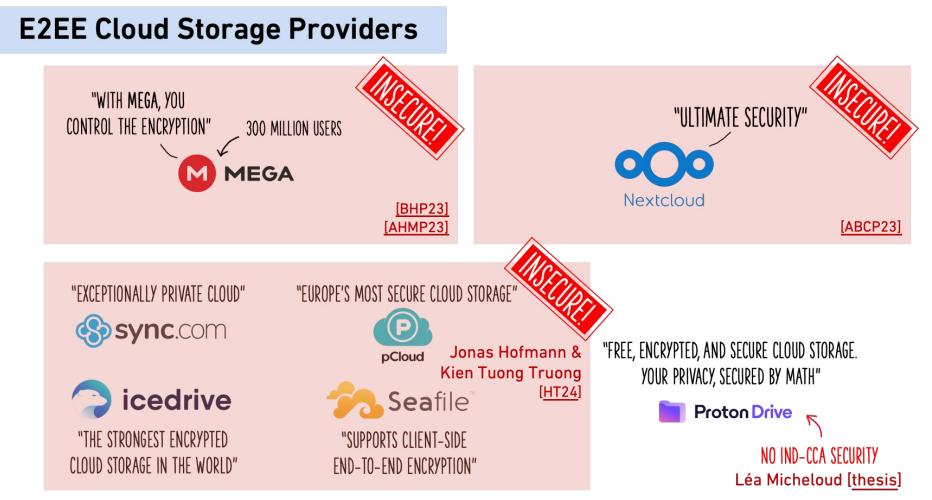
Mind the Gap! Secure File Sharing, from Theory to Practice

Matilda Backendal, David Balbás, Nicola Dardanis, Miro Haller, Matteo Scarlata

Sofia, 28 March 2025

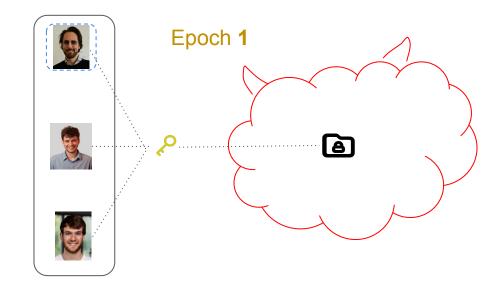




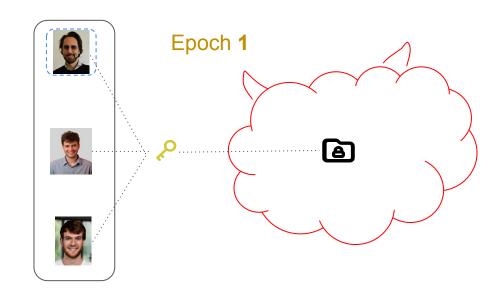
E2EE file sharing and secure shared folders (SSF) are particularly tricky.

• Uploader shares files pairwise

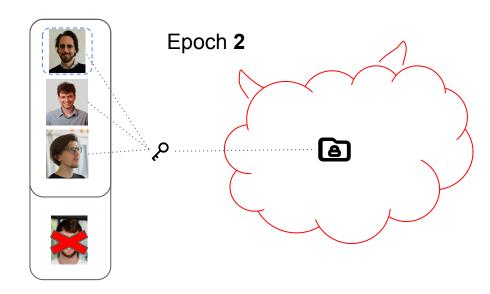
- Uploader shares files pairwise
- MEGA: shared static folder key



- Uploader shares files pairwise
- MEGA: shared static folder key
- **Dynamic members**: access rights change!



- Uploader shares files pairwise
- MEGA: shared static folder key
- **Dynamic members**: access rights change!

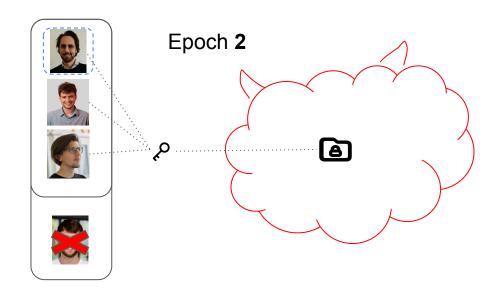


E2EE file sharing and secure shared folders (SSF) are particularly tricky.

- Uploader shares files pairwise
- MEGA: shared static folder key

Dynamic members: access rights change!

X State exposure



E2EE file sharing and secure shared folders (SSF) are particularly tricky.

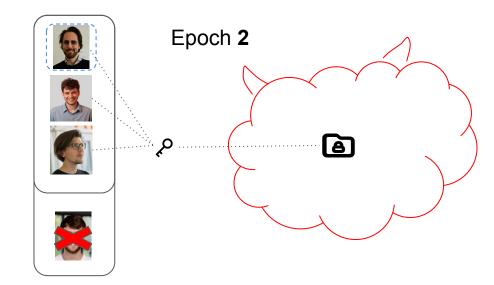
- Uploader shares files pairwise
- MEGA: shared static folder key

Dynamic members: access rights change!

X State exposure

Build SSF based on **group keys**:

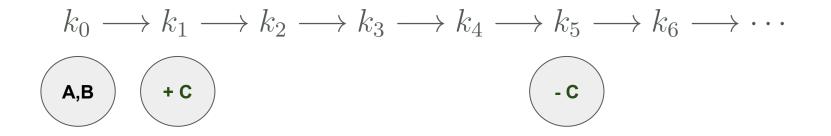
- Strong security
- Efficiency
- Real-world usability



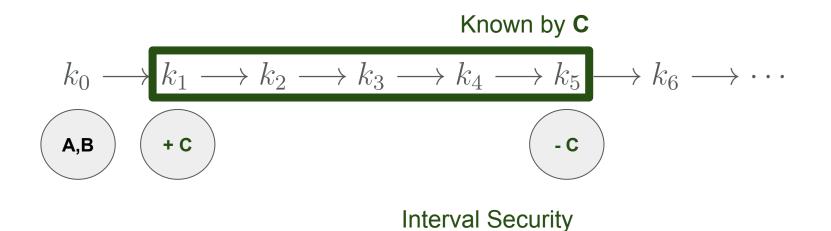
What epoch keys can we protect?

$$k_0 \longrightarrow k_1 \longrightarrow k_2 \longrightarrow k_3 \longrightarrow k_4 \longrightarrow k_5 \longrightarrow k_6 \longrightarrow \cdots$$

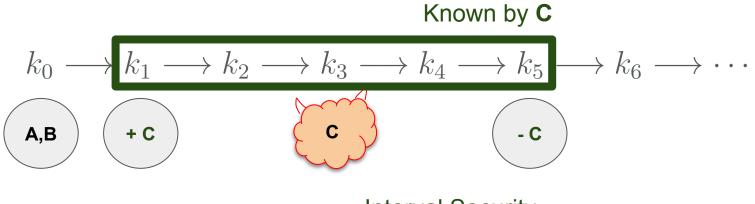
What epoch keys can we protect?



What epoch keys can we protect?

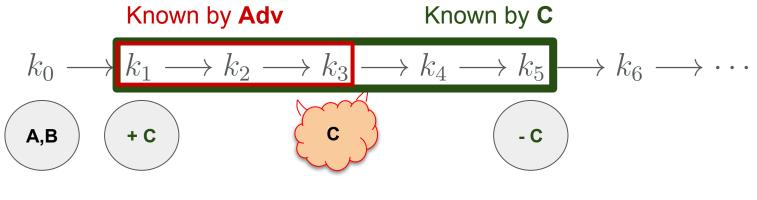


What epoch keys can we protect?



Interval Security

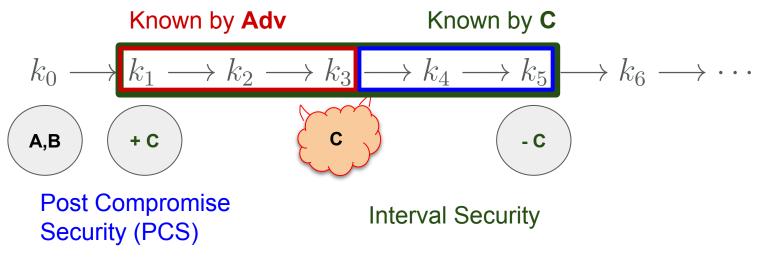
What epoch keys can we protect?



Interval Security

Persistency: leakage of k_1, k_2 unavoidable

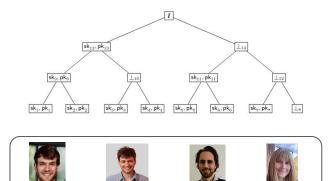
What epoch keys can we protect?



Persistency: leakage of k_1, k_2 unavoidable

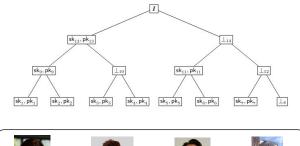
Natural attempt:

• Run group key agreement (e.g. CGKA as in MLS)



Natural attempt:

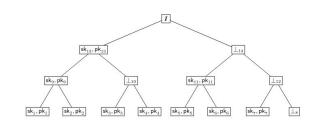
- Run group key agreement (e.g. CGKA as in MLS)
- Derive fresh k_i per epoch





Natural attempt:

- Run group key agreement (e.g. CGKA as in MLS)
- Derive fresh k_i per epoch
- Store all keys

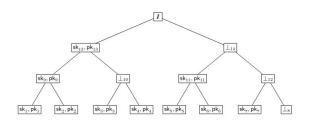




Natural attempt:

- Run group key agreement (e.g. CGKA as in MLS)
- Derive fresh k_i per epoch
- Store all keys

Security (from CGKA) & MLS implementation



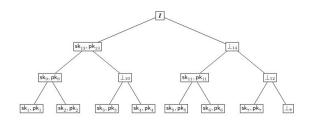


Natural attempt:

- Run group key agreement (e.g. CGKA as in MLS)
- Derive fresh k_i per epoch
- Store all keys

Security (from CGKA) & MLS implementation

 \mathbf{X} State grows linearly on number of epochs





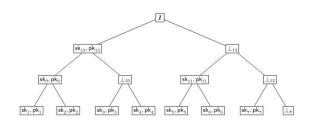
Natural attempt:

- Run group key agreement (e.g. CGKA as in MLS)
- Derive fresh k_i per epoch
- Store all keys

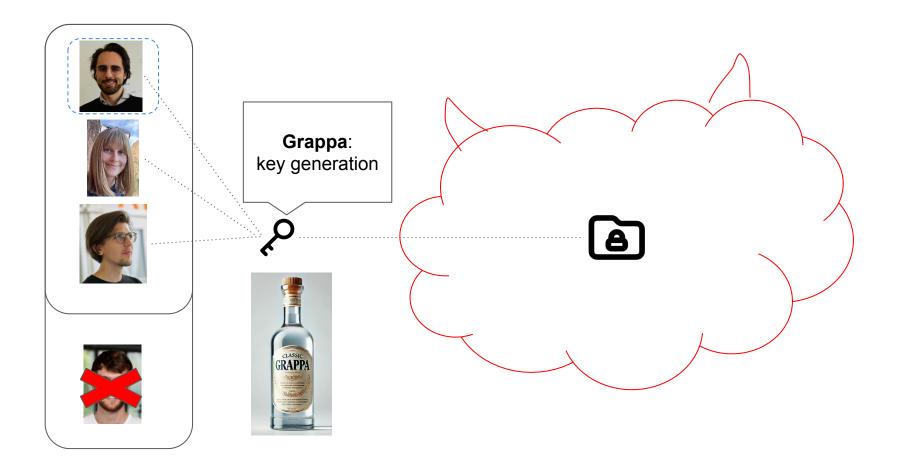
Security (from CGKA) & MLS implementation

X State grows linearly on number of epochs

Can we get a good trade-off?



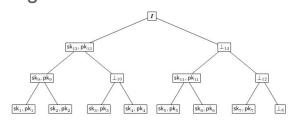




Epoch-based progression of keys for persistent use

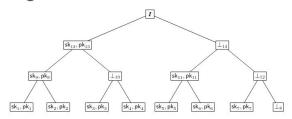
Epoch-based progression of keys for persistent use

CGKA: Continuous Group Key Agreement



Epoch-based progression of keys for persistent use

CGKA: Continuous Group Key Agreement

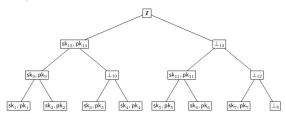


Interval scheme: compact symmetric-key primitive to produce interval keys



Epoch-based progression of keys for persistent use

CGKA: Continuous Group Key Agreement



Interval scheme: compact symmetric-key primitive to produce interval keys



CGKA keys encrypt interval scheme states – CGKA as transport layer

• Grappa: strong security for persistent data in group settings

- Grappa: strong security for persistent data in group settings
- Provable security, compact state

- Grappa: strong security for persistent data in group settings
- Provable security, compact state
- Novel use-case for CGKA beyond messaging

- Grappa: strong security for persistent data in group settings
- Provable security, compact state
- Novel use-case for CGKA beyond messaging

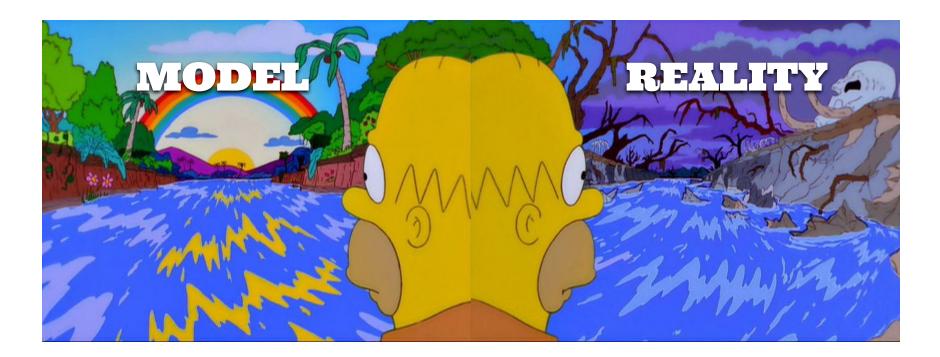
What are the challenges of building a secure shared folder system using CGKA?

- Grappa: strong security for persistent data in group settings
- Provable security, compact state
- Novel use-case for CGKA beyond messaging

What are the **challenges** of building a secure shared folder system using CGKA?



Implementation



Engineering Gaps

Model

1. Abstract client device / capabilities



Reality

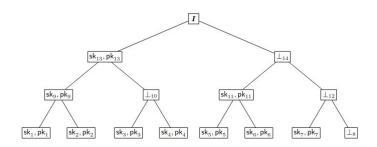
1. Unharmonized capabilities / portability



Engineering Gaps

Model

- 1. Abstract client device / capabilities
- 2. Crypto primitives as mathematical objects



Reality

- 1. Unharmonized capabilities / portability
- 2. Crypto primitives support in the execution platform / libraries

List of existing MLS implementations

- MLSpp (C++) <u>https://github.com/cisco/mlspp</u> (Status: RFC)
- OpenMLS (Rust) https://github.com/openmls/openmls (Status: RFC)
- mls-kotlin (Kotlin) https://github.com/Traderjoe95/mls-kotlin (Status: RFC)
- mls-rs (Rust) [<u>https://github.com/awslabs/mls-rs</u>] (Status: RFC)
- <u>RingCentral</u> proprietary implementation (C++) (Status: draft-11; RFC in progress)
- MLS* (F*) (Status: RFC in progress)
- BouncyCastle (Java) https://github.com/bcgit/bc-java (Status: RFC)
- go-mls (Go) (Status: RFC in progress)

Why crypto in Browsers?

• Browser: cross platform runtime to access applications in cloud





Why crypto in Browsers?

- Browser: cross platform runtime to access applications in cloud
- Standardised Web Crypto API (W3C) for JS Runtimes



Web Cryptography API W3C Recommendation 26 January 2017

This Version: https://www.w3.ora/TR/2017/REC-WebCrvptoAPI-20170126/ Latest Published Version: https://www.w3.org/TR/WebCryptoAPI/ Latest editor's draft: https://w3c.github.io/webcrypto/Overview.html Previous Version: https://www.w3.org/TR/2016/PR-WebCryptoAPI-20161215/ Editor: Mark Watson, Netflix <watsonm@netflix.com> Errata for this document will be gathered from issues. See also translations. Participate: We are on GitHub. Send feedback to public-web-security@w3.org (archives). File a bug (see existing bugs) Copyright © 2012-2017 W3C® (MIT, ERCIM, Keio, Beihang). W3C liability, trademark and document use rules apply.



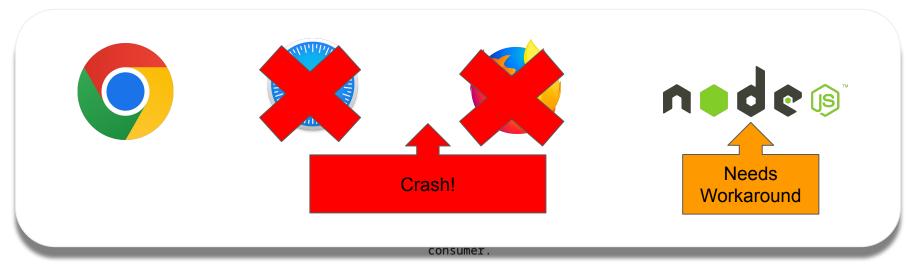
• Exhibits non-standard behaviours (loose specifications)

MLS implementation relies on public keys computed from private keys

o publicKey contains the elliptic curve public key associated with the private key in question. The format of the public key is specified in <u>Section 2.2 of [RFC5480]</u>. Though the ASN.1 indicates publicKey is OPTIONAL, implementations that conform to this document SHOULD always include the publicKey field. The publicKey field can be omitted when the public key has been distributed via another mechanism, which is beyond the scope of this document. Given the private key and the parameters, the public key can always be recomputed; this field exists as a convenience to the consumer.

• Exhibits non-standard behaviours (loose specifications)

MLS implementation relies on public keys computed from private keys



- Exhibits non-standard behaviours (loose specifications)
- Introduction of new primitives takes very long!



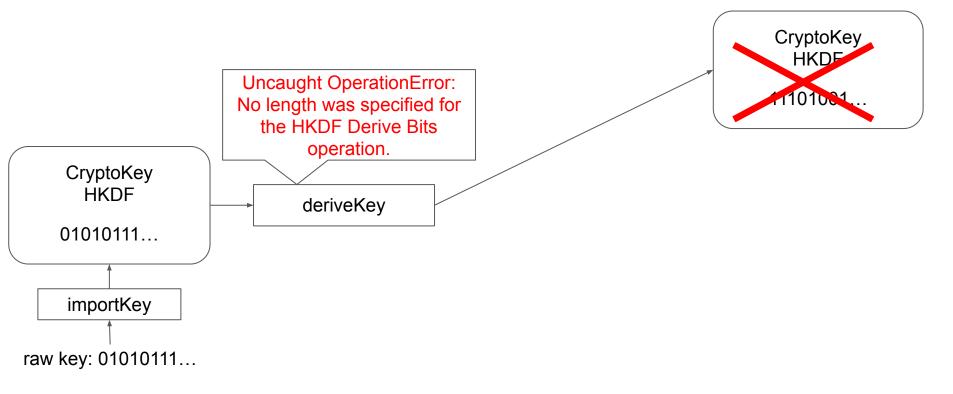
This is an unofficial proposal.

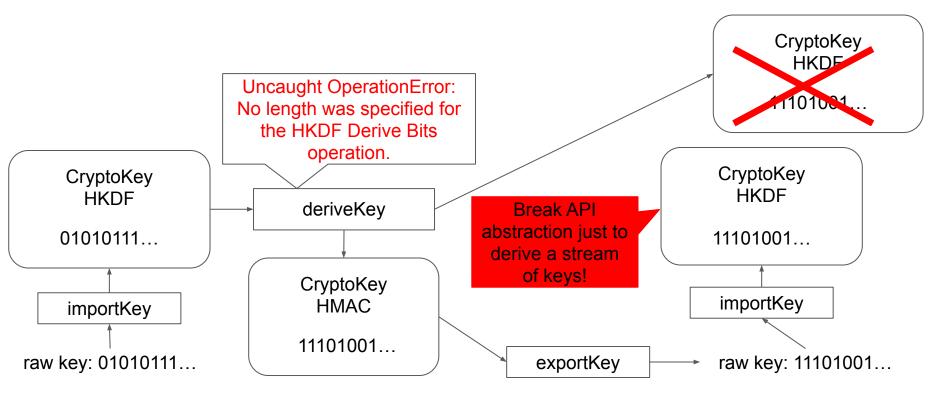
Groups.

> 8y, not a standard yet!!

.iptic curves over prime fields that offer security in cryptographic applications, including Transport Layer Security (TLS). These curves are intended to operate at the ~128-bit and ~224-bit security level, respectively, and are generated deterministically based on a list of required properties.

- Exhibits non-standard behaviours (loose specifications)
- Introduction of new primitives takes very long!
- Overprotective, too restrictive to implement advanced crypto





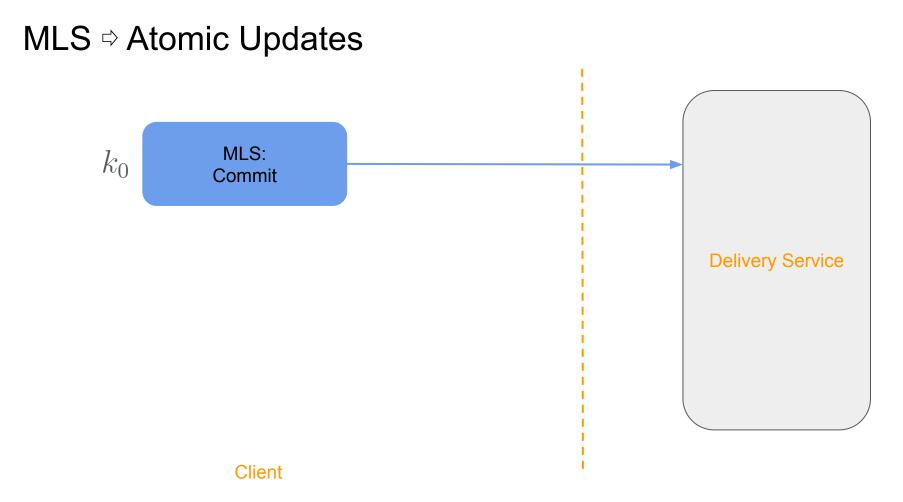
Engineering Gaps

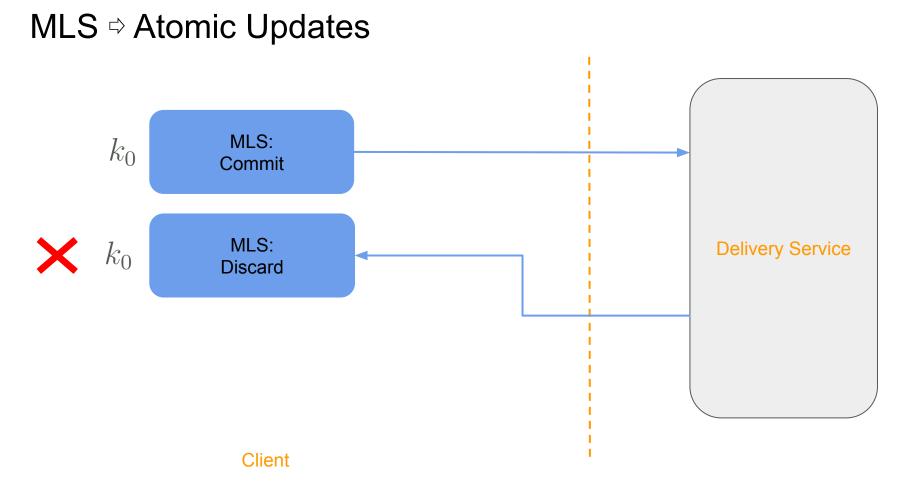
Model

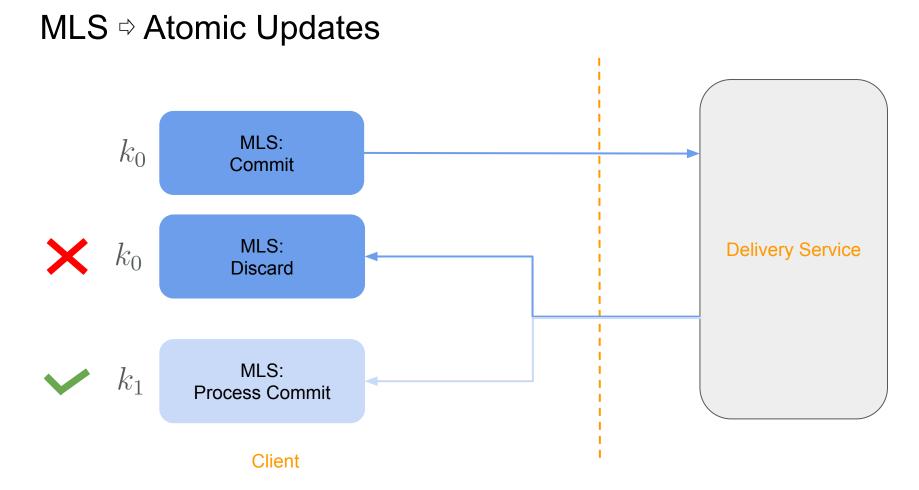
- 1. Abstract client device / capabilities
- 2. Crypto primitives as mathematical objects
- 3. Atomic operation of the scheme

Reality

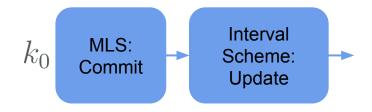
- 1. Unharmonized capabilities / portability
- 2. Crypto primitives support in the execution platform / libraries
- 3. Multiple schemes, non atomic interactions between components

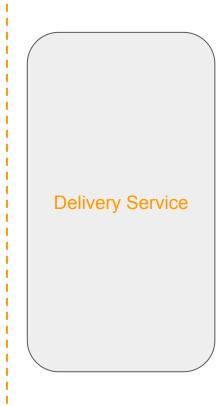




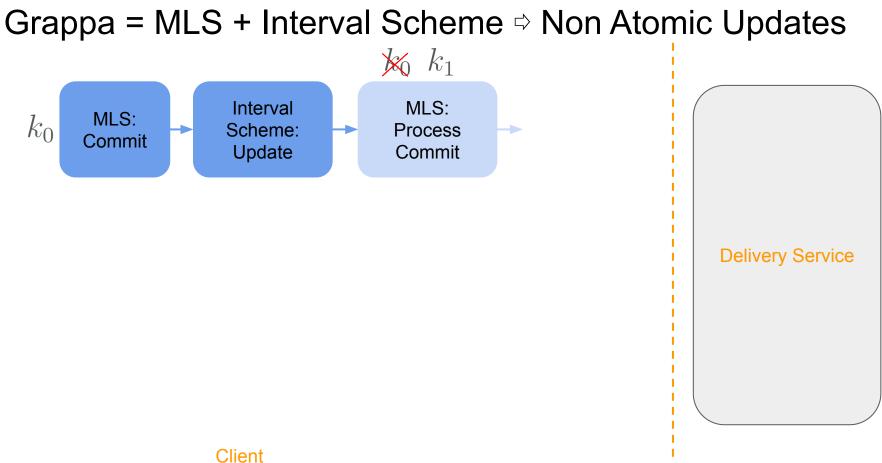


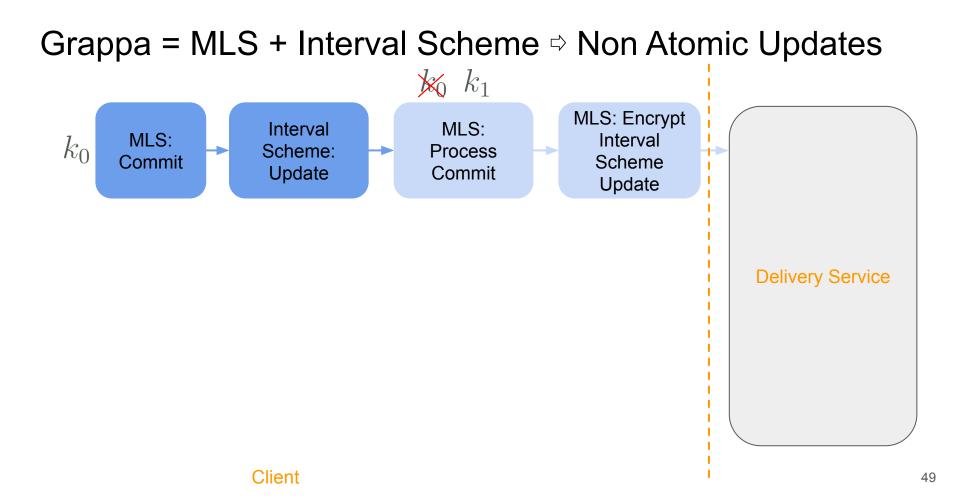
Grappa = MLS + Interval Scheme ⇒ Non Atomic Updates

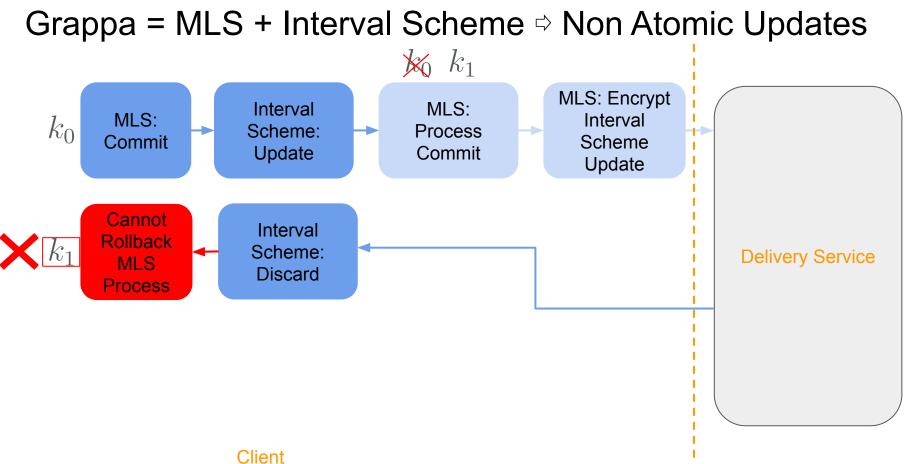


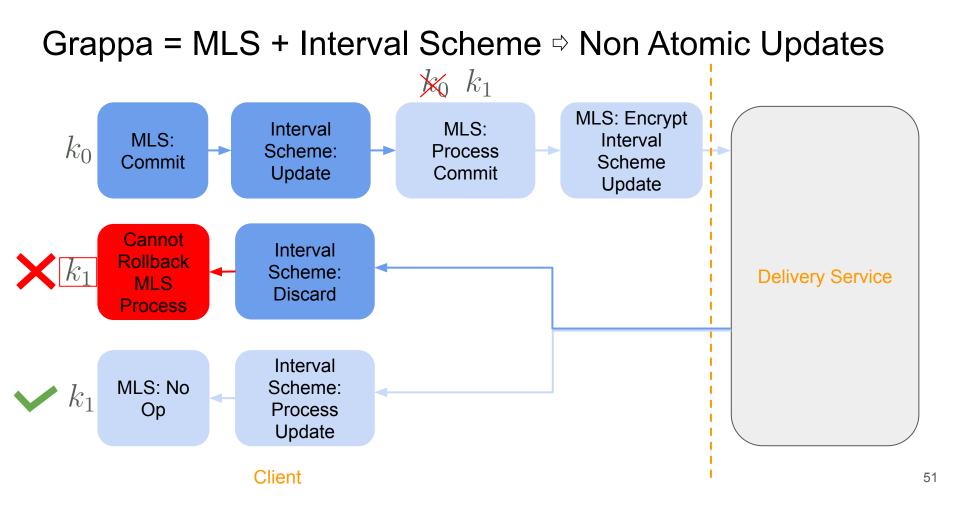


Client

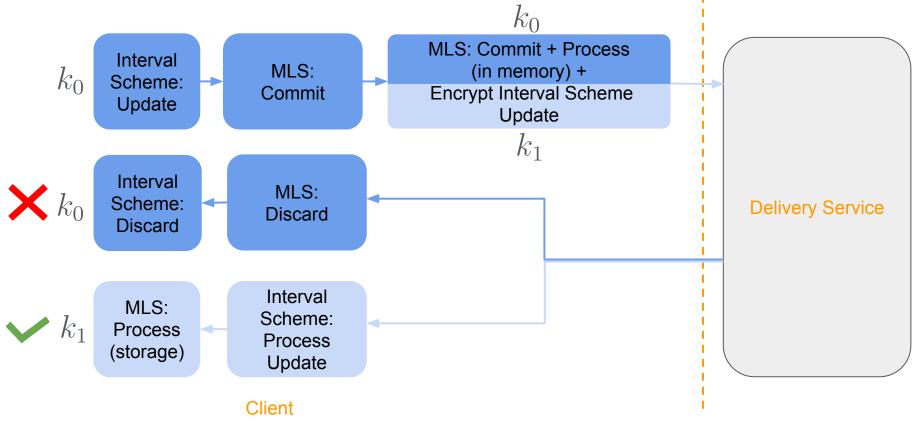






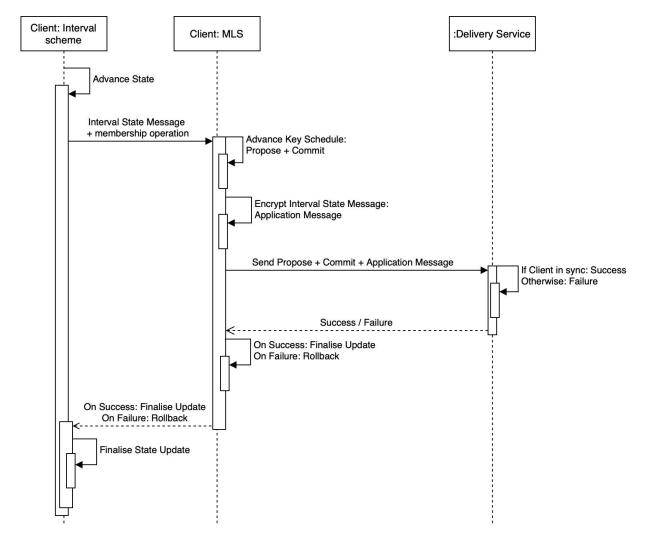


MLS as a Transport Layer Grappa Atomic Updates



Takeaways

- We introduce **Grappa** for dynamic groups of users to agree on a key progression for **persistent data**:
 - Builds on CGKA and Interval Scheme
 - Advanced security guarantees
 - Secure Shared Folders as a real-world application
- We implement Grappa for browser and desktop:
 - Relied on existing MLS implementation and WebCrypto API standard
 - Cryptography in JS environments is lagging behind
- MLS as transport layer to construct new protocols:
 - Gap between security modeling and functionality



56		
57	+	<pre>#[cfg(not(feature = "node"))]</pre>
58	+	async in import_with_public_info(&self, crypto: &SubtleCrypto, key: &Uint8Array, params: &EcKeyImportParams, key_usages: &Array) -> Result <promise, jsvalue=""> {</promise,>
59	+	crypto.import_key_with_object(self.format(), &key, ¶ms, true, &key_usages)
60	+	}
61	+	
62	+	<pre>#[cfg(feature = "node")]</pre>
63	+	async in import_with_public_info(&self, crypto: &SubtleCrypto, bytes: &Uint8Array, params: &EcKeyImportParams, key_usages: &Array) -> Result <promise, jsvalue=""> {</promise,>
64	+	<pre>let crypto_key_promise = crypto.import_key_with_object(self.format(), bytes, params, true, key_usages)?;</pre>
65	+	let crypto key - lcEuture: from(crypto key promise) await? into():
66	+	// Export the key to jwk to force the generation of the public key.
67	+	<pre>let jwk_promise = crypto.export_key(&"jwk", &crypto_key)?;</pre>
68	+	<pre>let iwk = JsFuture::from(iwk promise).await?:</pre>
69	+	// Re-import into the original requested format with the same usages from jwk.
70	+	crypto.import_key_with_object('jwk', @jwk.into(), params, true, @key_usages)
71	+	}
72	+	
ISSN- core for		

