



# SoK: Security of the Ascon Modes

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Radboud University

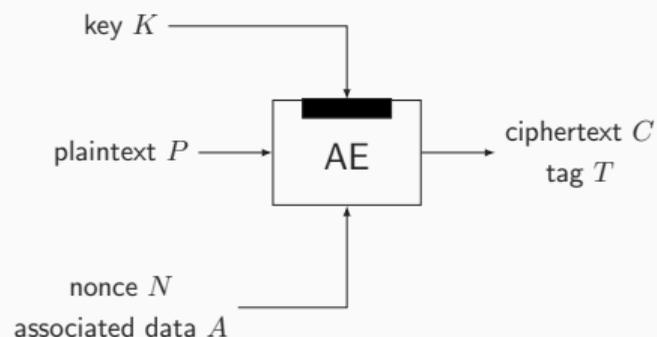
FSE 2025

March 17, 2025

# Introduction

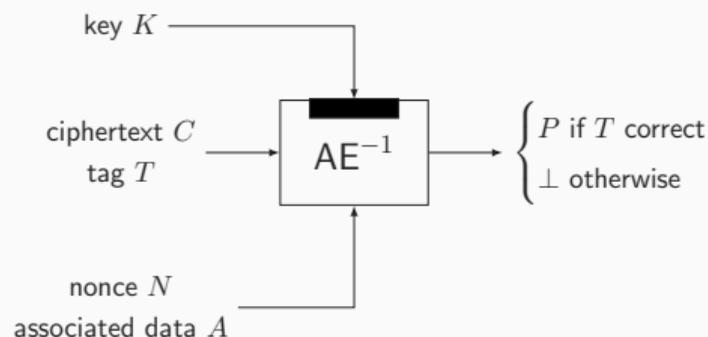
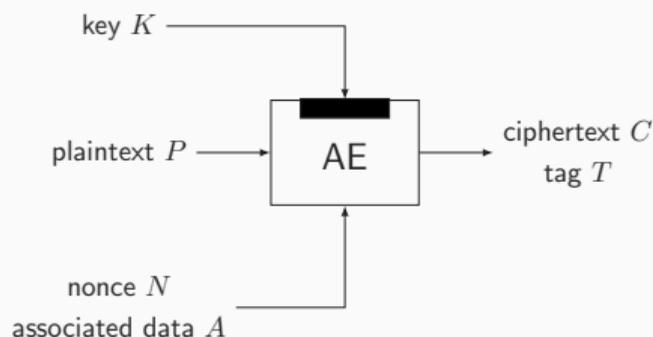
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# Authenticated Encryption



- Using key  $K$ :
  - Ciphertext  $C$  encrypts plaintext  $P$
  - Tag  $T$  authenticates  $(N, A, P)$

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  - Ciphertext  $C$  encrypts plaintext  $P$
  - Tag  $T$  authenticates  $(N, A, P)$
- Unwrapping needs to satisfy that
  - Plaintext disclosed if tag is **correct**
  - Plaintext is not leaked if tag is **incorrect**

## CAESAR Competition

- 2014–2019
- Call for authenticated encryption scheme
- 57 submissions (of which  $\approx 10$  sponge/duplex-based)
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## NIST Lightweight Cryptography Competition

- 2019–2023
- Call for authenticated encryption scheme and, optionally, hash function
- 57 submissions (of which  $\approx 22$  sponge/duplex-based)
- **Ascon selected as winner**





## Authenticated Encryption

- Duplex-based but with additional key blindings



## Authenticated Encryption

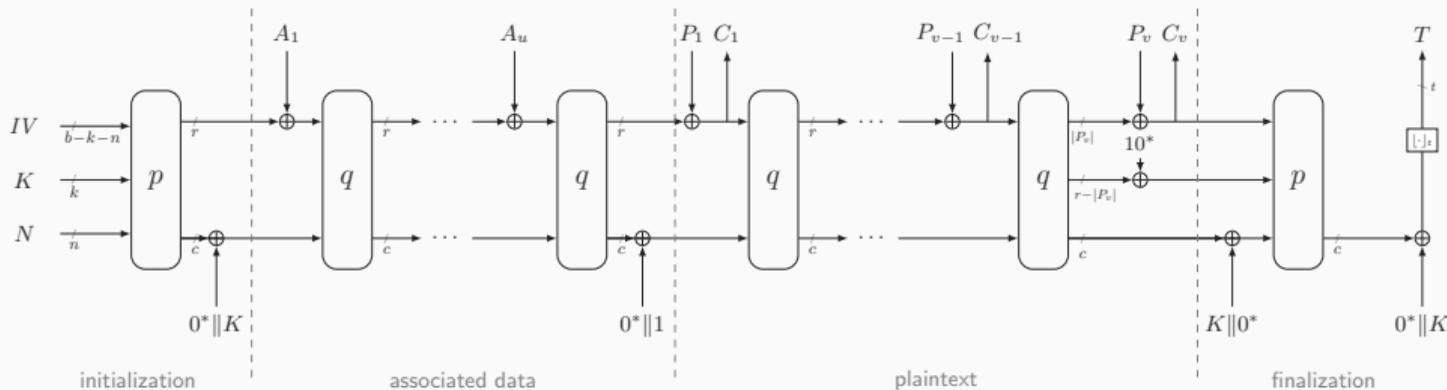
- Duplex-based but with additional key blindings

## Hashing

- Sponge-based hashing and XOFing
- Only included in NIST Lightweight Cryptography submission

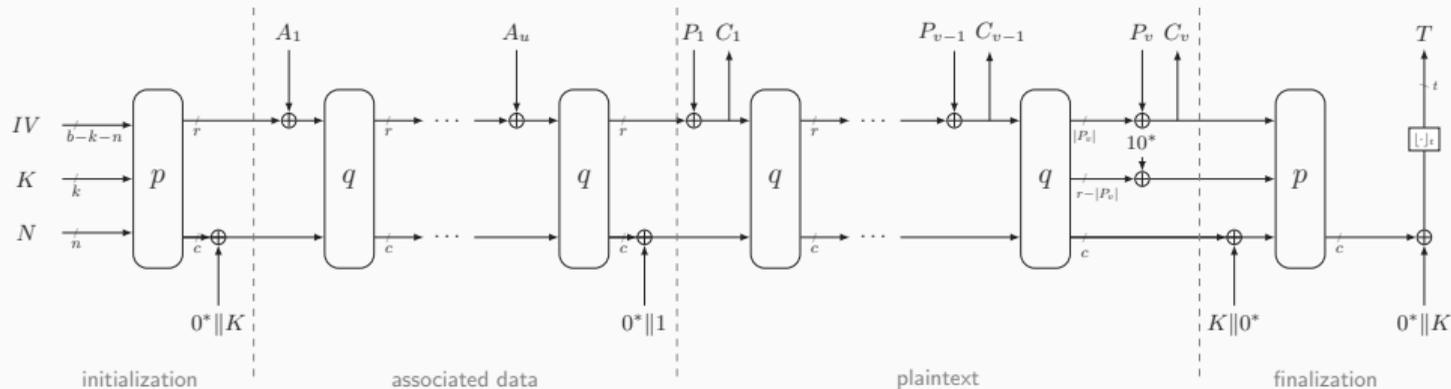
## Ascon-AE

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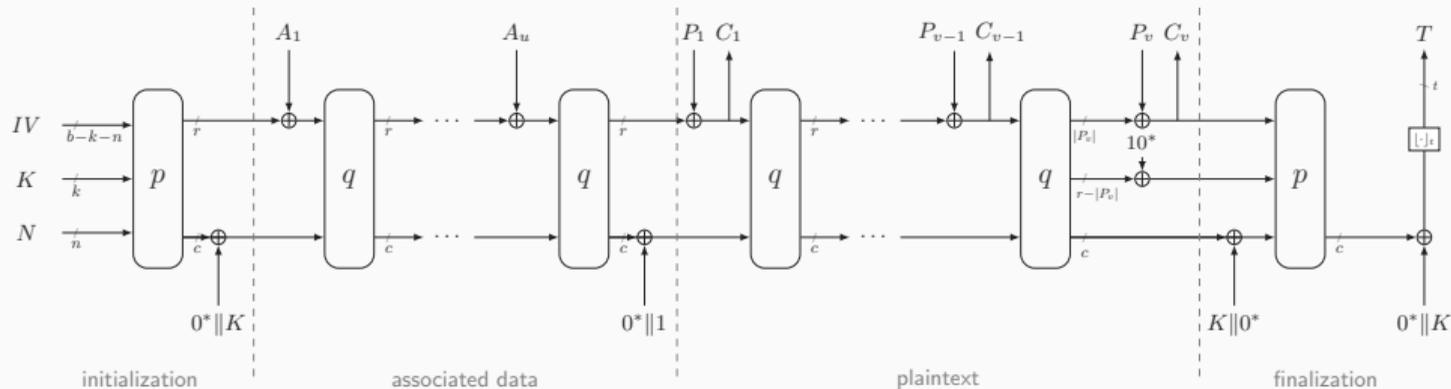
## Variant of SpongeWrap [BDPV11]

- Outer permutation  $p$  and inner permutation  $q$ , both on  $b$  bits
  - $r$  is the rate,  $c$  is the capacity (security parameter)



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- Additional key blindings around “outer” permutations
- Domain separation simplified and spilled-over into inner part

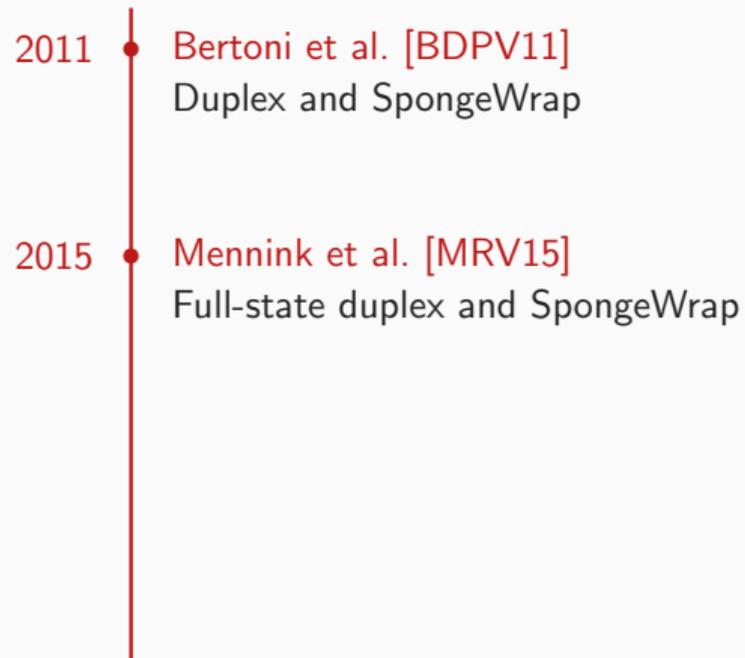
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none of these  
results deals with  
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### Dedicated Ascon Analysis



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- 2023 • Chakraborty et al. [CDN23]  
Single-user security in nonce-respecting setting

## Dedicated Ascon Analysis

- 
- A vertical red line on the left side of the slide serves as a timeline axis. Two red circular markers are placed on this line, corresponding to the years 2023 and 2024. The text for each year is positioned to the right of its respective marker.
- 2023 • Chakraborty et al. [CDN23]  
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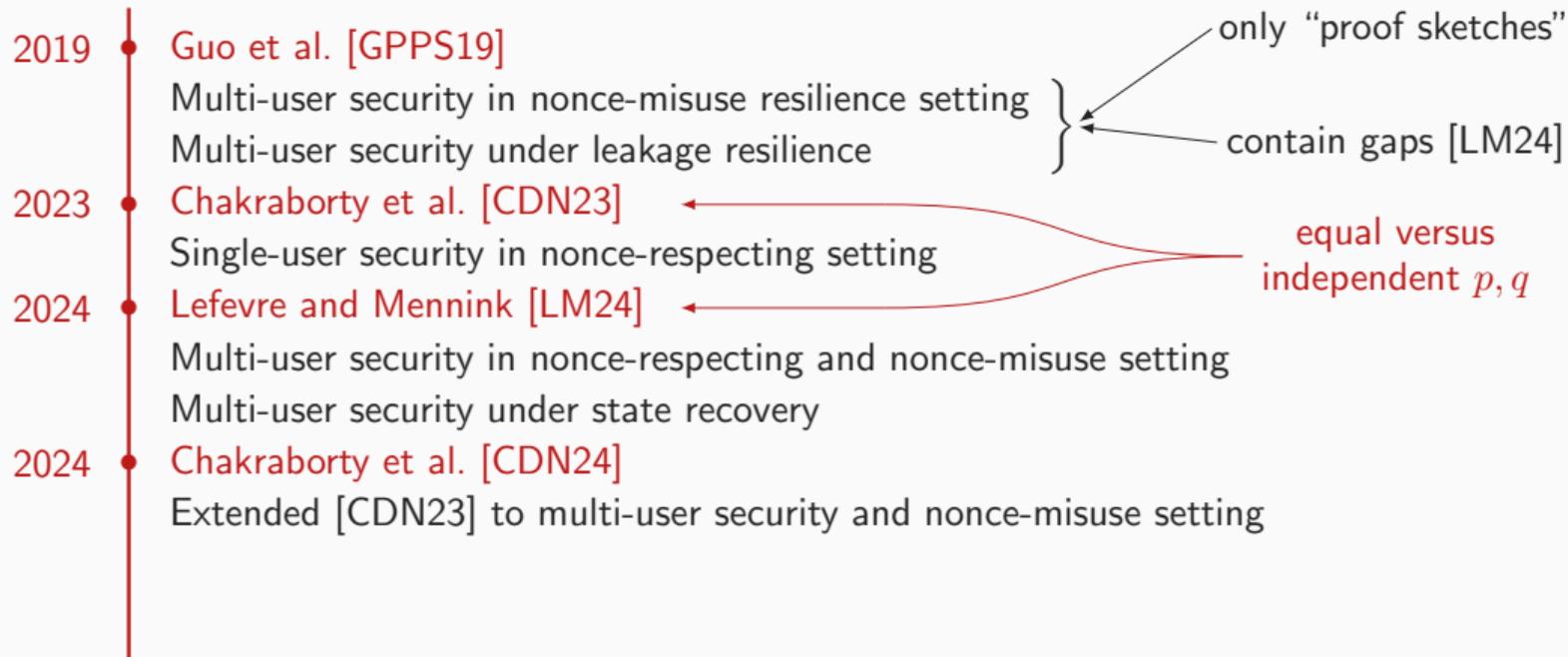
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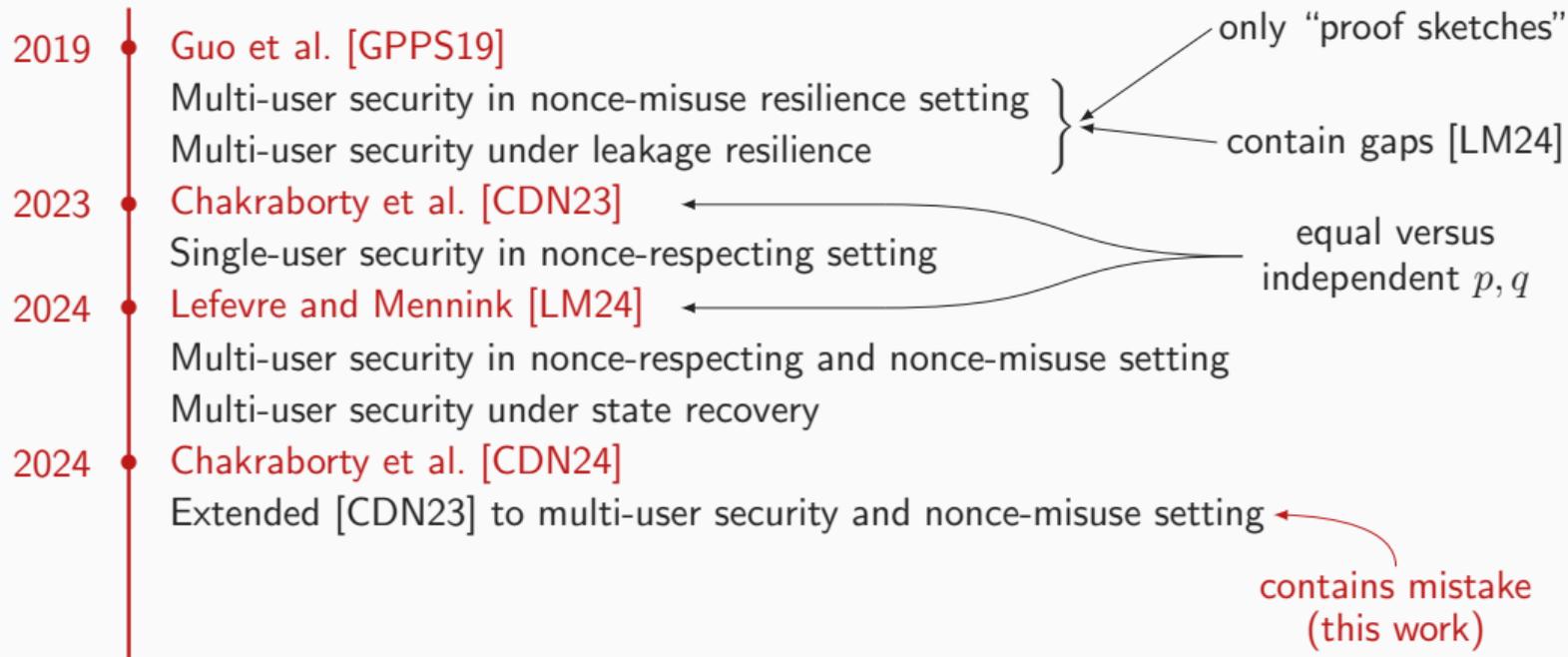
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- only "proof sketches"  
contain gaps [LM24]

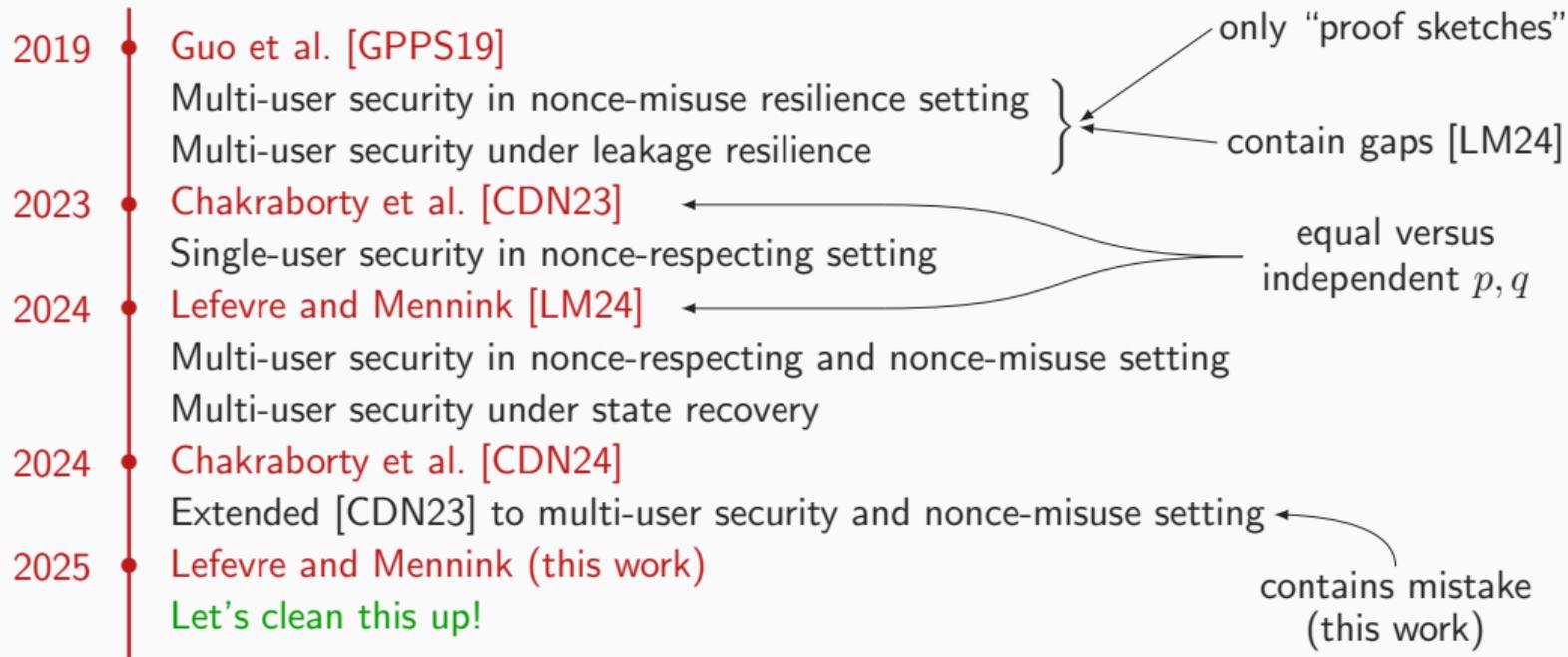
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### Complete Overview of Generic Security of the Ascon-AE Mode

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- Three flavors of conventional security:
  - ① Nonce-respecting security [BN00]
  - ② Nonce-misuse resistance [RS06]
  - ③ Nonce-misuse resilience [ADL17]

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- We **categorize** existing lower and upper bounds
- We **derive** new security bounds and matching attacks where needed
- All results assume that  $p = q$  is a random permutation

# Overview of Results on Ascon-AE

nonce-respecting security
confidentiality
authenticity

# Overview of Results on Ascon-AE

nonce-respecting security	
confidentiality	$\frac{\mu\mathcal{N}}{2^k} + \frac{\mathcal{M}\mathcal{N}}{2^b} + \frac{\mathcal{N}}{2^c}$
authenticity	$\frac{Q_D}{2^t} + \frac{\mu\mathcal{N}}{2^k} + \frac{\mathcal{M}\mathcal{N}}{2^b} + \frac{\mathcal{N}}{2^c}$

- $\mu$  number of users
- $Q_E/\mathcal{M}_E$  encryption queries/complexity
- $Q_D/\mathcal{M}_D$  decryption queries/complexity
- $Q/\mathcal{M}$  construction queries/complexity
- $\mathcal{N}$  permutation queries

bounds of [CDN23, CDN24]

carry over

new: matching attacks

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core term (★)

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nonce-misuse resistance	
confidentiality	1
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

new: flaw in proof of [CDN24]

new: transformation of [LM24]

to  $p = q$  setting

new: matching attacks

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core term (★)

nonce-misuse resilience	
confidentiality	(★) + $\frac{\mathcal{M}_E\mathcal{N}}{2^c}$
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analysis of [GPPS19] incomplete  
new: security bounds  
and matching attacks

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nonce-misuse resilience	
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leakage resilience, no leakage	
confidentiality	
authenticity	

leakage resilience, limited	
confidentiality	
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follows by equivalence



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leakage resilience, limited	
confidentiality	
authenticity	

leakage resilience, unlimited	
confidentiality	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$

analysis of [GPPS19] incomplete  
and in different model  
new: security bounds  
and matching attacks

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follows by implication

nonce-misuse resilience	
confidentiality	(*) + $\frac{\mathcal{M}_E\mathcal{N}}{2^c}$
authenticity	(*) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

leakage resilience, no leakage	
confidentiality	(*) + $\frac{\mathcal{M}_E\mathcal{N}}{2^c}$
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leakage resilience, limited	
confidentiality	(*) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$
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leakage resilience, unlimited	
confidentiality	(*) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$
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nonce-misuse resistance	
confidentiality	1
authenticity	(*) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

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new: transformation of [LM24]  
 to  $p = q$  setting  
 new: matching attacks

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nonce-misuse resistance	
confidentiality	1
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

state-recovery security	
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authenticity	(★) + $\frac{\mathcal{N}^2}{2^c}$

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core term (★)

$\mu$  number of users  
 $Q_E/\mathcal{M}_E$  encryption queries/complexity  
 $Q_D/\mathcal{M}_D$  decryption queries/complexity  
 $Q/\mathcal{M}$  construction queries/complexity  
 $\mathcal{N}$  permutation queries

new: security bounds and matching attacks

nonce-misuse resilience	
confidentiality	(★) + $\frac{\mathcal{M}_E\mathcal{N}}{2^c}$
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

leakage resilience, no leakage	
confidentiality	(★) + $\frac{\mathcal{M}_E\mathcal{N}}{2^c}$
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

leakage resilience, limited	
confidentiality	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$

leakage resilience, unlimited	
confidentiality	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c} + \min\left\{\frac{\mathcal{N}^2}{2^c}, \frac{Q\mathcal{N}}{2^k}\right\}$

nonce-misuse resistance	
confidentiality	1
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

RUP security	
confidentiality	1
authenticity	(★) + $\frac{\mathcal{M}\mathcal{N}}{2^c}$

state-recovery security	
confidentiality	1
authenticity	(★) + $\frac{\mathcal{N}^2}{2^c}$

# Simplified Numerical Interpretation

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setting	confidentiality as long as	authenticity as long as
nonce-respecting		
nonce-misuse resilience		
nonce-misuse resistance		
state-recovery security		

---

# Simplified Numerical Interpretation

setting	confidentiality as long as	authenticity as long as
nonce-respecting nonce-misuse resilience nonce-misuse resistance state-recovery security	$\mathcal{N} \ll \min\{2^k/\mu, 2^b/\mathcal{M}, 2^c\}$	$\mathcal{N} \ll \min\{2^k/\mu, 2^b/\mathcal{M}, 2^c\}, Q_D \ll 2^t$

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nonce-misuse resilience	$\mathcal{N} \ll \min\{2^k/\mu, 2^c/\mathcal{M}\}$	$\mathcal{N} \ll \min\{2^k/\mu, 2^c/\mathcal{M}\}, Q_D \ll 2^t$
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## Application to Ascon-AEAD Parameters

- $(k, b, c, r, t) = \begin{cases} (128, 320, 256, 64, 128) & \text{for Ascon-128} \\ (128, 320, 192, 128, 128) & \text{for Ascon-128a} \\ (160, 320, 256, 64, 128) & \text{for Ascon-80pq} \end{cases}$
- Assume online complexity of  $Q, \mathcal{M} \ll 2^{64}$  (could be taken higher)

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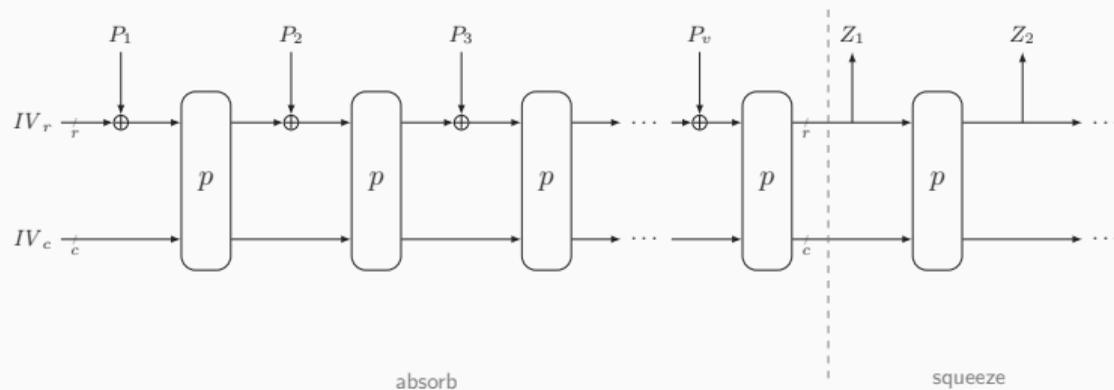
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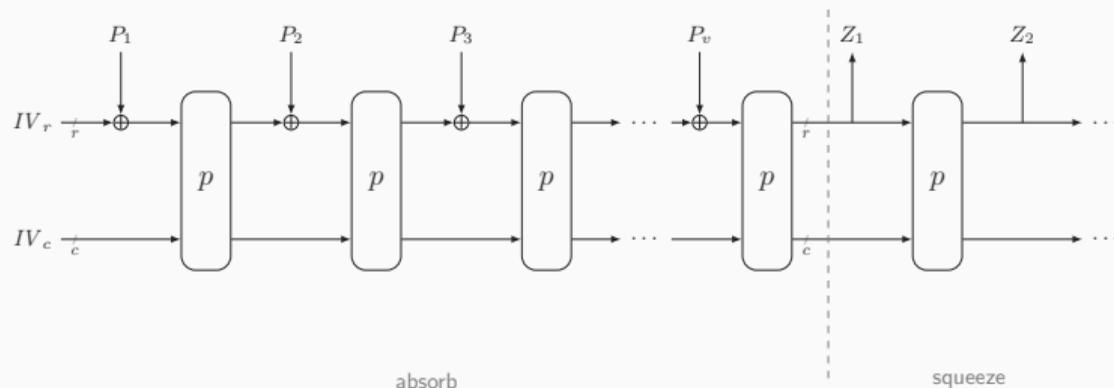
- $(k, b, c, r, t) = \begin{cases} (128, 320, 256, 64, 128) & \text{for Ascon-128} \\ (128, 320, 192, 128, 128) & \text{for Ascon-128a} \\ (160, 320, 256, 64, 128) & \text{for Ascon-80pq} \end{cases}$
- Assume online complexity of  $Q, \mathcal{M} \ll 2^{64}$  (could be taken higher)
- Generic** security as long as  $\mathcal{N} \ll 2^{128}/\mu$   
 (exceptions:  $\mathcal{N} \ll 2^{160}/\mu$  for Ascon-80pq;  $\mathcal{N} \ll 2^{96}$  for Ascon-128a under state-recovery)

# Ascon-Hash/Ascon-(C)XOF

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# Ascon-Hash/Ascon-(C)XOF





## Sponge [BDPV07]

- Permutation  $p$  on  $b$  bits
  - $r$  is the rate
  - $c$  is the capacity (security parameter)
- Output of  $\nu$  bits (256 for Ascon-Hash, unlimited for the XOFs)

- Sponge indifferentiable from random up to bound  $\mathcal{N}^2/2^c$  [BDPV08]

## Generic Security of the Sponge

- Sponge indifferentiable from random up to bound  $\mathcal{N}^2/2^c$  [BDPV08]
- Security of sponge truncated to  $\nu$  bits against classical attacks [AMP10]:

Collision resistance:  $\mathcal{N}^2/2^c + \mathcal{N}^2/2^{\nu+1}$

Second preimage resistance:  $\mathcal{N}^2/2^c + \mathcal{N}/2^\nu$

Preimage resistance:  $\mathcal{N}^2/2^c + \mathcal{N}/2^\nu$

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distance from sponge to RO  
( $\mathcal{N}$  is # primitive evaluations)

classical attacks against RO  
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Preimage resistance:	$\mathcal{N}^2/2^c + \mathcal{N}/2^\nu$	← attack in $\min\{2^{\nu-r} + 2^{c/2}, 2^\nu\}$
	↑	↑
	distance from sponge to RO ( $\mathcal{N}$ is # primitive evaluations)	classical attacks against RO ( $\mathcal{N}$ is # oracle evaluations)

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$\uparrow$  distance from sponge to RO       $\uparrow$  classical attacks against RO  
( $\mathcal{N}$  is # primitive evaluations)      ( $\mathcal{N}$  is # oracle evaluations)

- Attacks already described in [BDPV07]
- Tightened preimage resistance bound by Lefeuvre and Mennink [LM22]:

Preimage resistance:  $\min\{\mathcal{N}/2^{\nu-r}, \mathcal{N}/2^{c/2}\} + \mathcal{N}/2^\nu \quad \leftarrow$  attack in  $\min\{2^{\nu-r} + 2^{c/2}, 2^\nu\}$

## Application to Ascon-Hash and Ascon-(C)XOF Parameters

- $(b, c, r, \nu) = \begin{cases} (320, 256, 64, 256) & \text{for Ascon-Hash} \\ (320, 256, 64, \infty) & \text{for Ascon-XOF} \\ (320, 256, 64, \infty) & \text{for Ascon-CXOF} \end{cases}$

## Application to Ascon-Hash and Ascon-(C)XOF Parameters

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- **Generic** collision resistance as long as  $\mathcal{N} \ll \min\{2^{128}, 2^{\nu/2}\}$

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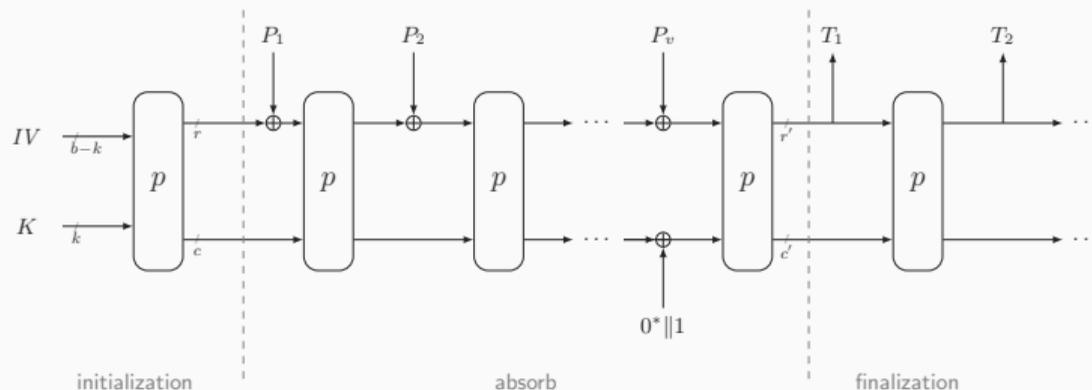
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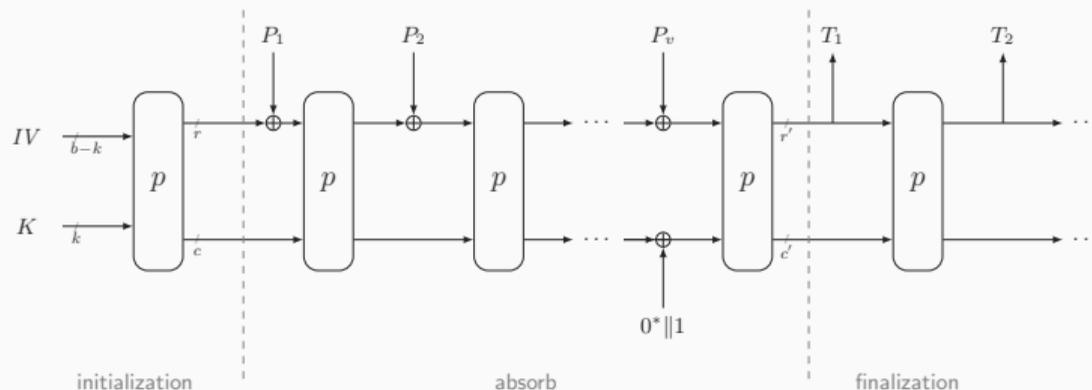
## Bonus: Ascon-PRF

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### Variant of Full-State Keyed Sponge [BDPV12, MRV15]

- Permutation  $p$  on  $b$  bits
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- Permutation  $p$  on  $b$  bits
  - $r$  is the rate,  $c$  is the capacity (security parameter)
- Domain separation to avoid squeezed tags being misused in absorption

## FSKS and Ascon-PRF



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- 2015 • Mennink et al. [MRV15]  
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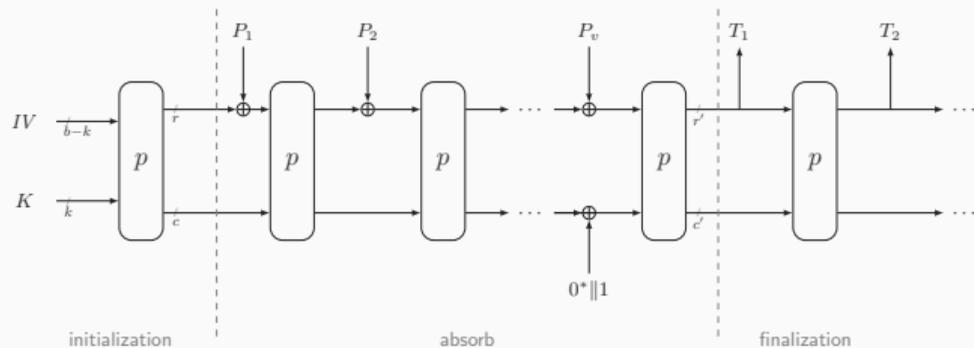
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Duplex guide and improved analysis of Ascon-PRF
- 2025 • Lefevre and Mennink (this work)  
Adapt bound of [Men23] with improved multicollision strategy

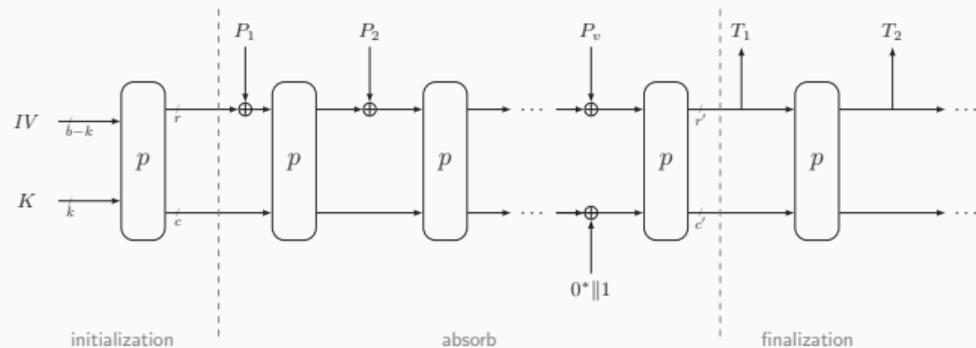
## Generic Security of Ascon-PRF (2/2)



### Generic Security Bound

- Ascon-PRF is multi-user secure up to bound  $\frac{\mu\mathcal{N}}{2^k} + \frac{\mathcal{N}}{2^{c'}} + \frac{\mathcal{M}\mathcal{N}}{2^b}$

## Generic Security of Ascon-PRF (2/2)



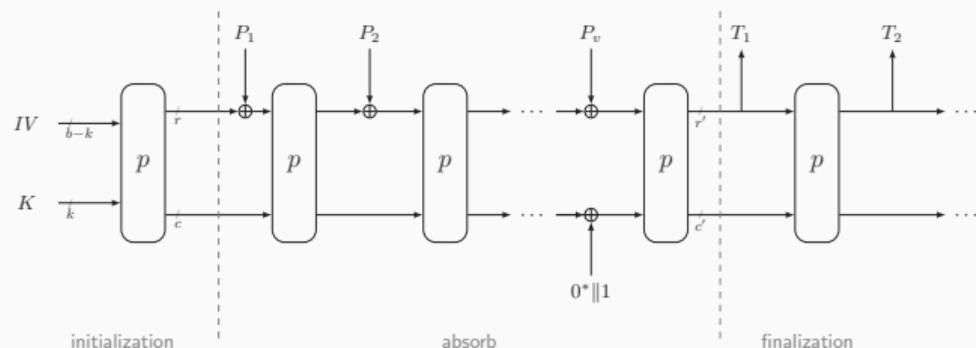
### Generic Security Bound

- Ascon-PRF is multi-user secure up to bound  $\frac{\mu\mathcal{N}}{2^k} + \frac{\mathcal{N}}{2^{c'}} + \frac{\mathcal{M}\mathcal{N}}{2^b}$

### Application to Ascon-PRF Parameters

- $(k, b, c, r, c', r', t) = (128, 320, 64, 256, 192, 128, \infty)$
- Assume online complexity of  $\mathcal{M} \ll 2^{64}$  (could be taken higher)

## Generic Security of Ascon-PRF (2/2)



### Generic Security Bound

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- Assume online complexity of  $\mathcal{M} \ll 2^{64}$  (could be taken higher)
- Generic security as long as  $\mathcal{N} \ll 2^{128}/\mu$

## Conclusion

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**More in Paper:** <https://eprint.iacr.org/2024/1969>

- Exact security models, settings, and discussions
- Discussion on multicollision bounding, assumption on  $p, q, \dots$
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## What We Did Not Cover

- Related-key security and security for arbitrary key distributions
- Security under fault attacks
- Variant with nonce masking [DM24]
- Committing security  $\rightarrow$  next talk!

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