

The Indifferentiability of the Duplex and its Practical Applications



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Asiacrypt 2023



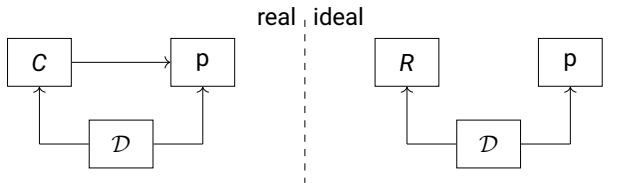
- 1 Background
- 2 Indifferentiability of the Duplex from the Online Random Oracle
- 3 Applications of the Indifferentiability of the Duplex

Background

Indistinguishability vs. Indifferentiability



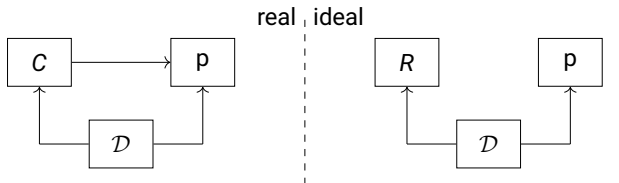
Indistinguishability



Indistinguishability vs. Indifferentiability



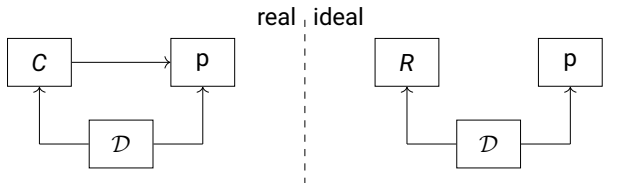
Indistinguishability



- Focused on a specific property



Indistinguishability

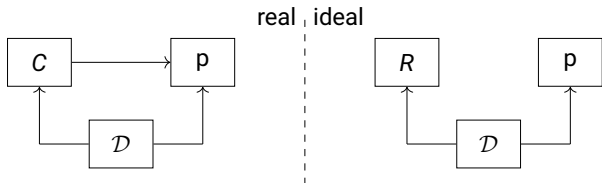


- Focused on a specific property
- Requires keyed constructions

Indistinguishability vs. Indifferentiability

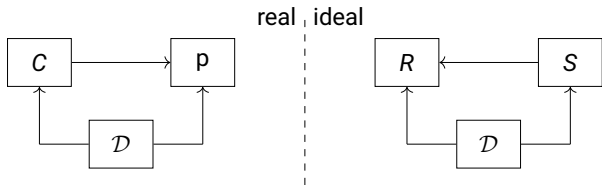


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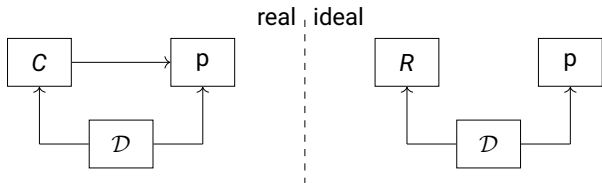
Indifferentiability [MRH04]



Indistinguishability vs. Indifferentiability

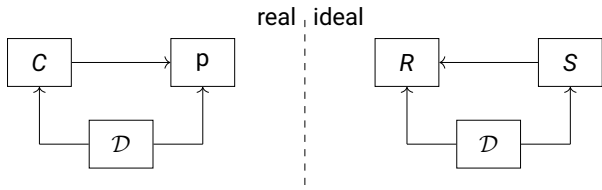


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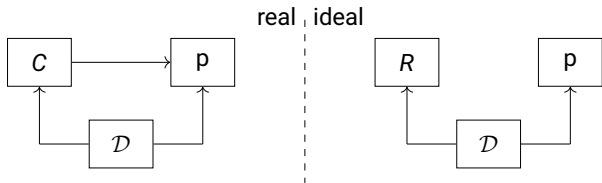
Indifferentiability [MRH04]



- Used to build an ideal primitive R

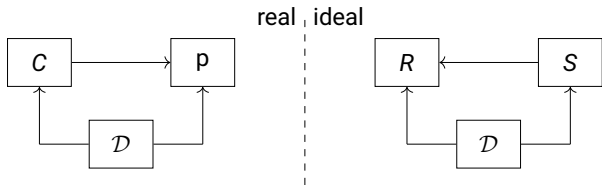


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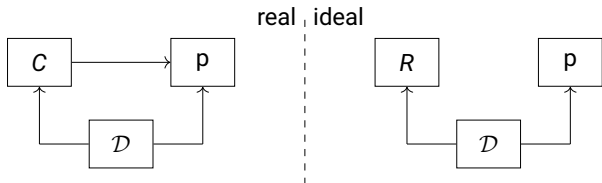
Indifferentiability [MRH04]



- Used to build an ideal primitive R
- For (un)keyed constructions

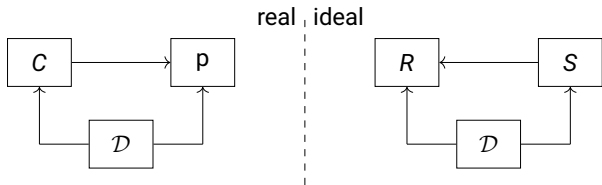


Indistinguishability



- Focused on a specific property
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Indifferentiability [MRH04]



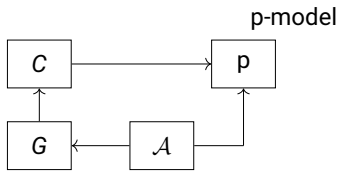
- Used to build an ideal primitive R
- For (un)keyed constructions
- Covers multiple security properties



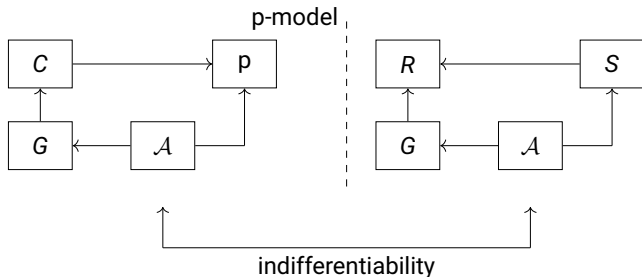
If C is **indifferentiable** from R ,
then, in the p -model, it **has the same security properties** as R .



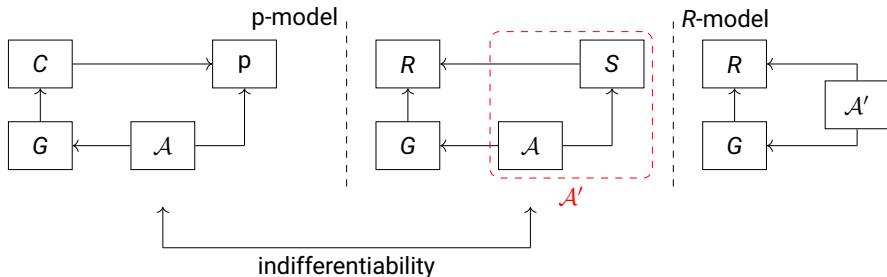
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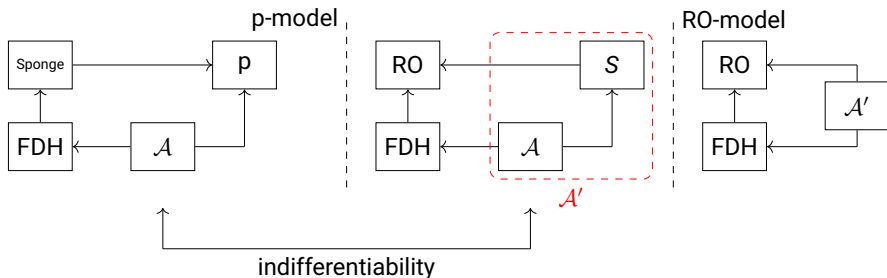


Composability [MRH04]

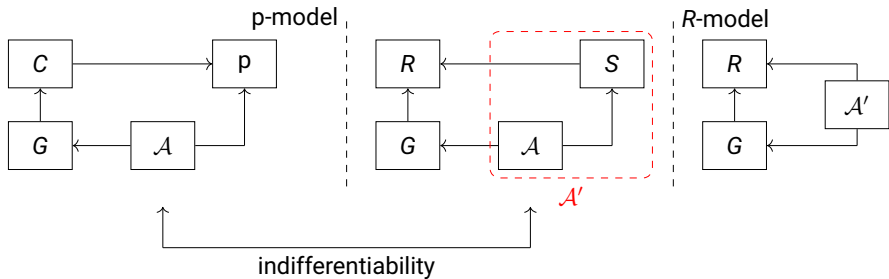
Full Domain Hash (FDH) Example



If C is indistinguishable from R ,
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If C is indistinguishable from R ,
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→ C and R need to share the **same interface**

Indifferentiability from Idealized Model



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Primitive C	Idealized Model R
Hash Function	Random Oracle [Bellare and Rogaway, <i>ACM CCS 93</i>]
Block Cipher	Ideal Cipher [Holenstein, Künzler, and Tessaro, <i>43rd ACM STOC</i>] [Andreeva et al., <i>CRYPTO 2013, Part I</i>]
Authenticated Encryption	Random Injection [Barbosa and Farshim, <i>CRYPTO 2018, Part I</i>]

Indifferentiability from Idealized Model



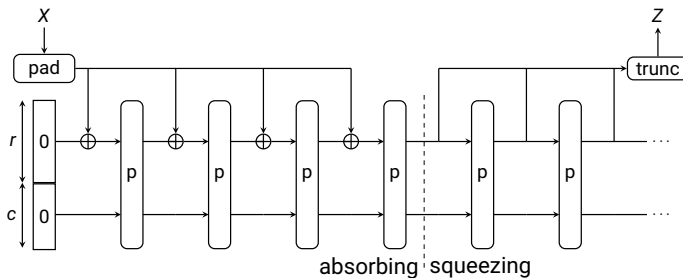
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Duplex	??

The Sponge Construction [Ber+08]



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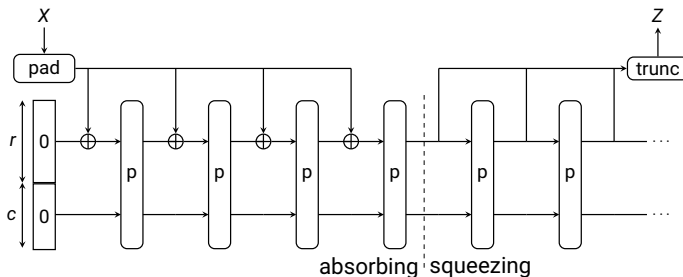


- Basis of multiple NIST standards:
 - ▣ SHA-3, cSHAKE, KMAC, TupleHash, ParallelHash

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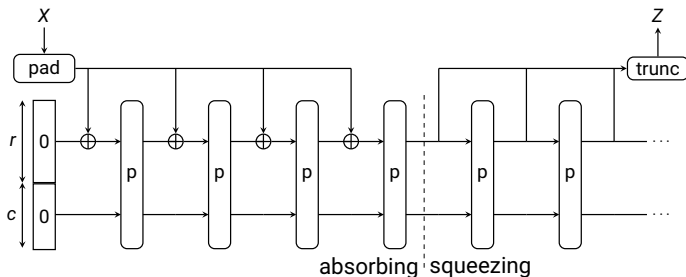


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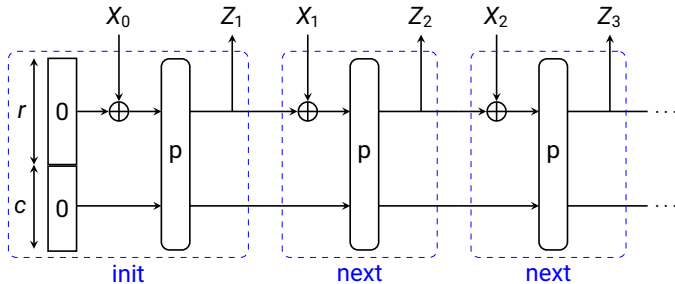


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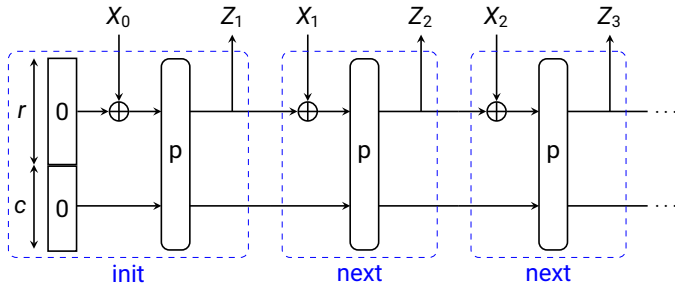
- Basis of multiple NIST standards:
 - ▣ SHA-3, cSHAKE, KMAC, TupleHash, ParallelHash
- Based on a public random permutation p
- Indifferentiable from a Random Oracle with bound $\mathcal{O}\left(\frac{q^2}{2^c}\right)$ [Ber+08]

The Duplex Construction [Ber+12]



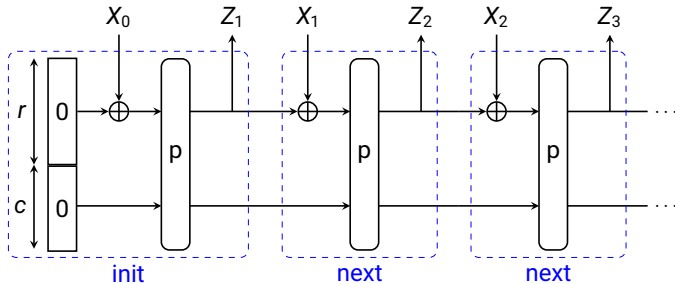
- Allows the construction of one-pass AEAD schemes

The Duplex Construction [Ber+12]



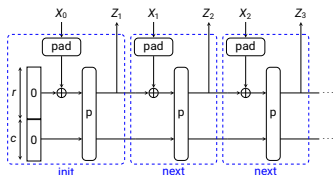
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- Basis of multiple AEAD candidates of the CAESAR & NIST competitions

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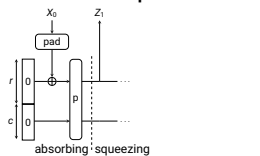
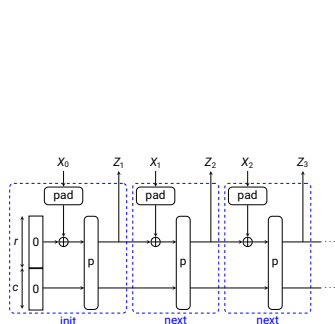


- Allows the construction of one-pass AEAD schemes
- Basis of multiple AEAD candidates of the CAESAR & NIST competitions
- Stateful construction that supersedes the Sponge

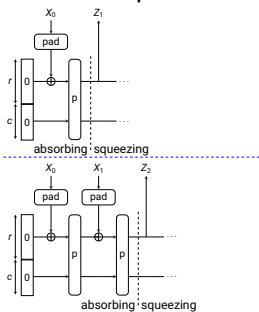
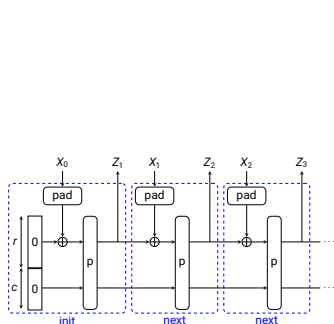
→ Reduces an instance of the Duplex to a sequence of Sponge calls



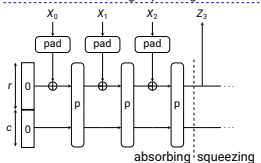
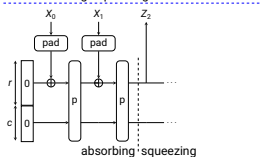
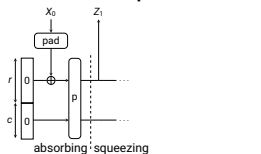
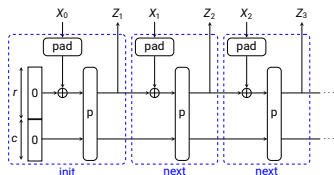
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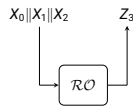
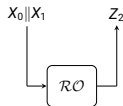
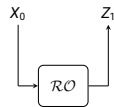
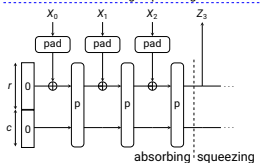
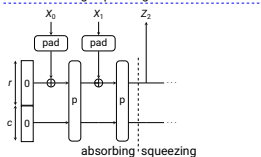
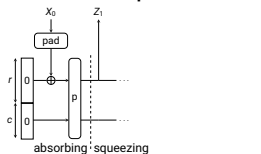
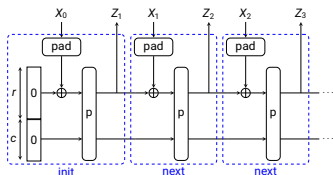
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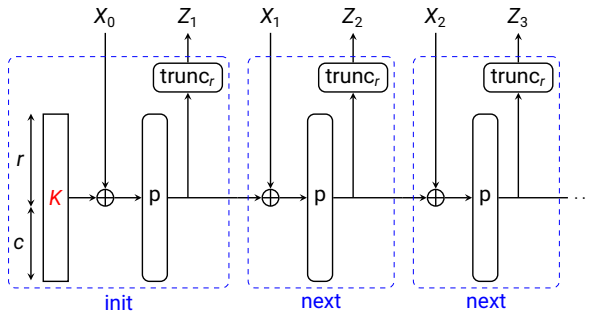
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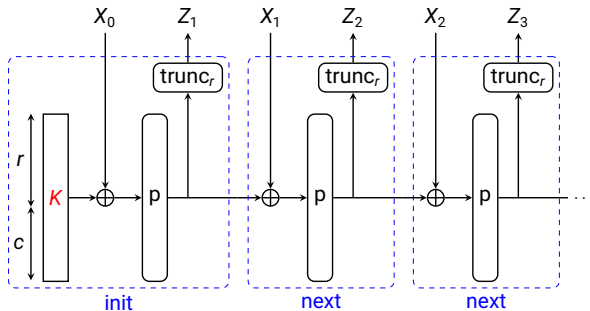


The Full-State Keyed Duplex [MRV15]



- Newer work focuses on the indistinguishability of the keyed Duplex

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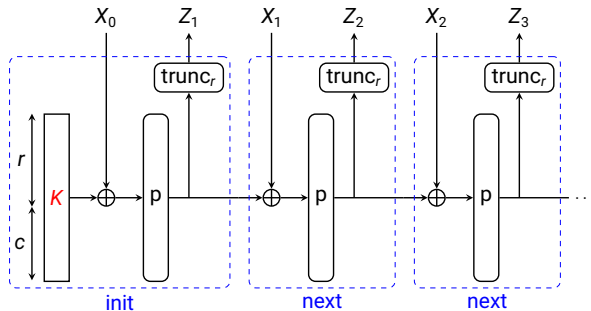


- Newer work focuses on the indistinguishability of the keyed Duplex
- Better bounds

The Full-State Keyed Duplex [MRV15]



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- Newer work focuses on the indistinguishability of the keyed Duplex
- Better bounds
- Improved absorption performance

The Duplex as a General-Purpose Primitive



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Used for:

- **Authenticated Encryption/AEAD**



Used for:

- **Authenticated Encryption/AEAD**
- **Online Hash**
 - Stateful Hash Object (SHO) within the Noise Protocol framework



Used for:

- **Authenticated Encryption/AEAD**
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- **MAC, Symmetric Ratcheting and Pseudorandomness Generation**
 - STROBE protocol framework (lib based only on the Duplex)



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→ Prior security analyses focused on specific usage, and not as a general-purpose primitive (keyed or unkeyed)



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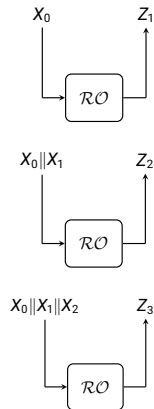
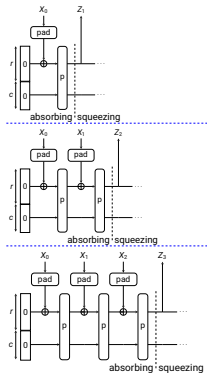
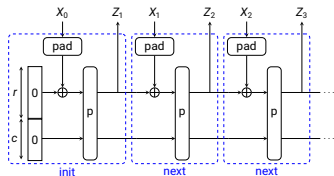
→ Need an idealized model for the Duplex

Indifferentiability of the Duplex from the Online Random Oracle

Limitations of [Ber+12] Security Analysis



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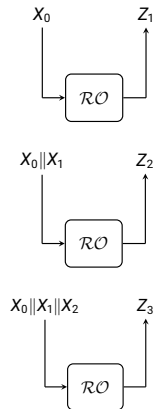
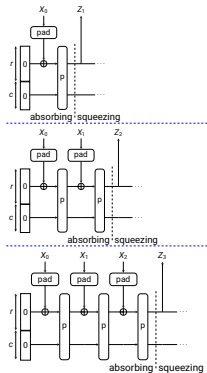
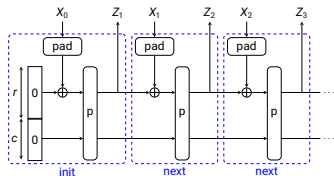


- Mismatching interface → cannot directly apply composition thm

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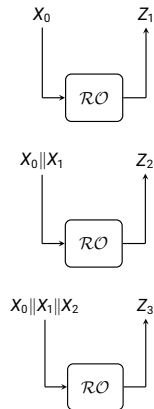
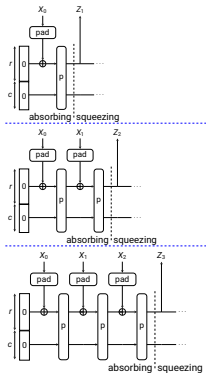
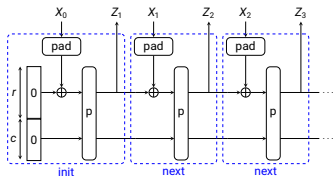


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Limitations of [Ber+12] Security Analysis



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- Mismatching interface → cannot directly apply composition thm
- Needs an extra step
- Needs sponge-compliant padding in every call to p within the Duplex

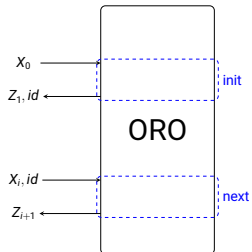
Indifferentiability from Idealized Model



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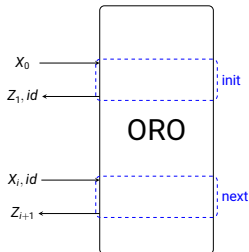
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Duplex	Online Random Oracle (ORO) [This work]

The Online Random Oracle (ORO)



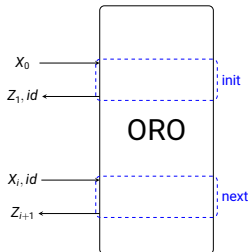
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The Online Random Oracle (ORO)



- Stateful & Online primitive
- To each query, we associate a path and we keep a table for mapping paths to answers
- The path corresponding to init is X_0
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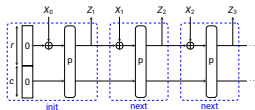
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- The answer associated with a path is sampled at random once
- We updated the syntax to supports multiple concurrent sessions

The Duplex is Indifferentiable from the ORO

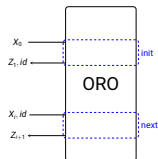


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The Duplex



The ORO



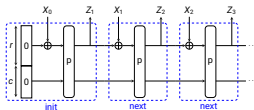
- We show that **the Duplex is indifferentiable from the ORO** with bound $\mathcal{O}\left(\frac{q^2}{2c}\right)$

The Duplex is Indifferentiable from the ORO

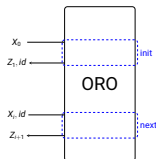


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The Duplex



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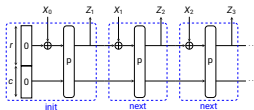
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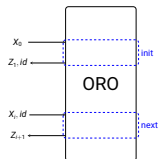


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The Duplex



The ORO



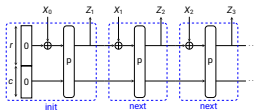
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- We obtain an **efficient simulator**

The Duplex is Indifferentiable from the ORO

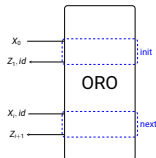


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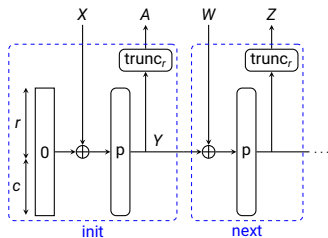


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- We give a **proof using the code-based framework** from Bellare–Rogaway
- We obtain an **efficient simulator**
- No padding required

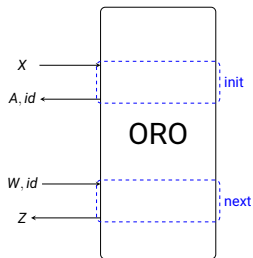
Full-State Duplex is Differentiable from the ORO



The Full-State Duplex



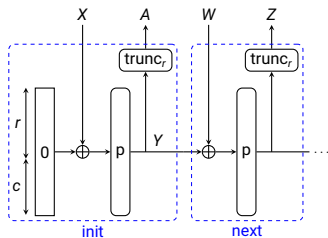
The ORO



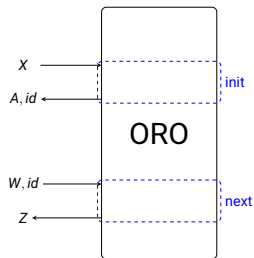
Full-State Duplex is Differentiable from the ORO



The Full-State Duplex



The ORO

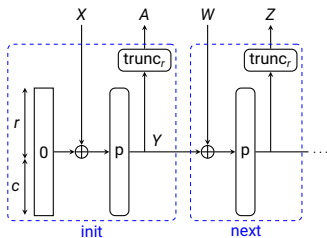


- The full string Y is recoverable in the real world through the access to p

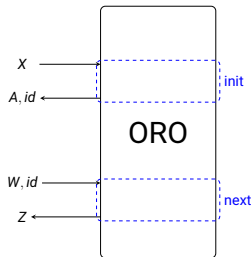
Full-State Duplex is Differentiable from the ORO



The Full-State Duplex



The ORO

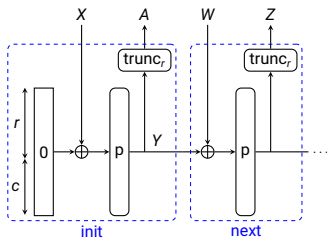


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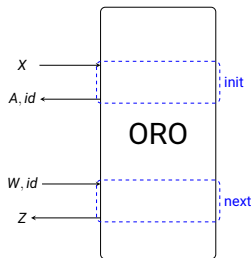
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The ORO



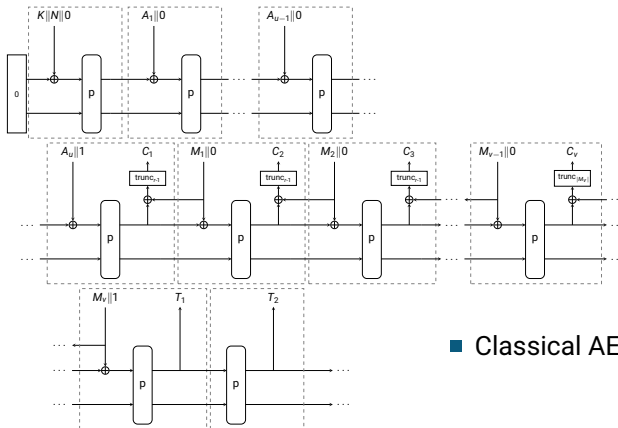
- The full string Y is recoverable in the real world through the access to p
- It is possible to mount a collision $Y \oplus W = Y' \oplus W'$ in the real world
- In the ideal world, the input path to the ORO will be different

Applications of the Indifferentiability of the Duplex

A Nonce-Based Variant of SpongeWrap



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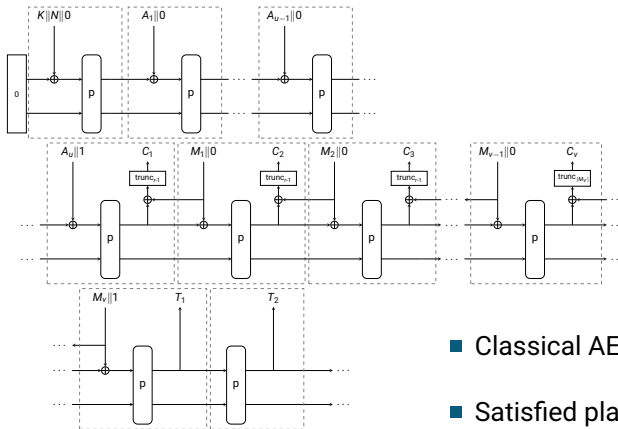


■ Classical AEAD template

A Nonce-Based Variant of SpongeWrap



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- Classical AEAD template
- Satisfied plain AEAD security



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→ useful for disk encryption, HSM, KMS
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→ models fault-injection attacks
- CMT-AEAD: **commitment** security
→ useful for message franking, key rotation



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- We use **composability** to translate the results in the random-permutation model
- Allow us to **bypass a complex analysis** in the random-permutation model
- We obtain the **first one-pass** AEAD scheme to **achieve KDM-AEAD, RKA-AEAD and CMT-AEAD** security



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