### Kleptographic Attacks against Implicit Rejection

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#### CPA Secure PKE



#### **CPA** Secure PKE

Analysis in the QROM





- Analysis in the QROM
- Used in NIST PQC Submissions, e.g., Kyber











**Explicit Rejection** 

Success	$K := f_{sk}(c)$
Failure	$K := \bot$

**Explicit Rejection** 

Success	$K := f_{sk}(c)$
Failure	$K := \bot$

#### Implicit Rejection

Success $K := f_{sk}(c)$ FailureK := H(c, s)

Explicit Rejection



#### Implicit Rejection

 $K := f_{sk}(c)$ Success K := H(c, s)Failure

**Explicit Rejection** 



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 $K := f_{sk}(c)$ Success K := H(c, s)Failure

#### • Implicit rejection: tighter bounds

**Explicit Rejection** 



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 $K := f_{sk}(c)$ Success K := H(c, s)Failure

- Implicit rejection: tighter bounds
- Implicit rejection in Kyber

**Explicit Rejection** 



#### Implicit Rejection

 $K := f_{sk}(c)$ Success K := H(c, s)Failure

- Implicit rejection: tighter bounds
- Implicit rejection in Kyber

#### **Our Observation:**

Implicit rejection can be less secure!

#### Implicit rejection in a kleptographic setting

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Kleptographic attacker can break it!

## Implicit rejection in a kleptographic setting

	Subvert	Memory	Time Offline	Time Online	Advantage
Attack I	Decaps	2 <sup>8</sup>	20	$2^{2}$	0.997
Attack 2	Key Gen	27	$2^{0}$	$2^{130}$	0.999
Attack 3	Key Gen	$2^{111}$	$2^{154}$	$2^{106}$	0.692

#### Kleptographic attacker can break it!

\* applied to Kyber



Cryptographic Algorithm





#### • Kleptographic attacker subverts algorithm

## Cryptographic Algorithm



#### • Kleptographic attacker subverts algorithm

Attacker's goals: Success and Undetectability













#### • Success: attacker breaks security for subverted user



# Kleptographic Attacker's Goals Cryptographic Algorithm



#### • Undetectability: user cannot detect subversion



 $K := f_{sk}(c)$ 



Leaks nothing about sk



Leaks nothing about sk

• Make seed s depend on sk



s = H(ak, truncate(sk))

Leaks nothing about sk

• Make seed s depend on sk



s = H(ak, truncate(sk))

Leaks nothing about sk

• Make seed s depend on sk

Undetectable if ak is random



s = H(ak, truncate(sk))

- Make seed s depend on sk
- Undetectable if *ak* is random
- Rejection keys K leak bits of sk

Leaks nothing about sk



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- Make seed s depend on sk
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- Rejection keys K leak bits of sk

Leaks nothing about *sk* 

#### • Mitigations?

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#### Conclusion for NIST PQC Standardization?

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Can we apply technique to other primitives, like PAKE?

### Thank you!



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