

Time-Lock Puzzles with Efficient Batch Solving



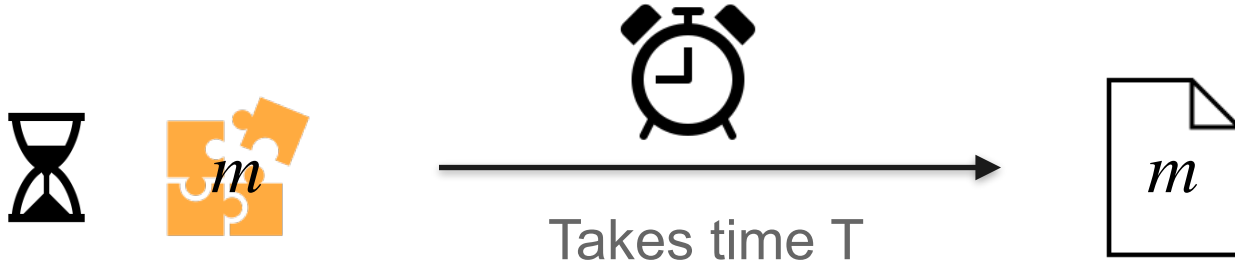
Jesko
Dujmovic
CISPA
and Saarland
University

Rachit
Garg
UT Austin

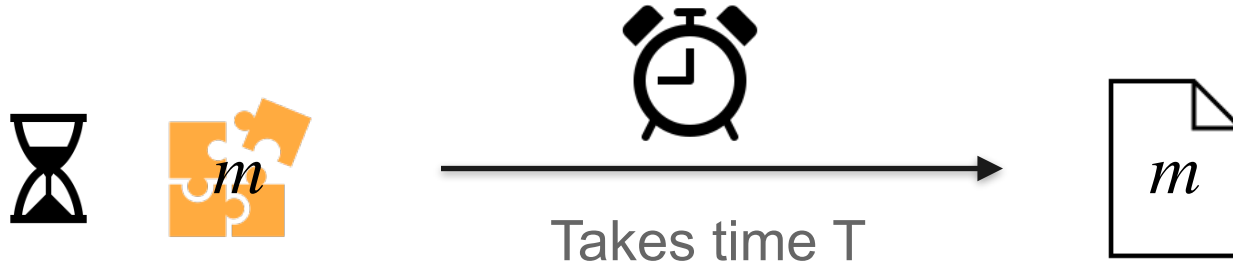
Giulio
Malavolta
Bocconi
University
and MPI-SP


Time-Lock Puzzles [May93, RSW96]

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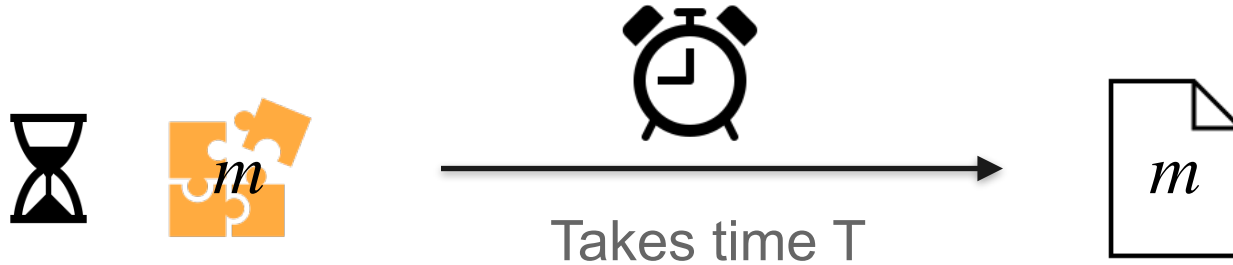




Time-Lock Puzzles [May93, RSW96]



- Fast puzzle generation - Time to generate  is much shorter than time T .

Time-Lock Puzzles [May93, RSW96]



- Fast puzzle generation - Time to generate  is much shorter than time T .
- Puzzle opening takes a long time - The circuit that opens  has depth at least T . Parallelism shouldn't help.

Applications

Applications



Encrypt to the future!

Applications



Encrypt to the future!



Sealed Bid Auctions

Applications



Encrypt to the future!



Sealed Bid Auctions



Non-Malleable Commitments

Applications



Encrypt to the future!



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Non-Malleable Commitments



Miner extractable value prevention

Applications



Encrypt to the future!



Sealed Bid Auctions



Non-Malleable Commitments

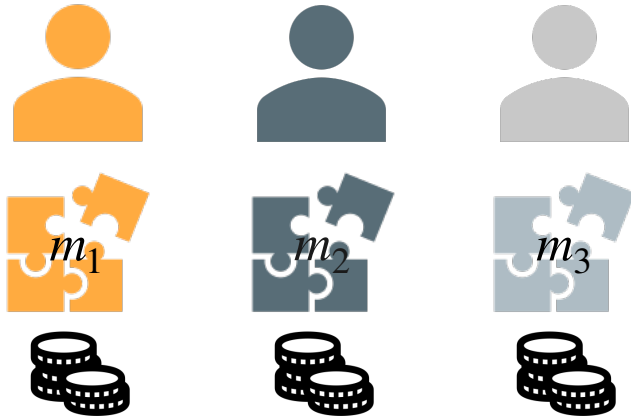


Miner extractable value prevention

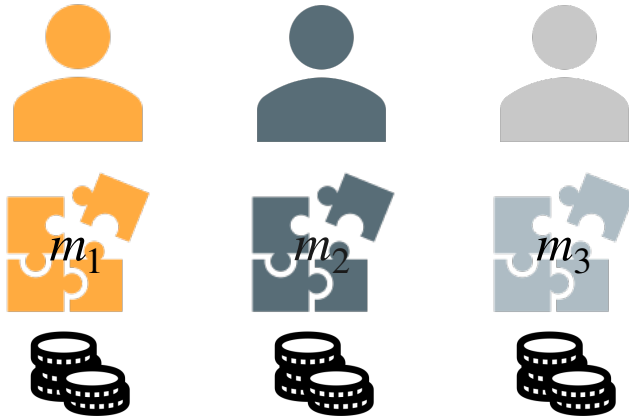
Blockchain front running prevention, fair contract signing, cryptocurrency payments, distributed consensus, more!

Applications - Batch Solving

Applications - Batch Solving



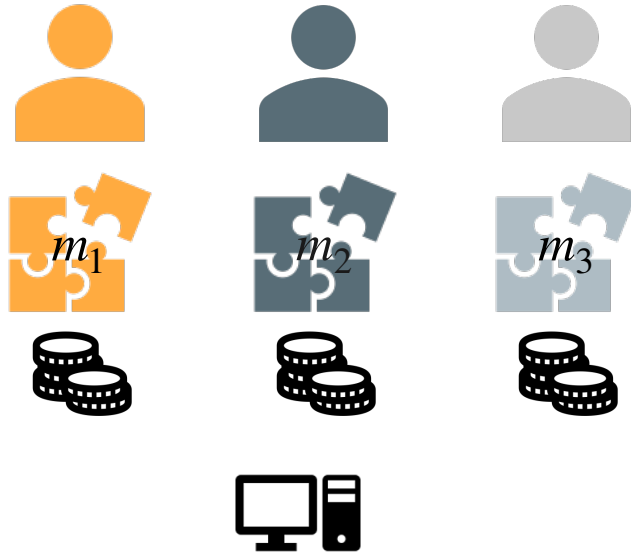
Applications - Batch Solving



Decrypt all transactions!
Solve all puzzles



Applications - Batch Solving



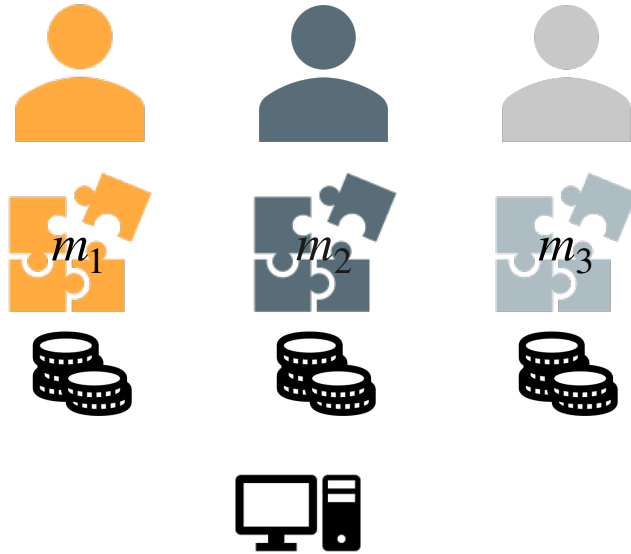
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Blockchains, byzantine broadcast

Scalability - Millions of users need solving

Applications - Batch Solving



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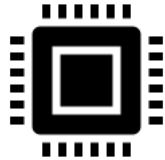
Denial of service attacks

Batching Complexity

- Fast batch solving - Time to solve    , multiple puzzles, grows with the time to solve a “single” puzzle.

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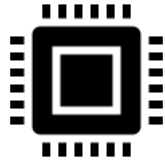


$$N \cdot \text{poly}(T)$$

Trivial solution

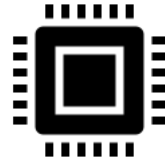
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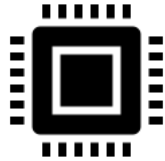
Trivial solution



$$o(N) \cdot \text{poly}(T) + \text{poly}(\log T, N)$$

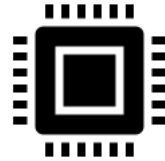
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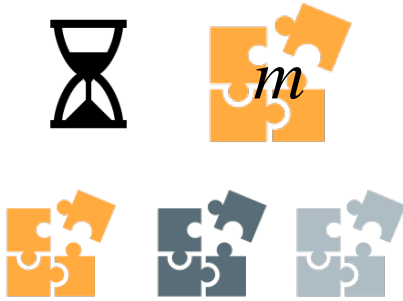


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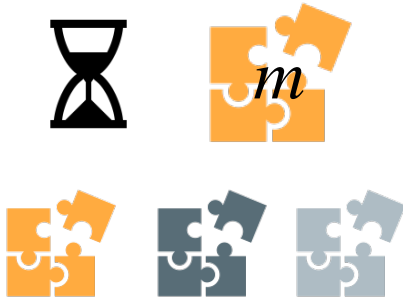
$$\boxed{\text{poly}(T)} + \text{poly}(\log T, N)$$

This work

Our Result

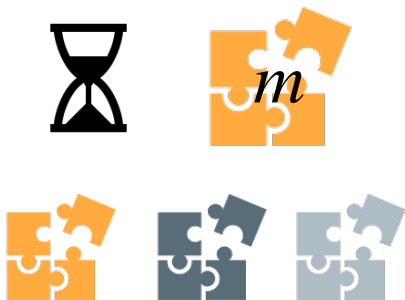


Our Result



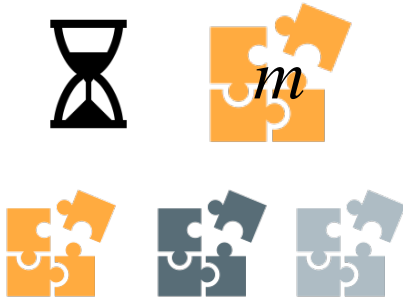
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Our Result



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- We give two concrete constructions and an implementation.

Our Result



- Generic template for constructing batchable TLPs.

Only prior solution was based on iO [SLM+23].

- We give two concrete constructions and an implementation.
- Introduce the notion of rogue batch solving.



Roadmap

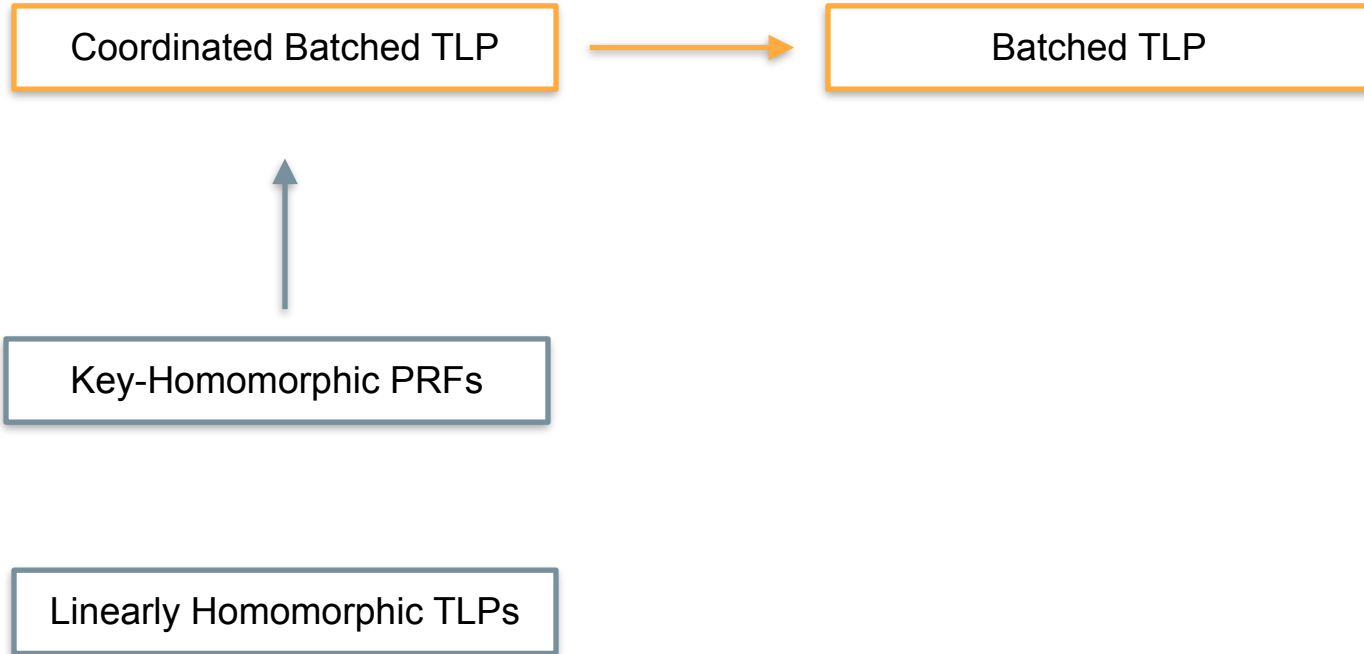
Roadmap

Batched TLP

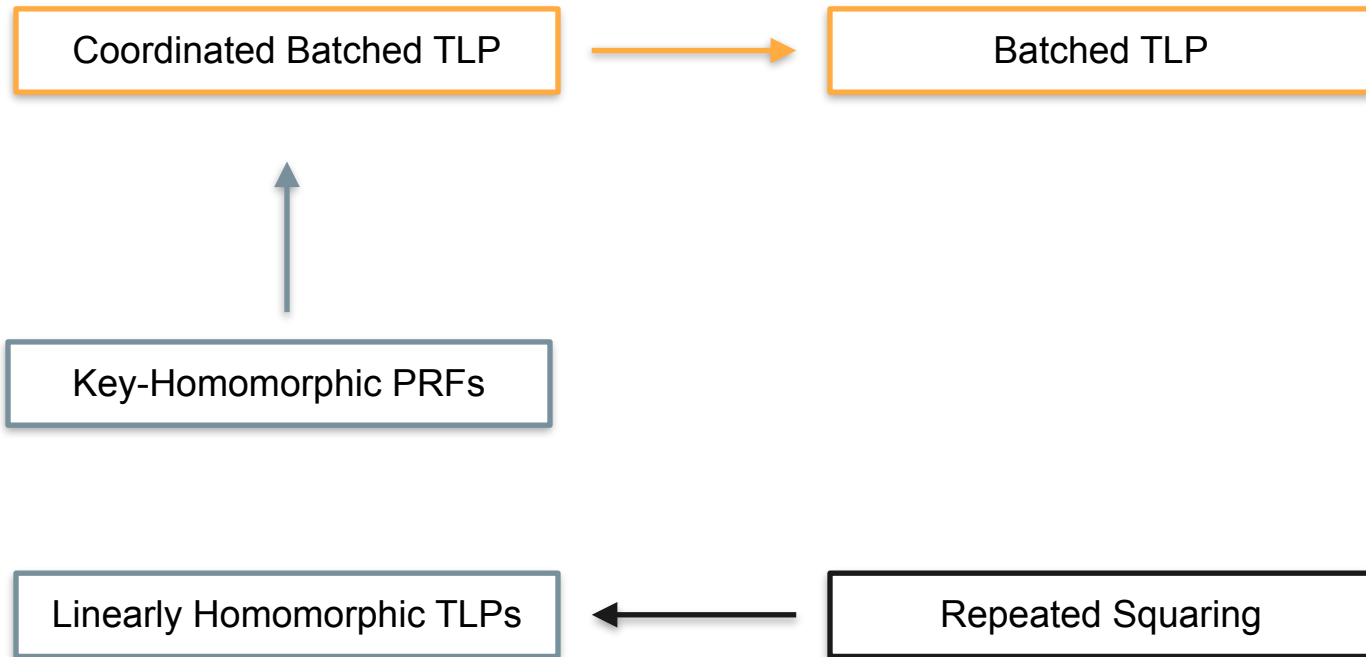
Roadmap



Roadmap

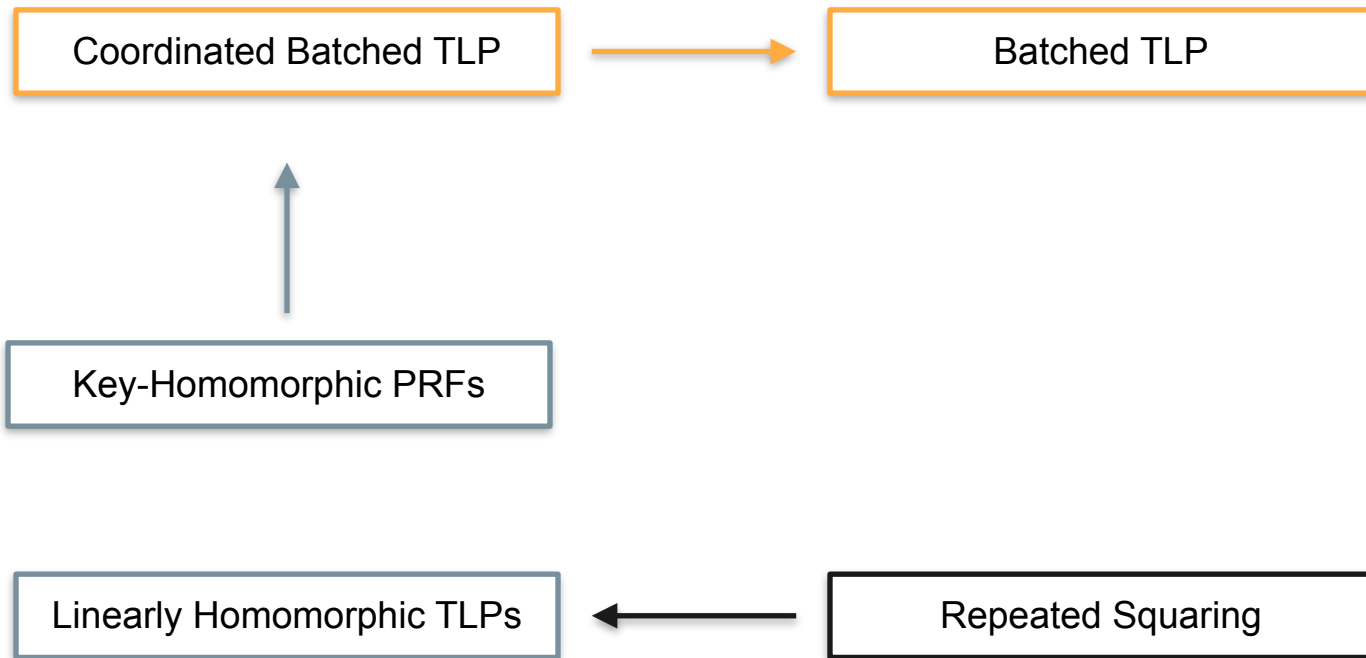


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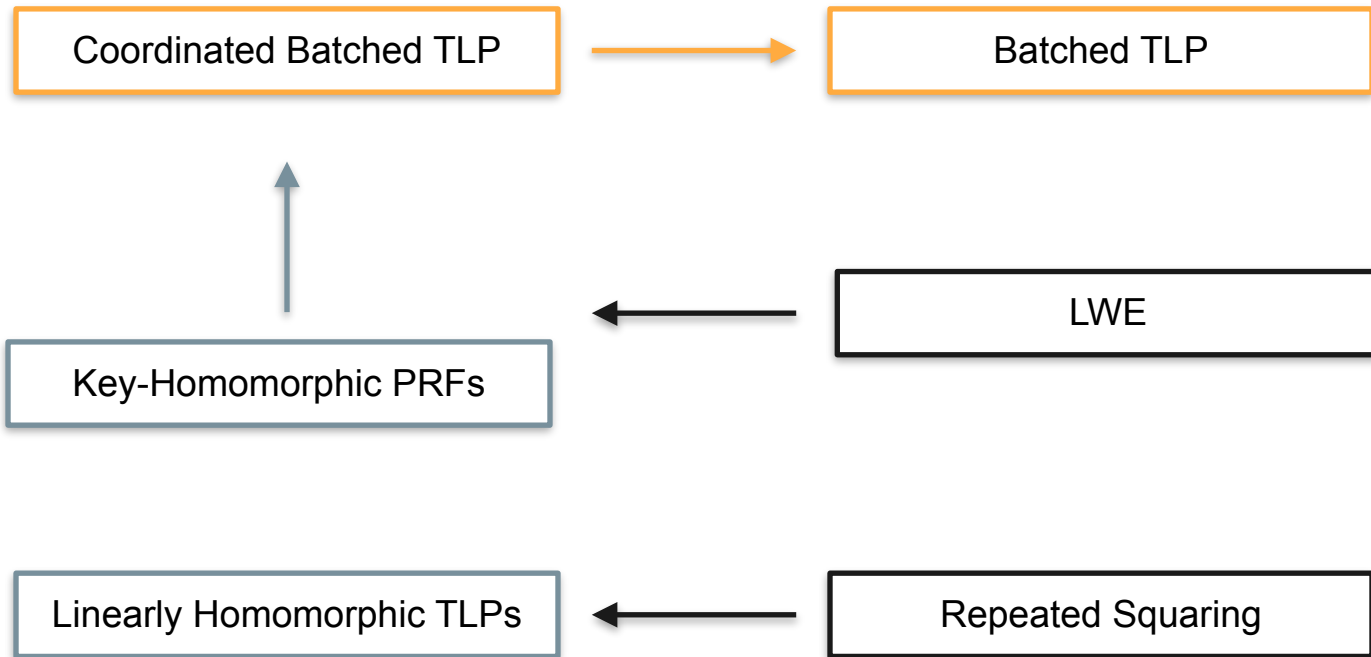
[MT19, TCLM21]

Roadmap



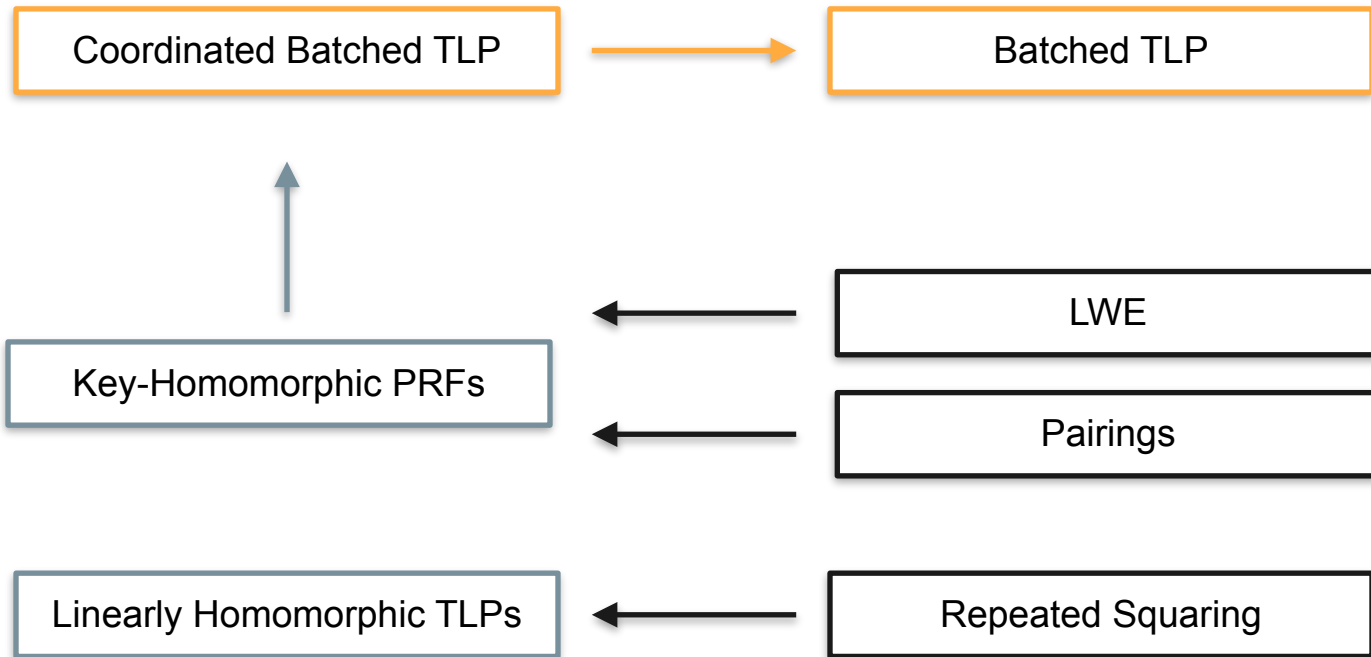
[MT19, TCLM21] $x, x^2, x^4 = (x^2)^2, x^8, \dots, x^{2^T}$

Roadmap



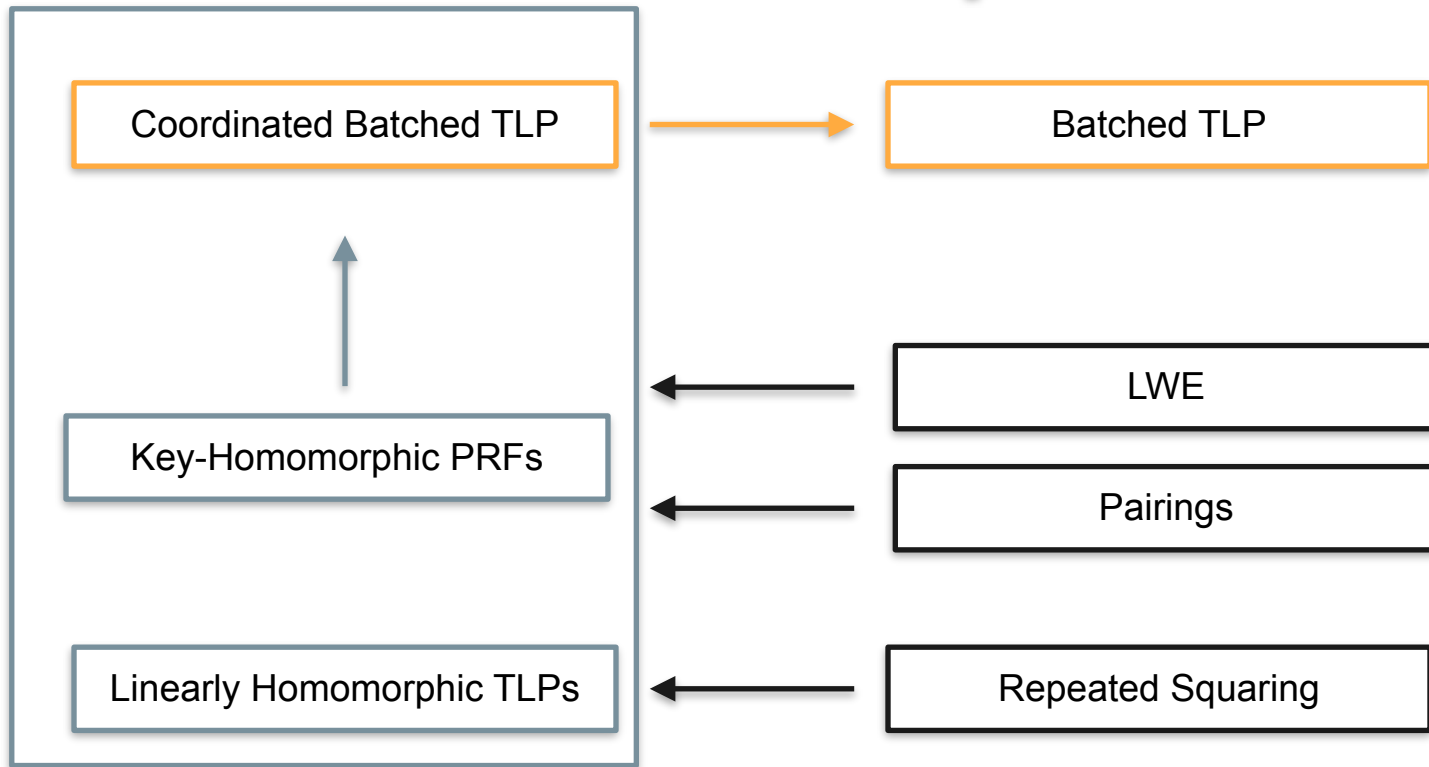
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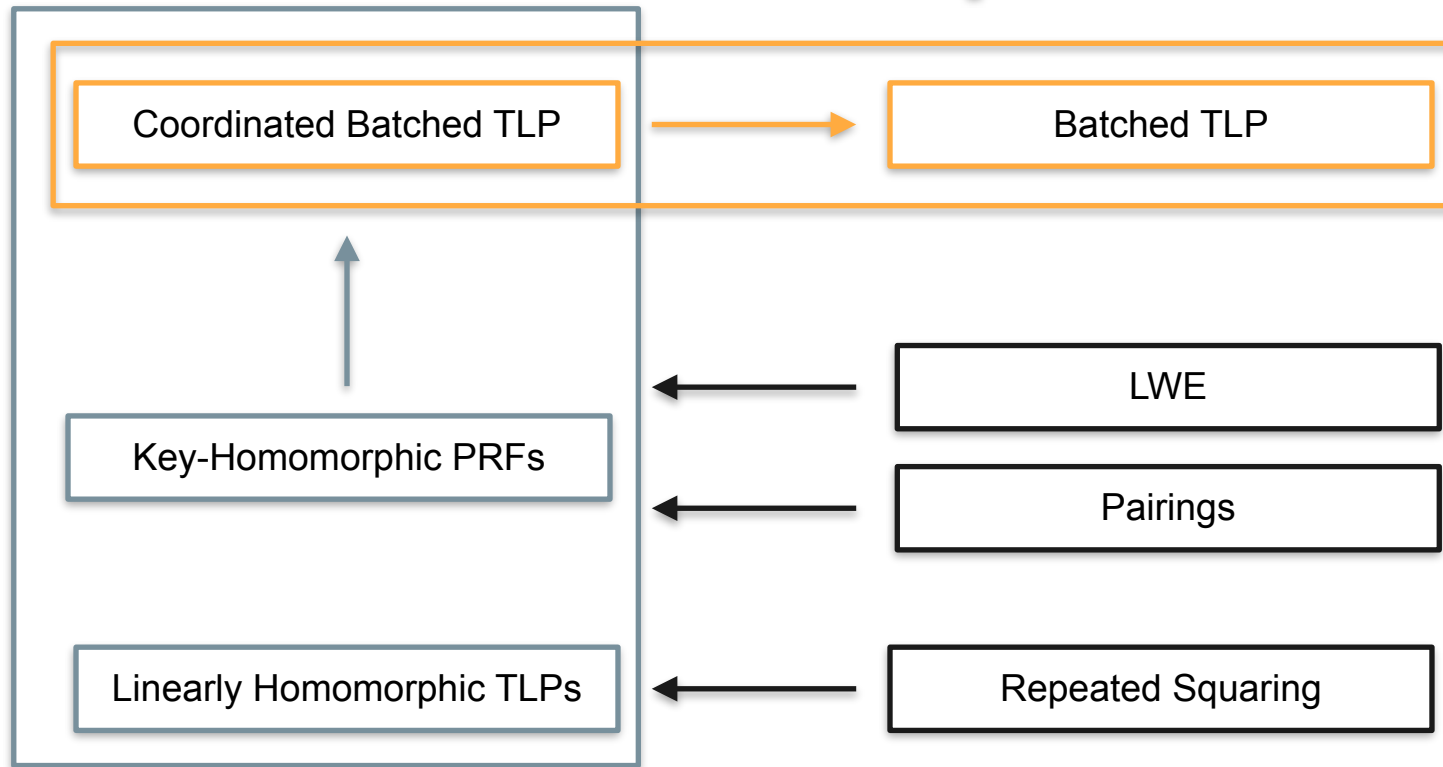
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Linearly Homomorphic TLP

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m_1



m_2

Linearly Homomorphic TLP



m_1



m_2

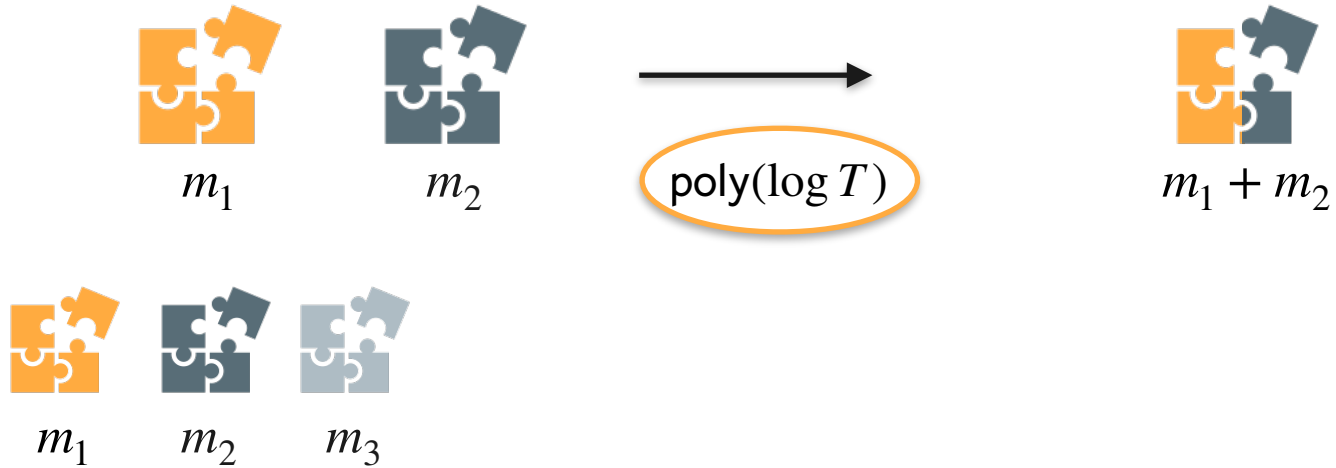


$\text{poly}(\log T)$

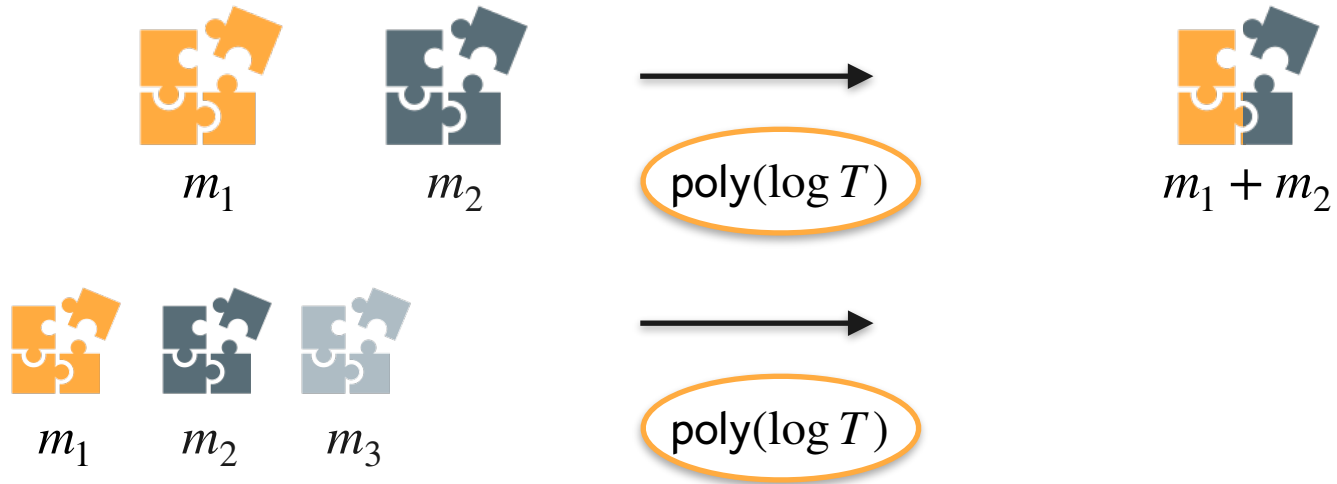


$m_1 + m_2$

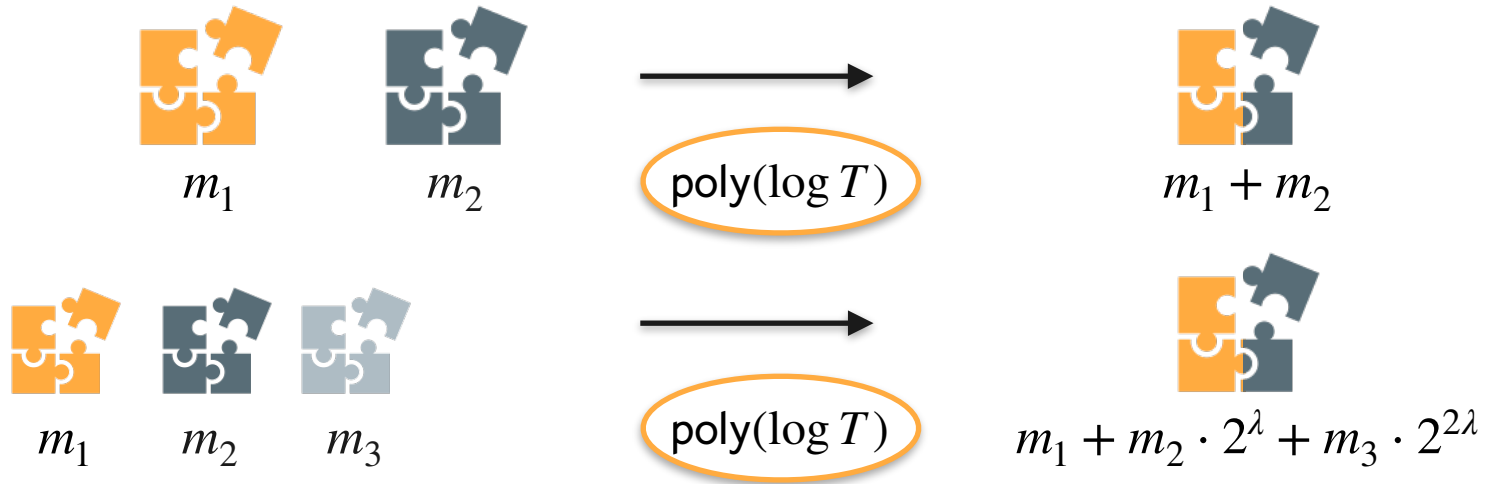
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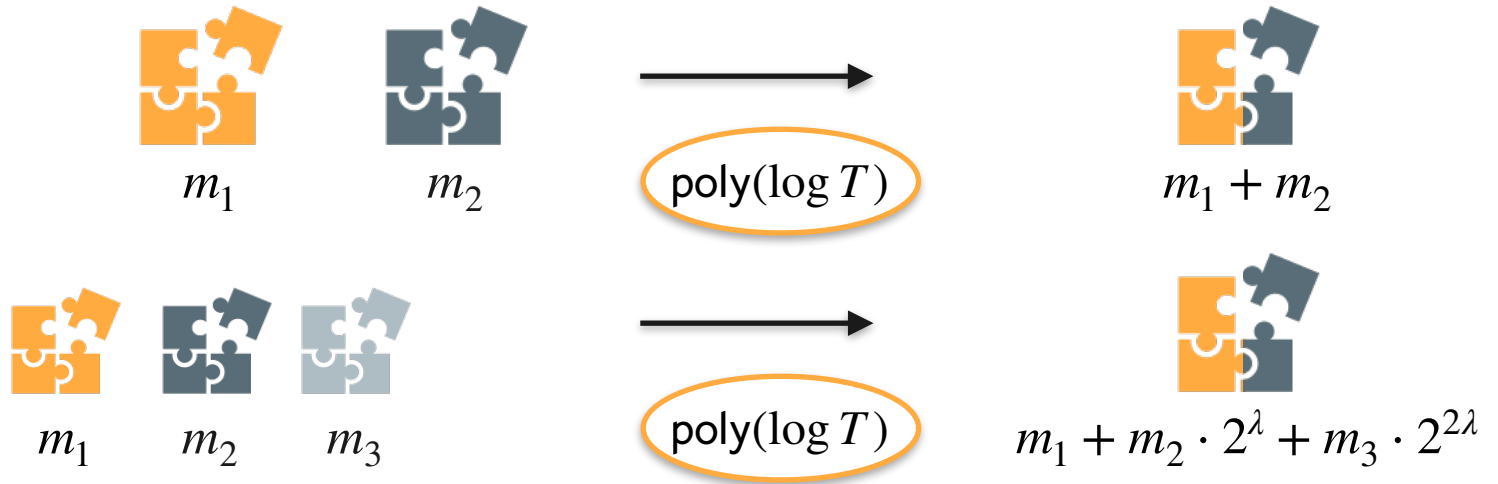
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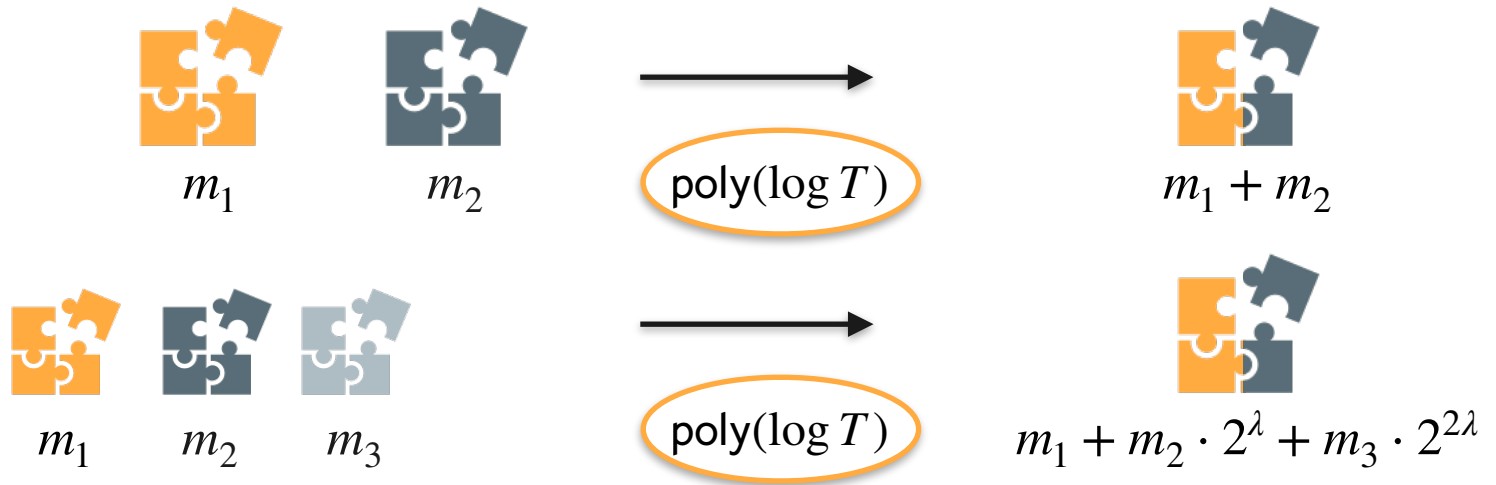


Linearly Homomorphic TLP



Bounded Batching only

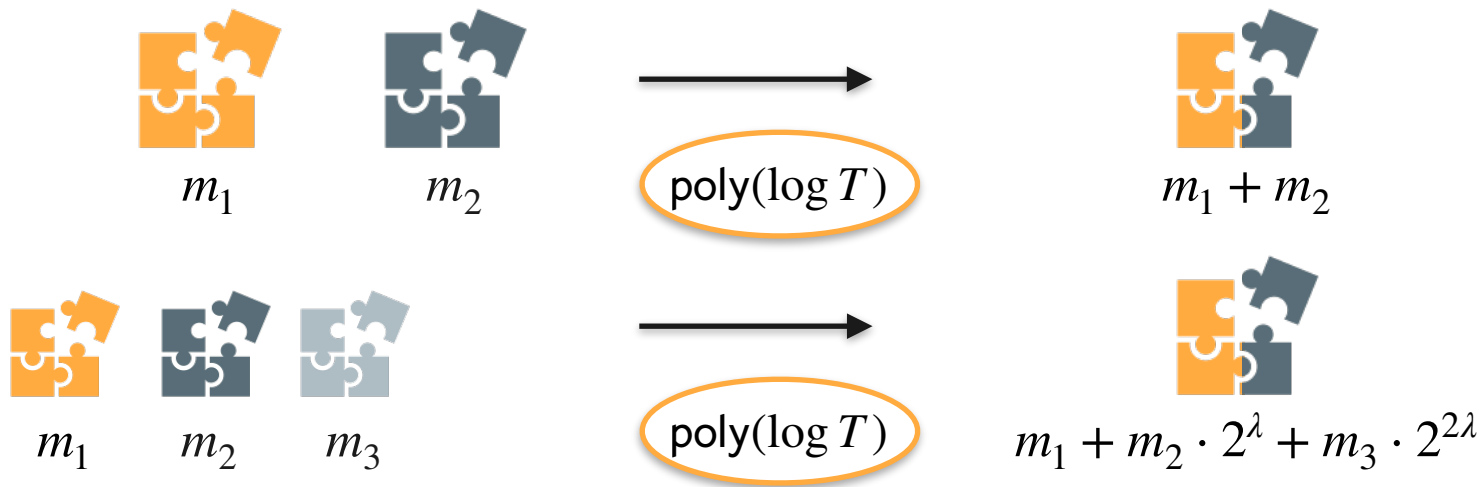
Linearly Homomorphic TLP



Bounded Batching only

Homomorphism over $\{0,1\}^{3\lambda}$

Linearly Homomorphic TLP



Bounded Batching only

Homomorphism over $\{0,1\}^{3\lambda}$



m_1



$O(N)$

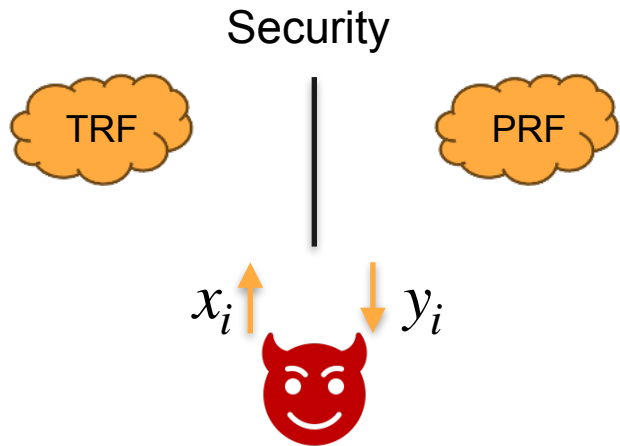
Key Homomorphic PRFs

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- PRF Setup - $\text{Setup}(1^\lambda) \rightarrow k$.
- PRF Evaluation - $\text{Eval}(k, x) \rightarrow y$.

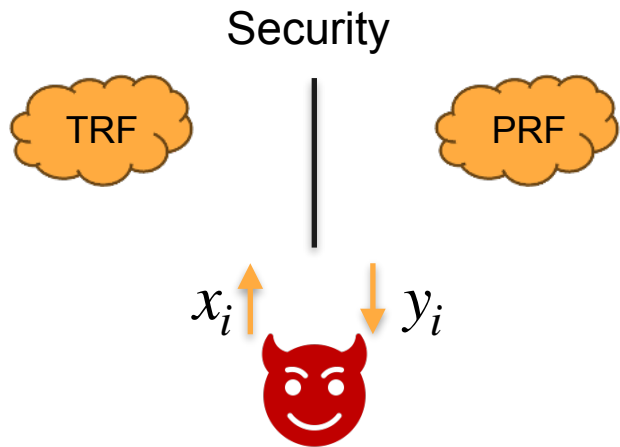
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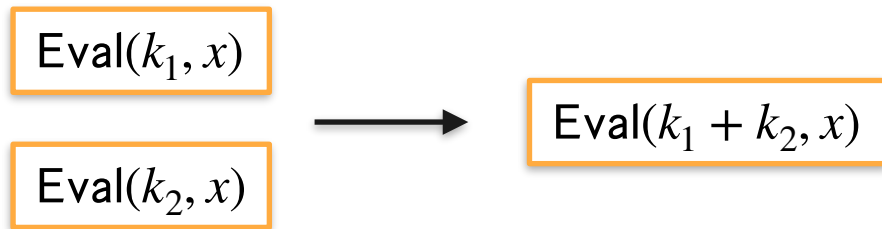
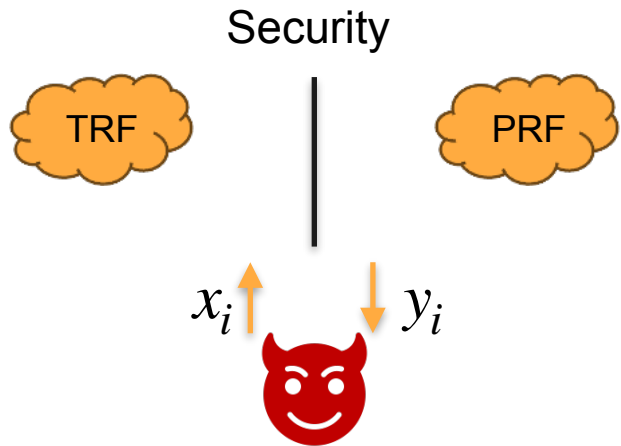


$\text{Eval}(k_1, x)$

$\text{Eval}(k_2, x)$

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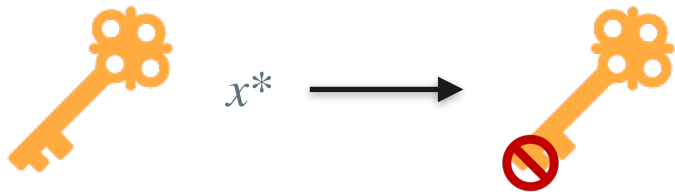
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$$\text{Eval}(\text{key}, x) = \text{Eval}(\text{key} \ominus x^*, x)$$

The equation shows that evaluating a PRF with a key and an input x yields the same result as evaluating it with a punctured key (indicated by a red circle with a slash) and the same input x .



Key Homomorphic PPRFs

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$$\text{Eval}(\text{key}, x) = \text{Eval}(\text{punctured key}, x)$$

$x \neq x^*$

Security

$$\text{Eval}(\text{key}, x^*) \mid \text{random}$$



Construction

Construction



Party i



m_i

Construction



Party i



m_i



k_i

$$\text{PRF_Eval}(k_i, i) + m_i$$

Construction



Party i



m_i



k_i

$$\text{PRF_Eval}(k_i, i) + m_i$$



Puncture k_i at i

Construction



Party i



m_i



k_i

$$\text{PRF_Eval}(k_i, i) + m_i$$



Puncture k_i at i



m_1



m_2



m_3

Construction



Party i



m_i



k_i

$$\text{PRF_Eval}(k_i, i) + m_i$$



Puncture k_i at i



m_1



m_2



m_3

$$\text{PRF_Eval}(k_1, 1) + m_1$$

Construction



Party i



m_i



k_i

$$\text{PRF_Eval}(k_i, i) + m_i$$



Puncture k_i at i



m_1



m_2



m_3



k_1



k_2



k_3

$$\text{PRF_Eval}(k_1, 1) + m_1$$

Construction



Party i



m_i



k_i

$\text{PRF_Eval}(k_i, i) + m_i$



Puncture k_i at i



m_1



m_2



m_3



k_1



k_2



k_3



$k_1 + k_2 + k_3$

$\text{PRF_Eval}(k_1, 1) + m_1$

Construction



Party i



m_i



k_i

$\text{PRF_Eval}(k_i, i) + m_i$



Puncture k_i at i



m_1



m_2



m_3

$\text{PRF_Eval}(k_1, 1) + m_1$



k_1



k_2



k_3



$k_1 + k_2 + k_3$



Takes time T

Construction



Party i



m_i



k_i

$\text{PRF_Eval}(k_i, i) + m_i$



Puncture k_i at i



m_1



m_2



m_3



k_1



k_2



k_3



$k_1 + k_2 + k_3$

$\text{PRF_Eval}(k_1, 1) + m_1$ $-$ $\text{Eval}(k_1 + k_2 + k_3, 1)$



Takes time T

Construction



Party i



m_i



k_i

$\text{PRF_Eval}(k_i, i) + m_i$



Puncture k_i at i



m_1



m_2



m_3



k_1



k_2



k_3



$k_1 + k_2 + k_3$

$\text{PRF_Eval}(k_1, 1) + m_1$

- $\text{Eval}(k_1 + k_2 + k_3, 1)$

+



Takes time T

Construction



Party i



m_i



k_i

$\text{PRF_Eval}(k_i, i) + m_i$



Puncture k_i at i



m_1



m_2



m_3



k_1



k_2



k_3



$k_1 + k_2 + k_3$

$\text{PRF_Eval}(k_1, 1) + m_1$

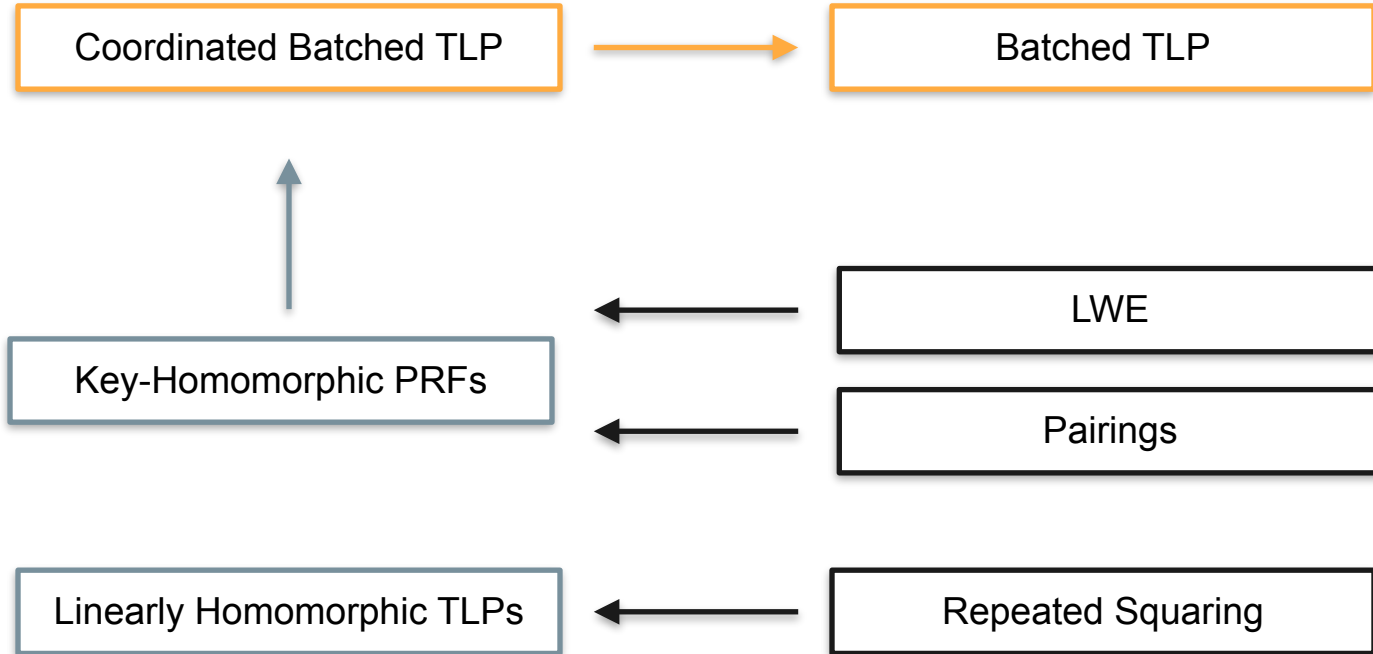
- $\text{Eval}(k_1 + k_2 + k_3, 1)$

+ $\text{Eval}(\text{key with red prohibition}, 1) + \text{Eval}(\text{key with red prohibition}, 1)$

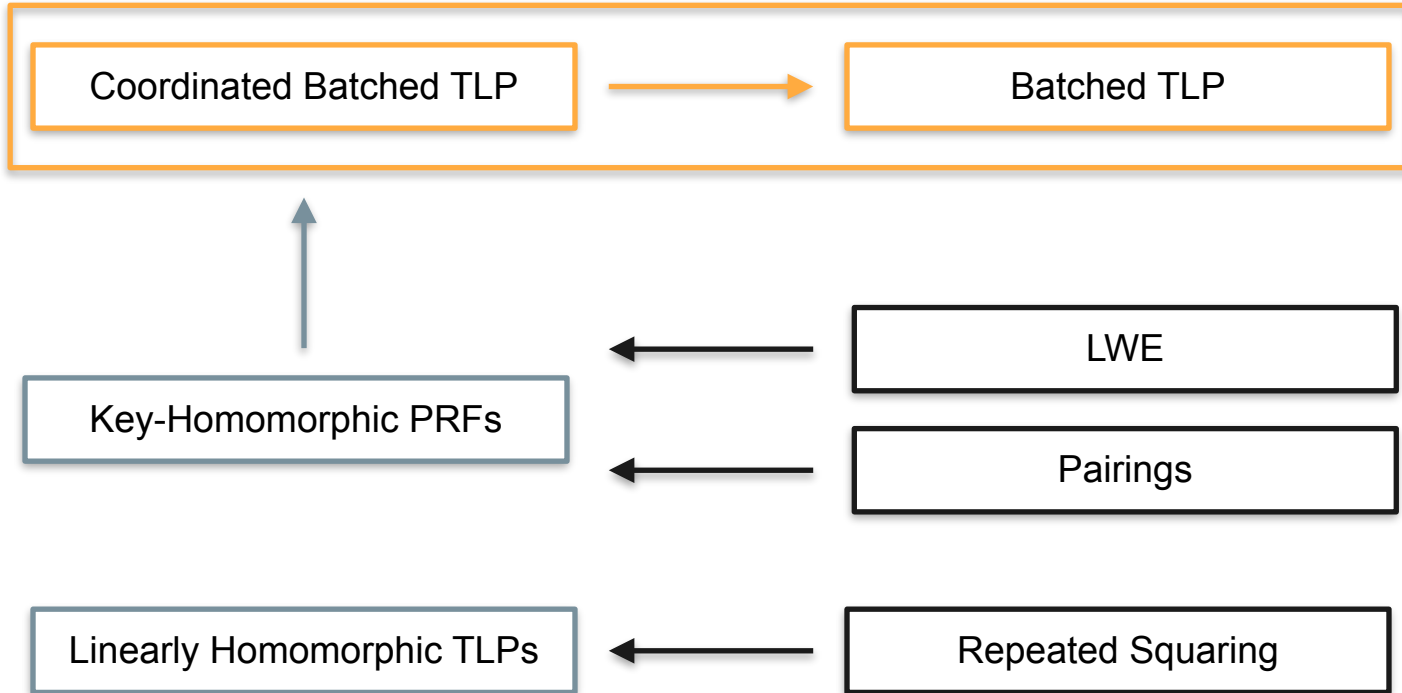


Takes time T

Roadmap



Roadmap



Transformation

Transformation

M - 5 - number of users

N - 3 - Batch to at-most 3 puzzles

D - 2 - degree

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Slots



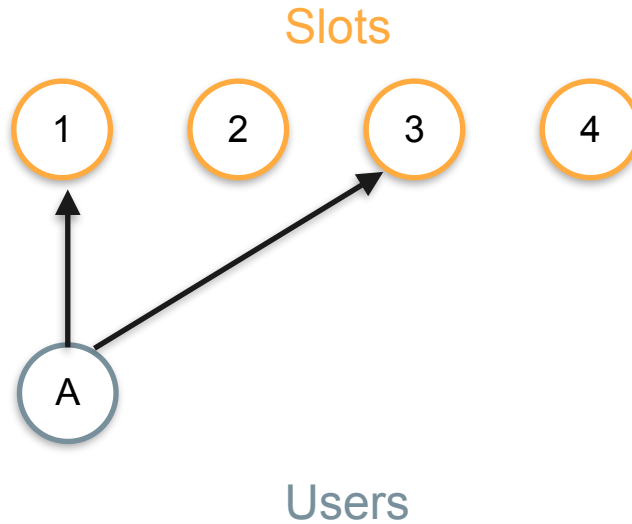
Setup($1^\lambda, T, 1^4$) \rightarrow pp

Users

Transformation

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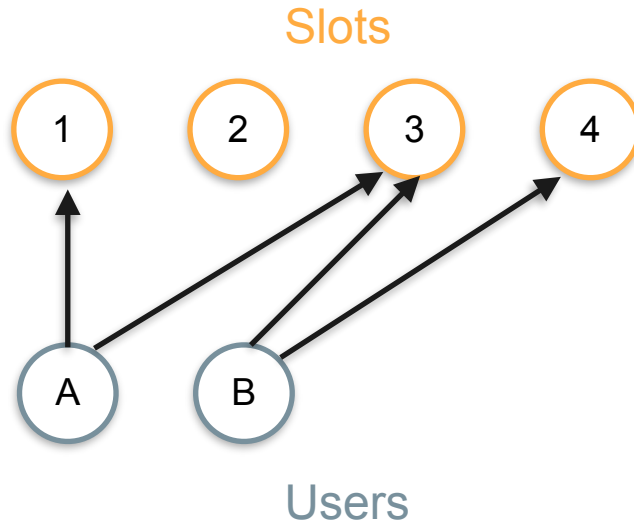
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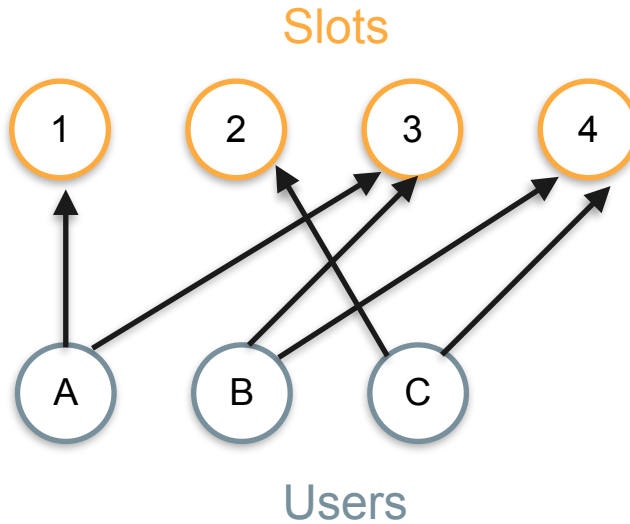
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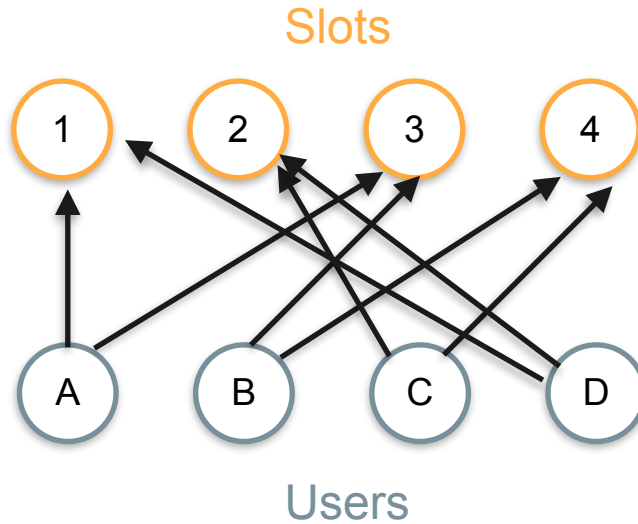
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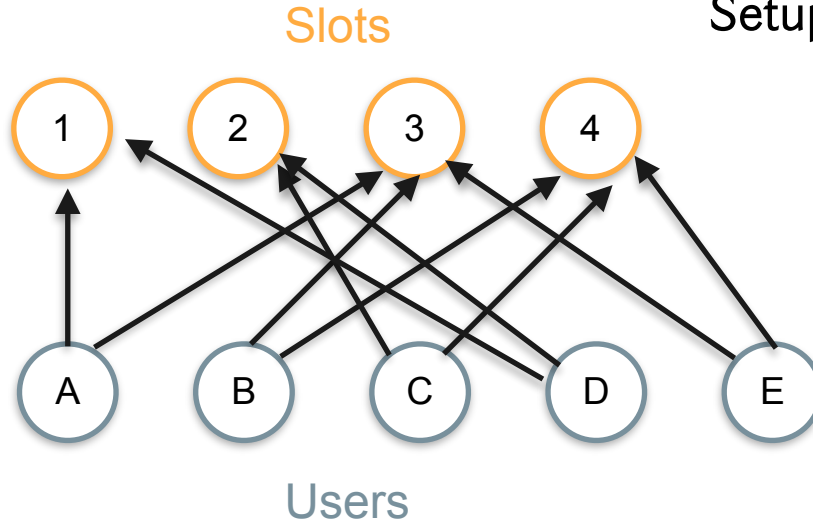
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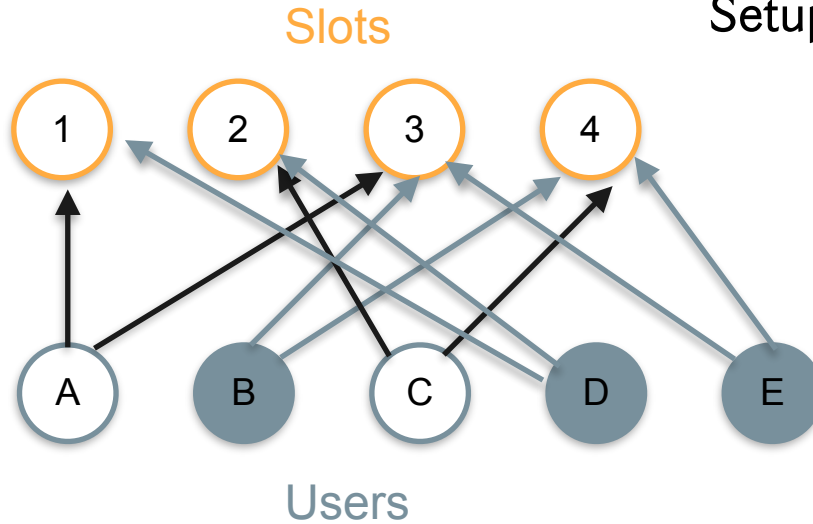
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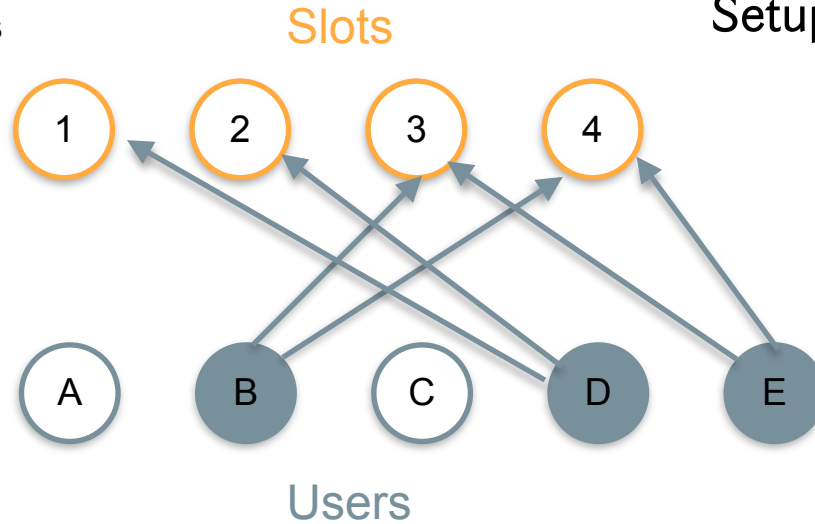
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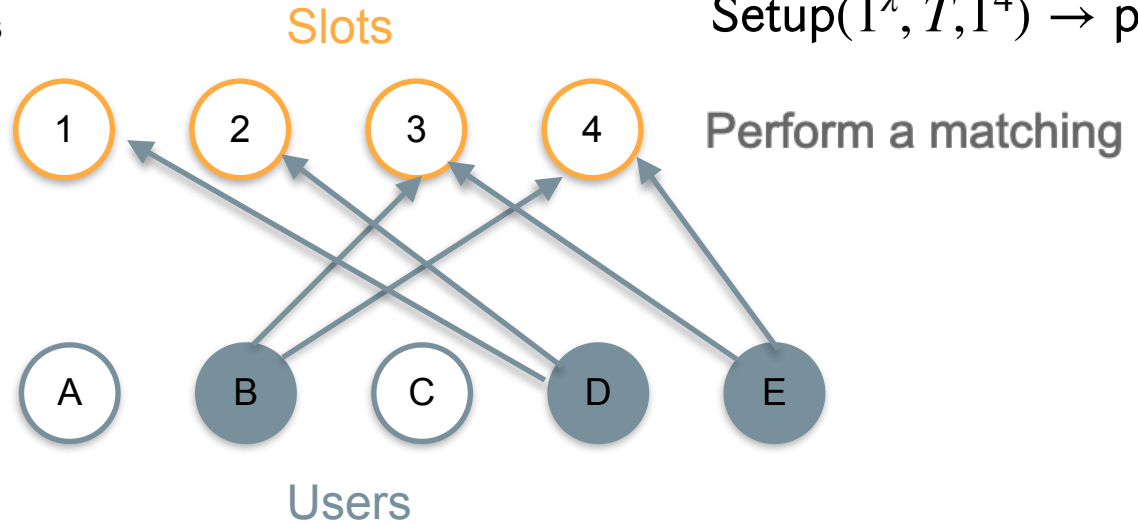
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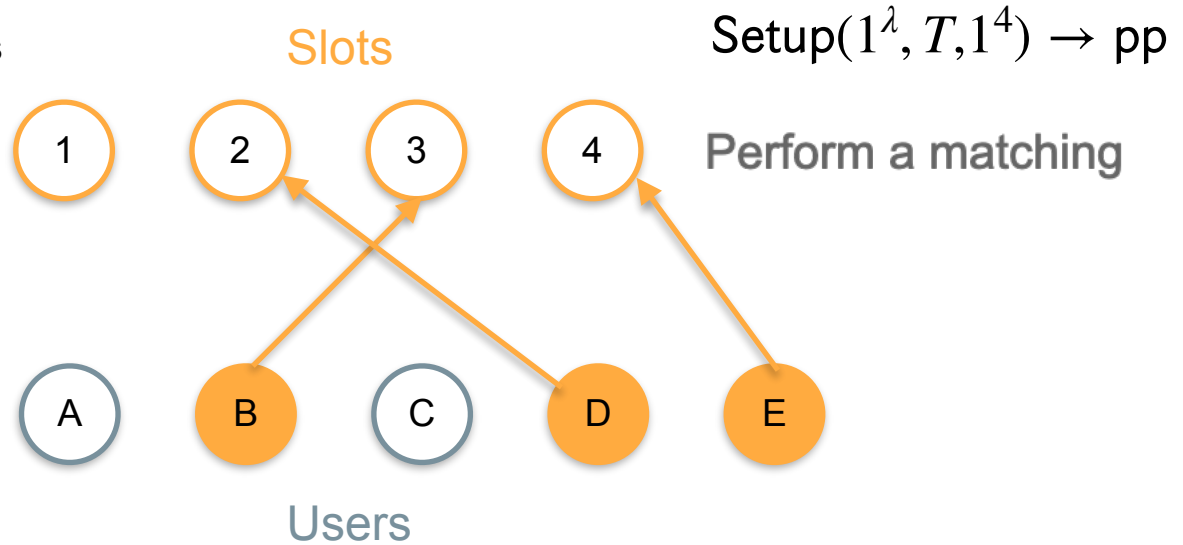
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$\text{Setup}(1^\lambda, T, 1^4) \rightarrow \text{pp}$



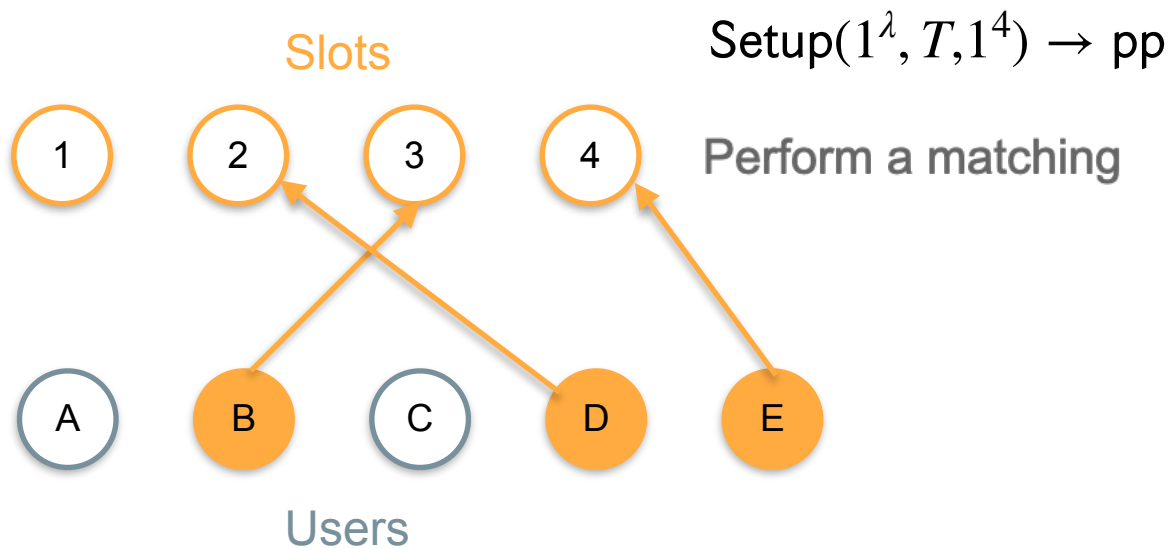
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Transformation

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D - 2 - degree



Theorem (informal) - Set slots $\geq 2e \cdot N$, degree $\geq \frac{\omega(\log \lambda)}{\log N}$, then we can build an uncoordinated batch TLP.

Rogue puzzles

Rogue puzzles



m_1



m_2



m_3



Rogue puzzles



m_1



m_2



m_3



BatchSolve



Rogue puzzles



m_1



m_2



m_3



BatchSolve

100



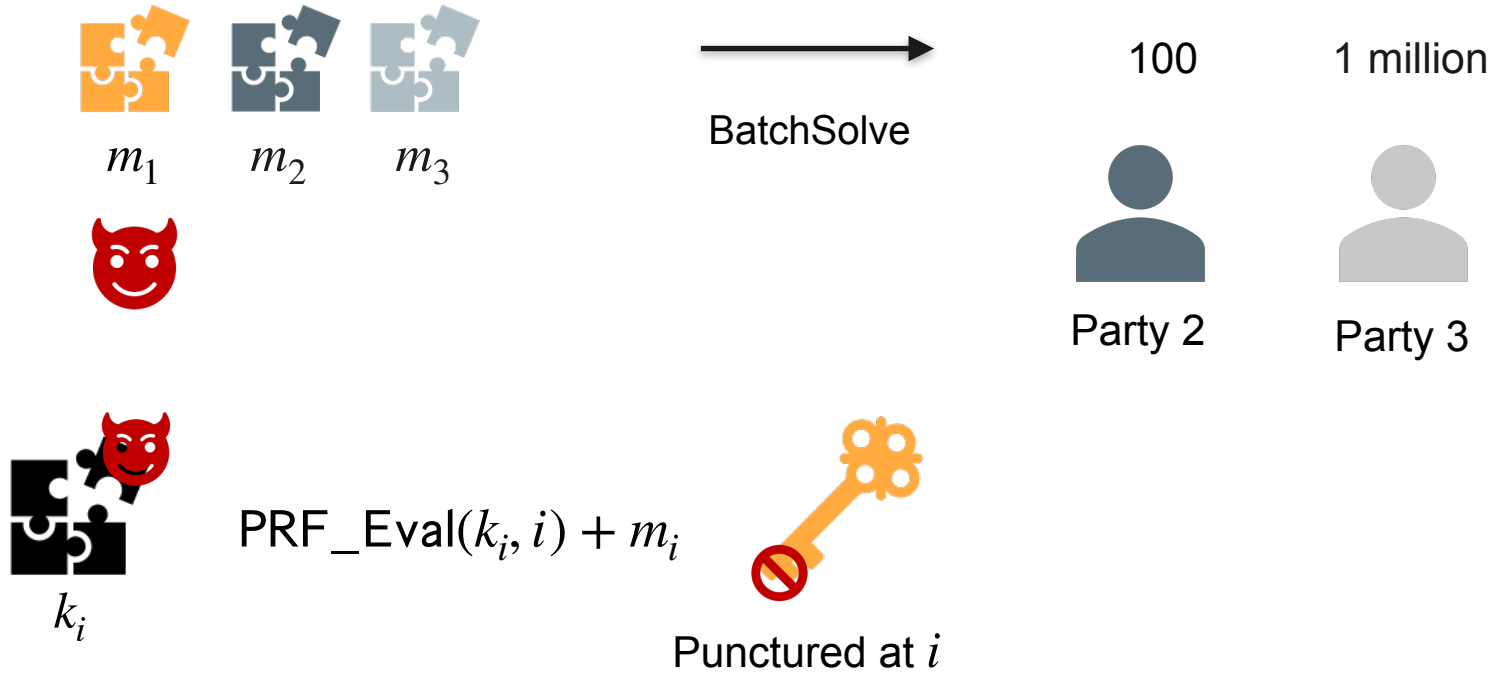
Party 2

1 million

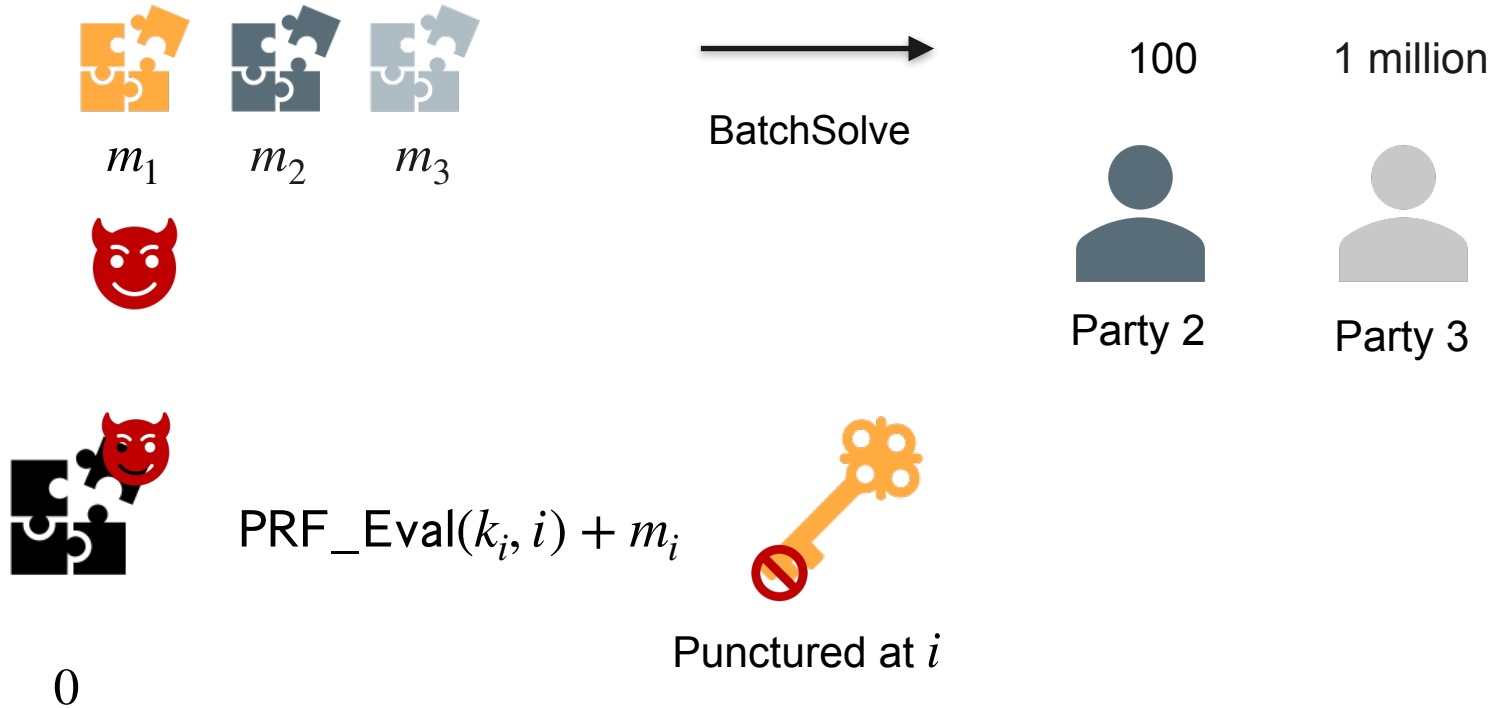


Party 3

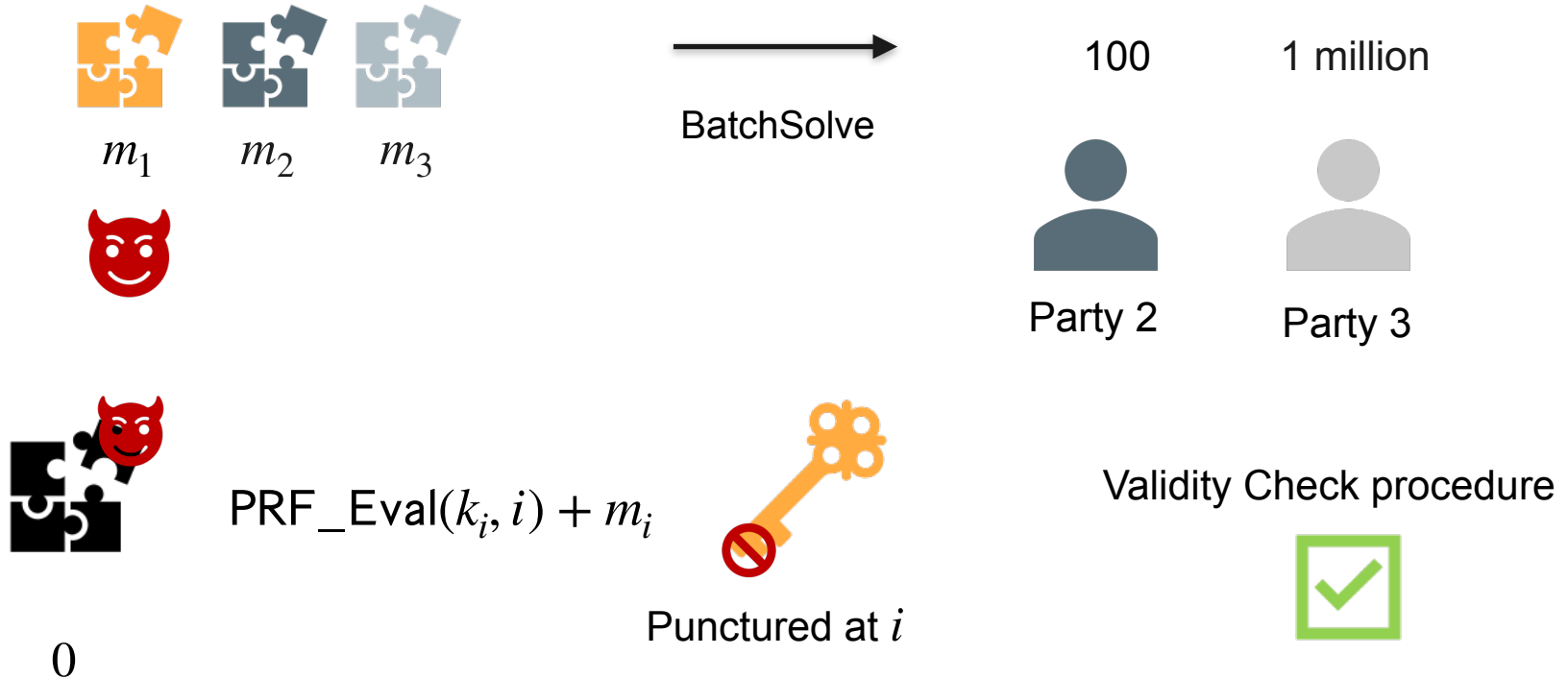
Rogue puzzles



Rogue puzzles



Rogue puzzles



Prototype Evaluation

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- For $T = 50$ million sequential computations*, and batching 500 puzzles, the batching time trivially would take 15 hours, while our solution takes close to 6 minutes (we did not use any parallelism for our experiments).

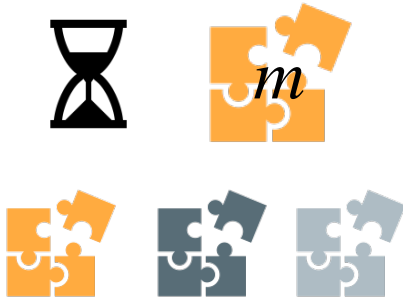
*the time to do 50 million sequential computations on the test machine is 5 minutes

Prototype Evaluation

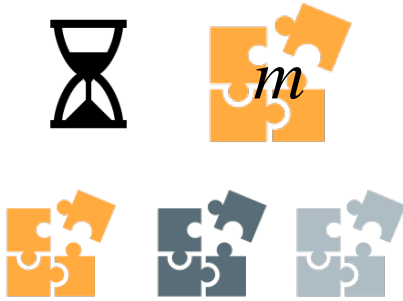
- For $T = 50$ million sequential computations*, and batching 500 puzzles, the batching time trivially would take 15 hours, while our solution takes close to 6 minutes (we did not use any parallelism for our experiments).
- For $T = 50$ million computations, and batching 7000 puzzles, the size of a single puzzle is 8 MB trivially, 37 MB using our solution and would be 790 MB using the linearly homomorphic solution.

*the time to do 50 million sequential computations on the test machine is 5 minutes

Conclusion!

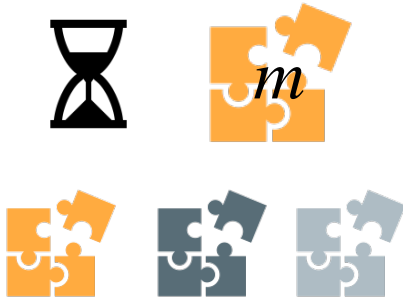


Conclusion!



- We gave a solution template for batch solving of time-lock puzzles.

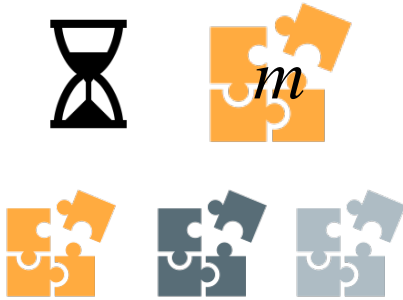
Conclusion!



- We gave a solution template for batch solving of time-lock puzzles.
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Conclusion!



- We gave a solution template for batch solving of time-lock puzzles.
- Introduction of rogue puzzle attacks.
- Give a concrete implementation and numbers.

