Round-Optimal Black-Box Protocol Compilers

Yuval Ishai (Technion) Dakshita Khurana (UIUC) Amit Sahai (UCLA) **Akshayaram Srinivasan (TIFR)**

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Common Input: Function f







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Function f



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Semi-Honest Adversary



• Follow the protocol specification but try to learn additional information.



Semi-Honest Adversary



- Follow the protocol specification but try to learn additional information.
- Easier to protect against.



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Malicious Adversary

Types of Adversary

- Follow the protocol specification but try to learn additional information.
- Easier to protect against.



Semi-Honest Adversary

• Could deviate from the protocol specification and break security.



Malicious Adversary



Types of Adversary

- Follow the protocol specification but try to learn additional information.
- Easier to protect against.



Semi-Honest Adversary

- Could deviate from the protocol specification and break security.
- Harder to protect against.



Malicious Adversary



Protocol for computing f' with security against Semi-Honest Adversary

Protocol Compiler



Protocol for computing *f* with security against Malicious adversary



Protocol for computing f'with security against Semi-Honest Adversary





Protocol for computing fwith security against Malicious adversary



Protocol for computing f'with security against Semi-Honest Adversary





Goal: Construct "efficient" Protocol Compilers.



Protocol for computing fwith security against Malicious adversary



• Preserve round complexity

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- Black-Box use

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• The compiled protocol makes black-box use of the initial protocol.

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- Making use of simple tools

• Preserve round complexity

• Black-Box use

• The compiled protocol makes black-box use of the initial protocol.

• Making use of simple tools

• The compiler avoids sophisticated and less efficient cryptographic tools.

Prior Approaches GMW Compiler





Semi-Honest Secure Protocol

Prior Approaches **GMW Compiler**









Protocol



Prior Approaches **GMW Compiler**

Non-Black-Box use



• A black-box protocol compiler.

- A black-box protocol compiler.
- Resulted in an increase in round-complexity.

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- A black-box protocol compiler.
- Resulted in an increase in round-complexity.
- namely, adaptive security with erasures.

• Required stronger cryptographic tools such as malicious secure oblivious transfer. • Required the semi-honest secure protocol to satisfy stronger security property,



Our Results


• A protocol compiler that is:

Our Results



A protocol compiler that is: \bullet

round-optimal malicious secure protocol.

Our Results

• Round-Optimal, i.e., transforms round-optimal semi-honest secure protocol to



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Our Results

• Round-Optimal, i.e., transforms round-optimal semi-honest secure protocol to

• Makes use of simple cryptographic tools/setup: ROM/1-out-2 OT correlations.





Applications **2PC**



Two-round Semi-Honest 2PC

Applications **2PC**







Applications **2PC**

Two-round Malicious OT

Two-round NISC Protocol

Two-round twosided NISC Protocol















Applications MPC



Applications MPC





Applications MPC

Three-round malicious MPC





Applications MPC

Three-round malicious MPC

Round-Optimality follows from [ABGIS 20]







Two-round Semi-Malicious* MPC

Applications MPC

Three-round malicious MPC

Round-Optimality follows from [ABGIS 20]







Two-round Semi-Malicious* MPC

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Round-Optimality follows from [ABGIS 20]

OT Correlations

Two-round malicious MPC







Two-round Semi-Malicious* MPC

Applications MPC

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OT Correlations

Two-round malicious MPC

Prior work [IKSS 21] required complex multiparty correlations.





In this talk

Two-round Semi-Honest 2PC



Two-round Malicious OT Protocol

In this talk





Two-round Malicious OT Protocol





















 $m_b = Dec(z_1, \dots, z_m)$



Malicious security against any A that corrupts one of the clients and a constant fraction of servers [IKP 10].





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Outer Protocol

Two-round Semi-Honest* 2PC for $S_1(\ \cdot \ , \ \cdot \)$

Two-round Semi-Honest* 2PC for $S_2(\ \cdot\ ,\ \cdot\)$

Two-round Semi-Honest* 2PC for $S_m(\ \cdot\ ,\ \cdot\)$

Inner Protocol

Malicious security against any A that corrupts one of the clients and a constant fraction of servers [IKP 10].





 $m_b = Dec(z_1, \dots, z_m)$



Outer Protocol

Two-round Semi-Honest* 2PC for $S_1(\ \cdot \ , \ \cdot \)$

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Two-round Semi-Honest* 2PC for $S_m(\ \cdot\ ,\ \cdot\)$

Inner Protocol

Watchlist Protocol





$$x = b$$



 (x_1,\ldots,x_m)



$$x = b$$

 (y_1, \ldots, y_m)
















$$(x_1,\ldots,x_m)$$



x = b



$$(x_1, \dots, x_m)$$

$$\{(x_i, r_i)\}_{i \in [m]}$$

$$x = b$$



$$(x_1, \dots, x_m)$$

$$\{(x_i, r_i)\}_{i \in [m]}$$

$$x = b$$



$$(x_{1}, ..., x_{m}) \qquad (y_{1}, ..., y_{m})$$

$$(y_{1}, ..., y_{m}) = K \subseteq [m] = \{(x_{i}, r_{i})\}_{i \in K} = b$$

$$(y_{1}, ..., y_{m}) = K \subseteq [m] = \{(x_{i}, r_{i})\}_{i \in K} = y = (m_{0}, m_{1})$$

Watchlist Protocol





$$(y_1, \dots, y_m)$$

$$K \subseteq [m]$$

$$\{(x_i, r_i)\}_{i \in K}$$

$$y = (m_0, m_1)$$

Check if for each $i \in K$, (x_i, r_i) is a valid input-randomness pair consistent with the transcript.











K need not be kept secret but needs to be randomly chosen after the receiver message.



 $(x_1, ..., x_m)$



$$x = b$$









$$Com((x_i, r_i))\}_{i \in [m]}$$

$$K \subseteq [m]$$

$$V = (m_0, m_1)$$







$$Com((x_{i}, r_{i}))\}_{i \in [m]}$$

$$K \subseteq [m]$$

$$en(c_{i}), (x_{i}, r_{i})\}_{i \in K}$$

$$y = (m_{0}, m_{1})$$







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$$(y_1, \ldots, y_m)$$

 $\{c_i = \mathsf{Com}((x_i, r_i))\}_{i \in [m]}$ $K \subseteq [m]$ $\{\mathsf{Open}(c_i), (x_i, r_i)\}_{i \in K}$

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Our Solution Fiat-Shamir paradigm



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Thank you!