
Batch Arguments to NIZKs from One-Way Functions

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Non-Interactive Zero Knowledge Argument (NIZK) [GMR85, BFM88]

CRS (common reference string)

Prover
 x, w

π

Verifier
 x

For an NP
Language L .



Non-Interactive Zero Knowledge Argument (NIZK) [GMR85, BFM88]

CRS (common reference string)

Prover
 x, w

π

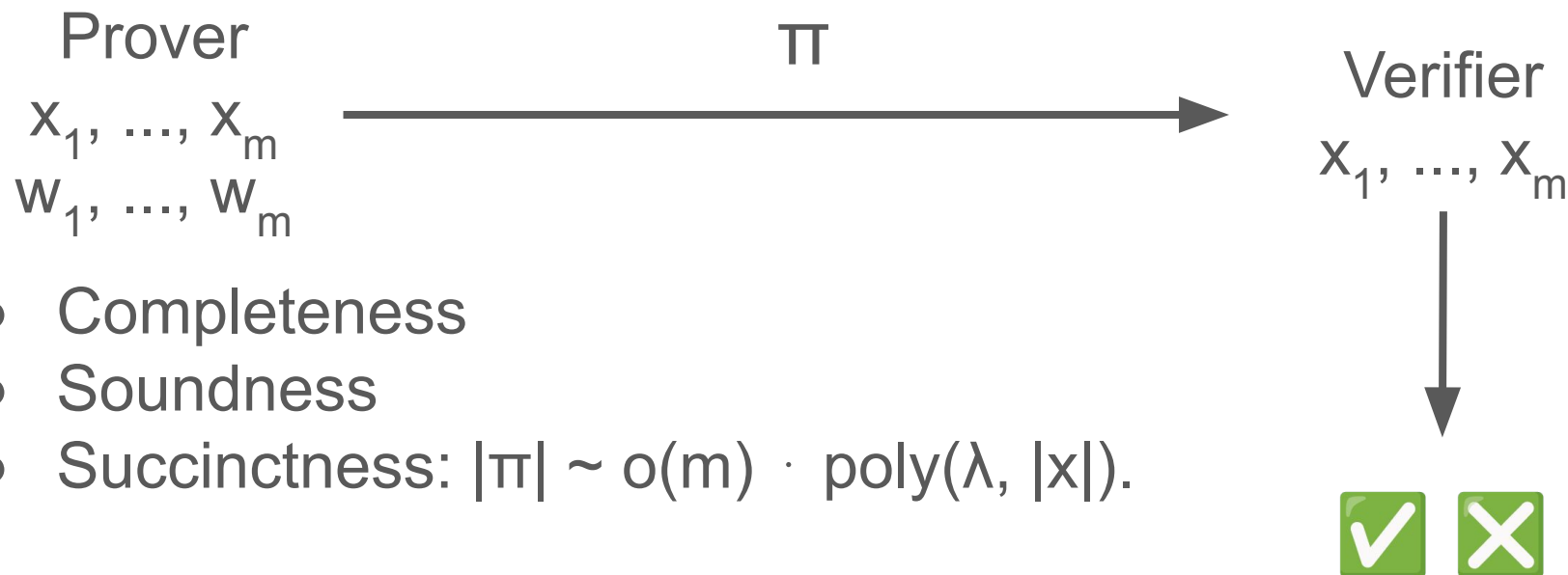
Verifier
 x

- Completeness: Honest proofs verify.
- Soundness: False statements don't verify.
- Zero-Knowledge: Proofs can be simulated without witnesses.



Batch Argument (BARG) [CJJ21a]

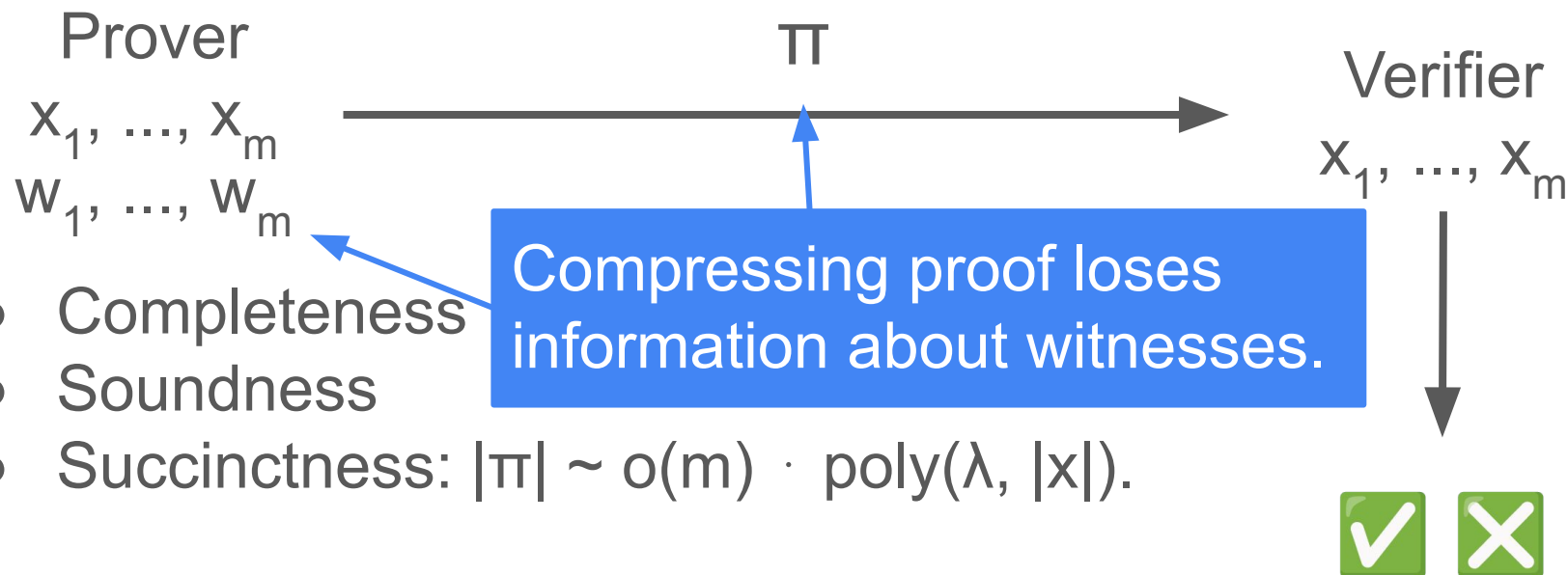
CRS (common reference string)



- Completeness
- Soundness
- Succinctness: $|\pi| \sim o(m) \cdot \text{poly}(\lambda, |x|)$.

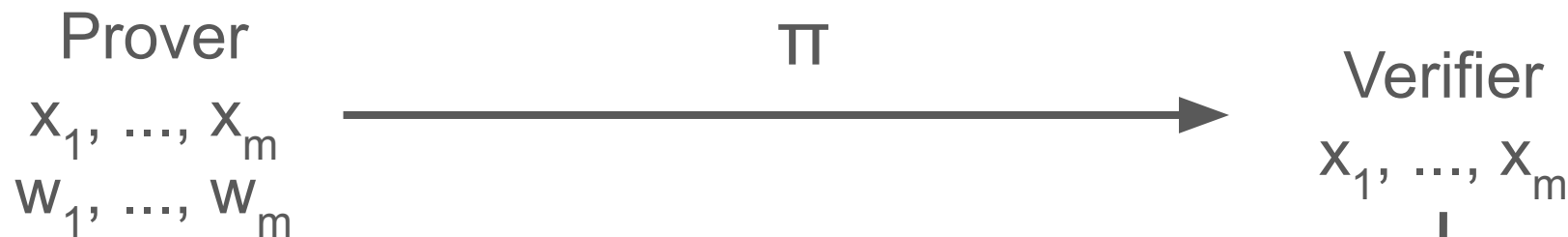
Batch Argument (BARG) [CJJ21a]

CRS (common reference string)



Batch Argument (BARG) [CJJ21a]

CRS (common reference string)



Existing constructions use NIZK techniques.
[CJJ21b, WW22, DGKV22, CGJ⁺23]



Can we get NIZKs from BARGs?

YES!

Previous and Concurrent Works

- [CW23]: NIZK from BARG, **Local PRG, dual-mode commitment.**
- [BKP⁺23]: NIZK from BARG, **dual-mode commitment.**
 - Concurrent update to [BKP⁺23]: NIZK from BARG, OWF.

Can we do it without **additional assumptions**?

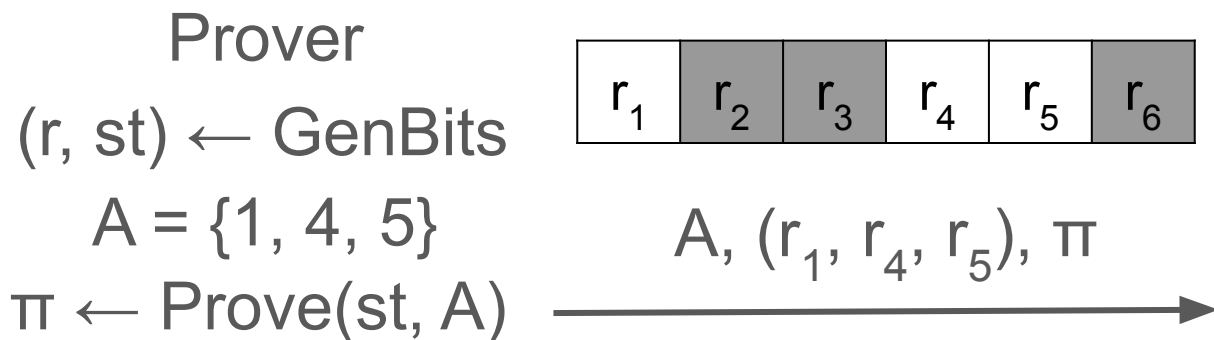
Main Technical Result

(Adaptively-sound) BARG + OWF \Rightarrow NIZK

Construct hidden-bits generator and apply
[KMY20] to get NIZK.

Hidden-Bits Generator

[FLS90, QRW19, KMY20]



- Binding: r_A corresponds to $r \in \text{supp}(\text{GenBits})$.
- Hiding: Hidden bits remain pseudorandom.
- Sparsity: Density of $\text{supp}(\text{GenBits})$ in $\{0, 1\}^m$ is low.

Verifier



Our Construction

Our Construction: Overview

PRG + BARG \Rightarrow HBG

Why should this work?

Our Construction: Overview

Idea: Use **PRG** and **BARG** to build HBG.

- Binding \Rightarrow BARG
- Hiding \Rightarrow PRG

Our Construction: Overview

Idea: Use **PRG** and **BARG** to build HBG.

- Seed: PRG seed
- Output: PRG output
- Proof: BARG proof + openings

Template of [KMY20] for SNARGs, [CW23] for BARGs.

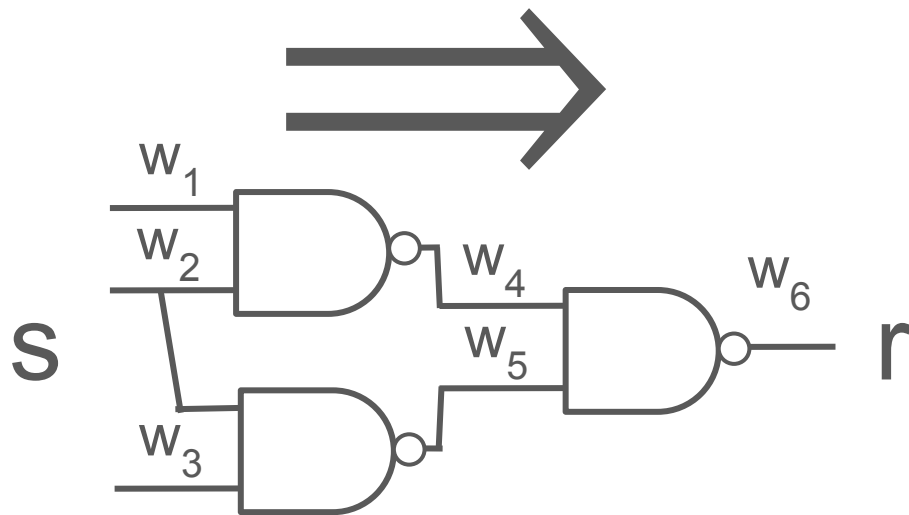
Our Construction: GenBits

Hidden-bits string

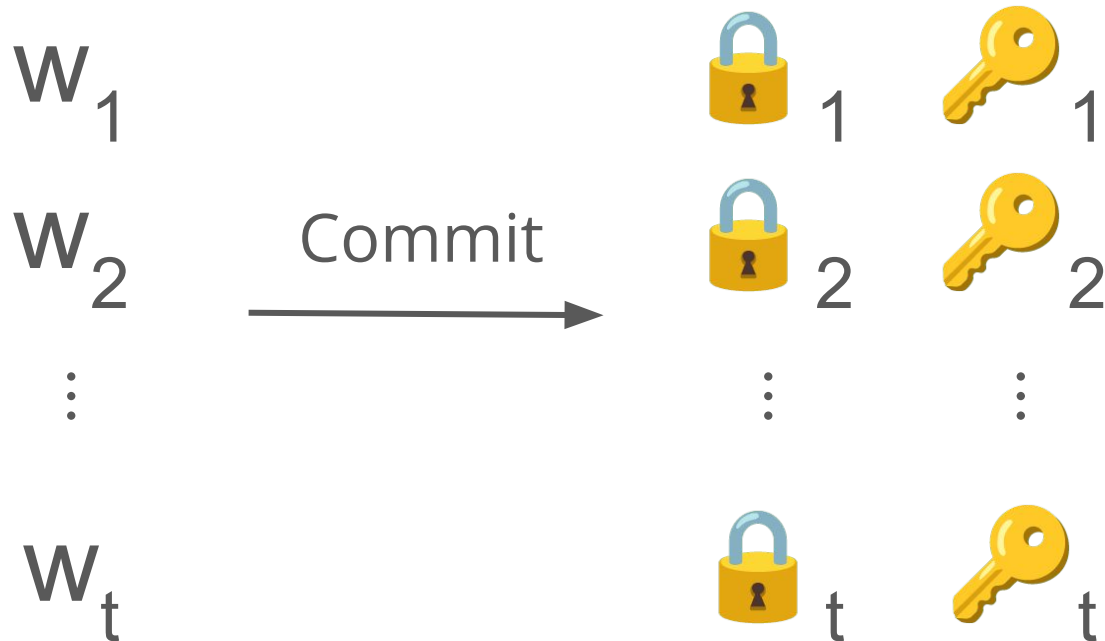
- $s \leftarrow \{0, 1\}^n$
- $r = \text{PRG}(s)$
- $st = s$

Prover's internal state

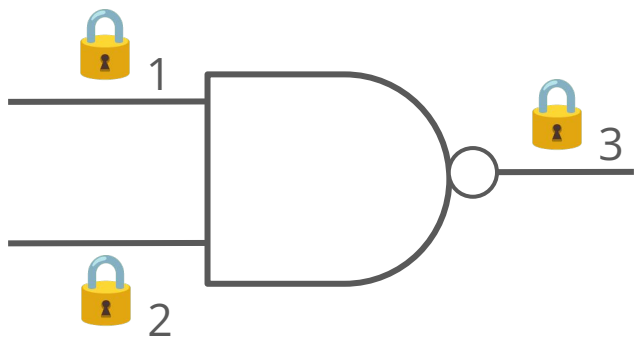
(LR) PRG



Our Construction: Prove



Our Construction: Prove




BARG over all
internal gates.

Statement:

- $\text{crs}_1 \text{ crs}_2 \text{ crs}_3$
- 



Witness:

- Wire values w_1, w_2, w_3
- 

Check commitments and
 $w_3 = \text{NAND}(w_1, w_2)$.

Our Construction: Prove

Output:

- Commitments  ₁ ...  _t
- Openings for revealed output wires.
- BARG proof π_{BARG}

Security: Binding

Binding: r_A corresponds to $r \in \text{supp}(\text{GenBits})$.

Security: Binding

Binding: r_A corresponds to $r = \text{PRG}(s)$.






Check the PRG circuit is consistently evaluated.

Recall prover gives:

- Bit commitments to wire values.
- BARG proof each gate is correctly evaluated.
- Openings for output bits.






Security: Hiding

Proof π :

- Commitments  ₁ ...  _t 
- Openings for revealed output wires. 
- BARG proof π_{BARG} 

Security: Hiding

Proof π :

- Commitments  ₁ ...  _t 
- Openings for revealed output wires. 
- BARG proof π_{BARG} 

PRG leakage-resilience

Summary

- We constructed a hidden-bits generator from:
 - Adaptively sound BARG
 - Leakage-resilient wPRF (\Leftarrow OWF)
 - One-time dual-mode bit commitment (\Leftarrow OWF)
- We get NIZK from the same assumptions [KMY20].

Additional Results

Stronger than
OWF.

Somewhere-sound BARG + PKE \Rightarrow NIZK

Weaker than
adaptive
soundness.

Open Problems

Can we get NIZK with weaker assumptions?

- NIZK from index BARG [CJJ21b]?
- NIZK from non-adaptively sound BARG?

Thank you!

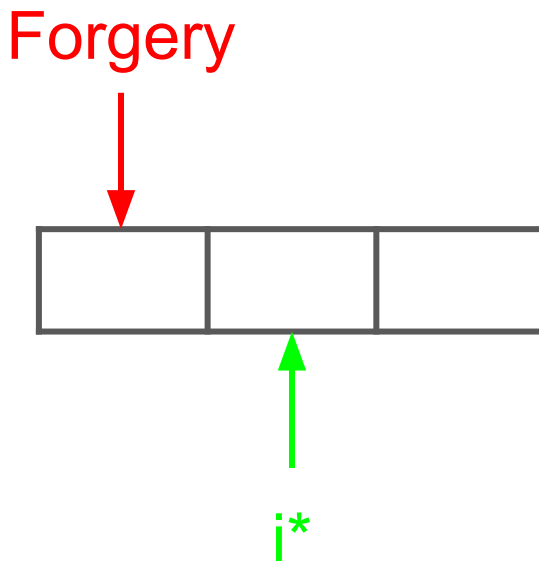
<https://eprint.iacr.org/2023/1938.pdf>

Getting Adaptively-Sound BARG

Somewhere-soundness

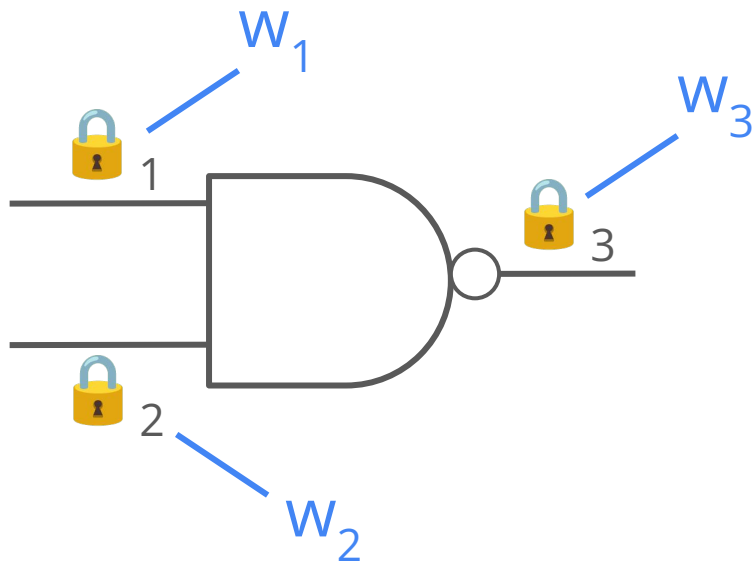
+ sub-exponential index hiding

⇒ Adaptive soundness.



Security: Binding

Binding: r_A corresponds to $r = \text{PRG}(s)$.

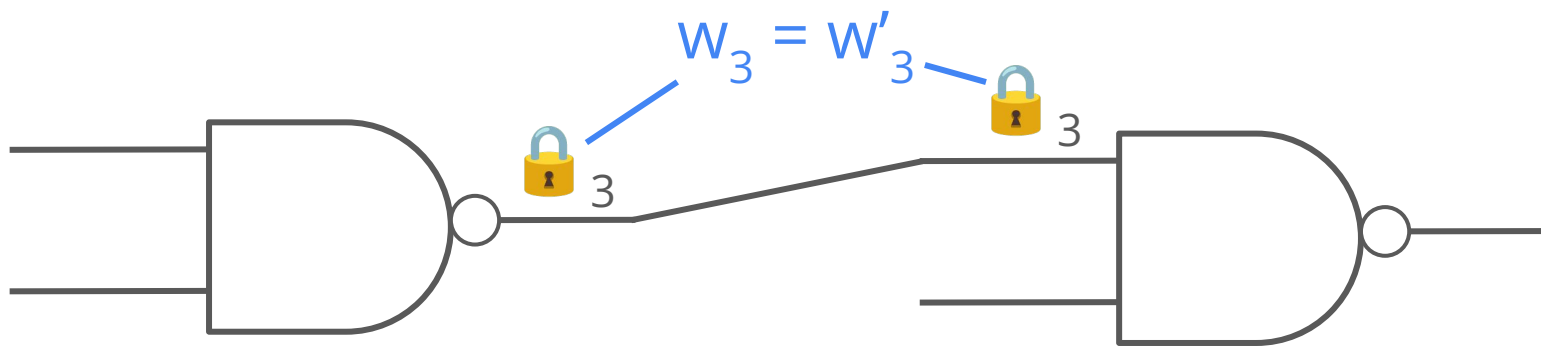


$$w_3 = \text{NAND}(w_1, w_2)$$

BARG soundness \Rightarrow gate consistency

Security: Binding

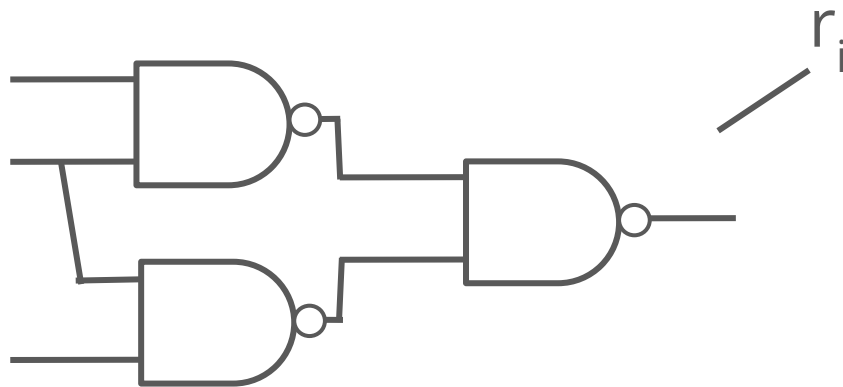
Binding: r_A corresponds to $r = \text{PRG}(s)$.



BC statistical binding \Rightarrow wire consistency

Security: Binding

Binding: r_A corresponds to $r = \text{PRG}(s)$.



Openings \Rightarrow output consistency

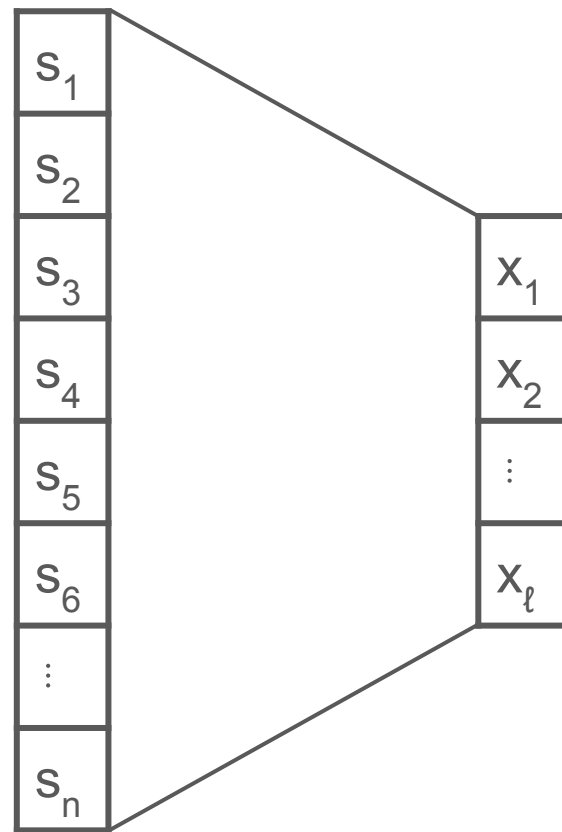
Leakage-Resilient PRG

$\forall \text{ leak: } \{0, 1\}^n \rightarrow \{0, 1\}^\ell$

$(\text{PRG}(s), \text{leak}(s))$

\approx_c

$(\text{Uniform}, \text{leak}(s))$



Leakage-Resilient PRG

$$\forall \text{leak}: \{0, 1\}^n \rightarrow \{0, 1\}^\ell$$

~~(PRG(s), leak(s))~~

Leakage Resilient wPRF

~~(Uniform, leak(s))~~

