

Key-Homomorphic and Aggregate Verifiable Random Functions

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<https://eprint.iacr.org/2024/643.pdf>



Verifiable Random Functions [MRV99]

$$\mathcal{F} : \mathcal{K} \times \{0,1\}^\lambda \rightarrow \mathcal{Y}$$

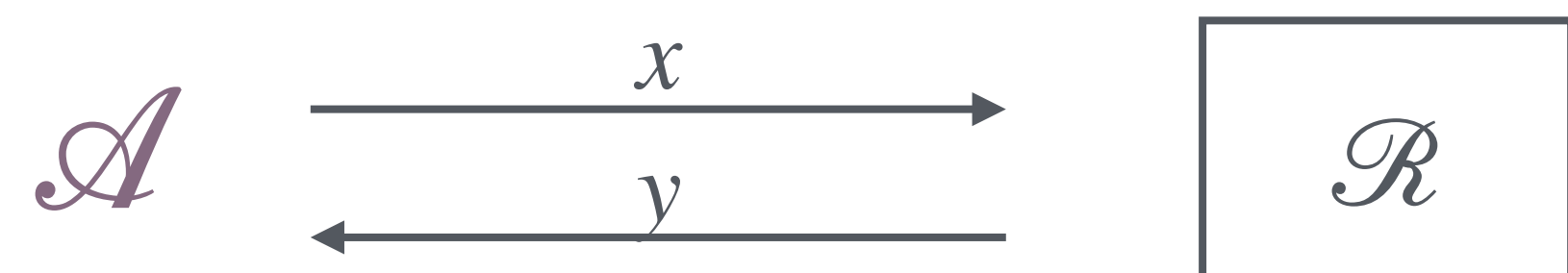
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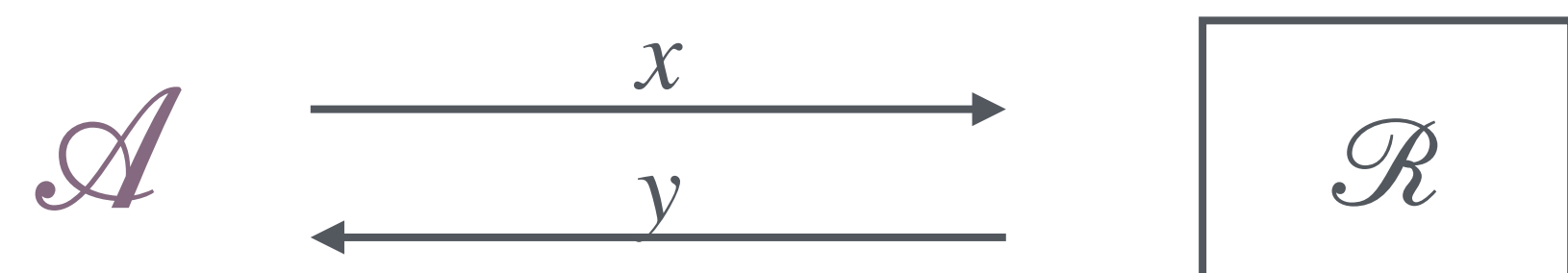
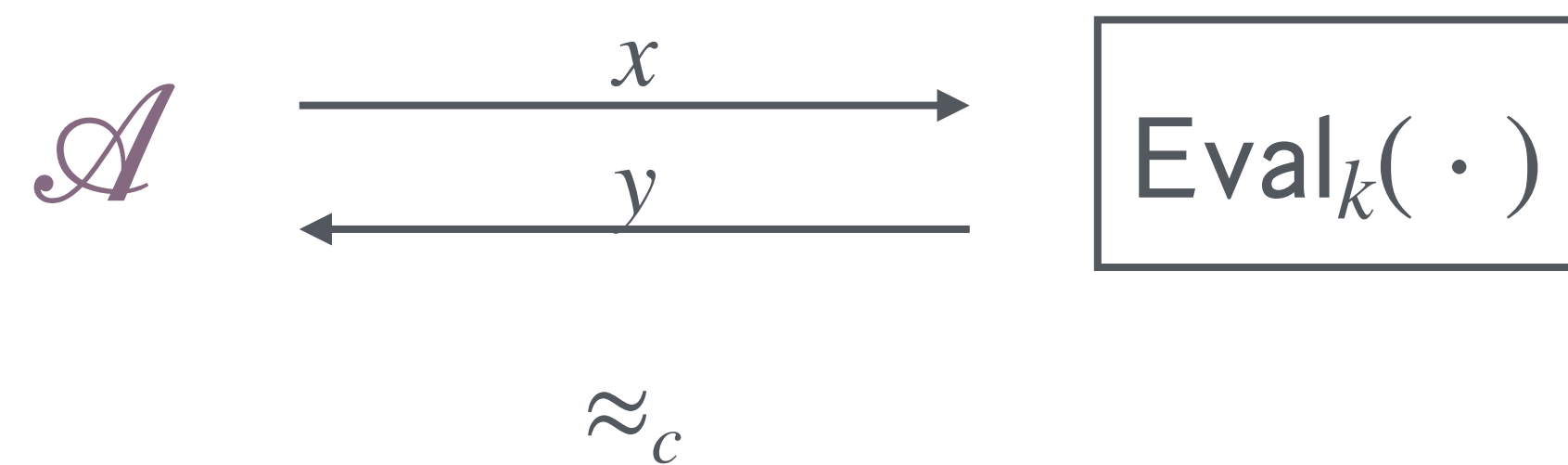
\approx_c



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UNIQUENESS

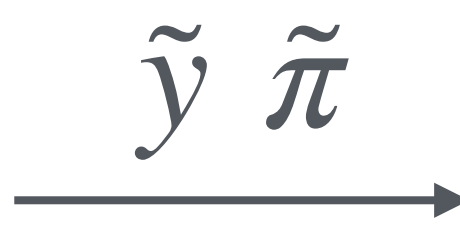
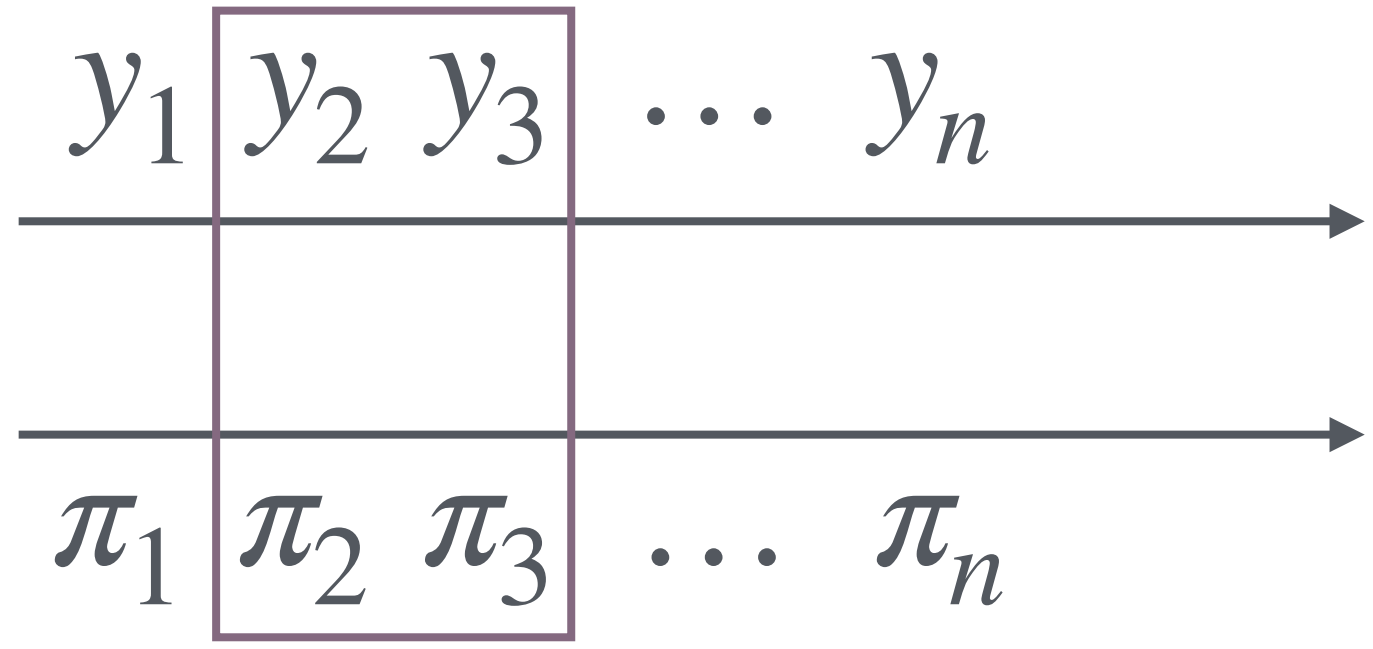
$$\text{Prove}_k(x) \rightarrow \pi$$

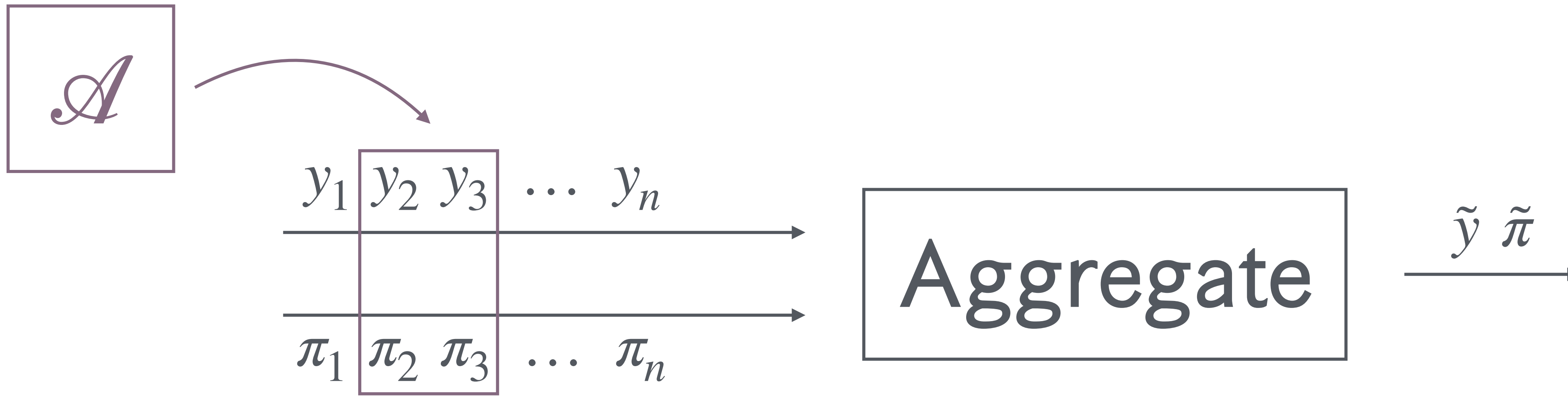
$$\text{Verify}(x, y, \pi) \rightarrow \{0,1\}$$

$$\Pr_{\mathcal{A}} [\text{Verify}(x^*, y^*, \pi^*) = 1 : y^* \neq \text{Eval}_k(x^*)] \approx 0$$

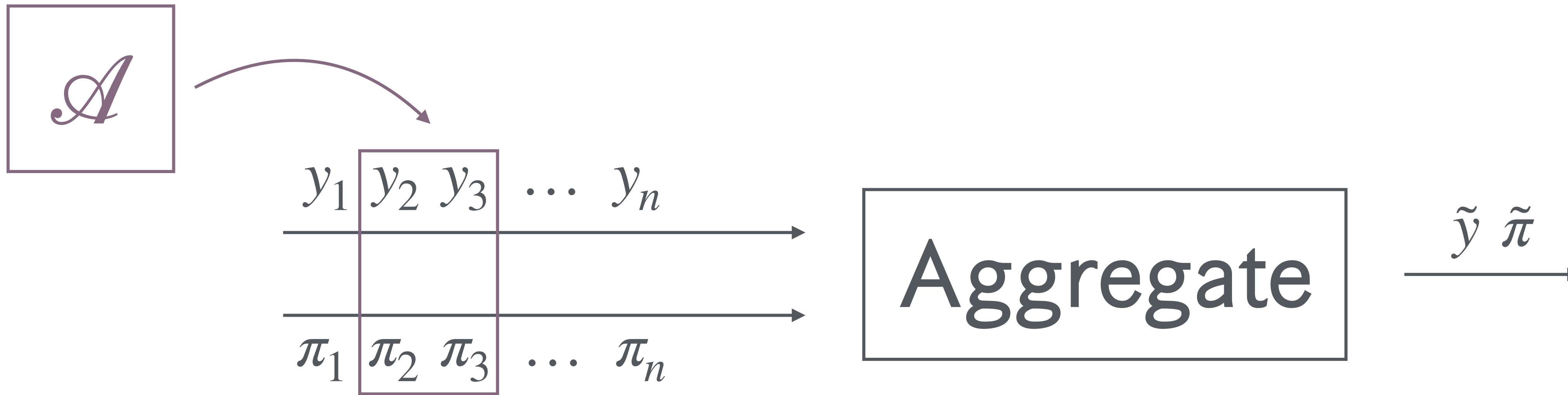
$$(x^*, y^*, \pi^*) \leftarrow \mathcal{A}$$



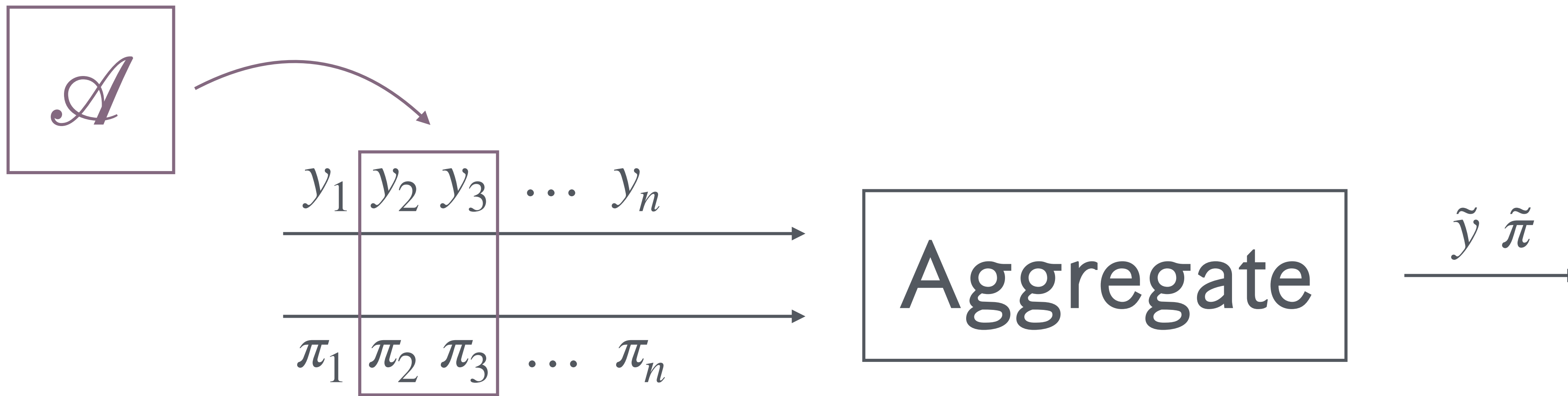




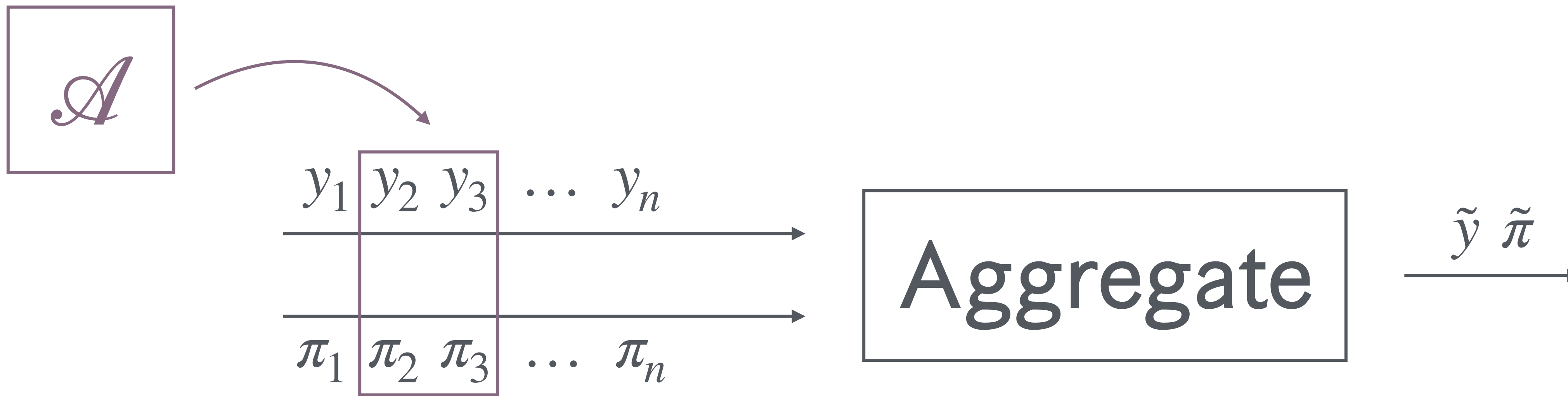
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- Simple **one-round** distributed VRF
- Many other conceivable applications

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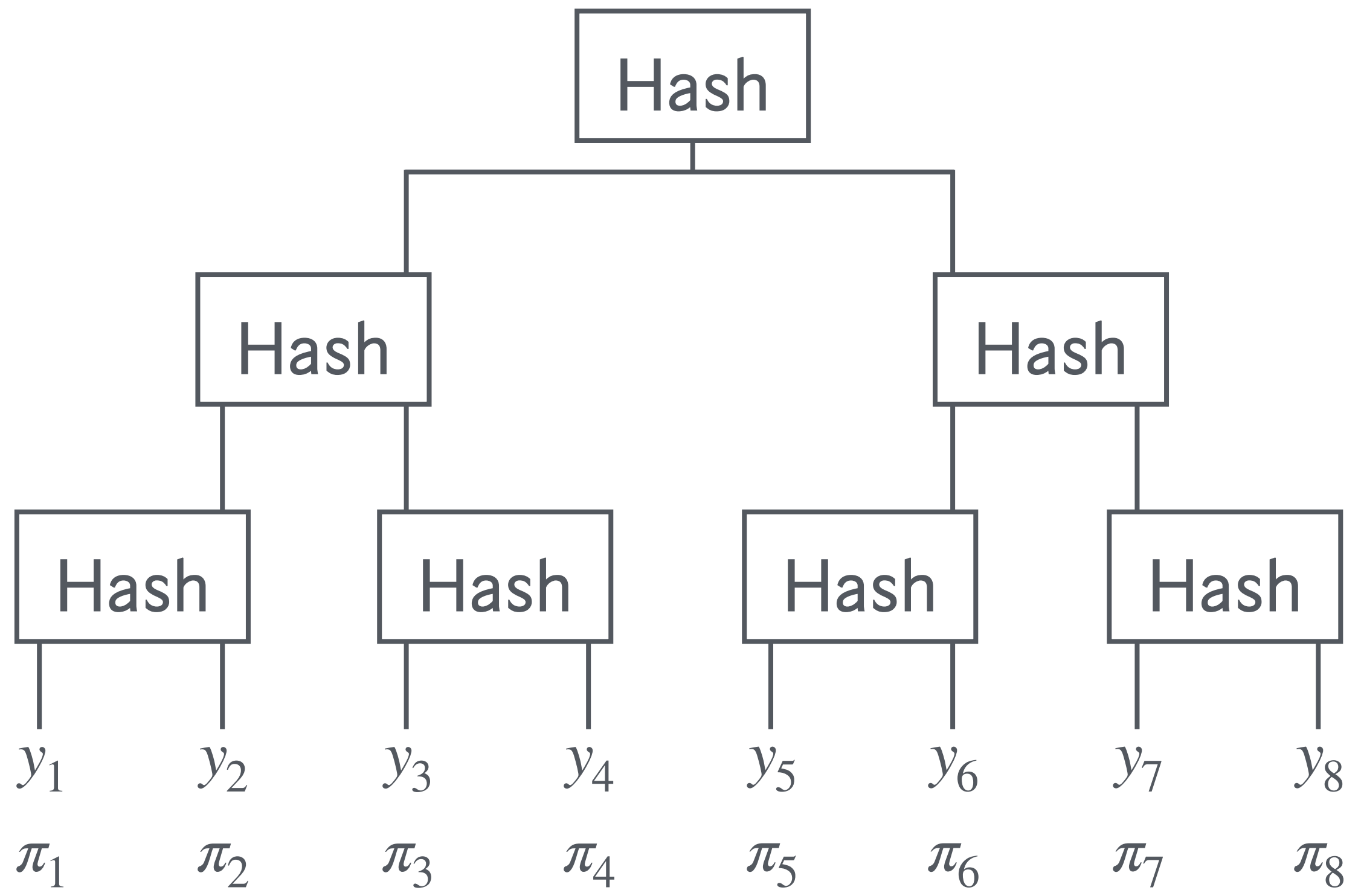
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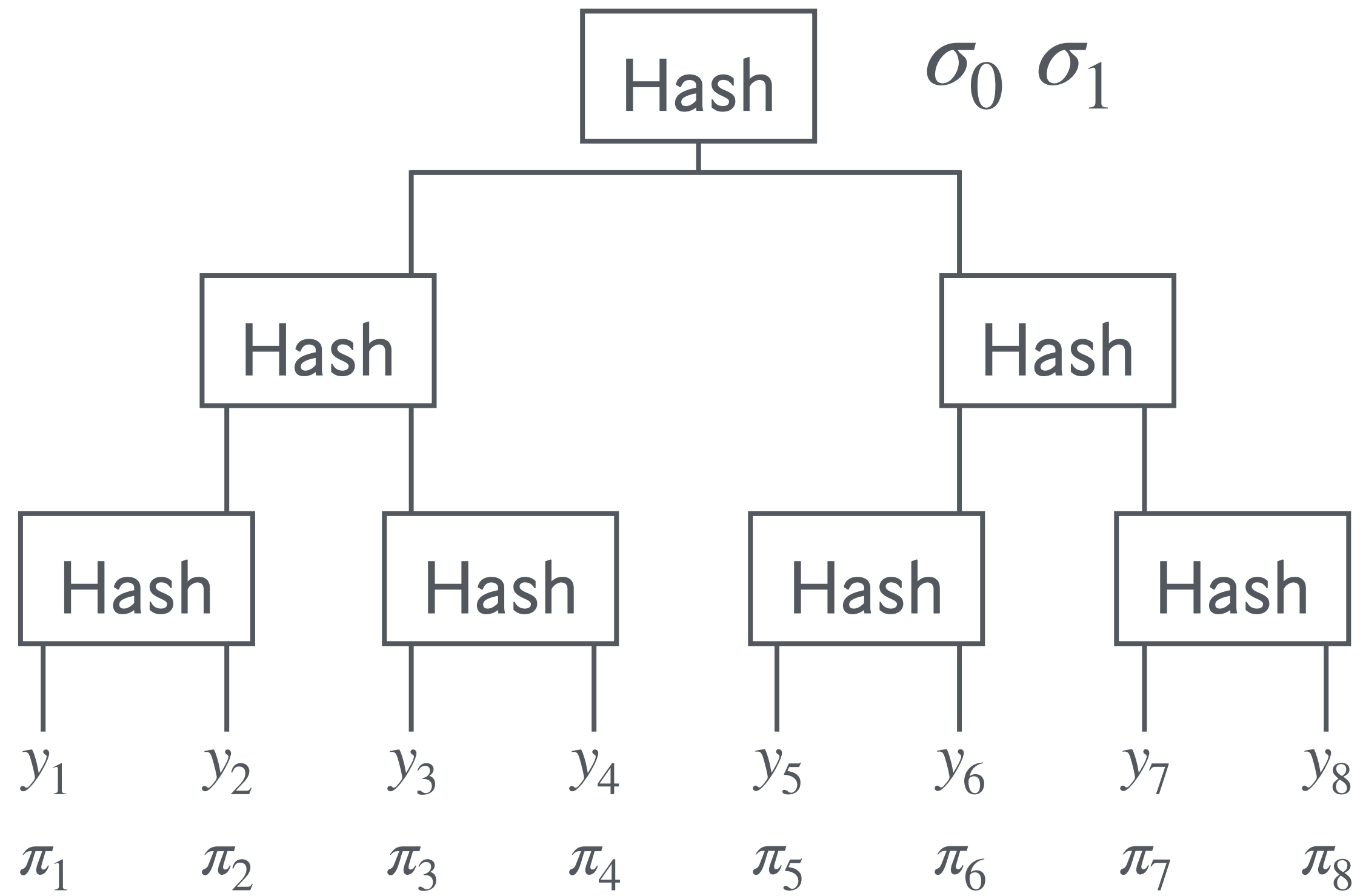
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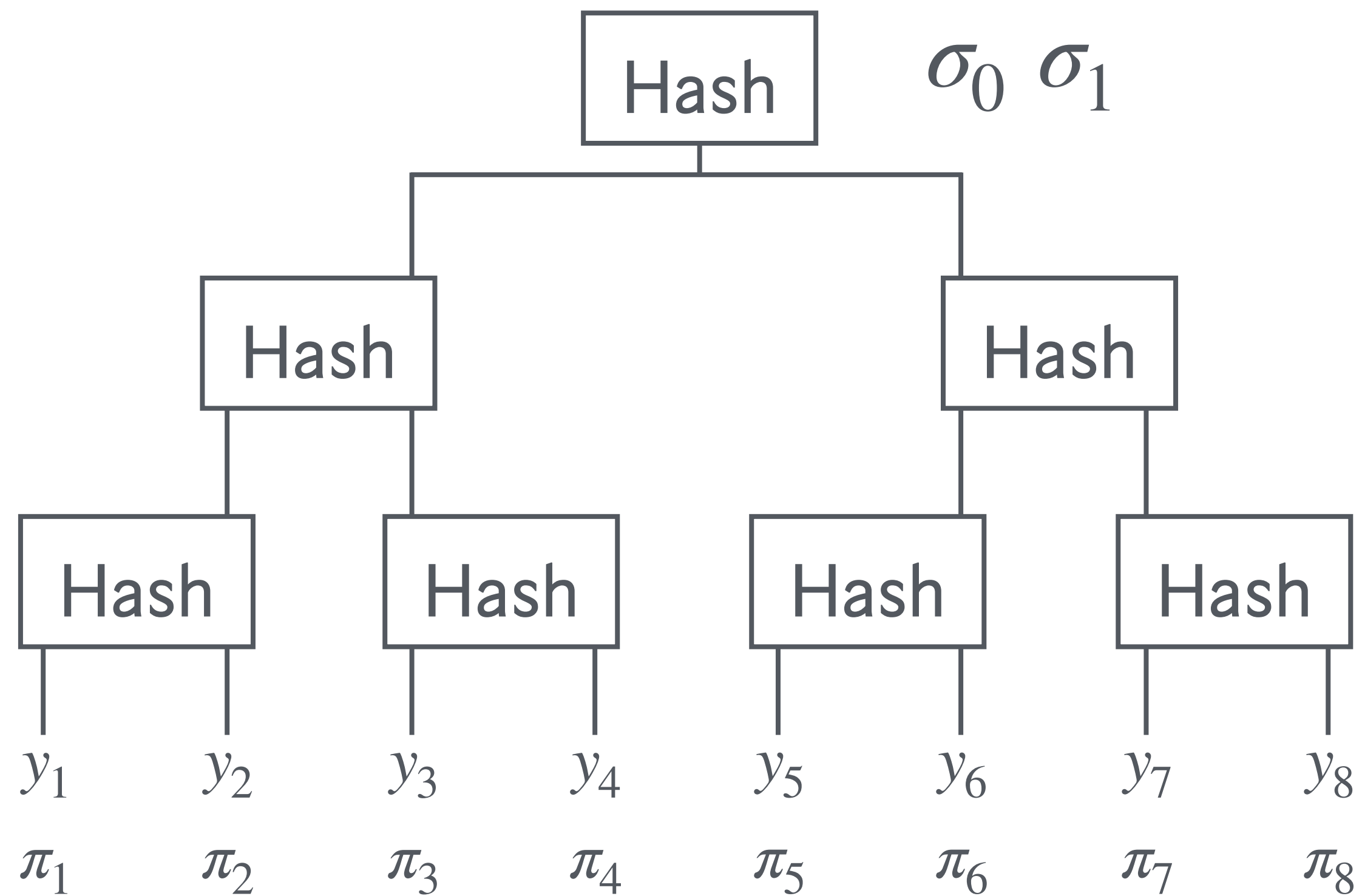


BATCH ARGUMENTS FOR NP



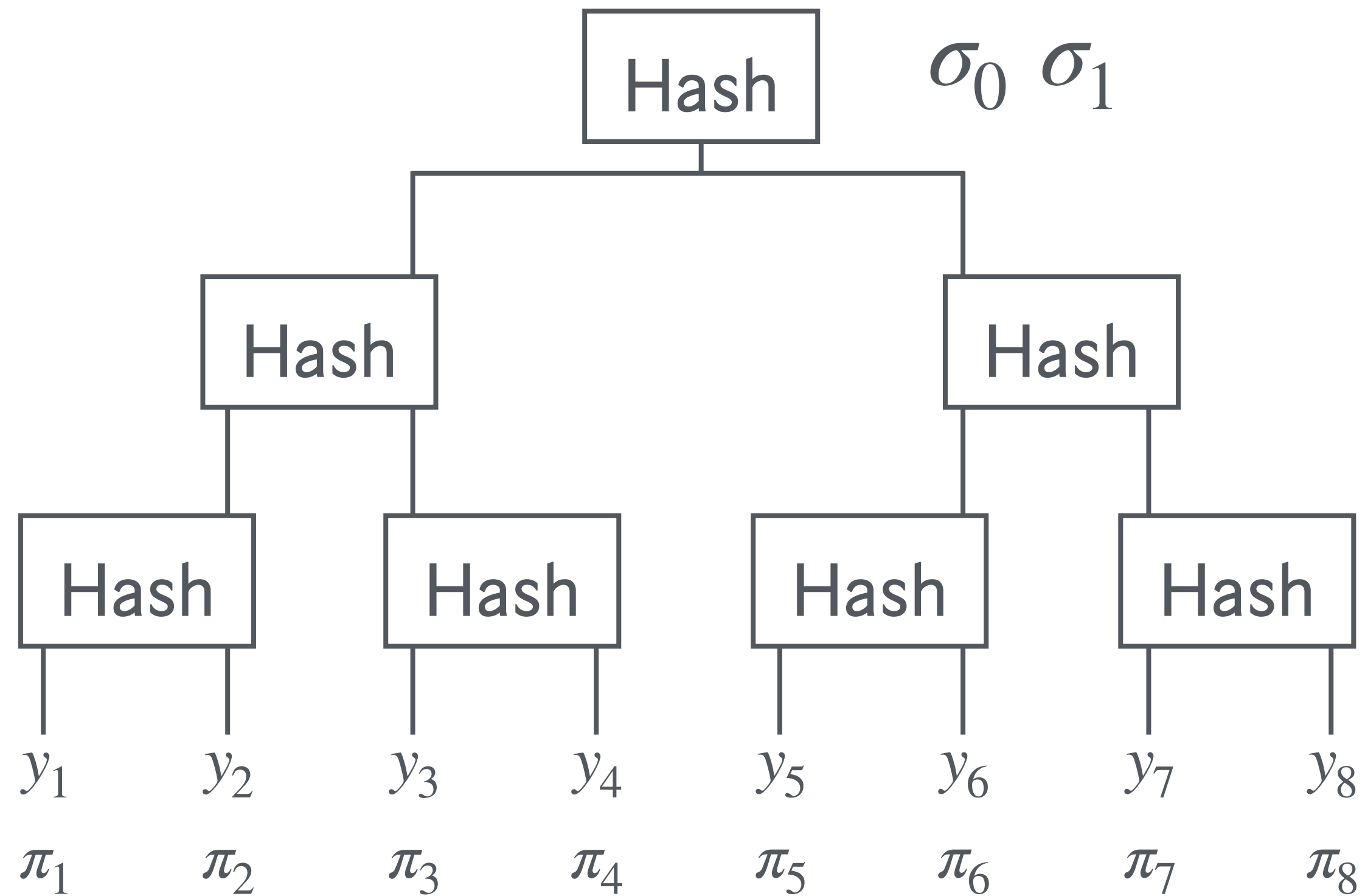
BATCH ARGUMENTS FOR NP

- Compute **two** BARGs for the **same** statement:



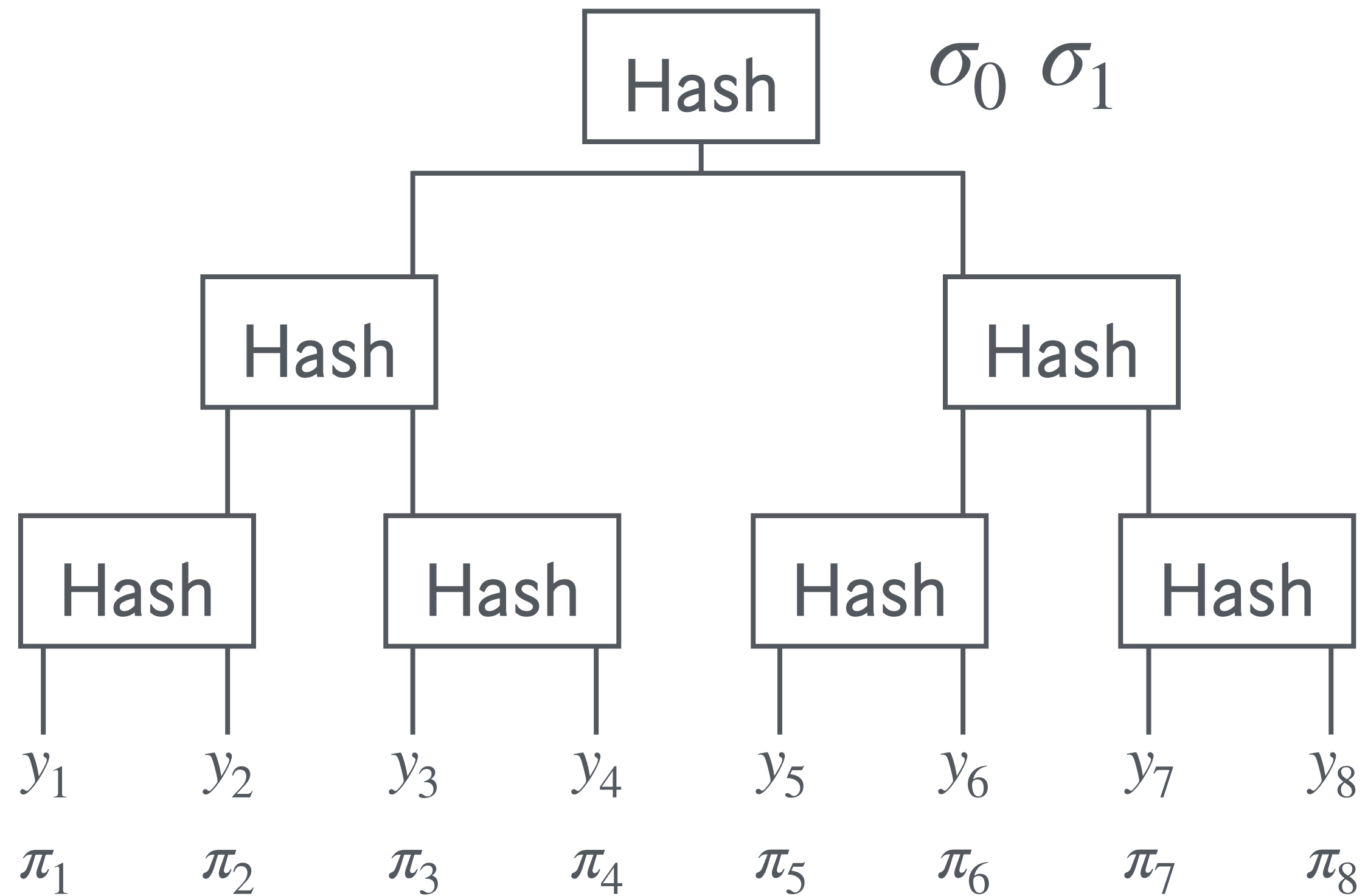
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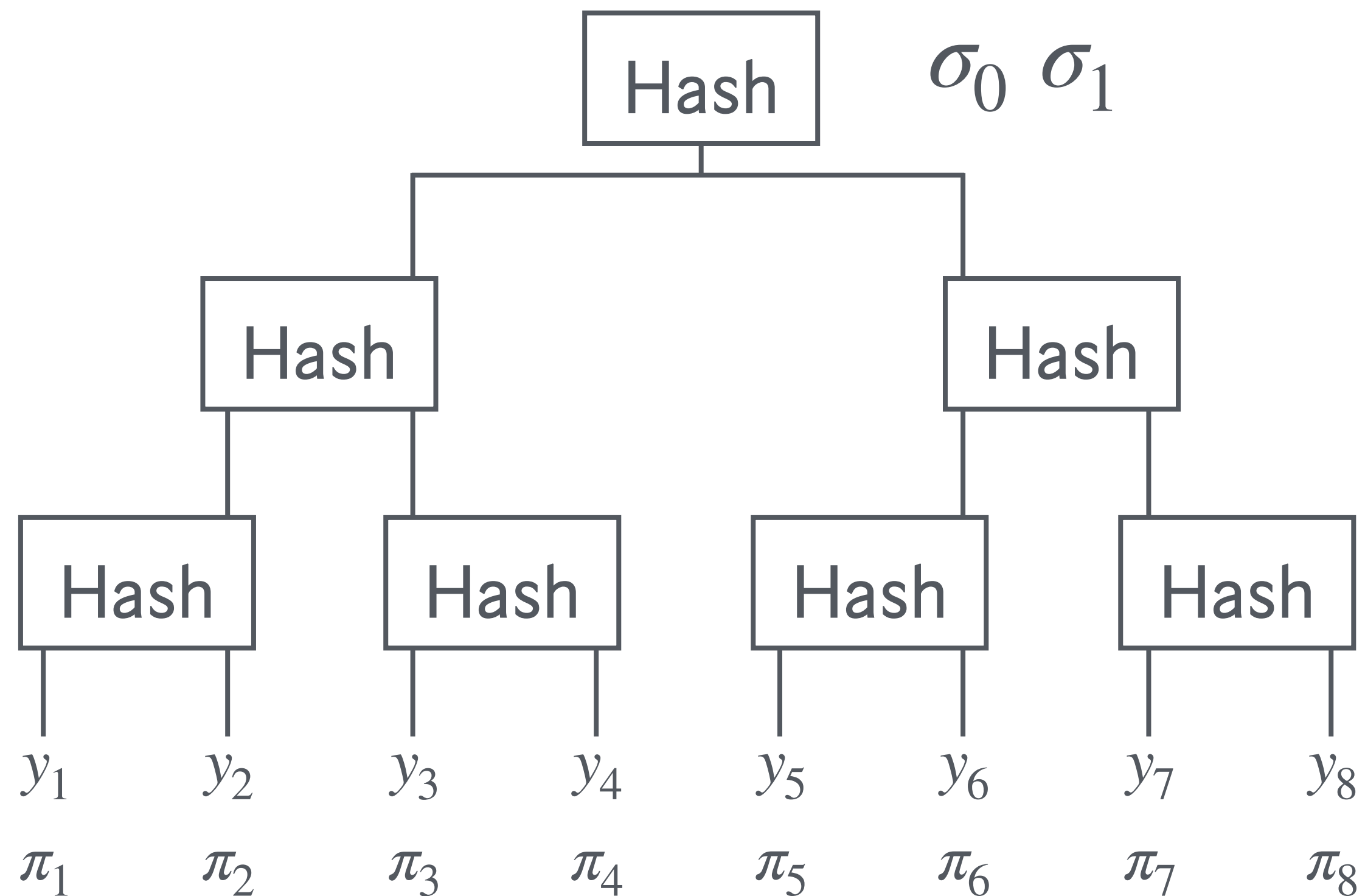


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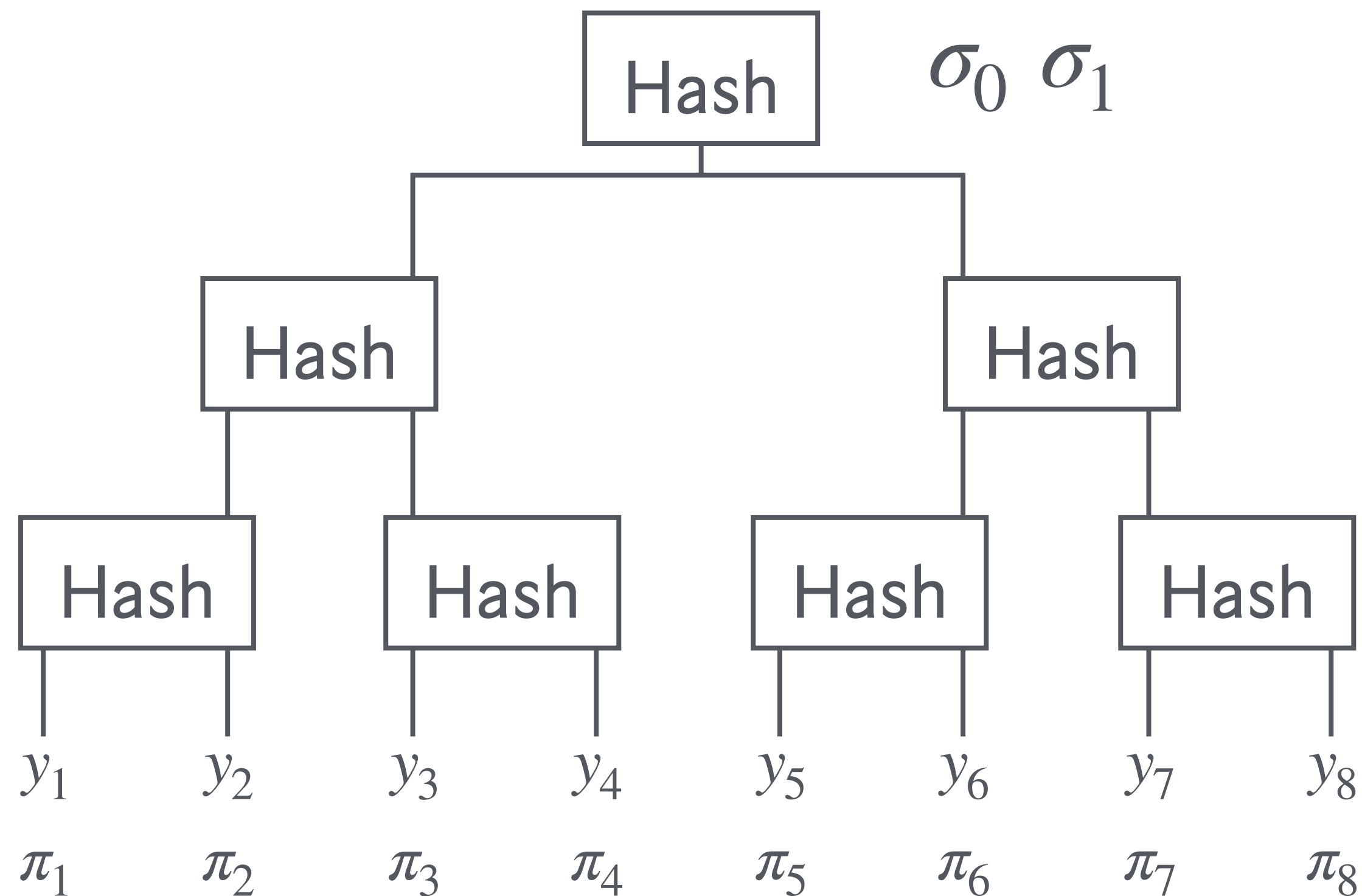
A. p_i is a valid root-to-leaf path

B. π_i is a valid proof for y_i



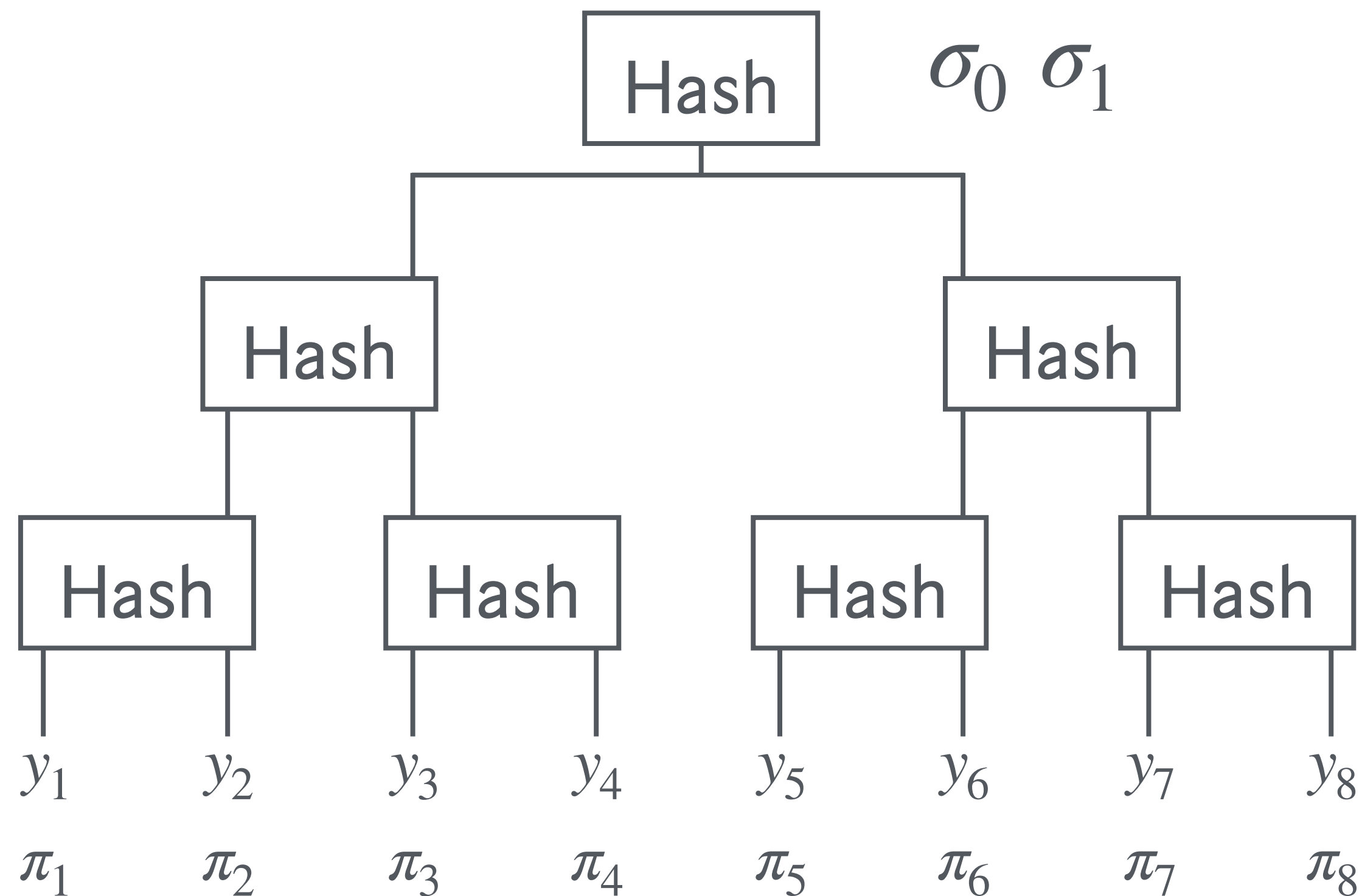
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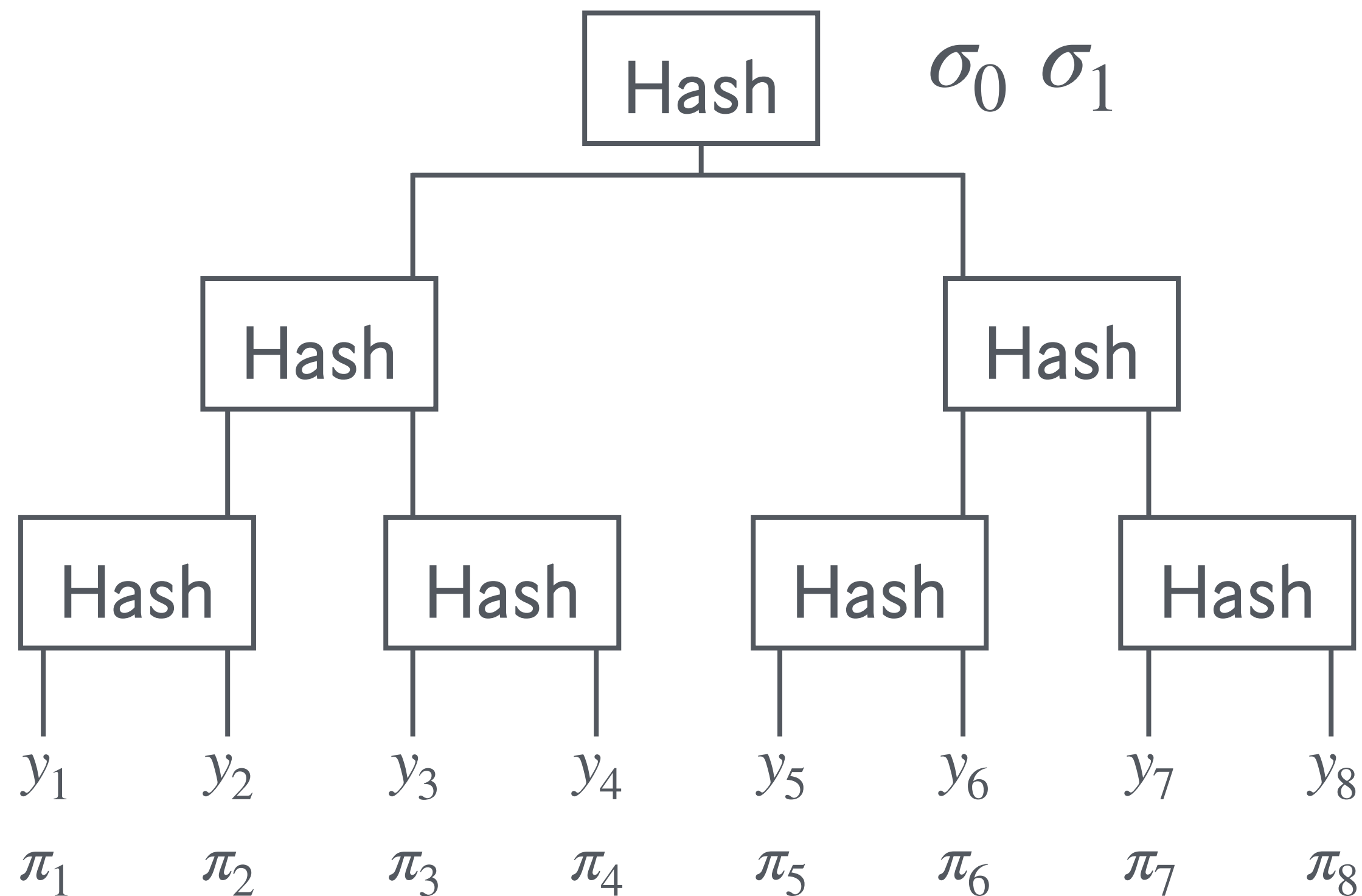
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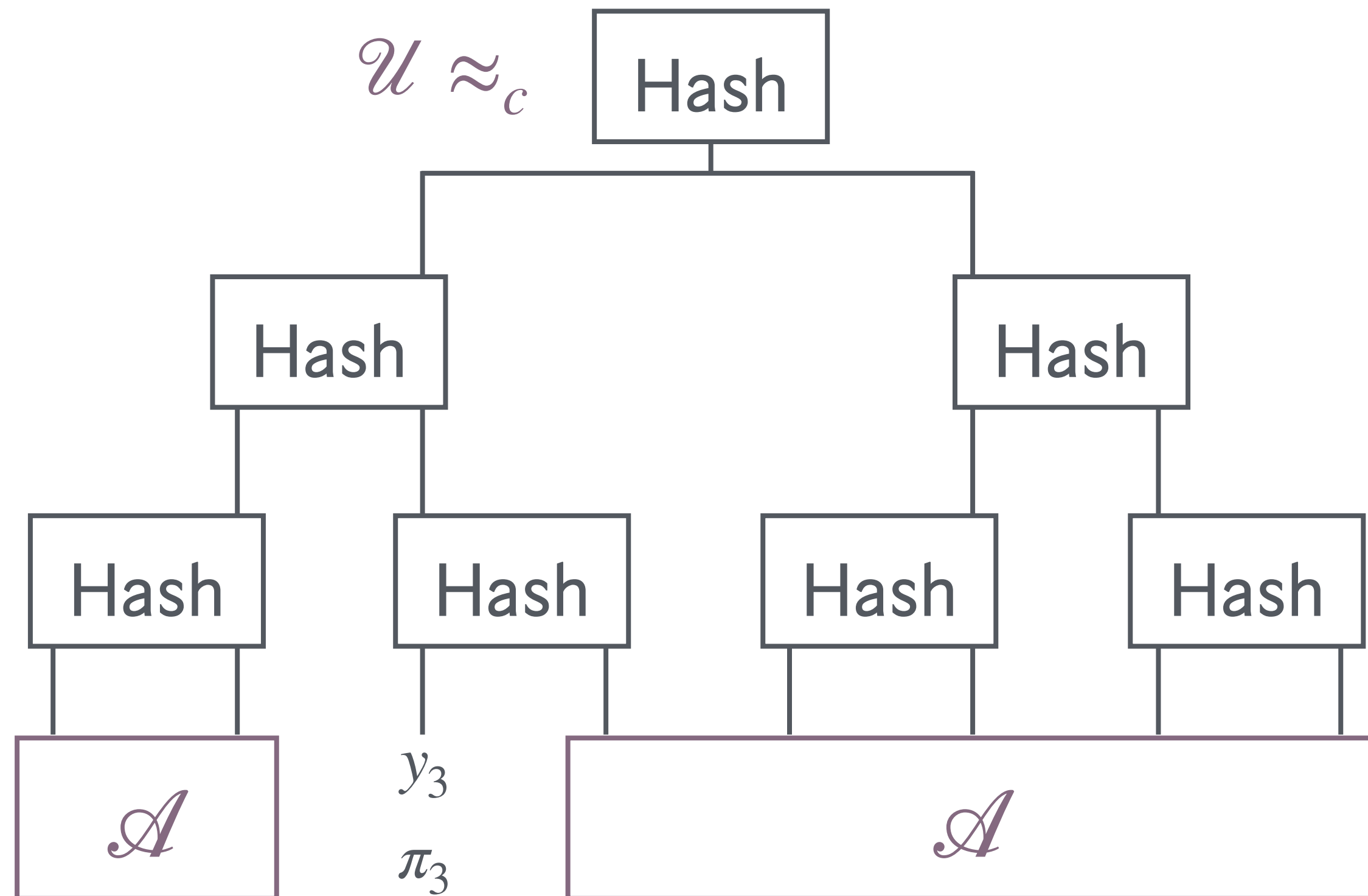


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 2. Index-**hiding**

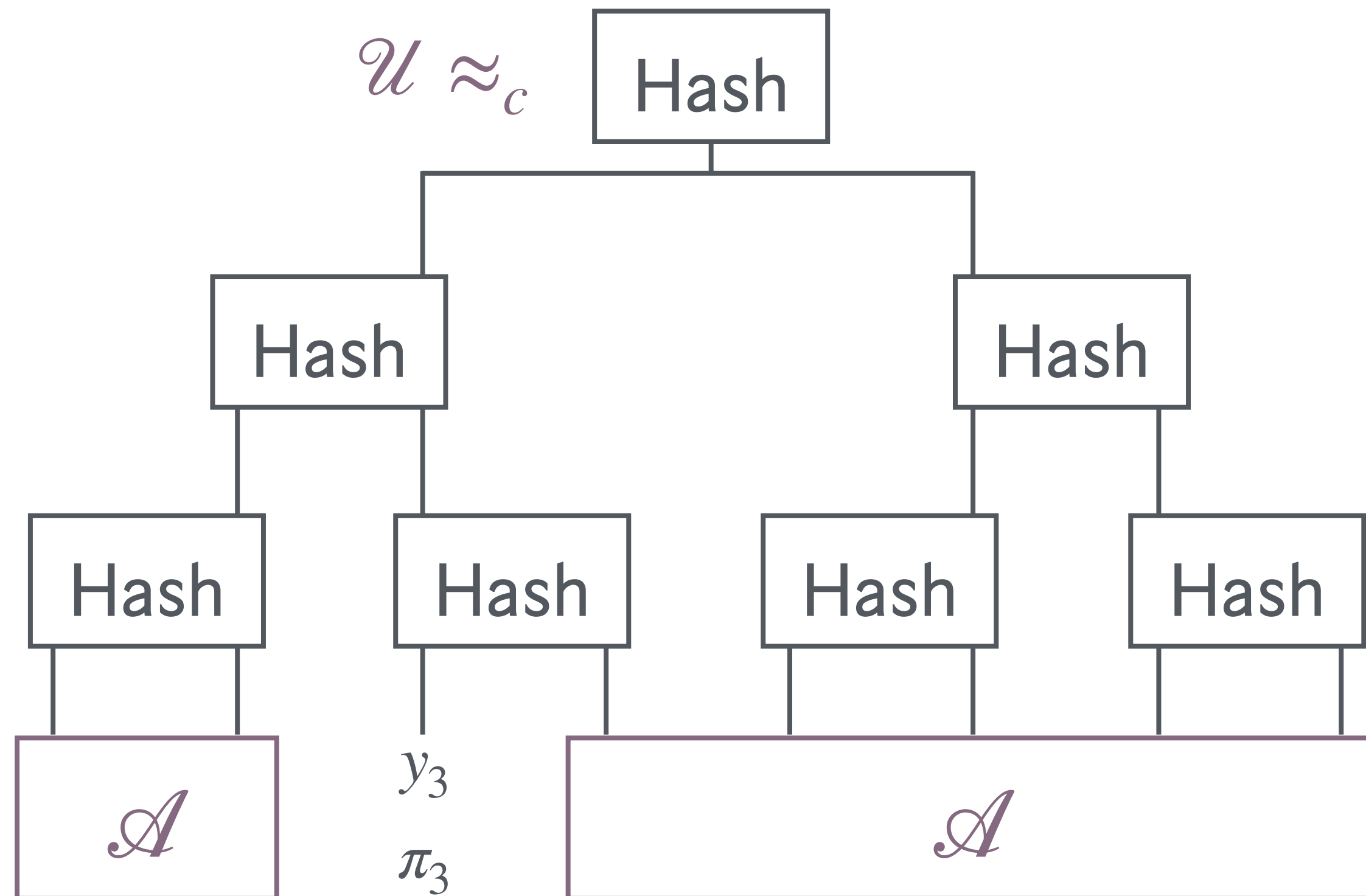


(AGGREGATE) PSEUDORANDOMNESS



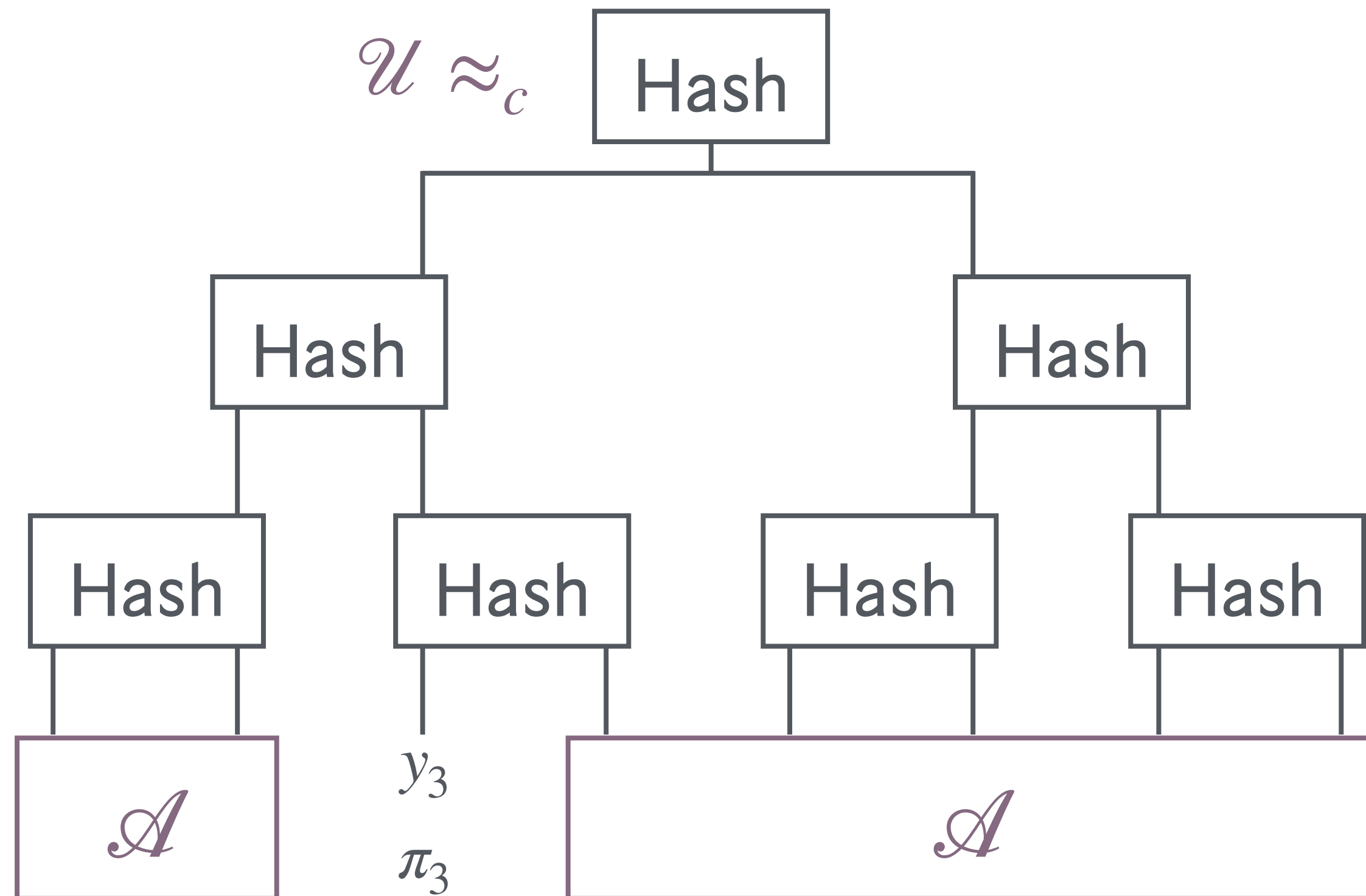
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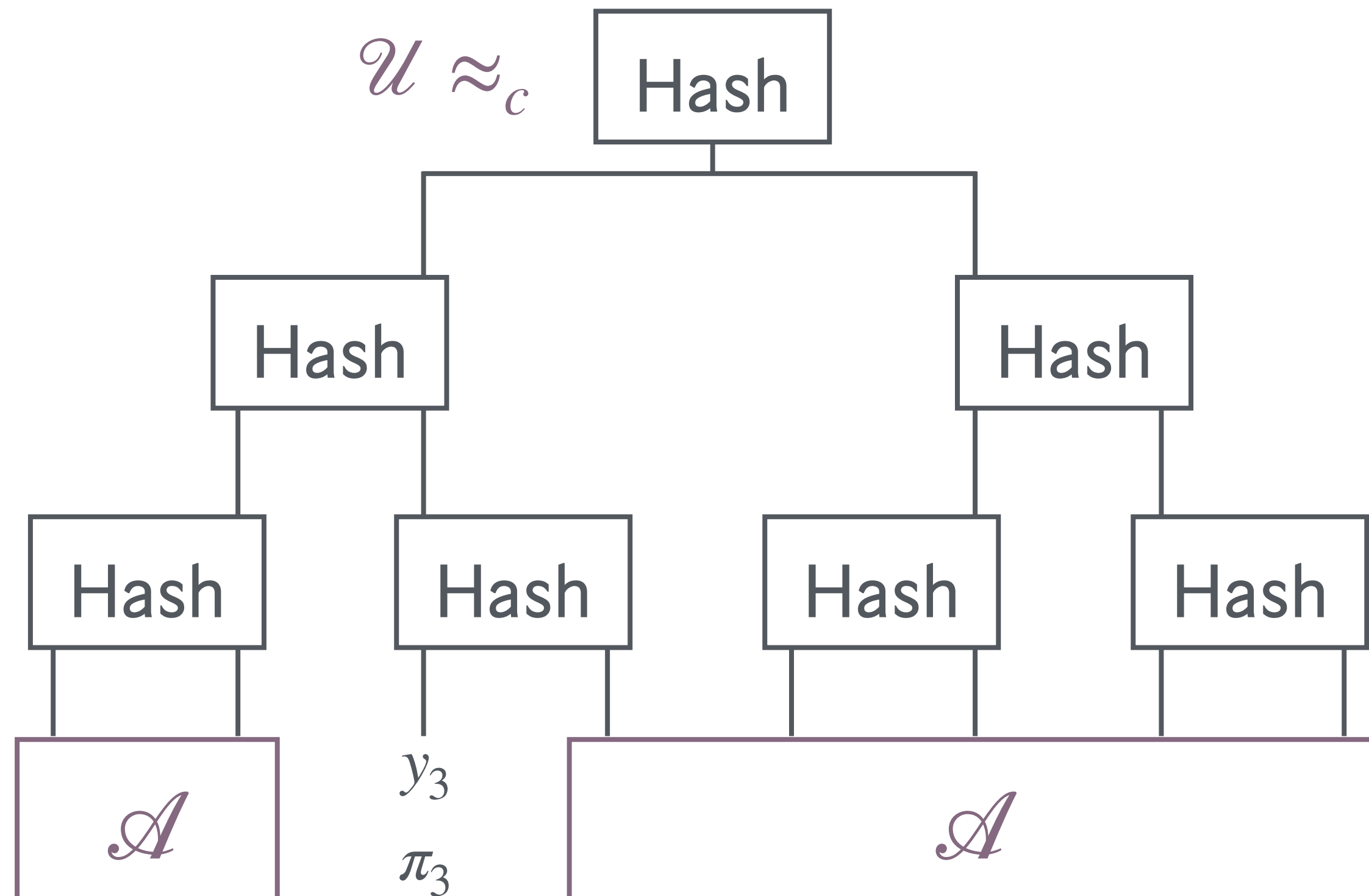
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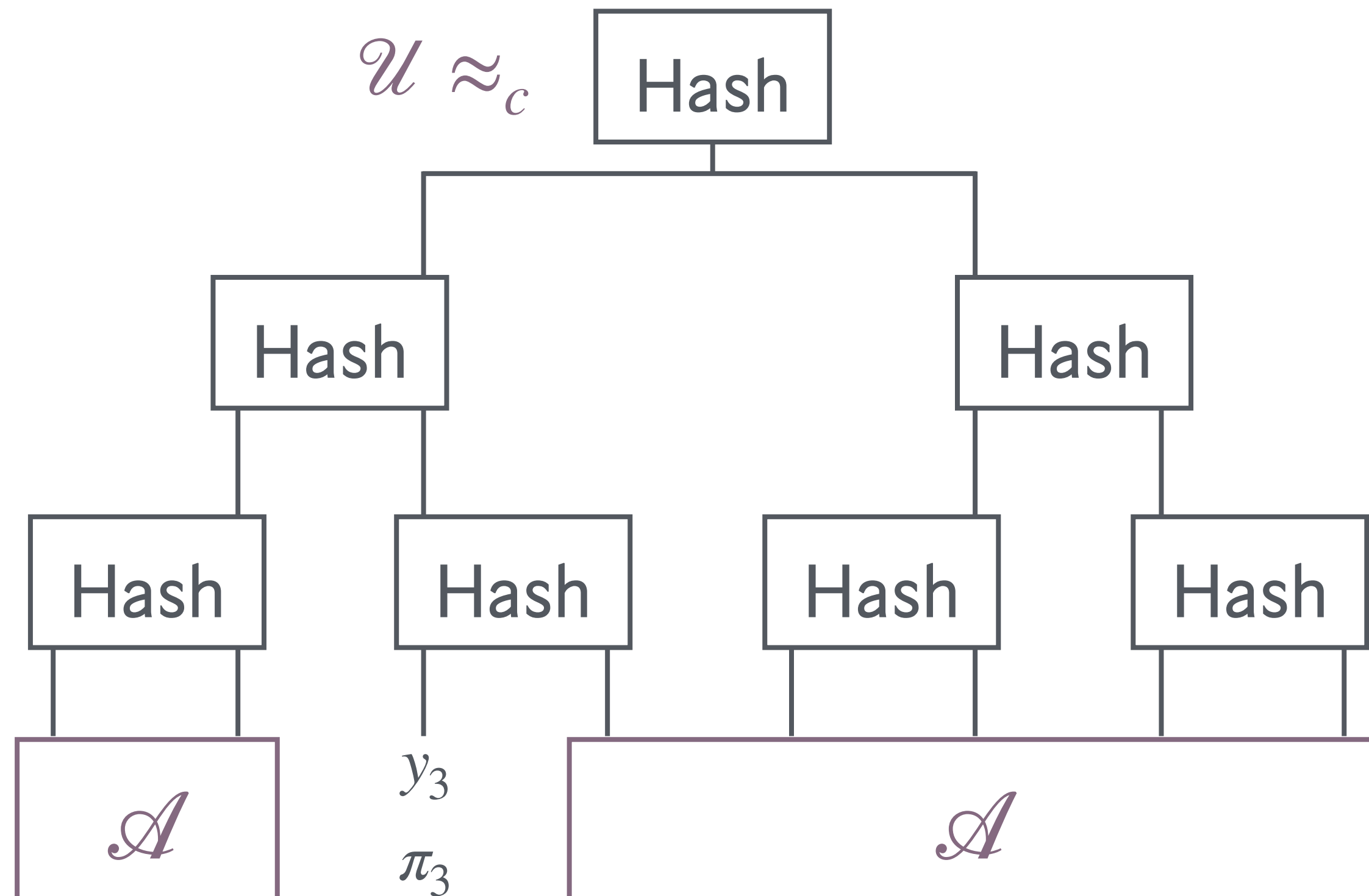
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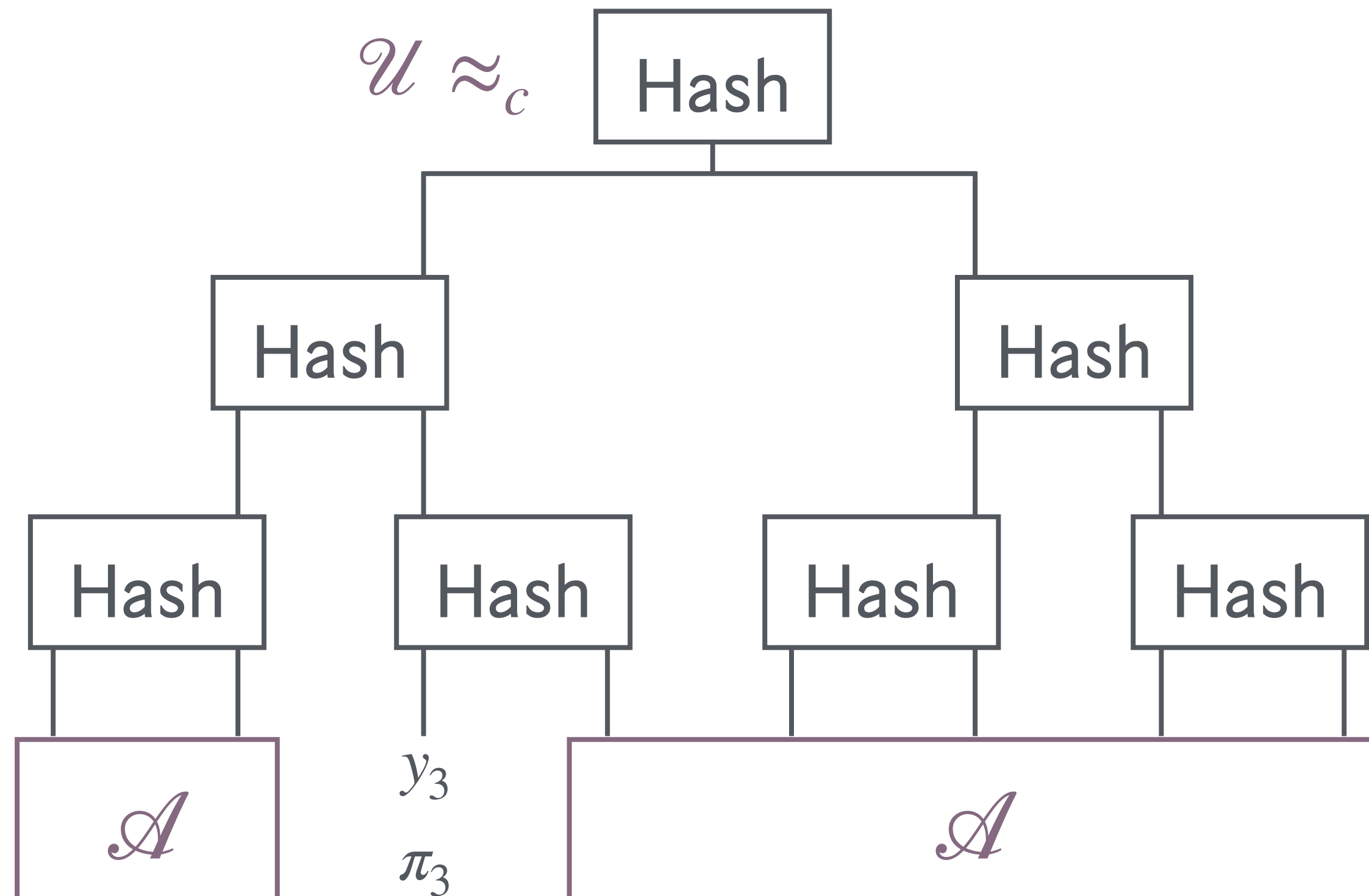
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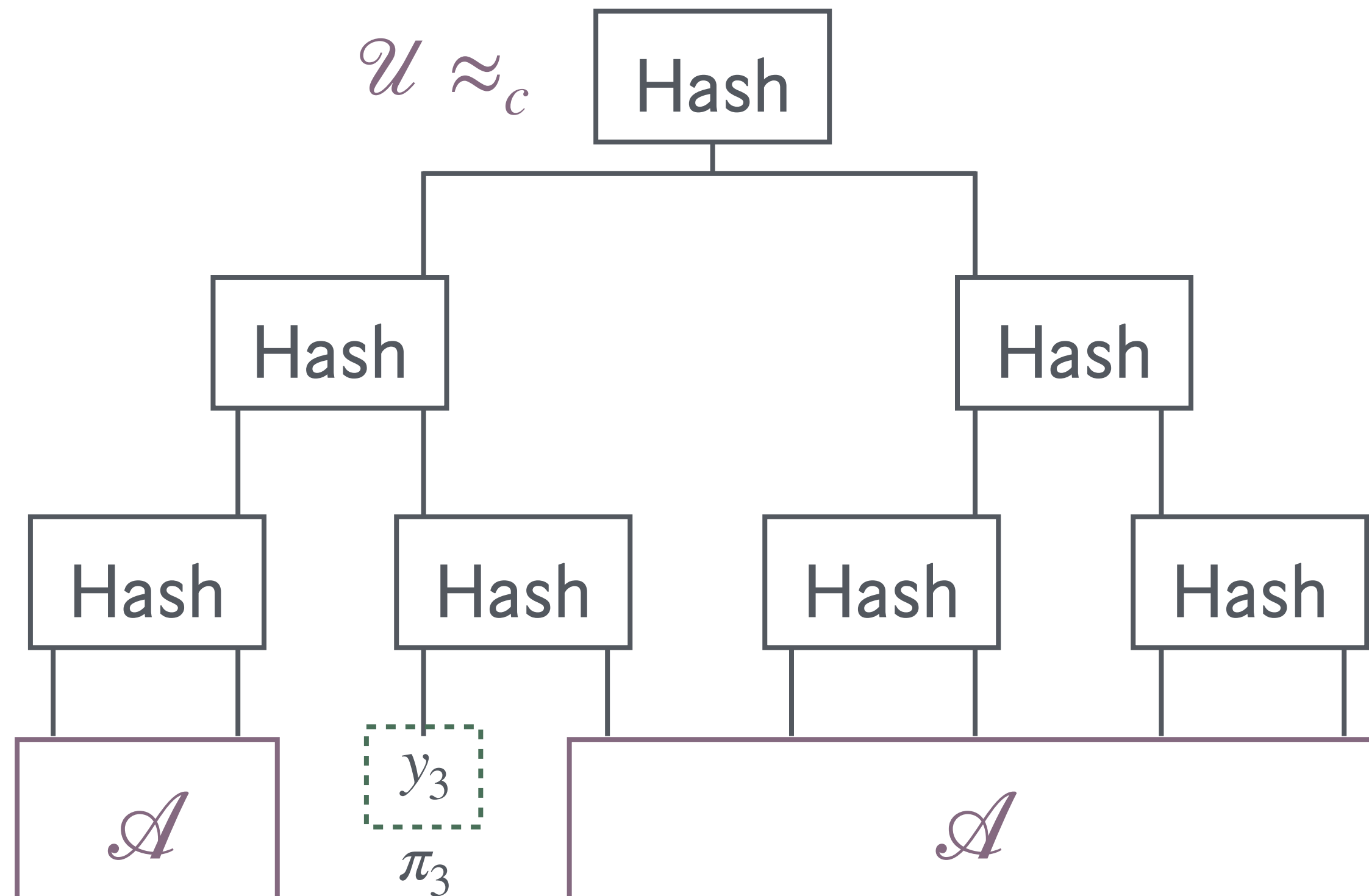
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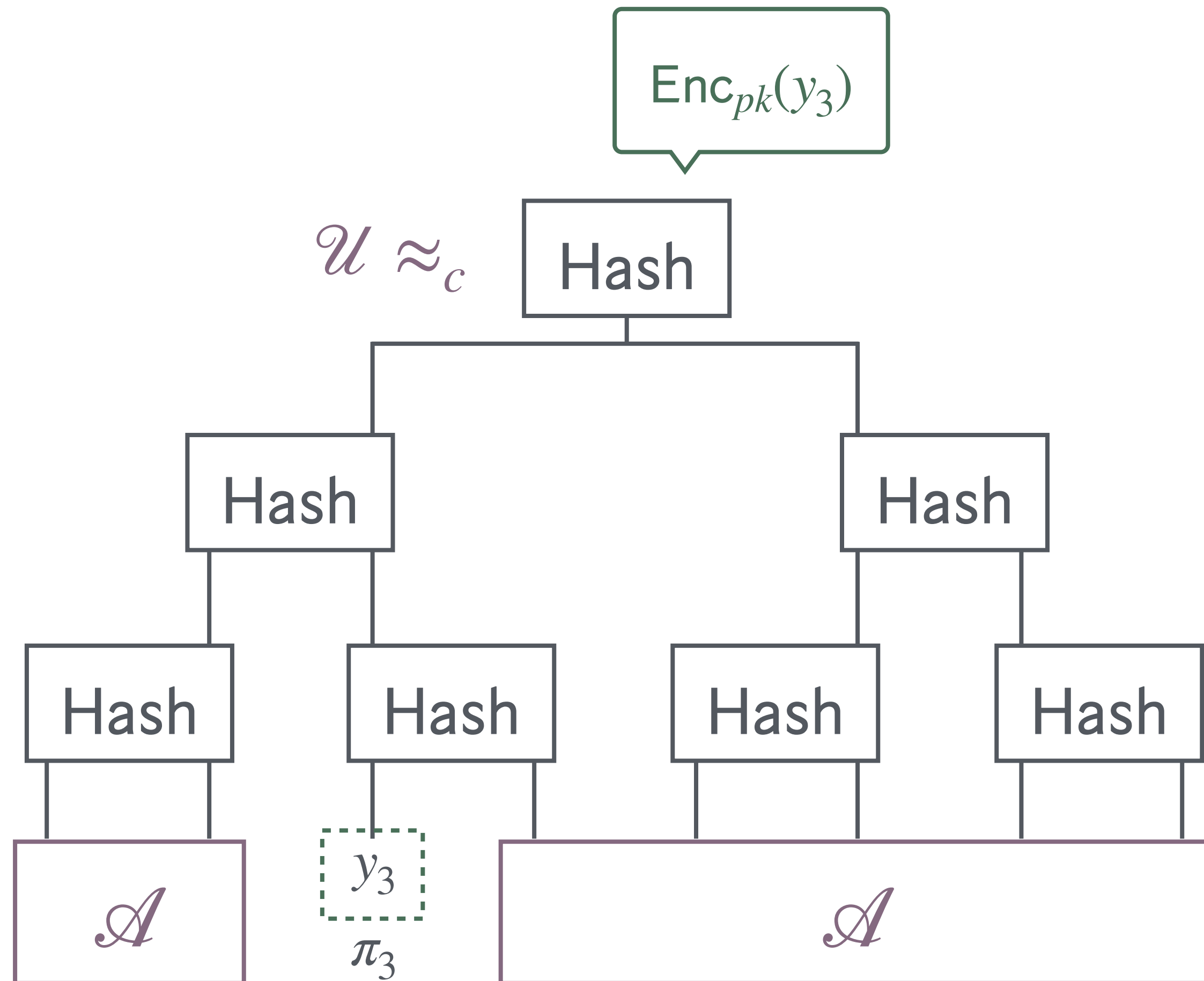
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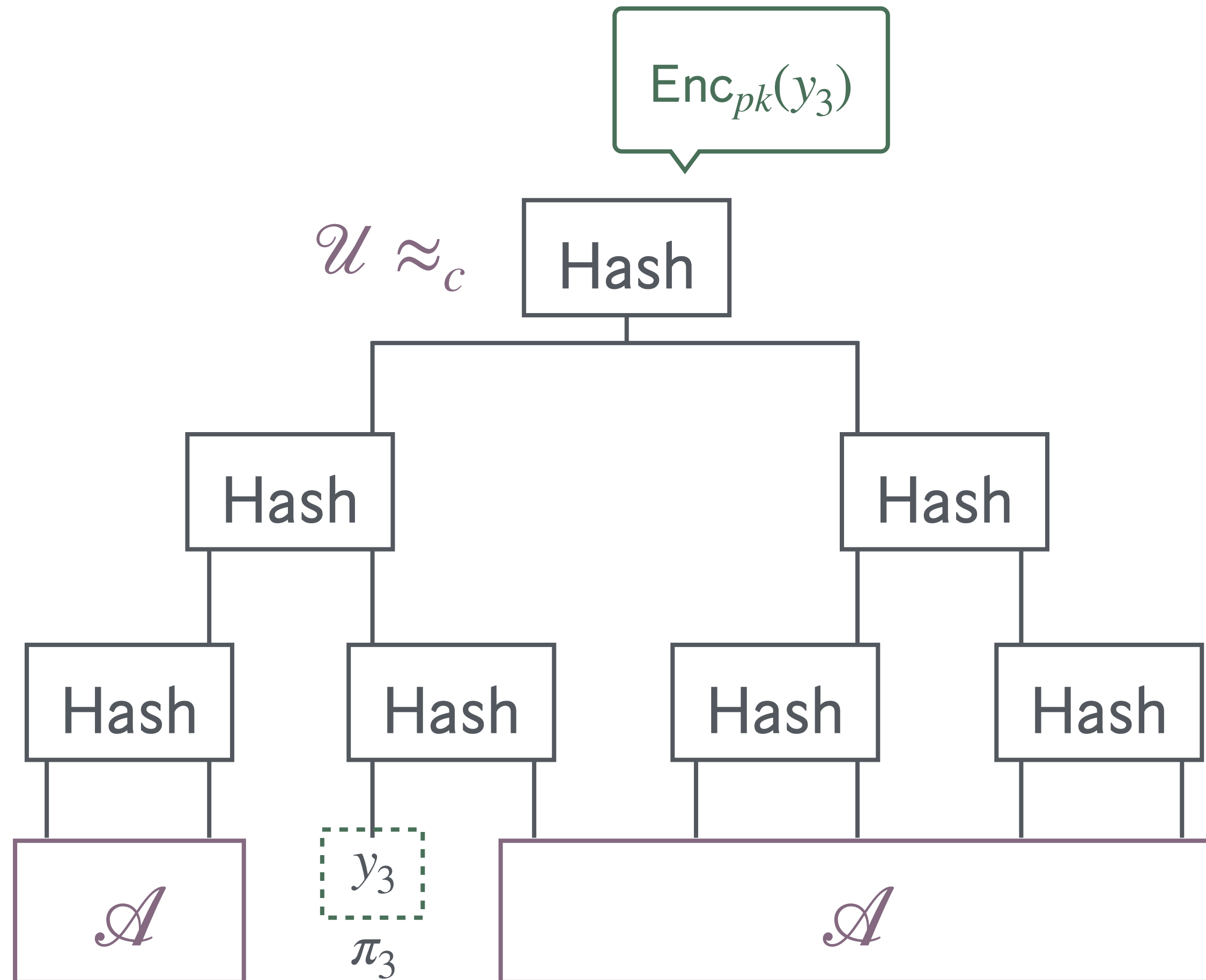


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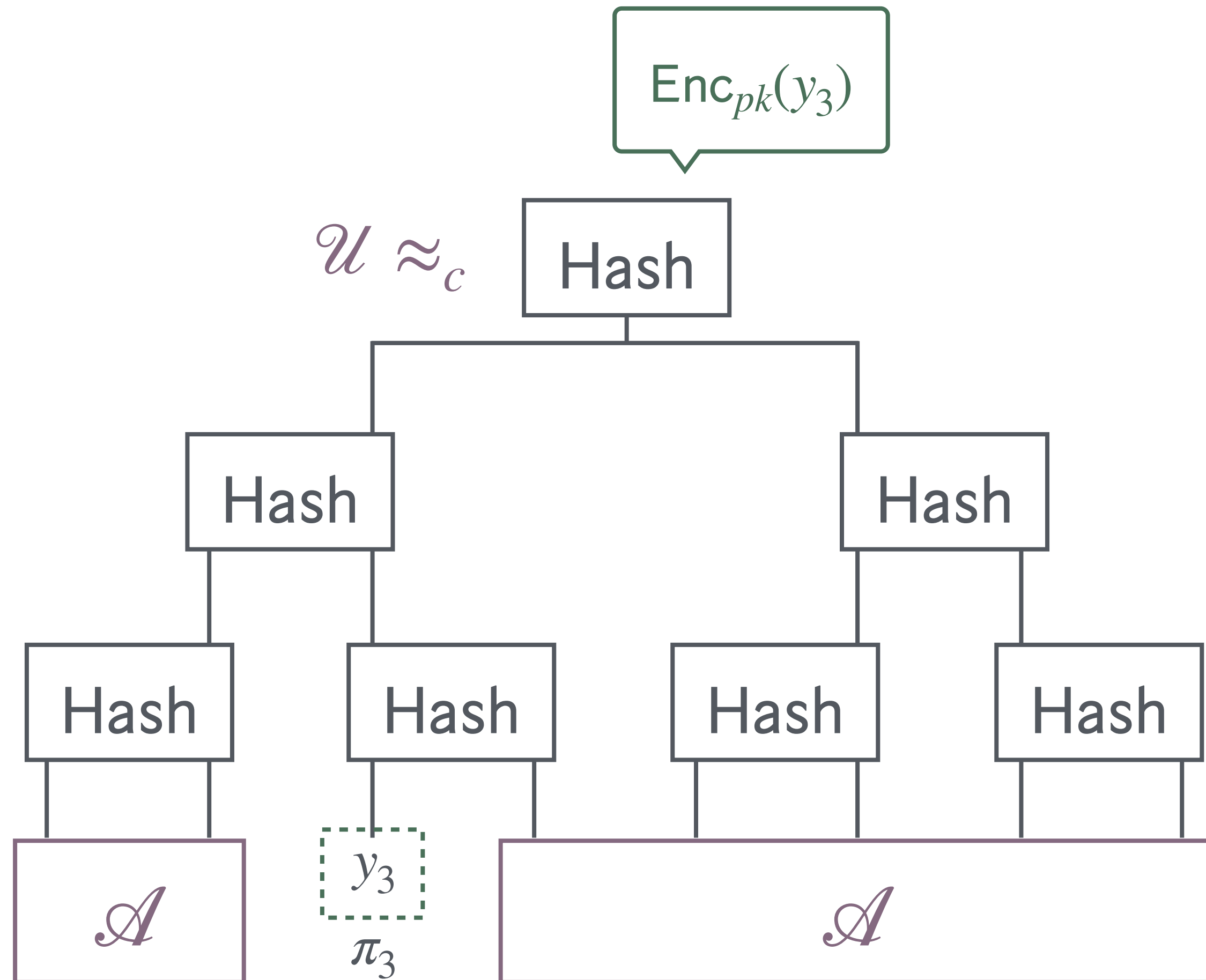


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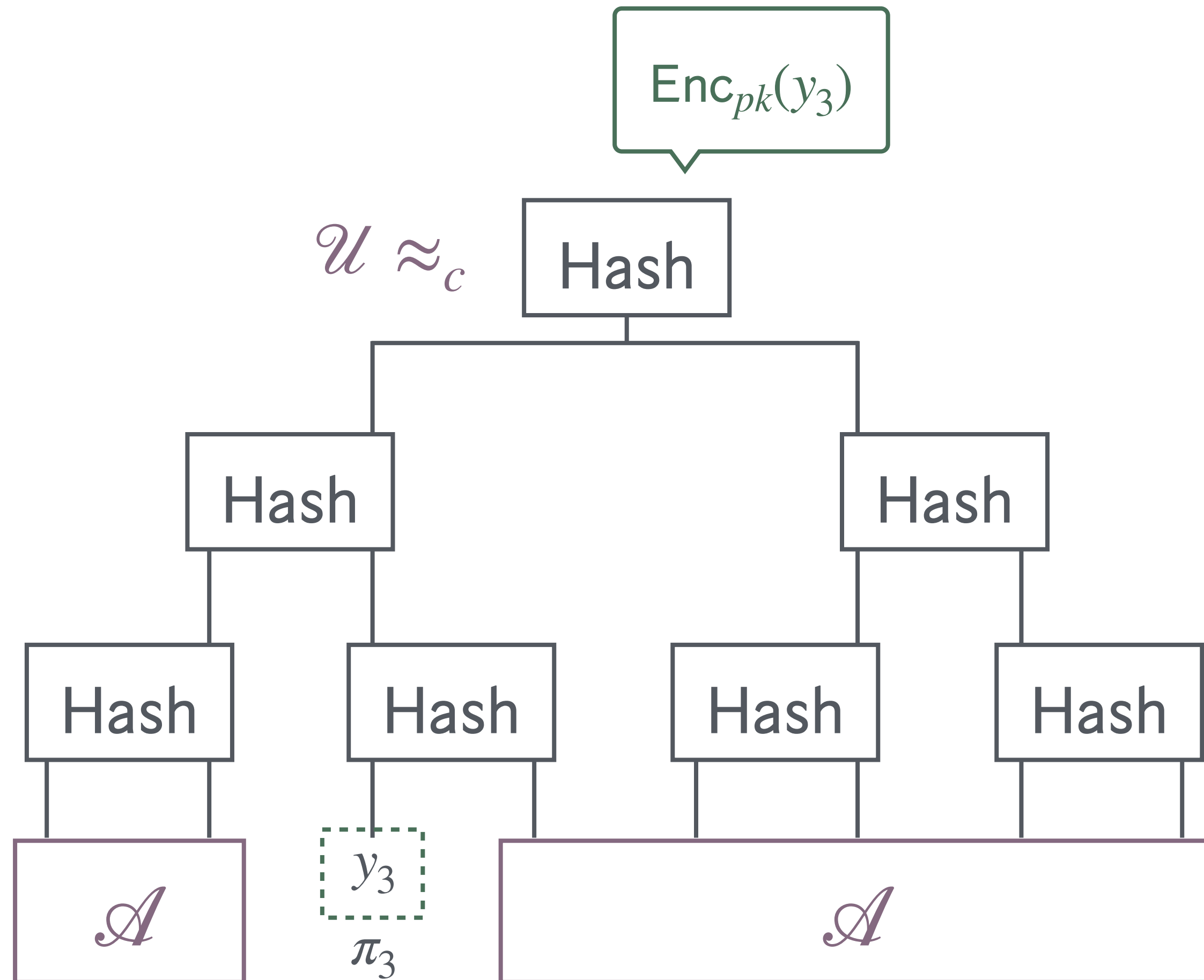
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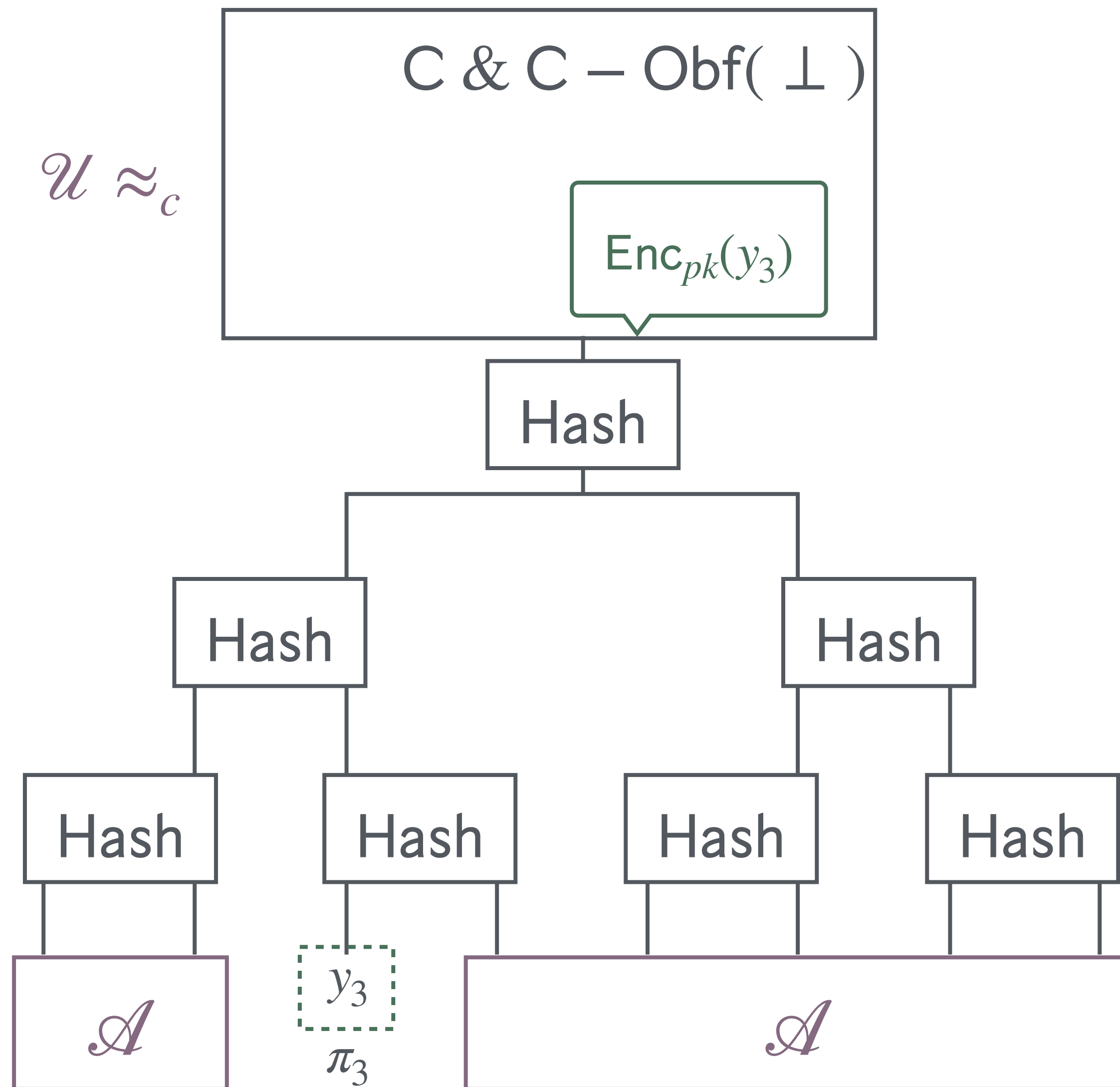
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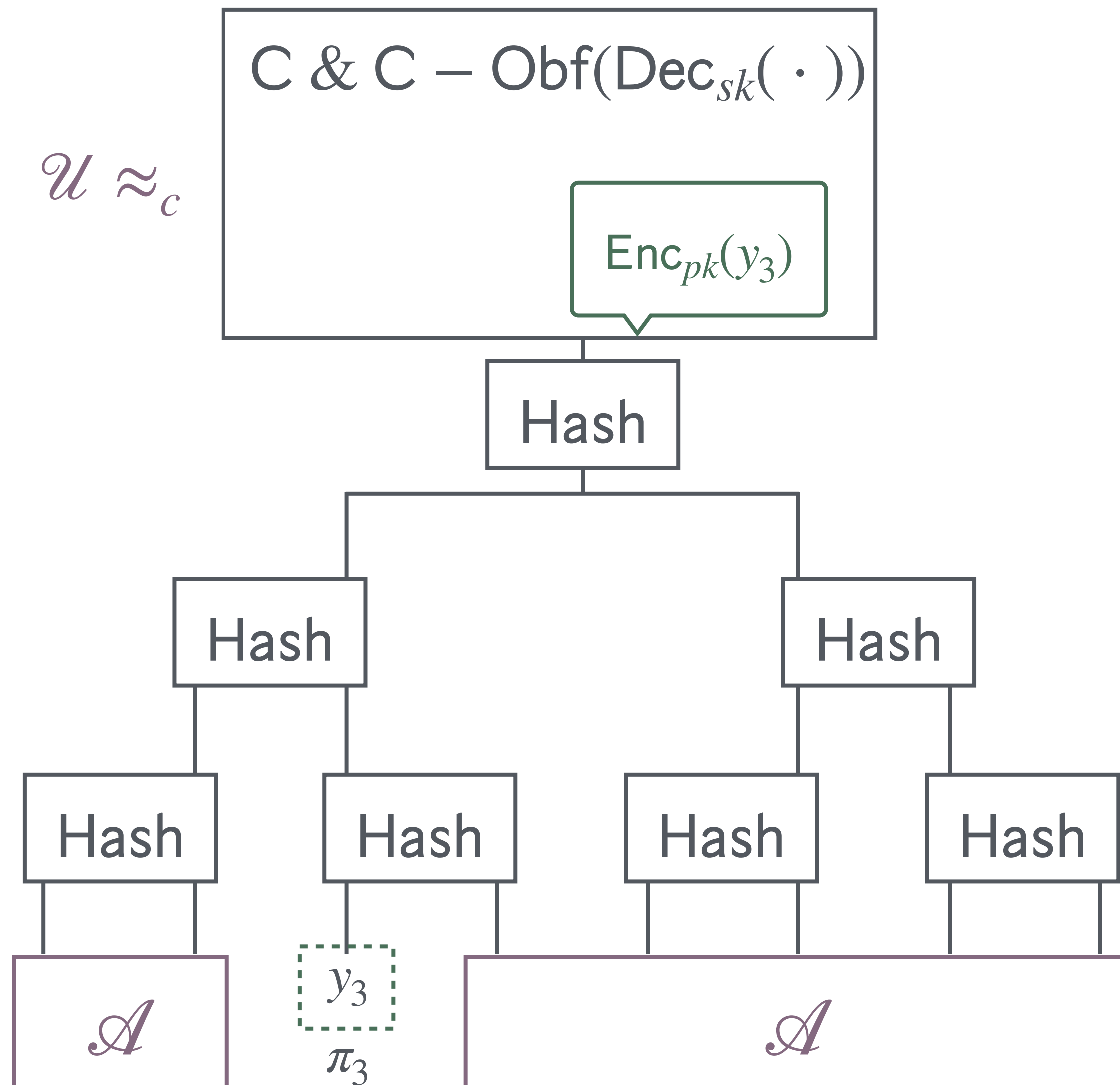
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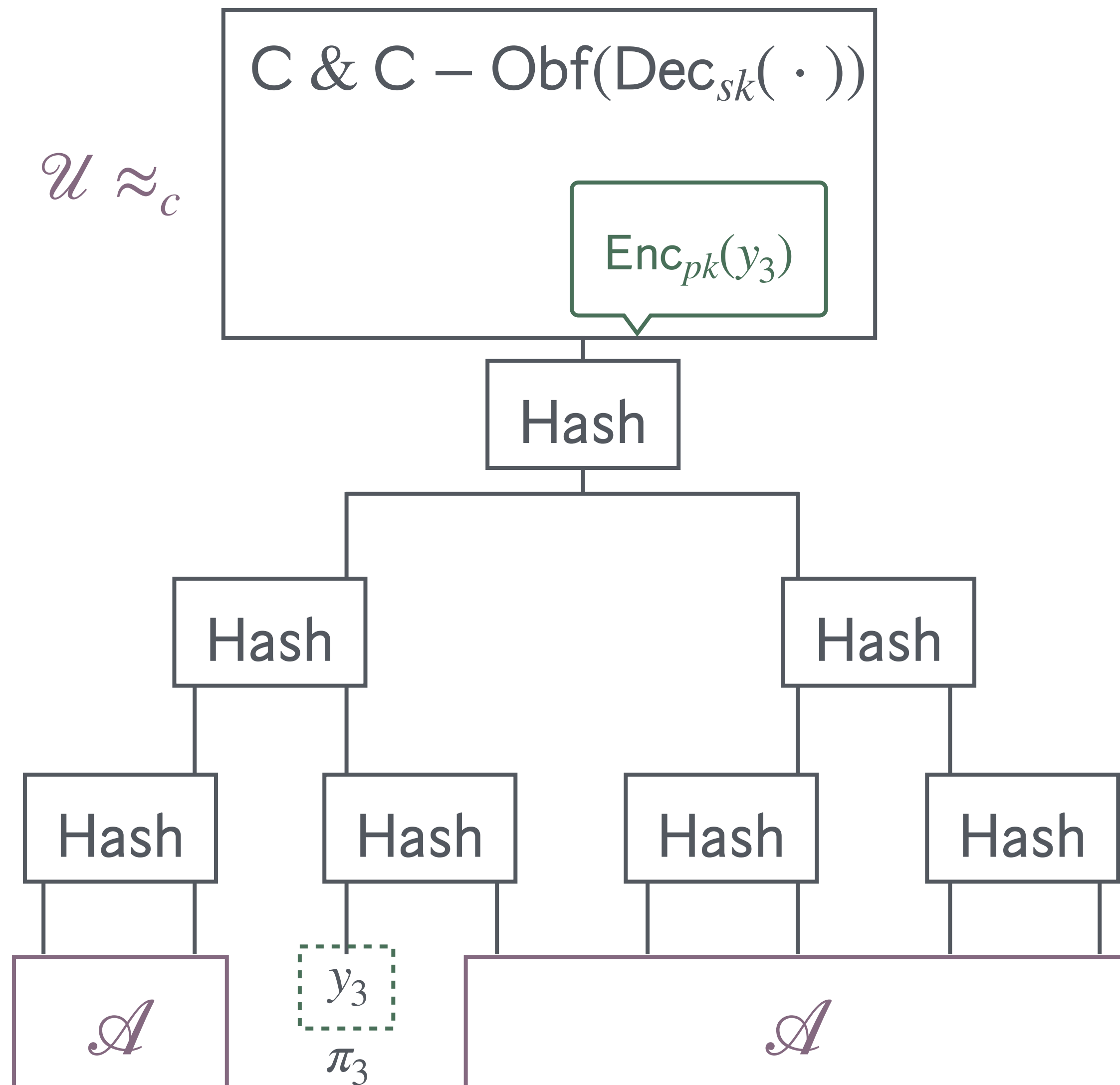
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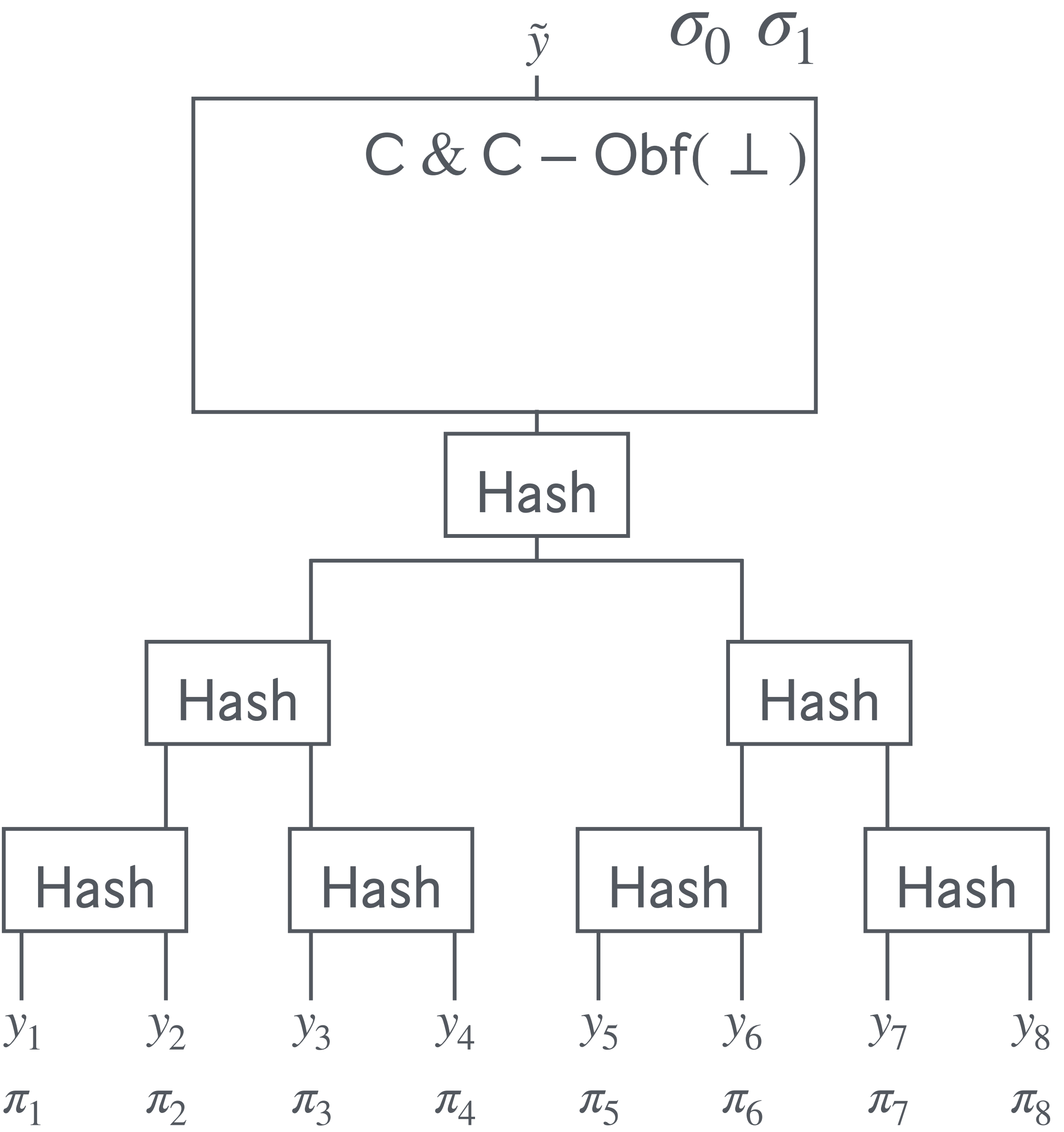
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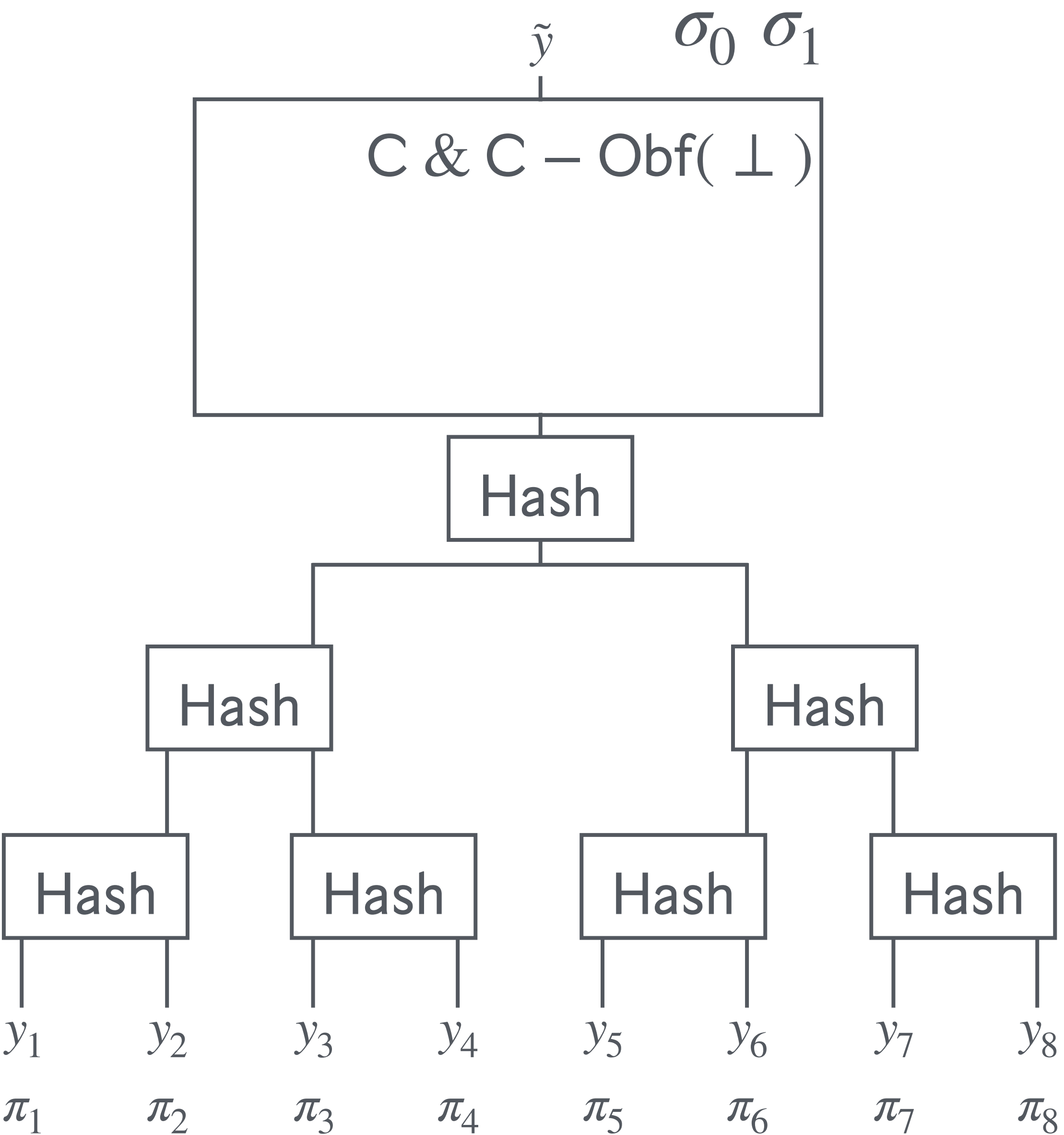
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 - Use the randomness extraction property of **random matrices**



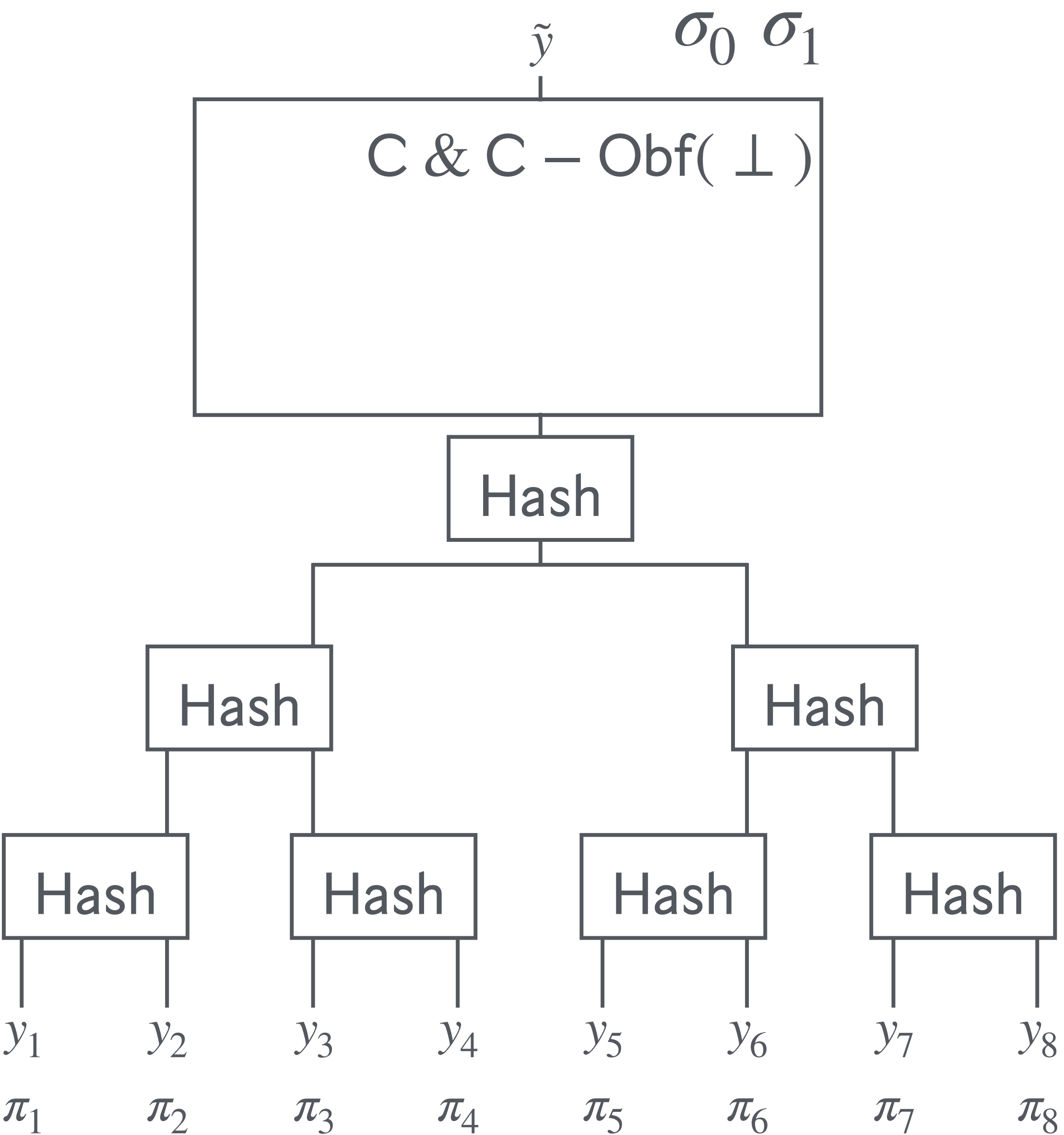


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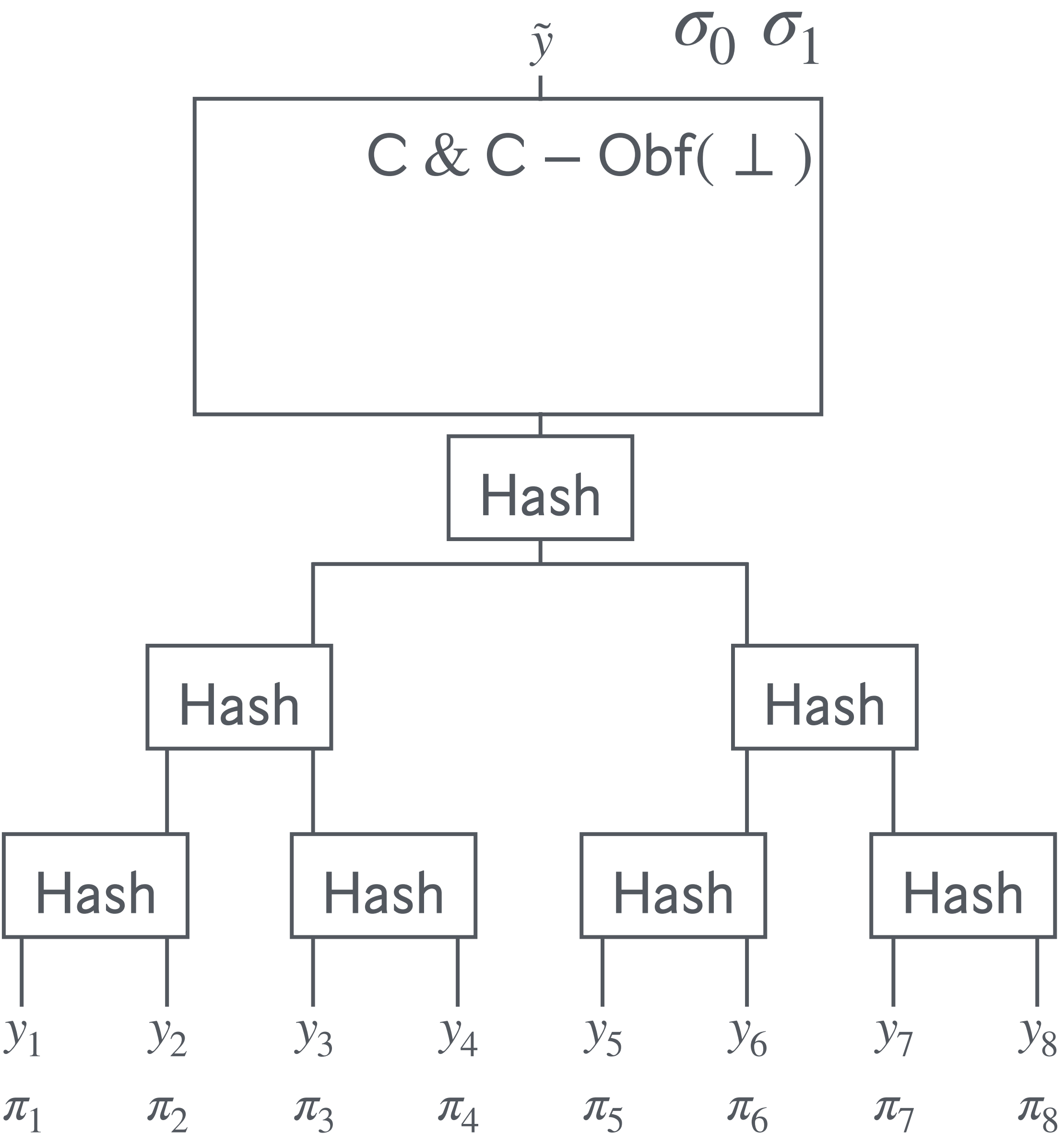
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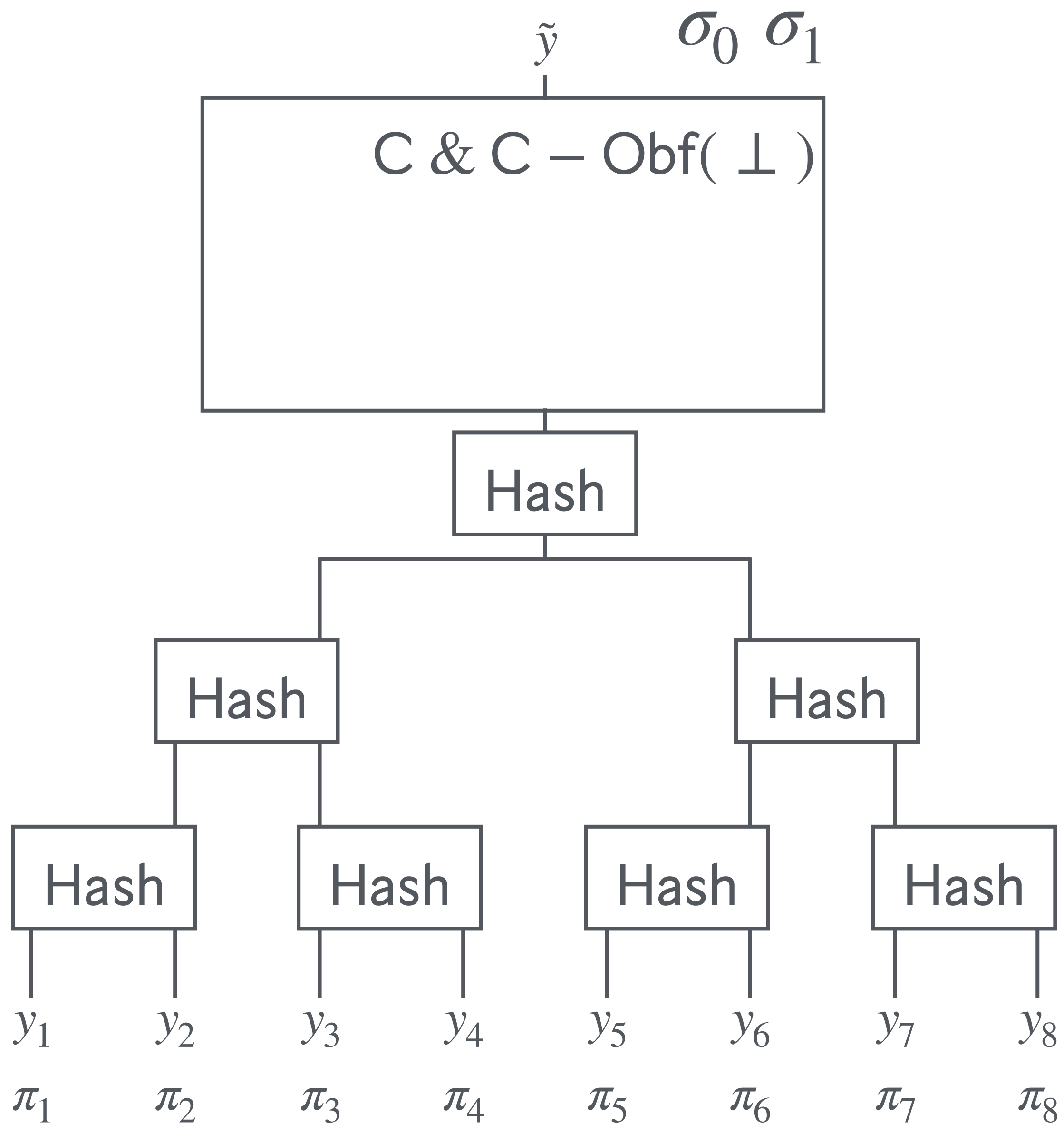
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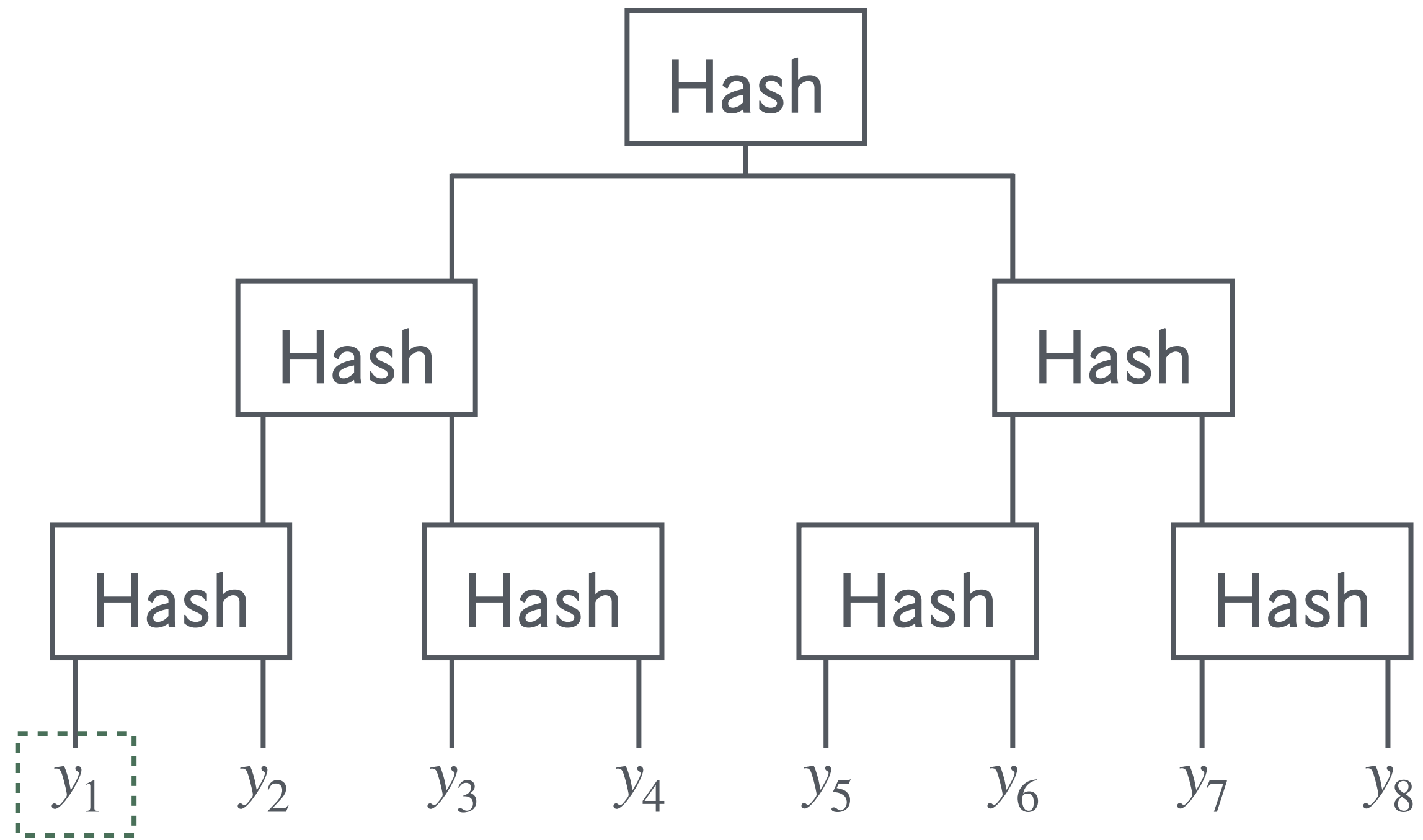
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- We want to derive a **contradiction** against the uniqueness of the (non-aggregate) VRF

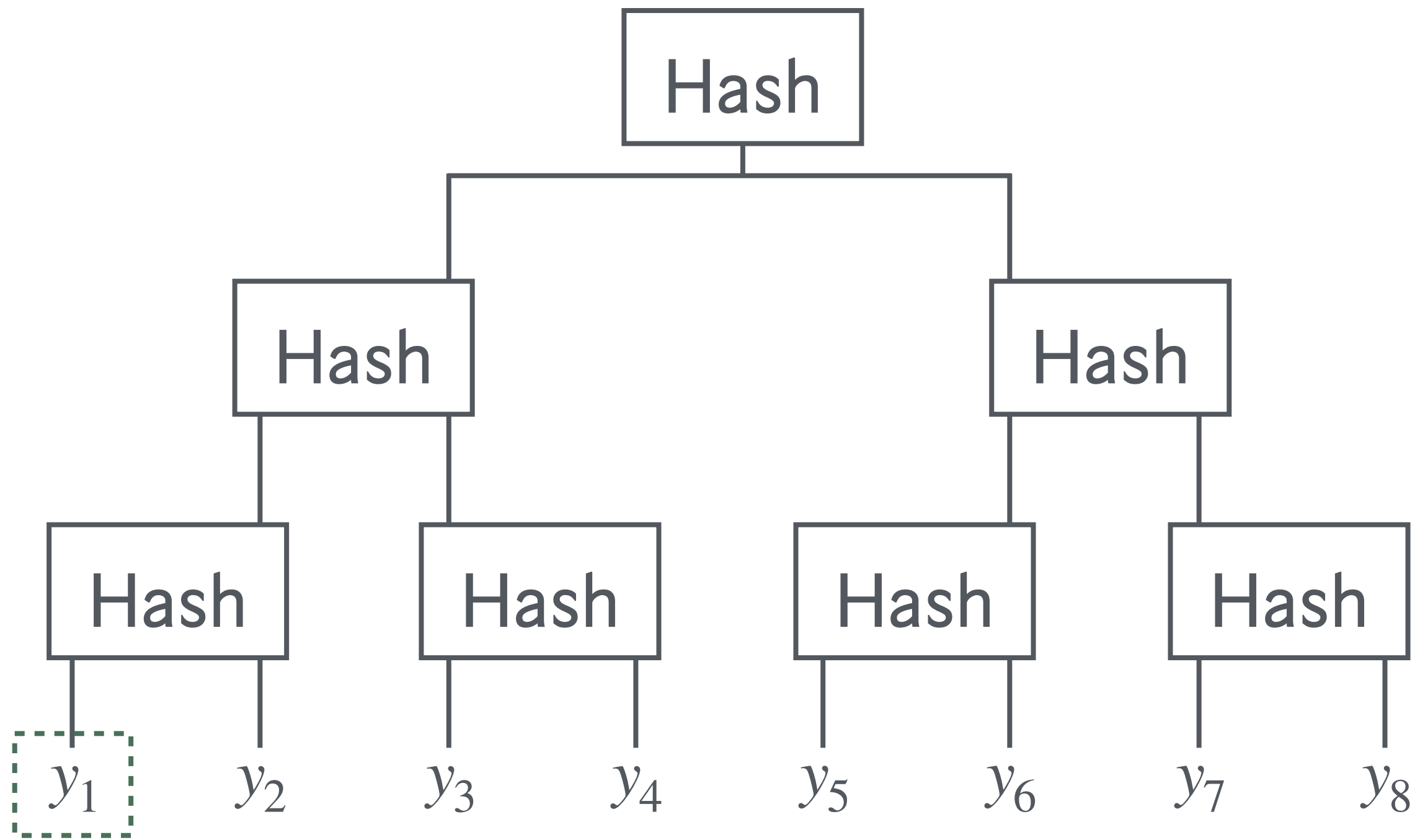
HONEST TREE



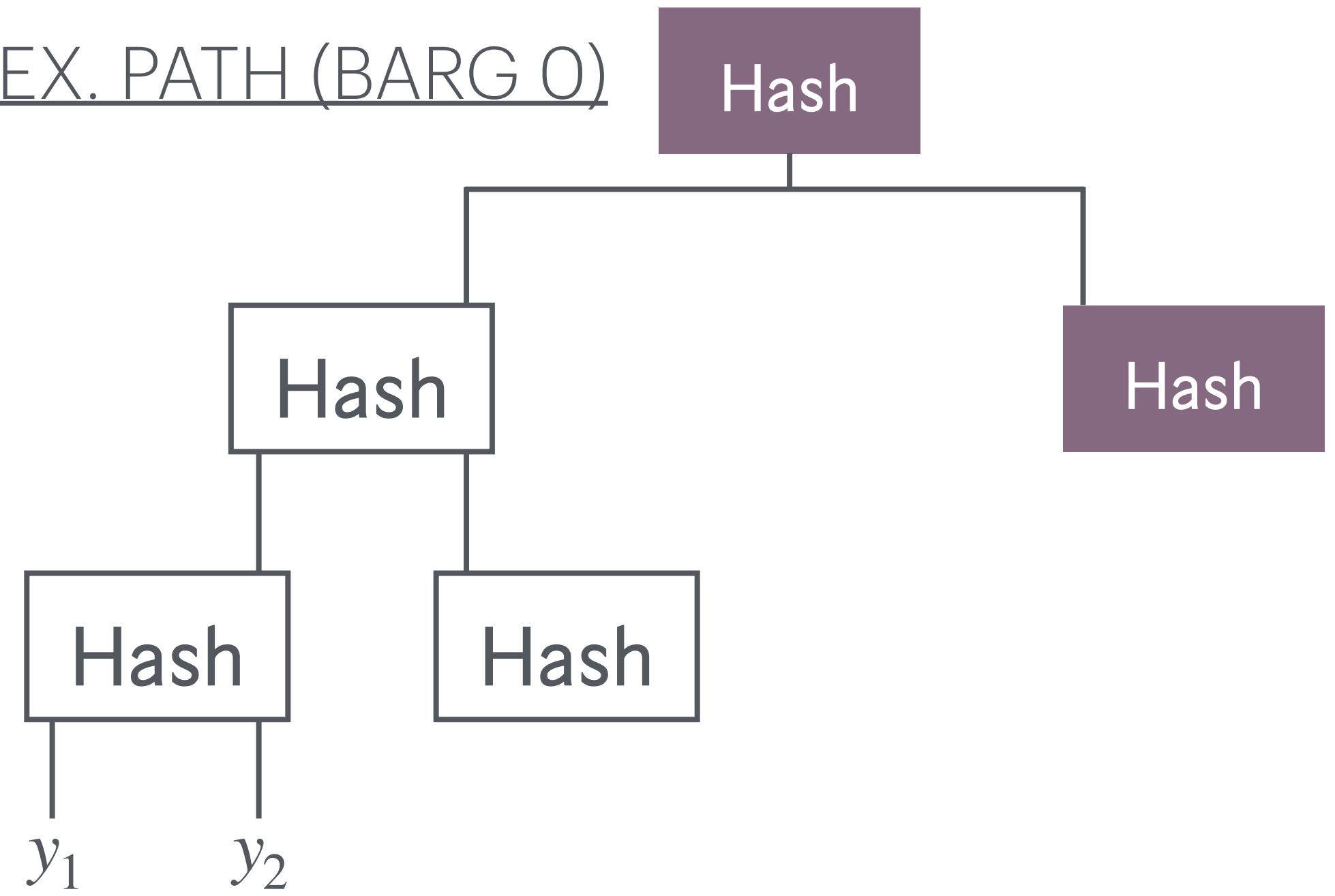
EX. PATH (BARG 0)

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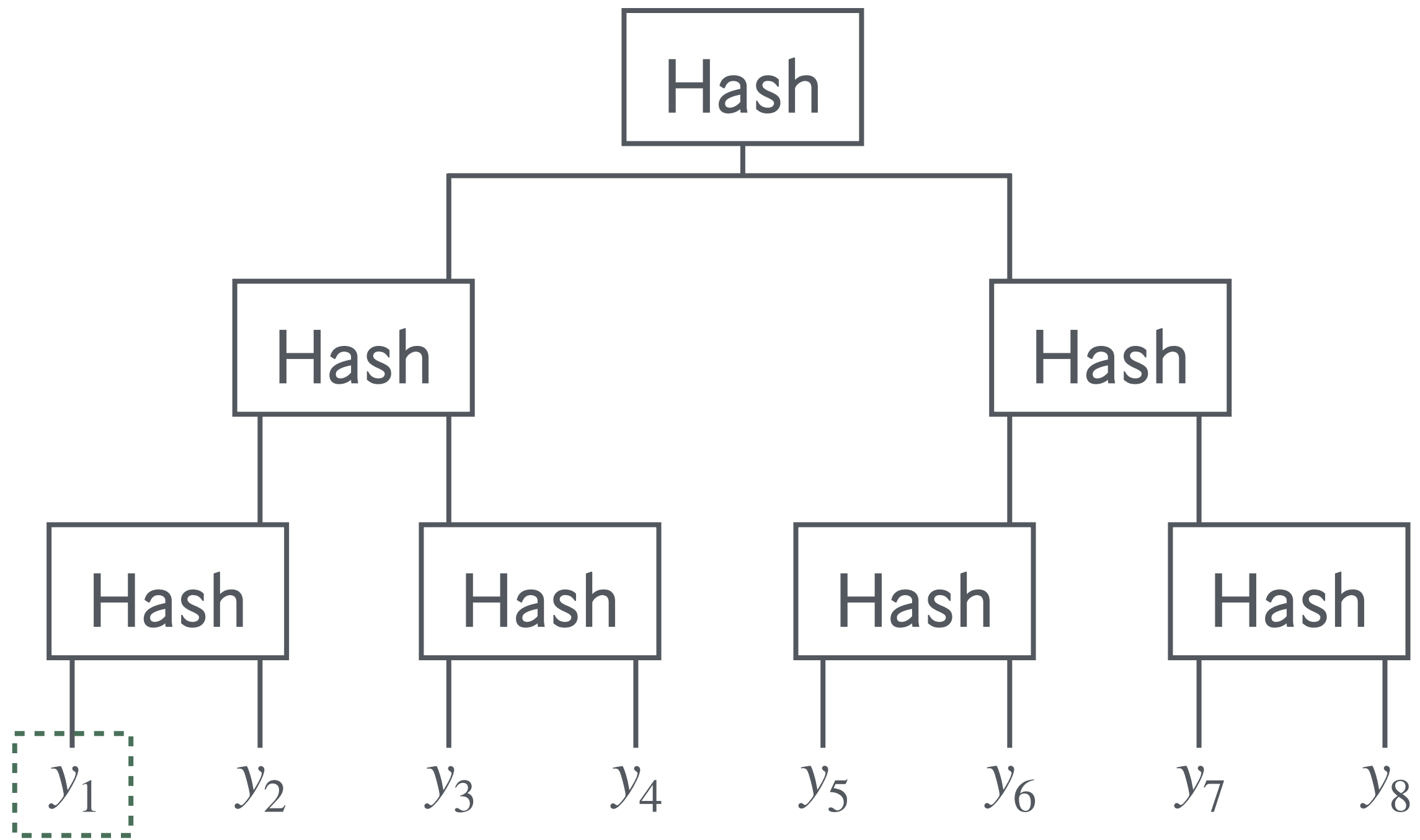


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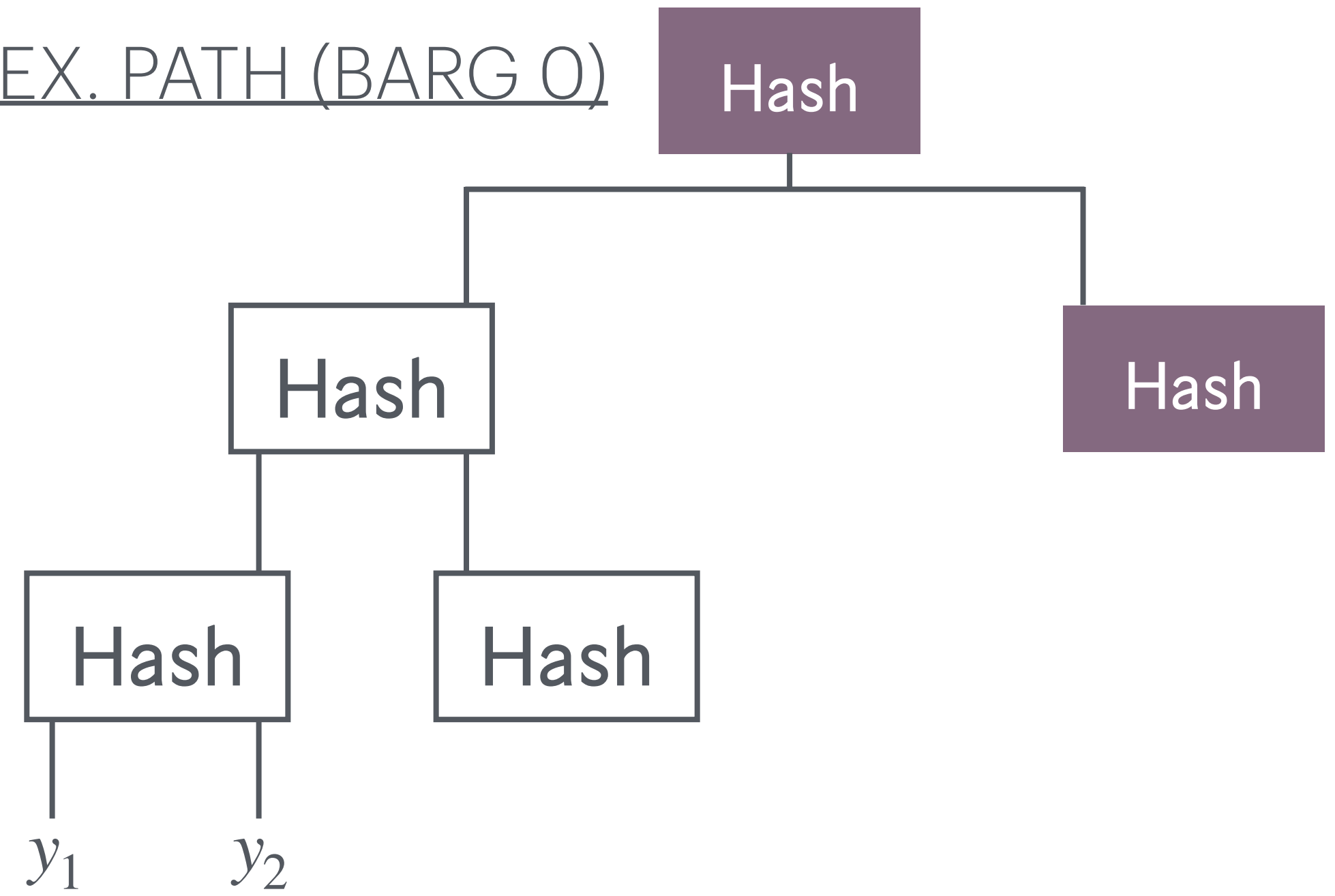


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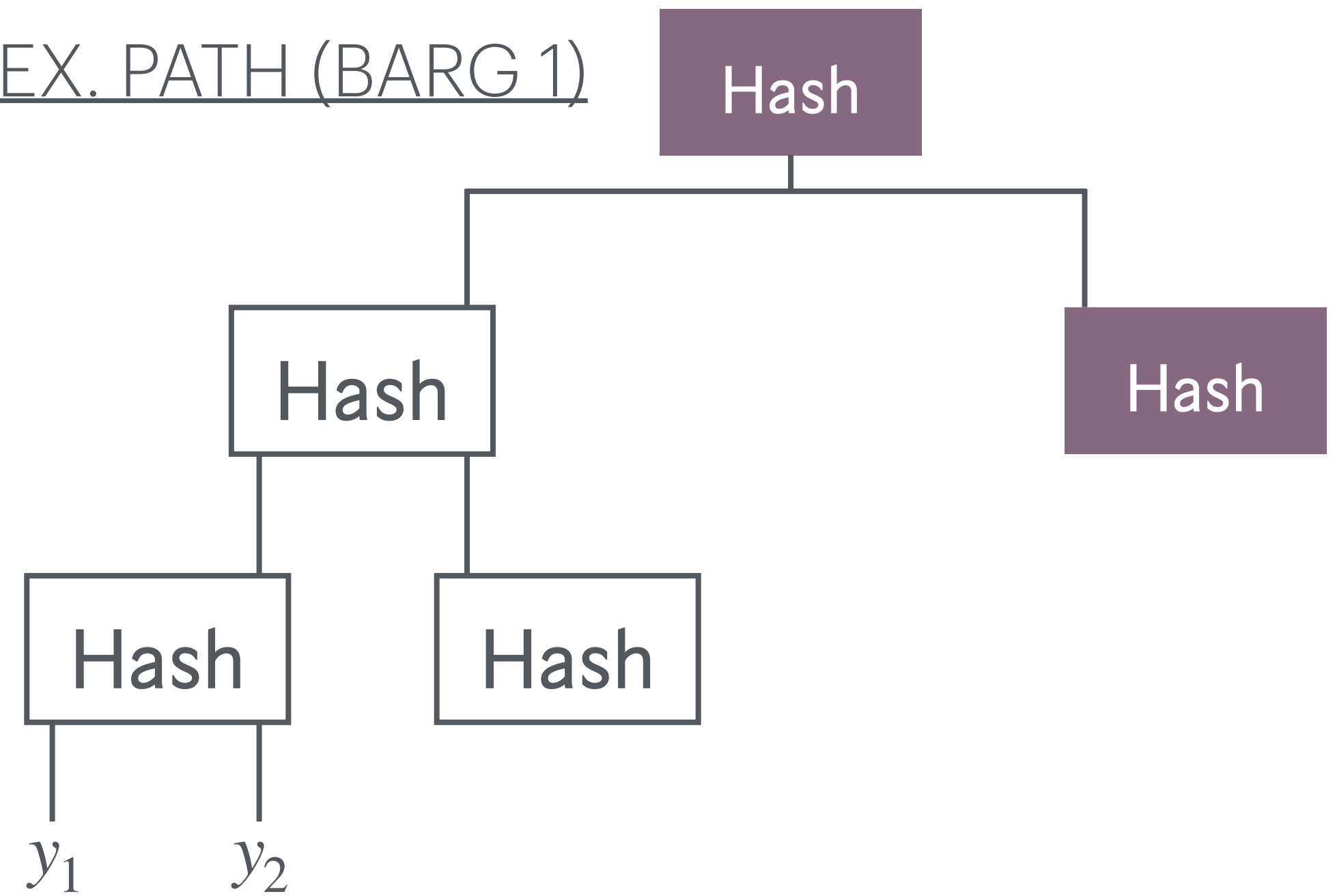
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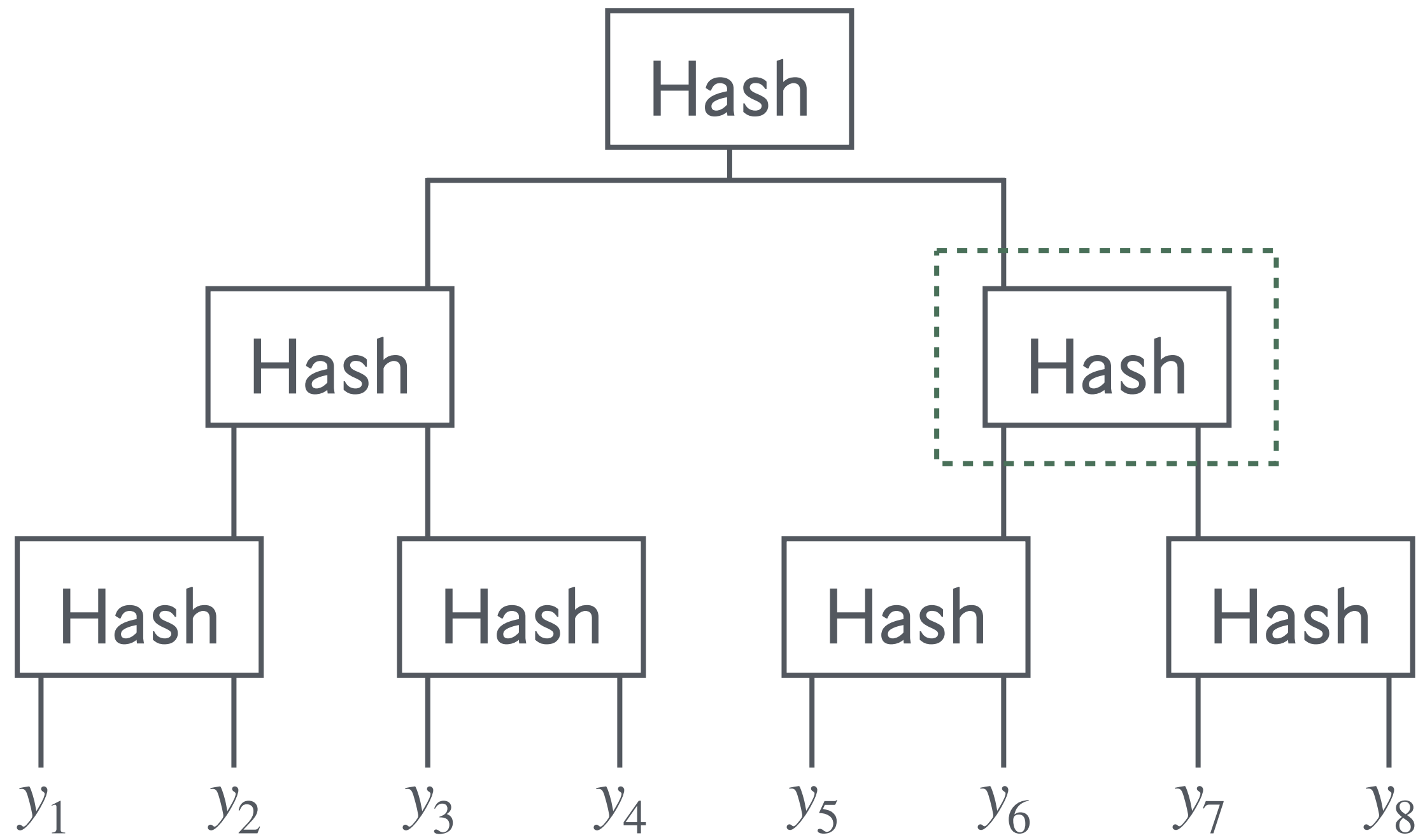
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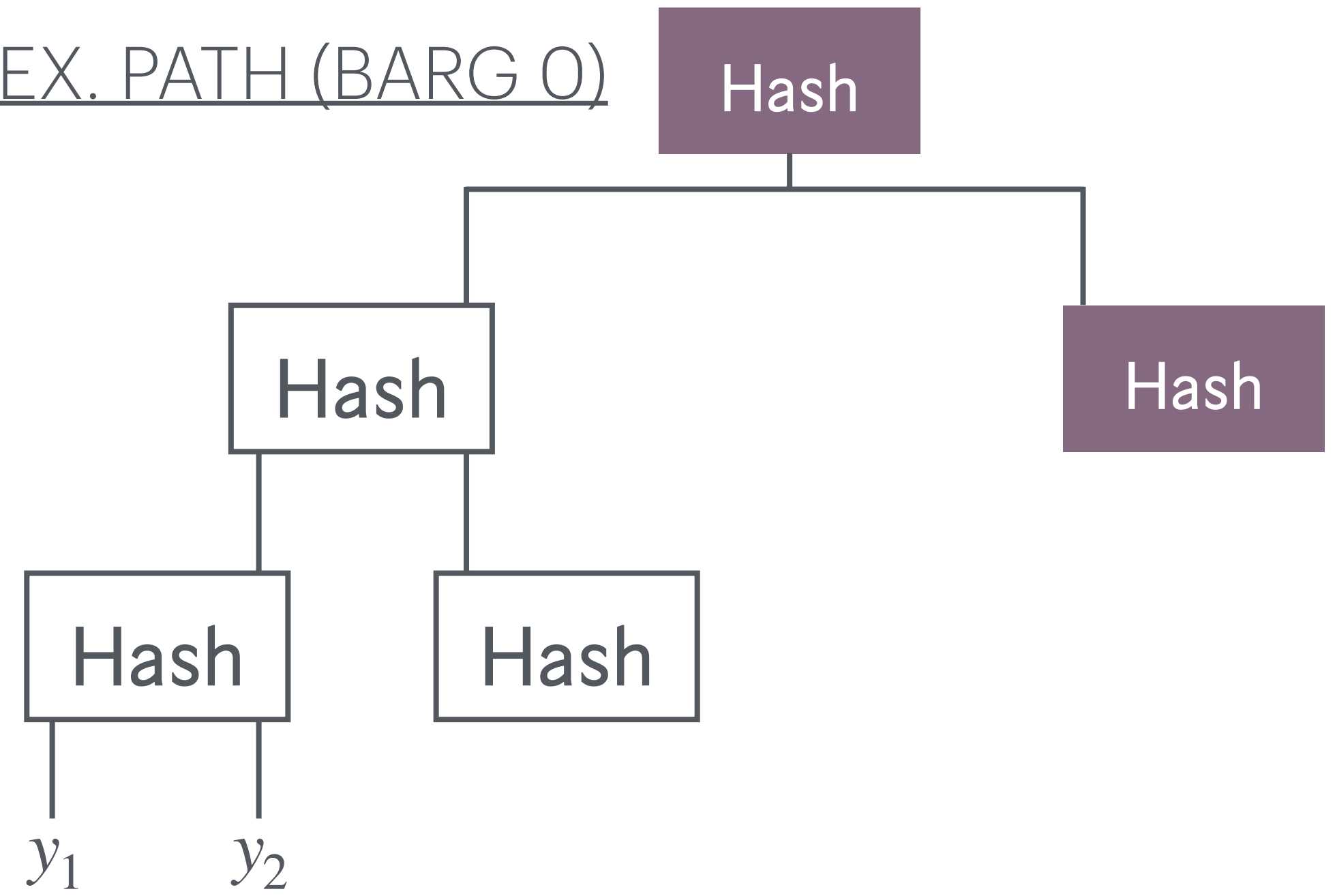
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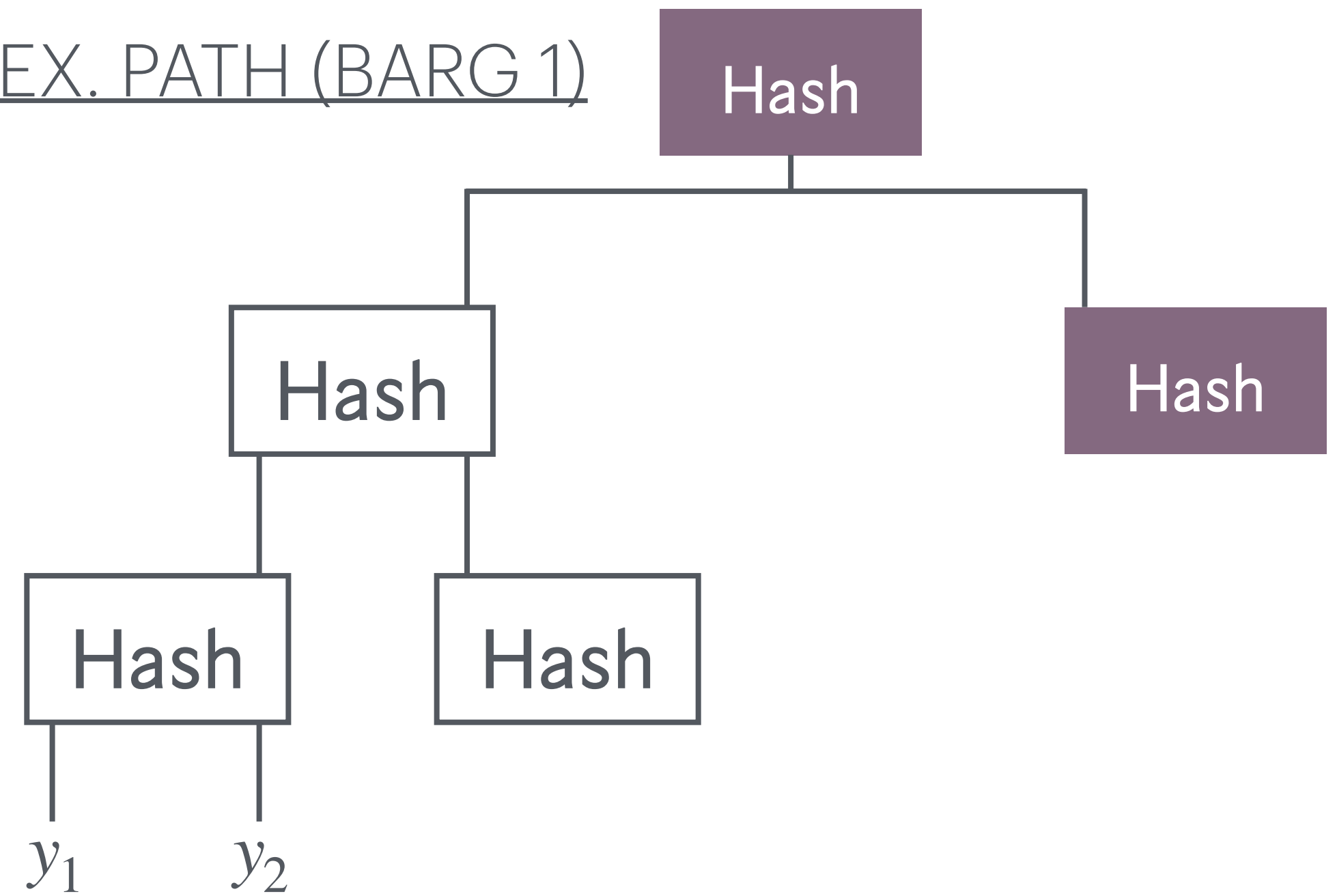
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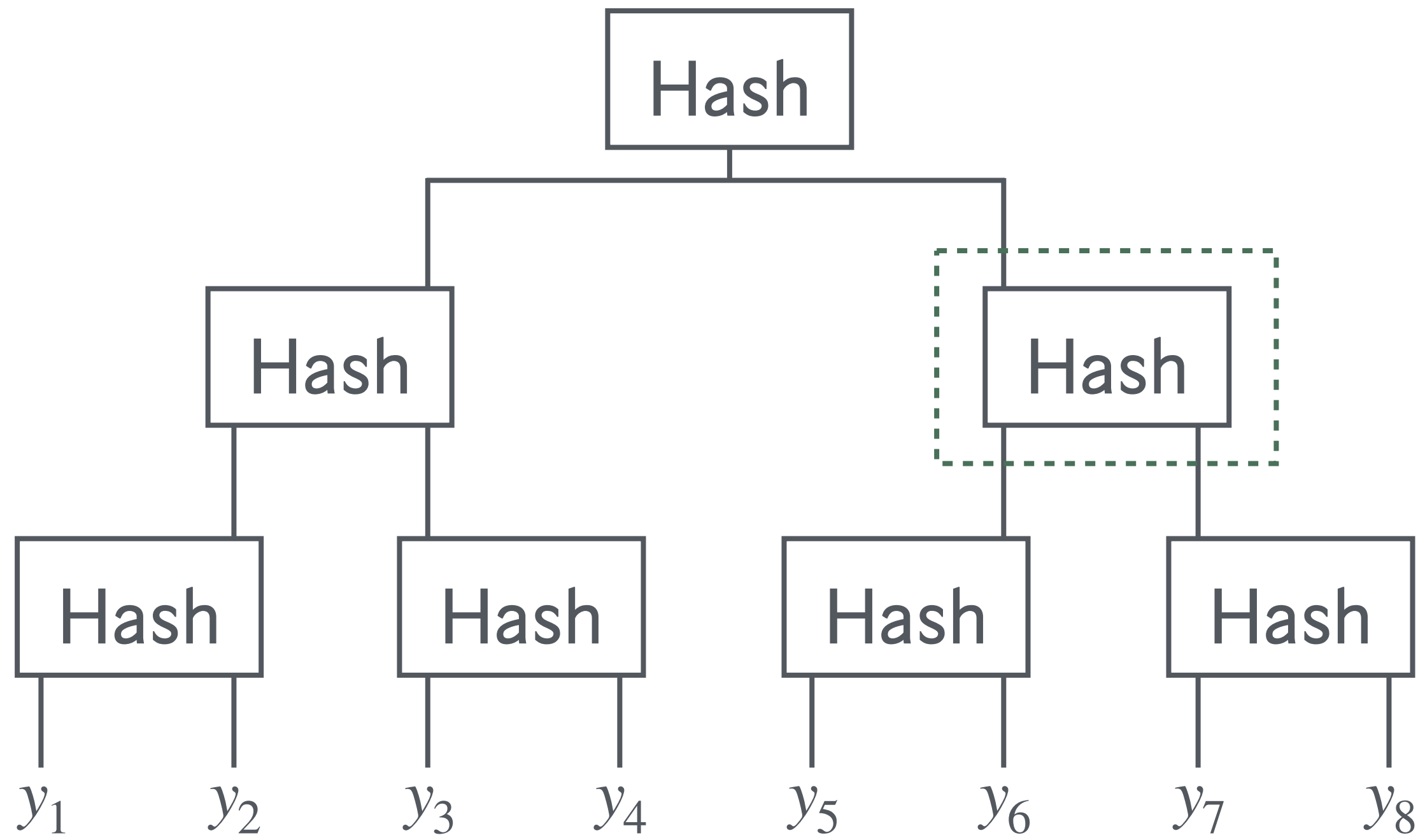
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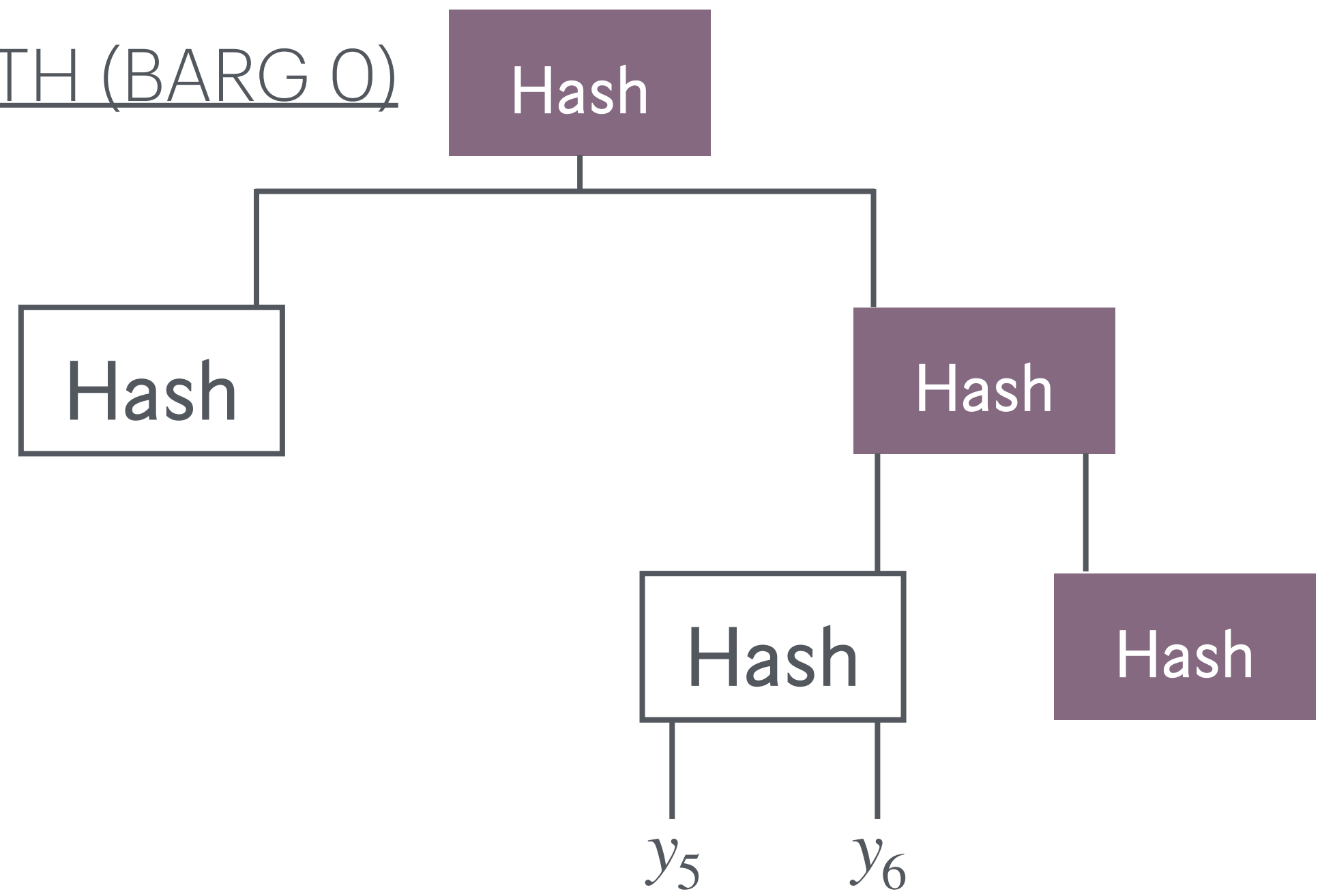
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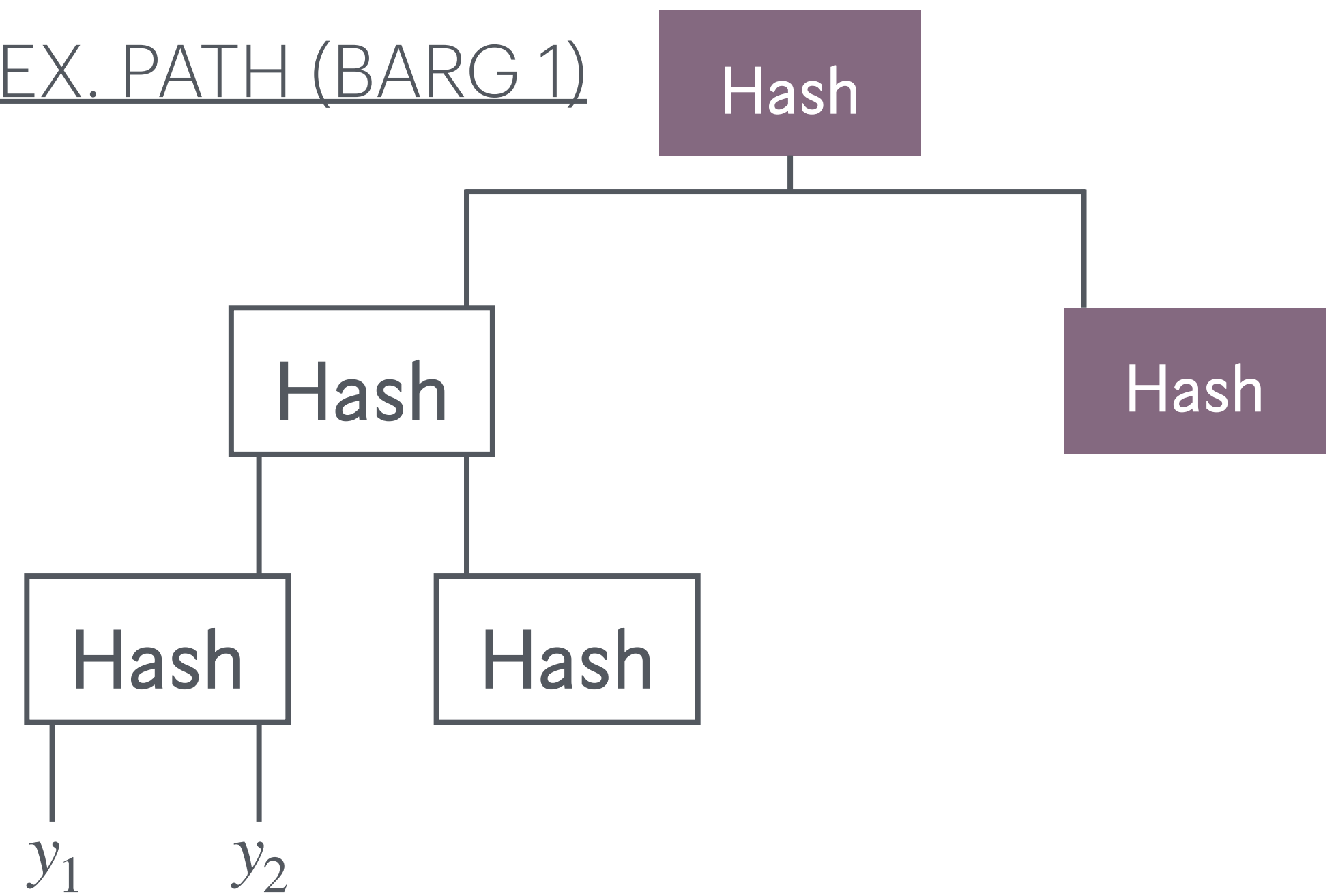
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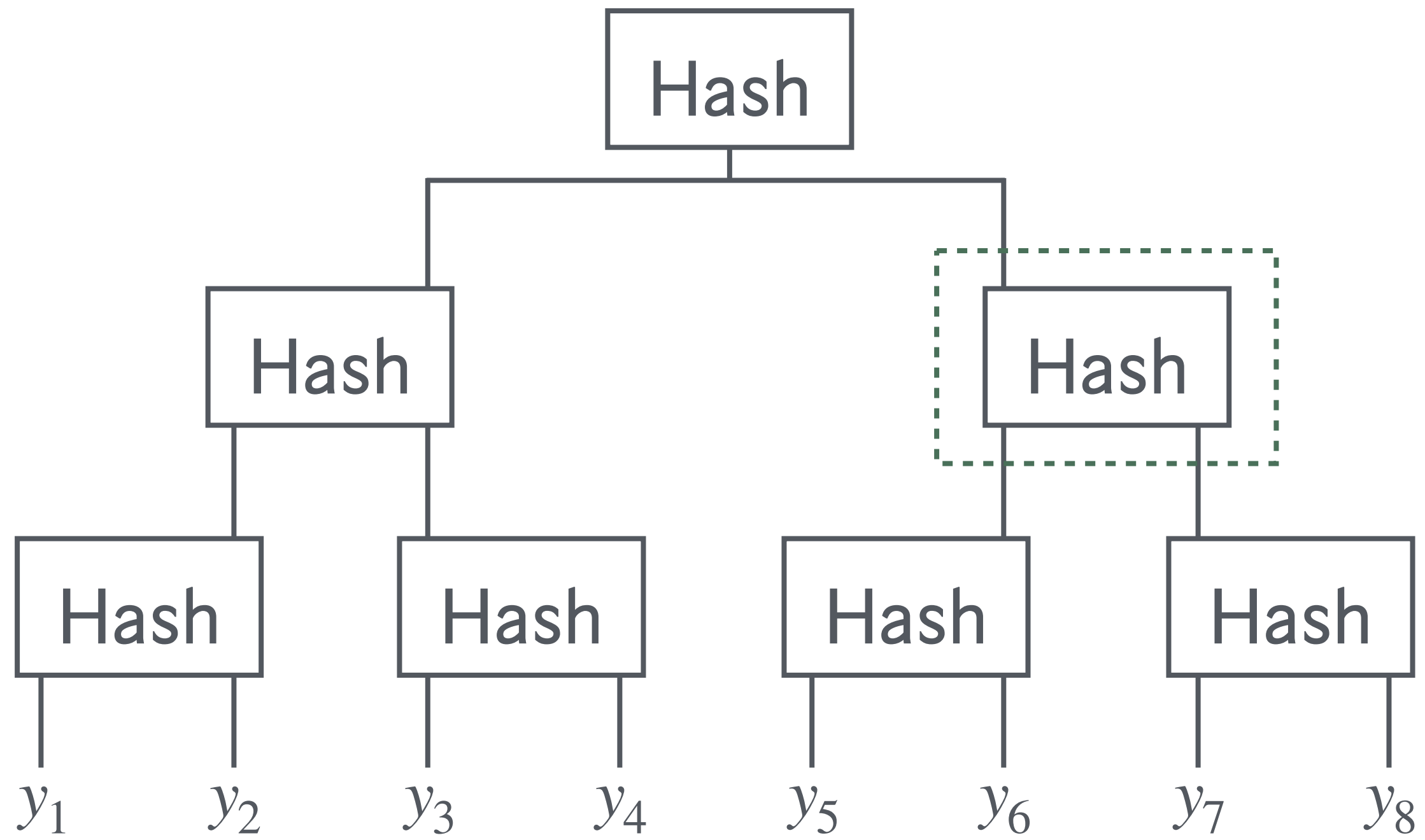
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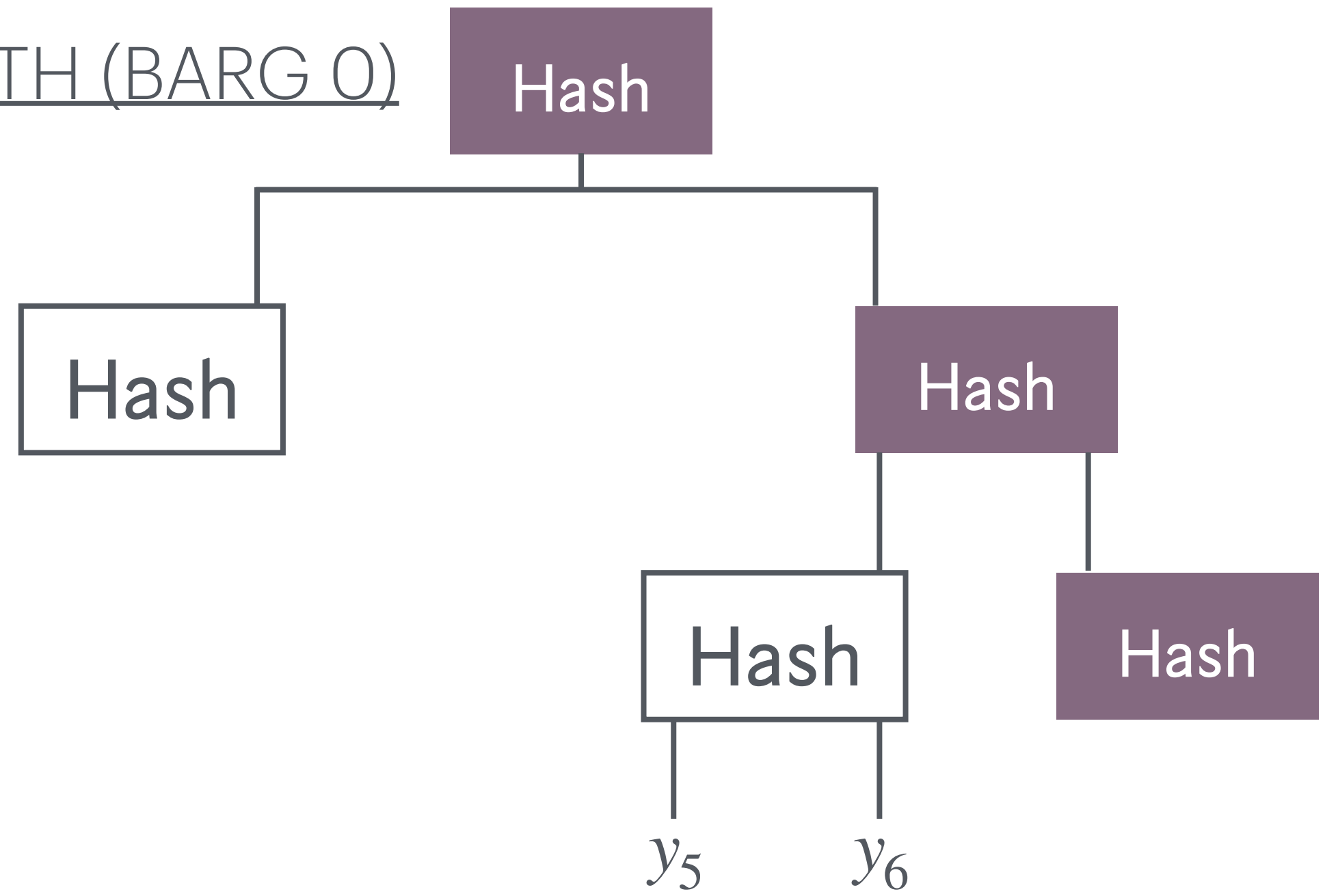
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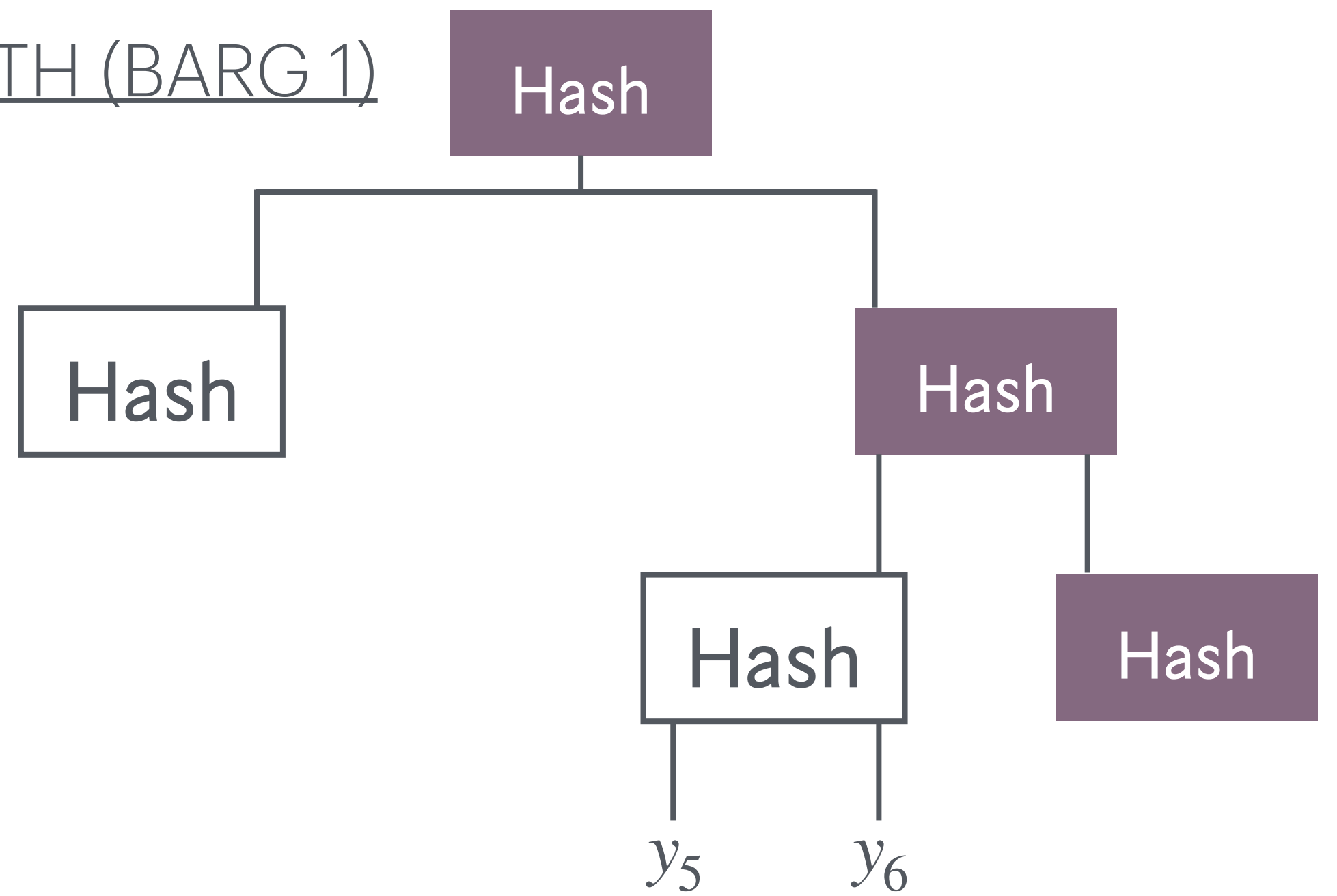
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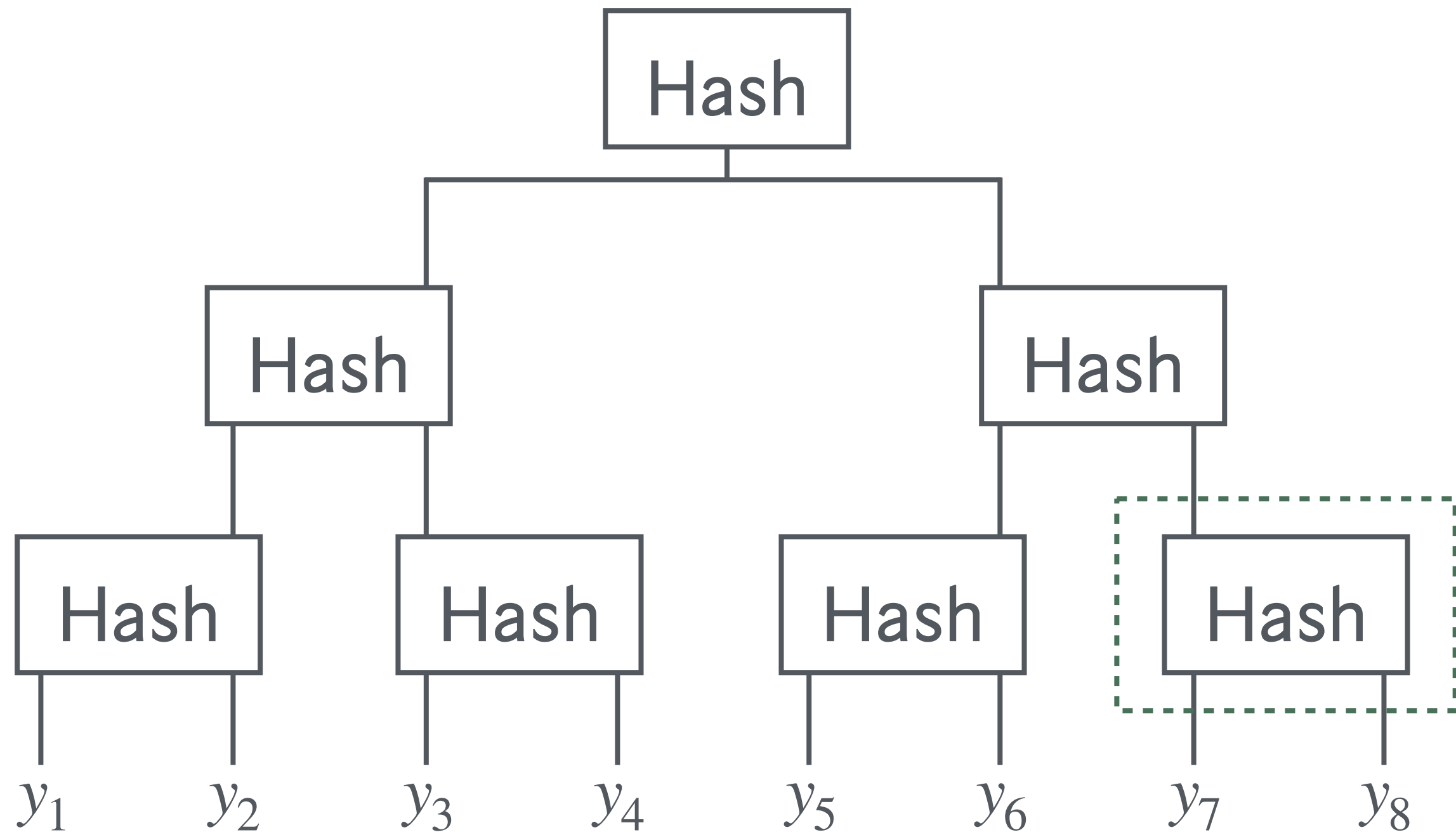
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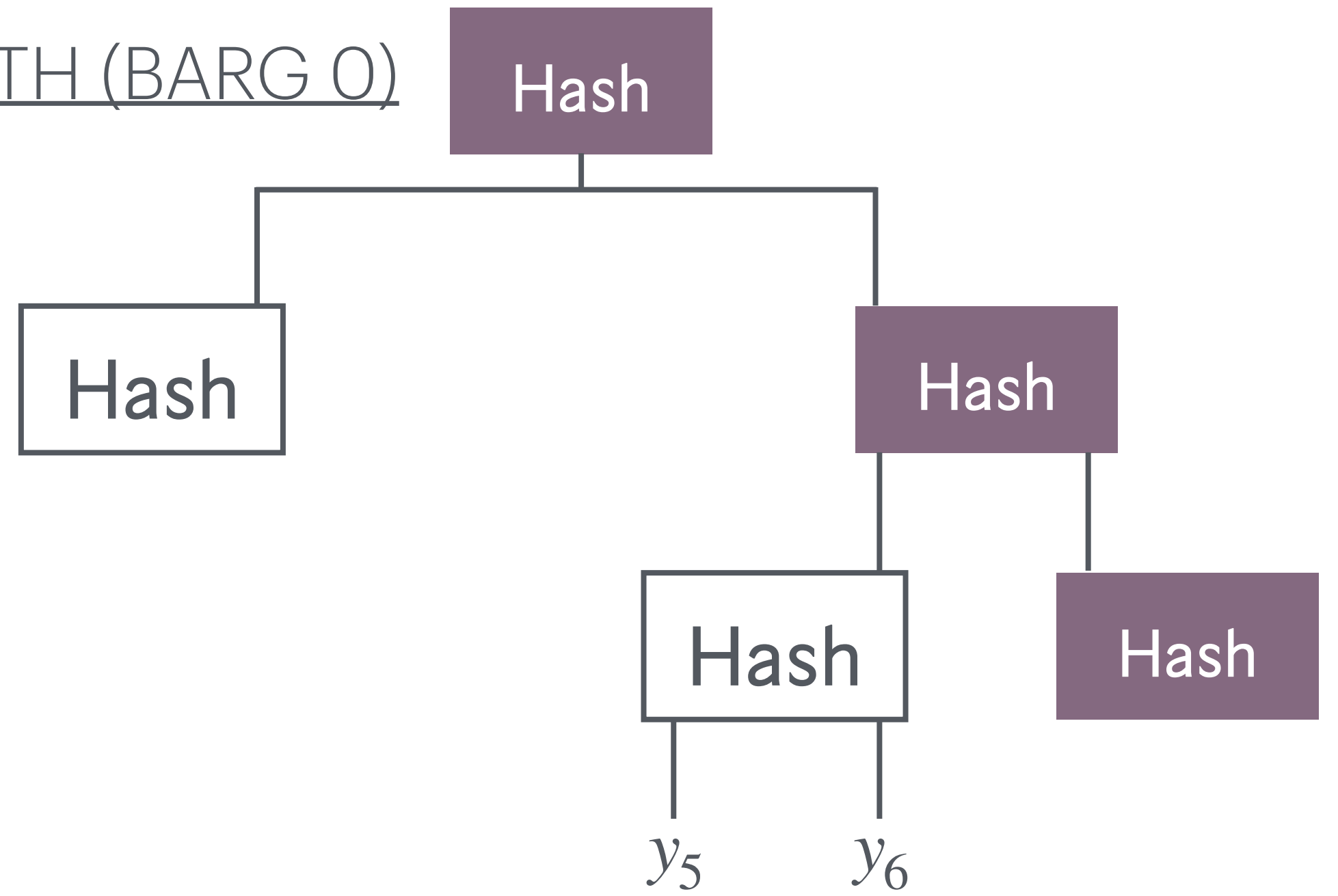
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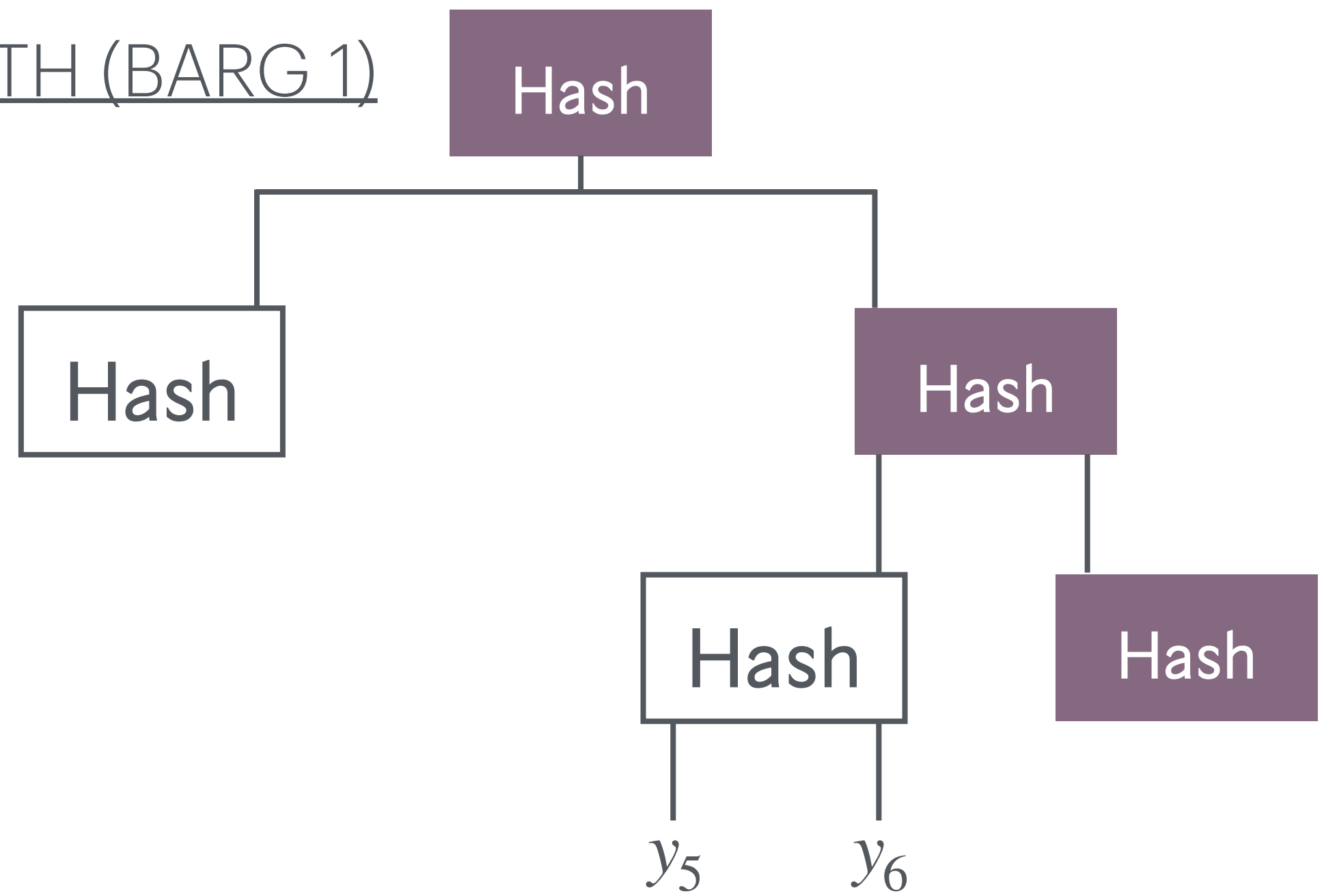
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Bonus

$$vk = (h, h^k)$$

$$sk = k$$

$$y = e(g, \mathcal{H}(x))^k$$

$$y \stackrel{?}{=} e(g, \mathcal{H}(x)^k) \quad \& \quad e(h^k, \mathcal{H}(x)) \stackrel{?}{=} e(h, \pi)$$

$\underbrace{\hspace{10em}}_{=: \pi}$

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Thank you!

