_x[j] ) \end{array} \xrightarrow{i=k} \ell_{j,i}$$

## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\begin{array}{lll} \mathbf{v}_{0,k} = ( \pi_k(k,1), \quad \dots ) & \mathbf{v}_{q,k} = ( \pi_k(k,1), \quad -\mathbf{r}_{M_k}[q], \quad (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] ) & \tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \quad \dots ) \\ \mathbf{u}_{0,i} = ( \rho_i(-1,i), \quad \dots ) & \mathbf{u}_{j,i} = ( \rho_i(-1,i), \quad \mathbf{r}_x[j-1], \quad \mathbf{r}_x[j] ) & \tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \quad \dots ) \end{array}$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

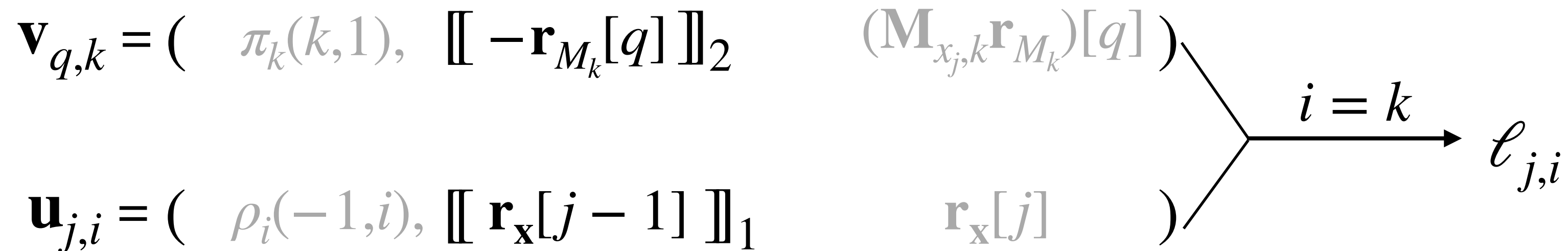
$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

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$$\begin{array}{lll} \mathbf{v}_{0,k} = ( \pi_k(k,1), \dots ) & \mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j, k} \mathbf{r}_{M_k})[q] ) & \tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots ) \\ \mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots ) & \mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] ) & \tilde{\mathbf{u}}_{j,i} = ( \rho_i(-1,i), \dots ) \end{array}$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), \llbracket 1 \rrbracket_2, (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

**FH-IPFE**

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \llbracket -\mathbf{r}_x[j-1] \mathbf{r}_{M_k}[q] \rrbracket_1, \mathbf{r}_x[j] )$$

$i = k \rightarrow \ell_{j,i}$

## Steps: Simulation Security

1. use **RevSamp** to hardwire  $\mu_i = z_i M_i(\mathbf{x}) + \beta_i$  in the *first* label

$$\ell_{0,i} \leftarrow \text{RevSamp}(M_i, \mathbf{x}, \mu_i, \dots)$$

2. use **Marginal Randomness** to randomize all *other* labels

$$\ell_{j,i} \leftarrow \$ \text{ for all } j > 1$$

$$\mathbf{v}_{0,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{v}_{q,k} = ( \pi_k(k,1), -\mathbf{r}_{M_k}[q], (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] )$$

$$\tilde{\mathbf{v}}_{q,k} = ( \pi_k(k,1), \dots )$$

$$\mathbf{u}_{0,i} = ( \rho_i(-1,i), \dots )$$

$$\mathbf{u}_{j,i} = ( \rho_i(-1,i), \mathbf{r}_x[j-1], \mathbf{r}_x[j] )$$

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# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

$$M = (M_1, M_2)$$

input

$$(\mathbf{x}, \mathbf{z} = (z_1, z_2, z_3))$$

output

$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$

$$\mathbf{v}_{q,k} = \left( \pi_k(k,1), \begin{bmatrix} 1 \\ \vdots \end{bmatrix}_2, (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] \right)$$

FH-IPFE

$$\mathbf{u}_{j,i} = \left( \rho_i(-1,i), \begin{bmatrix} -\mathbf{r}_{j-1}[q] \\ \vdots \end{bmatrix}_1, \mathbf{r}_x[j] \right)$$

$i = k \rightarrow \ell_{j,i}$

$$\mathbf{r}_{j-1} = \mathbf{r}_x[j-1] \cdot \mathbf{r}_{M_k}$$

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$$\begin{aligned} \mathbf{v}_{q,k} &= ( \pi_k(k,1), \llbracket 1 \rrbracket_2, (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] ) \\ \mathbf{u}_{j,i} &= ( \rho_i(-1,i), \llbracket -\mathbf{s}_{j-1}[q] \leftarrow \$ \rrbracket_1, \mathbf{r}_x[j] ) \end{aligned} \xrightarrow{i=k} \ell_{j,i}$$

$$\mathbf{r}_{j-1} = \mathbf{r}_x[j-1] \cdot \mathbf{r}_{M_k} \xrightarrow{\text{DDH in } \mathbb{G}_1} \mathbf{s}_{j-1} \leftarrow \$$$

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$\mathbf{r}_{M_k}$  may appear in *other* places of  $\mathbf{v}_{q,k}$  for any  $k$

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$$\mathbf{v}_{q,k} = \left( \pi_k(k,1), \llbracket 1 \rrbracket_2, (\mathbf{M}_{x_j,k} \mathbf{r}_{M_k})[q] \right)$$

$$\mathbf{u}_{j,i} = \left( \rho_i(-1,i), \llbracket -\mathbf{s}_{j-1}[q] \leftarrow \$ \rrbracket_1, \mathbf{r}_x[j] \right)$$

$i = k \rightarrow \ell_{j,i}$

$$\mathbf{r}_{j-1} = \mathbf{r}_x[j-1] \cdot \mathbf{r}_{M_k} \xrightarrow{\text{DDH in } \mathbb{G}_1} \mathbf{s}_{j-1} \leftarrow \$ \quad \times$$

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use distributed randomness mechanism

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# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

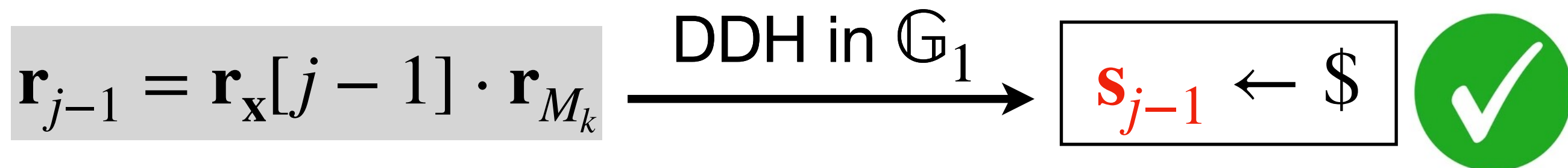
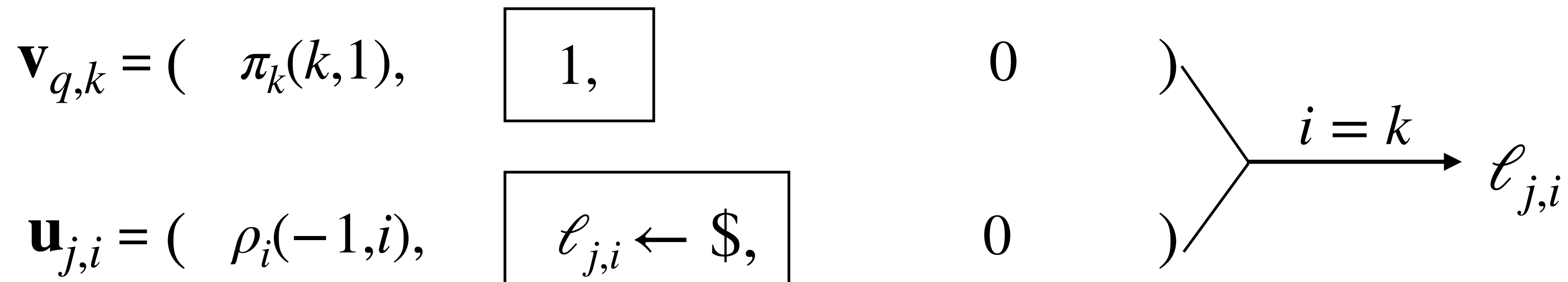
$$M = (M_1, M_2)$$

input

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$$z_1 M_1(\mathbf{x}) + z_2 M_2(\mathbf{x})$$



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 \end{array}$$

# Our Idea for FE-UAWS for TMs (DFA for concreteness)

function

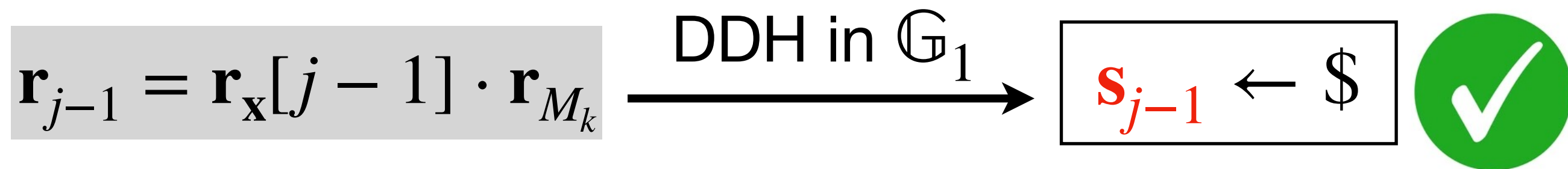
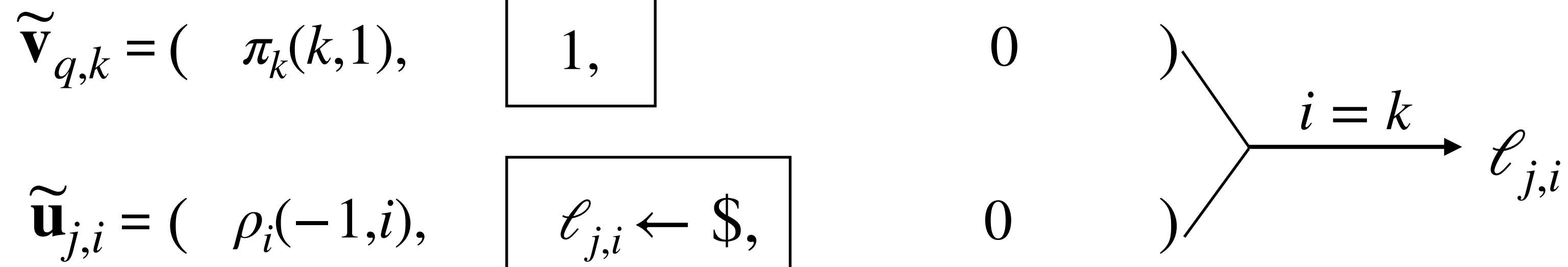
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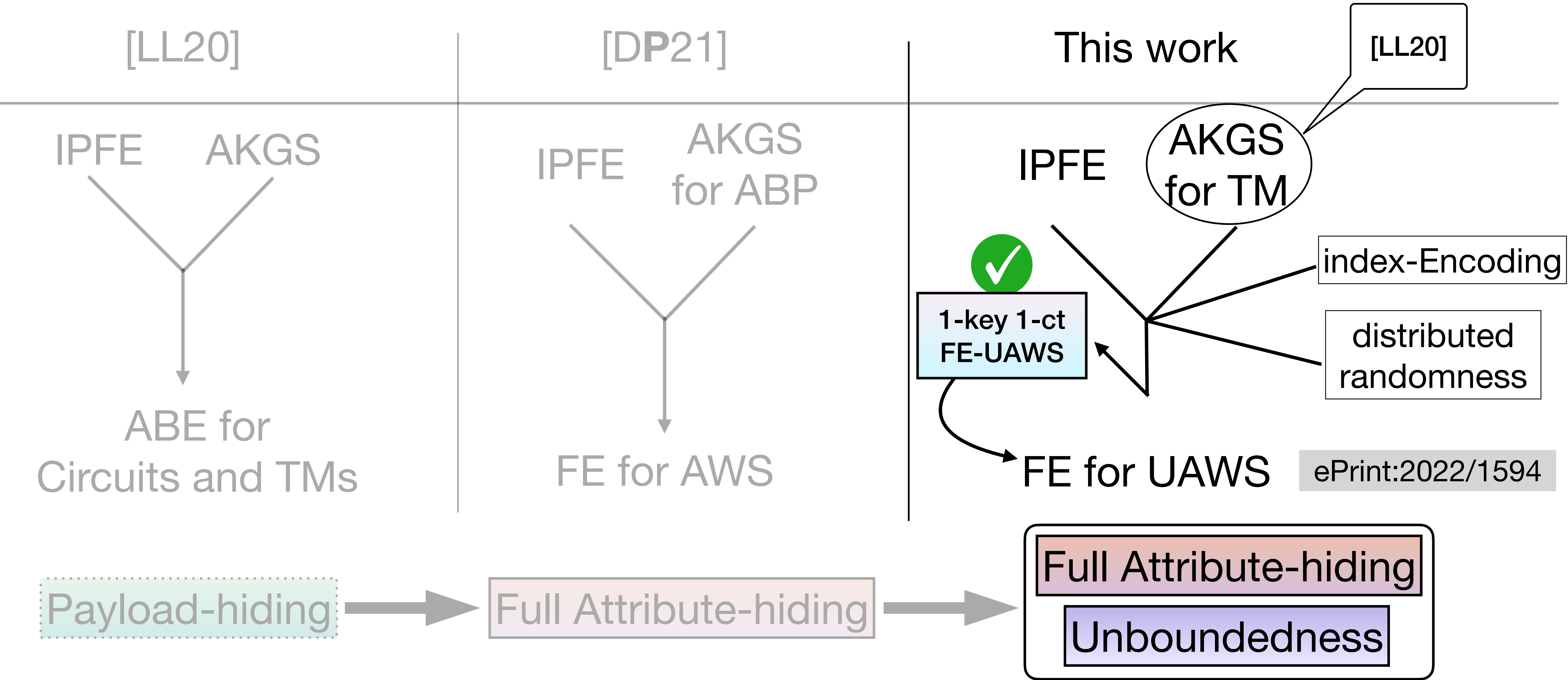
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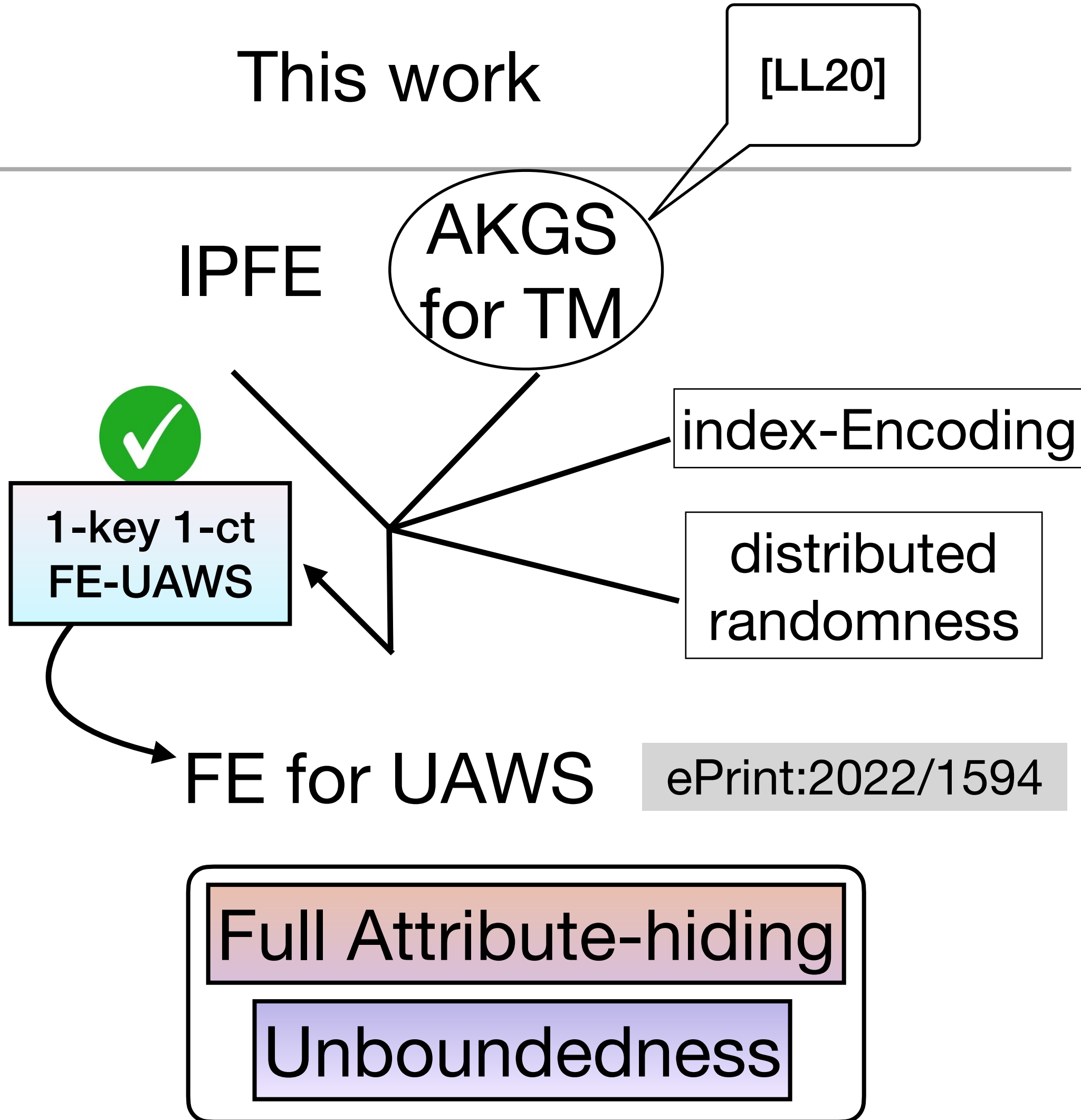
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# Roadmap towards FE for UAWS



# Conclusion

- Definition and Construction of FE for UAWS
  - ◆ **input-specific**  $|ct| = O(|\mathbf{x}|, |\mathbf{z}|)$
  - ◆ **Compact** ciphertext
  - ◆ Fully collusion-resistant **AD-SIM**
  - ◆ Standard assumption: **SXDH**



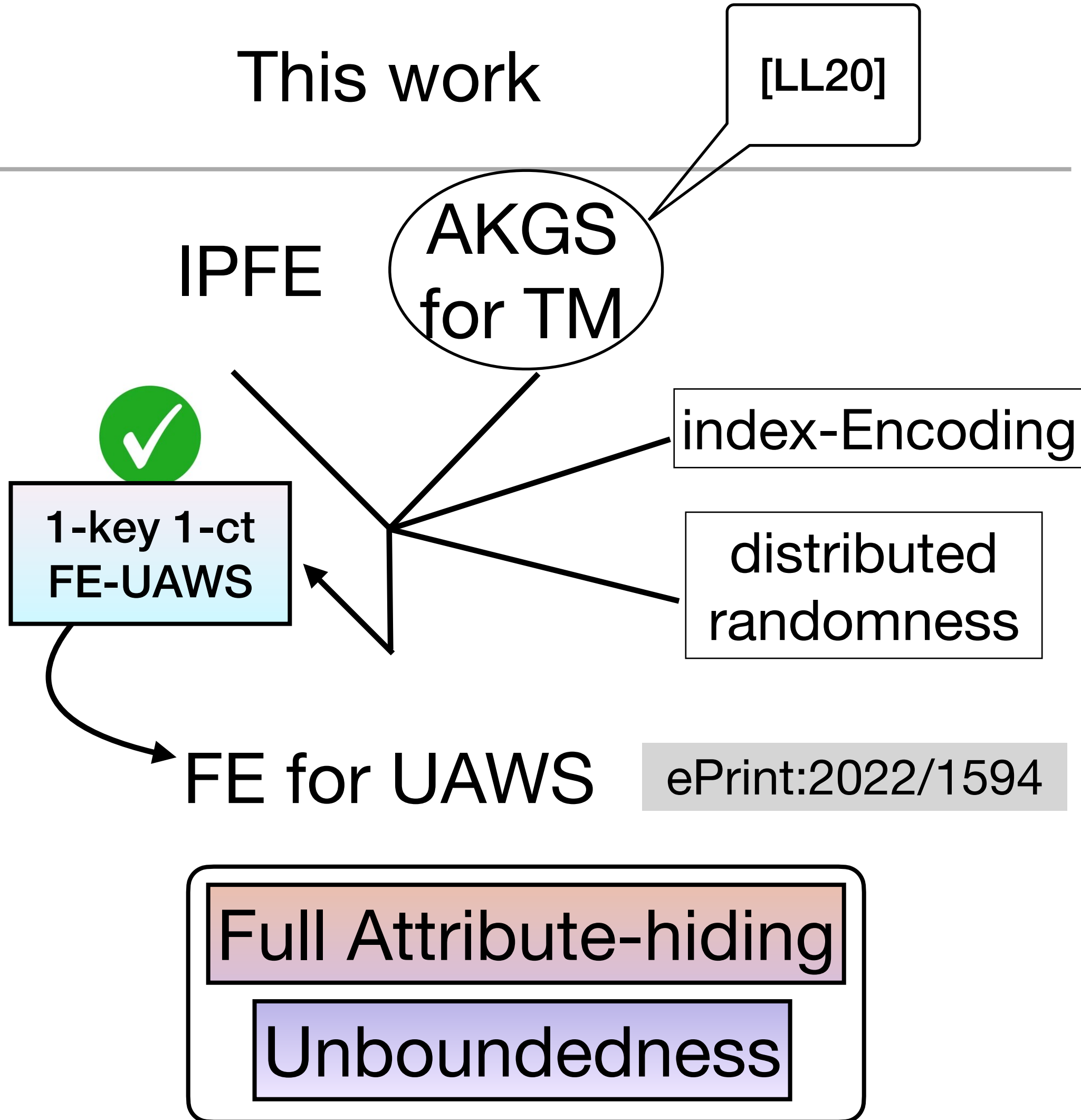
# Conclusion

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- Future research directions of FE for UAWS

- ◆ succinctness:  $|ct| = O(|\mathbf{z}|)$



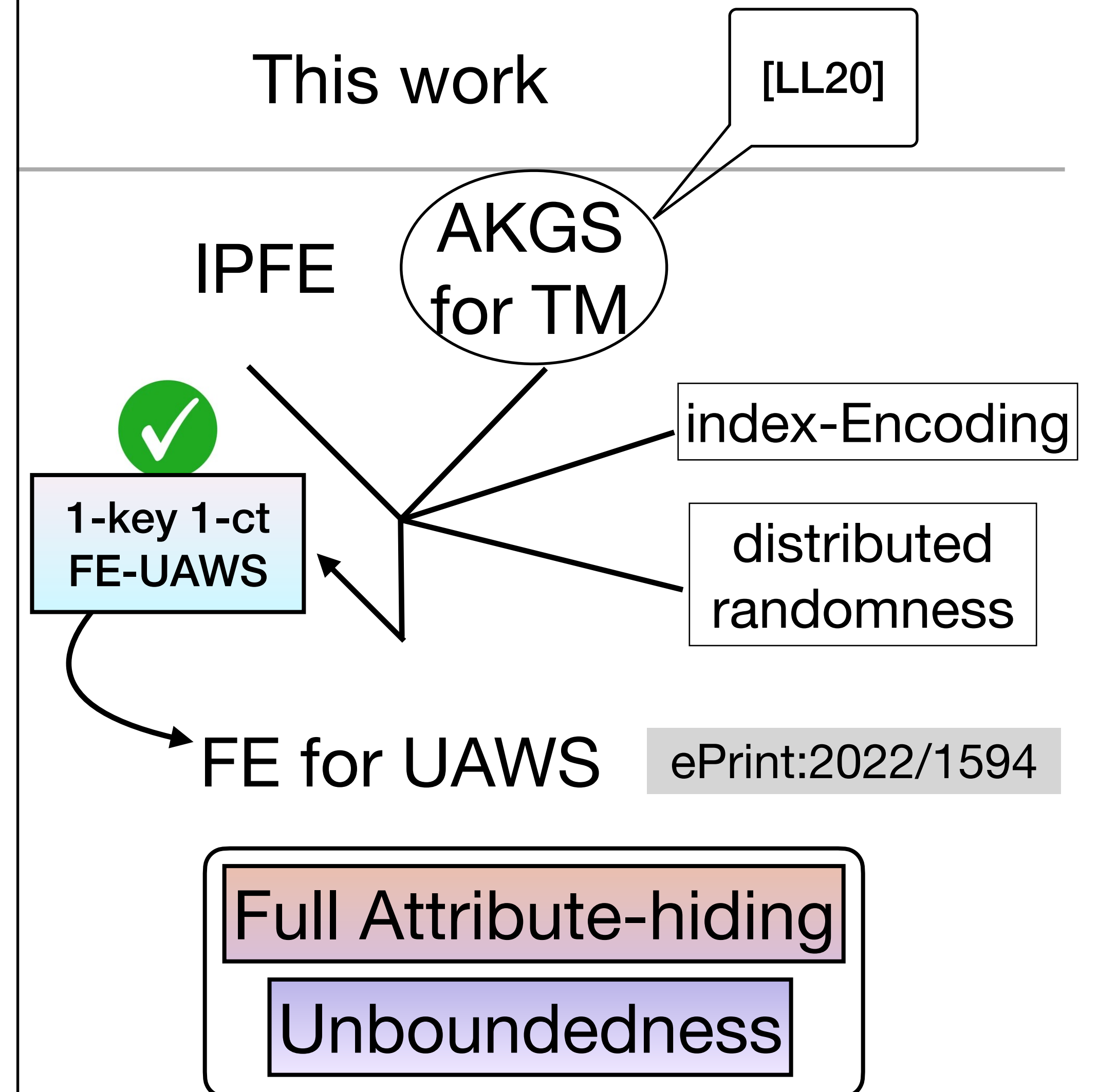
# Conclusion

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# Conclusion

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*Thank You!*

