

Two-Round Maliciously-Secure Oblivious Transfer with Optimal Rate

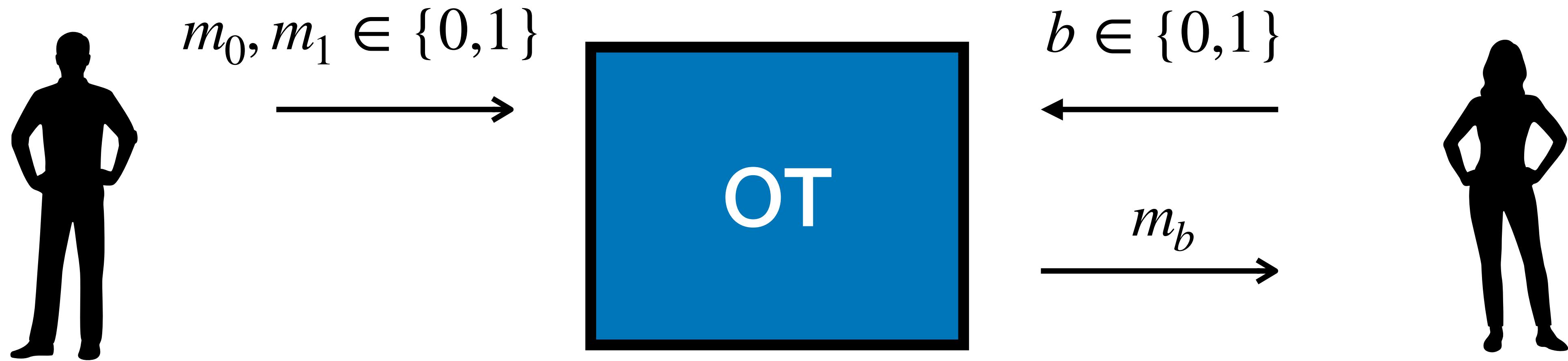
Pedro Branco *Max-Planck Institute for Security and Privacy*

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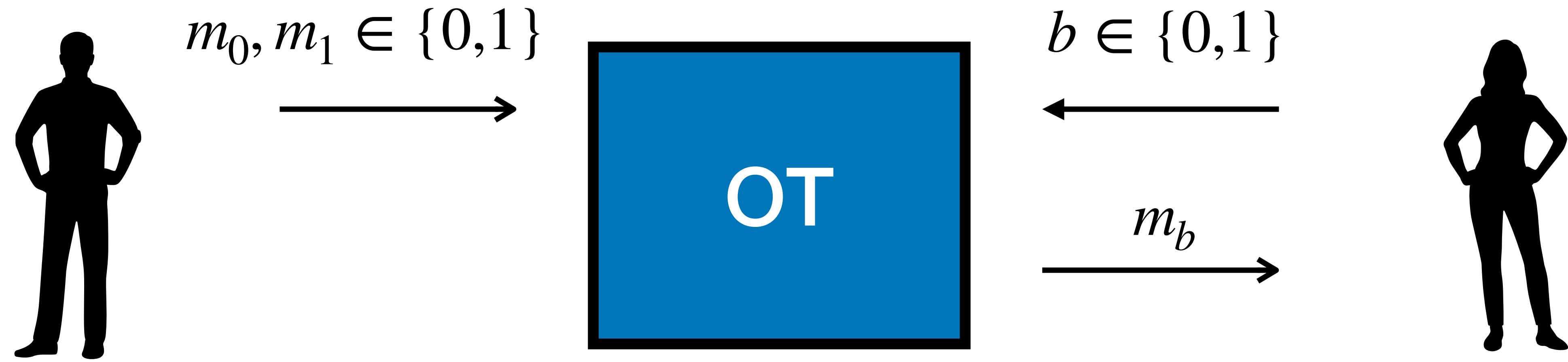
Akshayaram Srinivasan *University of Toronto*



Oblivious Transfer



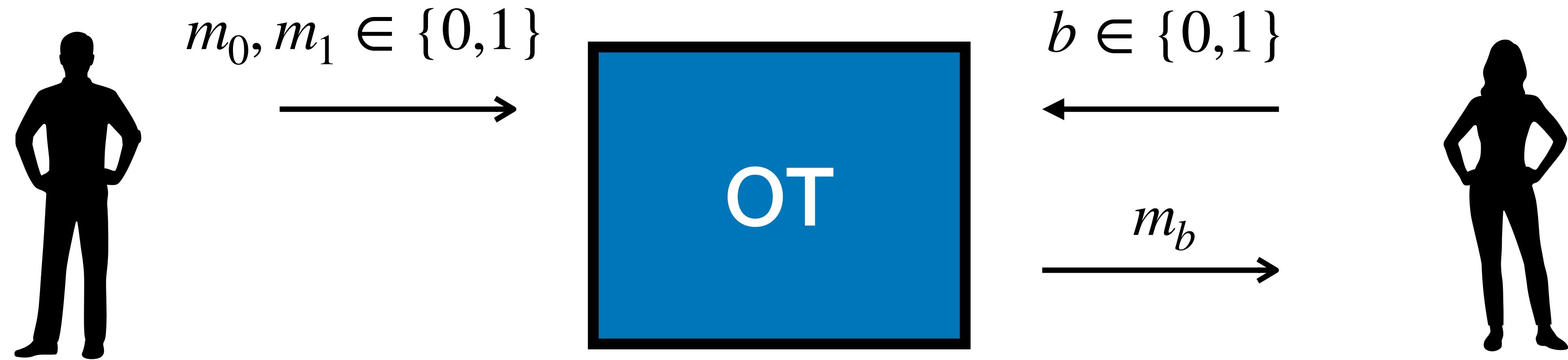
Oblivious Transfer



Receiver security: b is hidden from the sender

Sender security: m_{1-b} hidden from the receiver

Oblivious Transfer



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Main Application: OT is complete for 2PC/MPC

How much interaction needed for OT?

Rounds

≥ 2

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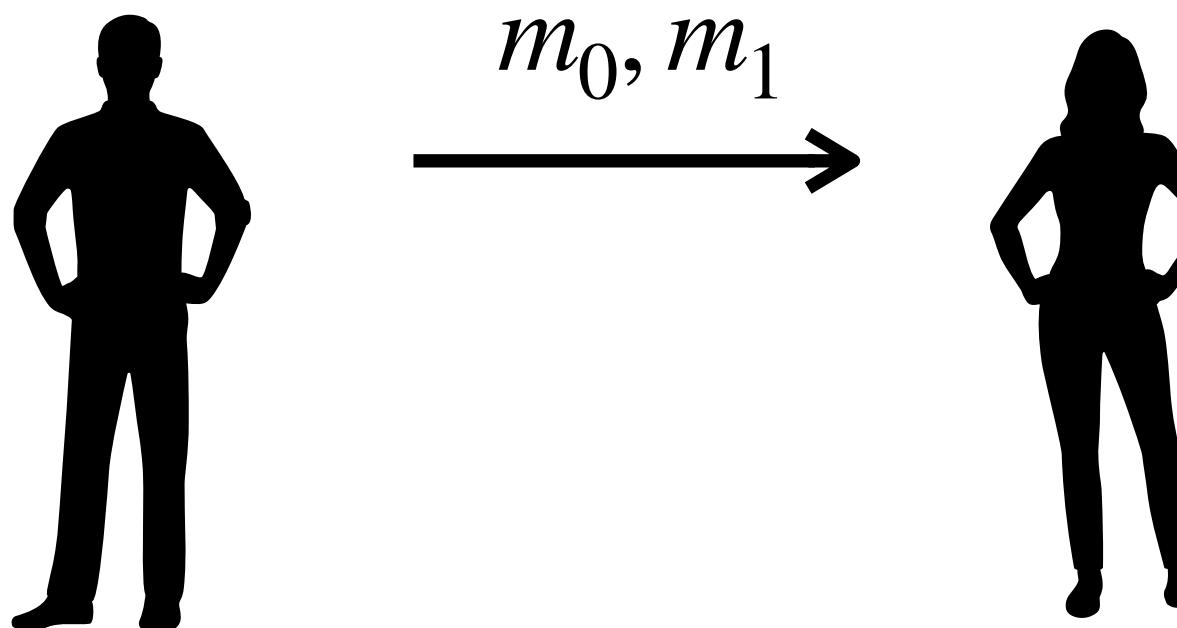
Communication complexity
(#bits exchanged)

≥ 2 bits per OT

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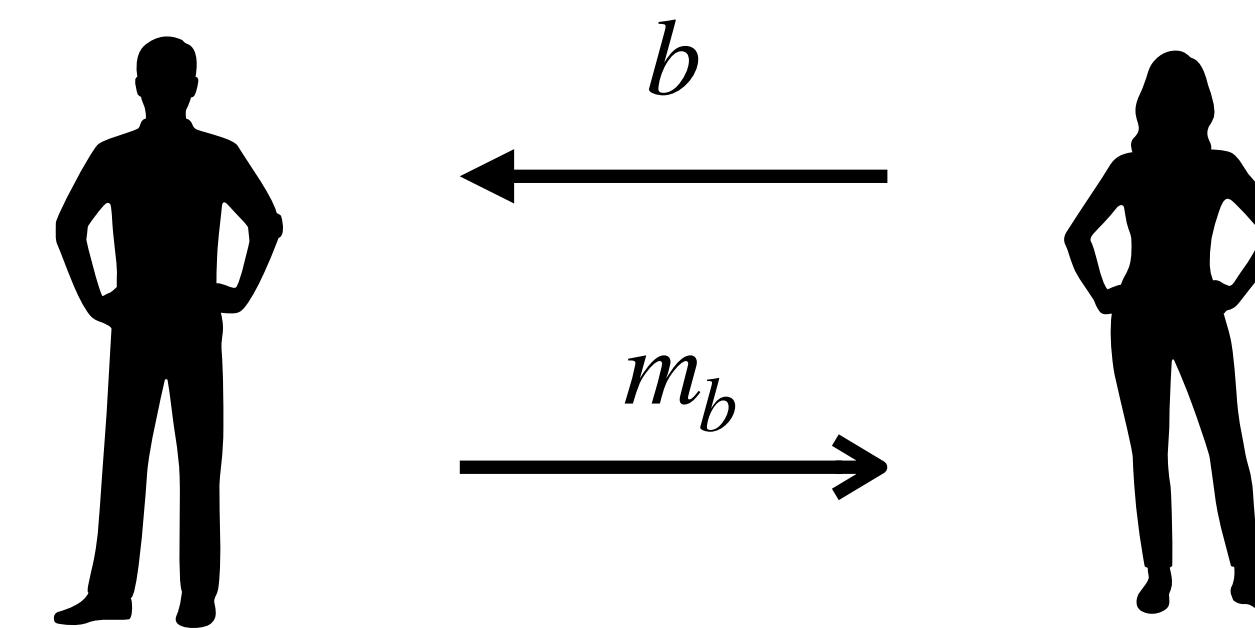
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Previous works: [GH19, BDGM19, BBDP22, BDS23]

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Security: Semi-honest or one-sided malicious (e.g. SSP)

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Malicious Security?

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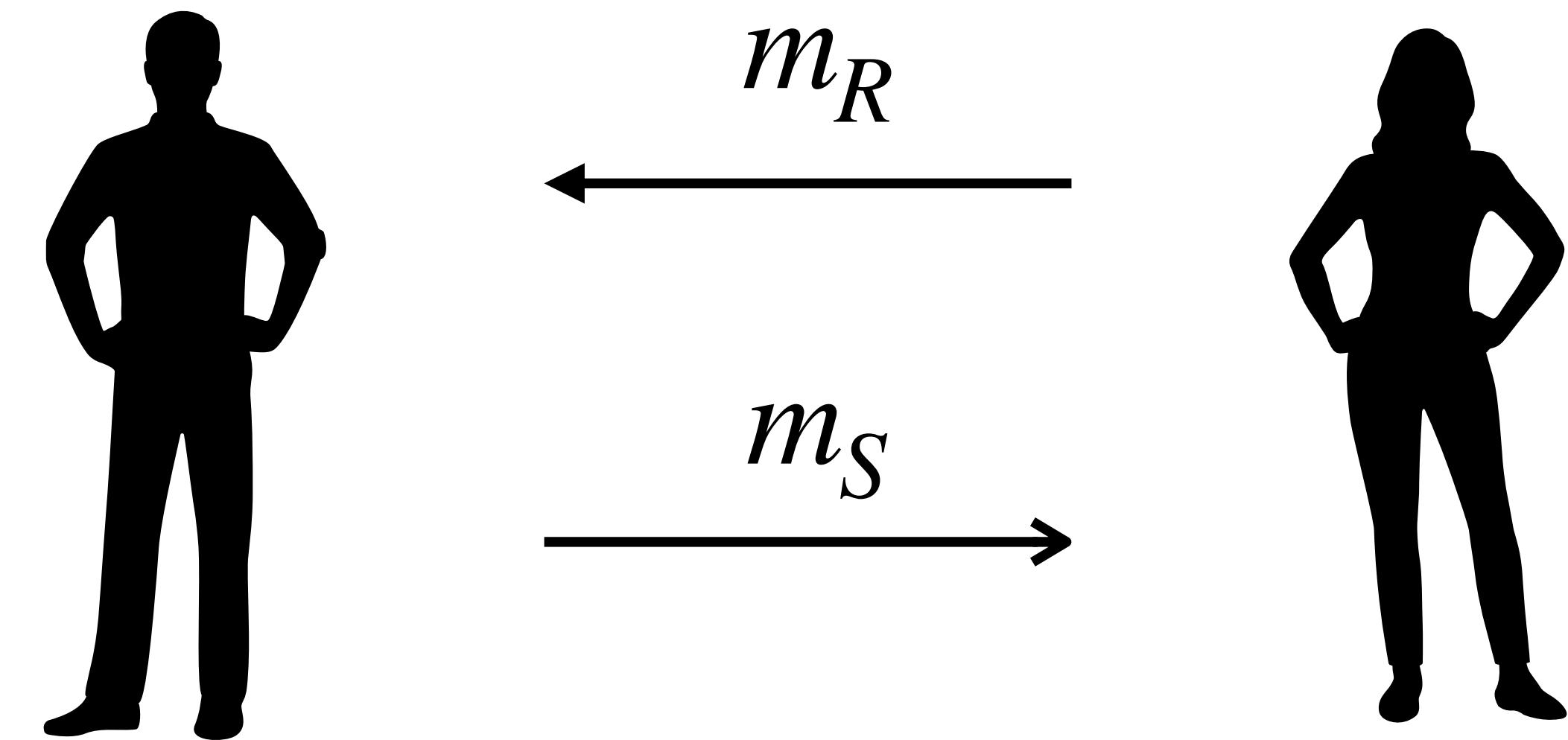
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Malicious	$\geq 2^*$	$\geq 3^{**}$

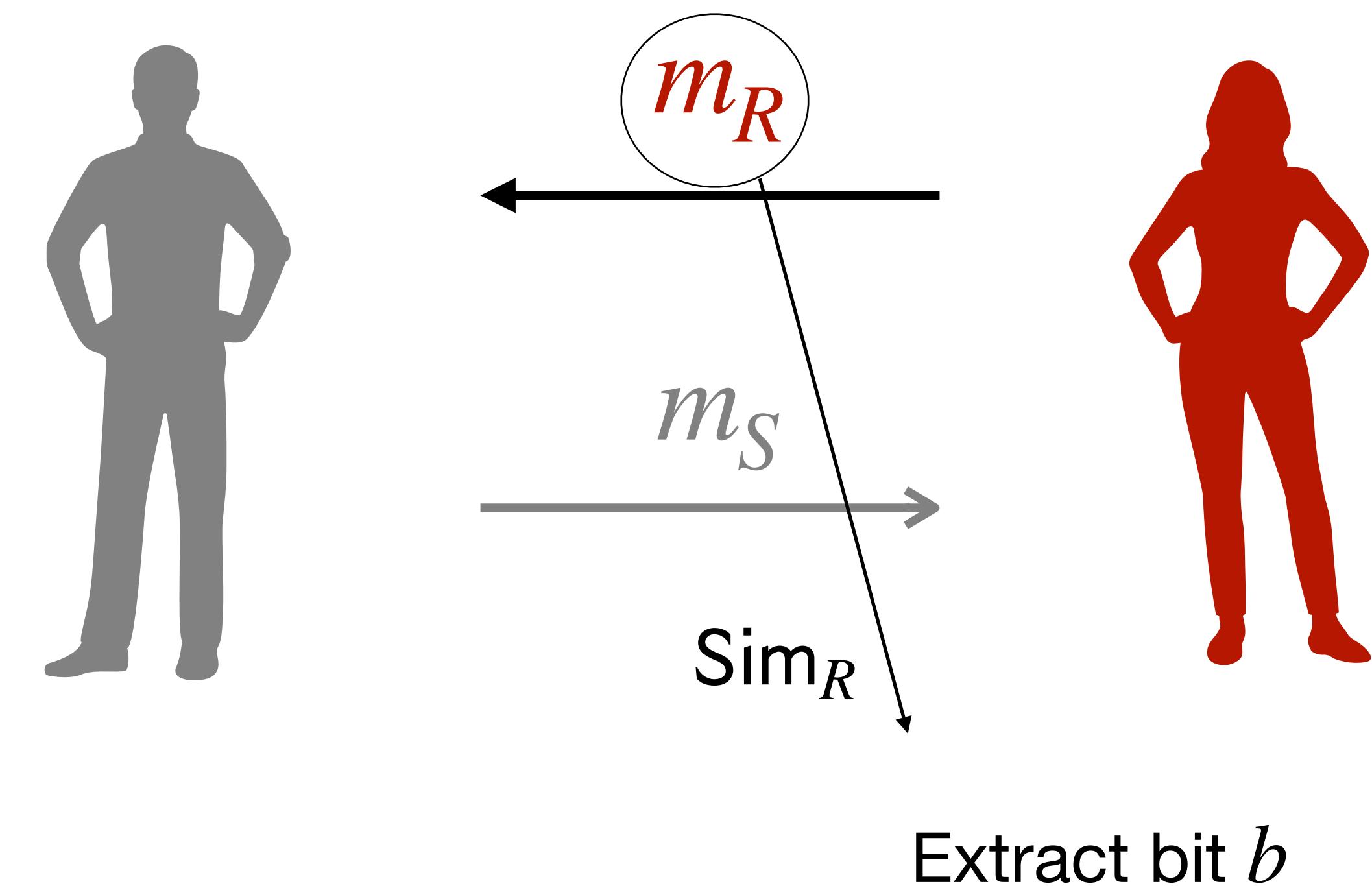
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** In two rounds

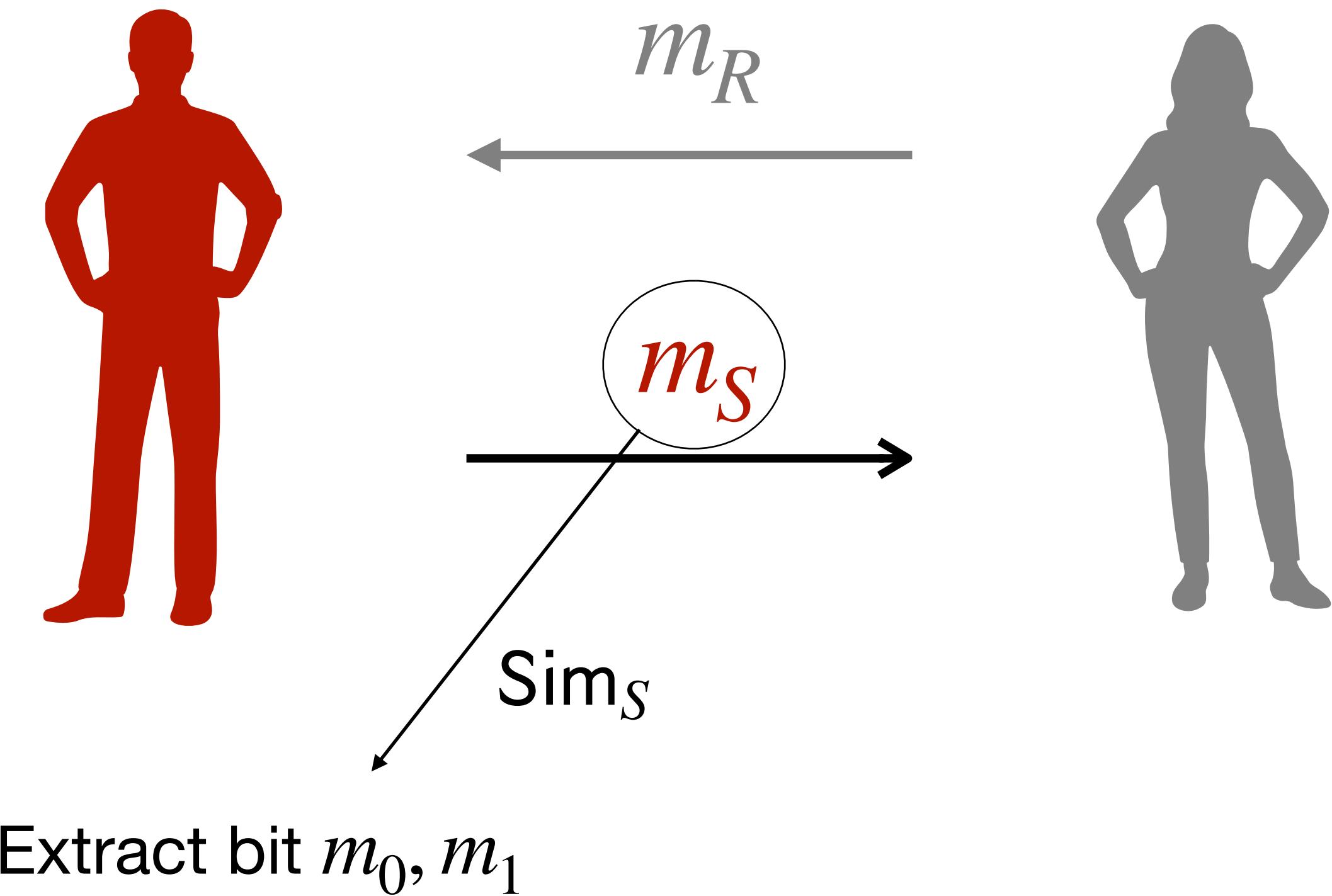
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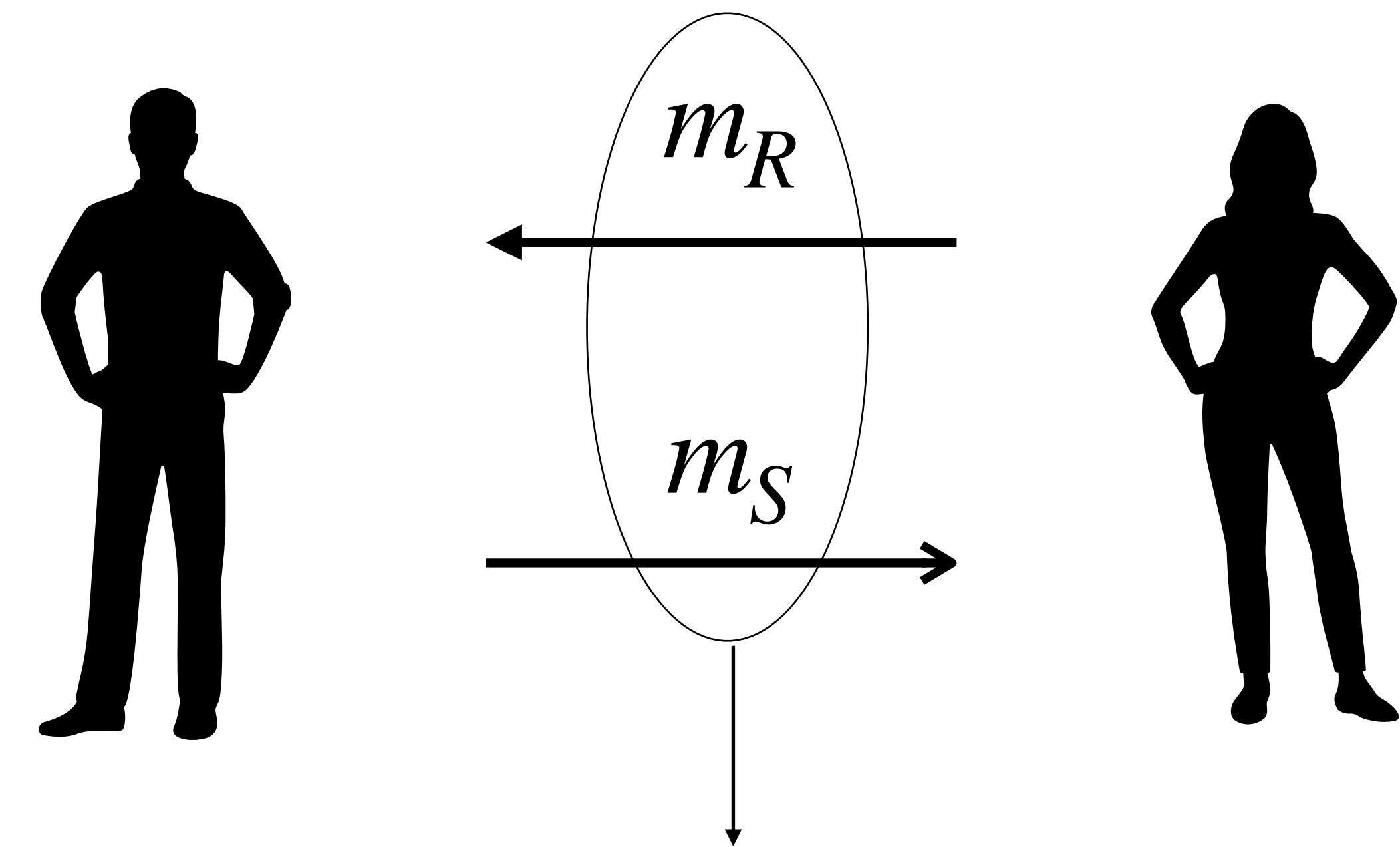
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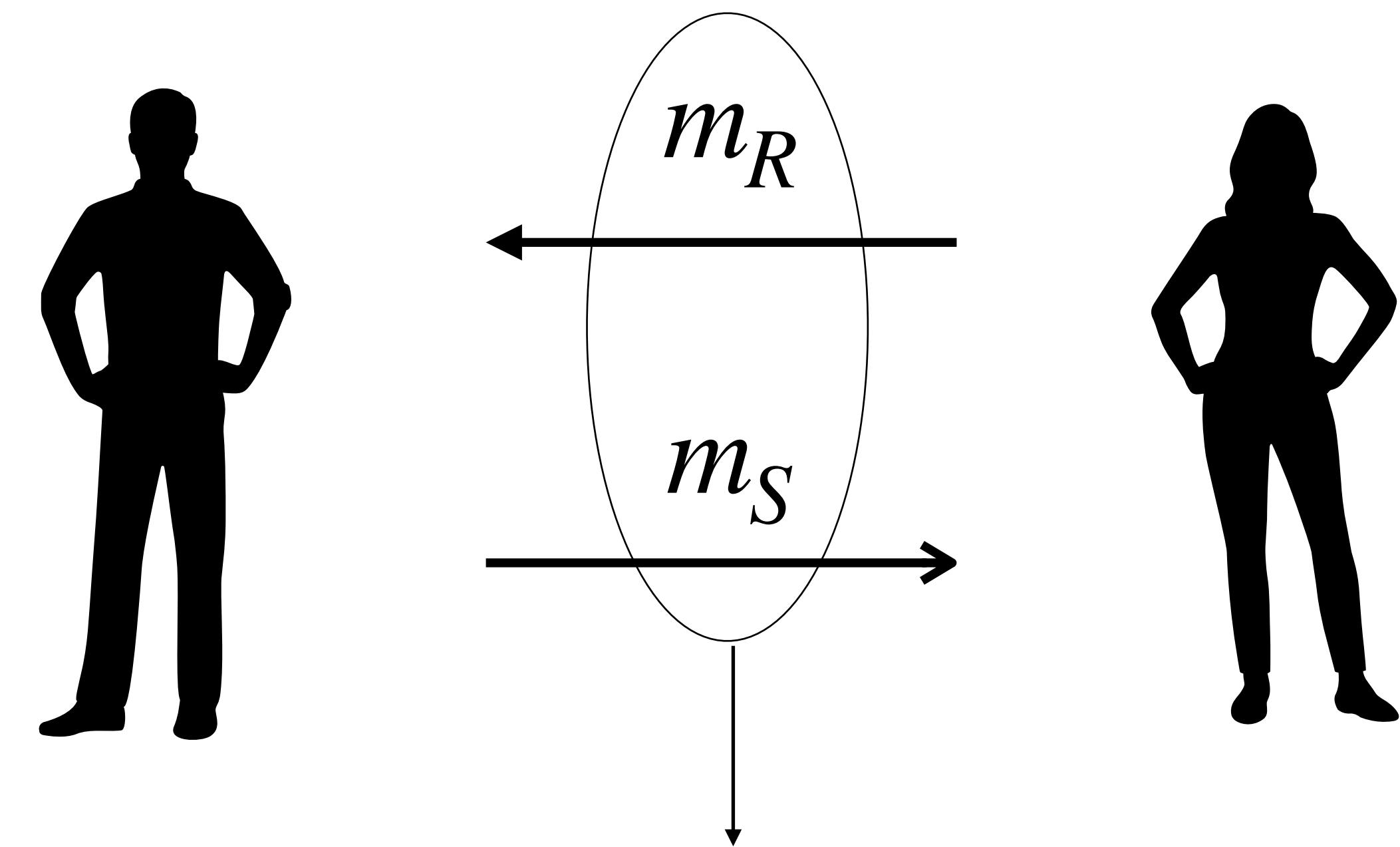


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If communication is ≈ 3 then it's **optimal rate**

Malicious OT schemes with Optimal Rate

Malicious OT with optimal rate in two rounds?

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Malicious OT with **optimal rate in two rounds?**

- **LWE** via spooky FHE [DHRW16]
- **QR/DCR with long CRS** via PCGs [OSY21]
- **LPN and Random Oracle** via silent OT extensions [BCG+19]

Malicious OT schemes with Optimal Rate

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Malicious two-round OT with optimal rate with a short CRS?

Our Results

Our Result: Two-round batch-OT scheme:

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- **UC-secure against malicious adversaries under QR+LPN.**
- **Short and reusable CRS**
- **Optimal rate:** Total communication of $3k + o(k) \cdot \text{poly}(\lambda)$ for a batch k

Blueprint of our Construction

OT from TDH “a la PVW”

Sender's message:
optimal size

Receiver's message:
large

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Not Sender Secure!!

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Hybrid Encryption
[BDP22,BDS23]

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$\text{Enc}(ek, x)$

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Perfect correctness

OT from TDH for linear functions

CRS: $hk, t \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

Receiver (b):

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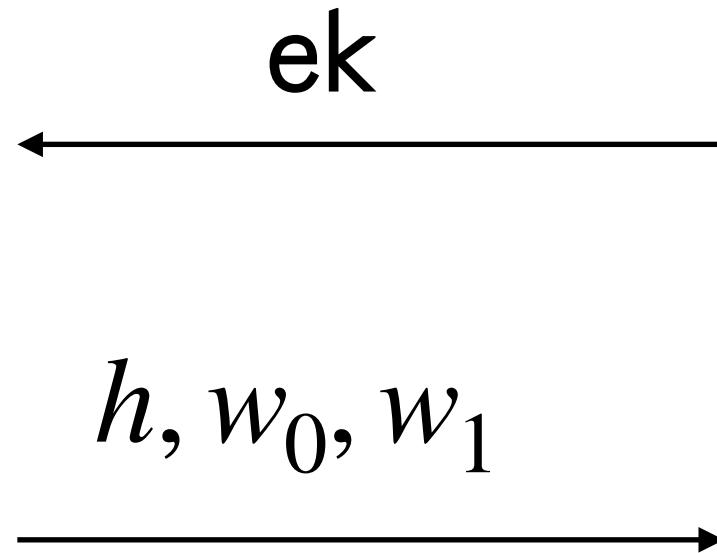
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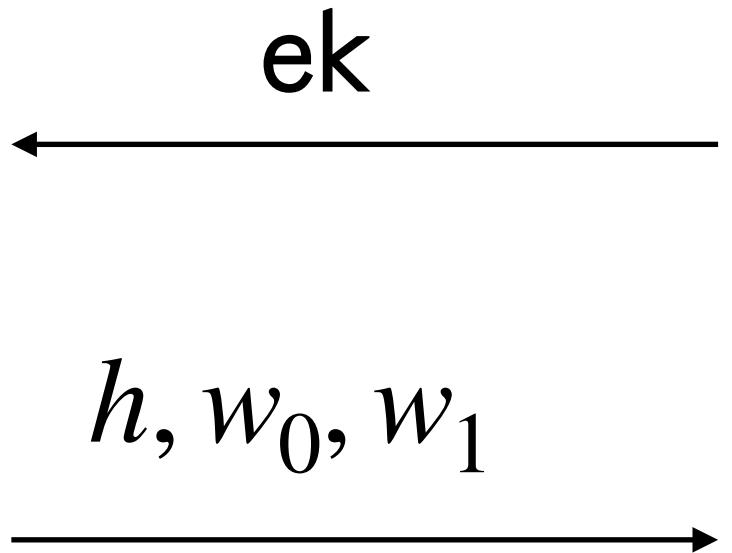
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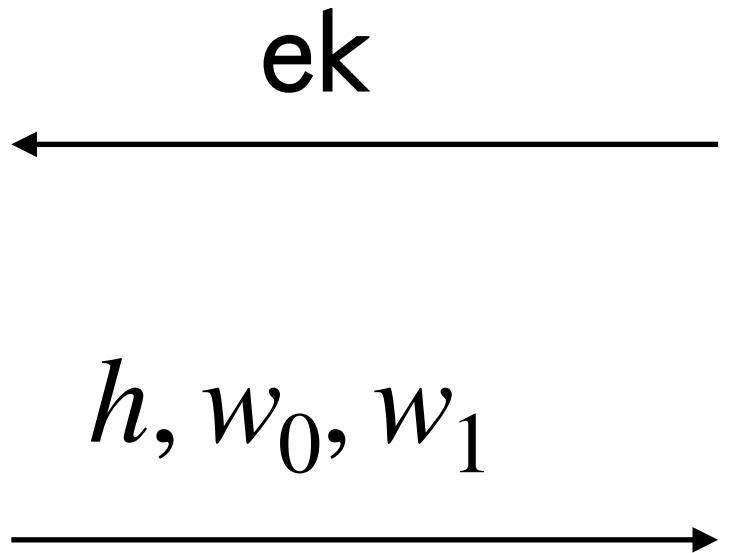
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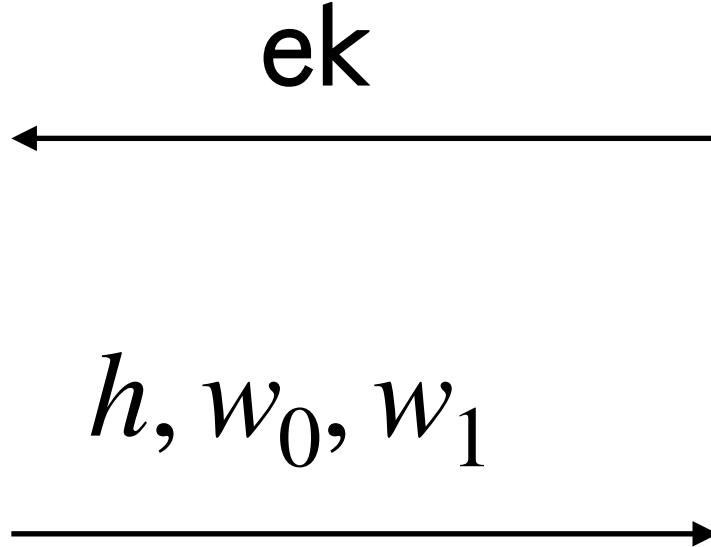
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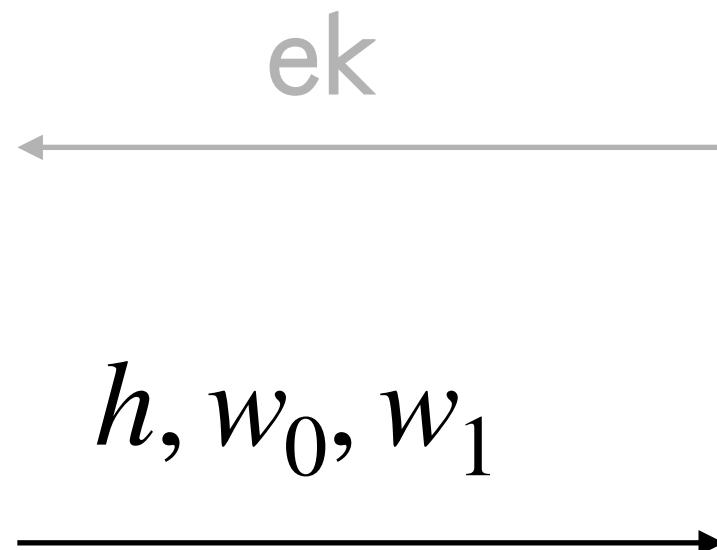
- $\text{Enc}(ek, r) \oplus \text{Dec}(h, td) = b \cdot \langle t, r \rangle$
- **Correct:** For w_b , the shift $\langle t, r \rangle$ cancels out.
- **Secure:** For w_{1-b} , the shift $\langle t, r \rangle$ hides μ_{1-b} (by LHL).

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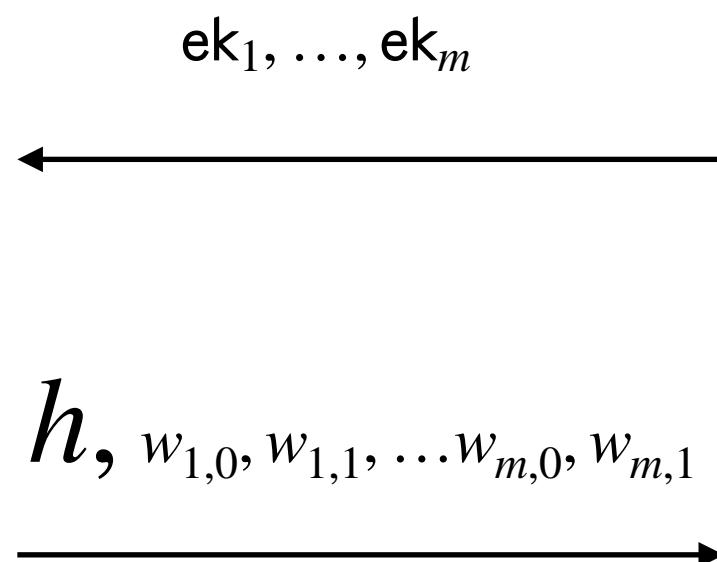
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OT from TDH for linear functions

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Receiver (b_1, \dots, b_m) :

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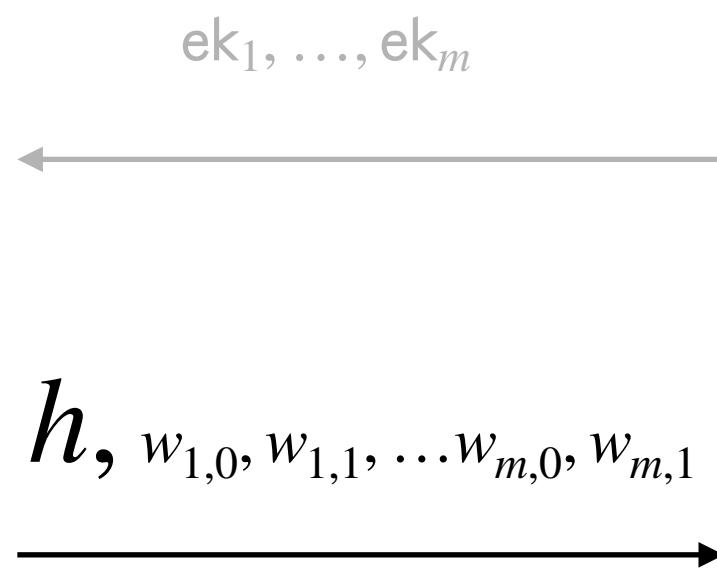
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Problem!

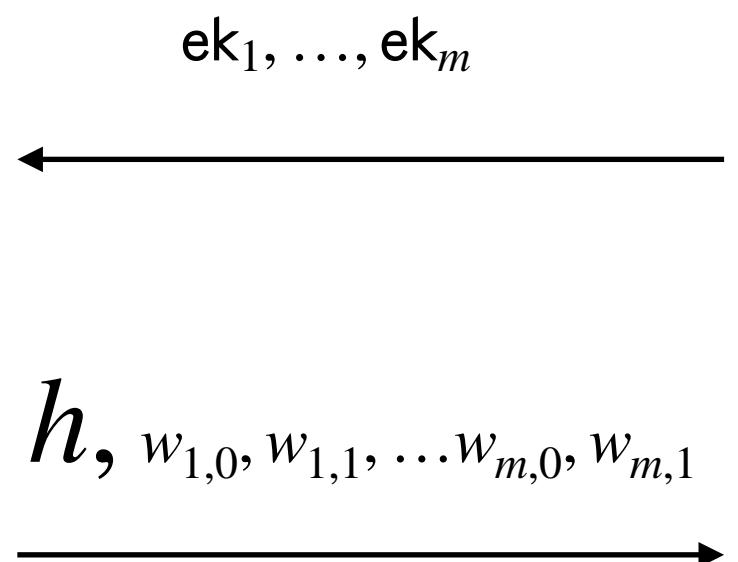
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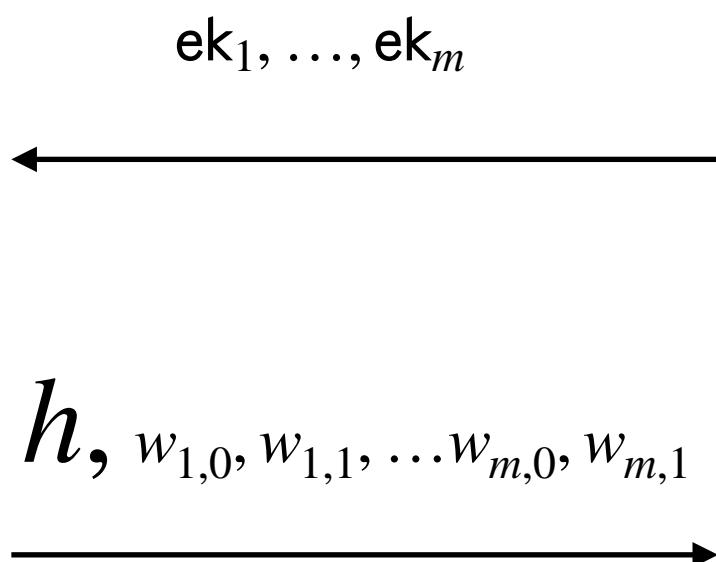
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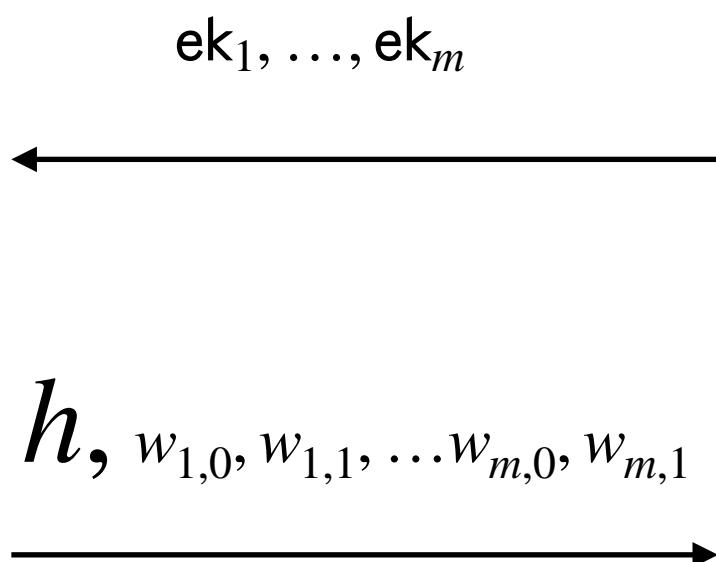
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- Each ek_i grows with $|t_i|$. That is $|ek_i| = \Omega(m)$.

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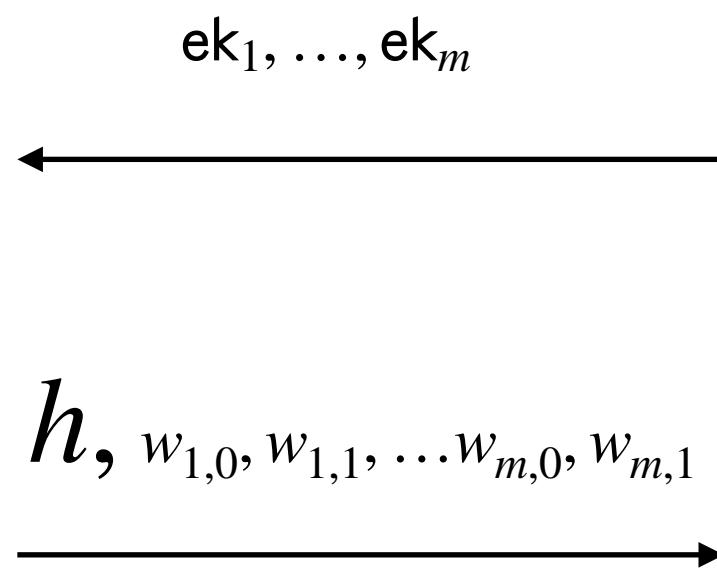
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- Receiver's rate will be a problem!
- **Solution:** Make L independent of m .

Learning Parity with Noise

$$\left(\boxed{T}, \boxed{r} \boxed{T} + \boxed{e} \right)$$

 \approx_c

$$\left(\boxed{T}, \boxed{u} \right)$$

$T \leftarrow \{0,1\}^{n \times m}$, $r \leftarrow \{0,1\}^n$, $u \leftarrow \{0,1\}^m$ and $e \leftarrow \text{Ber}(p)^m$

Learning Parity with Noise

$$\left(\boxed{T}, \boxed{r} \boxed{T} + \boxed{e} \right)$$

Expanding

$$\approx_c \left(\boxed{T}, \boxed{u} \right)$$

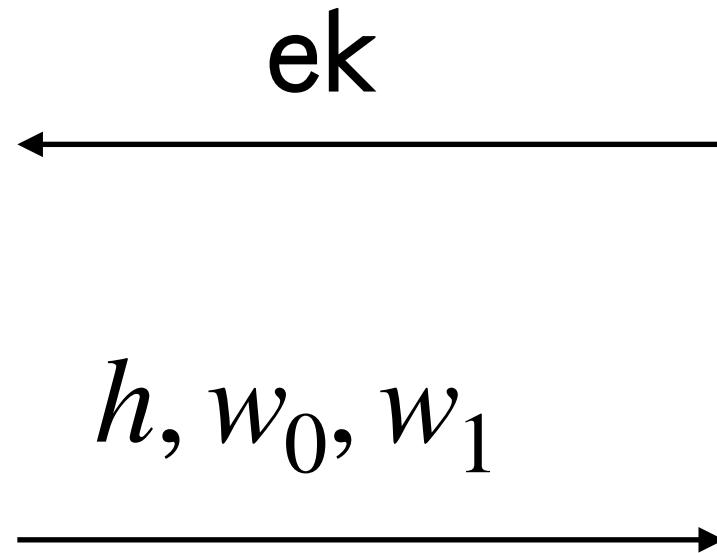
$$T \leftarrow \{0,1\}^{n \times m}, r \leftarrow \{0,1\}^n, u \leftarrow \{0,1\}^m \text{ and } e \leftarrow \text{Ber}(p)^m$$

OT from TDH for linear functions

CRS: $hk, t \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $r \leftarrow \{0,1\}^L$
 - $h \leftarrow H(hk, r)$
 - $w_0 = \text{Enc}(ek, r) \oplus e_0 \oplus \mu_0$
 - $w_1 = \text{Enc}(ek, r) \oplus \langle t, r \rangle \oplus e_1 \oplus \mu_1$
- $e_0, e_1 \leftarrow \text{Ber}(p)$



Receiver (b):

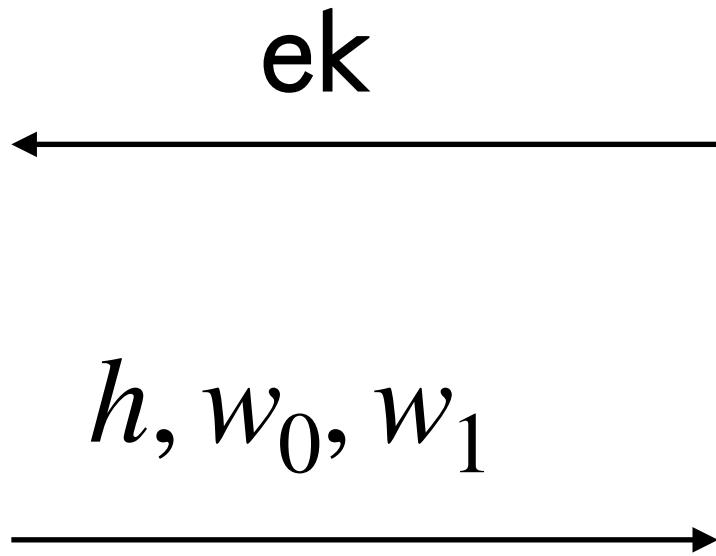
- ek encodes $b \cdot t$

OT from TDH for linear functions

CRS: $hk, t \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $r \leftarrow \{0,1\}^L$
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- $e_0, e_1 \leftarrow \text{Ber}(p)$



Receiver (b):

- ek encodes $b \cdot t$

Sender security: via

Since $(t_i, \langle t_i, r \rangle \oplus e_{i,1-b_i}) \approx_c (t_i, u_i)$

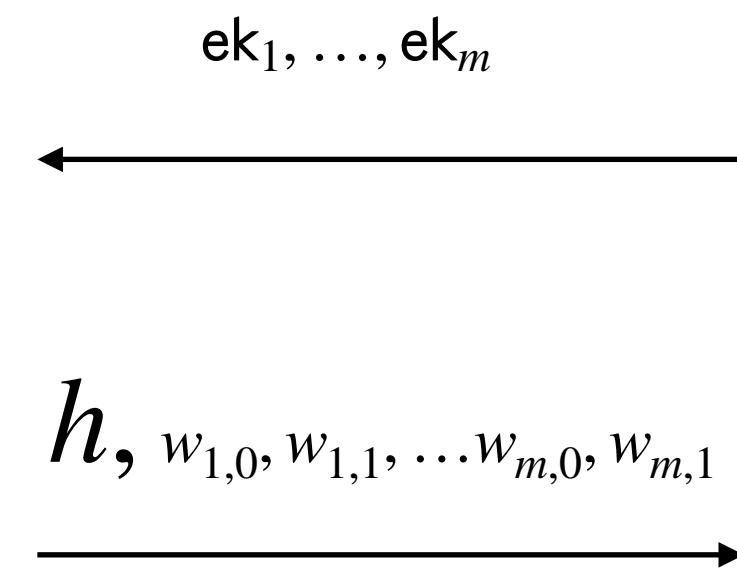
LPN

OT from TDH for linear functions

CRS: $\text{hk}, \mathbf{t}_1, \dots, \mathbf{t}_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $\mathbf{r} \leftarrow \{0,1\}^L$
- $\mathbf{h} \leftarrow \mathsf{H}(\text{hk}, \mathbf{r})$
- $w_{1,0} = \text{Enc}(\text{ek}_1, \mathbf{r}) \oplus \textcolor{red}{e_{1,0}} \oplus \mu_{1,0}$
- $w_{1,1} = \text{Enc}(\text{ek}_1, \mathbf{r}) \oplus \langle \mathbf{t}_1, \mathbf{r} \rangle \oplus \textcolor{red}{e_{1,1}} \oplus \mu_{1,1}$
- \vdots
- $w_{m,0} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus \textcolor{red}{e_{m,0}} \oplus \mu_{m,0}$
- $w_{m,1} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus \langle \mathbf{t}_m, \mathbf{r} \rangle \oplus \textcolor{red}{e_{m,1}} \oplus \mu_{m,1}$



Receiver (b_1, \dots, b_m):

- ek_1 encodes $b_1 \cdot \mathbf{t}_1$
- \vdots
- ek_m encodes $b_m \cdot \mathbf{t}_m$

Sender security: via

LPN

Since $(\mathbf{t}_i, \langle \mathbf{t}_i, \mathbf{r} \rangle \oplus e_{i,1-b_i}) \approx_c (\mathbf{t}_i, \mathbf{u}_i)$

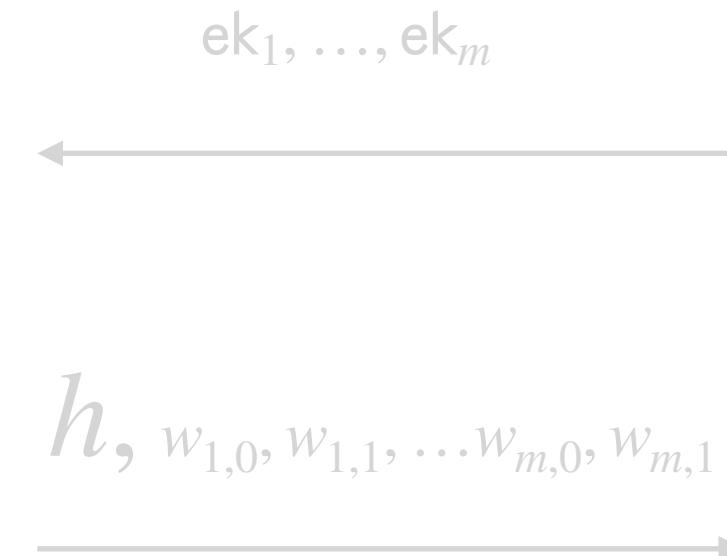
OT from TDH for linear functions

CRS: $\text{hk}, \mathbf{t}_1, \dots, \mathbf{t}_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1) :

- $\mathbf{r} \leftarrow \{0,1\}^L$
- $\mathbf{h} \leftarrow \mathsf{H}(\text{hk}, \mathbf{r})$
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- $w_{1,1} = \text{Enc}(\text{ek}_1, \mathbf{r}) \oplus \langle \mathbf{t}_1, \mathbf{r} \rangle \oplus e_{1,1} \oplus \mu_{1,1}$
- \vdots
- $w_{m,0} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus e_{m,0} \oplus \mu_{m,0}$
- $w_{m,1} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus \langle \mathbf{t}_m, \mathbf{r} \rangle \oplus e_{m,1} \oplus \mu_{m,1}$

Leak on r



Receiver (b_1, \dots, b_m) :

- ek_1 encodes $b_1 \cdot \mathbf{t}_1$
- \vdots
- ek_m encodes $b_m \cdot \mathbf{t}_m$

Sender security: via

LPN

Since $\left(\mathbf{t}_i, \langle \mathbf{t}_i, \mathbf{r} \rangle \oplus e_{i,1-b_i} \right) \approx_c \left(\mathbf{t}_i, \mathbf{u}_i \right)$

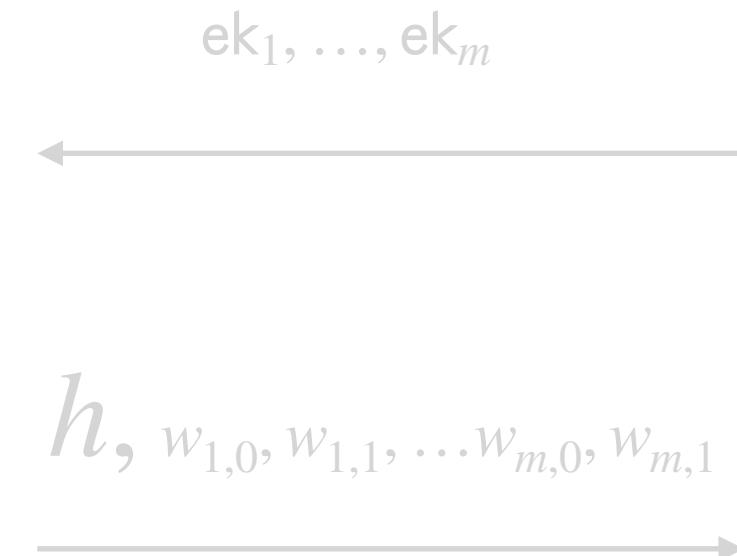
OT from TDH for linear functions

CRS: $\text{hk}, \mathbf{t}_1, \dots, \mathbf{t}_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1) :

- $\mathbf{r} \leftarrow \{0,1\}^L$
- $\mathbf{h} \leftarrow \mathsf{H}(\text{hk}, \mathbf{r})$
- $w_{1,0} = \text{Enc}(\text{ek}_1, \mathbf{r}) \oplus e_{1,0} \oplus \mu_{1,0}$
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- \vdots
- $w_{m,0} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus e_{m,0} \oplus \mu_{m,0}$
- $w_{m,1} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus \langle \mathbf{t}_m, \mathbf{r} \rangle \oplus e_{m,1} \oplus \mu_{m,1}$

Leak on r



Receiver (b_1, \dots, b_m) :

- ek_1 encodes $b_1 \cdot \mathbf{t}_1$
- \vdots
- ek_m encodes $b_m \cdot \mathbf{t}_m$

Sender security: via (entropic?) LPN

Since $\left(\mathbf{t}_i, \langle \mathbf{t}_i, \mathbf{r} \rangle \oplus e_{i,1-b_i} \right) \approx_c \left(\mathbf{t}_i, \mathbf{u}_i \right)$

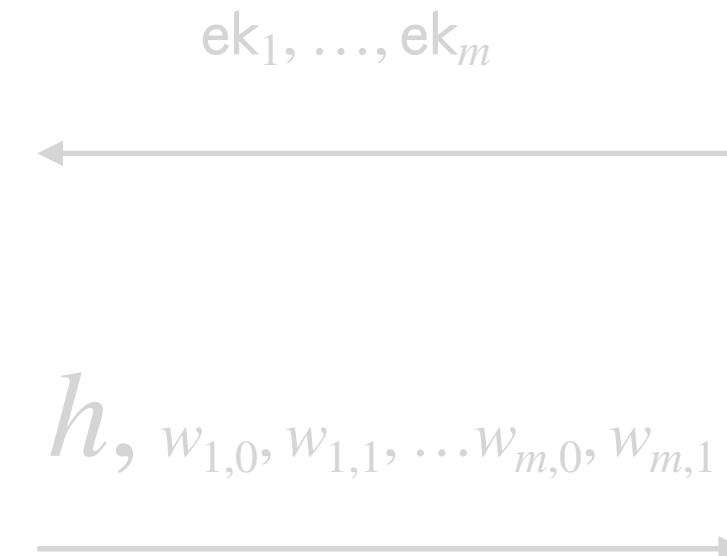
OT from TDH for linear functions

CRS: $hk, t_1, \dots, t_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1) :

- $r \leftarrow \{0,1\}^L$
- $h \leftarrow H(hk, r)$
- $w_{1,0} = \text{Enc}(ek_1, r) \oplus e_{1,0} \oplus \mu_{1,0}$
- $w_{1,1} = \text{Enc}(ek_1, r) \oplus \langle t_1, r \rangle \oplus e_{1,1} \oplus \mu_{1,1}$
- \vdots
- $w_{m,0} = \text{Enc}(ek_m, r) \oplus e_{m,0} \oplus \mu_{m,0}$
- $w_{m,1} = \text{Enc}(ek_m, r) \oplus \langle t_m, r \rangle \oplus e_{m,1} \oplus \mu_{m,1}$

Leak on r



Receiver (b_1, \dots, b_m) :

- ek_1 encodes $b_1 \cdot t_1$
- \vdots
- ek_m encodes $b_m \cdot t_m$

Sender security: via

LPN

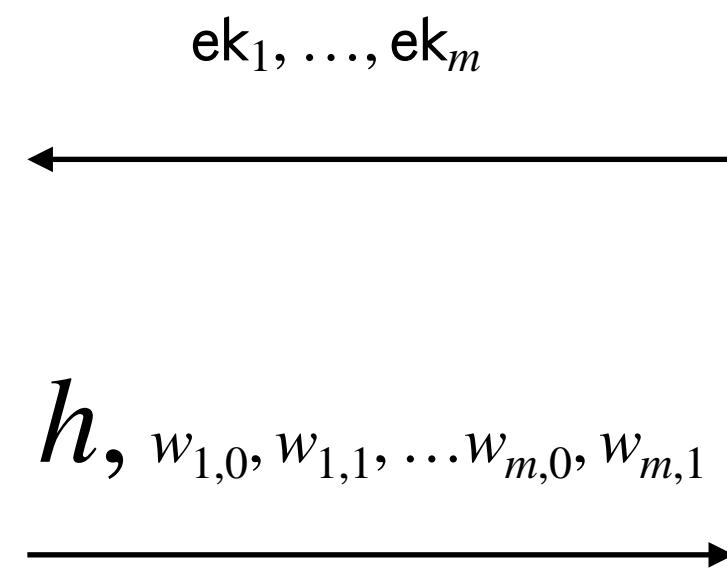
Since $\left(t_i, \langle t_i, r \rangle \oplus e_{i,1-b_i} \right) \approx_c (t_i, u_i)$

OT from TDH for linear functions

CRS: $hk, t_1, \dots, t_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $r \leftarrow \{0,1\}^L$
- $h \leftarrow H(hk, r)$
- $w_{1,0} = \text{Enc}(\text{ek}_1, r) \oplus e_{1,0} \oplus \mu_{1,0}$
- $w_{1,1} = \text{Enc}(\text{ek}_1, r) \oplus \langle t_1, r \rangle \oplus e_{1,1} \oplus \mu_{1,1}$
- \vdots
- $w_{m,0} = \text{Enc}(\text{ek}_m, r) \oplus e_{m,0} \oplus \mu_{m,0}$
- $w_{m,1} = \text{Enc}(\text{ek}_m, r) \oplus \langle t_m, r \rangle \oplus e_{m,1} \oplus \mu_{m,1}$



Receiver (b_1, \dots, b_m):

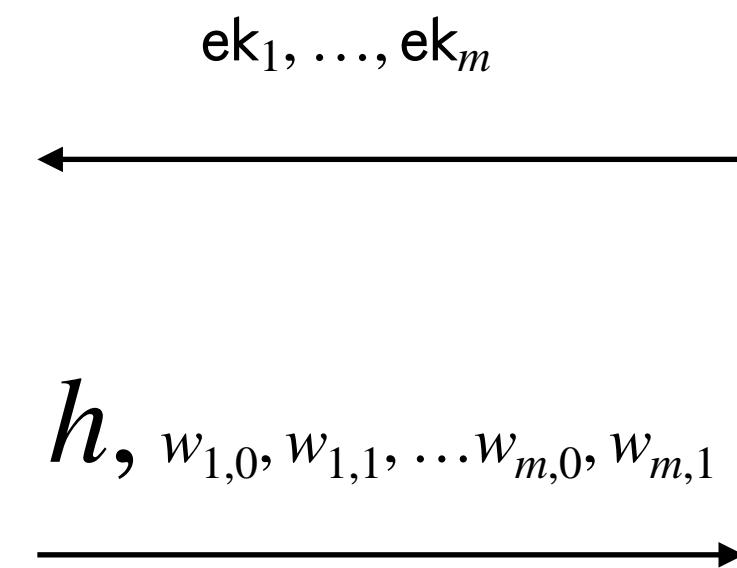
- ek_1 encodes $b_1 \cdot t_1$
- \vdots
- ek_m encodes $b_m \cdot t_m$

OT from TDH for linear functions

CRS: $hk, t_1, \dots, t_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $r \leftarrow \{0,1\}^L$
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Receiver (b_1, \dots, b_m):

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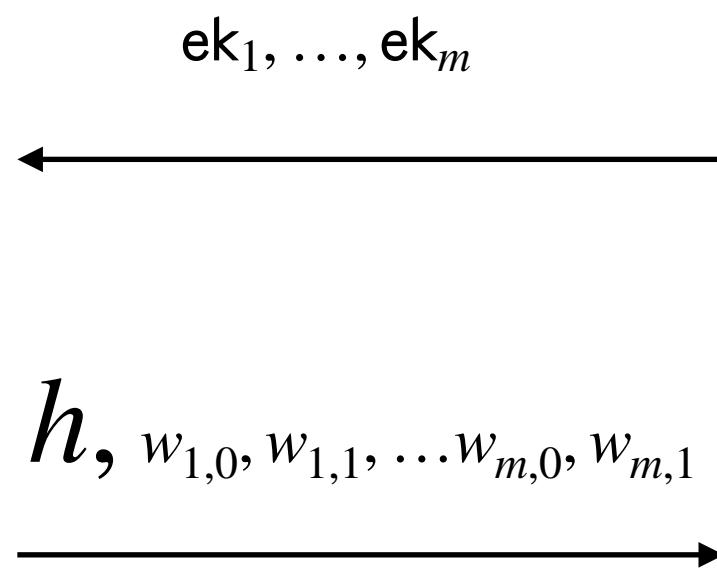
Faulty correctness

OT from TDH for linear functions

CRS: $\text{hk}, \mathbf{t}_1, \dots, \mathbf{t}_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $\mathbf{r} \leftarrow \{0,1\}^L$
- $\mathbf{h} \leftarrow \mathsf{H}(\text{hk}, \mathbf{r})$
- $w_{1,0} = \text{Enc}(\text{ek}_1, \mathbf{r}) \oplus \textcolor{red}{e}_{1,0} \oplus \mu_{1,0}$
- $w_{1,1} = \text{Enc}(\text{ek}_1, \mathbf{r}) \oplus \langle \mathbf{t}_1, \mathbf{r} \rangle \oplus \textcolor{red}{e}_{1,1} \oplus \mu_{1,1}$
- \vdots
- $w_{m,0} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus \textcolor{red}{e}_{m,0} \oplus \mu_{m,0}$
- $w_{m,1} = \text{Enc}(\text{ek}_m, \mathbf{r}) \oplus \langle \mathbf{t}_m, \mathbf{r} \rangle \oplus \textcolor{red}{e}_{m,1} \oplus \mu_{m,1}$



Receiver (b_1, \dots, b_m):

- ek_1 encodes $b_1 \cdot \mathbf{t}_1$
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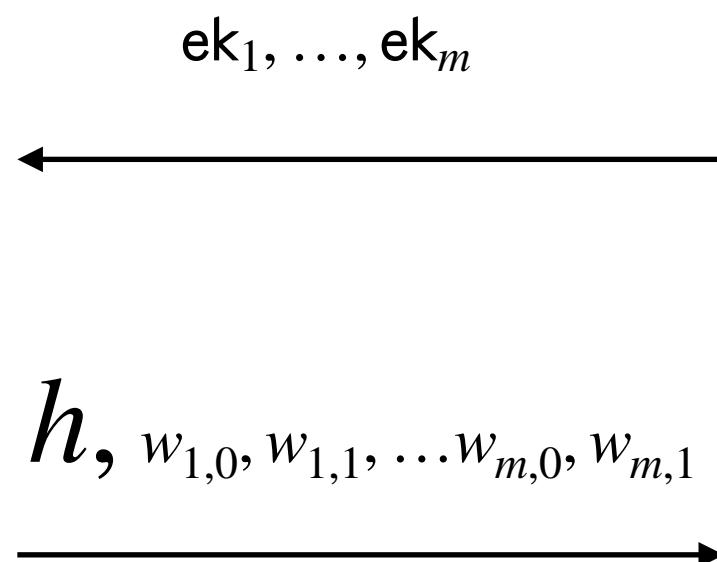


OT from TDH for linear functions

CRS: $hk, t_1, \dots, t_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $r \leftarrow \{0,1\}^L$
- $h \leftarrow H(hk, r)$
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- $w_{1,1} = \text{Enc}(ek_1, r) \oplus \langle t_1, r \rangle \oplus e_{1,1} \oplus \mu_{1,1}$
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- $w_{m,1} = \text{Enc}(ek_m, r) \oplus \langle t_m, r \rangle \oplus e_{m,1} \oplus \mu_{m,1}$



Receiver (b_1, \dots, b_m):

- ek_1 encodes $b_1 \cdot t_1$
- \vdots
- ek_m encodes $b_m \cdot t_m$



Communication overhead:

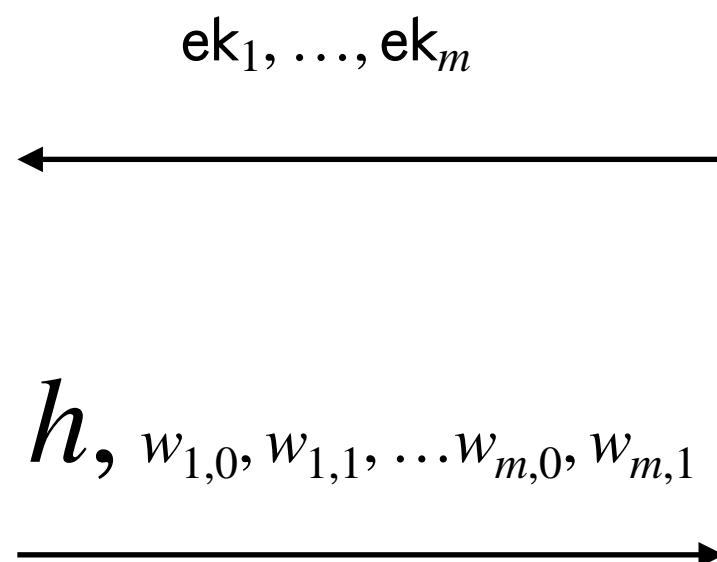
~ (Hamming weight of error)

OT from TDH for linear functions

CRS: $hk, t_1, \dots, t_m \leftarrow \{0,1\}^L$

Sender (μ_0, μ_1):

- $r \leftarrow \{0,1\}^L$
- $h \leftarrow H(hk, r)$
- $w_{1,0} = \text{Enc}(ek_1, r) \oplus e_{1,0} \oplus \mu_{1,0}$
- $w_{1,1} = \text{Enc}(ek_1, r) \oplus \langle t_1, r \rangle \oplus e_{1,1} \oplus \mu_{1,1}$
- \vdots
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- $w_{m,1} = \text{Enc}(ek_m, r) \oplus \langle t_m, r \rangle \oplus e_{m,1} \oplus \mu_{m,1}$



Receiver (b_1, \dots, b_m):

- ek₁ encodes $b_1 \cdot t_1$
- \vdots
- ek_m encodes $b_m \cdot t_m$



Communication overhead:

~ (Hamming weight of error)



Set as m^ϵ

Malicious security against senders

Malicious security:

- Almost for free
- 
- Perfect correctness of TDH

Malicious security against senders

Malicious security:

- Almost for free  Perfect correctness of TDH

- h of r are well-formed.



2PC

Communication overhead:

$\sim |r|$



Set as m^ϵ

Recap

- This talk: two-round maliciously-sender secure OT with optimal download rate from QR + LPN.

Recap

- **This talk:** two-round maliciously-sender secure OT with optimal download rate from QR + LPN.
- 
- **Main Result:** two-round maliciously-secure OT with optimal rate from QR + LPN.

Recap

- **This talk:** two-round maliciously-sender secure OT with optimal download rate from QR + LPN.

- **Main Result:** two-round maliciously-secure OT with optimal rate from QR + LPN.
- **Main insight:** how to use TDH with LPN for improved communication and security.

Recap

- **This talk:** two-round maliciously-sender secure OT with optimal download rate from QR + LPN.
- 
- **Main Result:** two-round maliciously-secure OT with optimal rate from QR + LPN.
 - **Main insight:** how to use TDH with LPN for improved communication and security.

Thanks!

Two-Round Maliciously-Secure Oblivious Transfer with Optimal Rate

Pedro Branco *Max-Planck Institute for Security and Privacy*

Nico Döttling *Helmholtz Center for Information Security (CISPA)*

Akshayaram Srinivasan *University of Toronto*

