

A Direct Key Recovery Attack on SIDH

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Contribution

- Merge of two papers:
 - M.–Martindale (2022), independent from Castryck–Decru (2022)
 - Wesolowski (2022)
- Proof-of-Concept Implementation — mainly by Panny and Pope — available at
<https://github.com/Breaking-SIDH/direct-attack>

Supersingular Isogeny with Torsion (SSI-T)

SSI-T

Let p be a large prime. There are **public** elliptic curves E_{start} and E_{Bob} , and a **secret** isogeny $\varphi_{\text{Bob}} : E_{\text{start}} \rightarrow E_{\text{Bob}}$ of degree 3^b .

Given $\varphi_{\text{Bob}}(E_{\text{start}}[2^a])$,

compute φ_{Bob} .

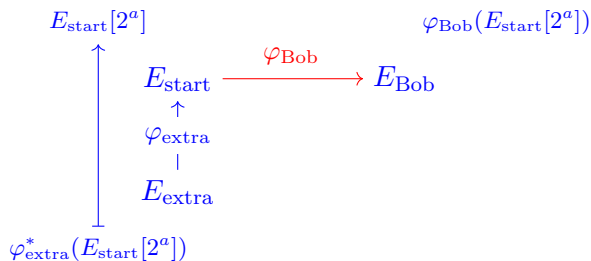
History of the SIDH problem

- 2011 De Feo, Jao, and Plût introduce SIDH
- 2016 Galbraith, Petit, Shani, Ti give an active attack
- 2017 Petit describes passive attacks on some parameter sets - SIDH not affected
- 2020 de Quehen, Kutas, Leonardi, Martindale, Panny, Petit, Stange give passive attacks on more parameter sets

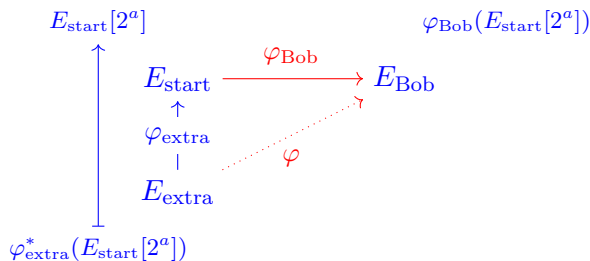
The attack

$$E_{\text{start}}[2^a] \qquad \varphi_{\text{Bob}}(E_{\text{start}}[2^a])$$
$$E_{\text{start}} \xrightarrow{\varphi_{\text{Bob}}} E_{\text{Bob}}$$

The attack

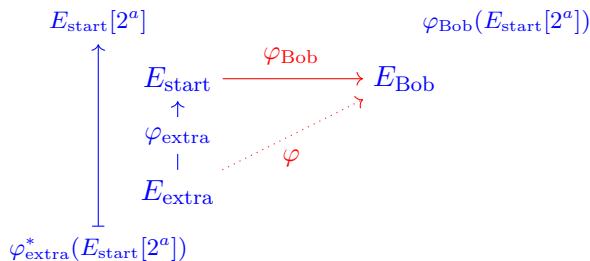


The attack



$$\Phi = \begin{pmatrix} \varphi_{\text{extra}} & -\widehat{\varphi}_{\text{Bob}} \\ * & * \end{pmatrix} : E_{\text{extra}} \times E_{\text{Bob}} \rightarrow E_{\text{start}} \times E_{\text{known}}$$

The attack

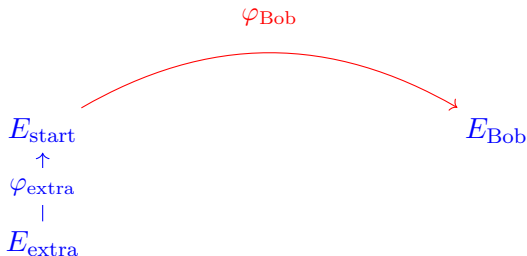


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$$\text{Ker}(\Phi) = \left(3^b E_{\text{extra}}[2^a], \varphi(E_{\text{extra}}[2^a]) \right)$$

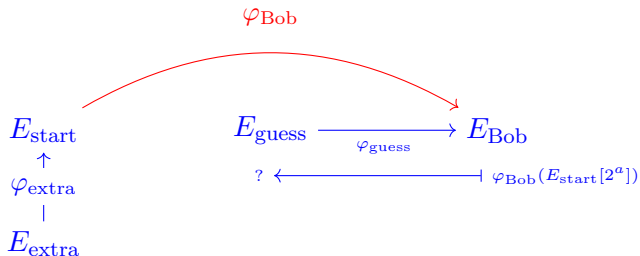
Unknown Endomorphism Ring

- $\deg(\varphi_{\text{extra}}) = 2^a - 3^b$ may not be smooth
- $\deg(\varphi_{\text{extra}}) = e2^{a-j} - 3^{b-i}$ for some small e, i , and j
- This determines the complexity of the attack ($L_{2^a}(1/2)$).



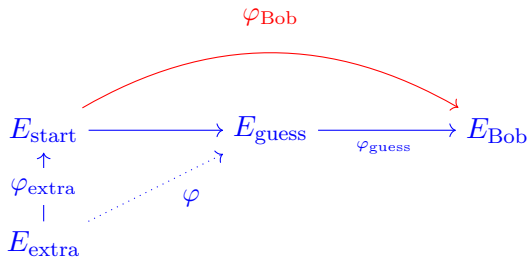
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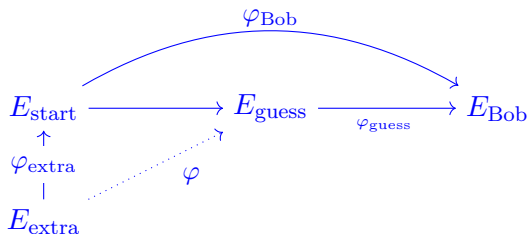
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Theorem

Assume the generalised Riemann hypothesis. Let E_{start} be a supersingular curve, together with a basis $(\alpha_i)_{i=1}^4$ of $\text{End}(E_{\text{start}})$. There exists a polynomial-time algorithm that finds an isogeny $\varphi_{\text{extra}} : E_{\text{extra}} \rightarrow E_{\text{start}}$ of degree $2^a - 3^b$.

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Thanks for your
attention!
Questions?