



Anonymity of NIST PQC Round 3 KEMs

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Backgrounds

Quantum Computers:

- Google: 54 qubits (2019)
- IBM Q53: 53 qubits (2019)
- USTC: 76 qubits (2020)

[Sho94] QC solves factoring/DL easily

[GE21] ERR 10^{-3} , 1ms/1cycle

20Mqubits QC solves RSA2048 in 8h.

→ Consider quantum !

cf. https://en.wikipedia.org/wiki/List_of_quantum_processors

[Sho94] Shor (FOCS 1994)

[GE21] Gidney, Ekerå (Quantum 2021)

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Countermeasures:

1. Looooooooong RSA/DL
2. Post-Quantum Cryptography (PQC)
3. Quantum Cryptography (e.g., QKD)

4 Finalists: KEM

Classic McEliece (code)
CRYSTALS-Kyber (lattice)
NTRU (lattice)
SABER (lattice)

5 Alternates: KEM

BIKE (code)
FrodoKEM (lattice)
HQC (code)
NTRU Prime (lattice)
SIKE (isogeny)

3 Finalists: Sig

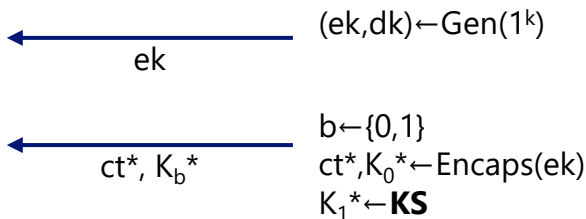
CRYSTALS-Dilithium (lattice)
Falcon (lattice)
Rainbow (MQ)

3 Alternates: Sig

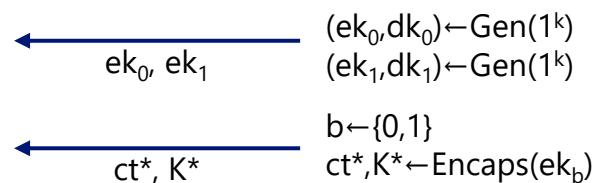
GeMSS (MQ)
Picnic (ZK)
SPHICS+ (Hash)

Anonymity of KEM

Indistinguishability (IND):



Anonymity (ANO):



Applications:

- Anonymous credential
- Auction
- Anonymous AKE, etc

Review of [GMP22]: ANON in QROM



NAME	KEM			PKE
	IND	ANO	CF	ANO
Classic McEliece	Y	?	N	?
Kyber	?	?	?	?
NTRU	Y	?	?	?
Saber	?	?	?	?
FrodoKEM (+Kyber'&Saber')	Y	Y	Y	Y

SPR=Strong Pseudorandomness, ANO=Anonymity
CF=Collision Freeness, ROB=Robustness

Show anonymity of FO^{\perp}

Problems:

1. Hard to simul. two dec. oracles
 - CME: not collision-free
 - NTRU: SXY uses $H(m)$
2. Need to avoid tweaks
 - Kyber/Saber: 'pre-key' and 'nest'

Our result

Our Results – 1



1. Strong Pseudorandom (SPR) → Anonymous
2. KEM/DEM framework for SPR HPKE
 1. Smooth SPR-CCA KEM $^{\pm}$ and SPR-otCCA DEM
 2. Sparse SPR-CCA KEM $^{\perp}$ and INT-CTXT, SPR-otCCA DEM
3. Strongly Disjoint-Simulatable PKE
–(SXY etc.)→ Smooth/Sparse SPR-CCA KEM
4. Apply them to NIST PQC Round 3 KEM

[CS03] Cramer, Shoup (SIAM J. Comput. 33(1), 2003)

[Moh10] Mohassel (ASIACRYPT 2010)

[GRP22] Grubbs, Maram, Paterson (EUROCRYPT 2022)

Our Results – 2



NAME	KEM			PKE
	IND	ANO	CF	ANO
Classic McEliece	Y	<u>Y</u>	N	<u>Y</u>
Kyber	?	?	?	?
NTRU	Y	<u>Y</u>	Y	<u>Y</u>
Saber	?	?	?	?
FrodoKEM (+Kyber'&Saber')	Y	Y	Y	Y

SPR=Strong Pseudorandomness, ANO=Anonymity
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NAME	KEM			PKE
	IND	ANO	CF	ANO
BIKE	Y	<u>Y</u>	<u>Y</u>	<u>Y</u>
HQC-128/196	Y	<u>Y</u>	<u>Y</u>	<u>Y</u>
HQC-256	Y	N	<u>Y</u>	N
sntrupr of NTRU Prime	<u>?</u>	<u>?</u>	<u>?</u>	<u>?</u>
ntrulpr of NTRU Prime	Y	<u>Y</u>	<u>Y</u>	<u>Y</u>
SIKE	Y	<u>Y</u>	<u>Y</u>	<u>Y</u>

SXY: SDS \rightsquigarrow SPR-CCA \rightarrow ANO-CCA

Deterministic PKE

Gen(1^k) \rightarrow (ek,dk)

Enc(ek,m) \rightarrow ct

Dec(dk,ct) \rightarrow m/ \perp

PKE is Strongly Disjoint-Siml.

1. **Enc**(ek,U(M)) \approx_c Sim(1^k)
2. $\Pr[c \leftarrow \text{Sim}(1^k): c \in \text{Enc}(ek,M)]$ is negl.



KEM = SXY[PKE,H,H']

Gen(1^k)

(ek,dk) \leftarrow **Gen**(1^k), seed \leftarrow $\{0,1\}^k$
return ek and (dk,ek,seed)

Encaps(ek;m)

ct \leftarrow **Enc**(ek,m), K \leftarrow H(m)
return ct and K

Decaps((dk,ek,seed),ct)

$m' \leftarrow$ **Dec**(dk,ct)

If ct = **Enc**(ek,m') then K \leftarrow H(m')

Else K \leftarrow H'(seed,ct)

return K

PKE is SDS

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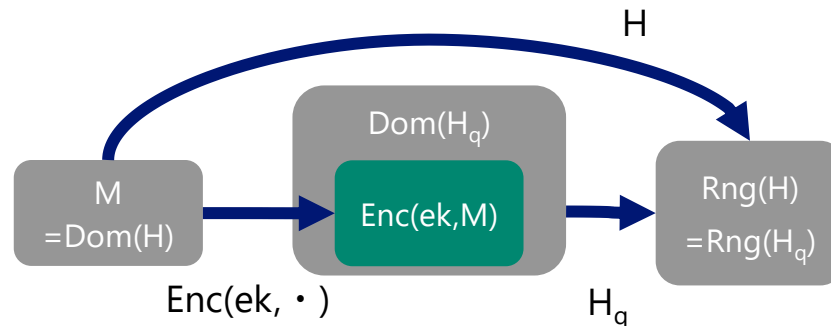
Encaps($ek; m$):
 $ct \leftarrow \text{Enc}(ek, m)$
 $K \leftarrow H(m)$

Decaps(dk, ct):
 $m' \leftarrow \text{Dec}(dk, ct)$
If $ct = \text{Enc}(ek, m')$
then $K \leftarrow H(m')$
else $K \leftarrow H'(\text{seed}, ct)$

💡 Simulate Decaps w/ H

Let $H(m) := H_q(\text{Enc}(ek, m))$

- **Decaps**(ct) = $H_q(ct)$
- $\text{Dom}(H_q) \setminus \text{Enc}(ek, M)$ is inaccessible
 $\rightarrow H_q(\text{Sim}(1^k))$ looks random



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SDS \rightsquigarrow IND-CCA

$\{(ek, ct, K): ct \leftarrow \mathbf{Enc}(ek, m), K \leftarrow H(m)\}$ w/ dk
 $\approx_s \{(ek, ct, K): ct \leftarrow \mathbf{Enc}(ek, m), K \leftarrow H_q(ct)\}$ w/ H_q
 $\approx_c \{(ek, ct, K): ct \leftarrow \text{Sim}, K \leftarrow H_q(ct)\}$ w/ H_q
 $\approx_s \{(ek, ct, K): ct \leftarrow \text{Sim}, K \leftarrow U\}$ w/ H_q
 $\dots \approx_c \{(ek, ct, K): ct \leftarrow \mathbf{Enc}(ek, m), K \leftarrow U\}$ w/ dk

Attempt: SDS \rightsquigarrow ANON-CCA



PKE is SDS

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If $ct = \mathbf{Enc}(ek, m')$
then $K \leftarrow H(m')$
else $K \leftarrow H'(\text{seed}, ct)$

Need to simul. two Decaps

Let $H(m) := H_q(\mathbf{Enc}(ek, m))$?

[GRM21] Use $H(m, c)$

$H(m, c) := H_i(c)$ if $\mathbf{Enc}(ek_i, m) = c$ or $H_q(c)$

$\mathbf{Decaps0}(ct) = H_0(ct)$

$\mathbf{Decaps1}(ct) = H_1(ct)$

Problem:

- Require collision-freeness
- Need to modify H

KEM is Strongly PR.

$$\{(ek, ct, K): (ct, K) \leftarrow \text{Encaps}(ek)\} \\ \approx_c \{(ek, ct, K): ct \leftarrow \text{Sim}(1^k), K \leftarrow U\}$$

KEM is Anonymous

$$\{(ek_0, ek_1, ct_0, K_0): (ct_0, K_0) \leftarrow \text{Encaps}(ek_0)\} \\ \approx_c \{(ek_0, ek_1, ct_1, K_1): (ct_1, K_1) \leftarrow \text{Encaps}(ek_1)\}$$

SPR-CCA → ANO-CCA

$$\{(ek_0, ek_1, ct_0, K_0): (ct_0, K_0) \leftarrow \text{Encaps}(ek_0)\} \\ \approx_c \{(ek_0, ek_1, ct, K): ct \leftarrow \text{Sim}(1^k), K \leftarrow U\}$$

☞ reduced to SPR-CCA on ek_0

$$\approx_s \{(ek_0, ek_1, ct, K): ct \leftarrow \text{Sim}(1^k), K \leftarrow U\}$$

$$\approx_c \{(ek_0, ek_1, ct_1, K_1): (ct_1, K_1) \leftarrow \text{Encaps}(ek_1)\}$$

☞ reduced to SPR-CCA on ek_1

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Summary

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sntrupr of NTRU Prime	<u>?</u>	<u>?</u>	<u>?</u>	<u>?</u>
ntrulpr of NTRU Prime	Y	<u>Y</u>	<u>Y</u>	<u>Y</u>
SIKE	Y	<u>Y</u>	<u>Y</u>	<u>Y</u>

Open Problems:

1. Show SPR/ANO of Kyber/Saber/sntrupr of NTRU Prime w/o quantum indiff.
2. Show SPR/ANO of FO[±] *tightly* as IND [BHHP19]+[KSSSS20]