

# Public-Coin, Complexity-Preserving, Succinct Arguments for NP from Collision Resistance

**Cody Freitag**<sup>1</sup>, Omer Paneth<sup>2</sup>, Rafael Pass<sup>2,3</sup>

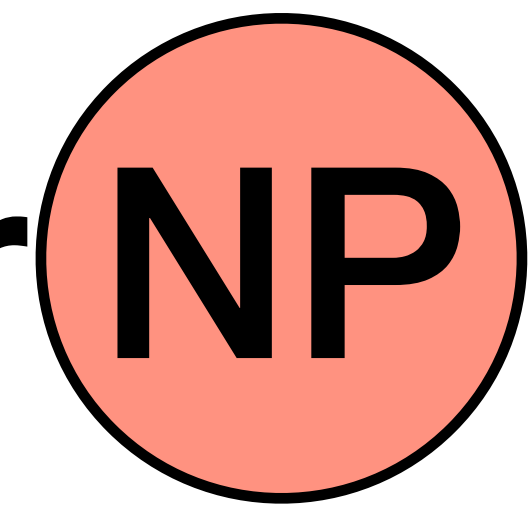
<sup>1</sup>Northeastern University

<sup>2</sup>Tel Aviv University

<sup>3</sup>Cornell Tech

# Cryptographic Proofs for NP

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# Cryptographic Proofs for **NP**

NP Prover

**P**

$x \in L?$

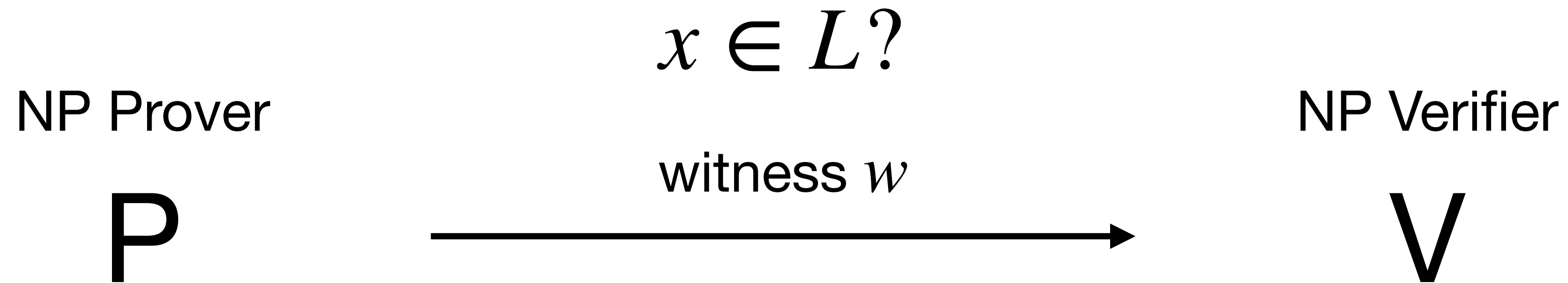
witness  $w$



NP Verifier

**V**

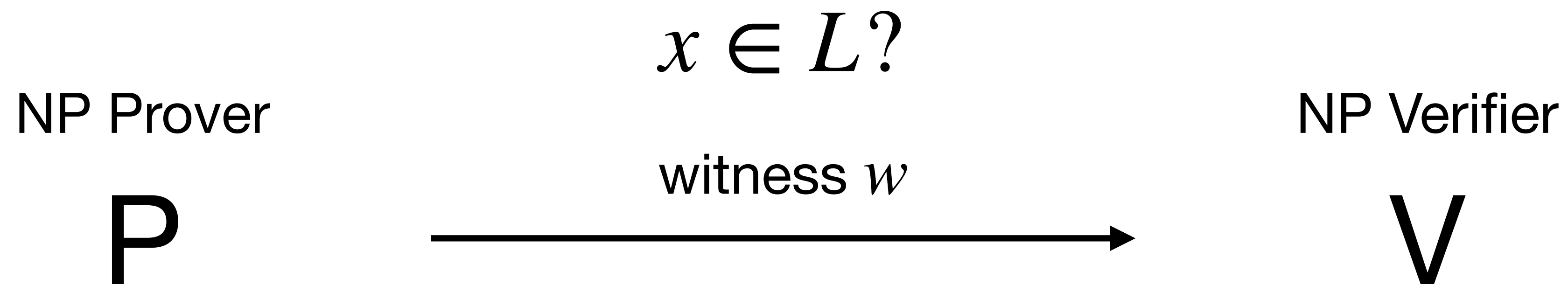
# Cryptographic Proofs for **NP**



**Completeness:** If  $x \in L$ ,  $\exists w$  s.t.  $V(x, w) = 1$ .

**Soundness:** If  $x \notin L$ ,  $\forall w$   $V(x, w) = 0$ .

# Cryptographic Proofs for NP

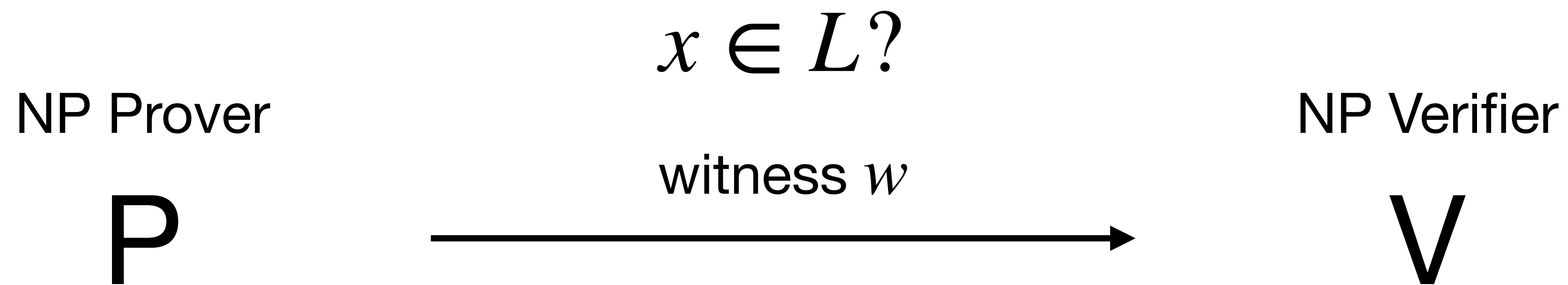


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Equivalent to  
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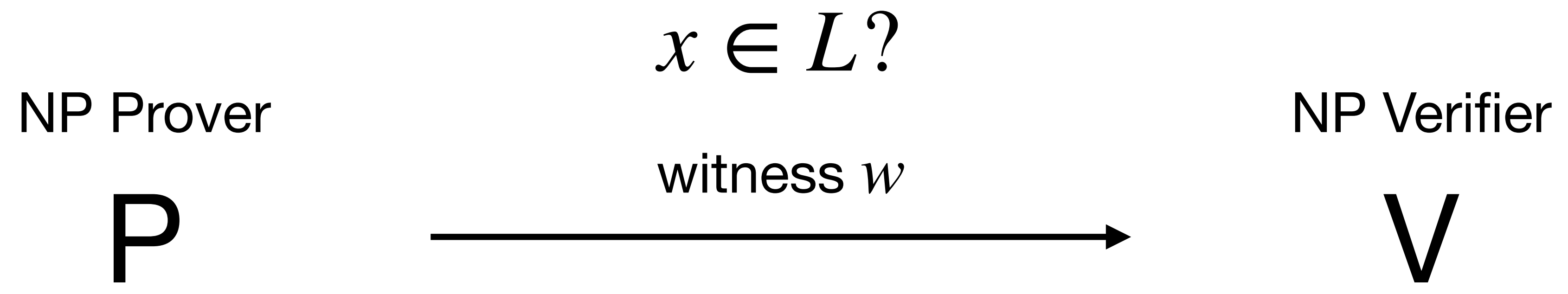
**Soundness:** If  $x \notin L$ ,  $\forall w V(x, w) = 0$ .

Equivalent to  
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## Assumptions:

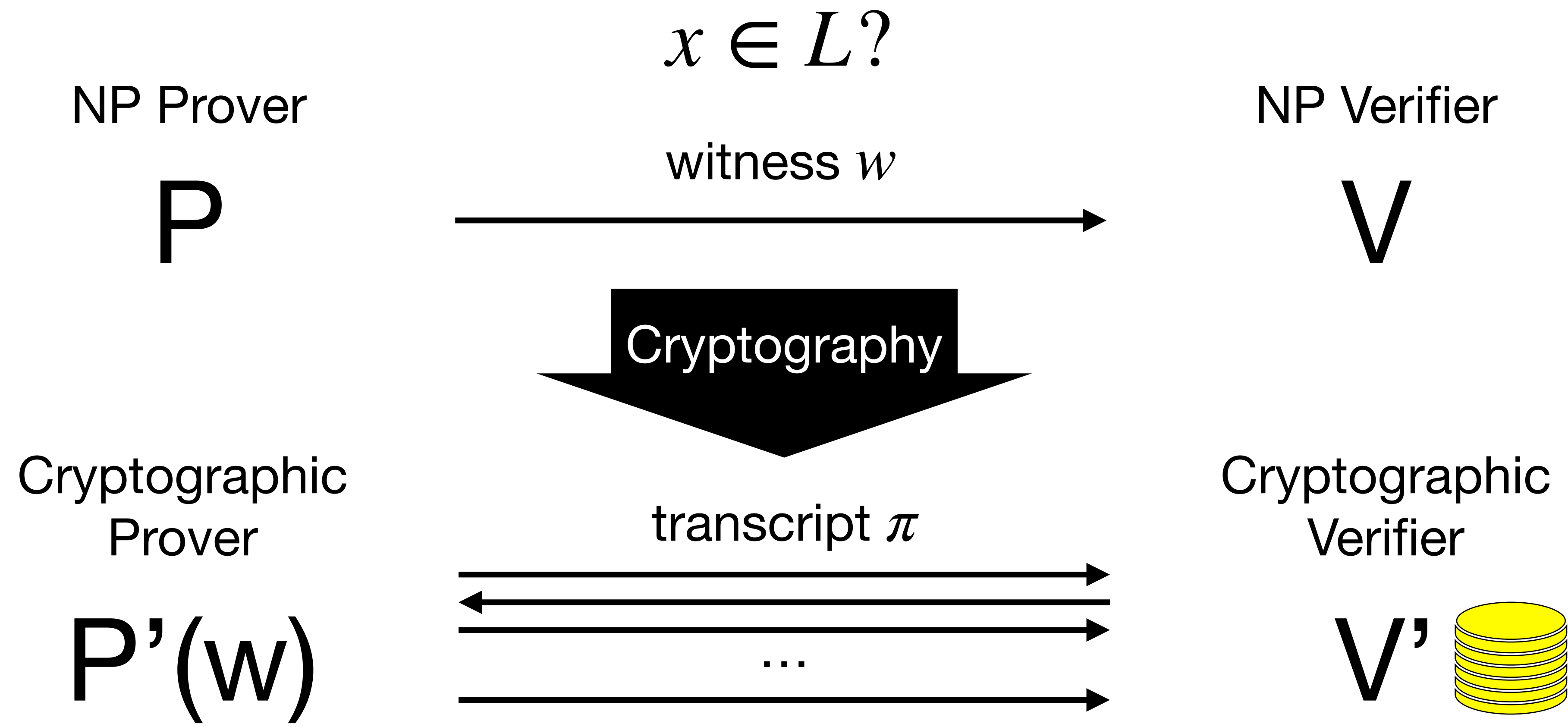
- $V(x, w)$  runs in **time  $T$**  and **space  $S$**   
as a **RAM program**

# Cryptographic Proofs for NP

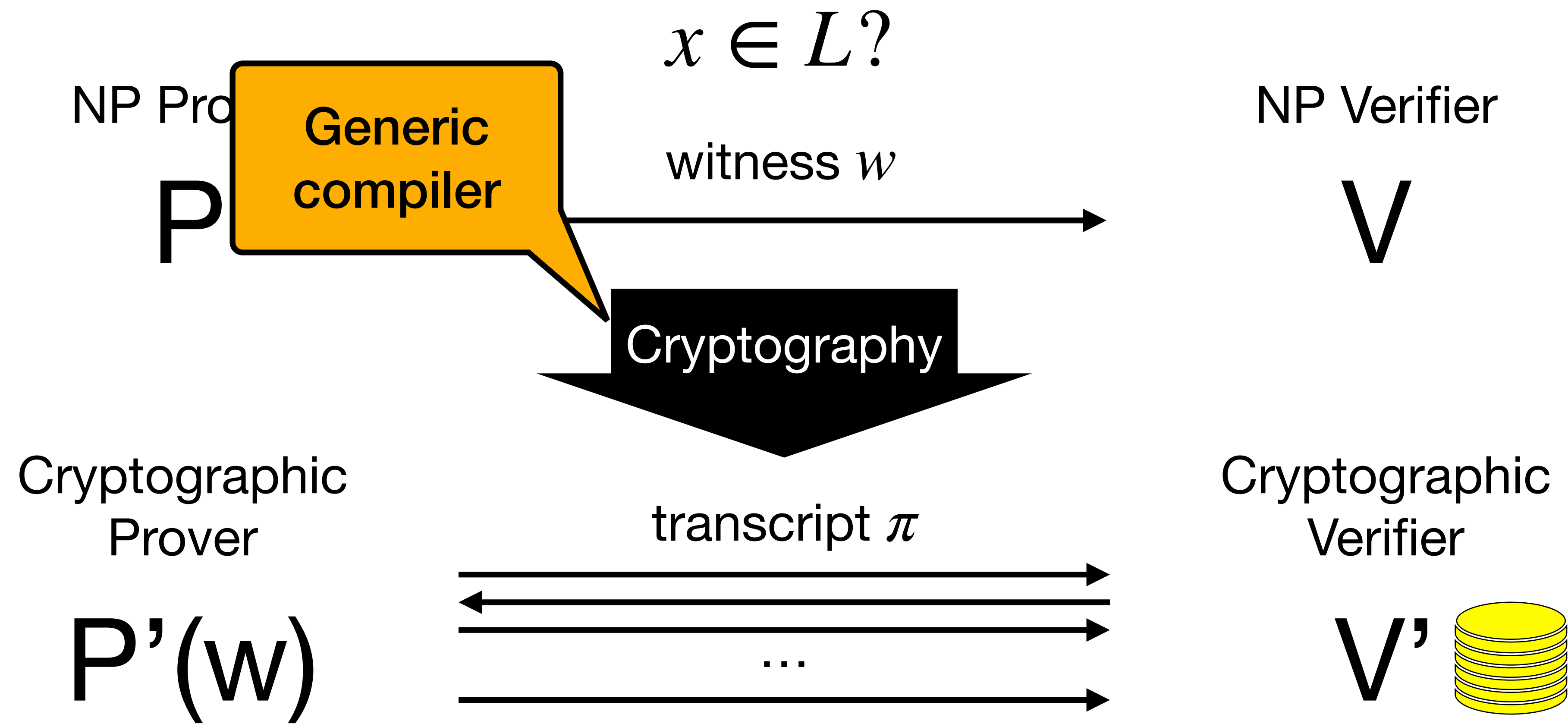




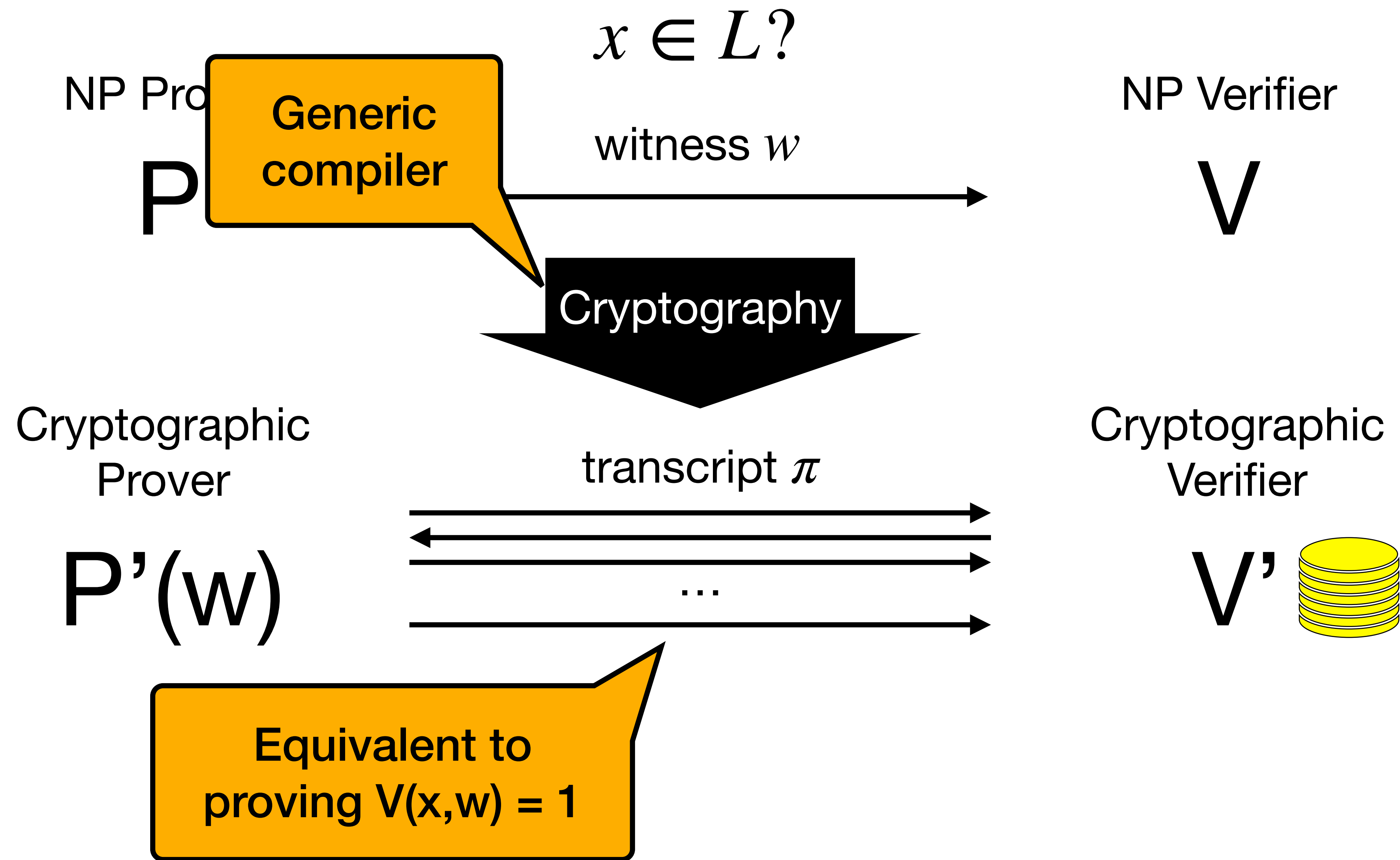
# Cryptographic Proofs for NP



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# Cryptographic Proofs for NP

Cryptographic  
Prover

$P'(w)$

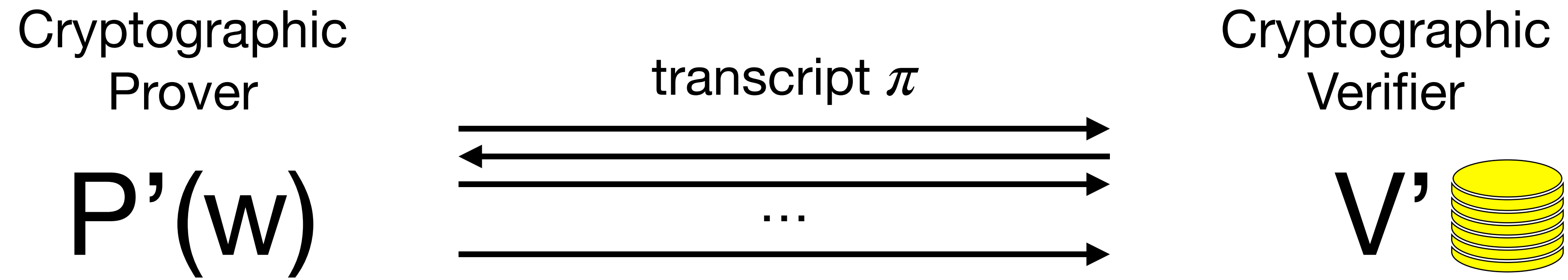
transcript  $\pi$



Cryptographic  
Verifier

$V'$  

# Cryptographic Proofs for NP



## Computational Soundness (Argument):

If  $x \notin L$ , no **poly-time** prover  $P^*$  can convince  $V'$  w.h.p.

## Argument of Knowledge:

If poly-time  $P^*$  convinces  $V'$ , it **“knows”** a witness.

## Public-coin (optional):

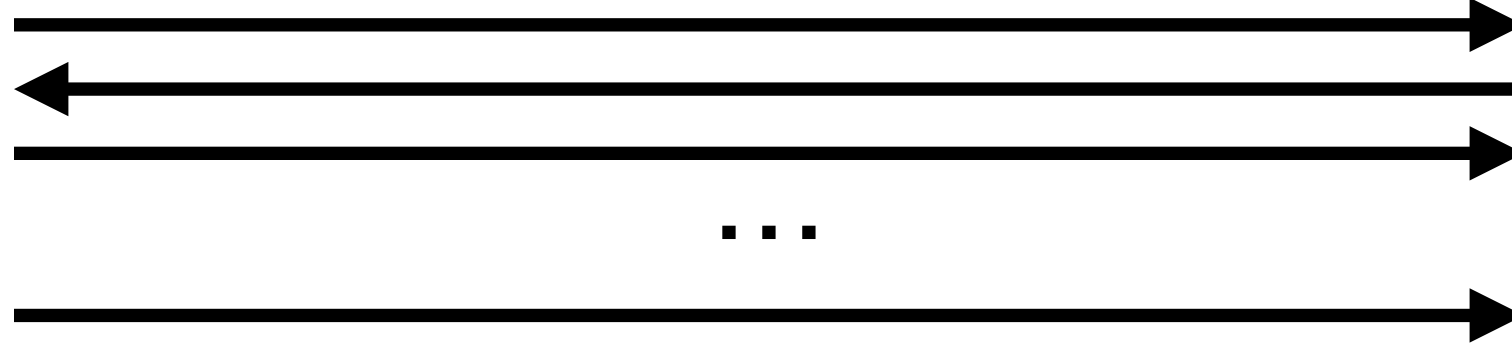
$V'$  sends random coins and has **no private state**.

# Cryptographic Proofs for NP

Cryptographic  
Prover

$P'(w)$

transcript  $\pi$



Cryptographic  
Verifier

$V'$  

**Goals:**

**Privacy**

**Efficiency**

# Cryptographic Proofs for NP

Cryptographic Prover

$P'(w)$

transcript  $\pi$



Cryptographic Verifier

$V'$  

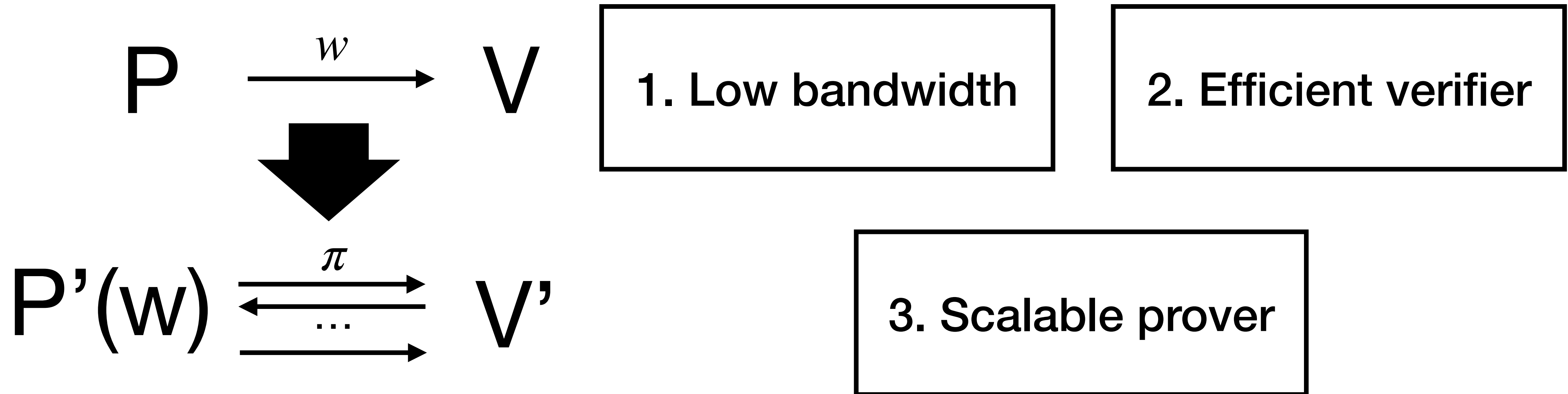
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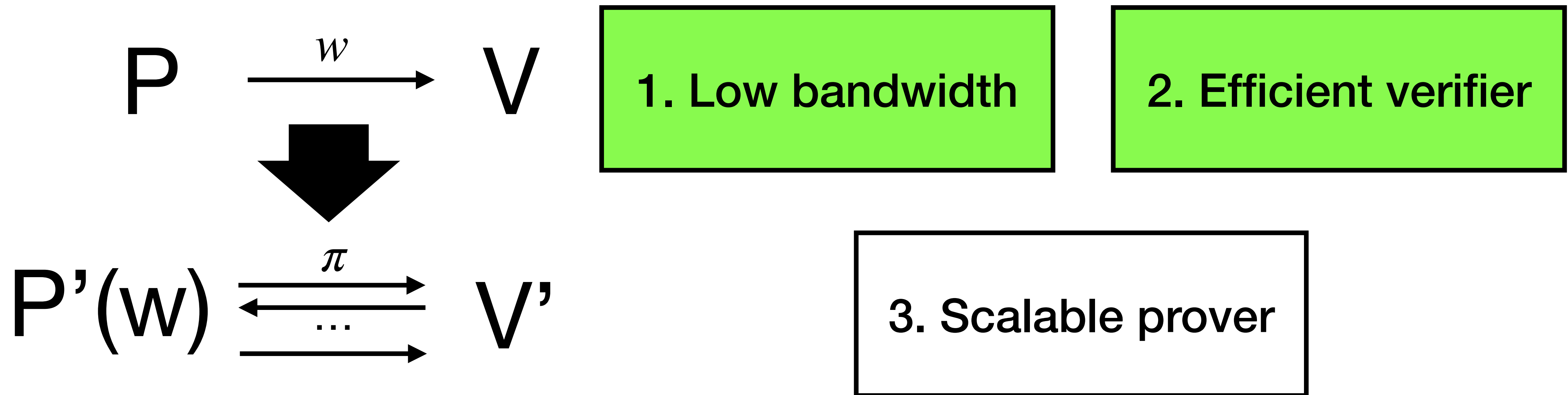
**Focus of this talk**

# Efficiency Goals





# Efficiency Goals



## Succinct Argument:

- $|\pi| = \text{polylog}(T)$
- $\text{Time}(V') = \text{poly}(|x|, \log T)$

# Efficiency Goals



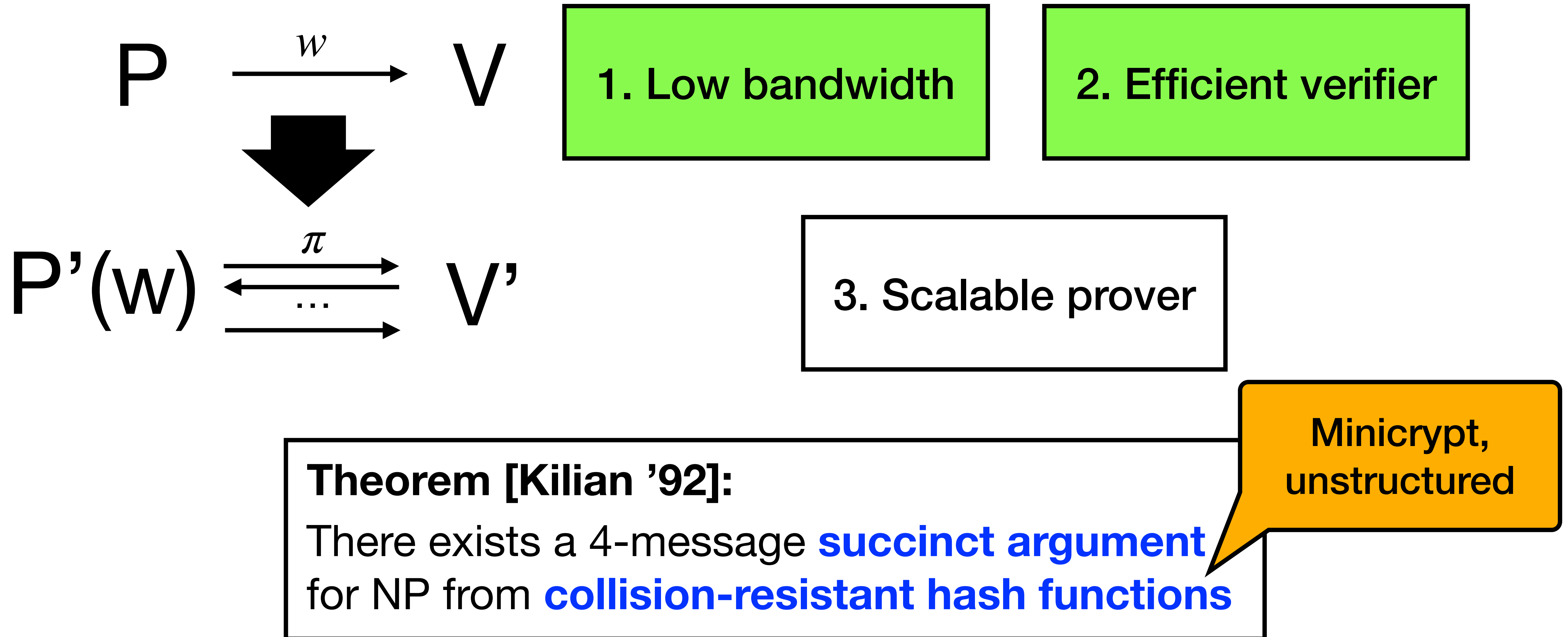
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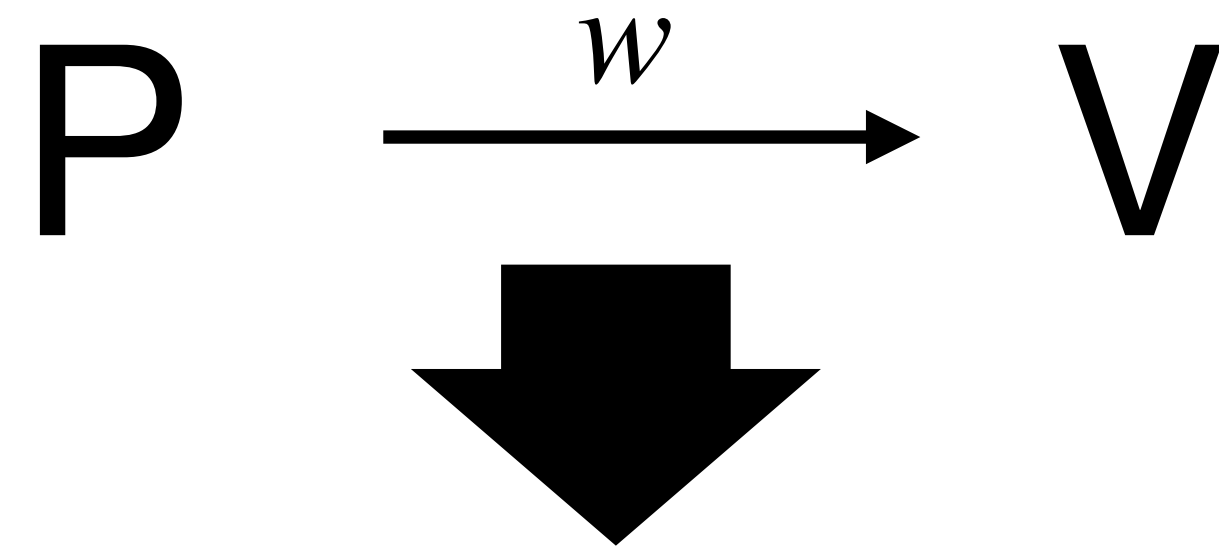
**Theorem [Kilian '92]:**

There exists a 4-message **succinct argument** for NP from **collision-resistant hash functions**

# Efficiency Goals



# Efficiency Goals



1. Low bandwidth

2. Efficient verifier

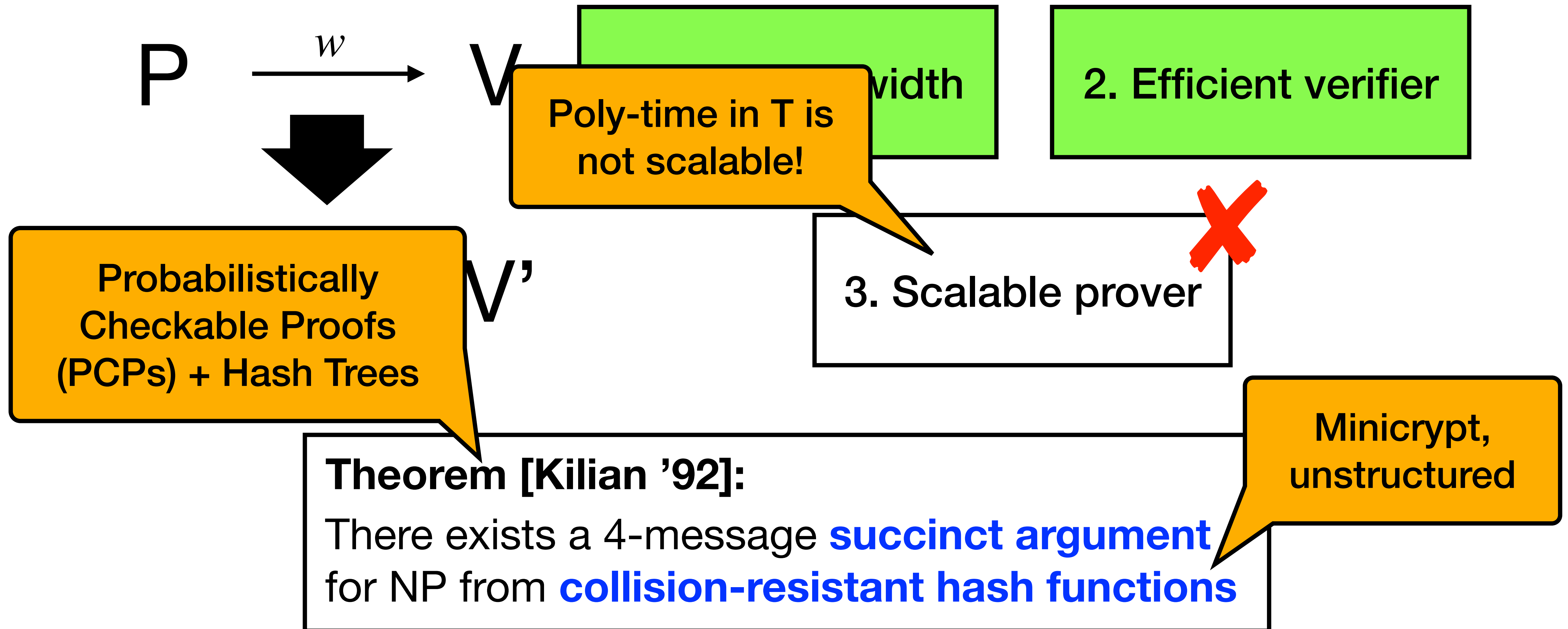
3. Scalable prover

Probabilistically Checkable Proofs (PCPs) + Hash Trees

**Theorem [Kilian '92]:**  
There exists a 4-message **succinct argument** for NP from **collision-resistant hash functions**

Minicrypt, unstructured

# Efficiency Goals



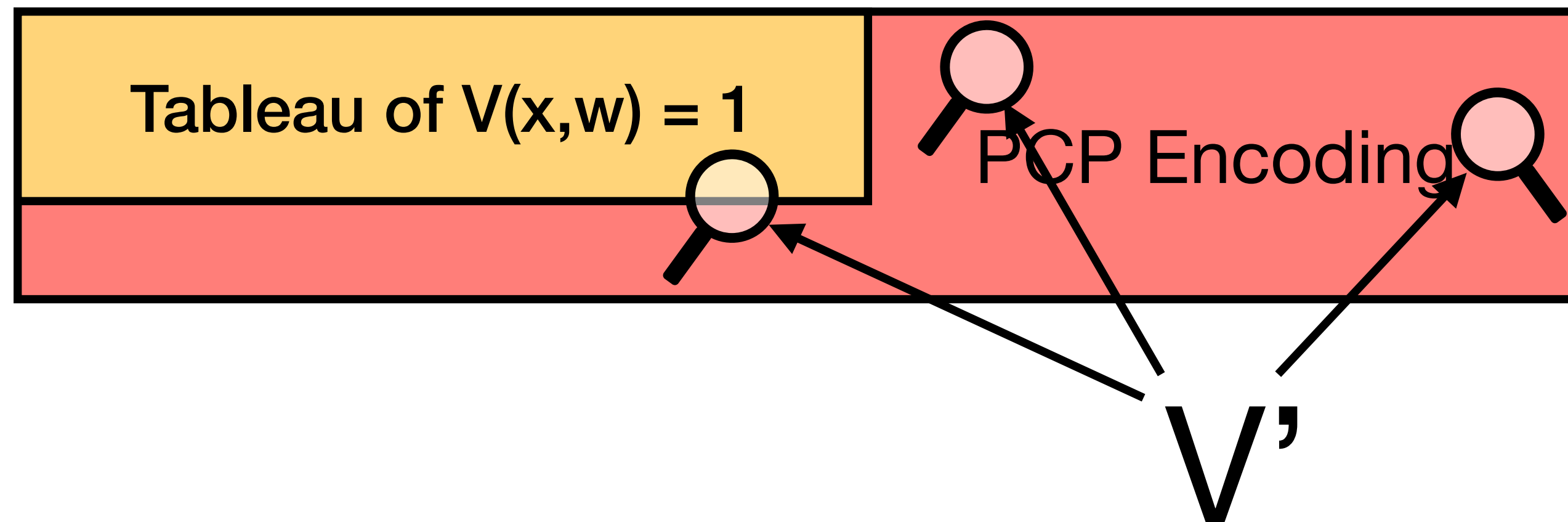
# The Quest for a Scalable Prover

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## More Efficient PCPs

[BFLS91],[BGHSV05],  
[D07],[BS08],[M09],  
[BCGT13],  
[BKKMS16],...

$P'(w)$  →



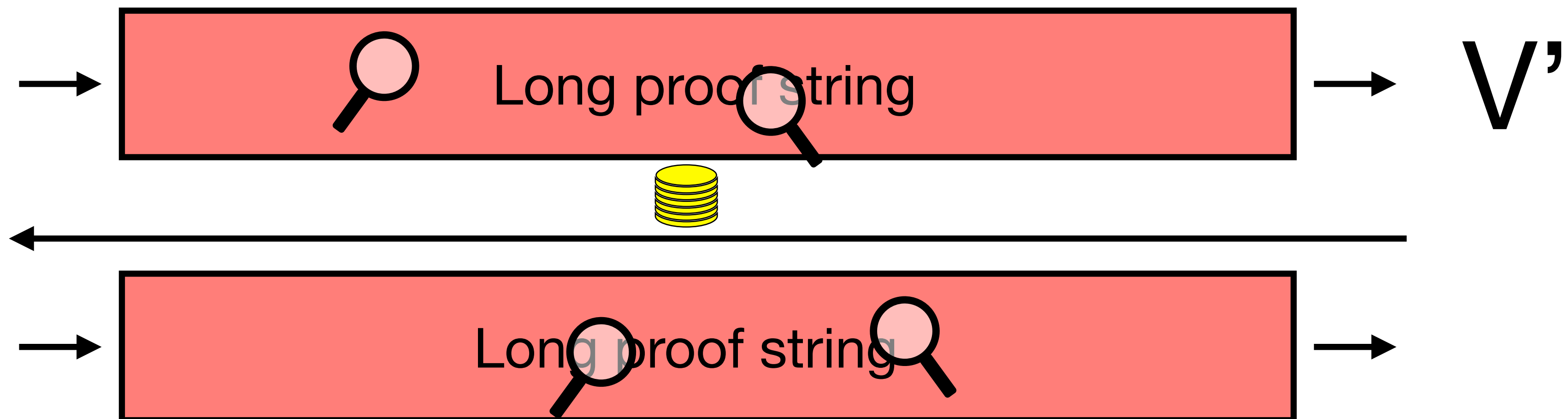


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**More Efficient PCPs**  
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**Interactive Oracle Proofs (IOPs)**  
[BCS16],[RRR16],  
[BCGRS17],[RR20],...

$P'(w)$



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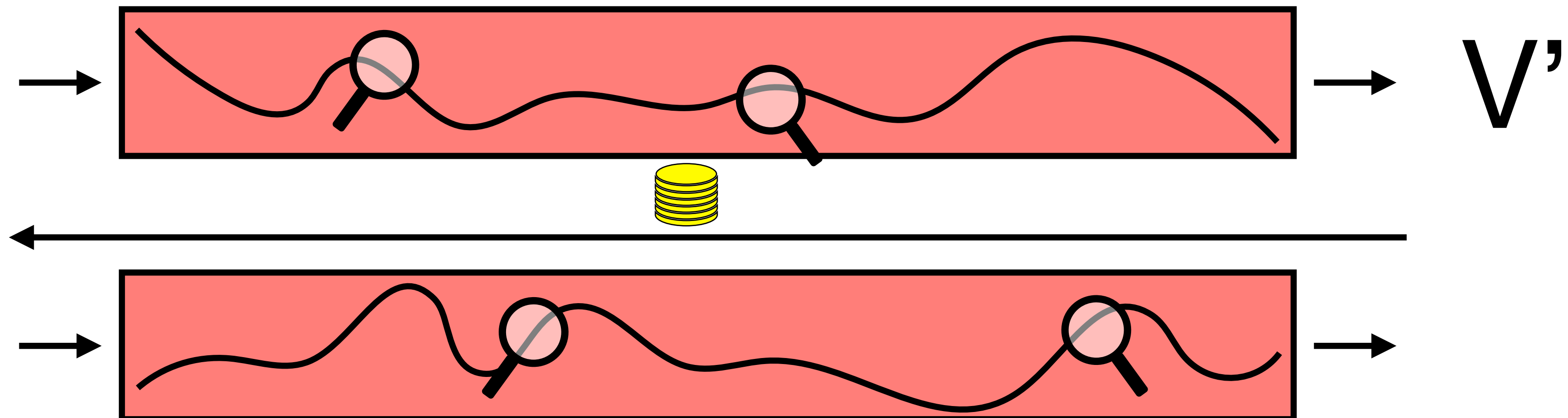
## Interactive Oracle Proofs (IOPs)

[BCS16],[RRR16],  
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## Structured Oracle Proofs

[IKO07],[KZG10],[BCIOP13],[GGPR13],  
[BCCGP16],[BBBPWM18],[BBHR18],  
[BBCGI19],[BBHR19],[MBKM19],  
[GWC19],[S20],[CHMMVW20],[BFS20],  
[BHRRS20],[BHRRS21],[SZ22],...

$P'(w)$



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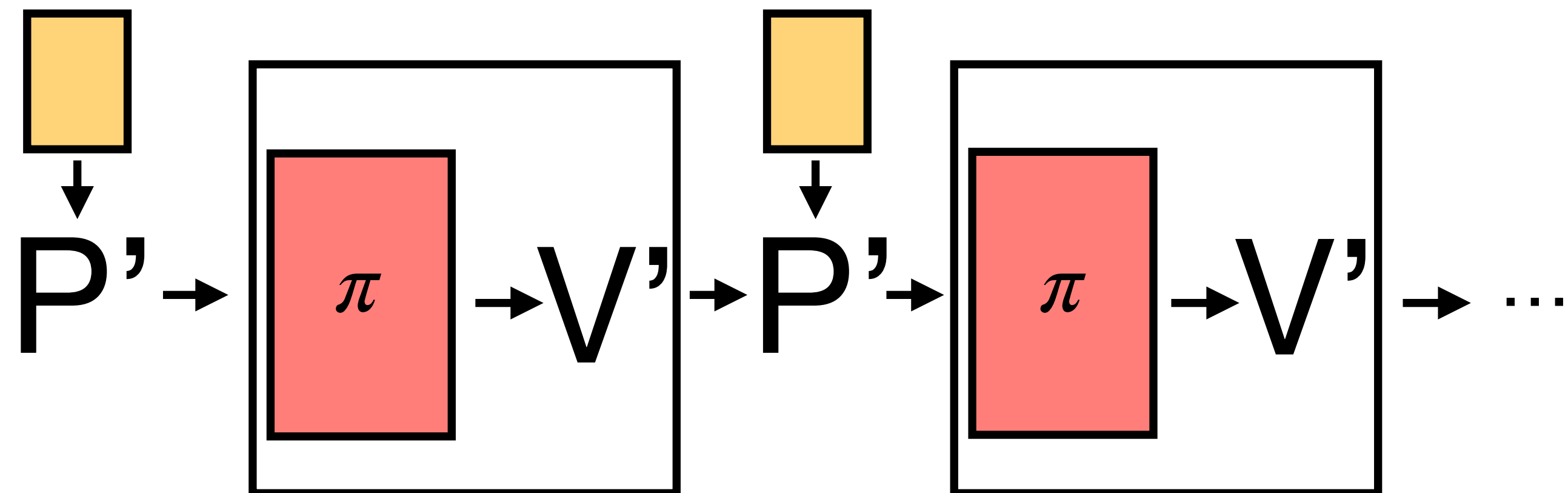
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## Recursive Composition

[V08],[BCCT12],...



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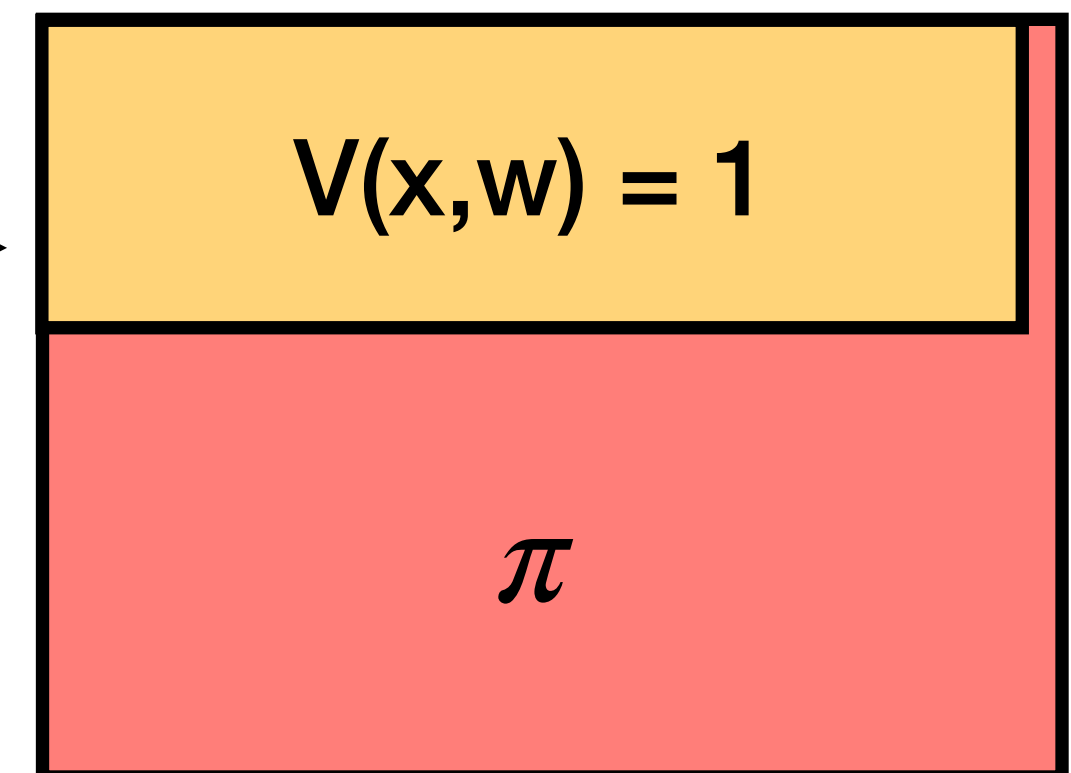
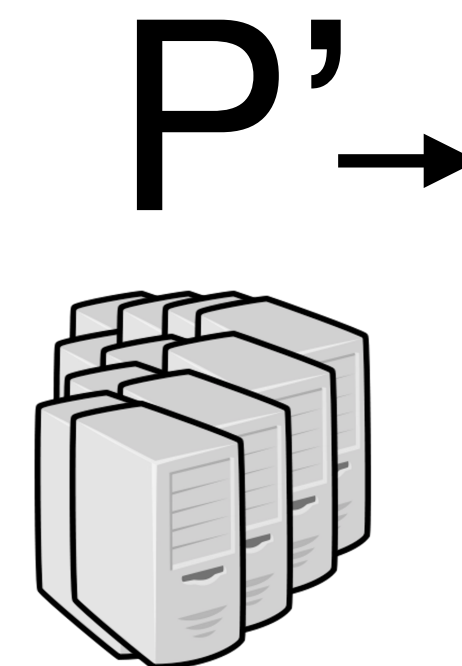
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[EFKP20],[FPS22]



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[V08],[BCCT12],...

## Parallelizable Provers

[E**F**KP20],[**F**PS22]

## Space-Efficient/ Complexity-Preserving

[BC12],[BCCT12],[BHRRS20],  
[BHRRS21],[BBHV22]

# Blueprint for Kilian-Based Arguments

$$P'(w) \begin{array}{c} \xrightarrow{\pi} \\ \xleftarrow{\quad} \\ \dots \\ \xrightarrow{\quad} \end{array} V'$$

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1. Write down  
computation tableau

Tableau of  $V(x,w) = 1$

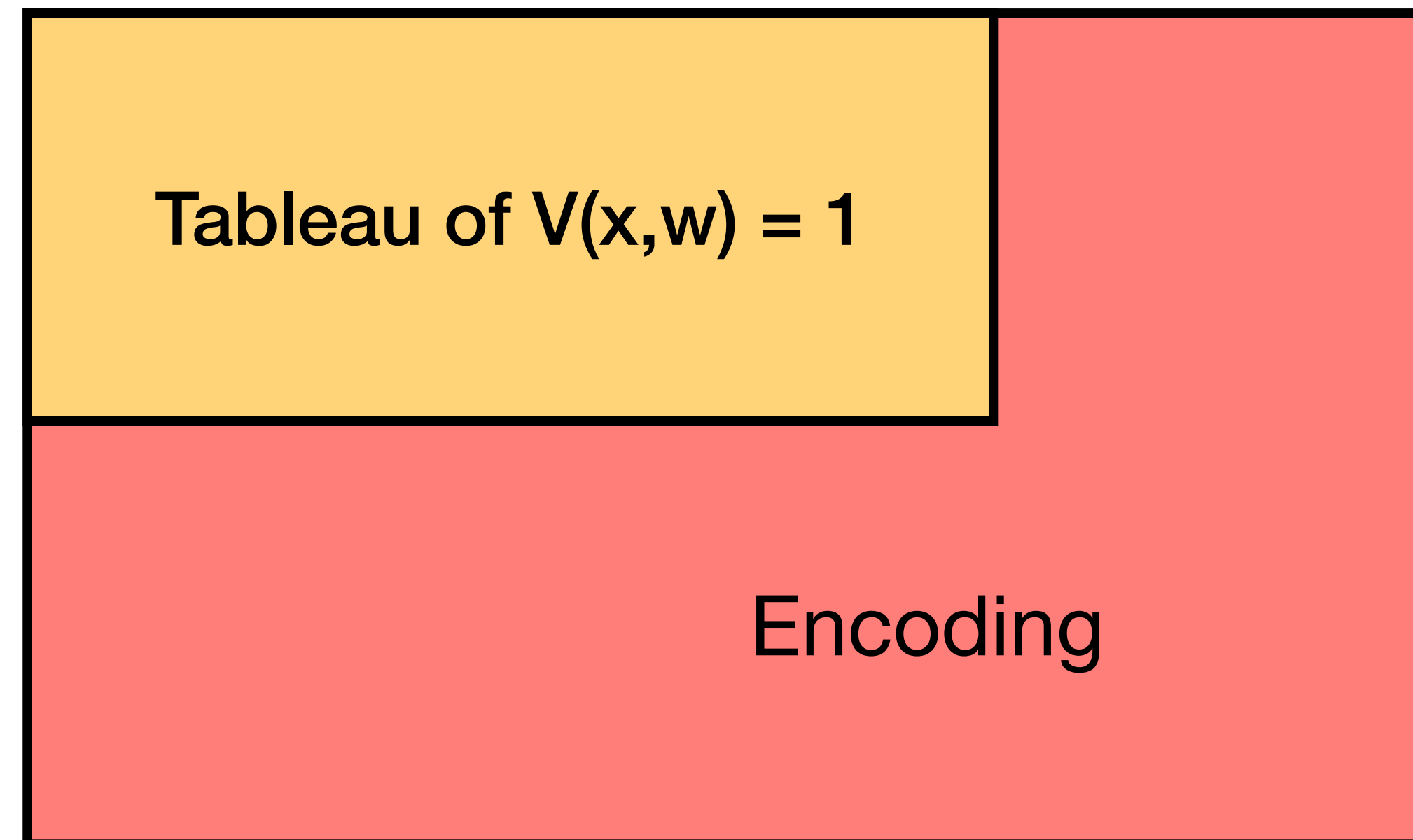


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2. Encode tableau



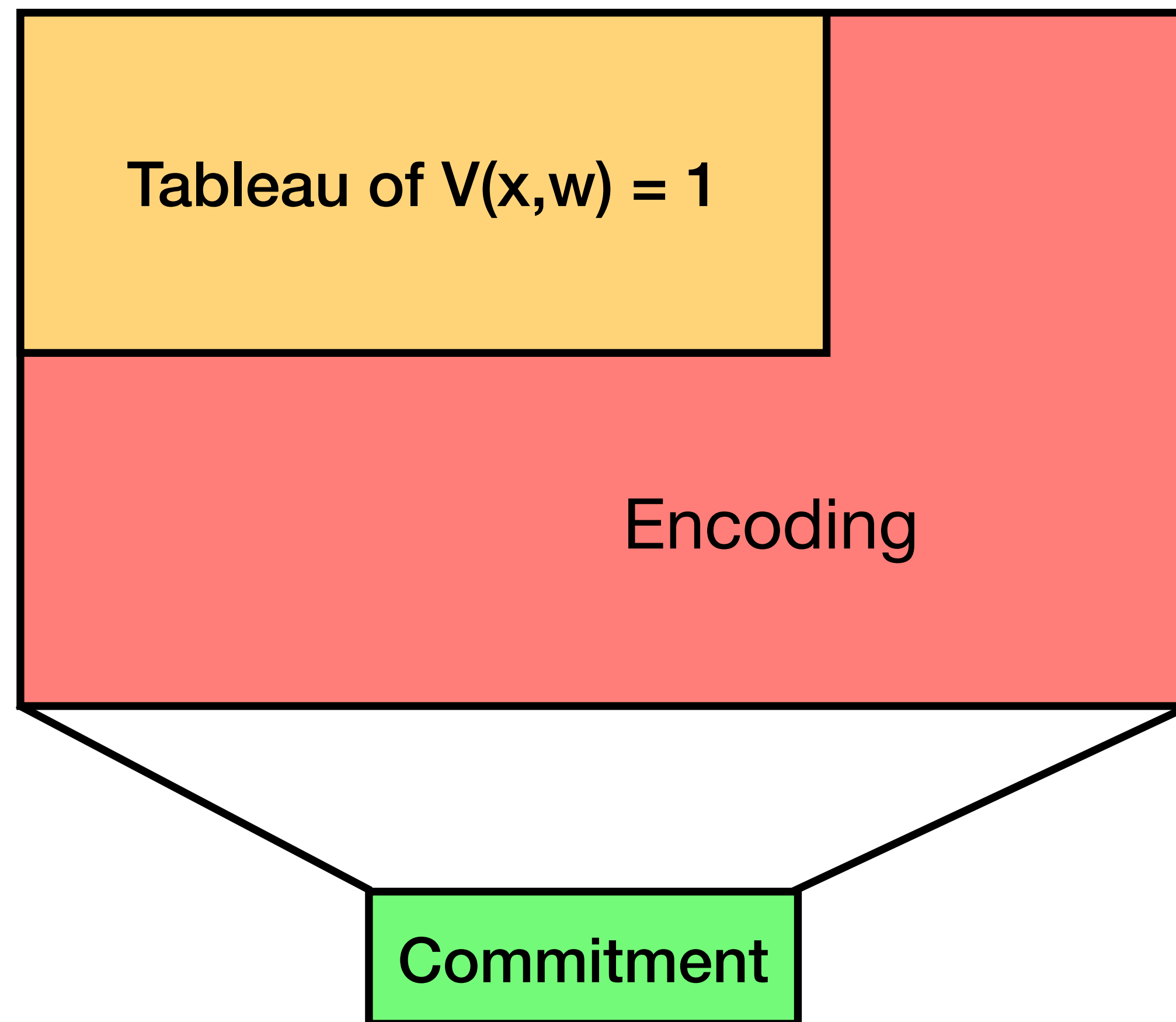
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3. Cryptographically commit to encoding



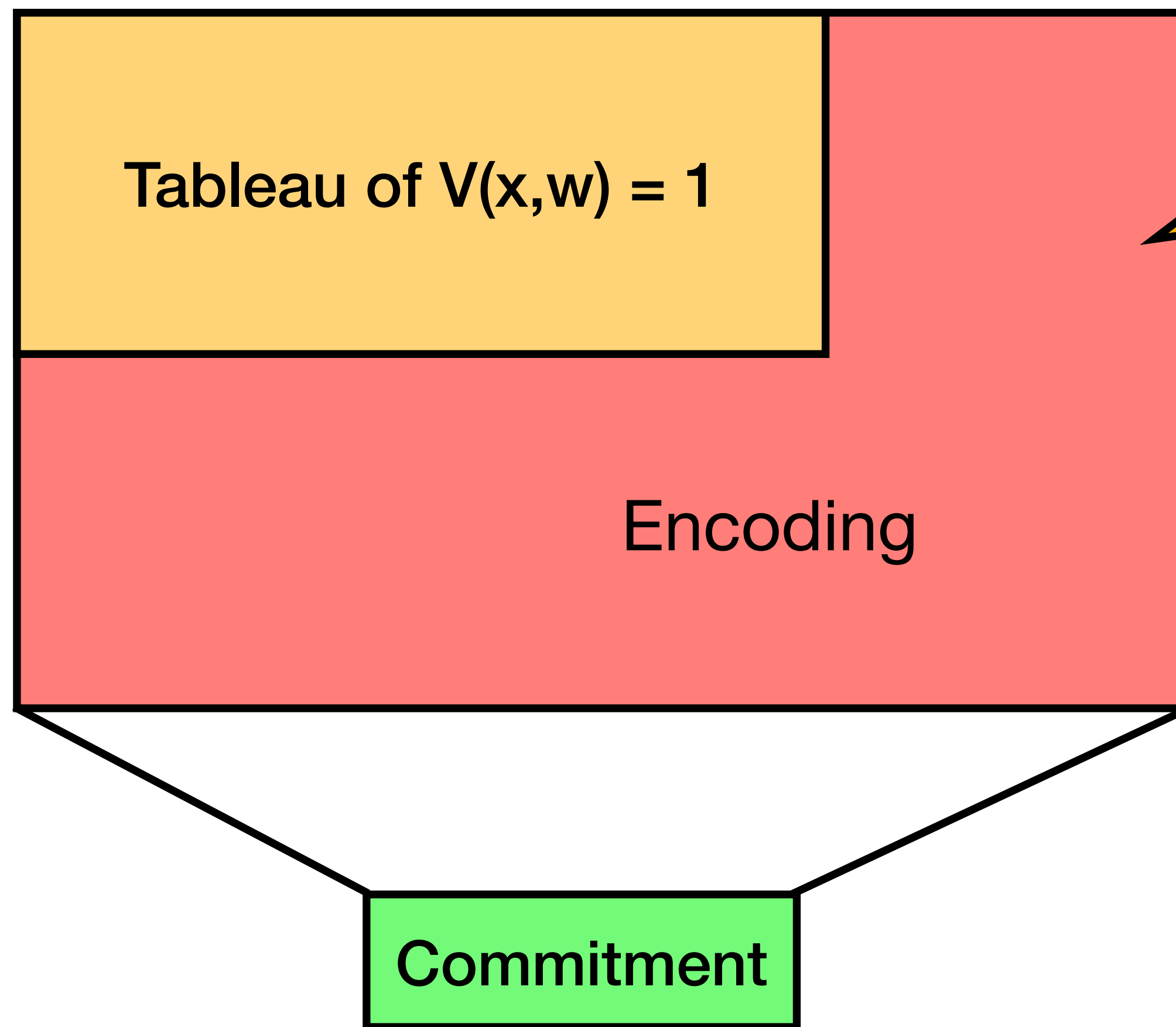
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Requires space  $\geq T!$

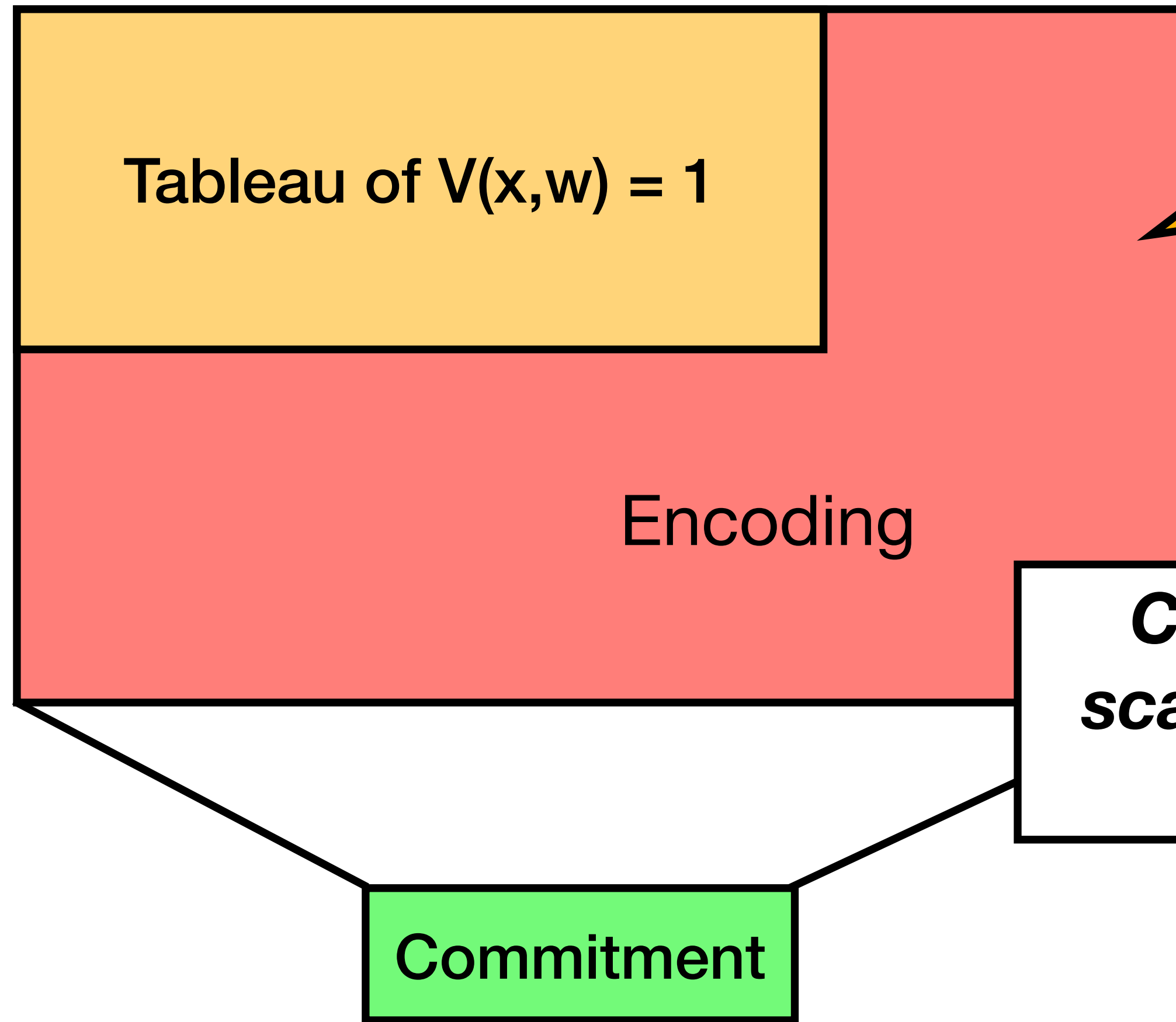
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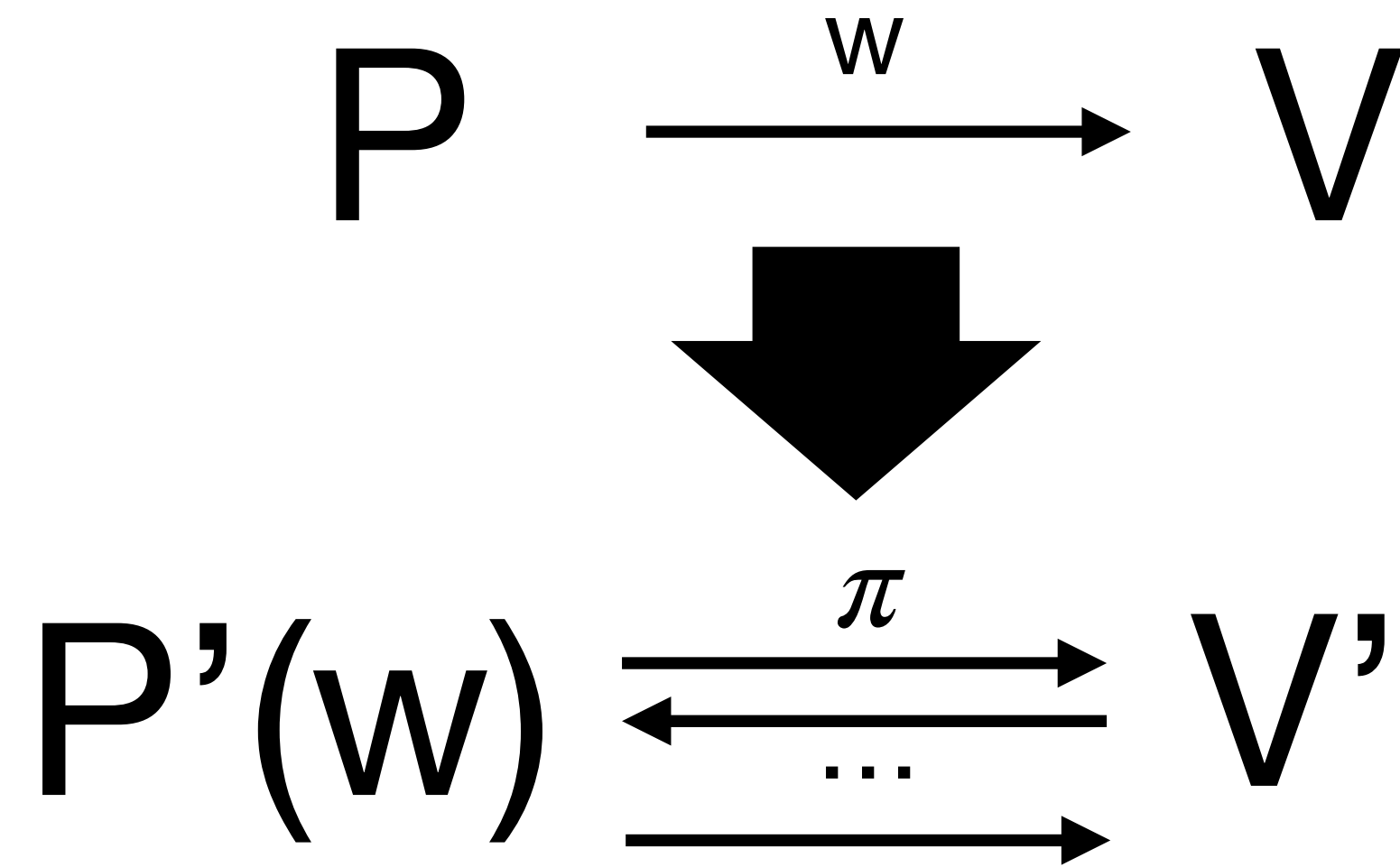
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Requires space  $\geq T!$

Can space of  $P'$  scale with **space  $S$**  of  $V(x,w)$ ?

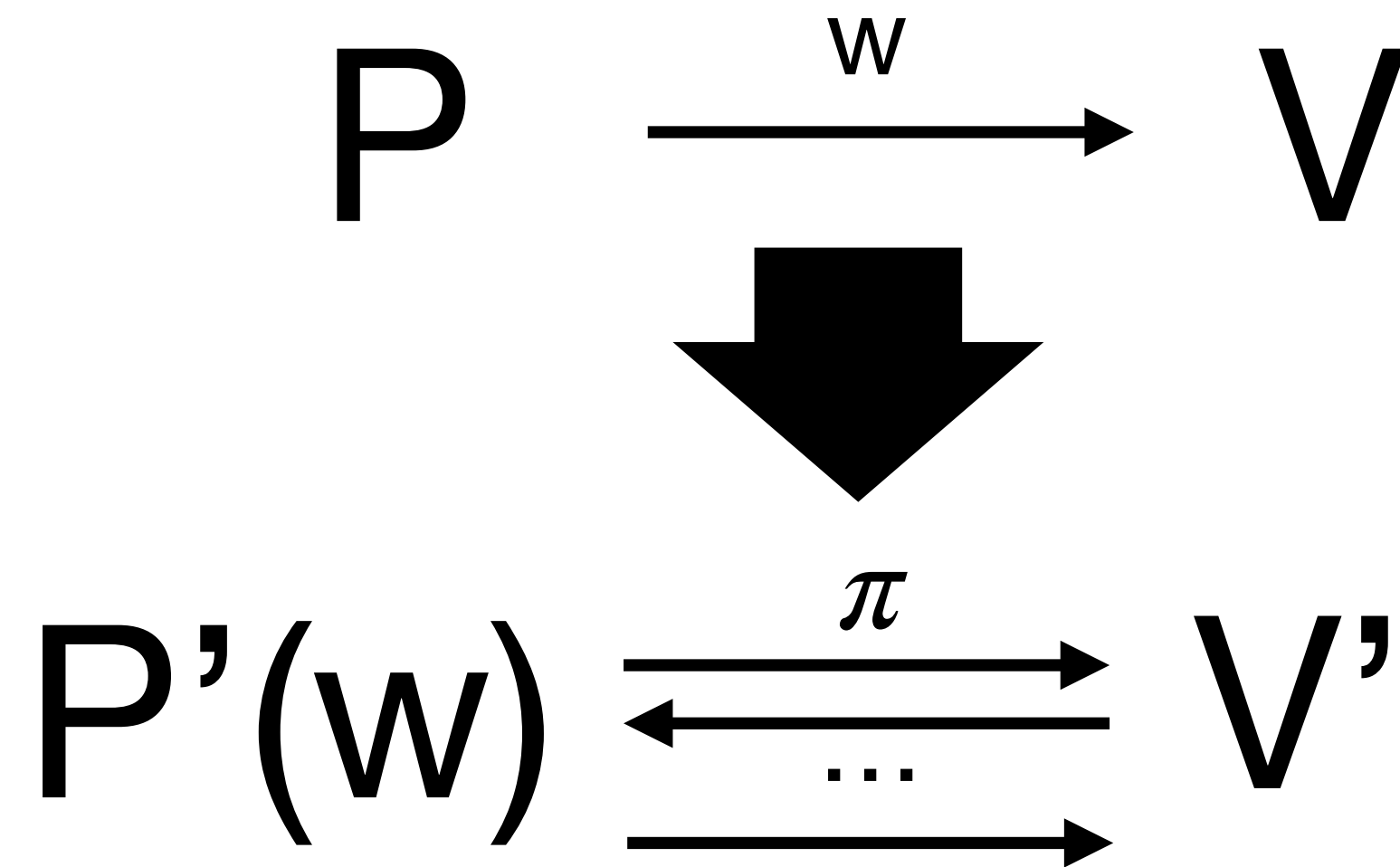
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**Complexity-Preserving:**  
-  $P'(w)$  runs in time  $\tilde{O}(T)$   
and space  $\tilde{O}(S)$

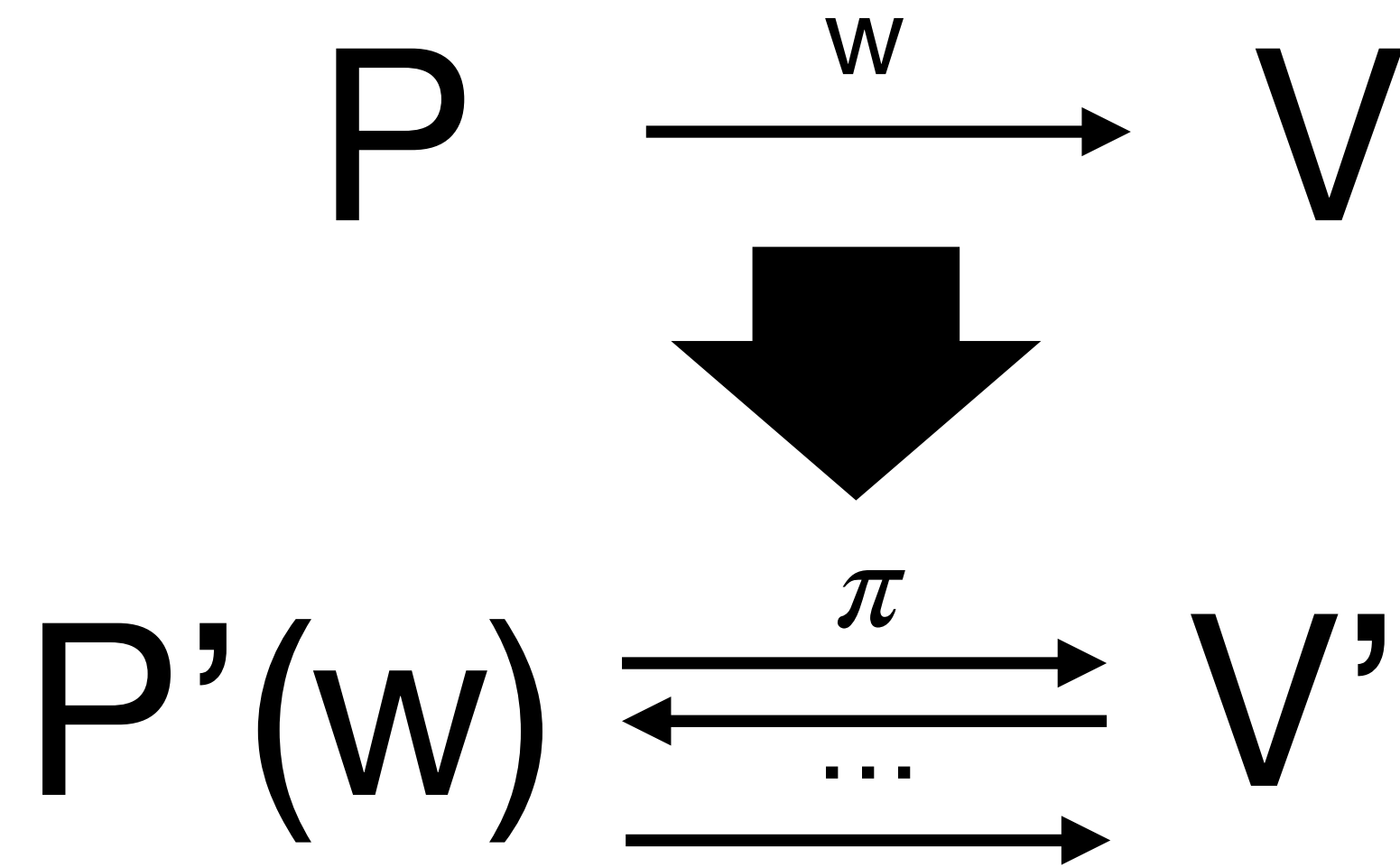
# Complexity-Preserving An

Best possible up to polylog factors!



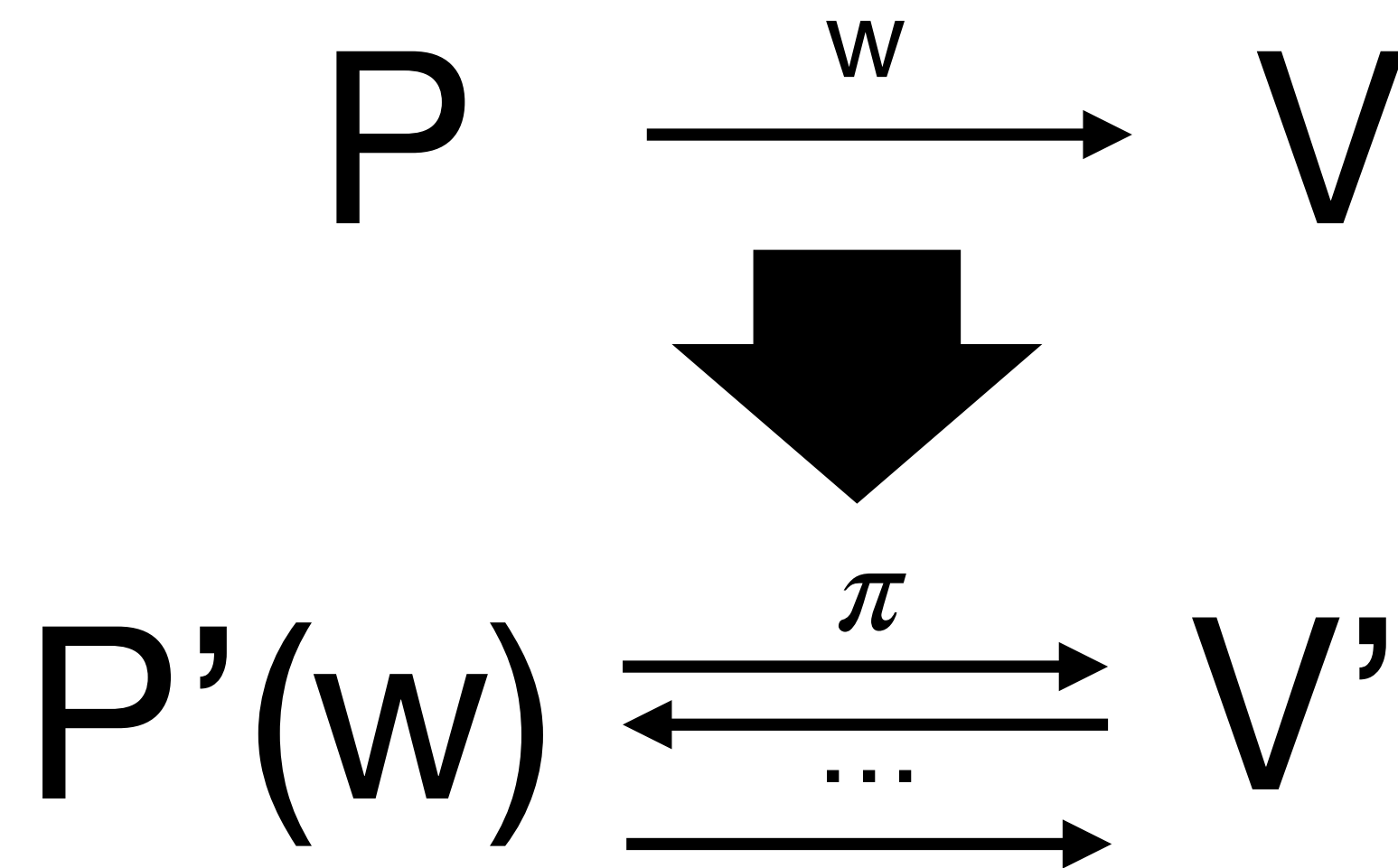
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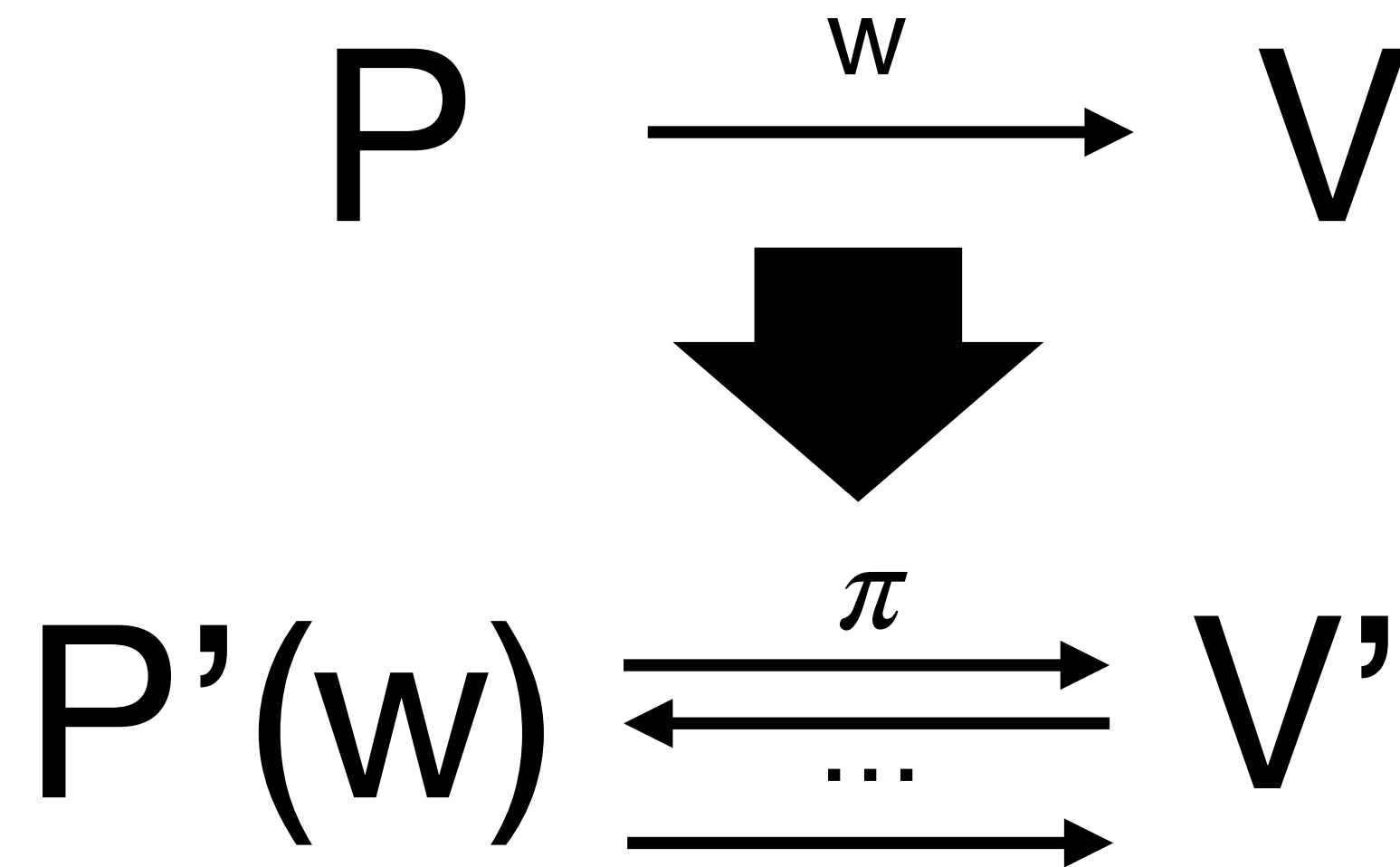


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**Theorem [BitanskyChiesa '12]:**  
There exists a 4-message **complexity-preserving** succinct argument for NP from **fully homomorphic encryption**



# Complexity-Preserving Arguments



**Complexity-Preserving:**  
-  $P'(w)$  runs in time  $\tilde{O}(T)$   
and space  $\tilde{O}(S)$

**Open Questions:**

Public  
verifiability?

Minicrypt?

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# Towards resolving these questions

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**Publicly-verifiable,  
Complexity-Preserving,  
Succinct Arguments**

\*not fully succinct  
\*\*only privately verifiable

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**Publicly-verifiable,  
Complexity-Preserving,  
Succinct Arguments**

**Non-standard  
assumptions**

BHRRS21

BHRRS20\*

BCCT12

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# Towards resolving these questions

**Publicly-verifiable,  
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**Non-standard  
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Class Groups

BHRRS21

ROM + DLOG

BHRRS20\*

SNARKs

BCCT12

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**Non-interactive  
for P**

KLW23

DGKV22

PP22

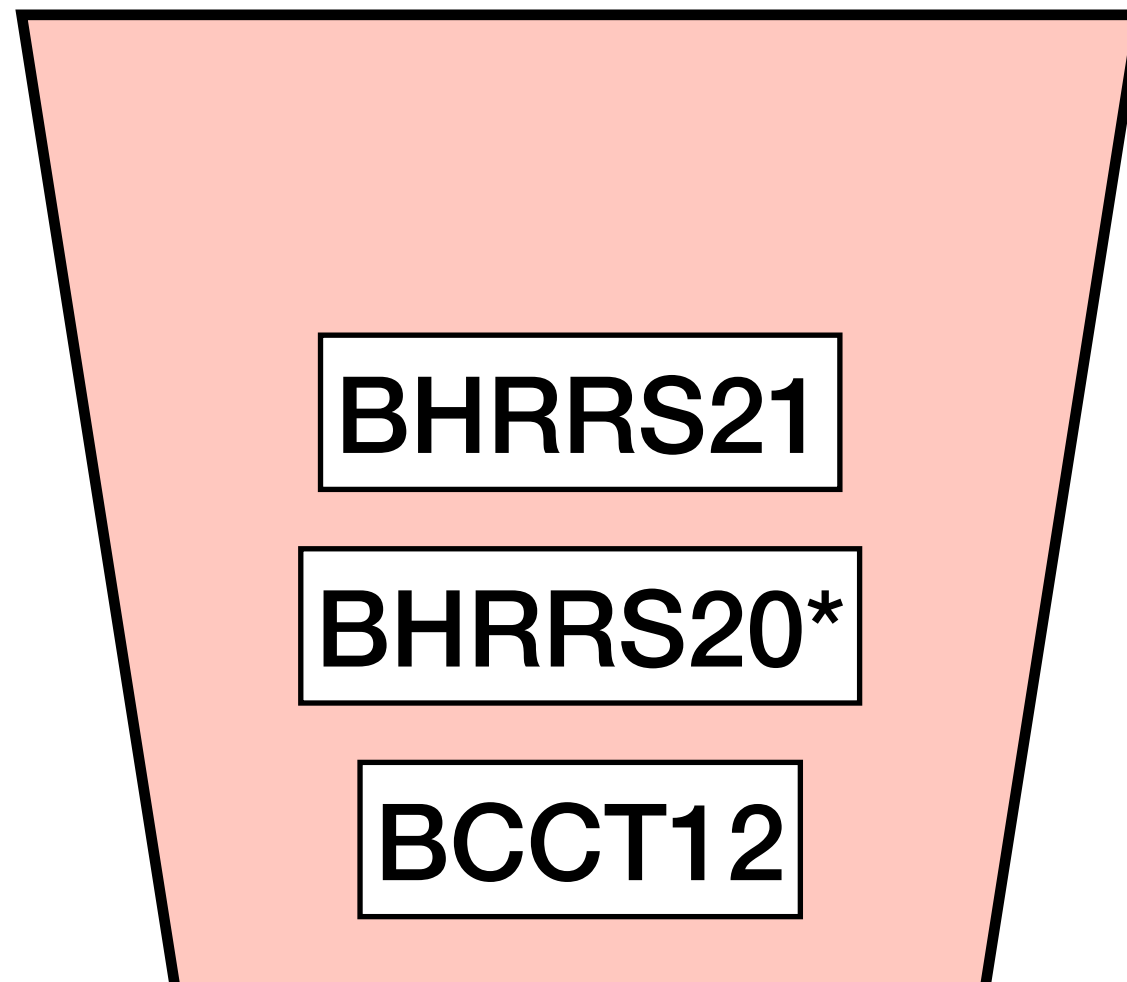
HR18\*\*

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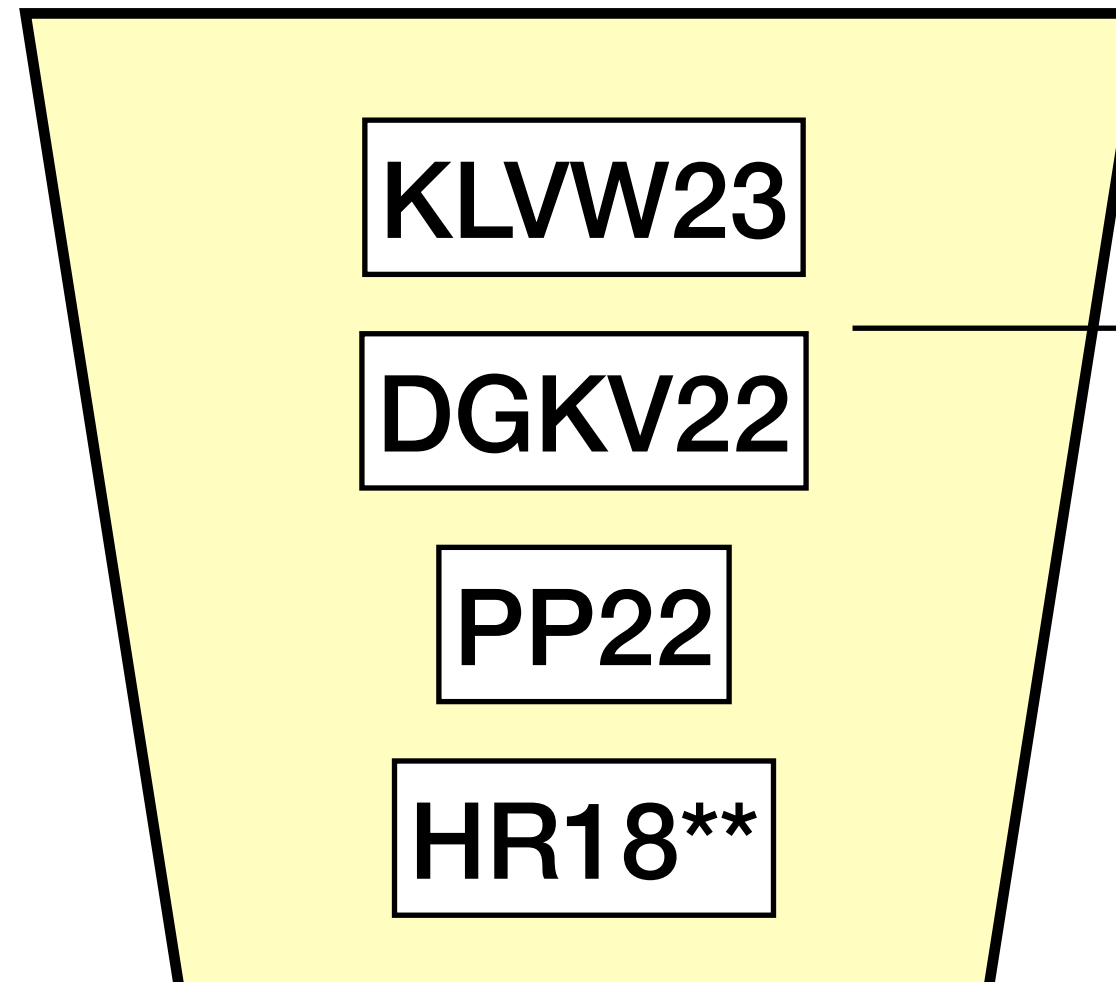
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BARGs for NP:

- LWE
- Bilinear Maps
- Sub-Exp DDH

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**Still all  
Cryptomania  
assumptions!**

**BARGs for NP:**

- LWE
- Bilinear Maps
- Sub-Exp DDH

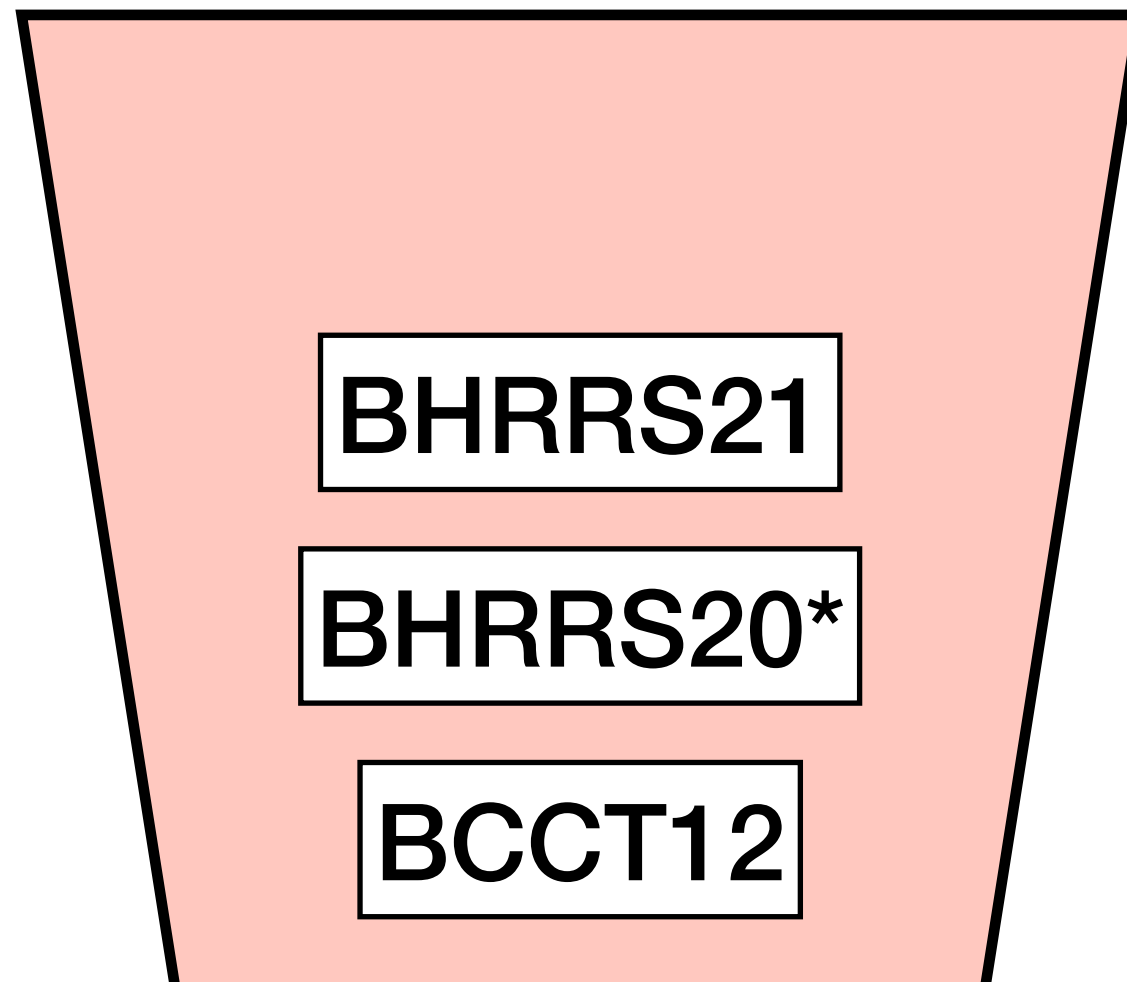
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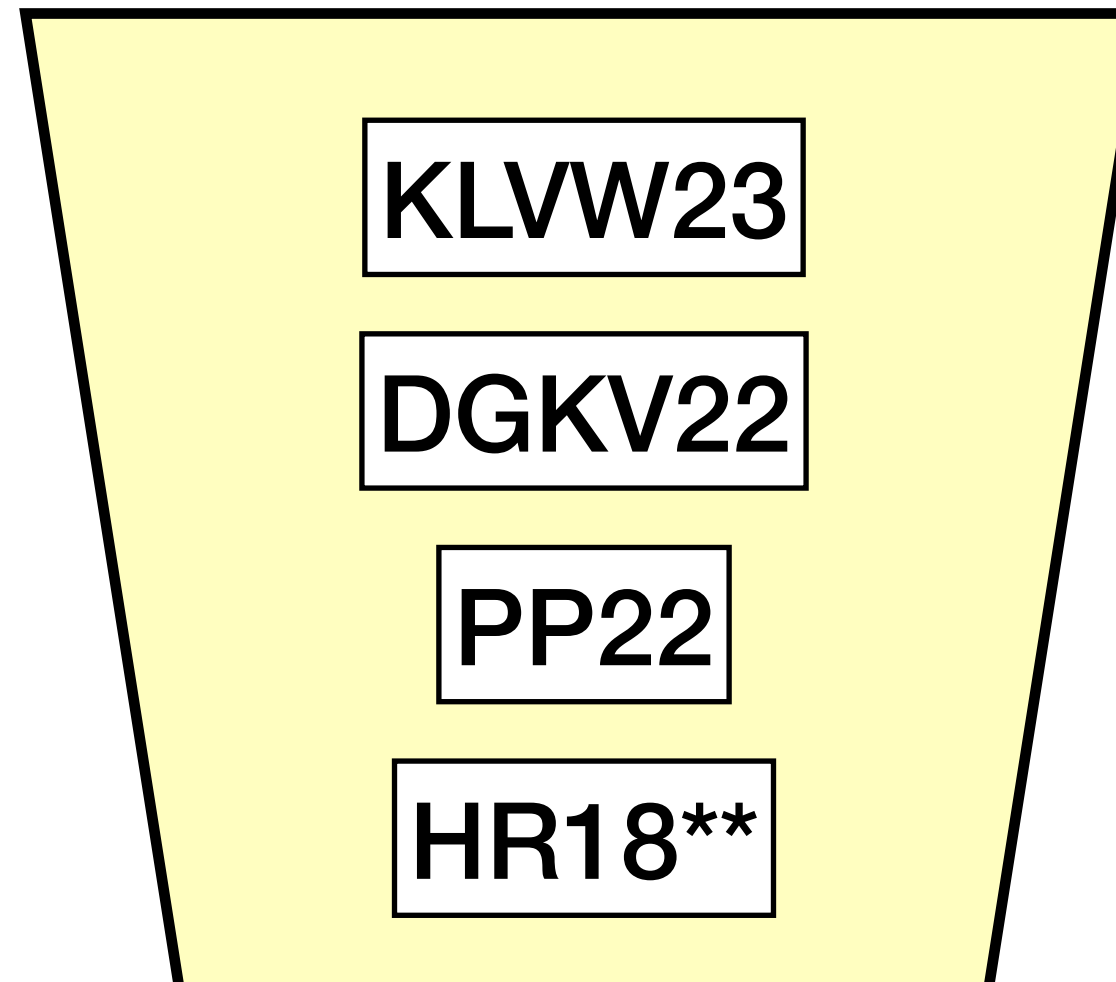
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**Black-Box Use  
of CRH**

BBHV22\*

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Black-Box Use  
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Comm:  $\tilde{O}(T/S)$

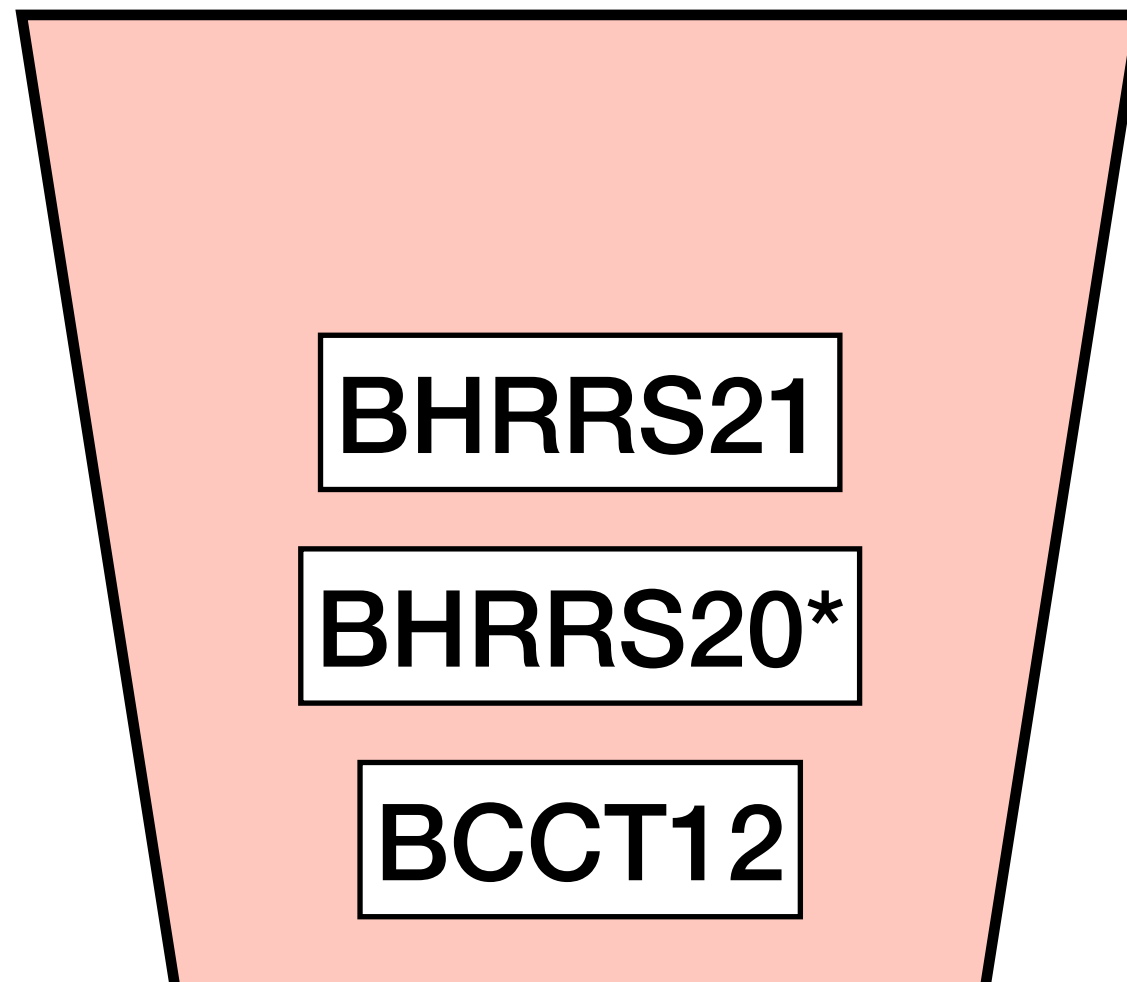
Ver:  $\tilde{O}(T/S + S)$

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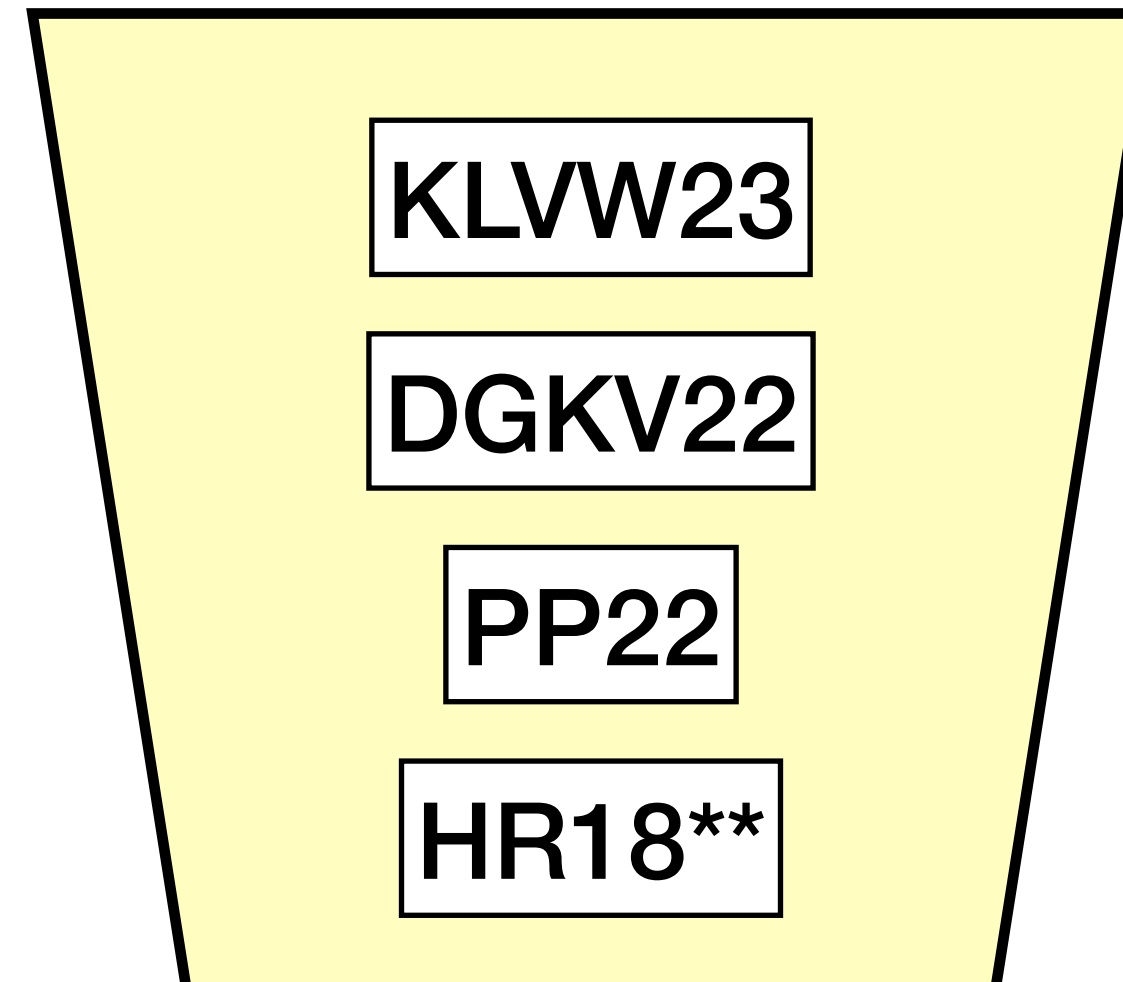
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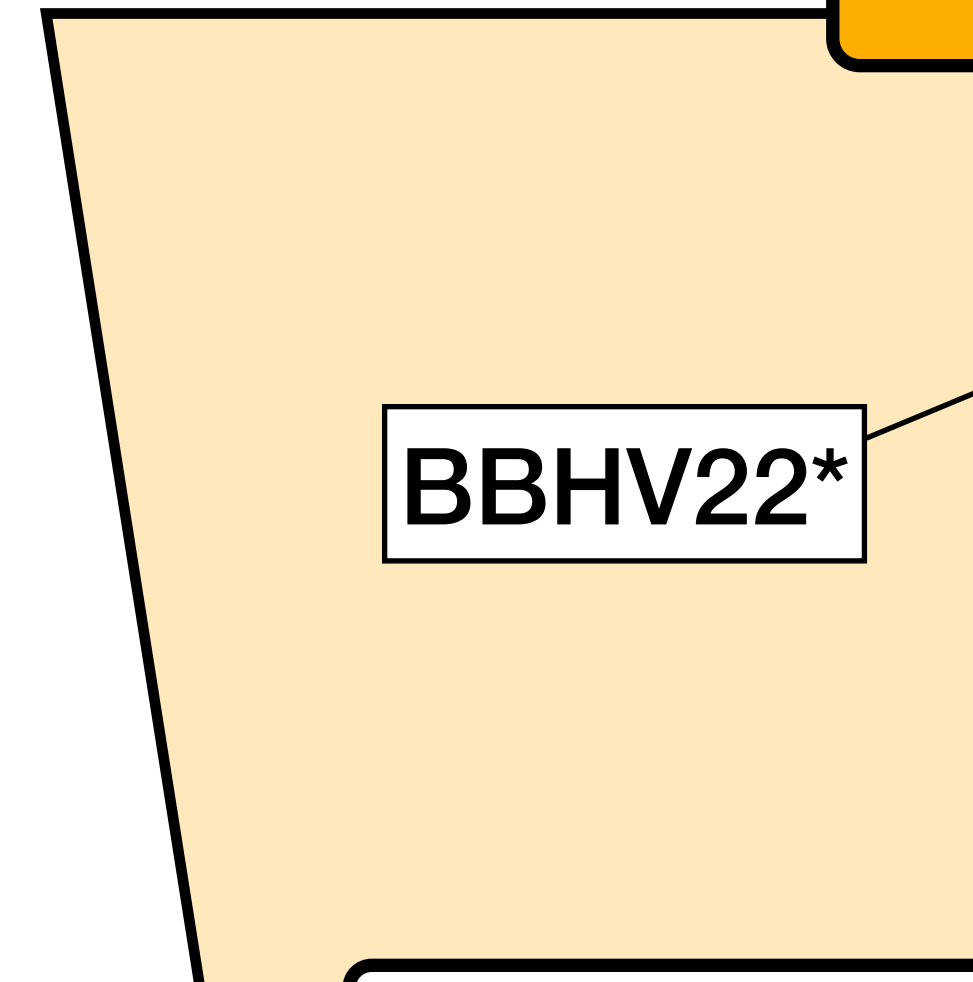
Non-standard  
assumptions



Non-interactive  
for P



Black-Box  
of CRH



“Tight” wrt  
black-box  
constructions

Comm:  $\tilde{O}(T/S)$   
Ver:  $\tilde{O}(T/S + S)$

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**Theorem:**

There exists a **public-coin**  $O(\log_\lambda T)$ -**message complexity-preserving** succinct argument for NP from **collision-resistant hashing**

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## [BC12] Questions:

Publicly  
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Minicrypt 

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## Theorem:

There exists a **public-coin**  $O(\log_\lambda T)$ -**message complexity-preserving** succinct argument for NP from **collision-resistant hashing**

## [BC12] Questions:

Publicly verifiable 

Minicrypt 

## Remaining Open Questions:

Fixed constant rounds?

Universal argument?

Turing machines?



# Corollaries to Our Main Result

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**Zero-Knowledge**

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Zero-Knowledge

“Commit-and-prove”  
via [BGGHKMR88]

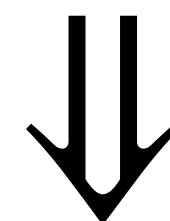
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Zero-Knowledge



Parallelizability

“Commit-and-prove”  
via [BGGHKMR88]

# Corollaries to Our Main Result

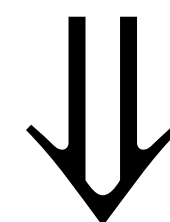
**Theorem:**

There exists a **public-coin**  $O(\log_\lambda T)$ -**message complexity-preserving** succinct argument for NP from **collision-resistant hashing**



Zero-Knowledge

“Commit-and-prove”  
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Parallelizability

Generic transformation  
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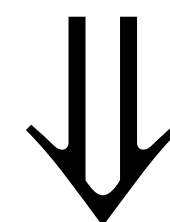
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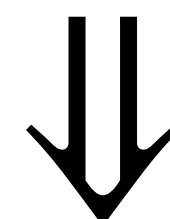
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Fiat-Shamir transform  
[FS86]

# Roadmap for Our Construction

**Theorem:**

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1. A  $\tilde{O}(T^{1/2}S)$ -space protocol

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Reduction via hash trees [KP16]

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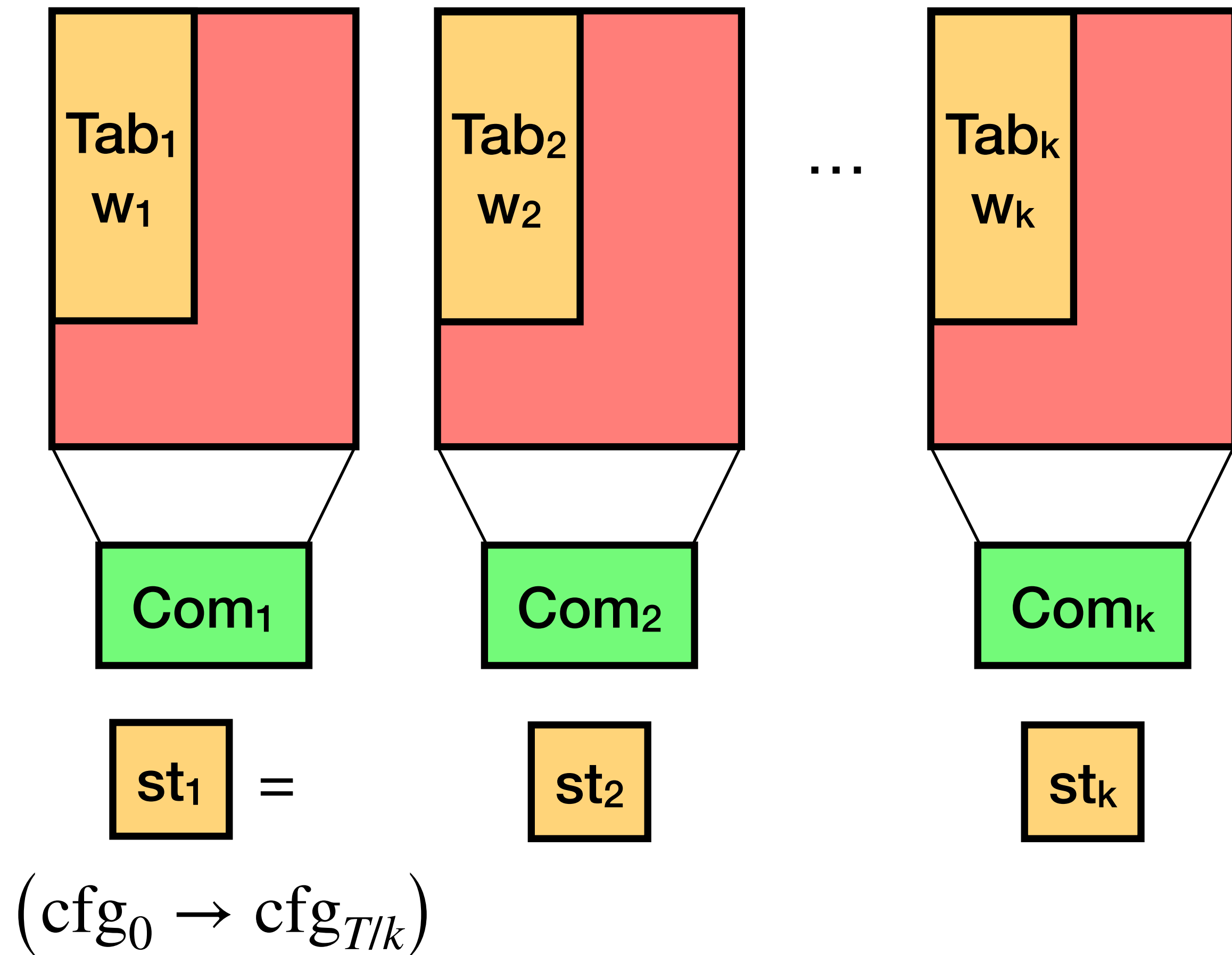
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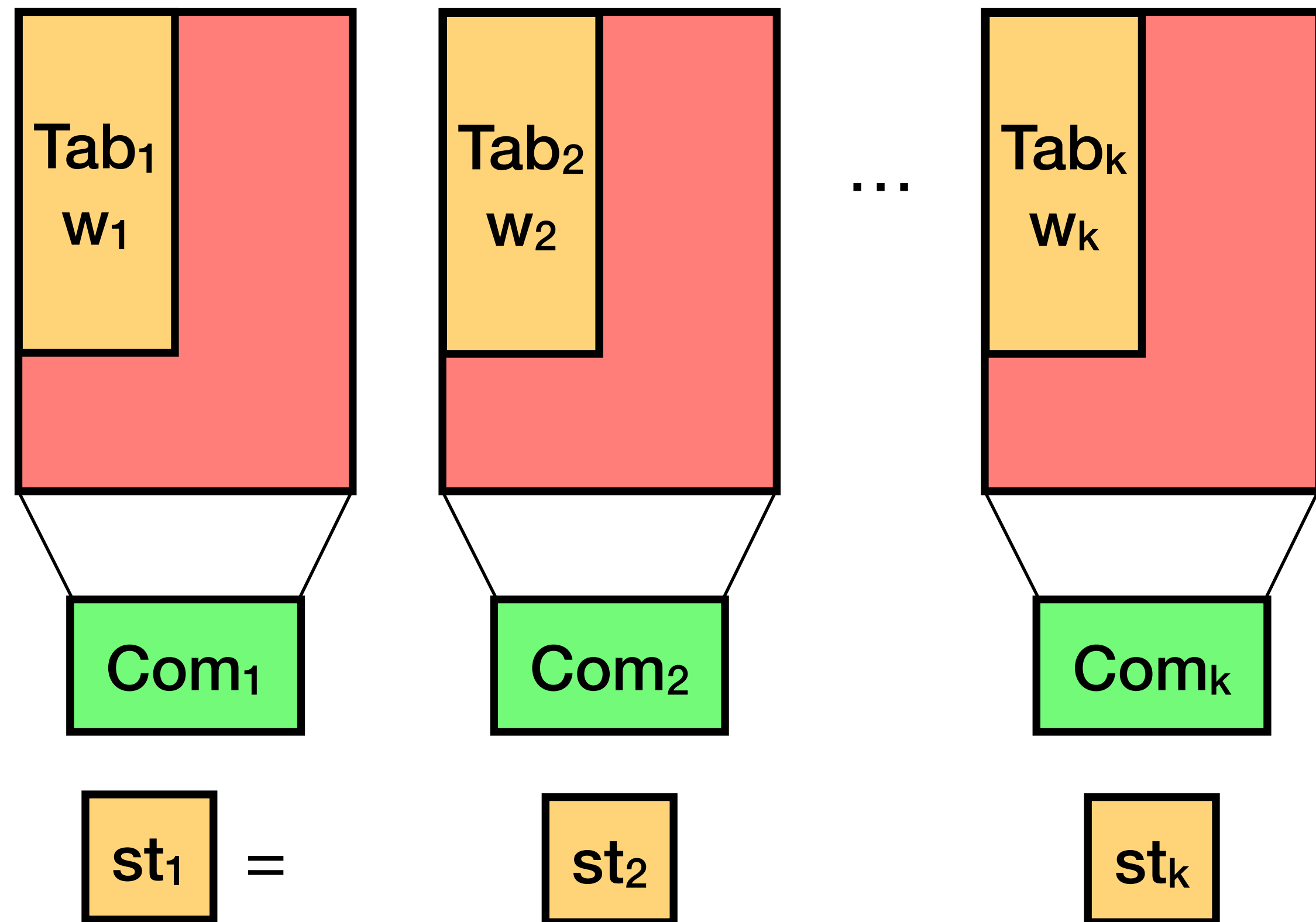
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$P'(x, w)$



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$(\text{cfg}_0 \rightarrow \text{cfg}_{T/k})$

**Prover efficiency:**

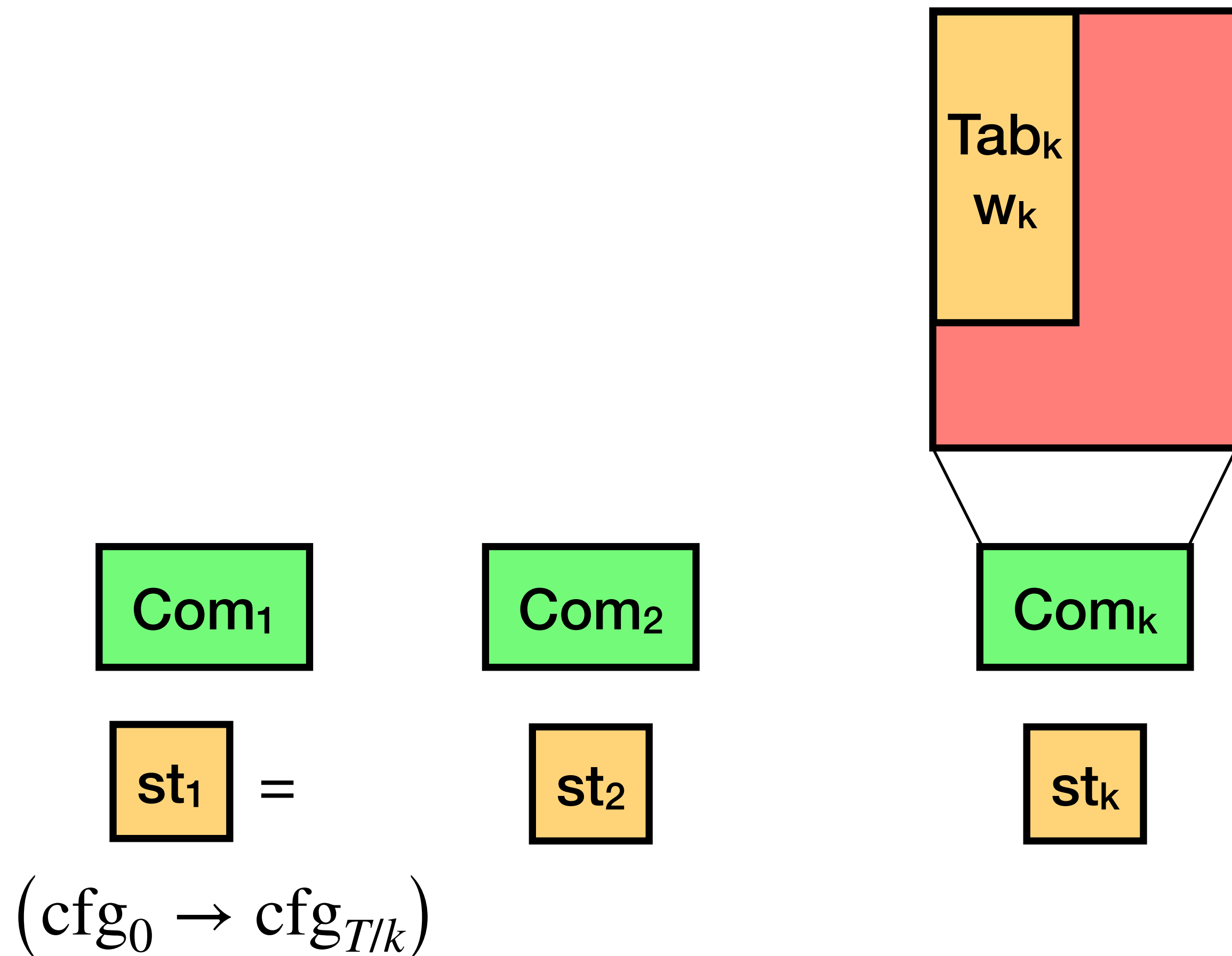
- Time  $k \cdot \tilde{O}(ST/k) = \tilde{O}(ST)$
- Space  $\tilde{O}(ST/k + Sk)$   
 $k = T^{1/2} \Rightarrow \tilde{O}(T^{1/2}S)$

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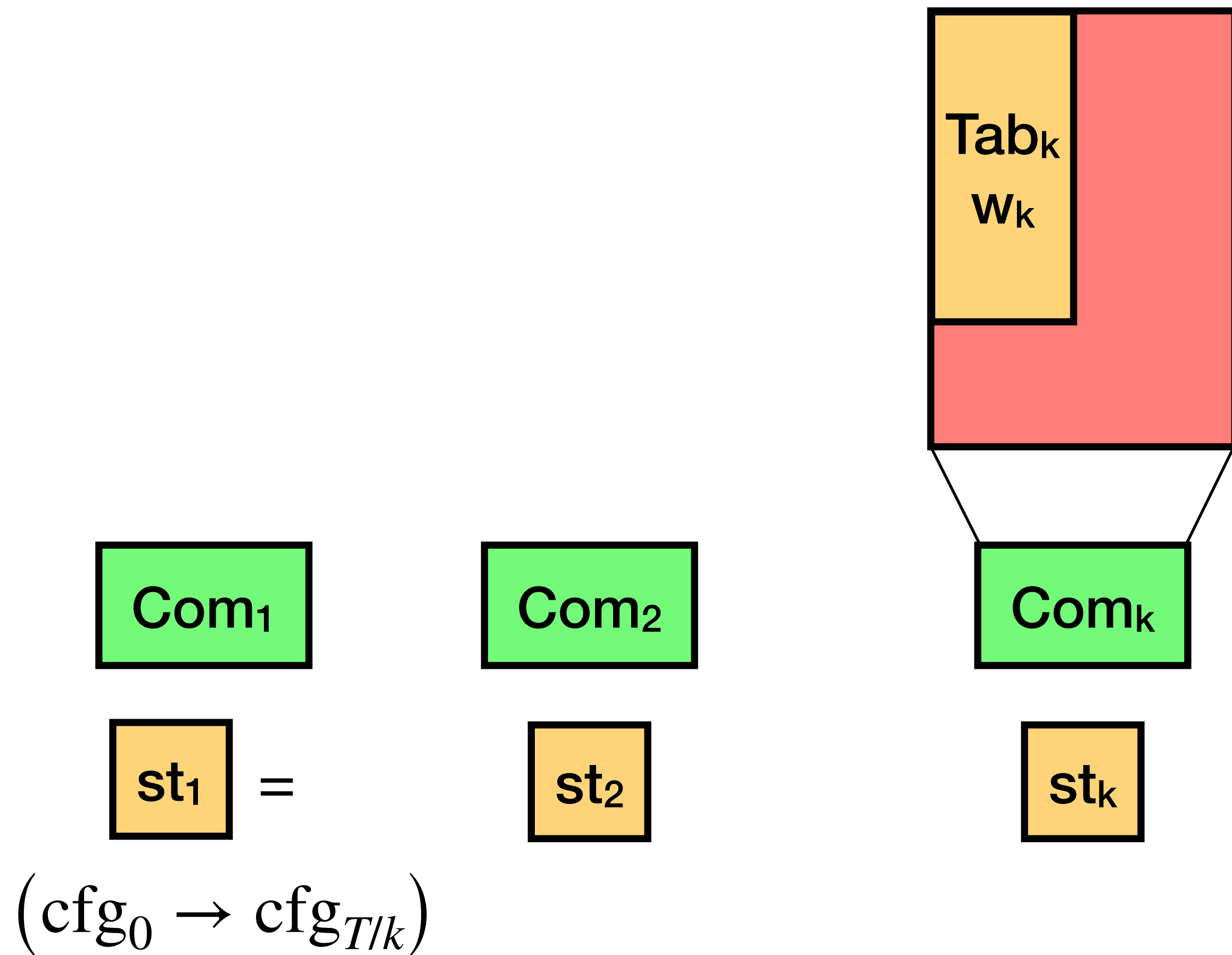
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## Communication/ Verification:

- $\tilde{O}(kS) = \tilde{O}(T^{1/2}S)$




# A $\tilde{O}(T^{1/2}S)$ -space protocol

## Intermediate Result:

There exists a public-coin 4-message argument for NP from CRH with:

- prover **time**  $\tilde{O}(TS)$  and **space**  $\tilde{O}(T^{1/2}S)$
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Recursion!

# Recursive Protocol

$P_r(w)$

$V_r$

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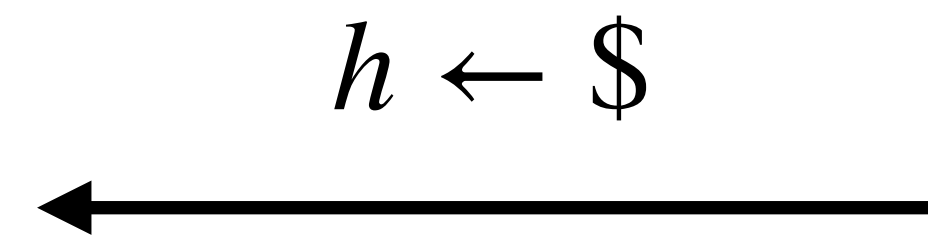
$h \leftarrow \$$



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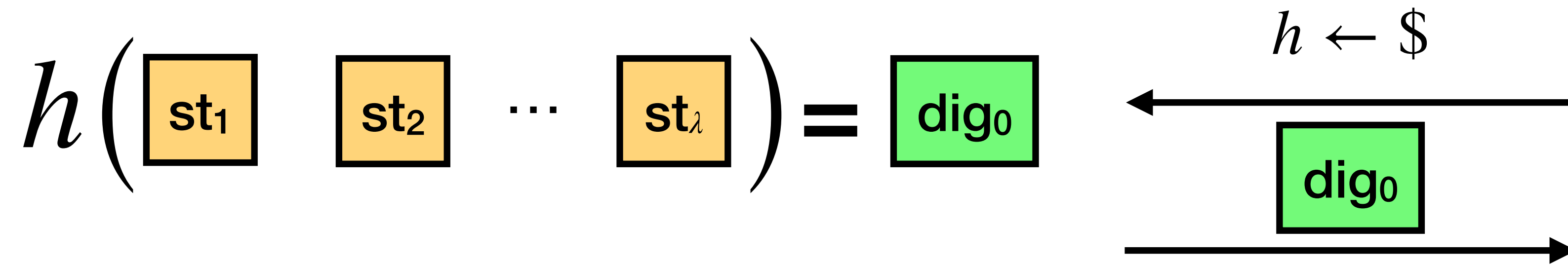




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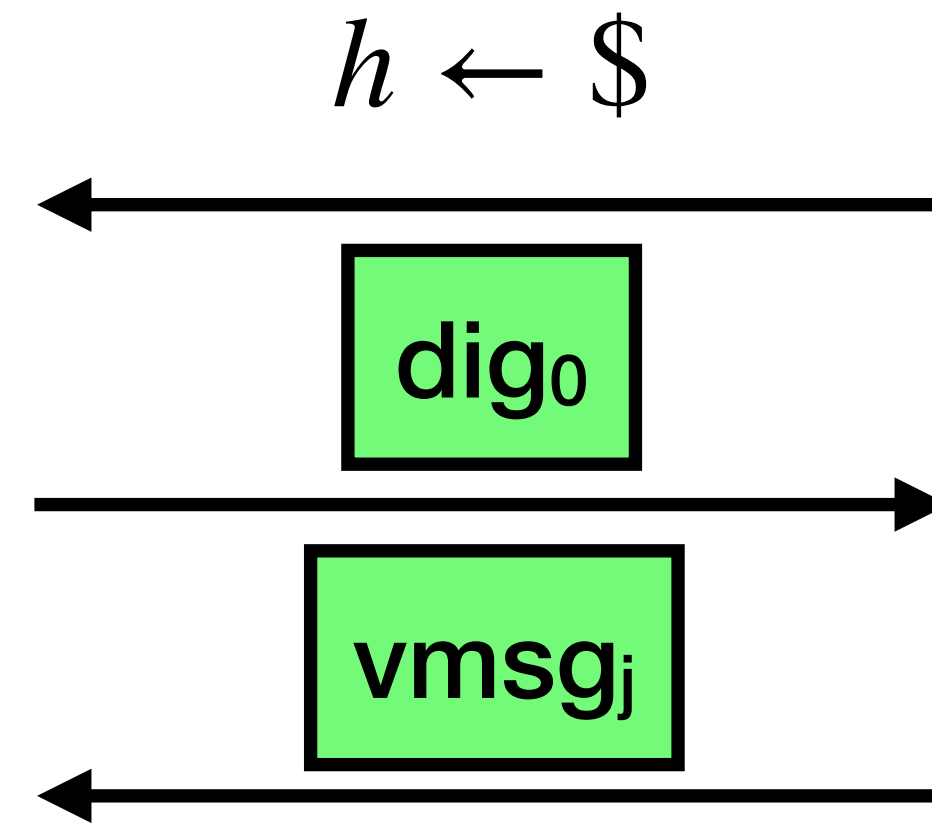


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$$h\left(\boxed{st_1} \quad \boxed{st_2} \quad \dots \quad \boxed{st_\lambda}\right) = \boxed{dig_0}$$



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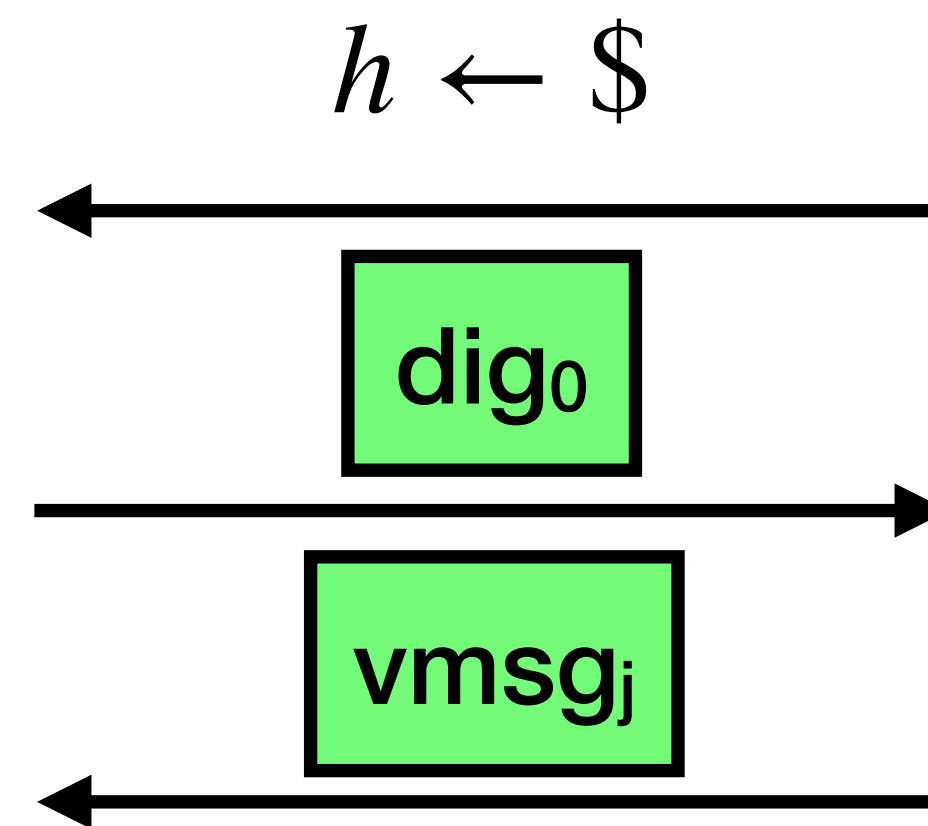
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$\boxed{msg_{j,1}} \quad \boxed{msg_{j,2}} \quad \dots \quad \boxed{msg_{j,k}}$



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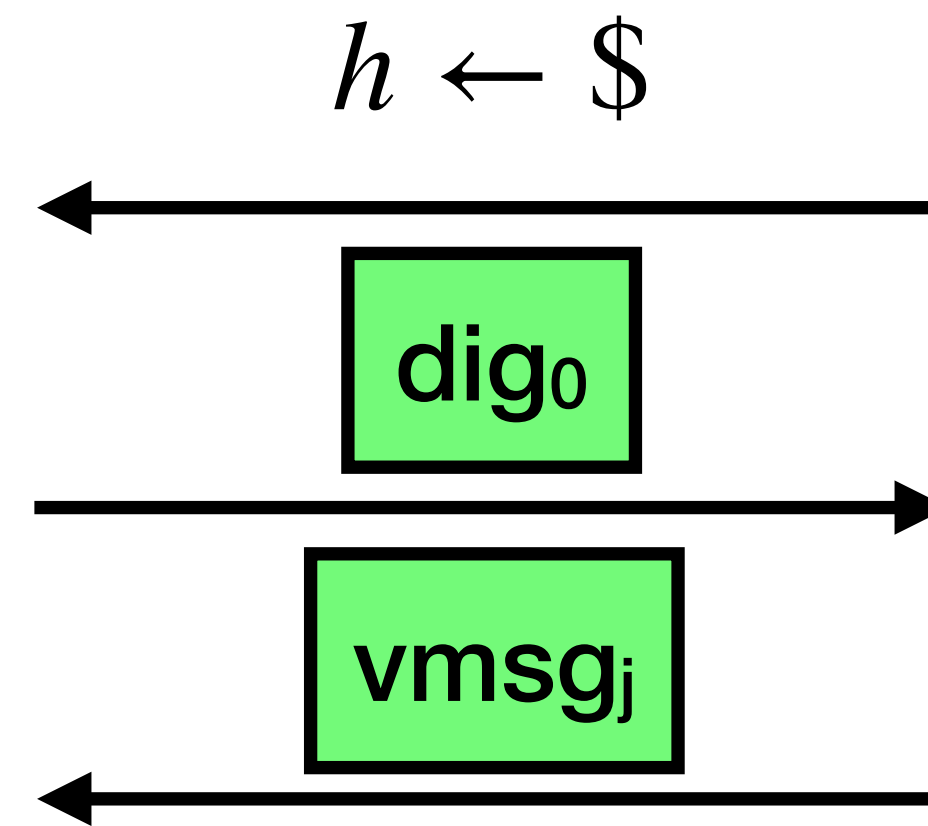
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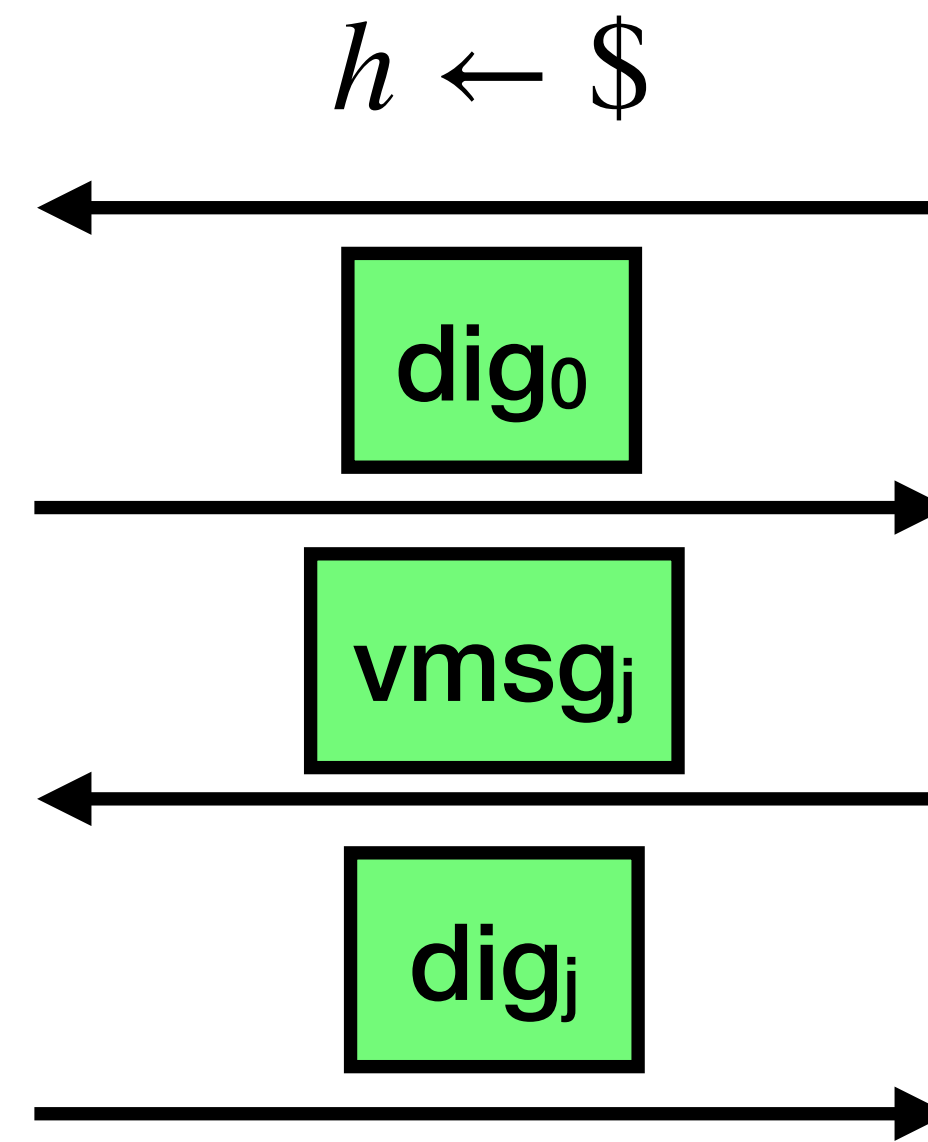
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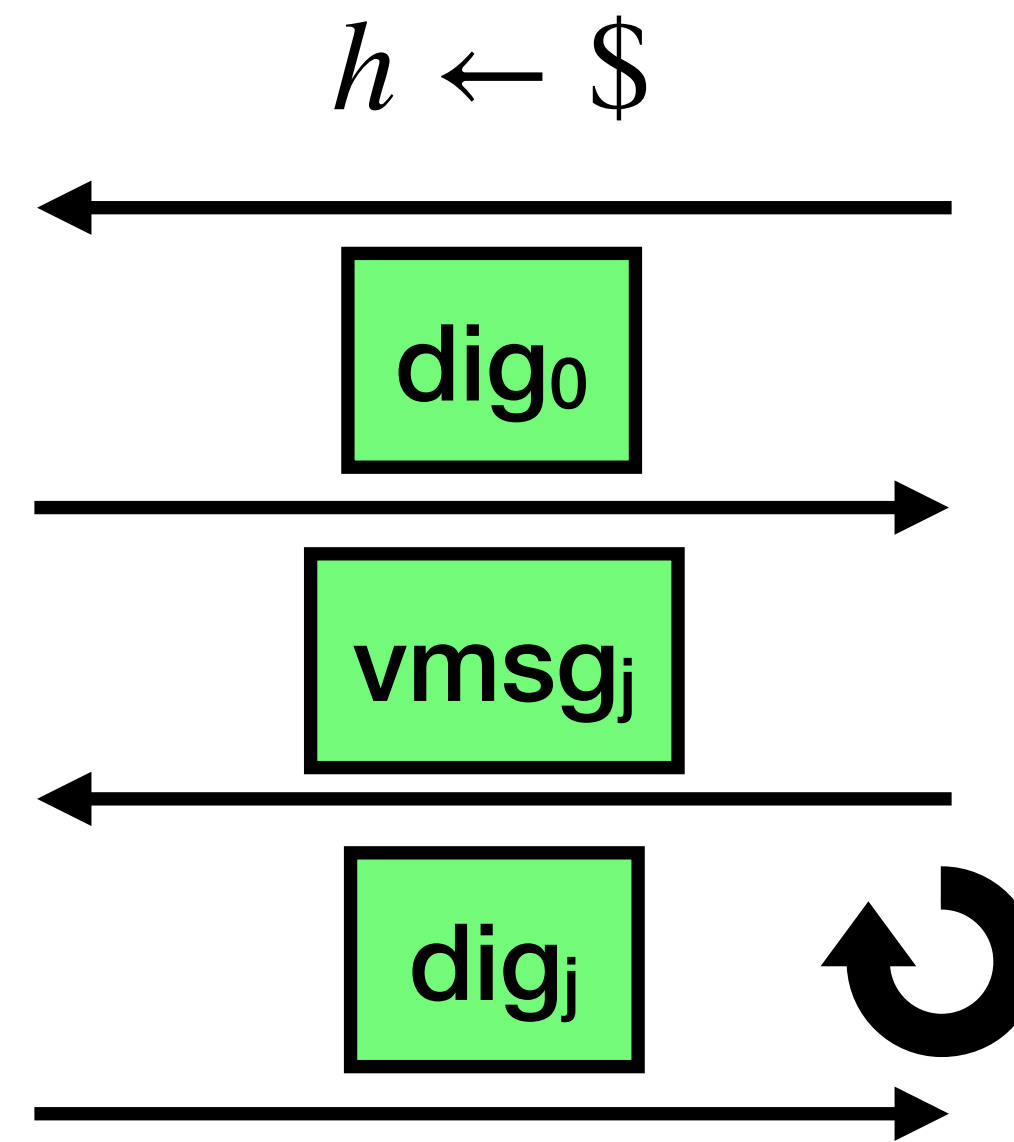
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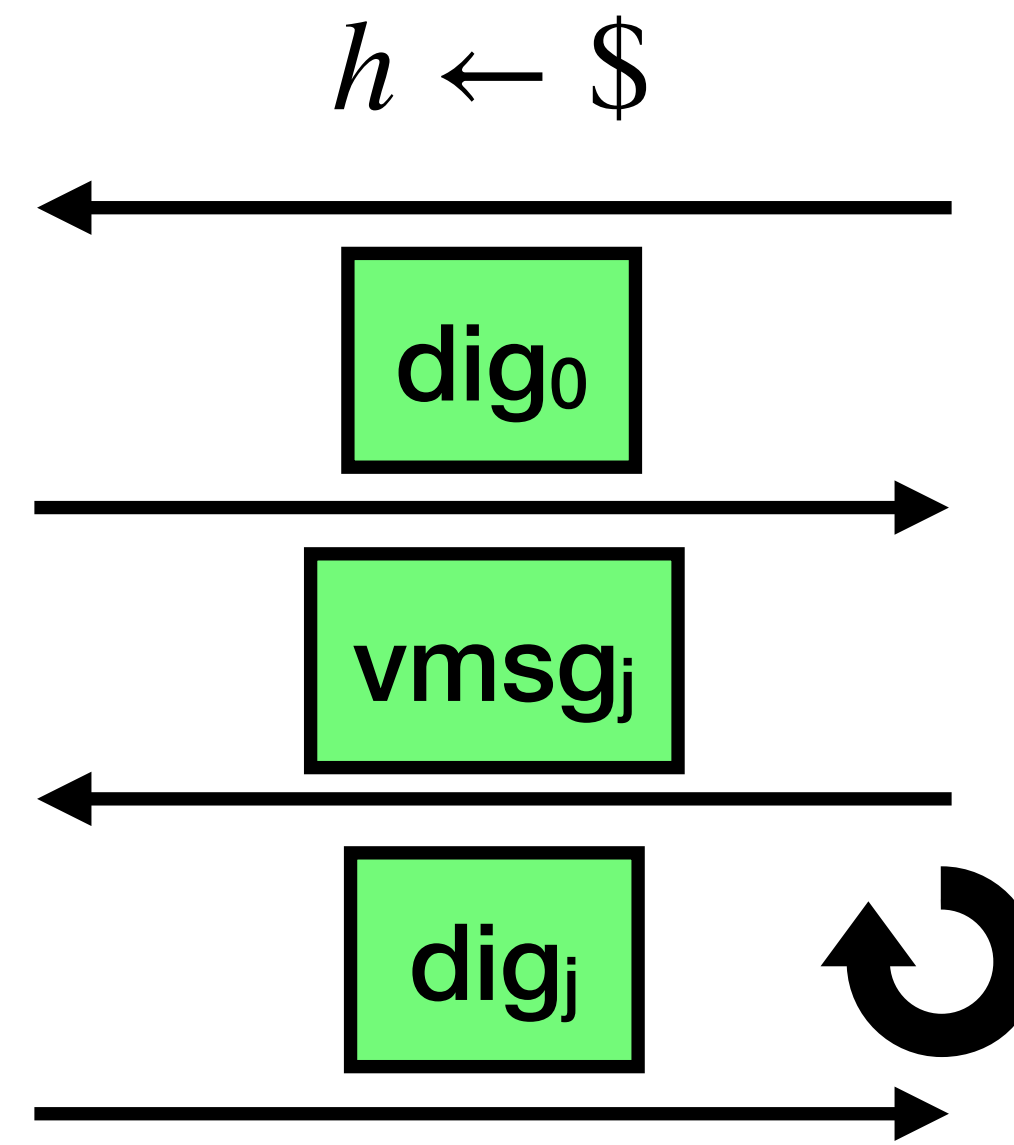
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$$R_{\text{merge}} = \left\{ \begin{array}{l} ((\text{cfg}, \text{cfg}', \vec{dig}), (\vec{st}, \text{Msg})) : \\ 1. \text{dig}_0 = h(\vec{st}) \\ 2. \forall j \geq 1, \text{dig}_j = h(\text{Msg}_{j,*}) \\ 3. \forall i \in [k], V_{r-1}(st_i, \text{Msg}_{*,i}) = 1 \end{array} \right\}$$

Succinct argument  
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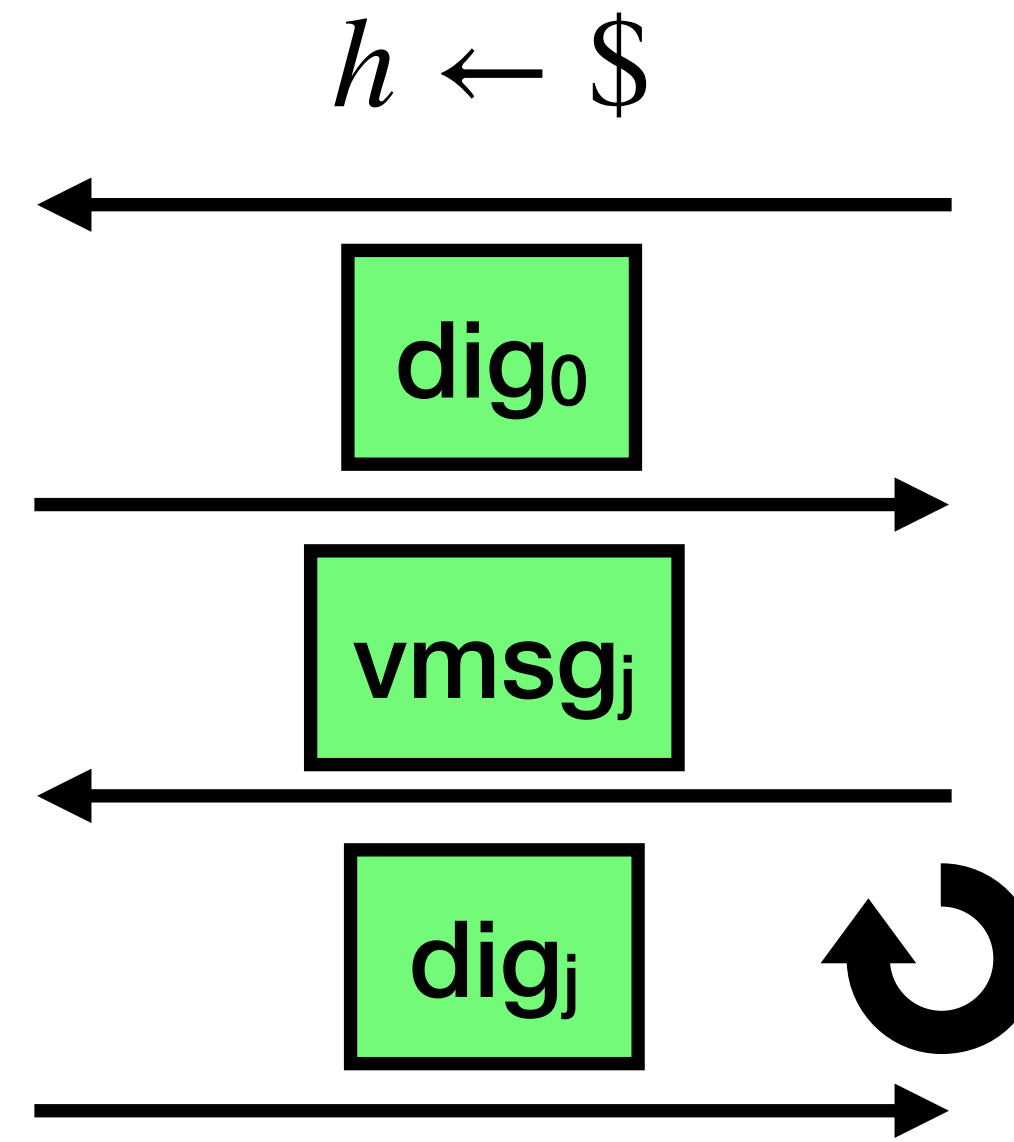
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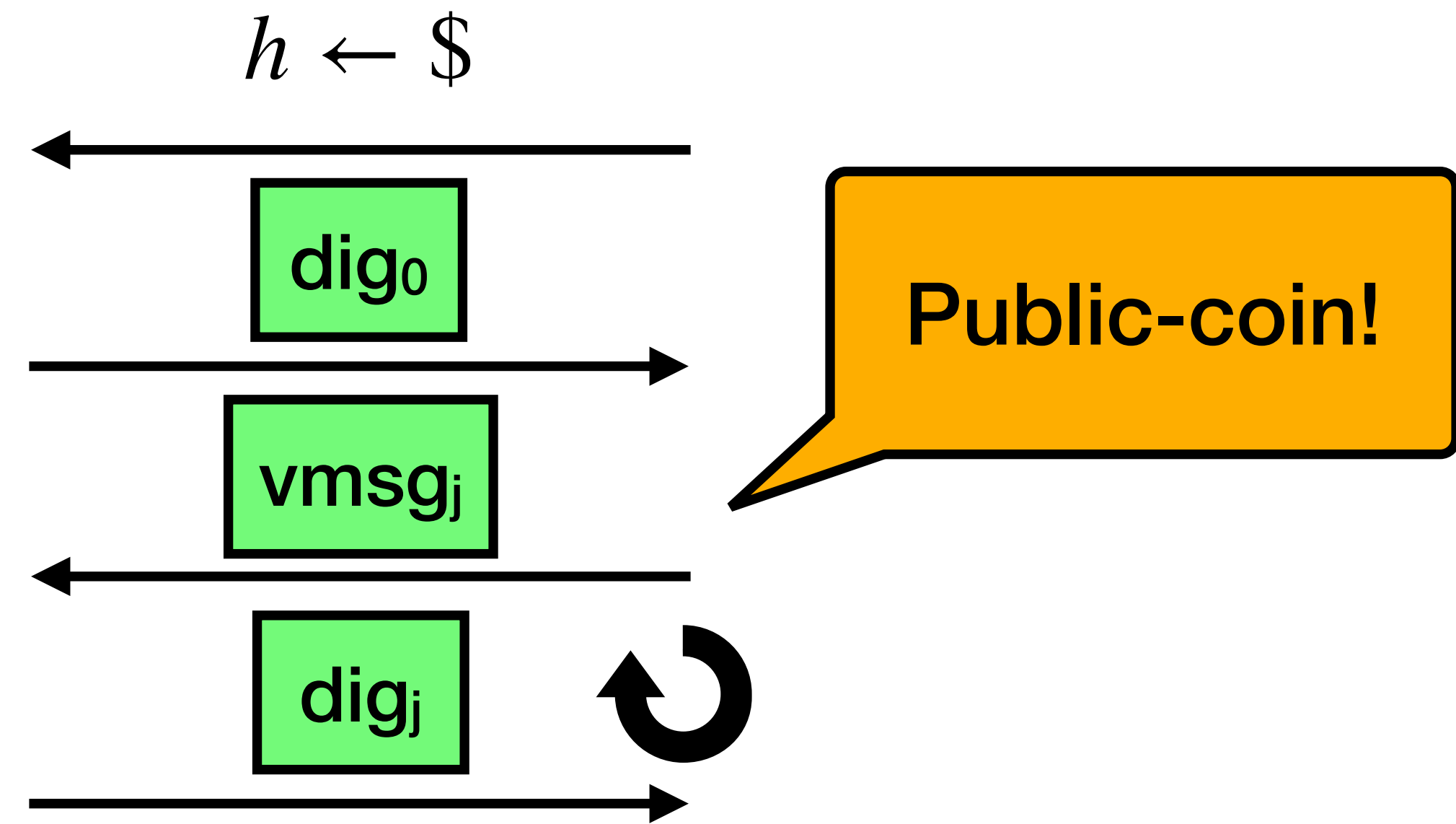
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# Questions?

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## Open Questions:

Fixed constant  
rounds?

Universal  
argument?

Turing  
machines?