Composable Long-Term Security with Rewinding

Robin Berger¹, Brandon Broadnax, <u>Michael Klooß^{1 \rightarrow 2}</u>, Jeremias Mechler¹, Jörn Müller-Quade¹, Astrid Ottenhues¹, Markus Raiber¹ 2023-12-02 ©TCC



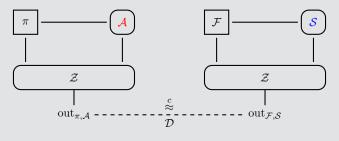


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Computational UC in a nutshell [Can01; Can20]

Security experiment (Computional UC)

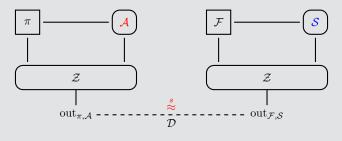
- PPT \mathcal{A} , \mathcal{S} , \mathcal{Z} where \mathcal{Z} outputs a string out.
- **PPT** distinguisher \mathcal{D} gets out.



Statistical UC in a nutshell [Can01; Can20]

Security experiment (Statistical UC)

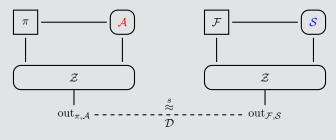
- Unbounded \mathcal{A} , \mathcal{S} , \mathcal{Z} where \mathcal{Z} outputs a string out.
- **Unbounded** distinguisher \mathcal{D} gets out.



Long-Term UC in a nutshell [MU07]

Security experiment (Long-Term UC)

- PPT \mathcal{A} , \mathcal{S} , \mathcal{Z} where \mathcal{Z} outputs a string out.
- \bullet Unbounded distinguisher ${\cal D}$ gets ${\rm out.}$



→ hardness assumptions hold (only) during protocol execution.

Long-Term UC commitments

Possibility results

 \mathcal{F}_{Com} from hardware assumptions (signature card [MU07], PUF+CRS [Mag+22]).

Impossibility result [MU07]

 \mathcal{F}_{Com} is impossible to realize in the CRS-hybrid model or any long-term revealing setup.

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Core problem

- $\bullet\,$ If CRS is not stat. hiding, ${\cal D}$ can extract.
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- ...but what about rewinding?

Our contribution



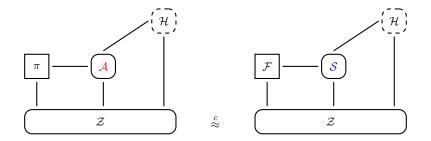
- New notion: Long-term rewinding UC (LTR-UC).
- New possibilities/protocols:
 - LTR-UC-secure \mathcal{F}_{Com} in the CRS-hybrid model (and commit-and-prove ZK).
 - One-sided LTR-UC-secure OT.
- New impossibilities: No full LTR-UC-secure OT from long-term revealing assumptions.
- New tools: Pseudo-oracles and their properties.

Angel-based UC security



Angel-based UC [PS04]

- Global entity, helper or angel \mathcal{H} with "special power".
- E.g.: \mathcal{H} brute-forces commitments under *judiciously chosen circumstances*.



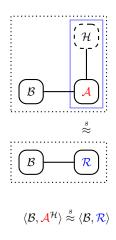
Rewinding-simulatable angels

[CLP10; Goy+15]

- ${\mathcal H}$ is a CCA commitment oracle:
 - $\circ~\mathcal{A}$ can run COM with $\mathcal{H}.$
 - $\circ~\mathcal{H}$ will brute-force extracts accepting commitments.
- ${\mathcal H}$ is simulatable in PPT via rewinding through ${\mathcal R}.$

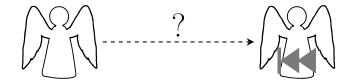
Robust rewinding

- UC simulation is straightline \rightsquigarrow use ${\mathcal H}$
- Security reductions \rightsquigarrow use \mathcal{R} .
- k-robust rewinding: Exempt k-round "left side" from being rewound.



Rewinding-based angels/oracles?

- \bullet LTR-UC also based on a CCA commitment oracle $\mathcal{H}.$
- But what is an "angel/oracle that rewinds"?



Pseudo-Oracles

Oracle/ITM

Stateful $\mathcal O$ gets message from $\mathcal A$, returns output.

 \rightsquigarrow Inherently unable to rewind $\mathcal{A}.$

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Pseudo-Oracle

Stateful ${\mathcal O}$ gets message and view of ${\mathcal A},$ returns output.

Properties of pseudo-oracles

Black-box

 \mathcal{O} only uses \mathcal{A} black-box (instead of view (\mathcal{A})).

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k-robust pseudo-PPT ($\hat{=}$ rewinding simulatable)

For any k-round \mathcal{B} :

$$\exists \mathsf{PPT} \ \mathcal{R} \colon \quad \langle \mathcal{B}, \mathcal{A}^{\mathcal{O}} \rangle \stackrel{s}{\approx} \langle \mathcal{B}, \mathcal{R} \rangle$$

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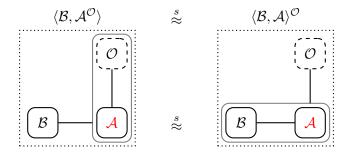
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k-robust composition-order invariant

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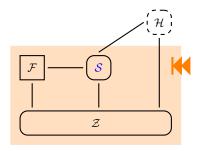
 $\langle \mathcal{B}, \mathcal{A}^{\mathcal{O}} \rangle \stackrel{s}{\approx} \langle \mathcal{B}, \mathcal{A} \rangle^{\mathcal{O}}$

Composition-order invariance (COI)



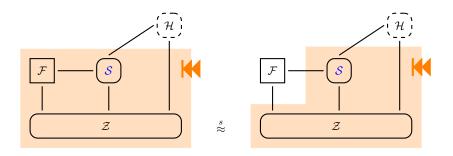
Why is LTR-UC meaningful at all?

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- What remains of the ideal guarantees of \mathcal{F} ?



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- LTR-UC angel $\mathcal H$ rewinds environment and ideal functionalities!
- What remains of the ideal guarantees of \mathcal{F} ?
- k-robust COI \implies meaningful for k-round functionalities.



Conclusion

- LTR-UC brings rewinding-based simulation to UC.
 - New possibilities: Com, ZK, one-sided-secure OT from CRS
 - Old impossibilities: (fully secure) OT from long-term revealing assumptions.
- Pseudo-Oracles \neq Oracles: Basic properties need non-trivial proofs.



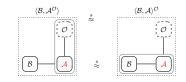
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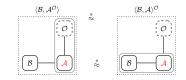
COI for our CCA-Com $\ensuremath{\mathcal{O}}$



Core difference to [CLP10; Goy+15]:

- [CLP10; Goy+15]: COI holds unconditionally due to bruteforce extraction.
- This work: COI via reduction to hardness assumption.

COI for our CCA-Com ${\cal O}$



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Proof idea (based on [PRS02] rewinding schedule):

- Given same randomness, **main thread** execution is **identical**, **unless** different committed values extracted (during look-ahead).
- Reduce different extracted values to binding break of COM.



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Image sources

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