

Unbiasable Verifiable Random Functions

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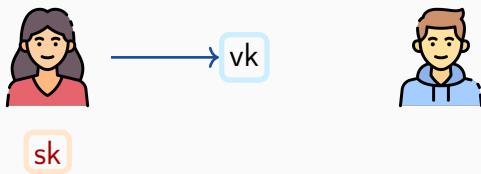
`emanuele.giunta@imdea.org`

Universidad Politecnica de Madrid, Spain.

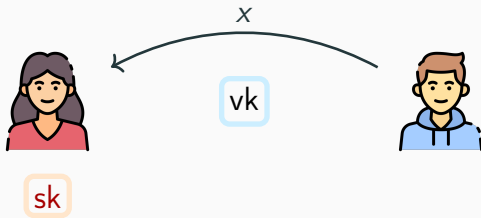
VRF.Gen, VRF.Eval, VRF.Vfy



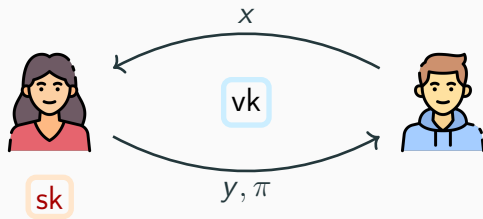
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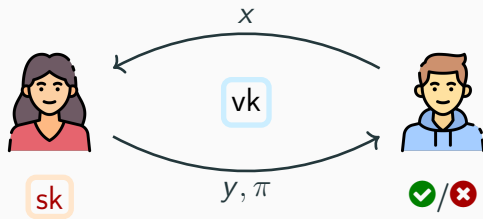
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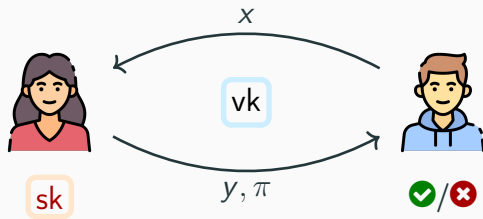
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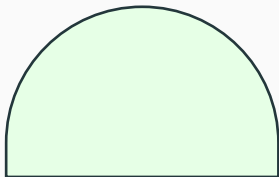


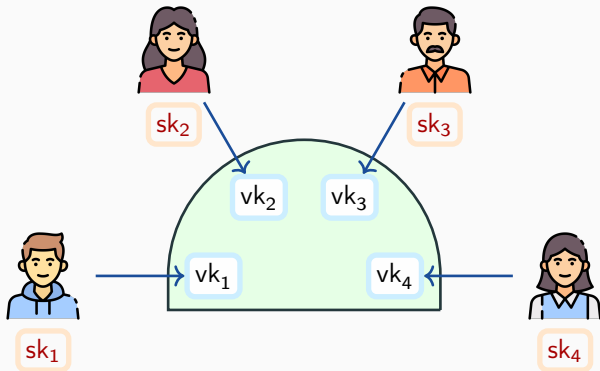
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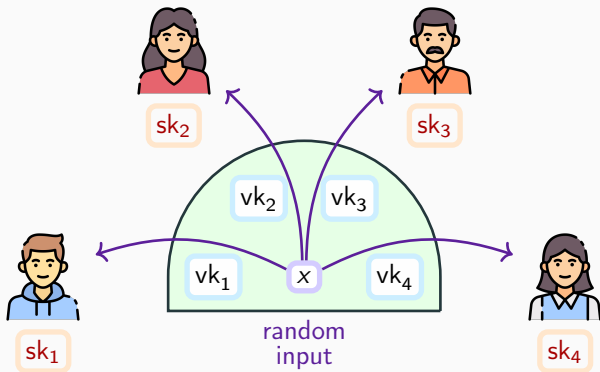


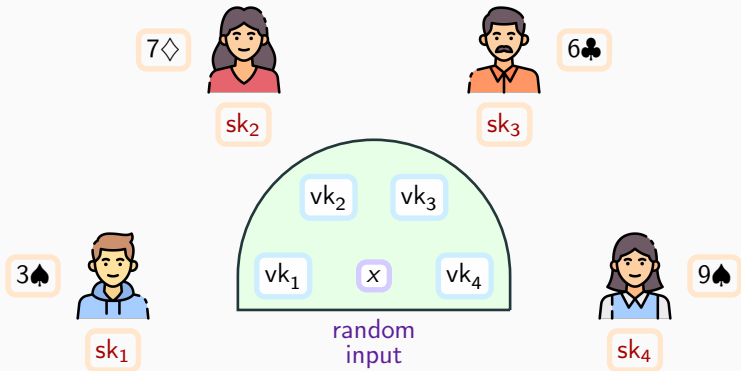
Uniqueness

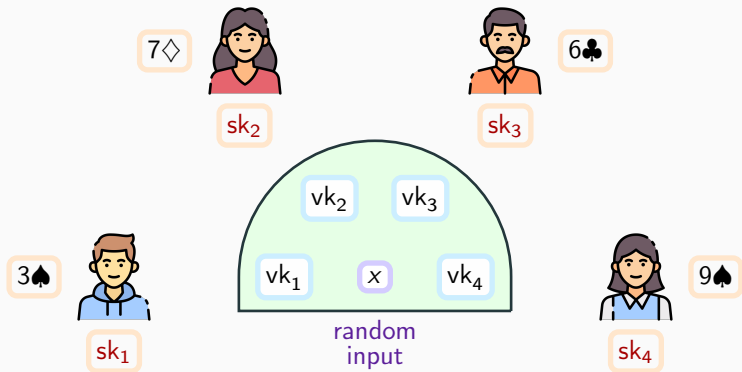
Pseudorandom





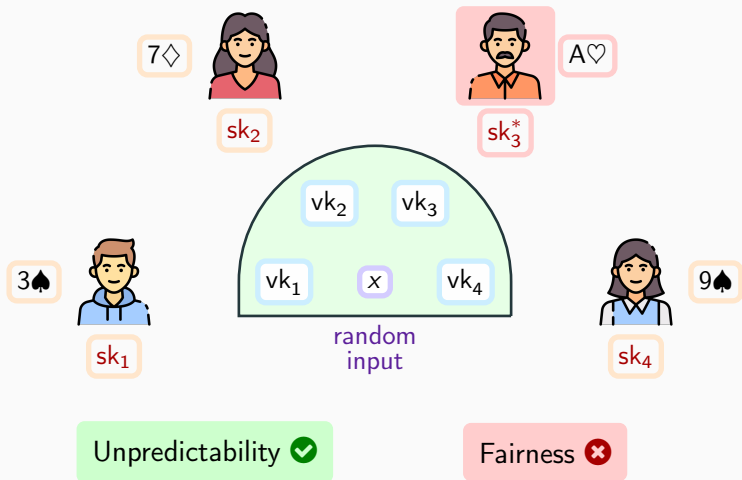


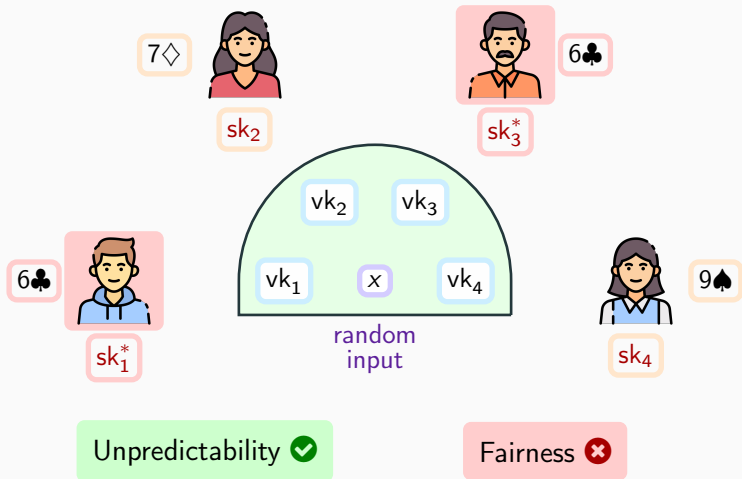




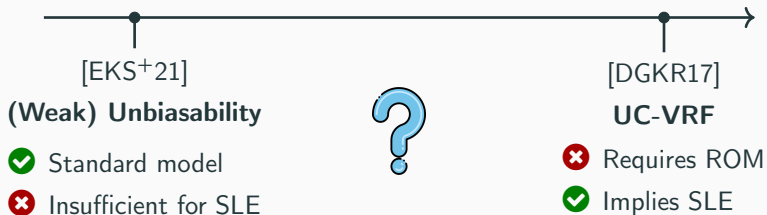
Unpredictability

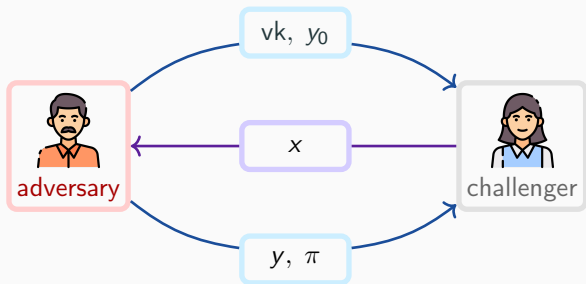
Fairness



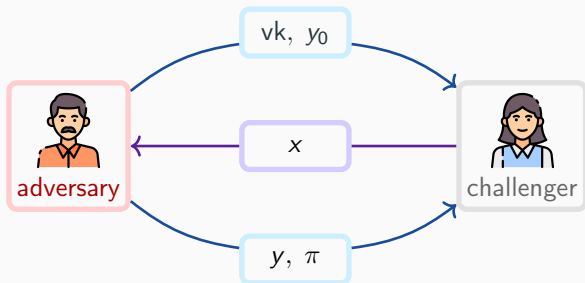


VRF Unbiasability





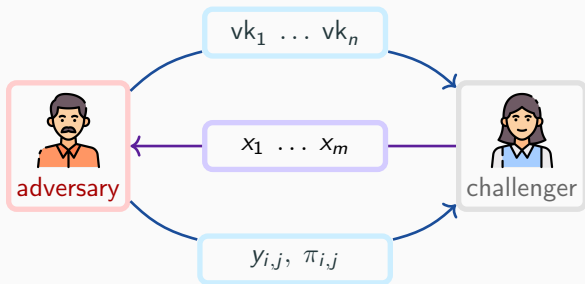
Win if $y = y_0$ and π is correct.



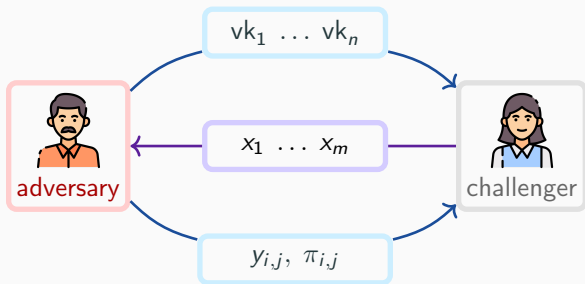
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✘ Some output bits could be biased

✘ Different key's output could be correlated

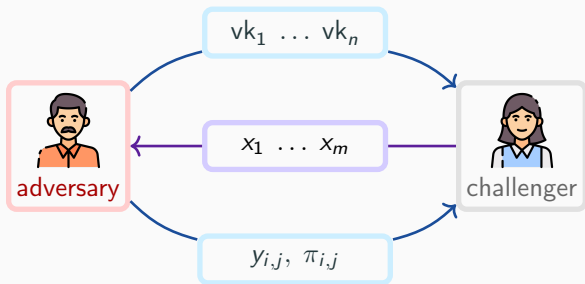


Win if $\pi_{i,j}$ are **correct** $y_{i,j}$ are **far** from random.

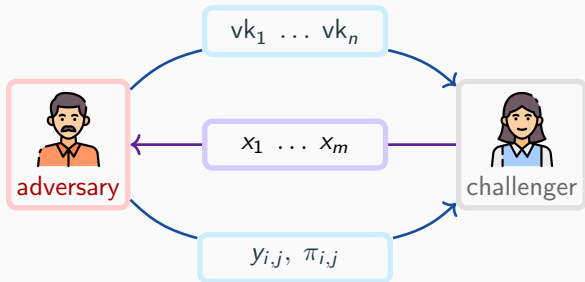


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- ✘ Does not exclude biased vk^* that cannot be evaluated on the full domain

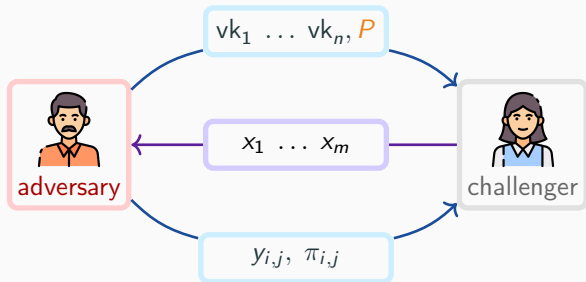


1. If $\pi_{i,j}$ is incorrect, **erase** $y_{i,j} \leftarrow \perp$
2. Win if **non-erased** $y_{i,j}$ are **far** from random



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✘ Selective openings
are always biased



1. P is **monotone**¹.
2. If $\pi_{i,j}$ is incorrect, **erase** $y_{i,j} \leftarrow \perp$.
3. Win if $\Pr [P(y) = 1] \gg \Pr [P(z) = 1]$ for a **random** z .

¹Let x^* be x with some positions erased. Then $P(x^*) \leq P(x)$.

Constructions

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3. **Unbiasable VRF** from any VRF + DL* + PRF*.

VUF to VRF transform in the ROM

$F_{sk} : X \rightarrow Y$ Verifiable Unpredictable Function with public key vk .

1st Transform: $F_{sk}^*(x) = H(F_{sk}(x), x, vk)$.

2nd Transform: $F_{sk}^*(x) = H(F_{sk}(x))$.

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- Both are VRF in the ROM.
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 1. $F_{sk}(x_1) \neq F_{sk}(x_2)$ for random x_1, x_2 .
 2. $F_{sk_1}(x) \neq F_{sk_2}(x)$ for random x and $sk_1 \neq sk_2$.

Verifiable Random Bijection

\mathbb{G} prime order group where DL is hard.

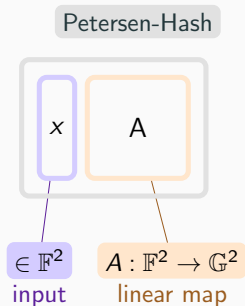
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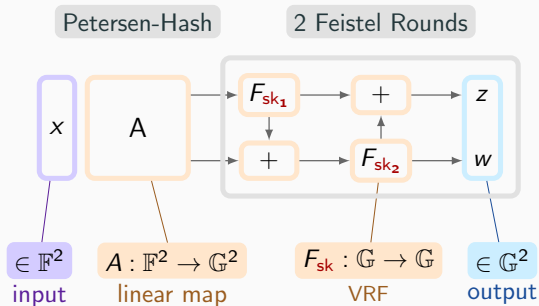
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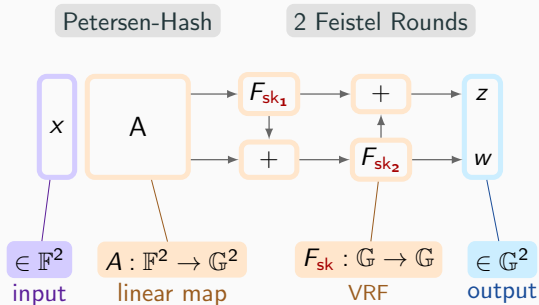
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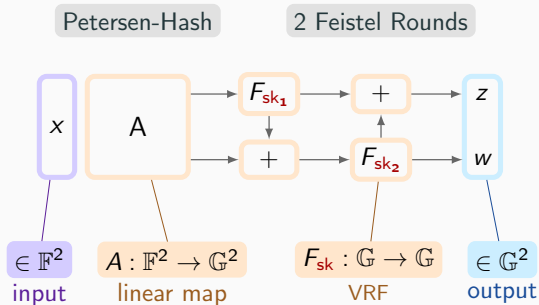
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✓ Certified bijection

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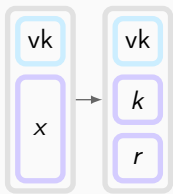
- ✓ Certified bijection
- ✗ Unbiasable only for a single vk!

Let F_{sk}^* be a VRB with verification key vk and f a PRF.

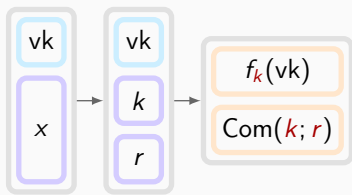
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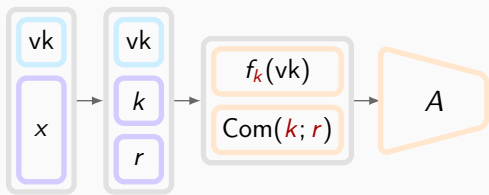
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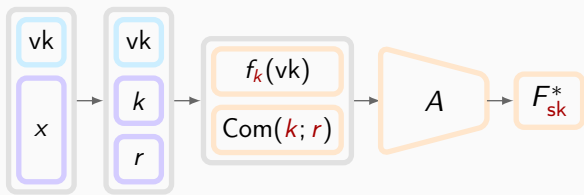
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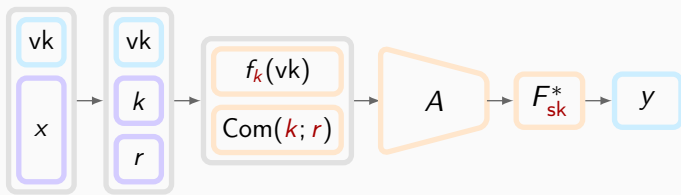
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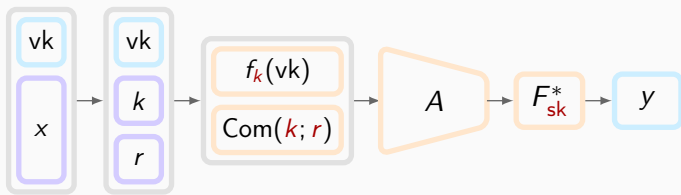
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Unbiasable if **DL** and the **PRF** are secure with **preprocessing**.

Conclusions

We provided a new notion of **unbiasability** that is:

- **Sufficient** for applications (e.g. leader election).
- **Satisfied** by existing constructions in the NPRM.
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Open questions:

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Thanks for your attention!