

Quantum State Group Actions

Saachi Mutreja (Columbia)

Mark Zhandry (NTT Research and Stanford University)

Cryptographic Groups

- A pair (G, \star) consisting of a set G and a binary operation $\star : G \times G \rightarrow G$ is an abelian group if the following properties hold:
 - Identity: $\exists e \in G : e \star a = a \star e = a, \forall a \in G.$
 - Inverse: $\forall a \in G, \exists b : b \star a = a \star b = e.$
 - Associativity: $(a \star b) \star c = a \star (b \star c)$
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- [Diffie-Hellman '76]: Hardness assumptions on (cyclic) groups \mathbb{G} .
 - Discrete log: $g, g^a \rightarrow a$
 - CDH: $g, g^a, g^b \rightarrow g^{ab}$
 - DDH: g, g^a, g^b, g^{ab} vs g, g^a, g^b, g^c

Quantum Computers break Cryptographic Groups

- Finding discrete log is easy:

Suppose $h = g^a$, want to find a

Define $F(x, y) = g^x h^y$

F is periodic: $F((x, y) + (-a, 1)) = F(x, y)$

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How can we use group-theoretic problems in a quantum secure manner?

Cryptographic Group Actions [Brassard-Yung'91]

- (Abelian) group \mathbb{G} acting on set X via an action \star .

- Identity:

If e is identity in \mathbb{G} , then $e \star x = x, \forall x \in X$.

- Compatibility:

$$g \star (h \star x) = (g \cdot h) \star x, \forall g, h \in \mathbb{G}, \forall x \in X.$$

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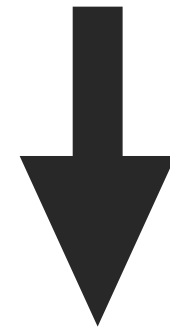
$$(x, g \star x) \rightarrow g$$

- Groups are a special case of group actions:

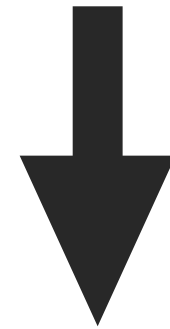
$$\mathbb{Z}_p \text{ acts on } \mathbb{G} \text{ via } a \star x = x^a.$$

Understanding the security of Group Actions

Justifying security in classical group + classical attack



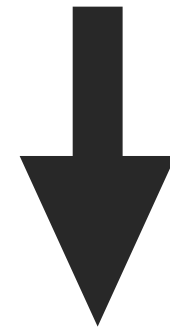
Prove security under hardness of some computational problem on group.



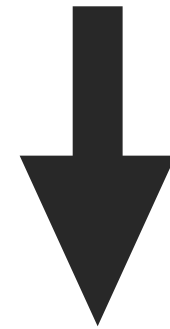
Justify hardness of the problem. (Can be justified in the *generic black box model*)

Understanding the security of Group Actions

Justifying security in classical group action + quantum attack



Prove security under hardness of some computational problem on group action.



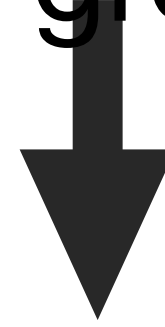
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[EH00,EHK04]:

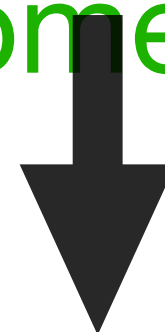
- there exist algorithms w/ polynomial query complexity (but super poly run time) that break crypto assumptions on GAs.

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Prove security under hardness of some computational problem on group.



Justify hardness of the problem?

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- there exist algorithms w/ polynomial query complexity (but super poly run time) that break crypto assumptions on GAs.
- Unconditional lower bounds not possible!

Prove unconditional hardness for similar assumptions on **Quantum State Group Actions.**

Quantum State Group Actions

- A quantum state group action will consist of:
 - A classical group action (\mathbb{G}, X, \star) .
 - A collection of states $\psi = (|\psi_x\rangle \in \mathcal{H})_{x \in X}$.
 - Distinguished starting element $x_* \in X$.
 - QPT procedure Start: produces $|\psi_{x_*}\rangle$.
 - QPT procedure Act ($g \in \mathbb{G}, |\psi_x\rangle \in \mathcal{H}$) produces $|\psi_{g \star x}\rangle = g \star |\psi_x\rangle$.

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 - $\forall x, y \in X : \langle \psi_x | \psi_y \rangle = 0$.
- Remark: $|\psi_{g \star x_*}\rangle$ is not necessarily clonable, hard problems on quantum group action parametrized by number of copies.

Quantum State Group Actions

- ℓ -DDH:

$$|\psi_a\rangle^{\otimes \ell} |\psi_b\rangle^{\otimes \ell} |\psi_{a+b}\rangle^{\otimes \ell} \approx |\psi_a\rangle^{\otimes \ell} |\psi_b\rangle^{\otimes \ell} |\psi_c\rangle^{\otimes \ell}$$

- ℓ - Discrete log:

$$\Pr[g \leftarrow \mathcal{A}(|\psi_{g \star x_*}\rangle^{\otimes \ell})] \leq \text{negl}$$

Hard (without parametrization by ℓ) if ℓ -hard for all polynomials ℓ .

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- DDH- special case of GMP with $M = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & 1 \end{bmatrix}$
- Linear Hidden Shift assumption: Generalizing GMP to non-uniform $s \in \{0,1\}^m$, uniformly random M .
- Extended Linear Hidden Shift assumption: Generalizing GMP to structured M .

Main Results

- A Hash based construction of Quantum State group actions

- Generalized matrix assumption
- DDH assumption
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Orthogonal

Main Results

- An attack in the “many copy” regime:
 - When hash based construction is orthogonal: **query bounded** , **computationally inefficient** quantum coset sampling attacks [EH00, EHK04] on classical group actions generalize given multiple copies.

Main Results

- **Unifying Quantum Money:**
- [Zha24]: constructed Quantum Money from abelian GA.

Generalize construction to quantum state group actions + instantiate
w/ Hash based construction



[Zha19]: Quantum Money construction from non collapsing hash
functions.

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 - Act ($g \in \mathbb{Z}_N$, $|\psi\rangle = \sum_{r \in R} \alpha_r |r\rangle$): $P_g |\psi\rangle$, $P_g : |r\rangle \rightarrow \omega_N^{g \cdot H(r)} |r\rangle$

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 - $P_g P_h = P_{g+h}$ (additive notion for group operation).

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$$\sum_x \omega_N^{H(x) \cdot g_1} |x\rangle \sum_y \omega_N^{H(y) \cdot g_2} |y\rangle \sum_z \omega_N^{H(z) \cdot (g_1+g_2)} |z\rangle$$

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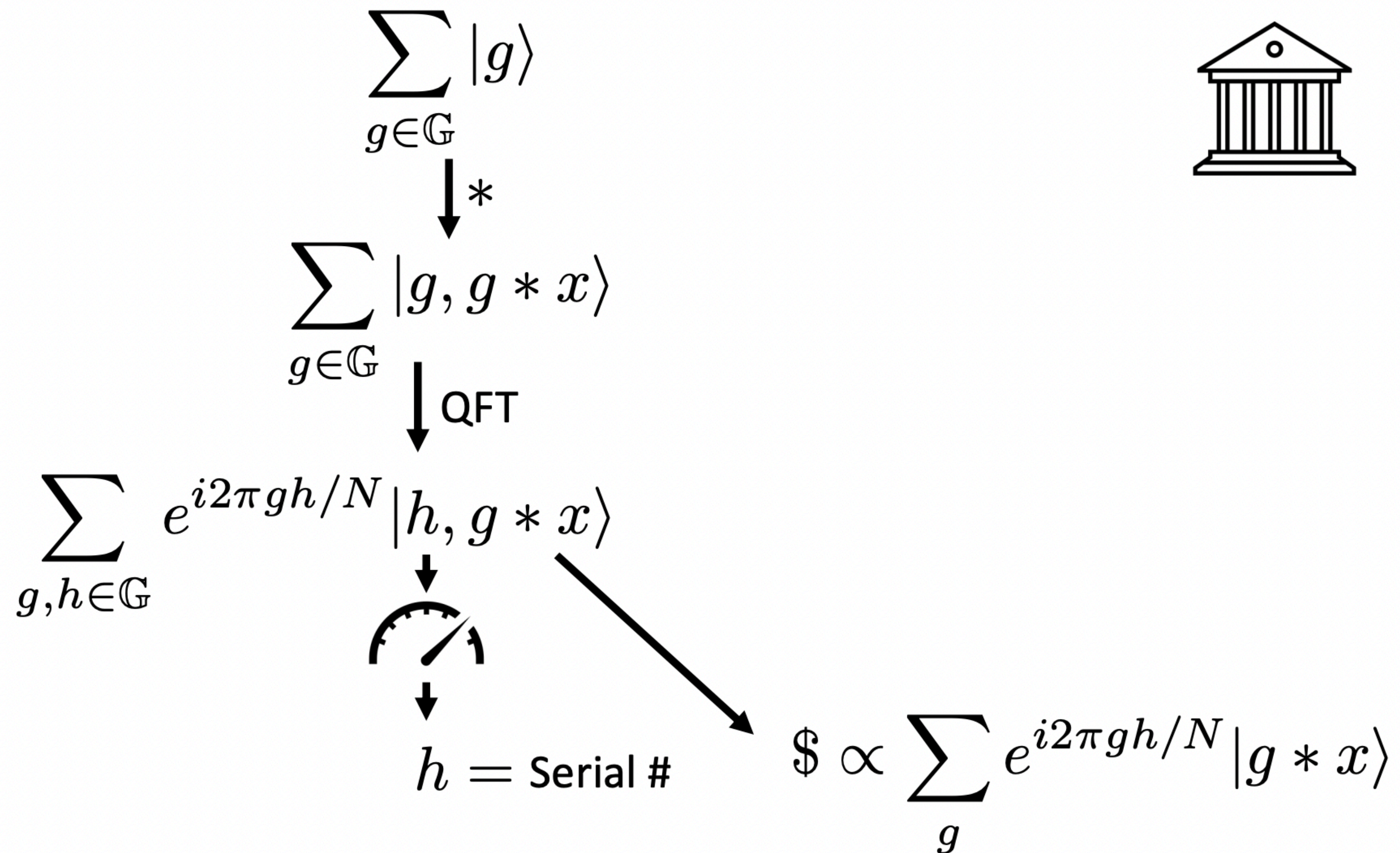
Averaging over choice of g_1, g_2, g_3 - resulting mixed states close if:

- $f'(x, y, z) = (H(x), H(y), H(z))$ and
- $f(x, y, z) = (H(x) + H(z), H(y) + H(z))$

are almost injective.

Unifying Quantum Money

**[Zha24] Q-Money
Scheme from abelian
group actions.**



Unifying Quantum Money

Plugging in hash-based quantum state group action:

$$| \$_h \rangle = \sum_g \omega_N^{gh} | \psi_{g \star x_*} \rangle$$

↓

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Money state from [Zha19] with non collapsing CRHF H .