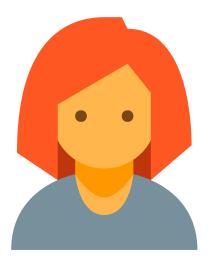


#### The Power of NAPs

Compressing OR-Proofs via Collision-Resistant Hashing

Katharina Boudgoust & <u>Mark Simkin</u>





Prover



Verifier

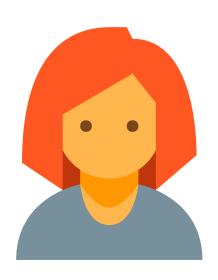




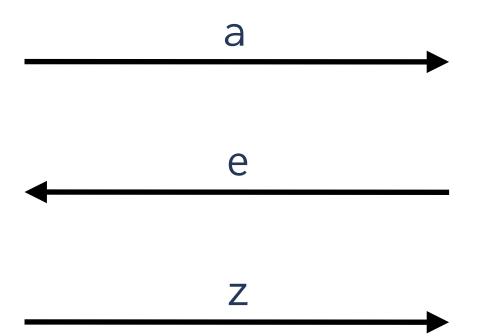
Prover







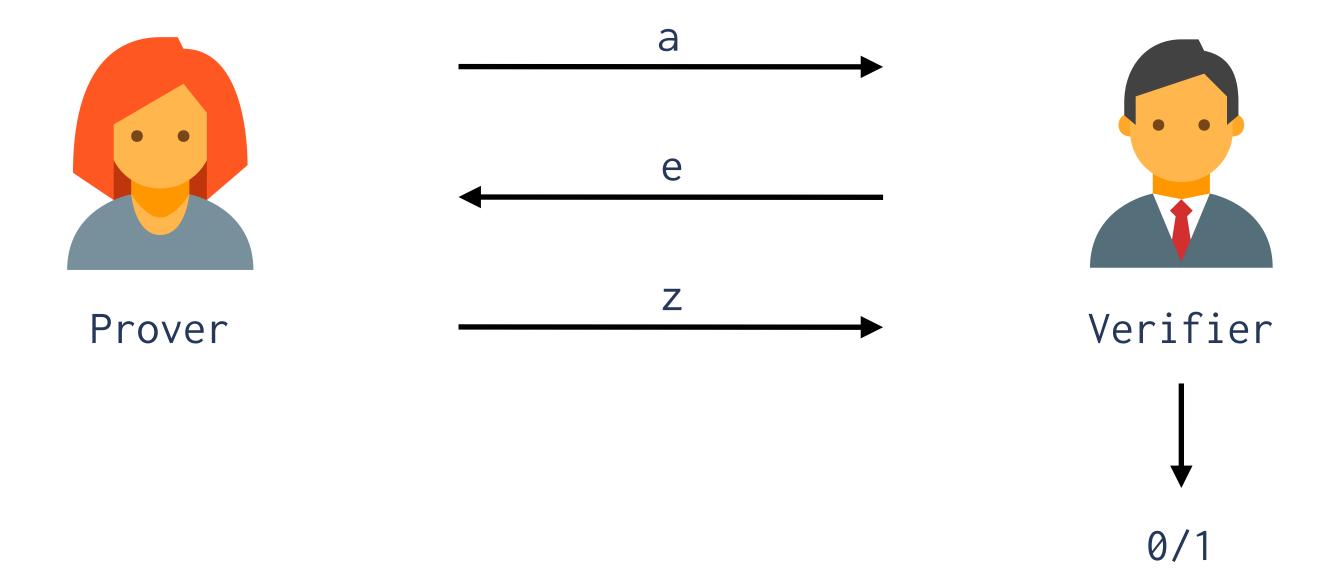




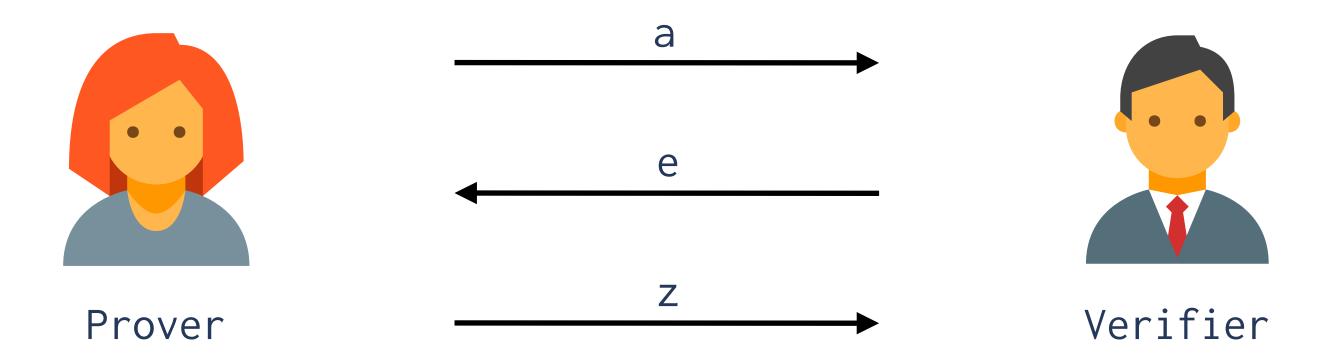


Verifier



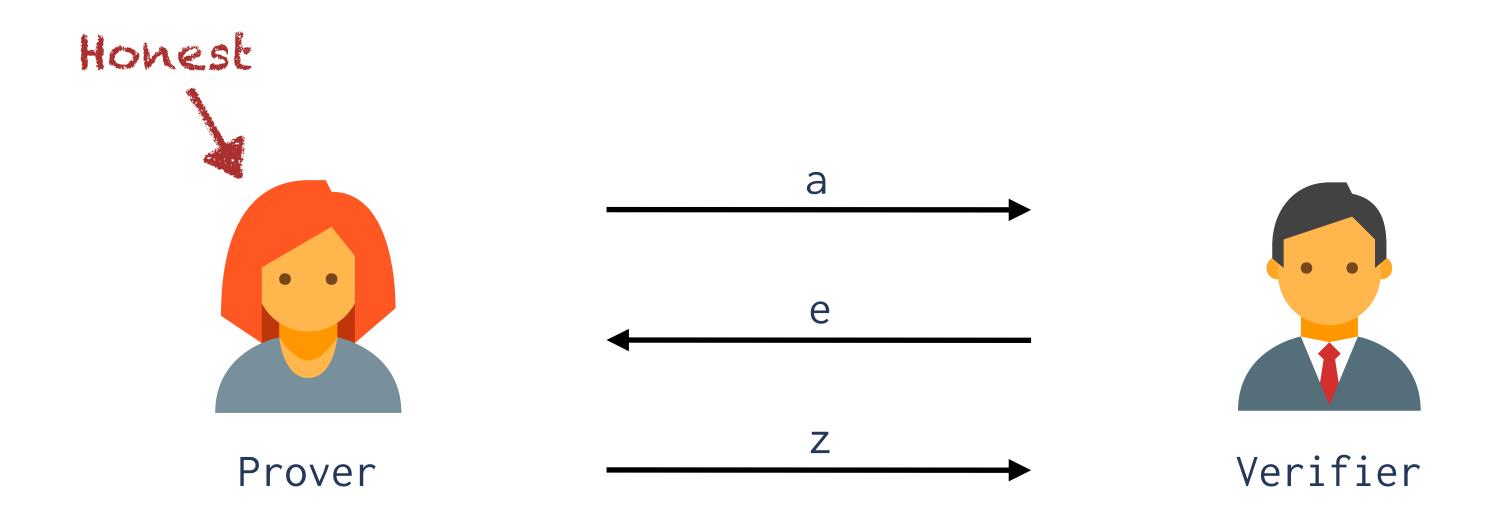




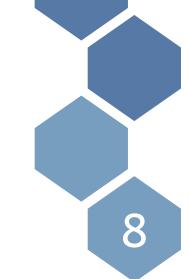


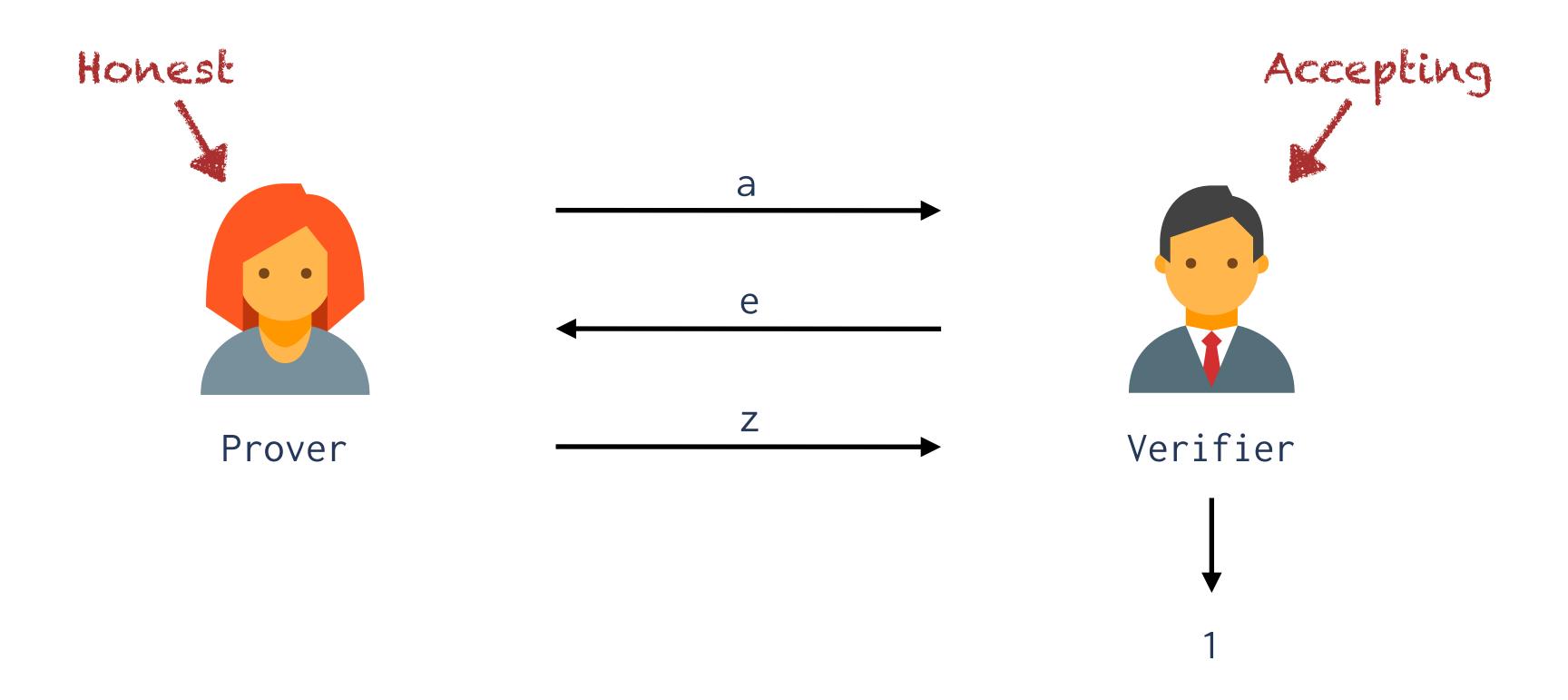
Completeness





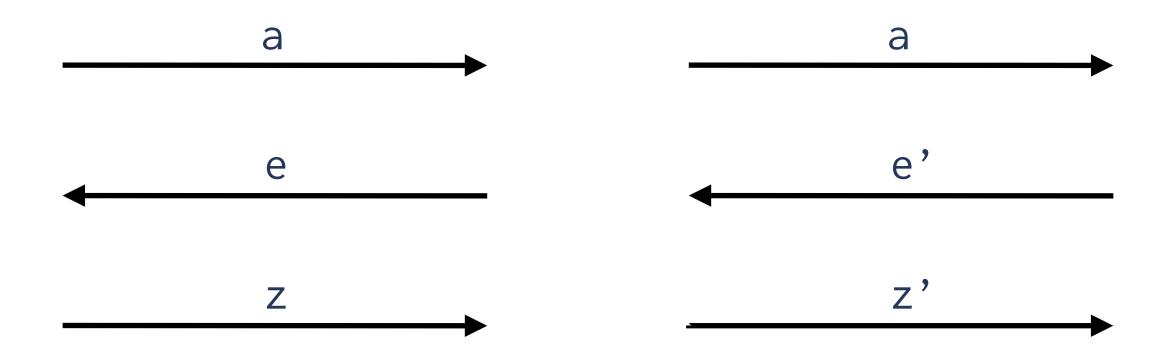
Completeness





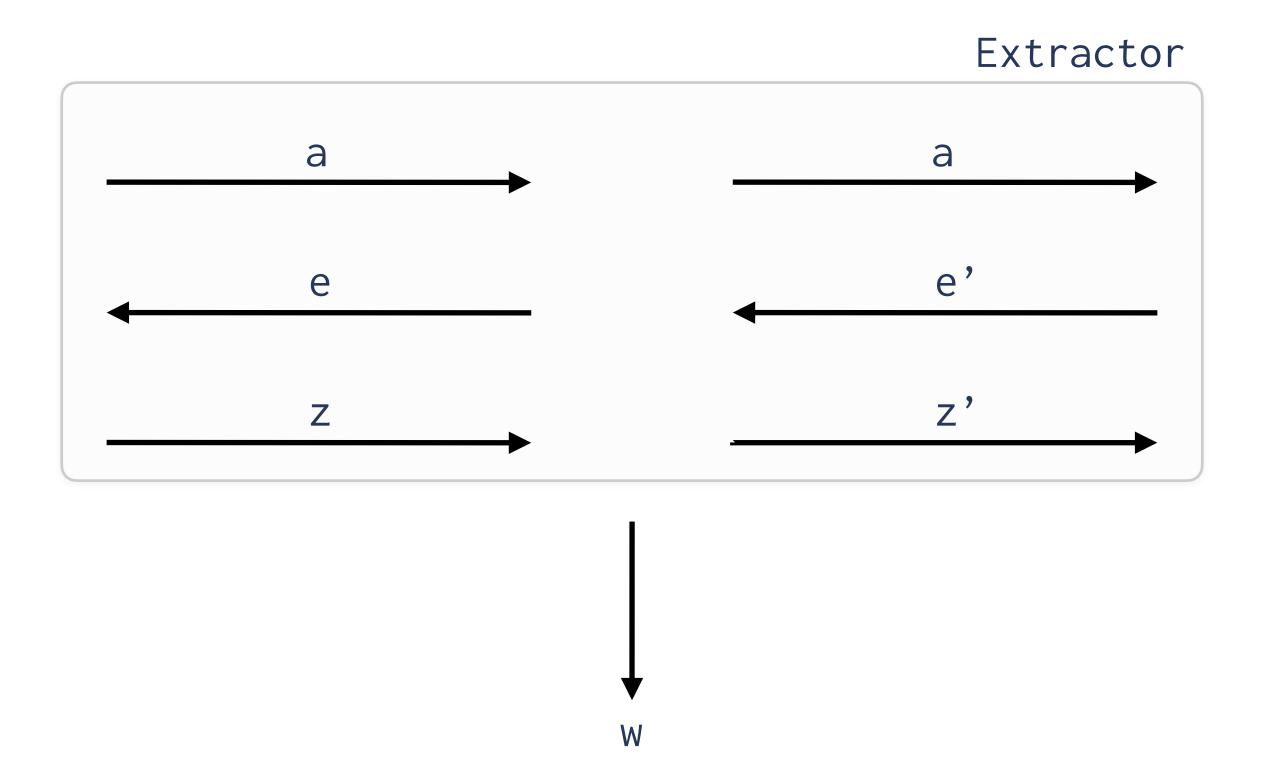
Completeness





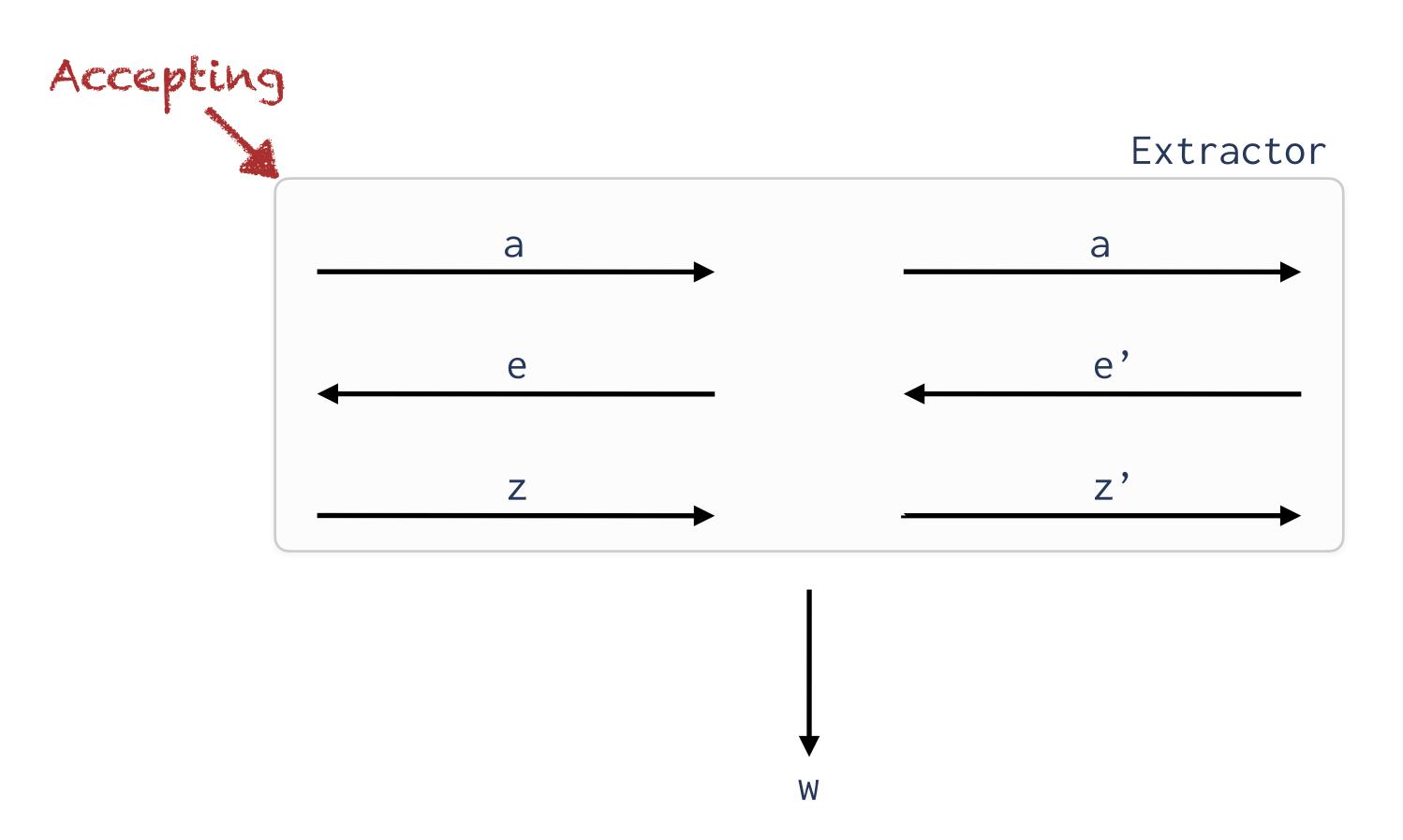
- Completeness
- Special soundness





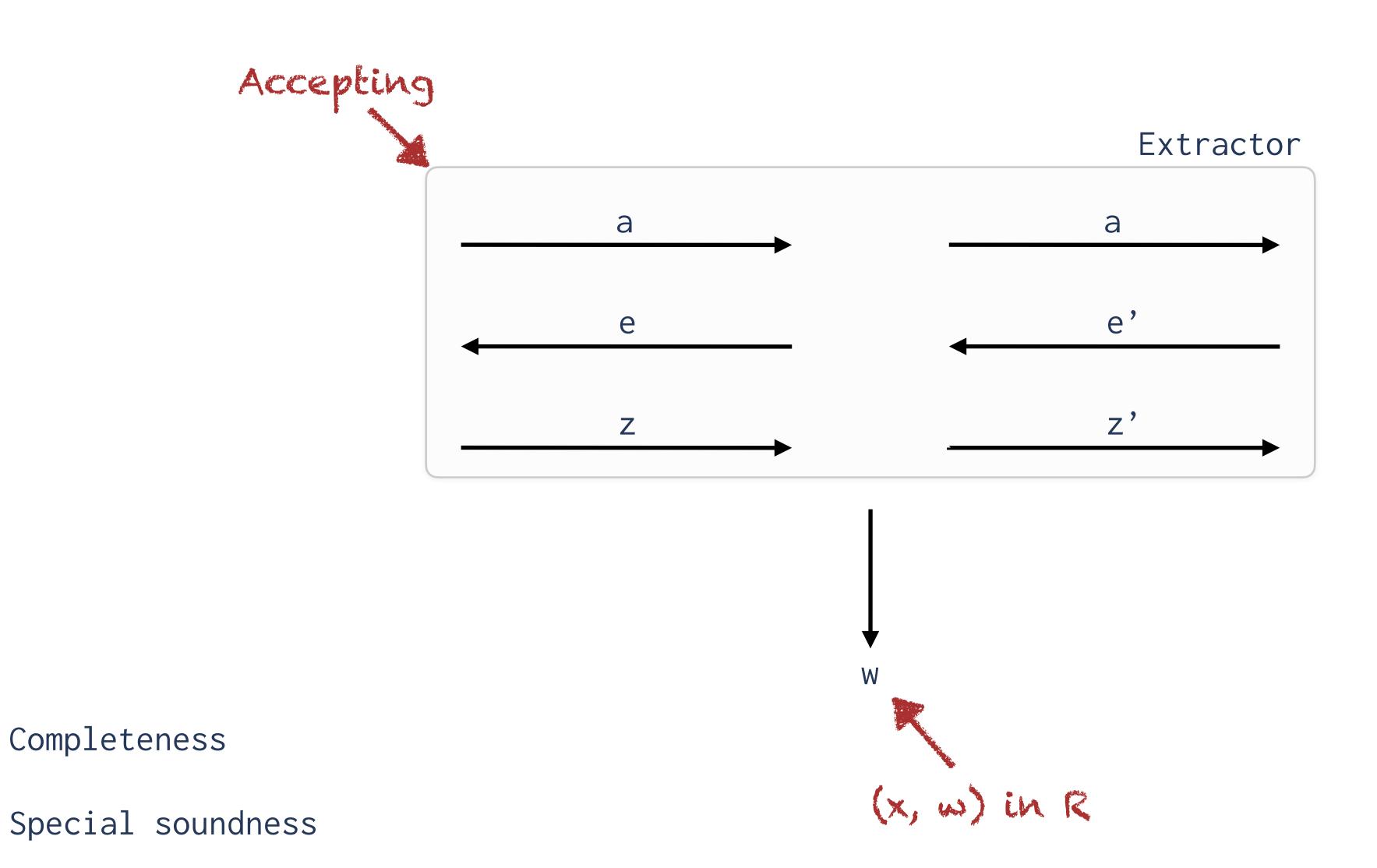
- Completeness
- Special soundness





- Completeness
- Special soundness







Simulator

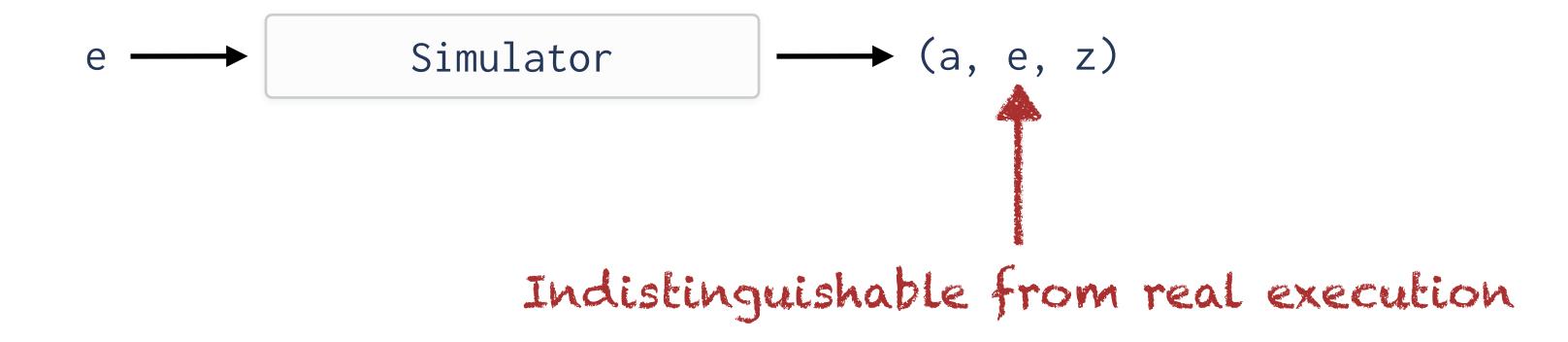
- Completeness
- Special soundness
- Honest verifier zero-knowledge





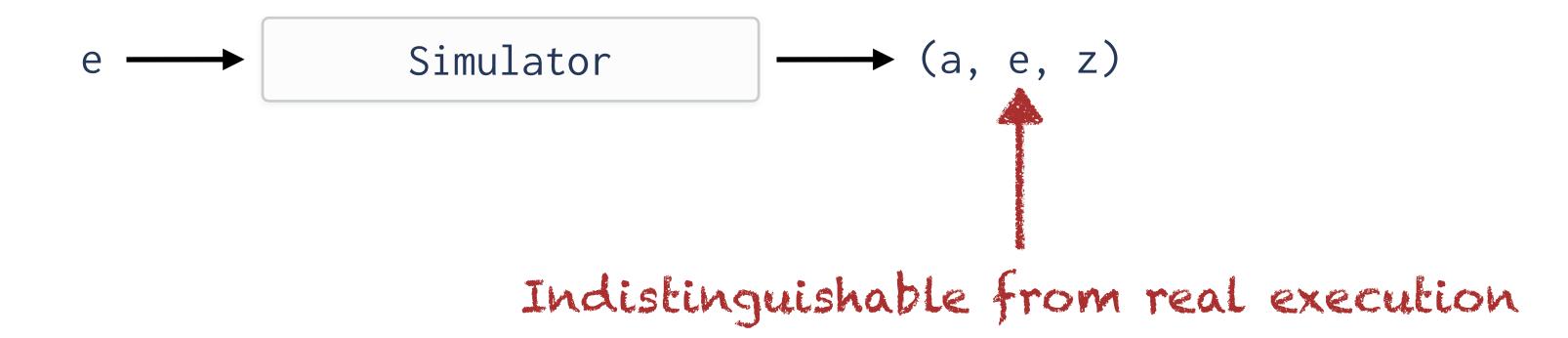
- Completeness
- Special soundness
- Honest verifier zero-knowledge





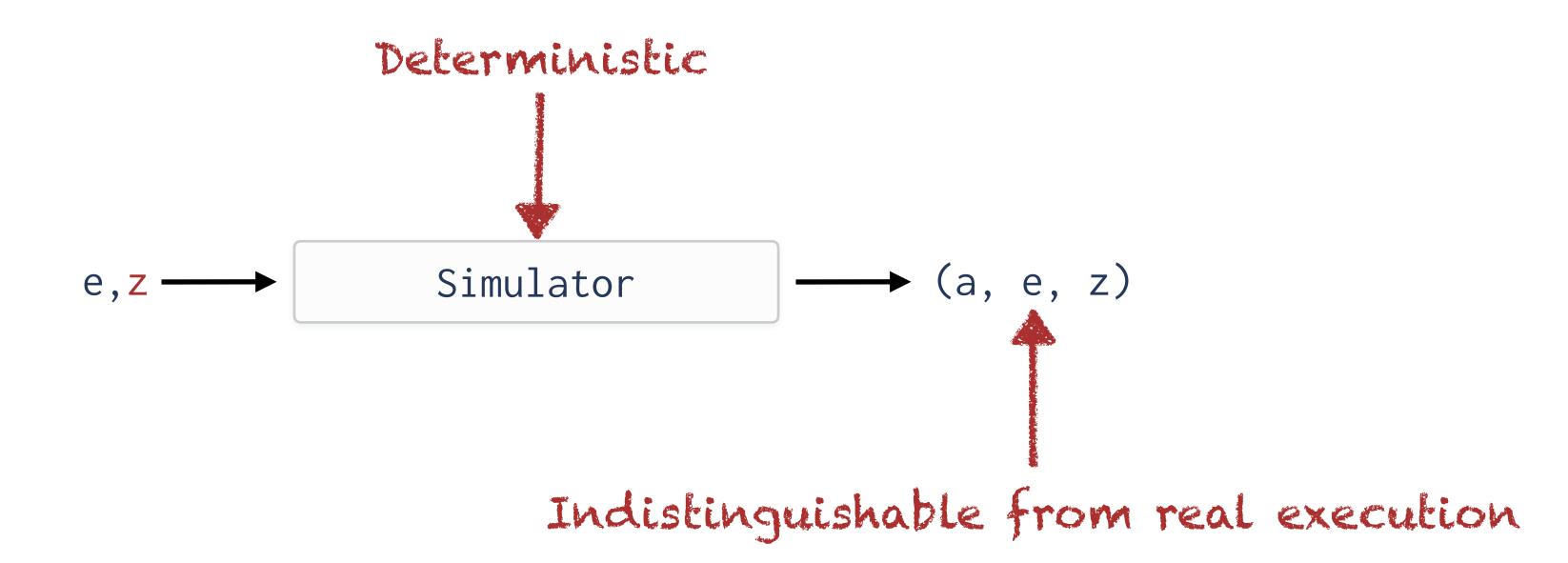
- Completeness
- Special soundness
- Honest verifier zero-knowledge





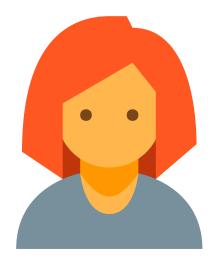
- Completeness
- Special soundness
- Strong honest verifier zero-knowledge



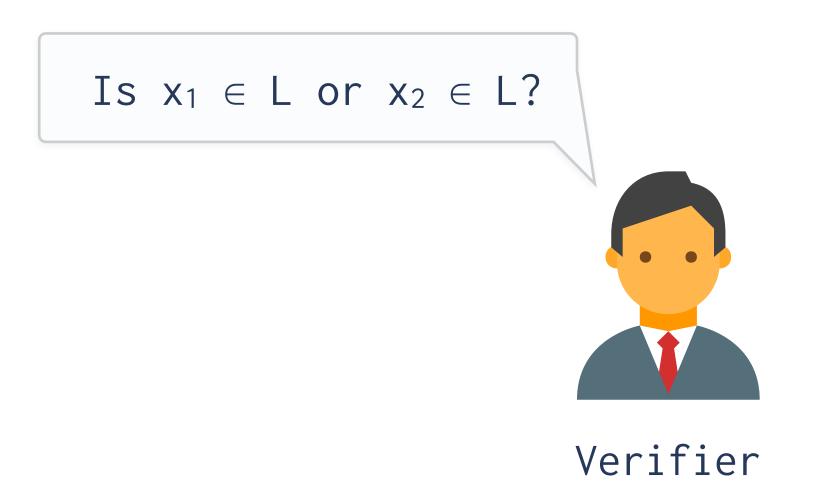


- Completeness
- Special soundness
- Strong honest verifier zero-knowledge





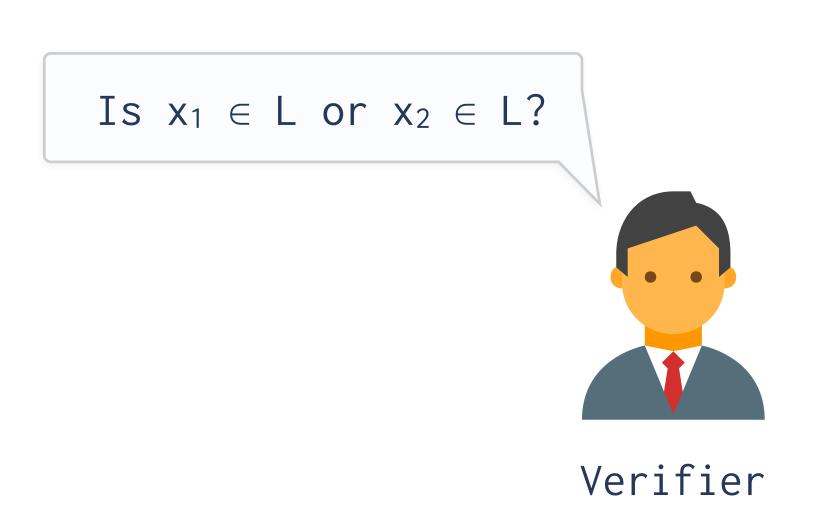
Prover









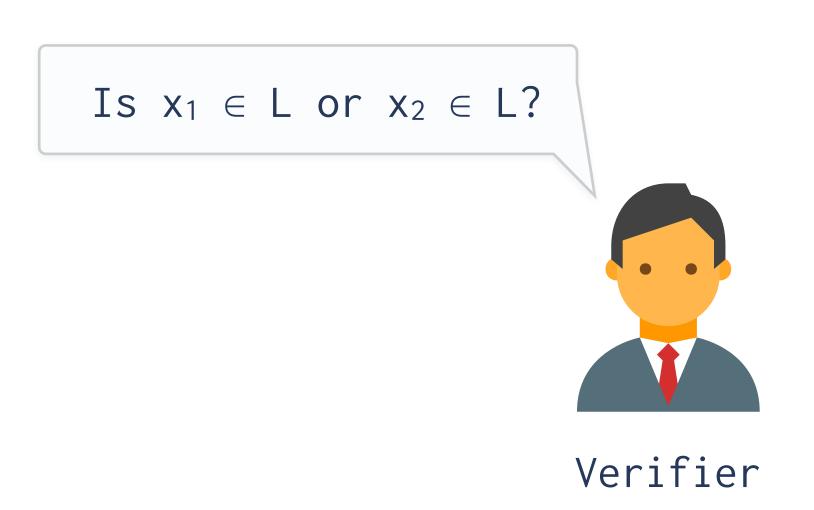


Prove that one out of several statements is true







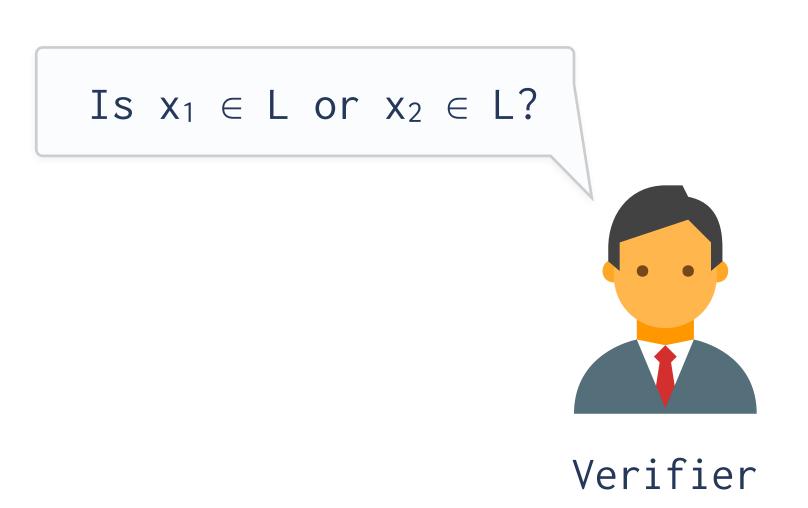


- Prove that one out of several statements is true
- Specific enough to be solved more efficiently



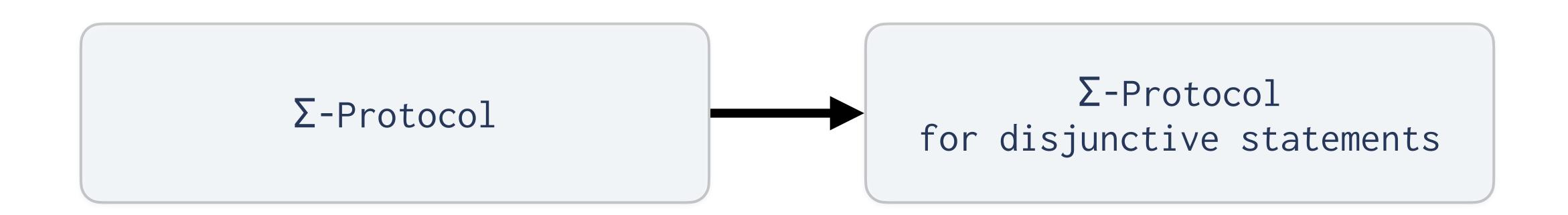




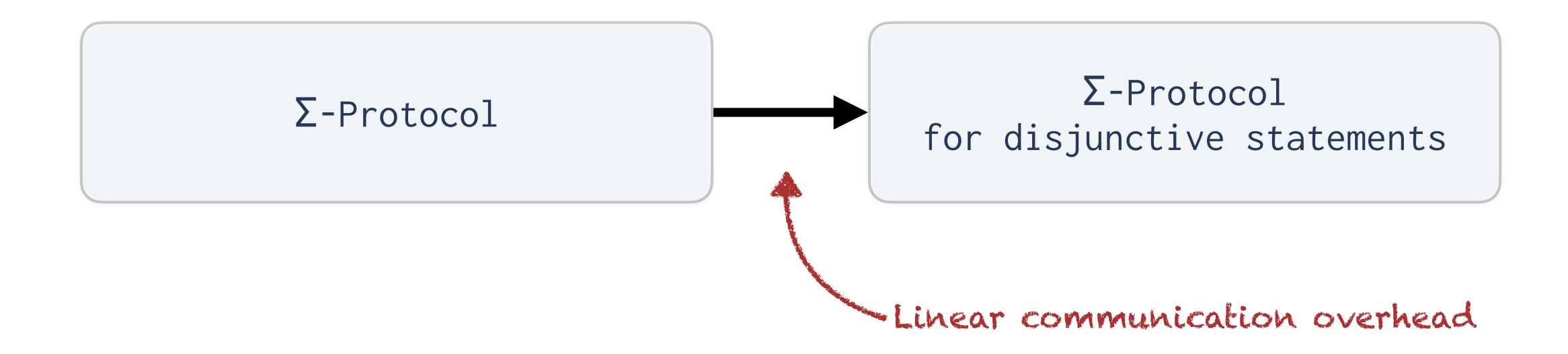


- Prove that one out of several statements is true
- Specific enough to be solved more efficiently
- General enough to be useful (ring signatures, electronic voting, ...)

19

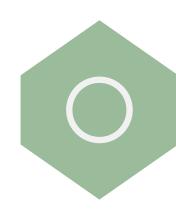








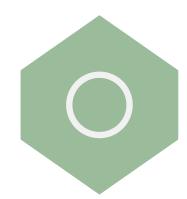




From structured hardness assumptions

Small communication complexity, but stronger assumptions





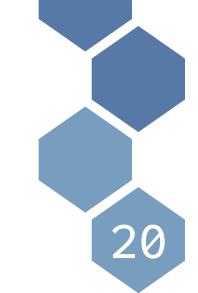
From structured hardness assumptions

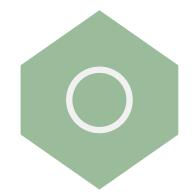
Small communication complexity, but stronger assumptions



Via MPC-in-the-head

Unstructured assumptions, but large communication complexity





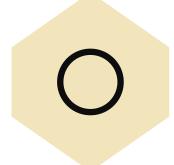
From structured hardness assumptions

Small communication complexity, but stronger assumptions



Via MPC-in-the-head

Unstructured assumptions, but large communication complexity



Via PCPs/IOPs and collision-resistant hashing Complex, heavy machinery with polylog communication

### Our Contribution



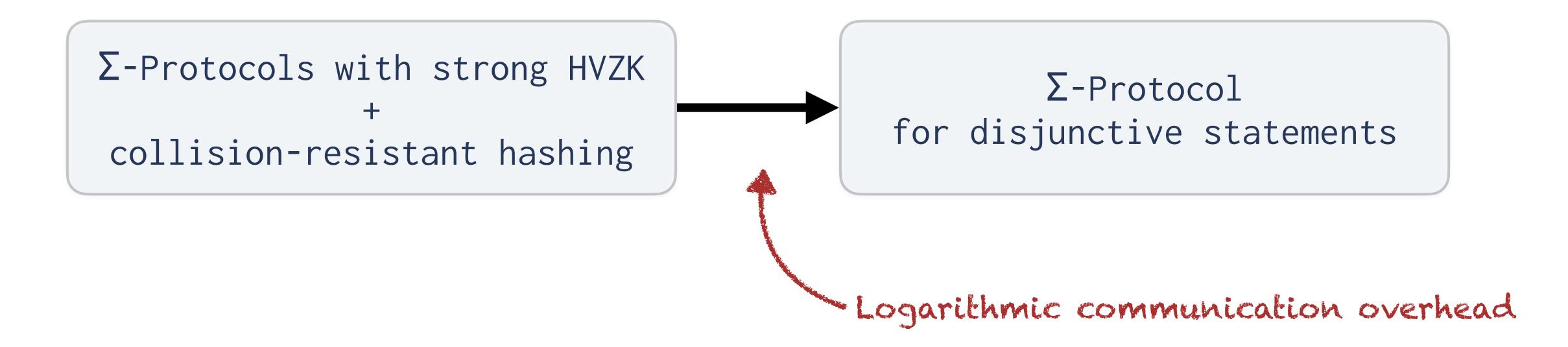
Σ-Protocols with strong HVZK
+
collision-resistant hashing

 $\begin{array}{c} \Sigma\text{-Protocol} \\ \text{for disjunctive statements} \end{array}$ 

\*Logarithmic communication overhead

#### Our Contribution

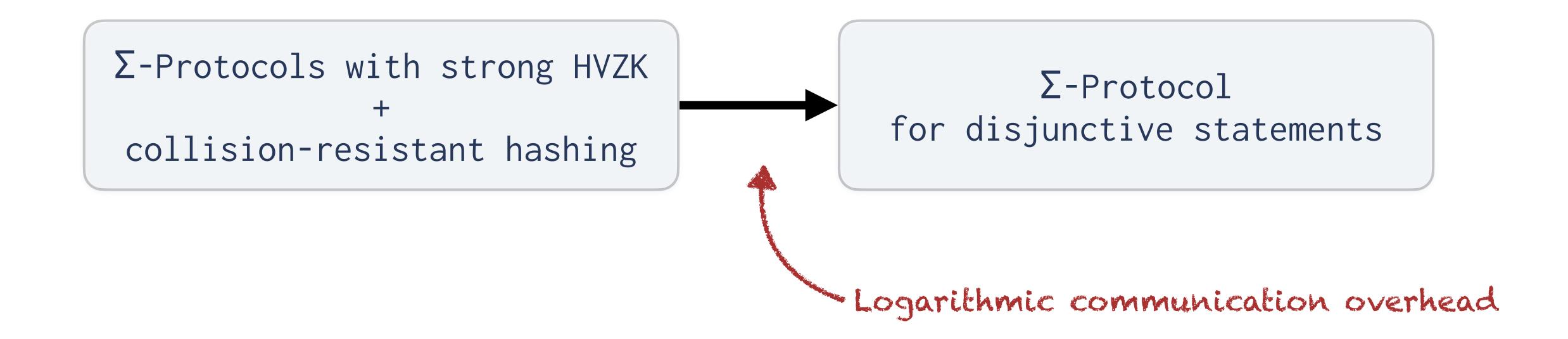




New Notion of non-adaptively programmable functions (NAPs)

#### Our Contribution





- New Notion of non-adaptively programmable functions (NAPs)
- Rejection sampling can be explained





Prover



Verifier





Prover









Prover



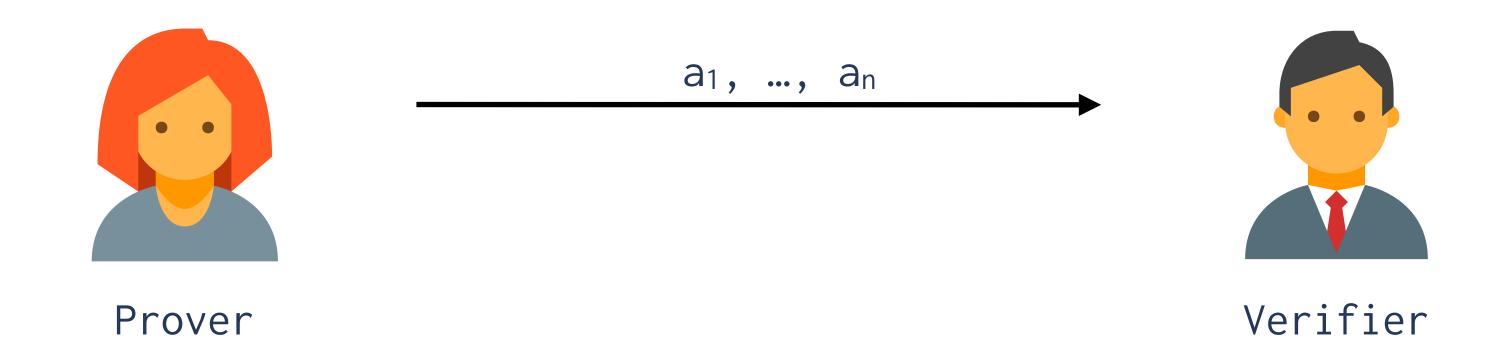
Verifier

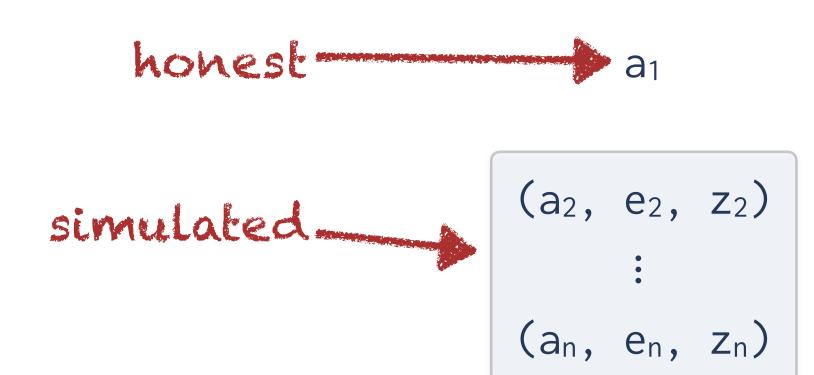
simulated 
$$(a_2, e_2, z_2)$$

$$\vdots$$

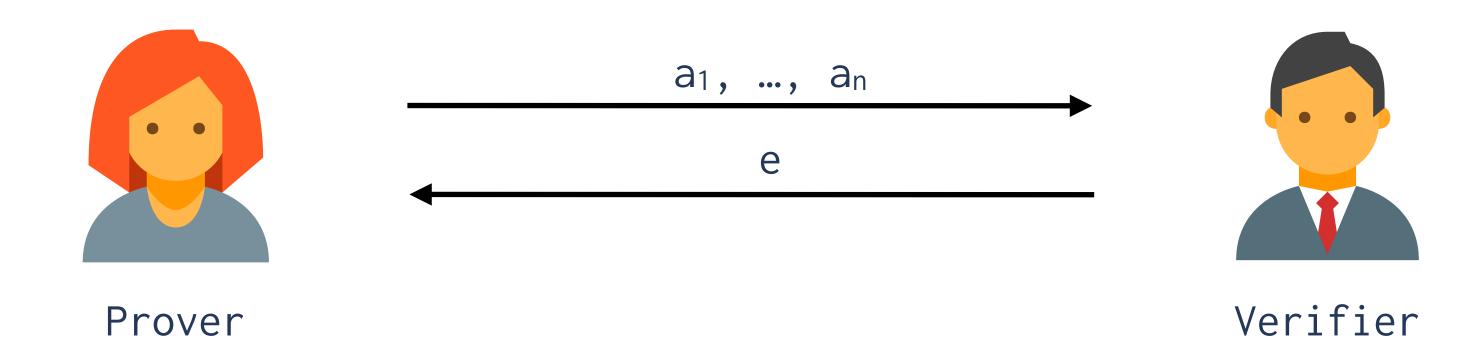
$$(a_n, e_n, z_n)$$

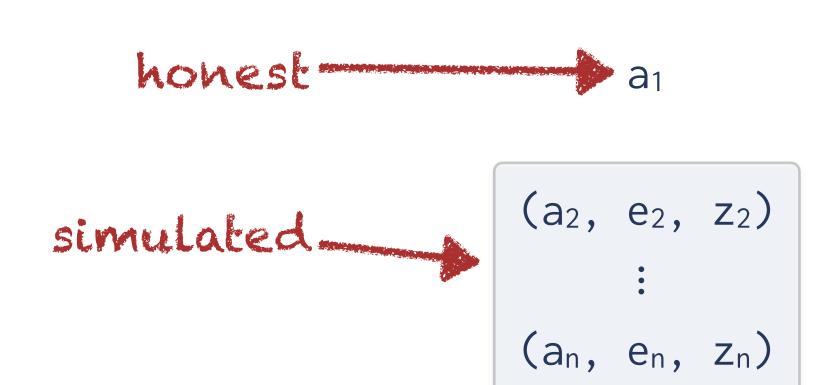


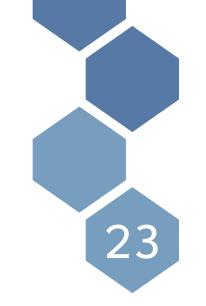


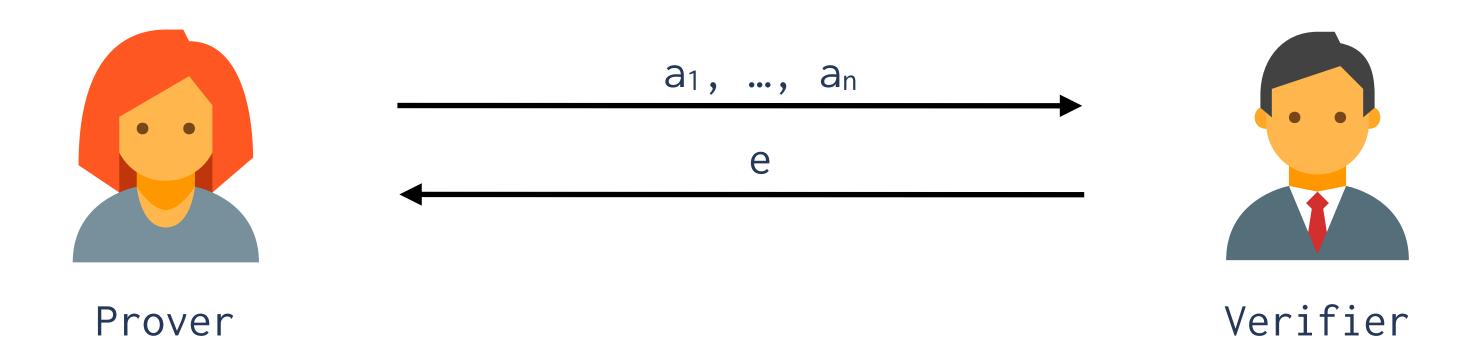


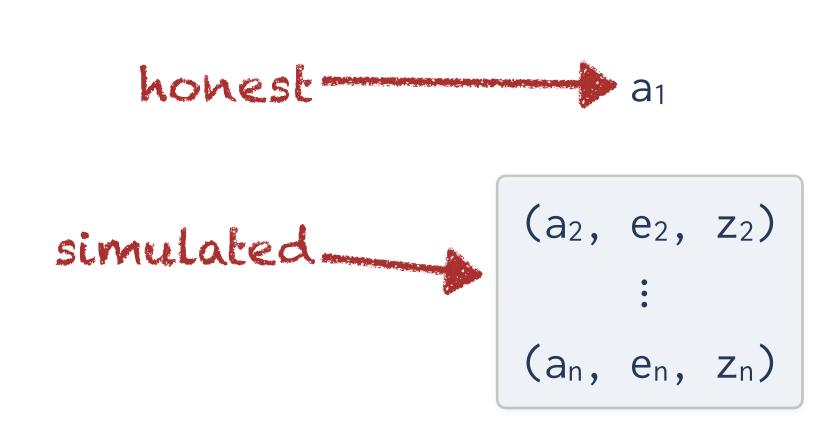






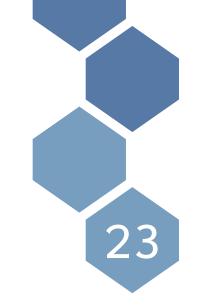


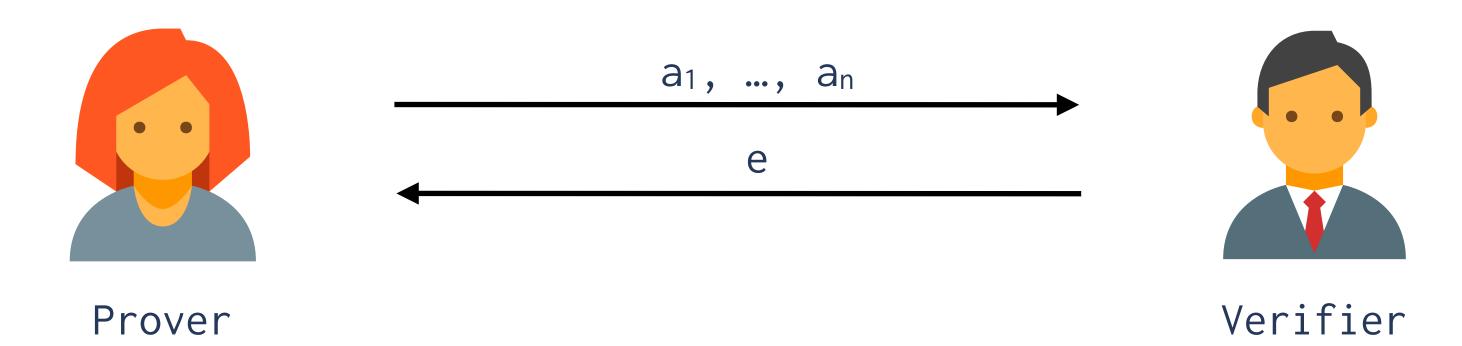




$$e = e_1 + ... + e_n$$

[Cramer, Damgård, Schoenmakers 1994]





honest

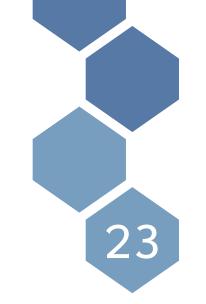
simulated 
$$(a_2, e_2, z_2)$$

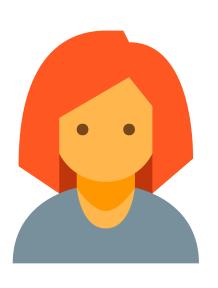
$$\vdots$$

$$(a_n, e_n, z_n)$$

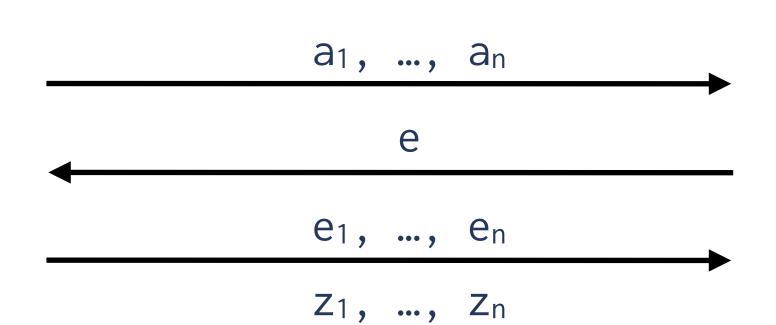
$$e = e_1 + ... + e_n$$

[Cramer, Damgård, Schoenmakers





Prover





honest.



$$(a_2, e_2, z_2)$$
 $\vdots$ 
 $(a_n, e_n, z_n)$ 

$$e = e_1 + ... + e_n$$

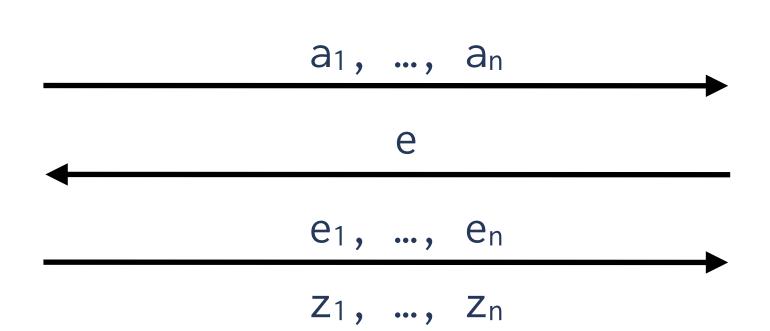


[Cramer, Damgård, Schoenmakers





Prover





Verifier

#### honest.



$$(a_2, e_2, z_2)$$
 $\vdots$ 
 $(a_n, e_n, z_n)$ 

$$e = e_1 + ... + e_n$$

$$(a_1, e_1, z_1)$$

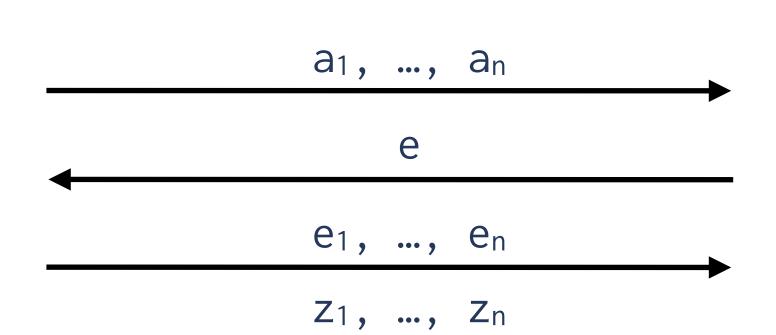
$$(a_n, e_n, z_n)$$

[Cramer, Damgård, Schoenmakers 1994]





Prover





Verifier

#### honest a1

simulated 
$$(a_2, e_2, z_2)$$

$$\vdots$$

$$(a_n, e_n, z_n)$$

$$e = e_1 + ... + e_n$$

Check each transcript

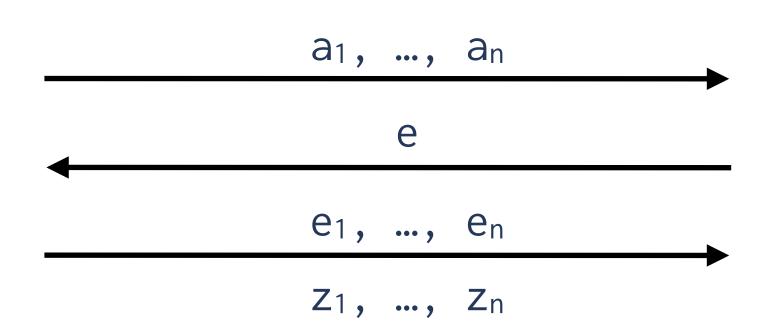
$$(a_n, e_n, z_n)$$

Our Construction











Verifier

**a**<sub>1</sub>

$$(a_2, e_2, z_2)$$
 $\vdots$ 
 $(a_n, e_n, z_n)$ 

$$e = e_1 + ... + e_n$$

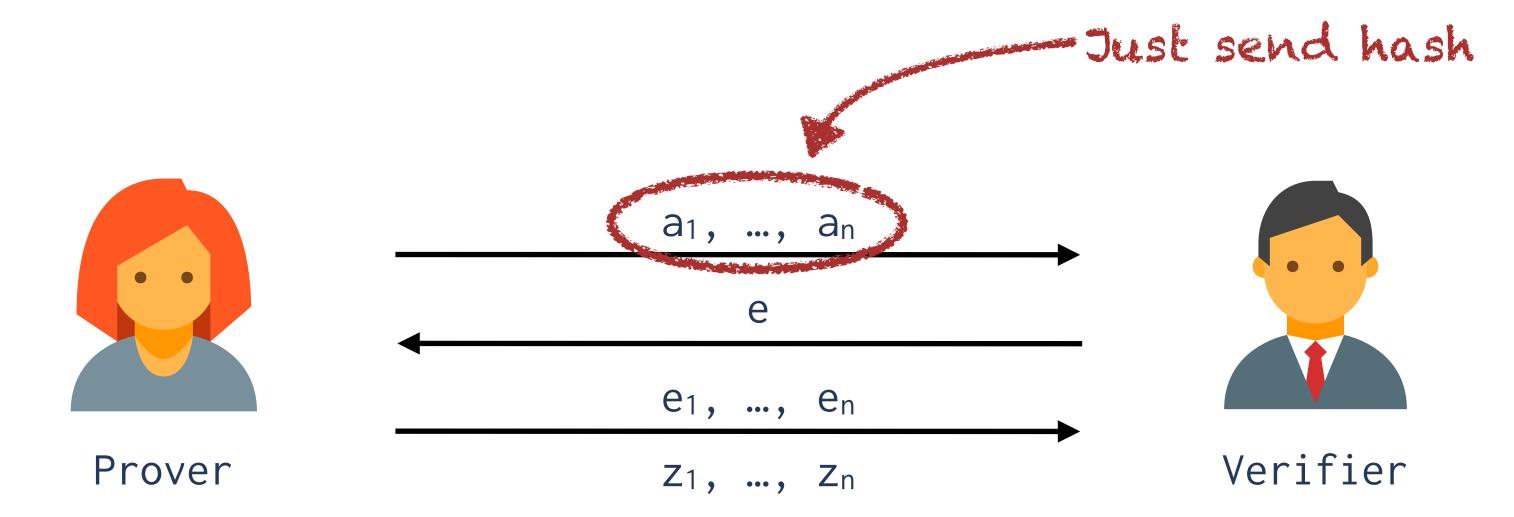
$$(a_1, e_1, z_1)$$

•

$$(a_n, e_n, z_n)$$

26

Our Construction



**a**<sub>1</sub>

$$(a_{2}, e_{2}, z_{2})$$

$$\vdots$$

$$(a_{n}, e_{n}, z_{n})$$

$$(a_{1}, e_{1}, z_{1})$$

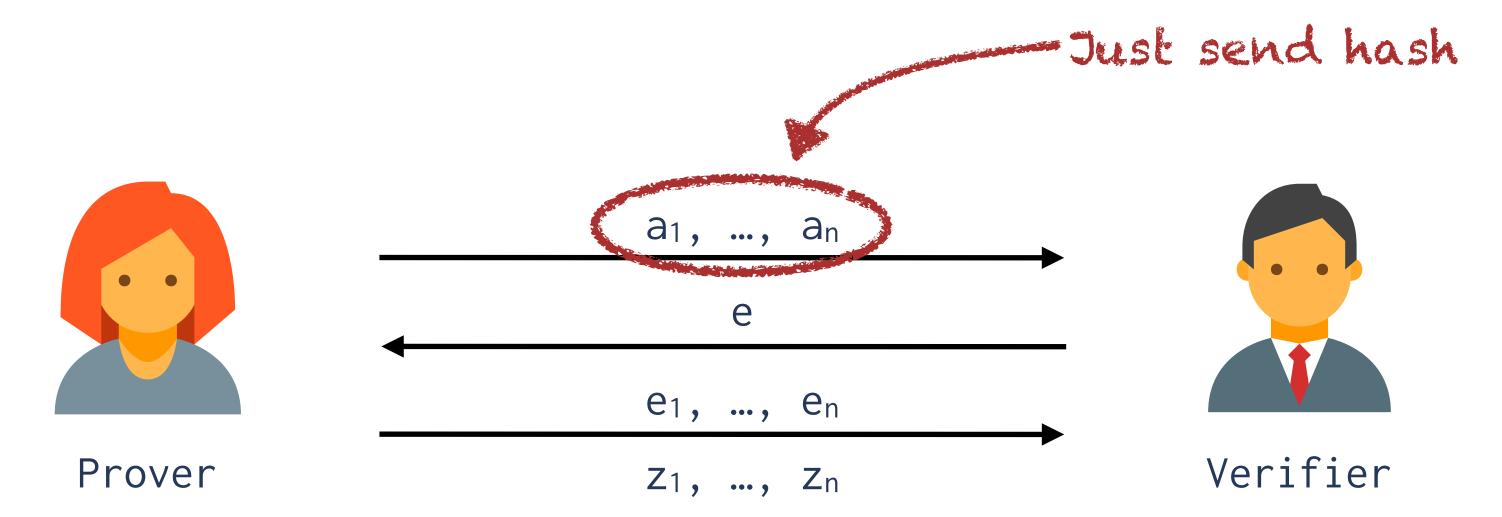
$$\vdots$$

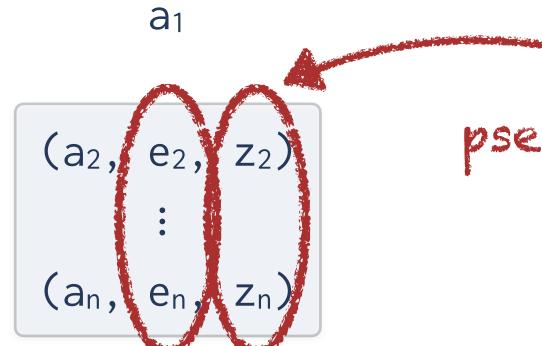
$$(a_{n}, e_{n}, z_{n})$$

$$e = e_1 + ... + e_n$$

27

Our Construction





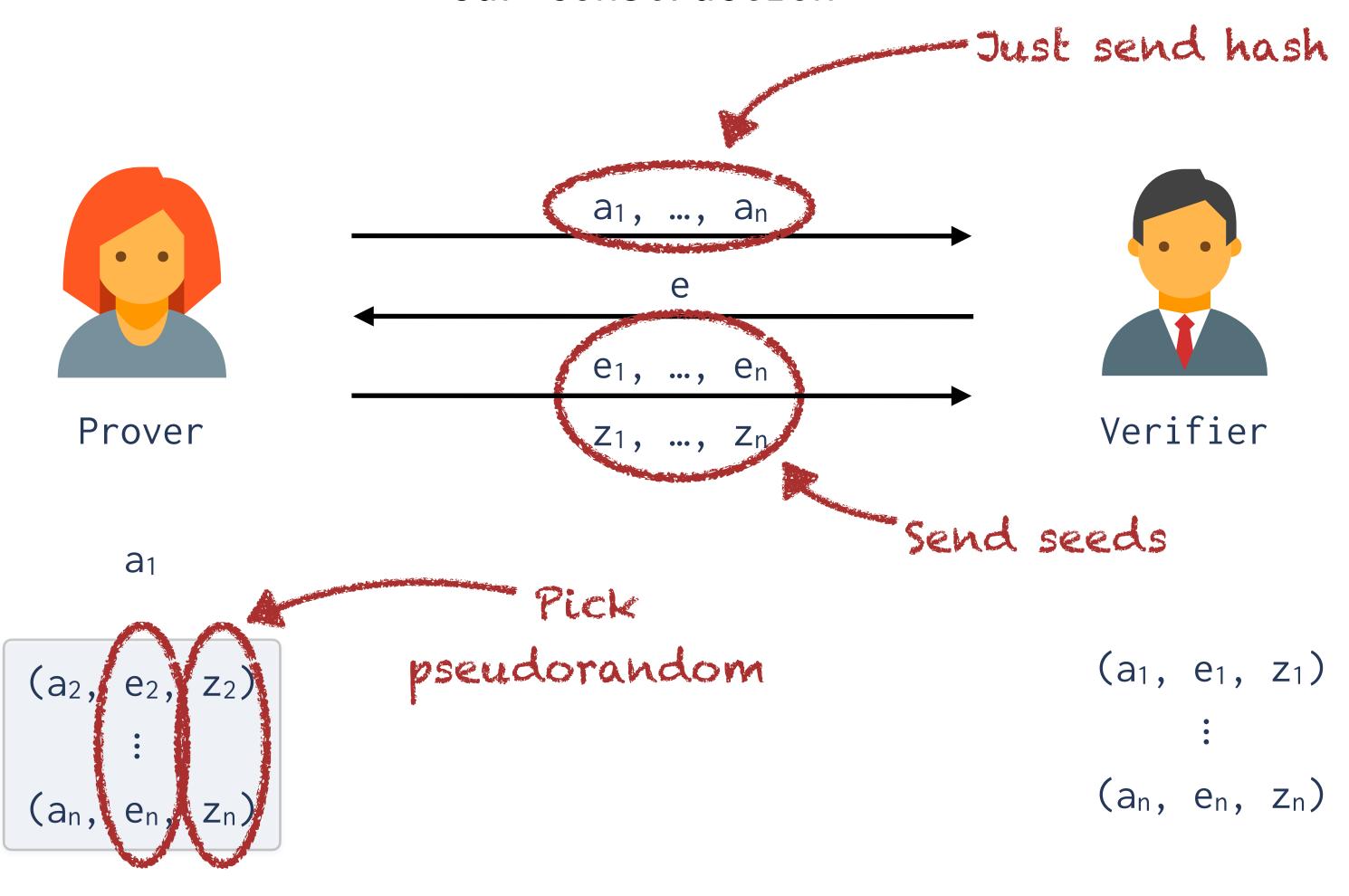
Pick
pseudorandom

$$(a_1, e_1, z_1)$$
 $\vdots$ 
 $(a_n, e_n, z_n)$ 

$$e = e_1 + ... + e_n$$

28

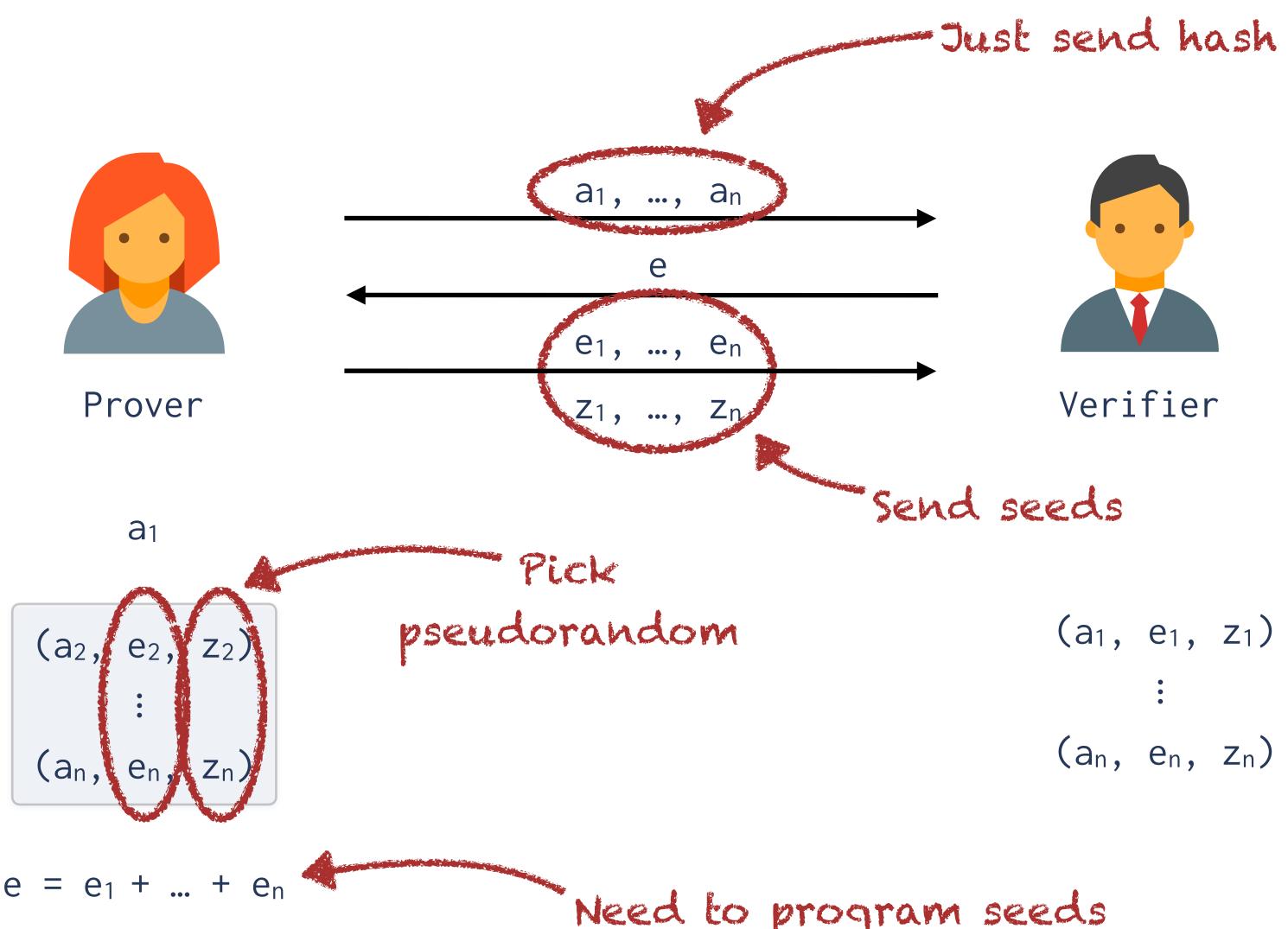
Our Construction



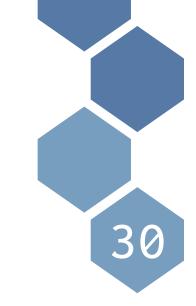
 $e = e_1 + ... + e_n$ 

Our Construction

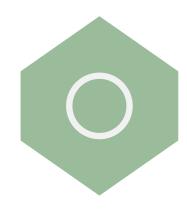




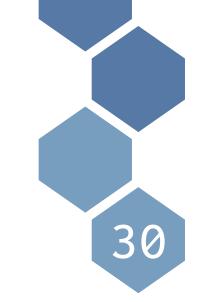
Need to program seeds

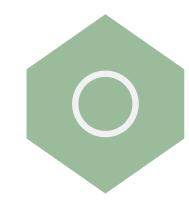






Privately programmable PRFs are hard Require iO or FHE

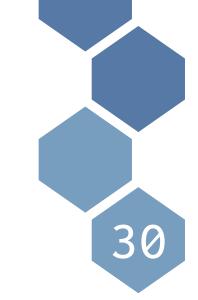


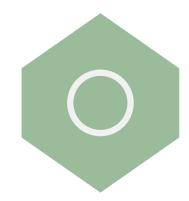


Privately programmable PRFs are hard Require iO or FHE



Know programming location during key generation Massively simplifies the problem

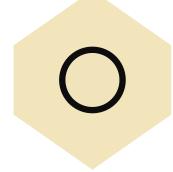




Privately programmable PRFs are hard Require i0 or FHE



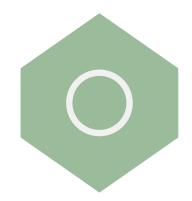
Know programming location during key generation
Massively simplifies the problem



Can be constructed from one-way functions

Via distributed point functions

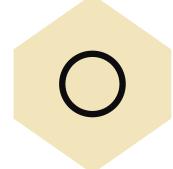




Privately programmable PRFs are hard Require i0 or FHE



Know programming location during key generation
Massively simplifies the problem



Can be constructed from one-way functions

Via distributed point functions (and explainable samplers)



### Questions?