

Round-Optimal Black-Box Secure Computation from Two-Round Malicious OT

Yuval Ishai, **Dakshita Khurana**, Amit Sahai, Akshayaram Srinivasan

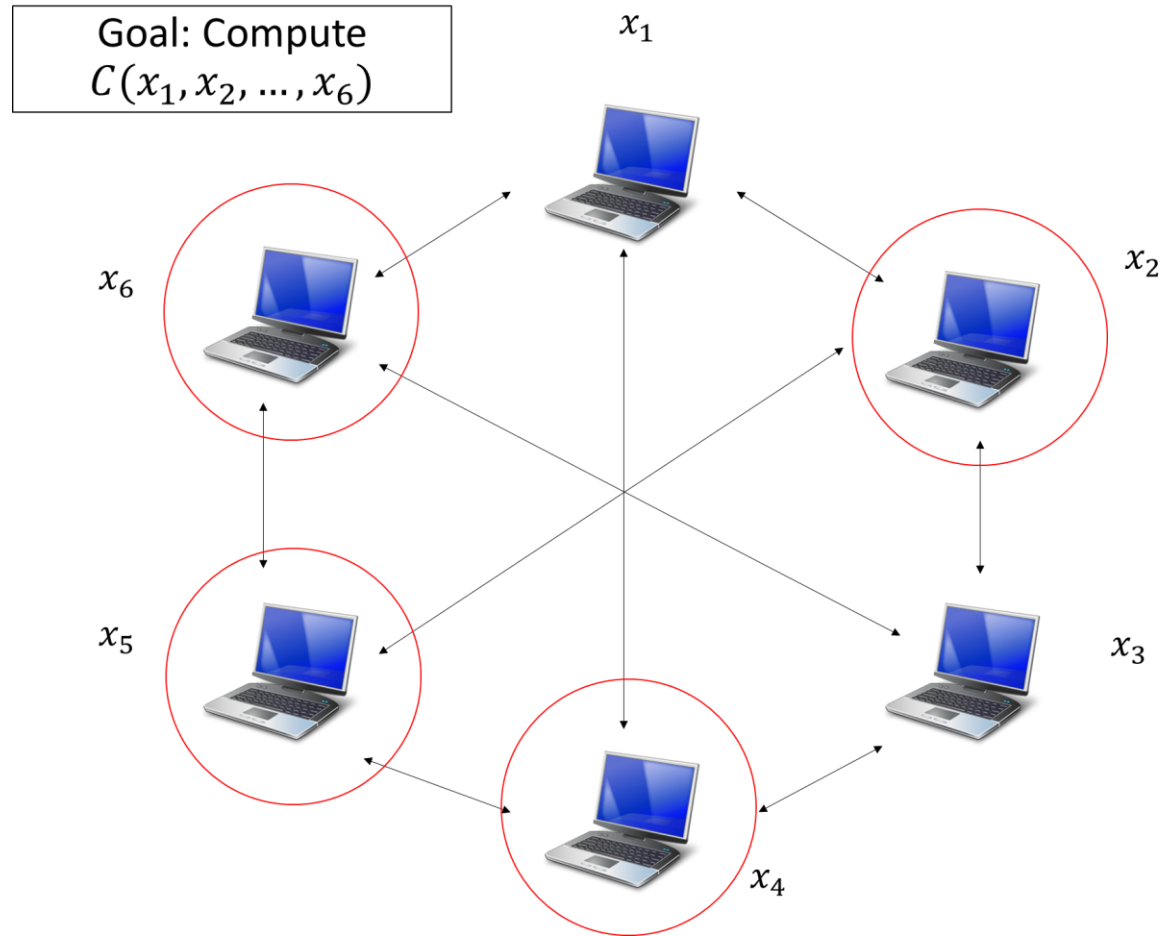
Technion

UIUC

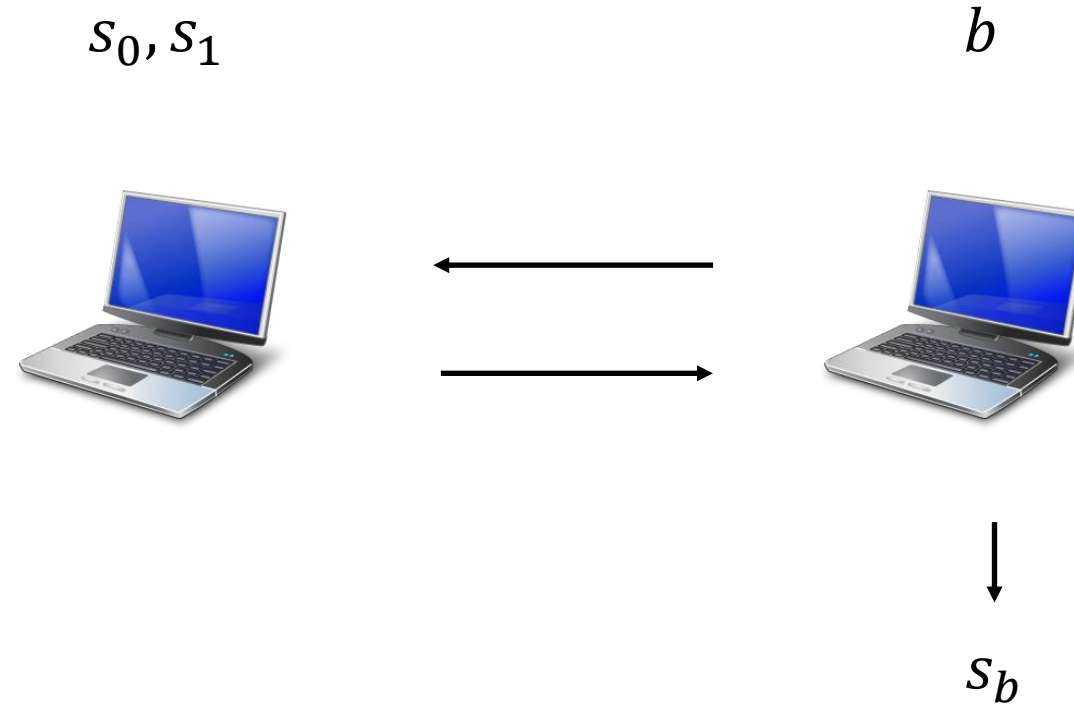
UCLA

TIFR

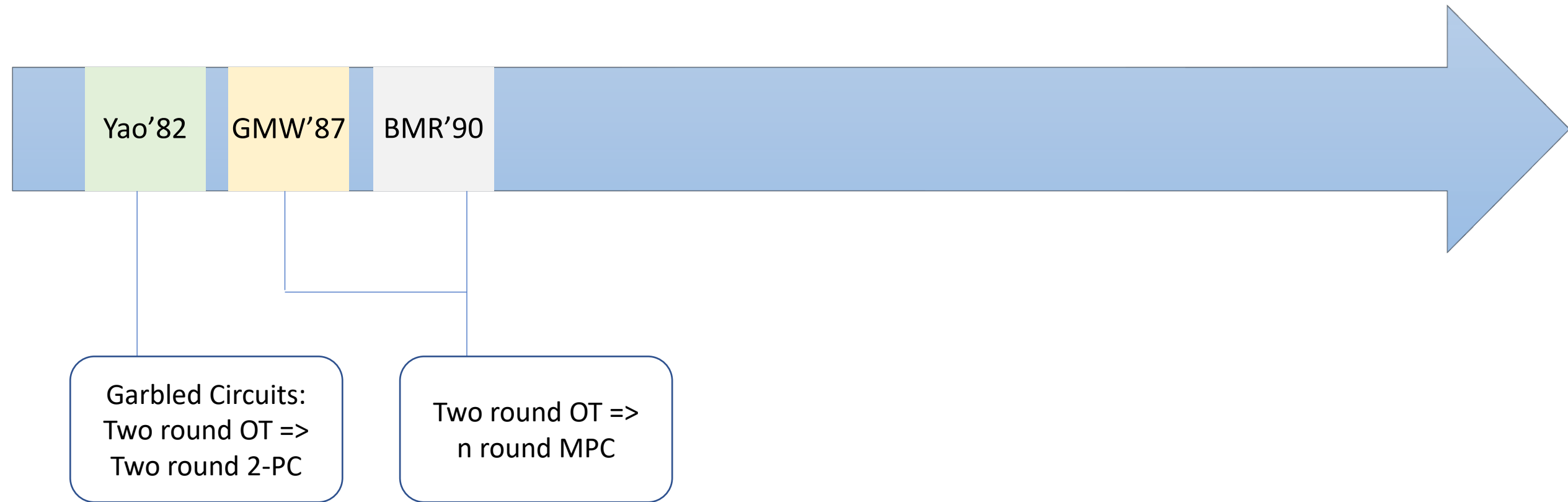
Multiparty Computation (MPC)



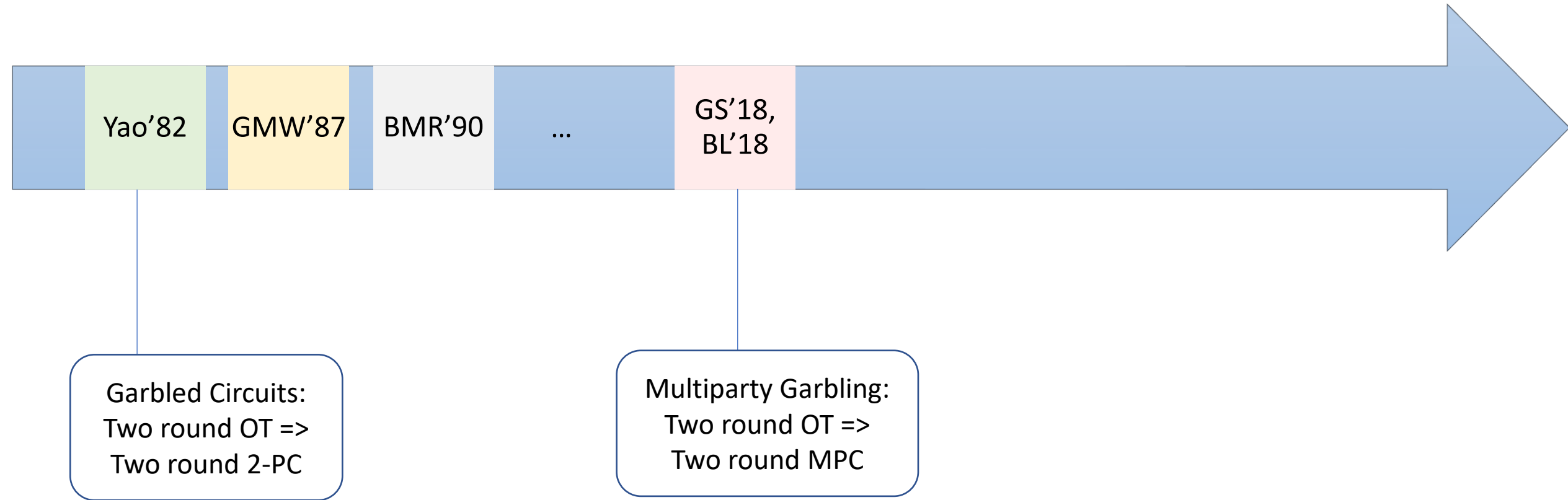
Building Block: Oblivious Transfer



Progress in Understanding MPC

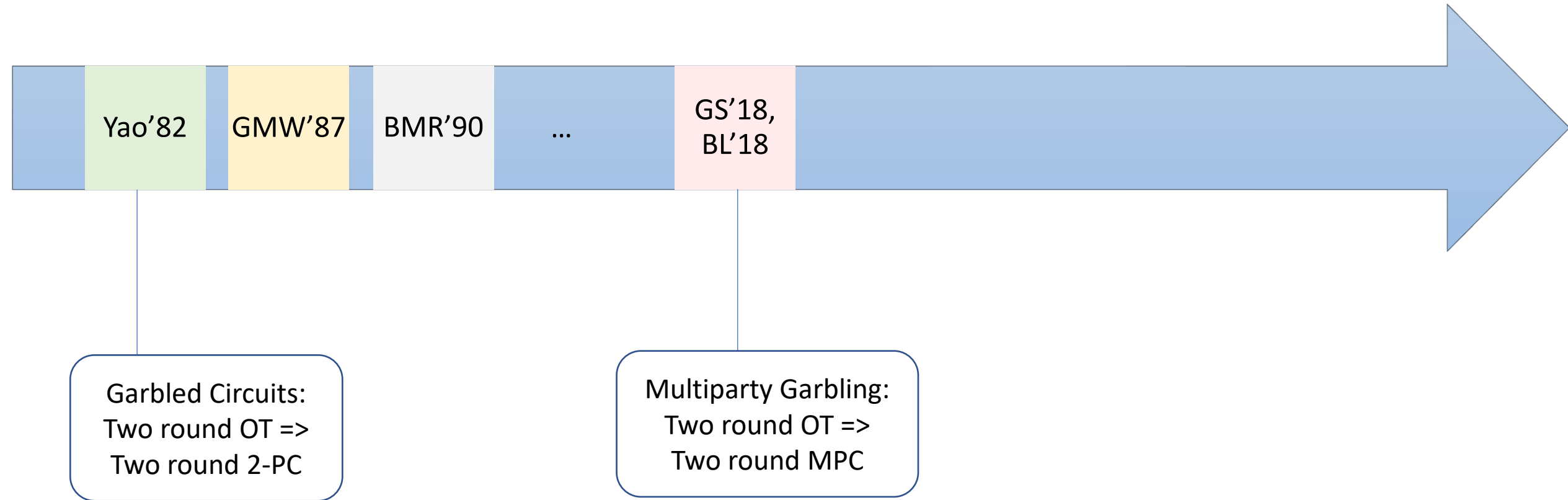


Progress in Understanding MPC



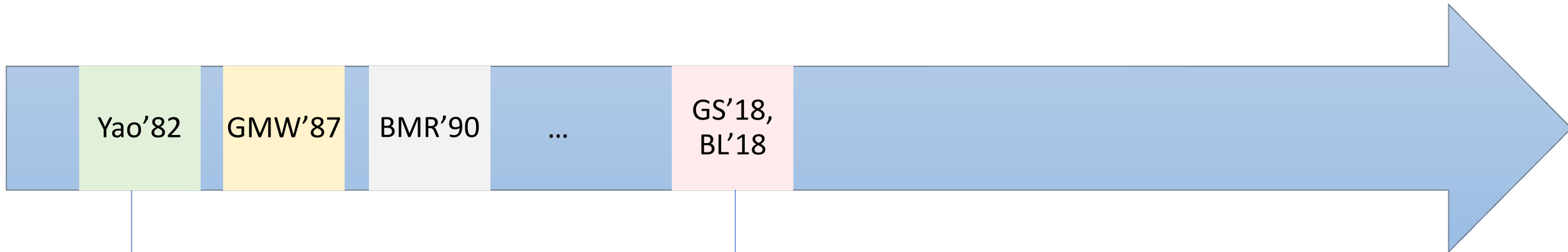
Progress in Understanding MPC

Malicious Security in the CRS Model



Progress in Understanding MPC

Malicious Security in the CRS Model



Garbled Circuits:
Two round OT =>
Two round 2-PC

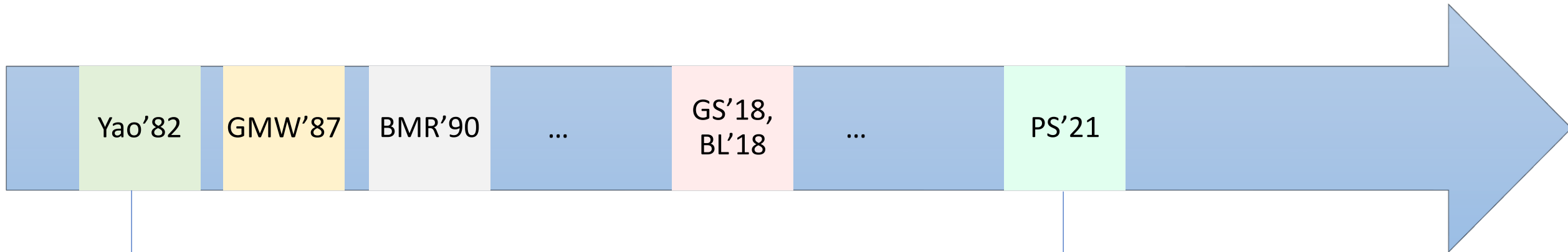
BLACK-BOX: [IKOPS'11]

Multiparty Garbling:
Two round OT =>
Two round MPC

NON-BLACK-BOX

Progress in Understanding MPC

Malicious Security in the CRS Model



Garbled Circuits:
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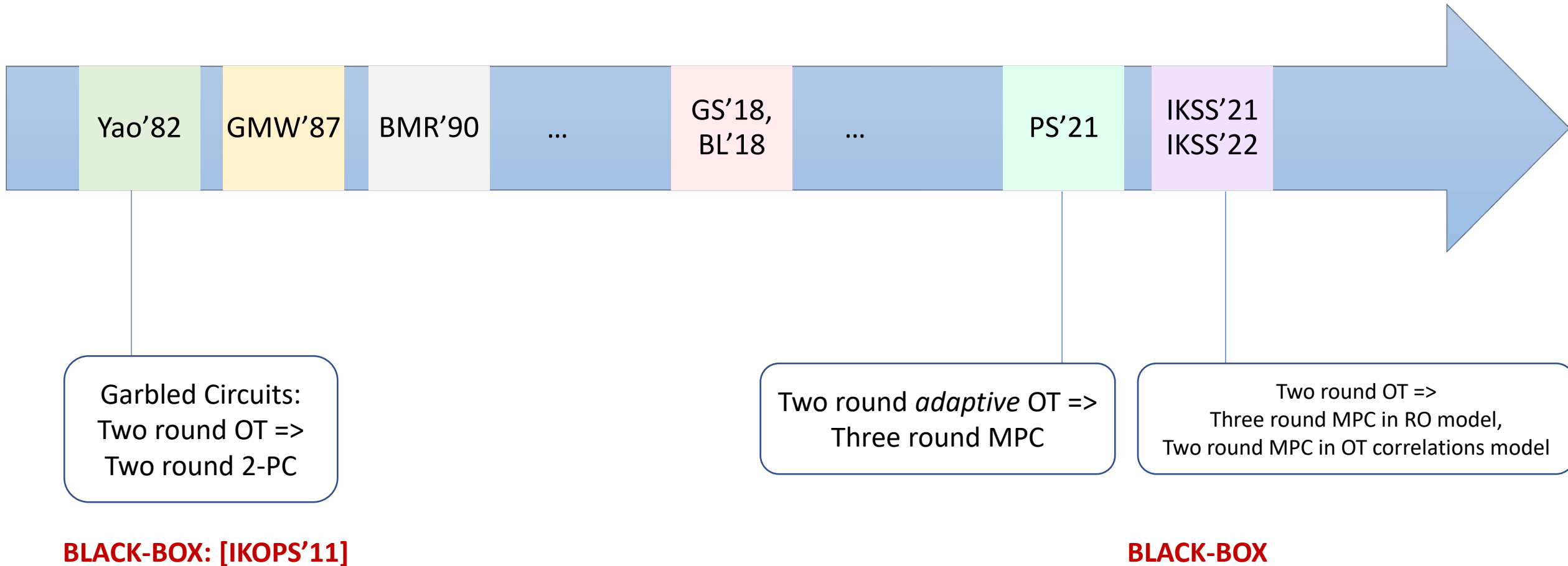
BLACK-BOX: [IKOPS'11]

Two round *adaptive* OT =>
Three round MPC

BLACK-BOX

Progress in Understanding MPC

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Progress in Understanding MPC

Malicious Security in the CRS Model



Garbled Circuits:
Two round OT =>
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BLACK-BOX: [IKOPS'11]

Two round OT =>
Three round MPC in CRS model

Matches known lower bound [ABGIS'20]

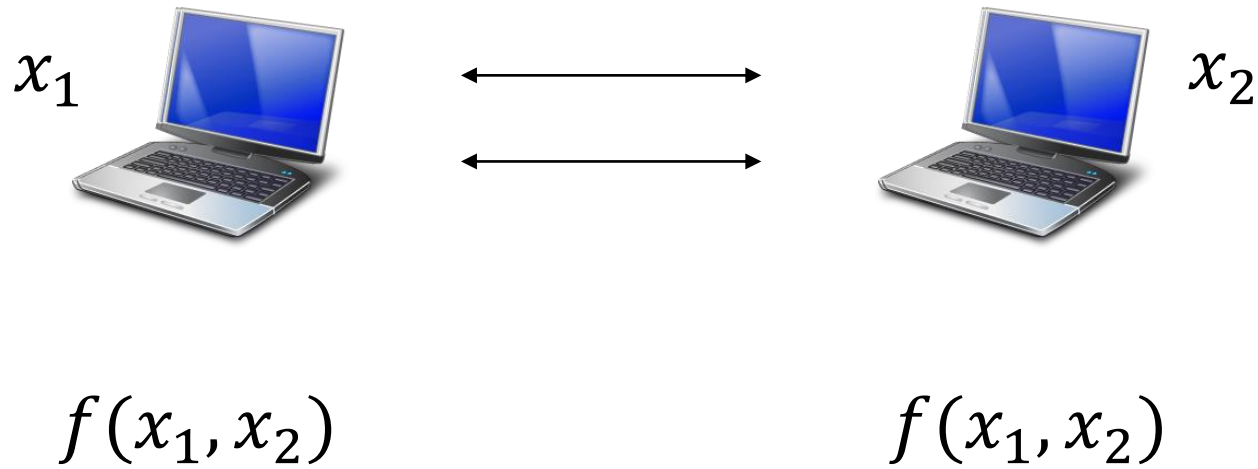
BLACK-BOX

Our Results

In the common reference string (CRS) model,

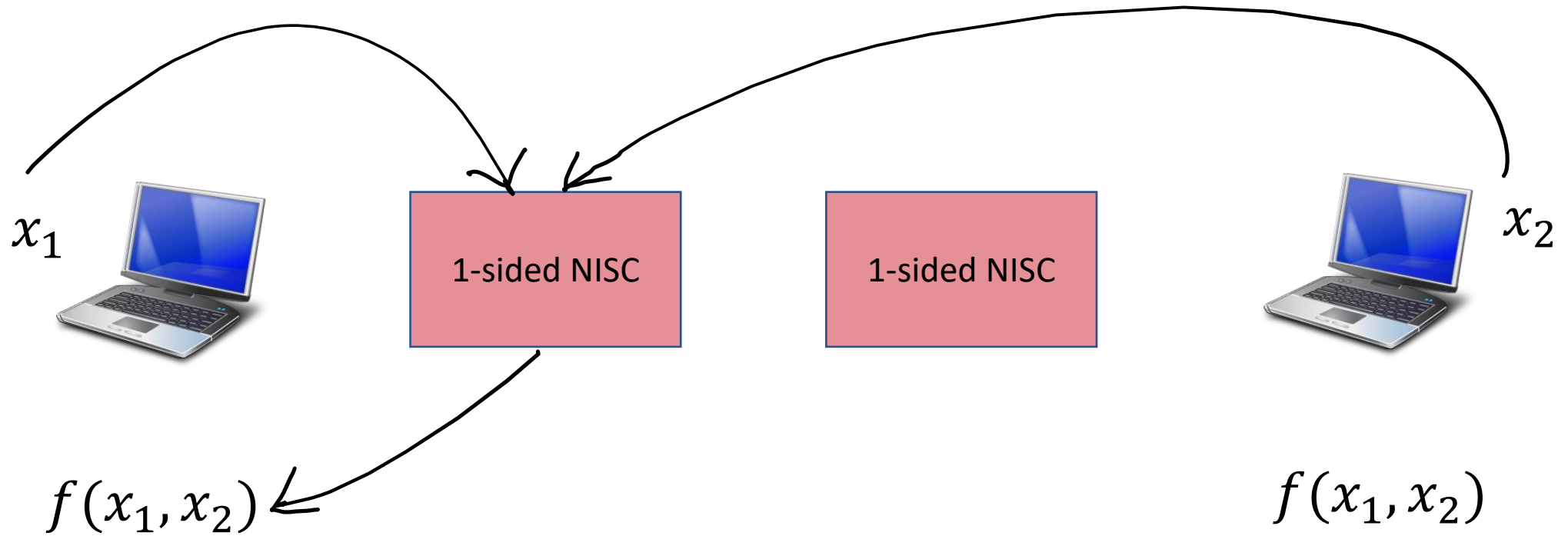
- 2 round maliciously secure OT \rightarrow 3 round maliciously secure MPC
- 2 round maliciously secure OT \rightarrow 2 round 2-sided NISC

Two-Sided Non-Interactive Secure Computation (NISC)



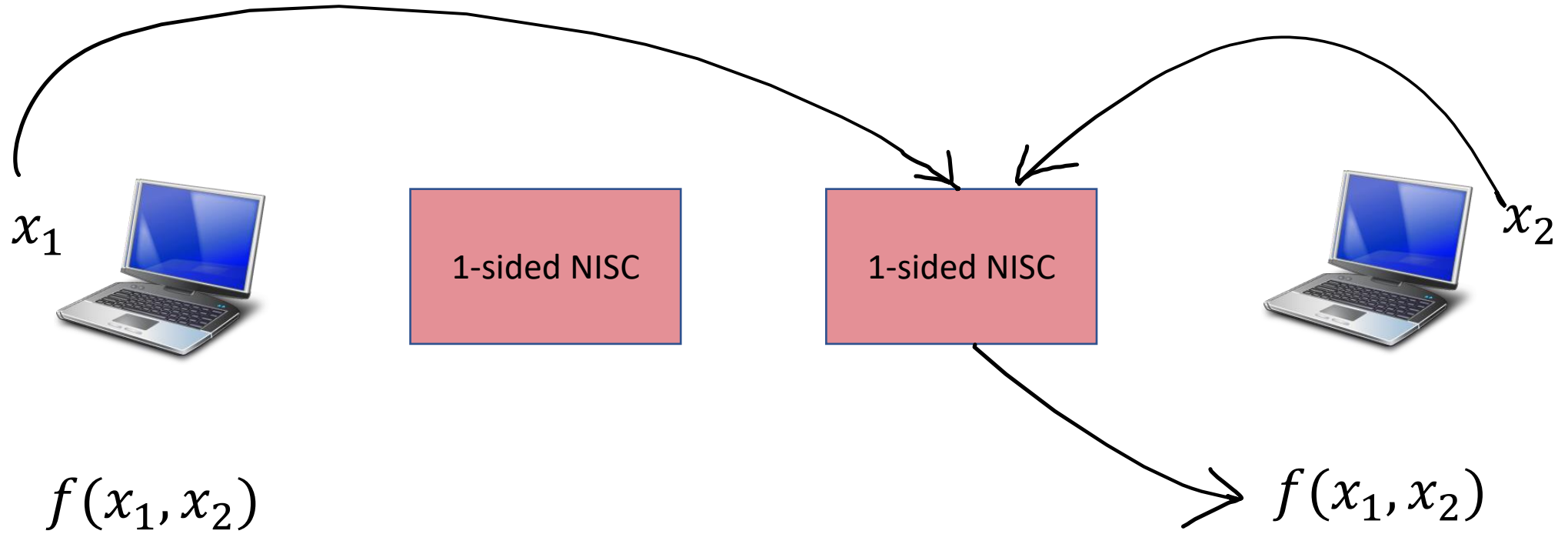
Known result: Two-round One-sided NISC from black-box use of two-round OT [IKOPS'11]

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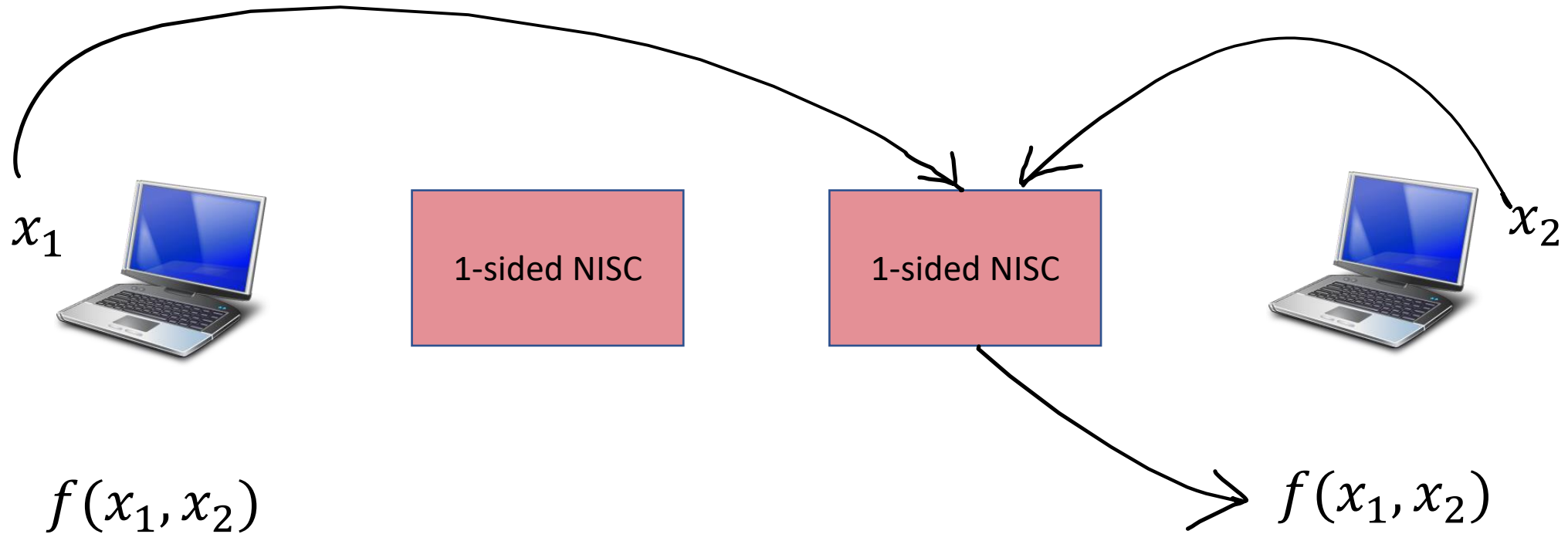


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Two-Sided Non-Interactive Secure Computation (NISIC)



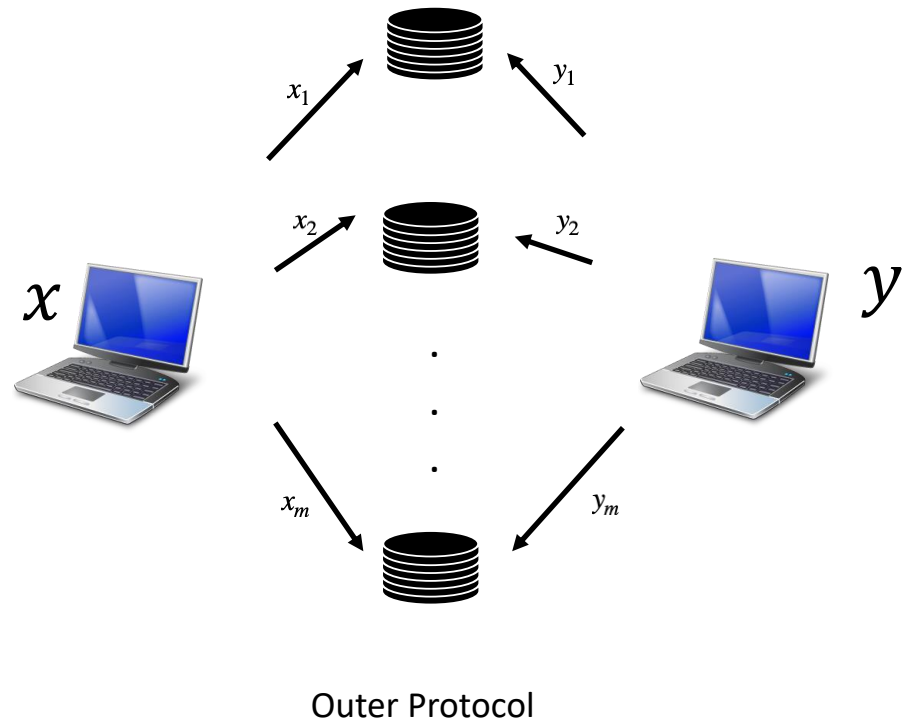
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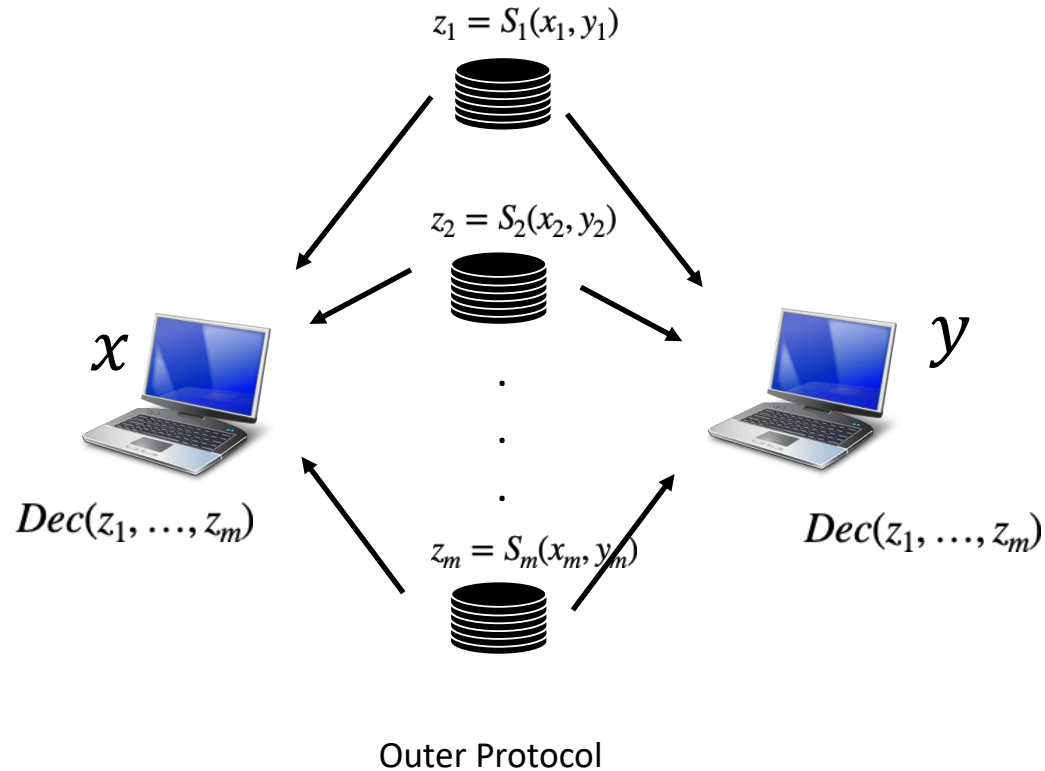
Key Technical Challenge: Ensure input consistency across executions

Use MPC-in-the-head techniques [IPS'08] to check input consistency

Quick Recap of the IPS Compiler

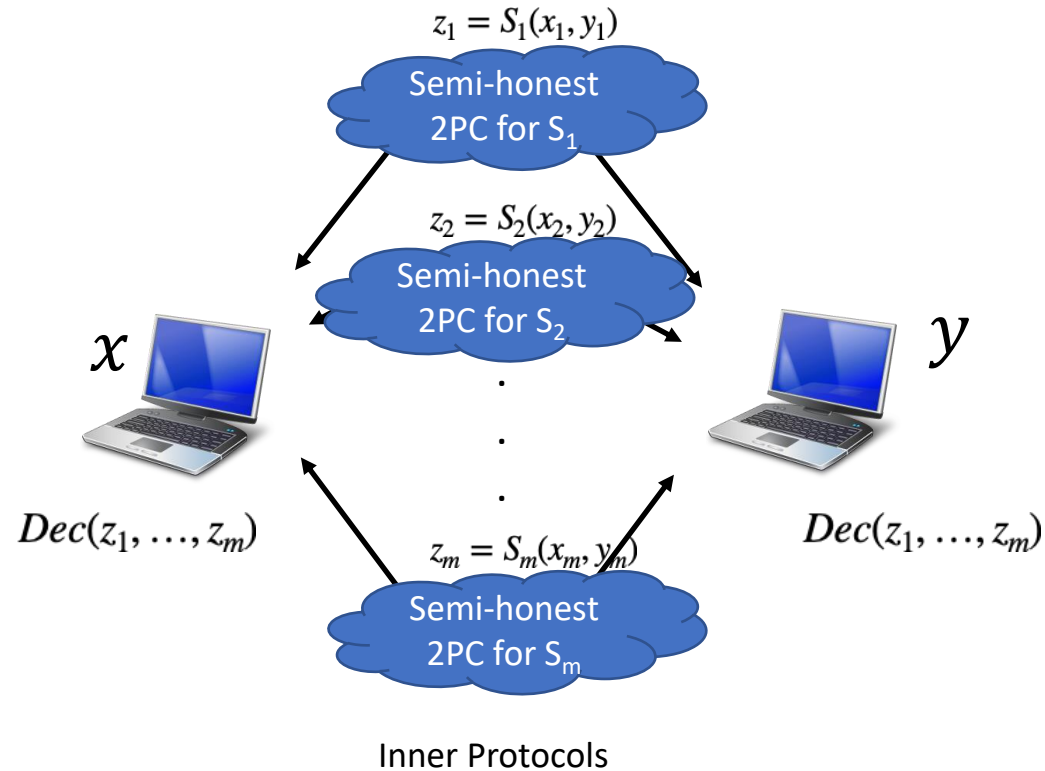


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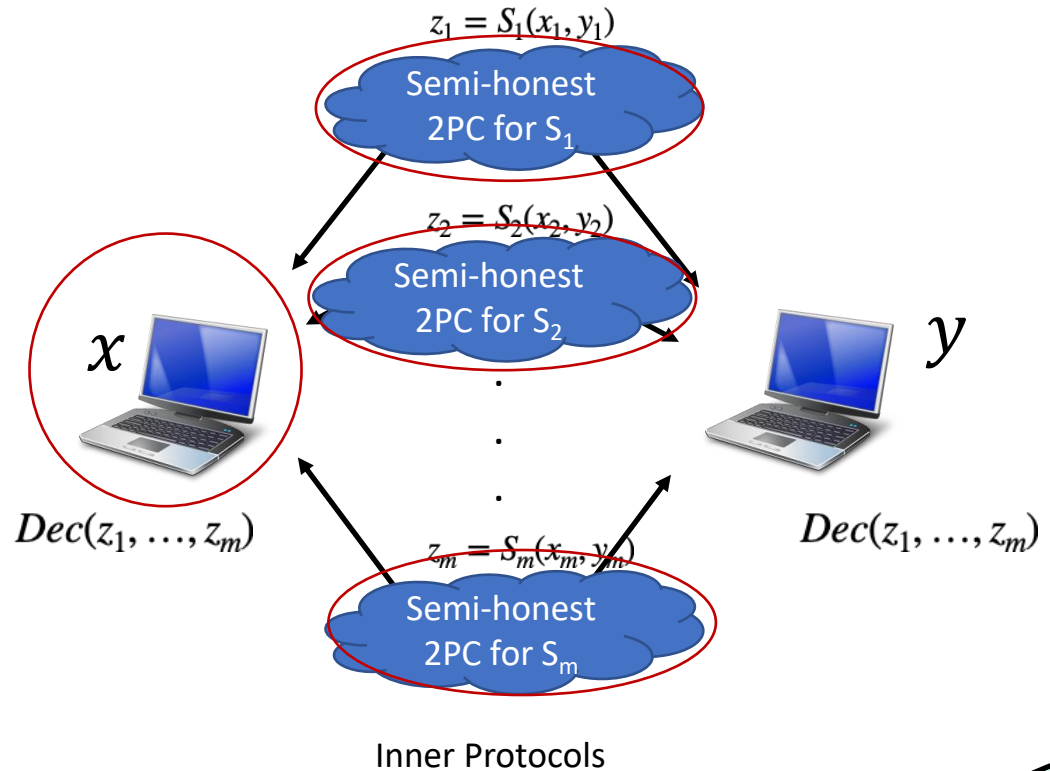


Malicious security against corruption of one of the clients and a constant fraction of servers [IKP 10].

Quick Recap of the IPS Compiler

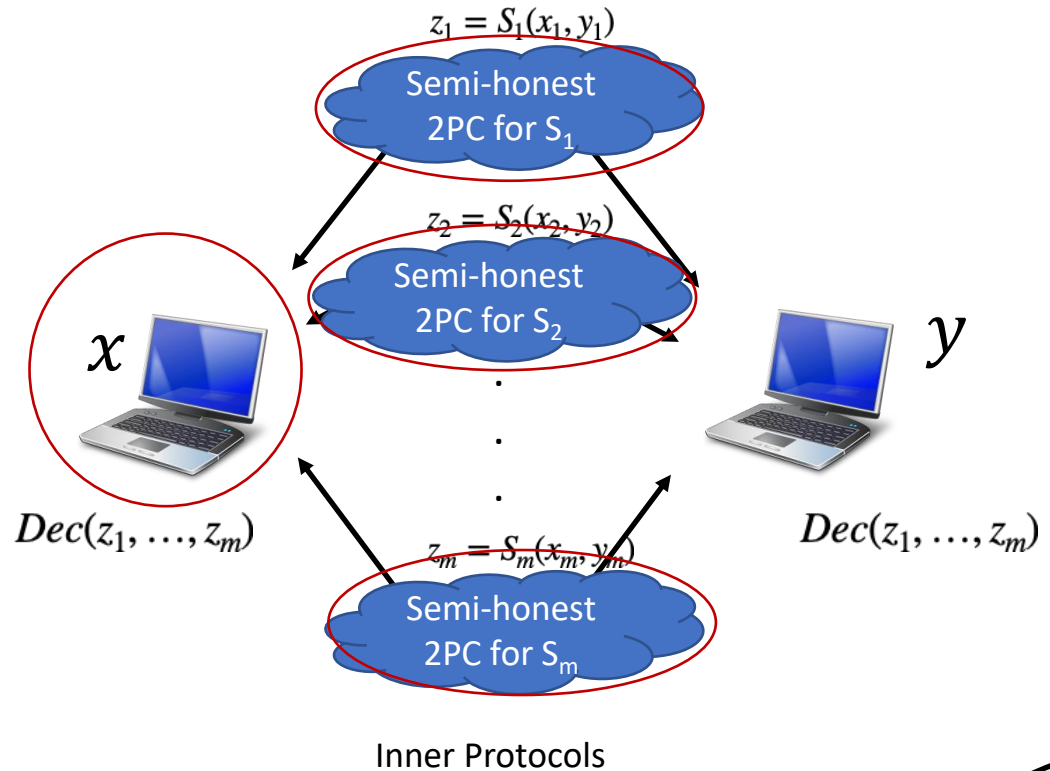


Quick Recap of the IPS Compiler



Need to ensure that adversary only cheats in a constant fraction

Quick Recap of the IPS Compiler



WATCHLISTS:
Cut-and-choose checks on a
random subset of inner protocols

Need to ensure that adversary
only cheats in a constant fraction

Specific choice of inner protocol and watchlists

- Inner Protocol should satisfy an additional output equivocality property
 - Two-sided NISC: Any one-sided NISC (which can be based on OT [IKOPS'11])
Alternatively Garbled circuits + OT
 - MPC: Robust MPC building on [PS'21]

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Alternatively Garbled circuits + OT
 - MPC: Robust MPC building on [PS'21]
- Watchlist mechanism should allow the simulator to decide and program which executions will be watched by malicious parties
 - Build this using ideas from [IKOPS'11]

Another Perspective

- Start with simple round-optimal realizations that BB use *semi-honest OT*
semi-honest 2-sided NISC (Garbled circuits + OT)
semi-honest MPC [PS'21]

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- Start with simple round-optimal realizations that BB use *semi-honest OT*
semi-honest 2-sided NISC (Garbled circuits + OT)
semi-honest MPC [PS'21]
- Plug in maliciously secure OT
- The result is an inner protocol, which is bootstrapped to full malicious security via the IPS compiler

Open Questions

- Obtain our results from BB use of two-round *semi-honest* OT
- Base 3 round MPC on 3 *round* maliciously secure OT
- Can similar results be obtained in the *OT-hybrid model*?