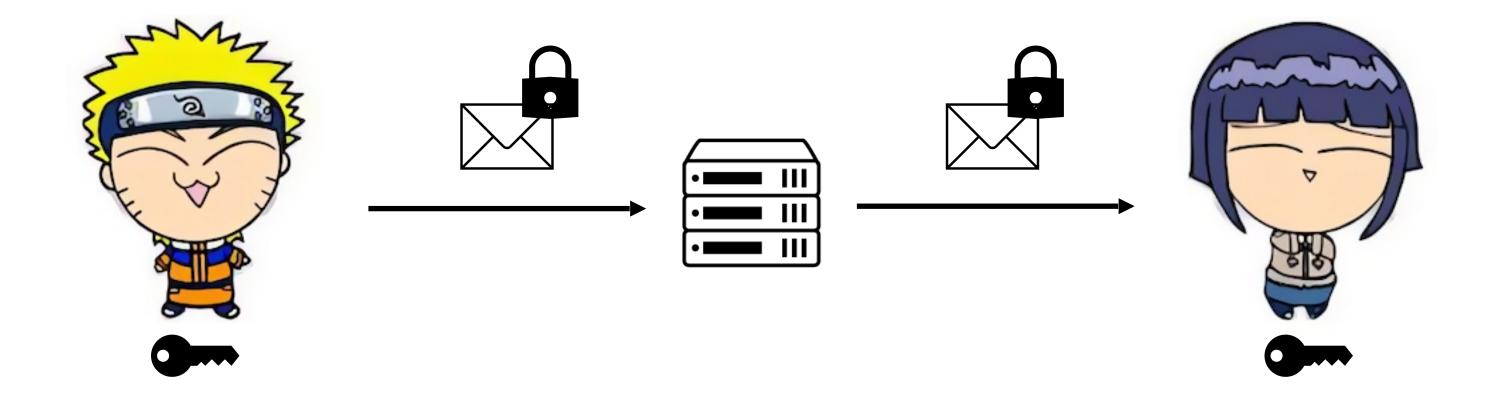
Analyzing Chat Encryption in Group Messaging Applications

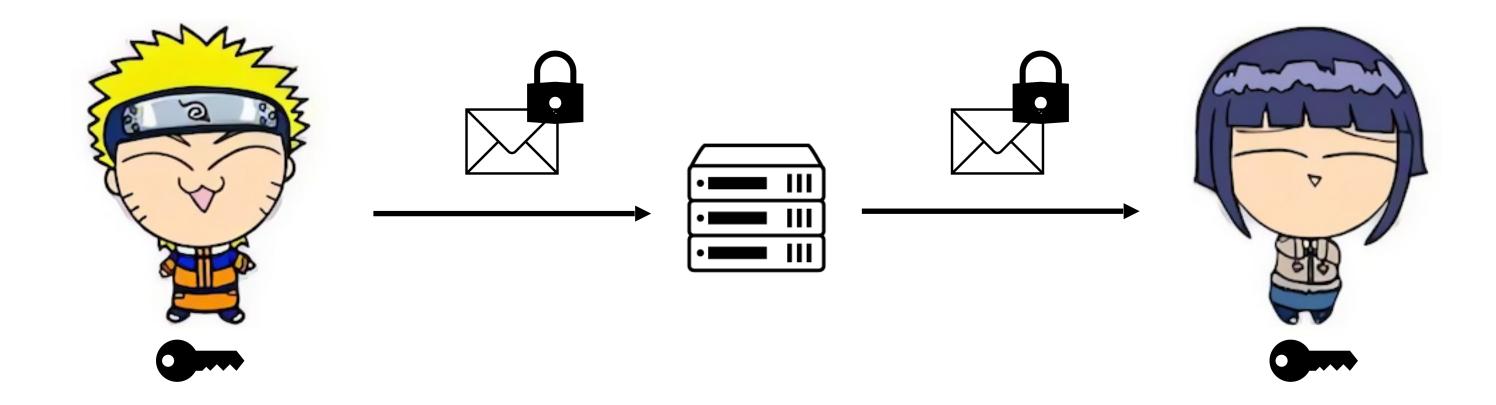
Joseph Jaeger*, <u>Akshaya Kumar</u>*, and Igors Stepanovs



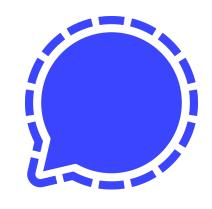
E2EE/Secure Messaging



E2EE/Secure Messaging







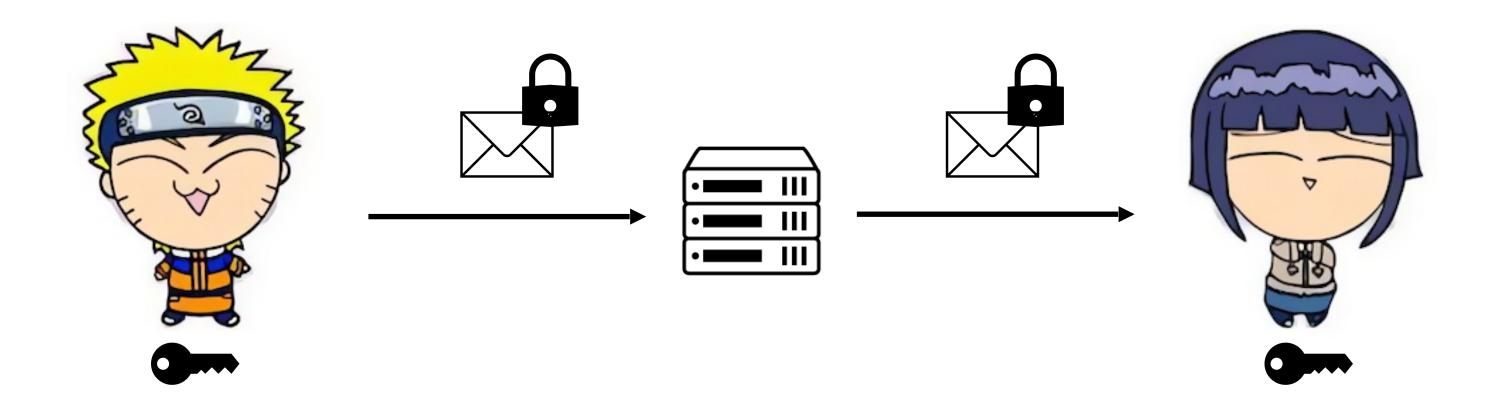




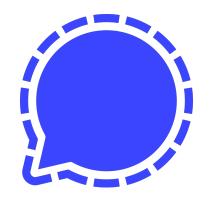




E2EE/Secure Messaging







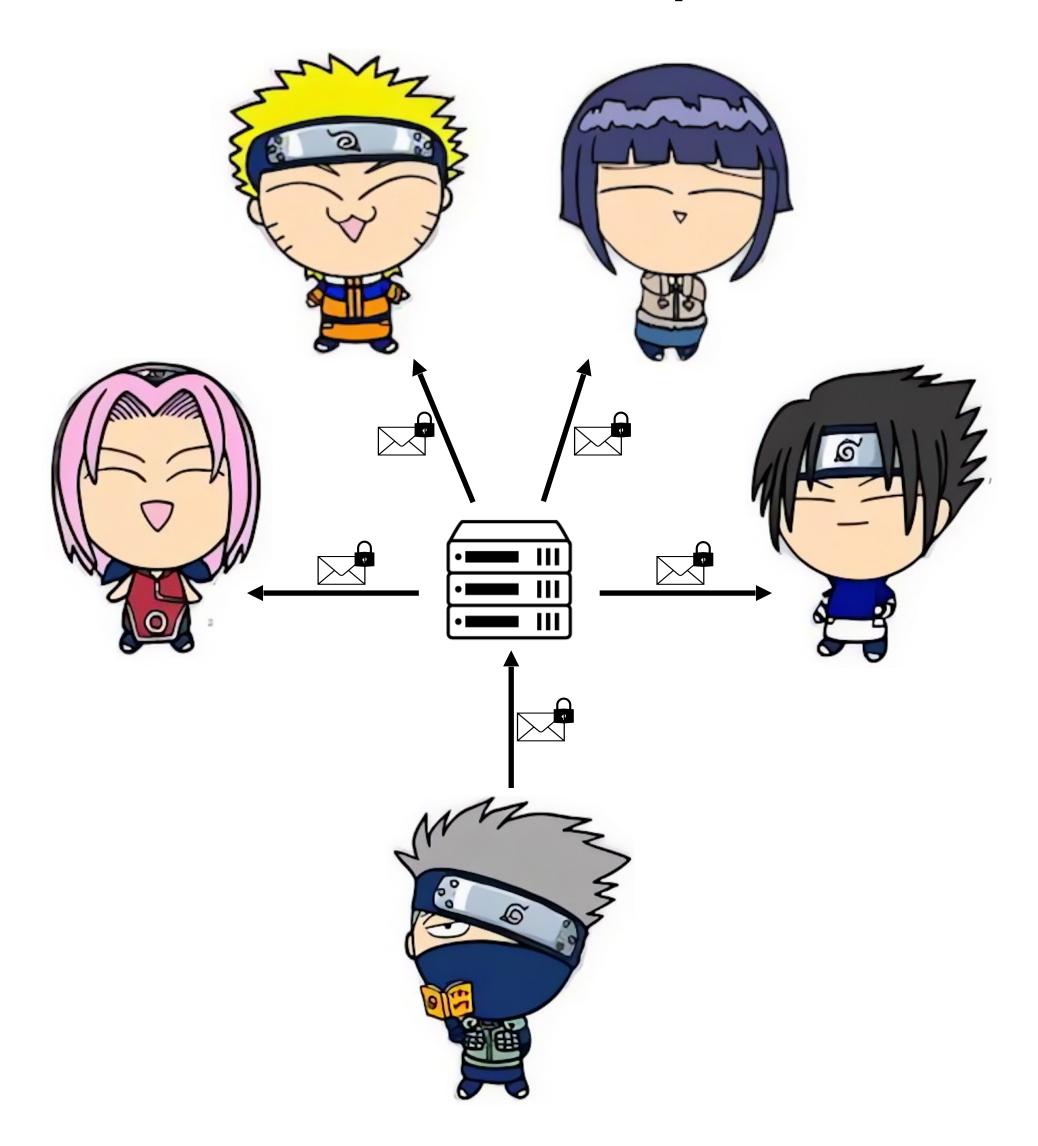


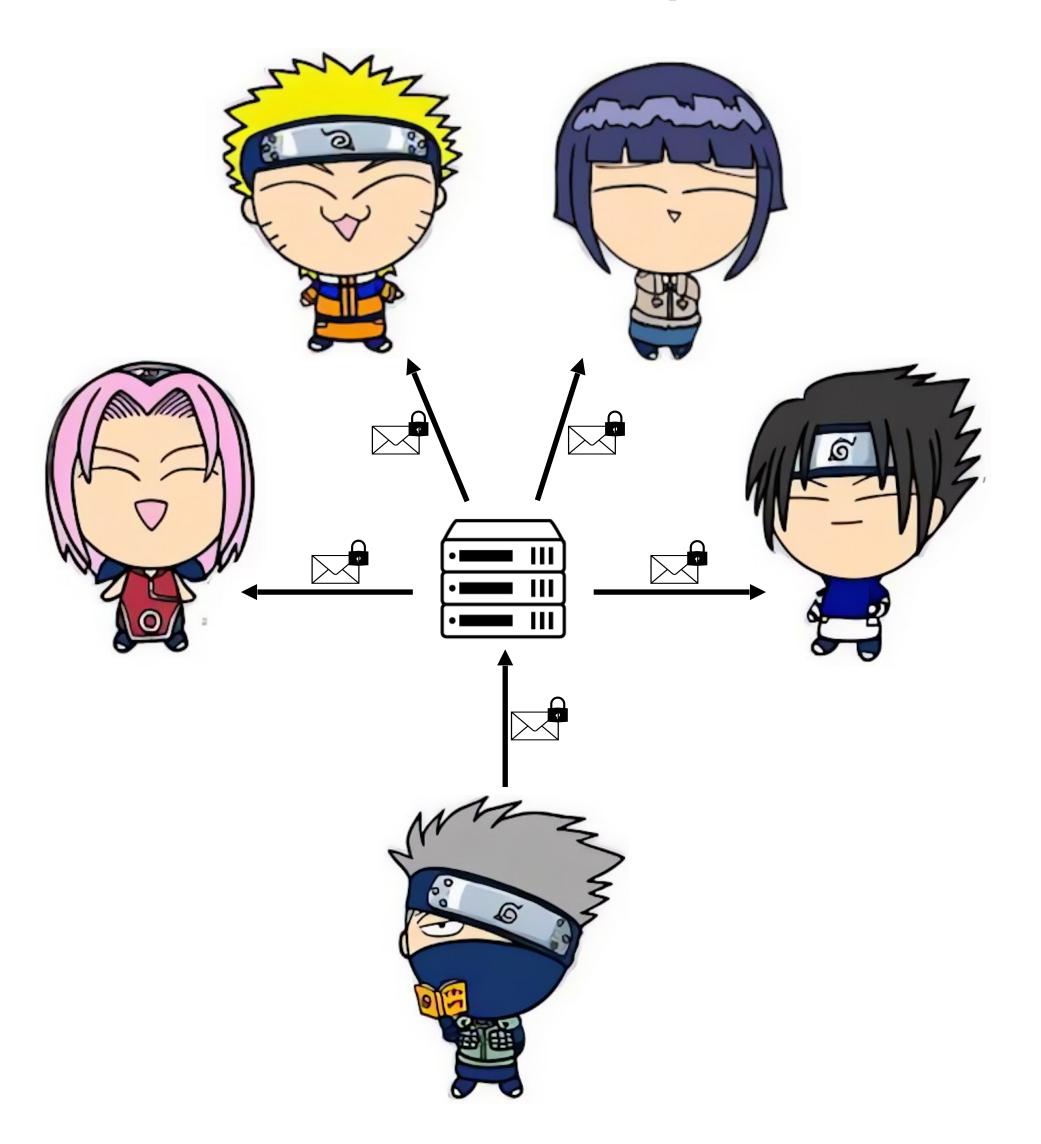


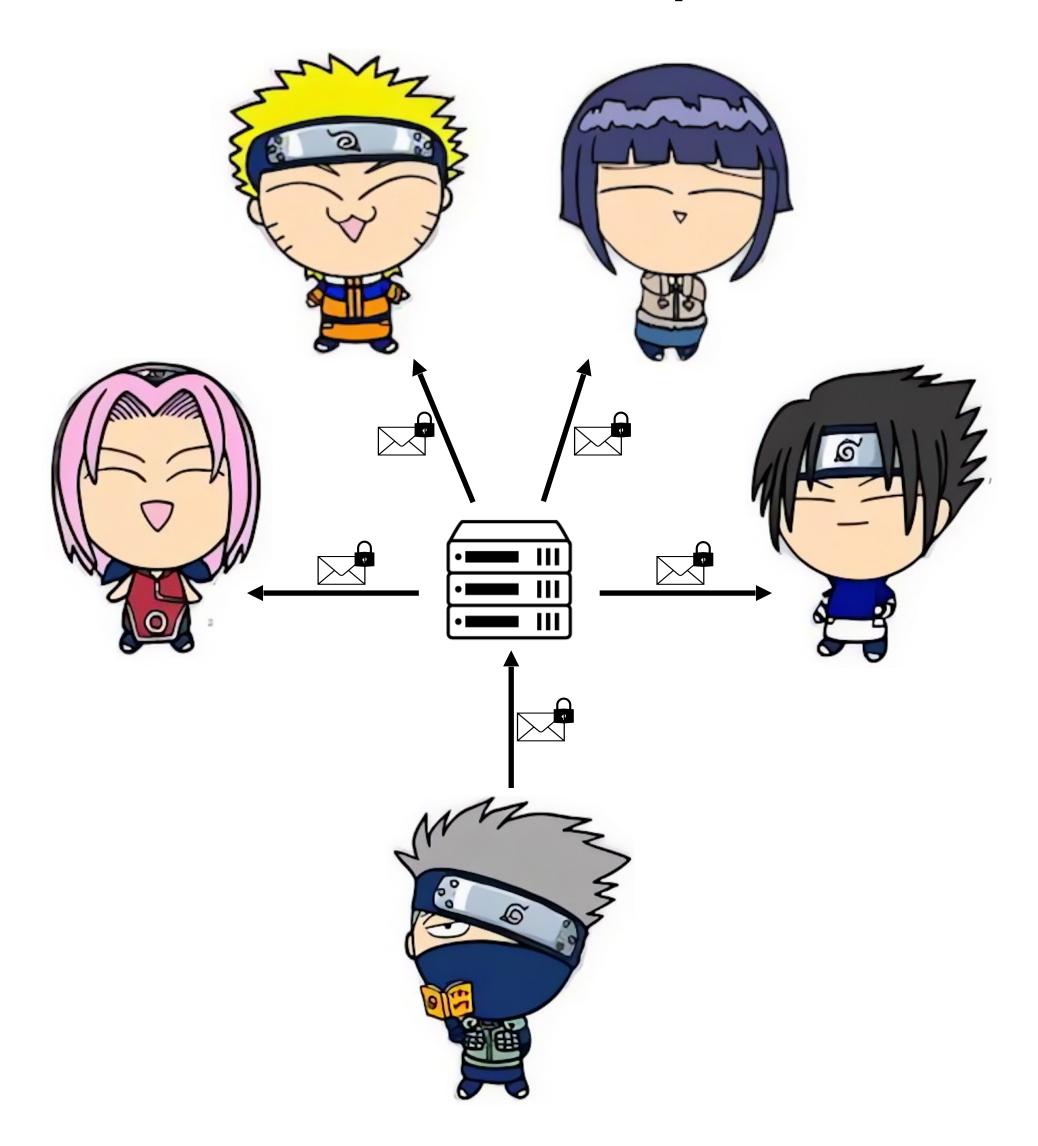


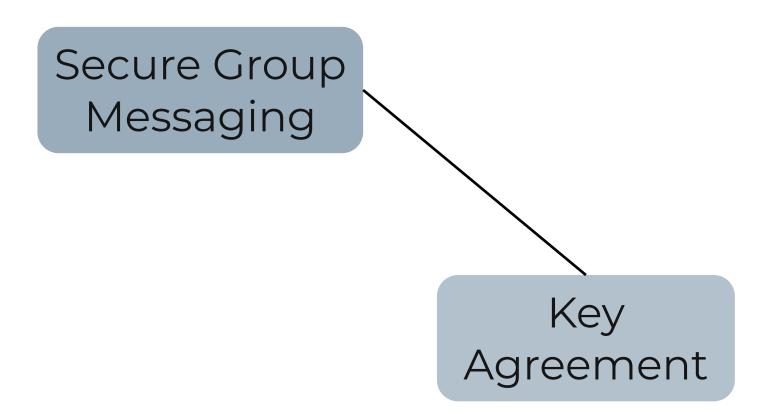


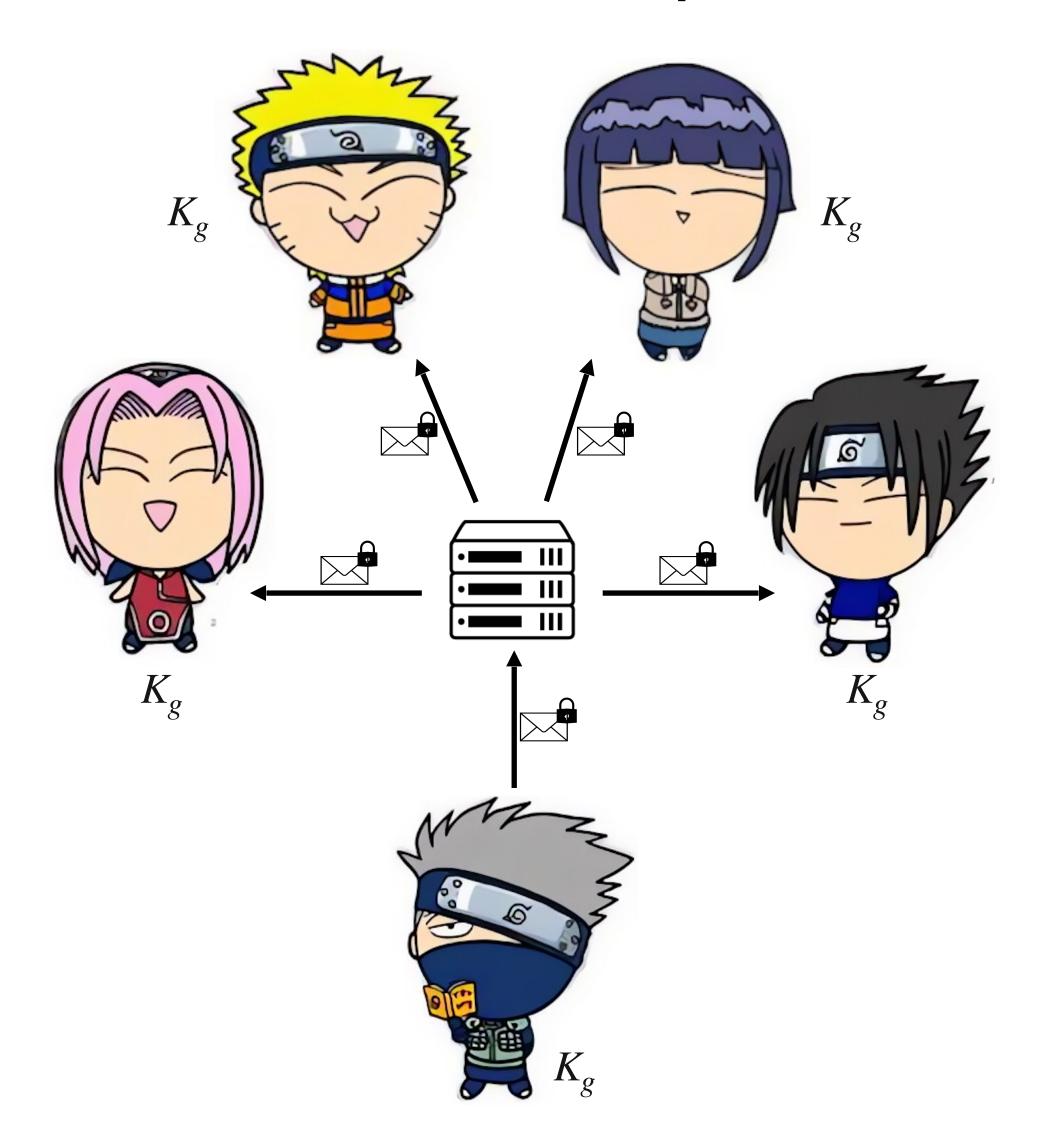


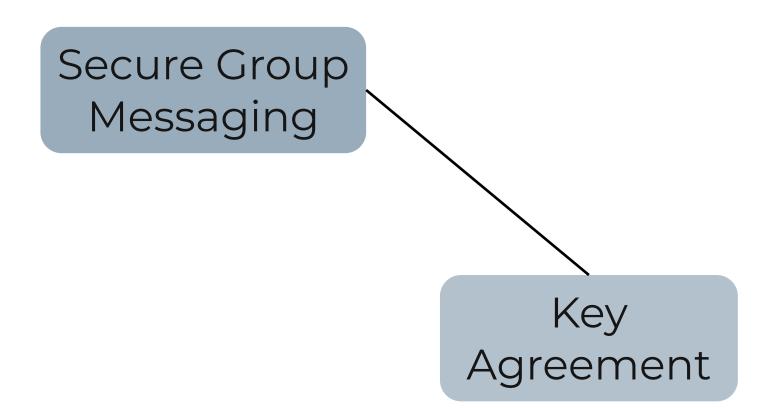


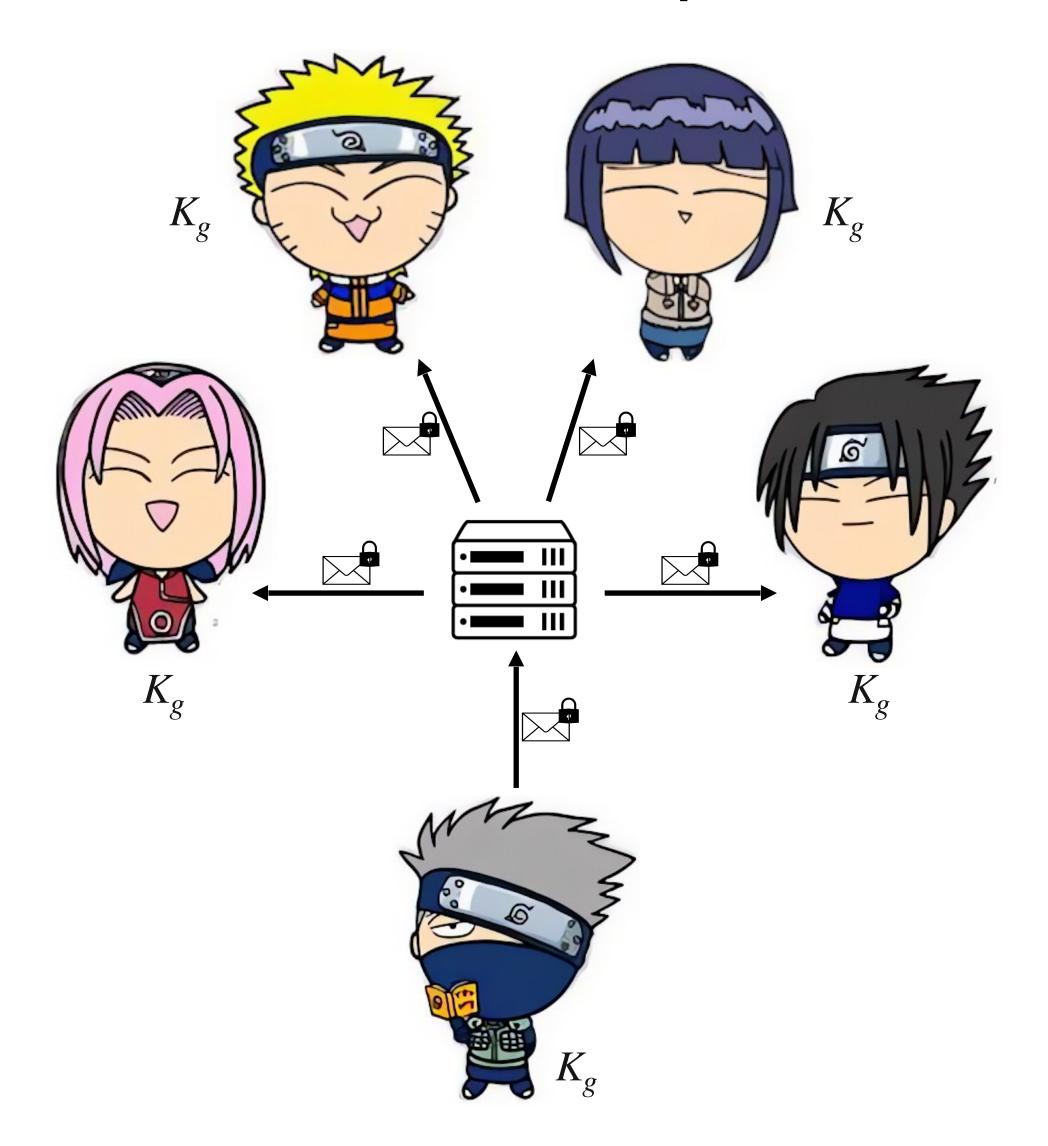


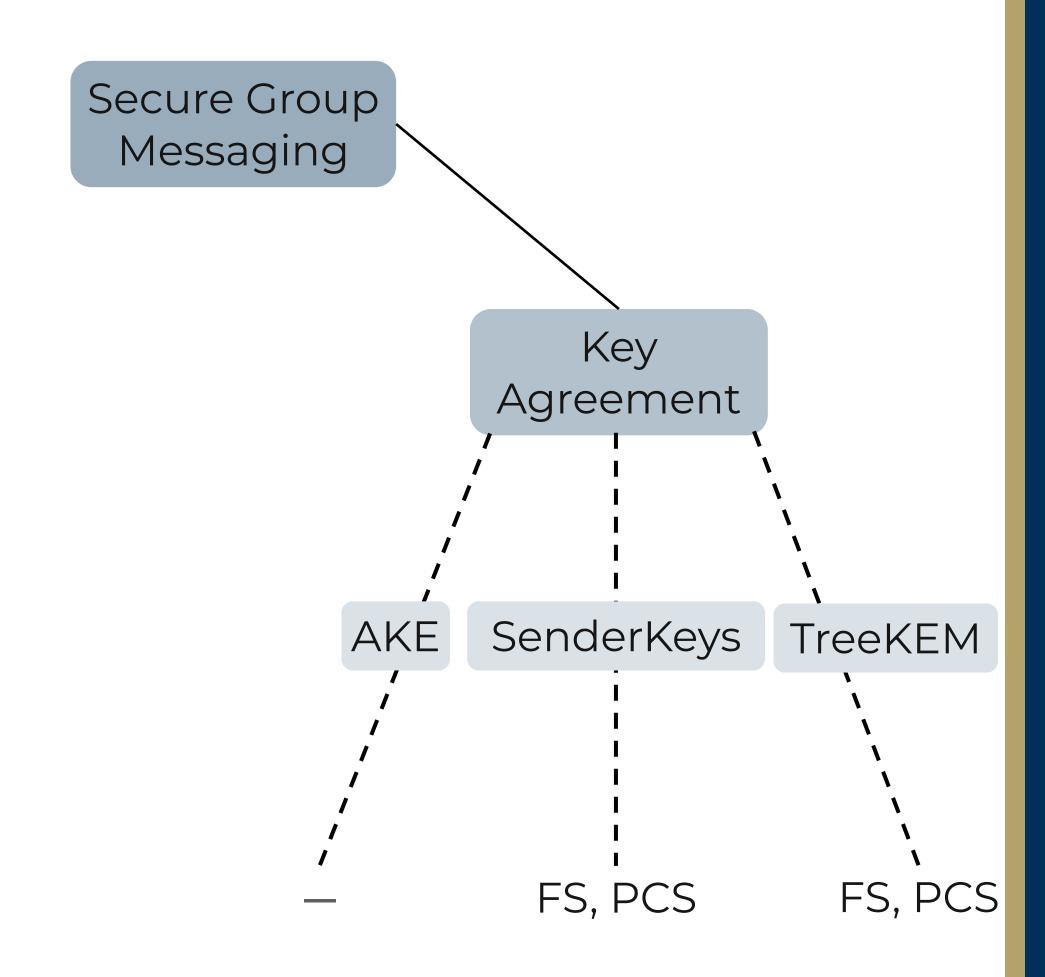


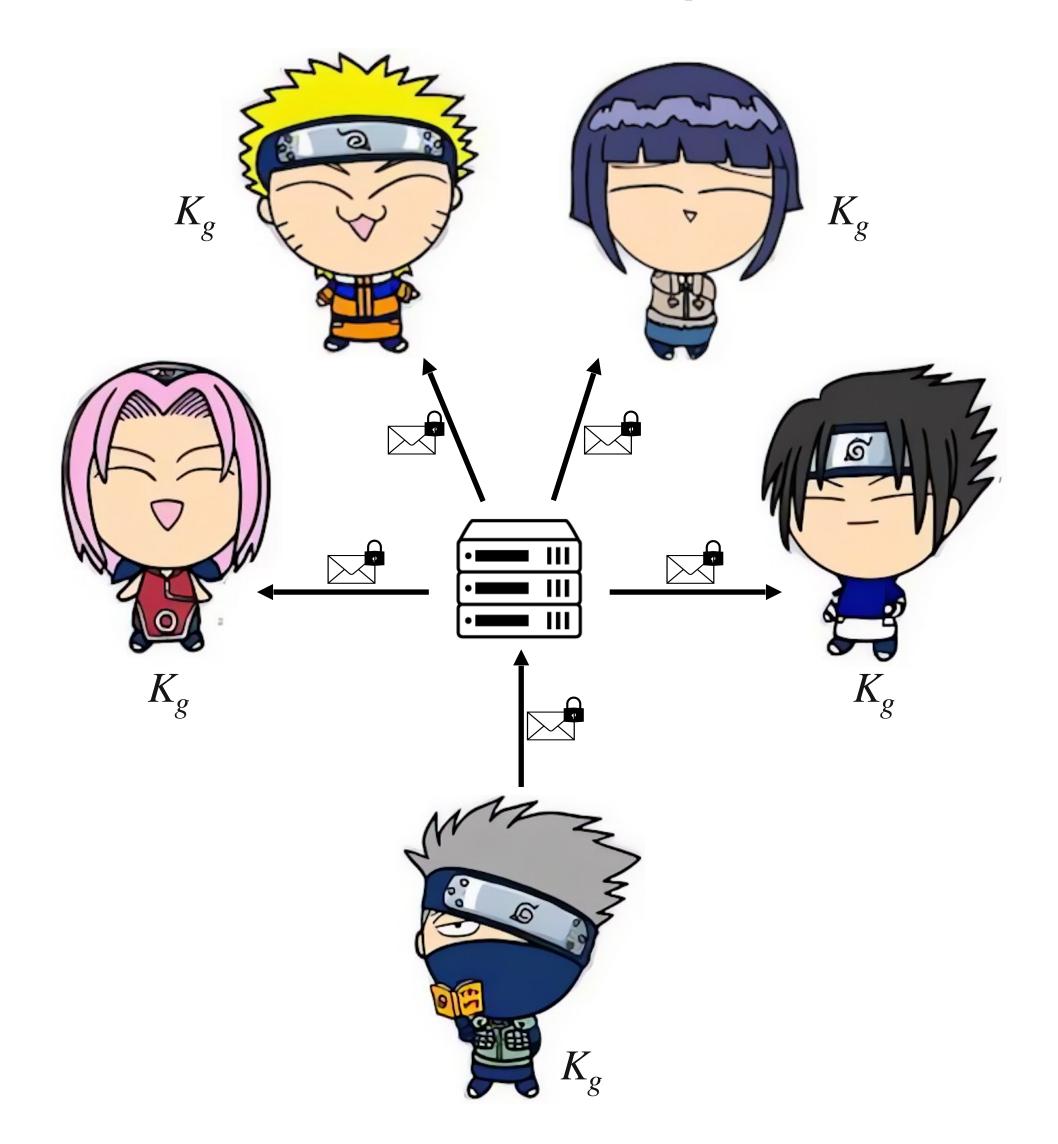


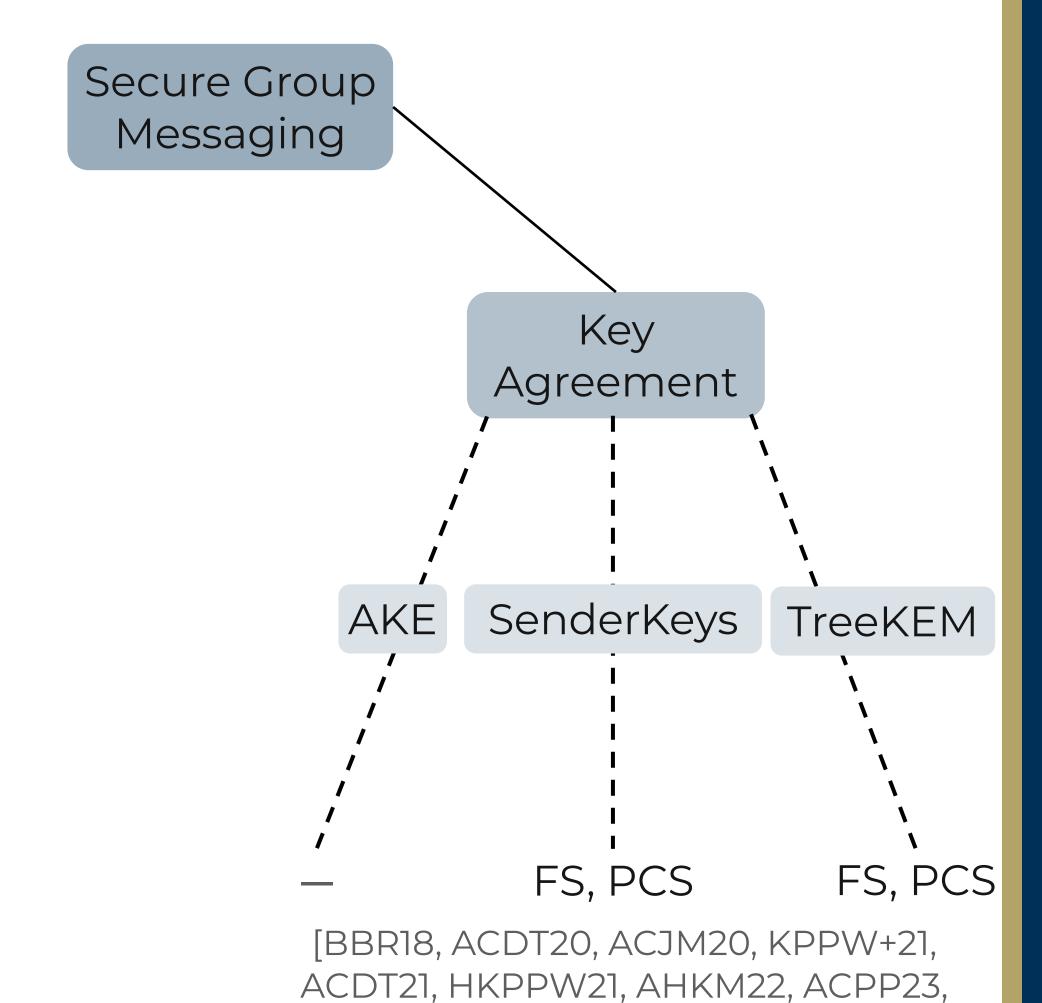




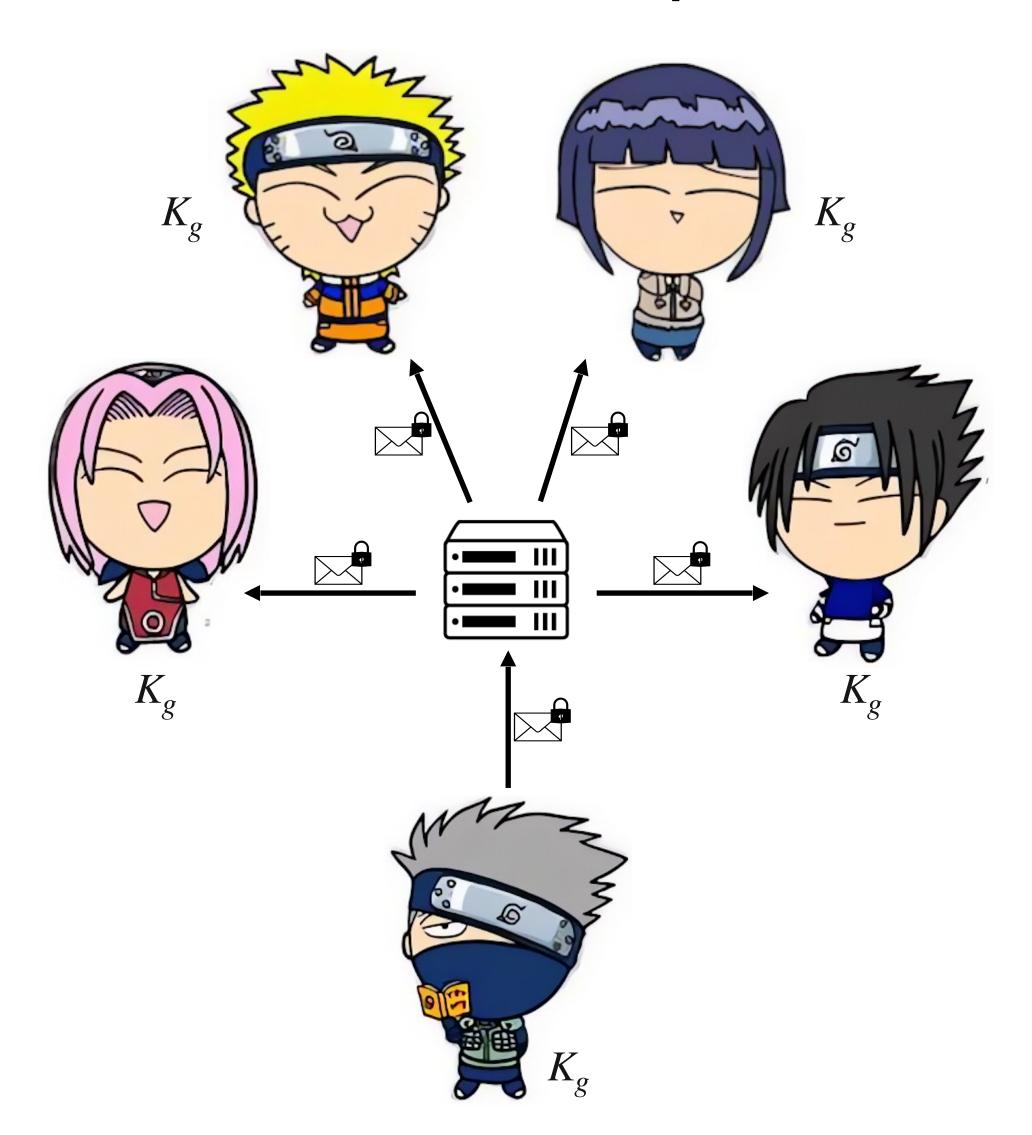






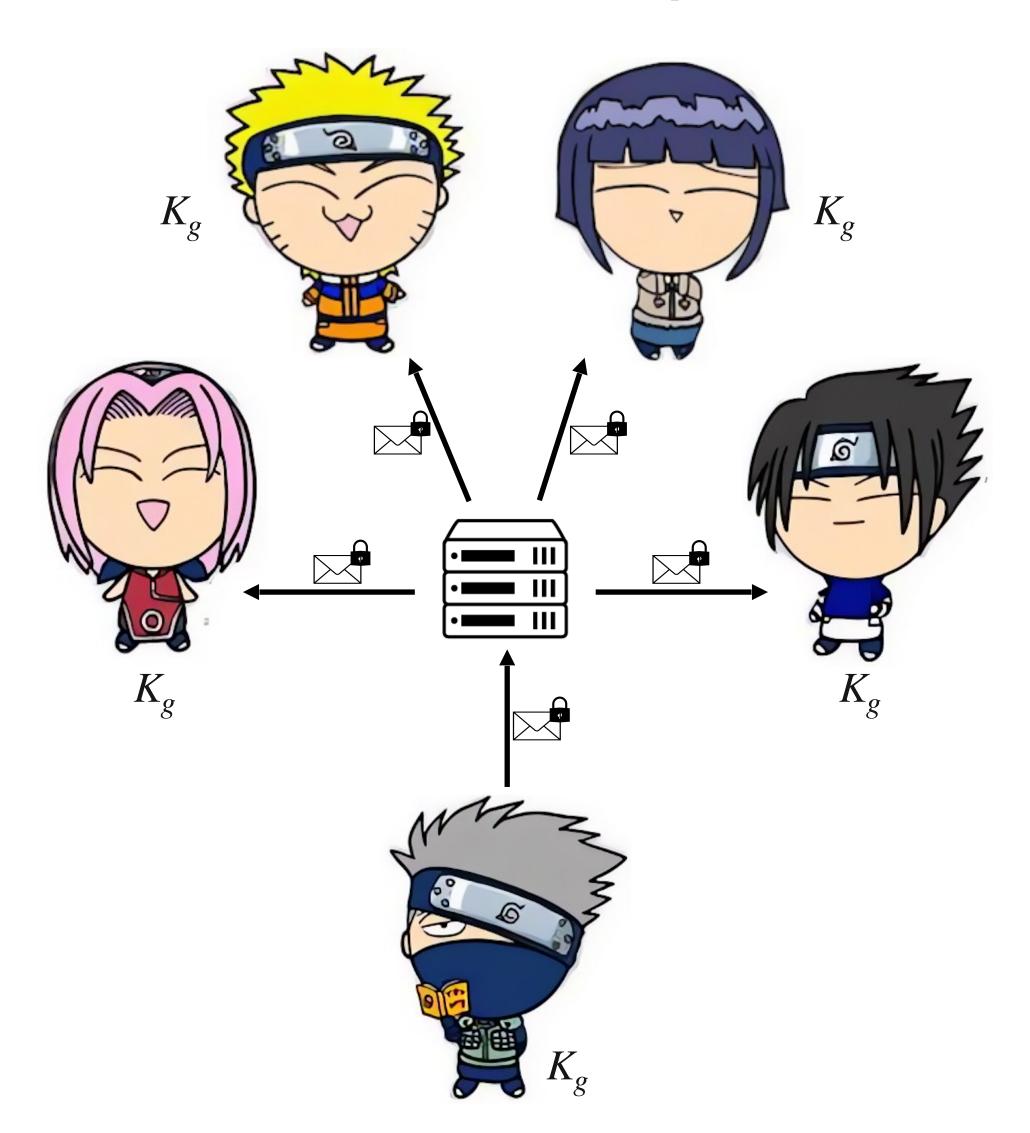


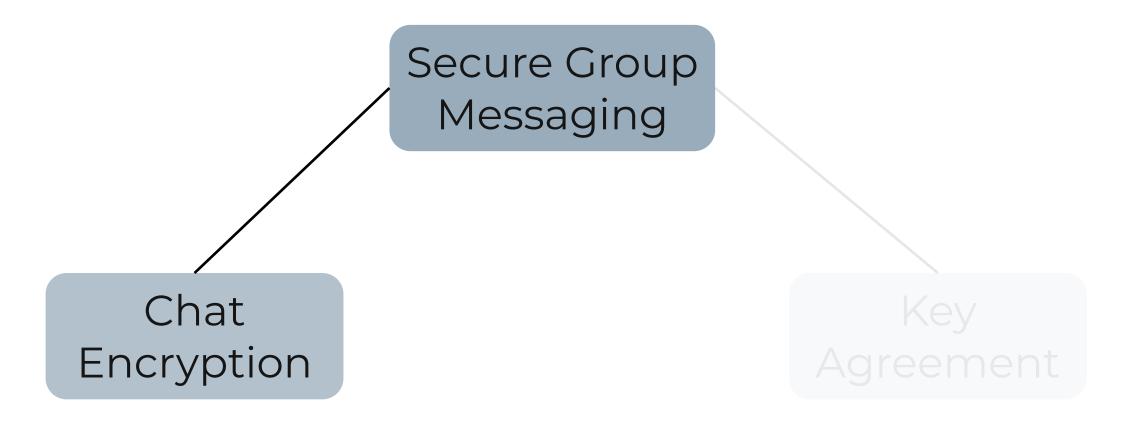
AWT23, BCG23, CLM23, ...]

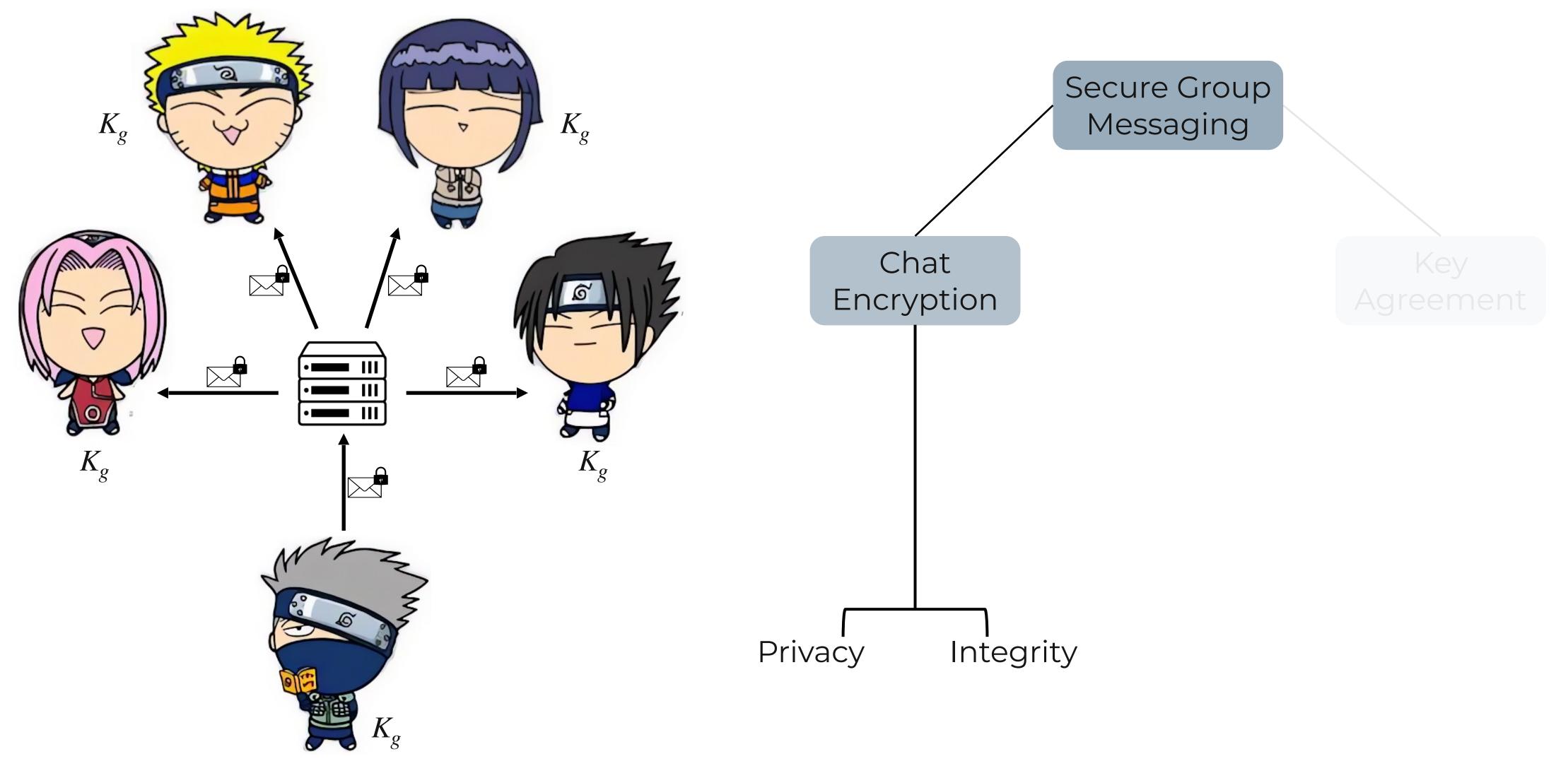


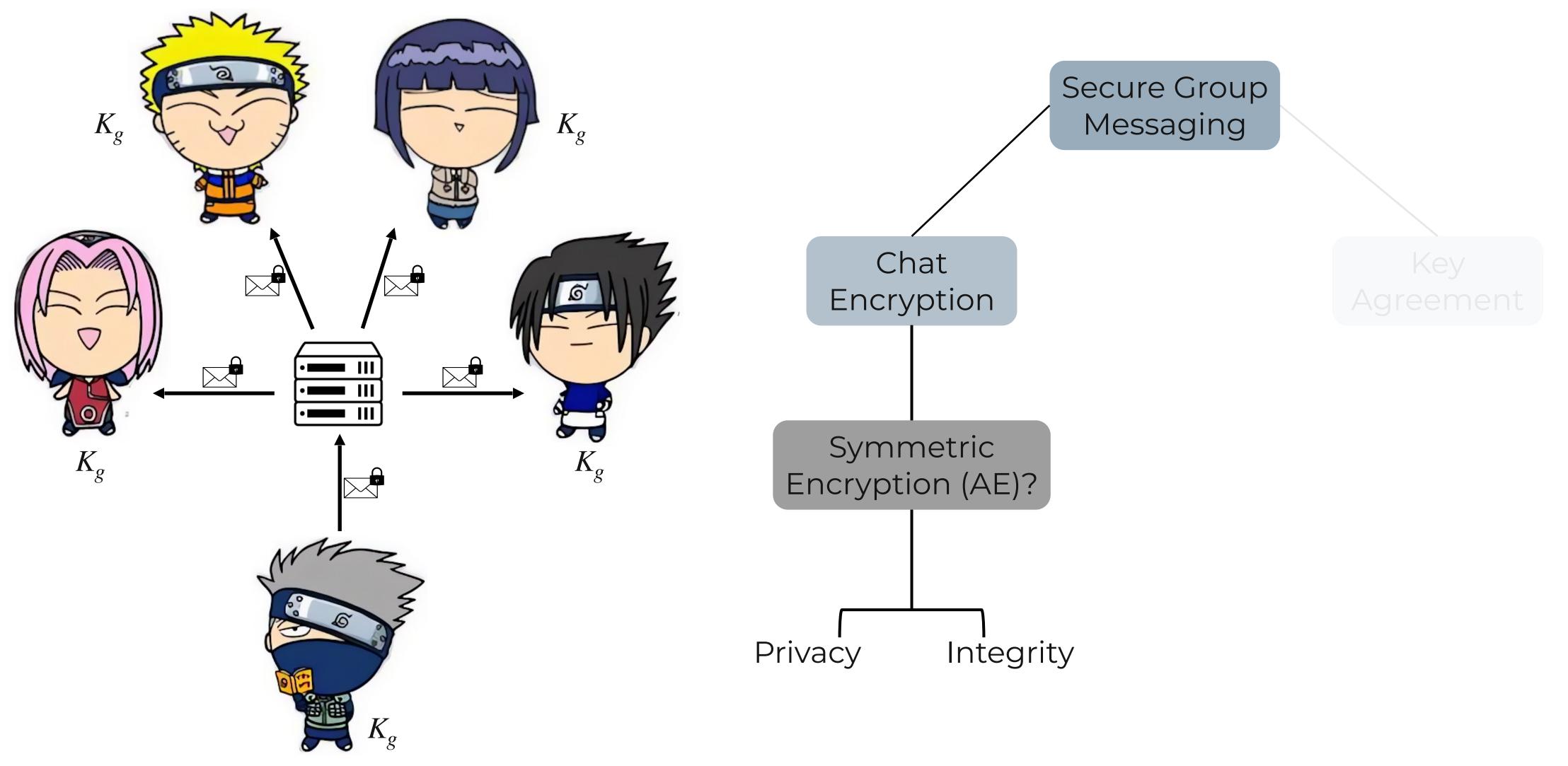
Secure Group Messaging

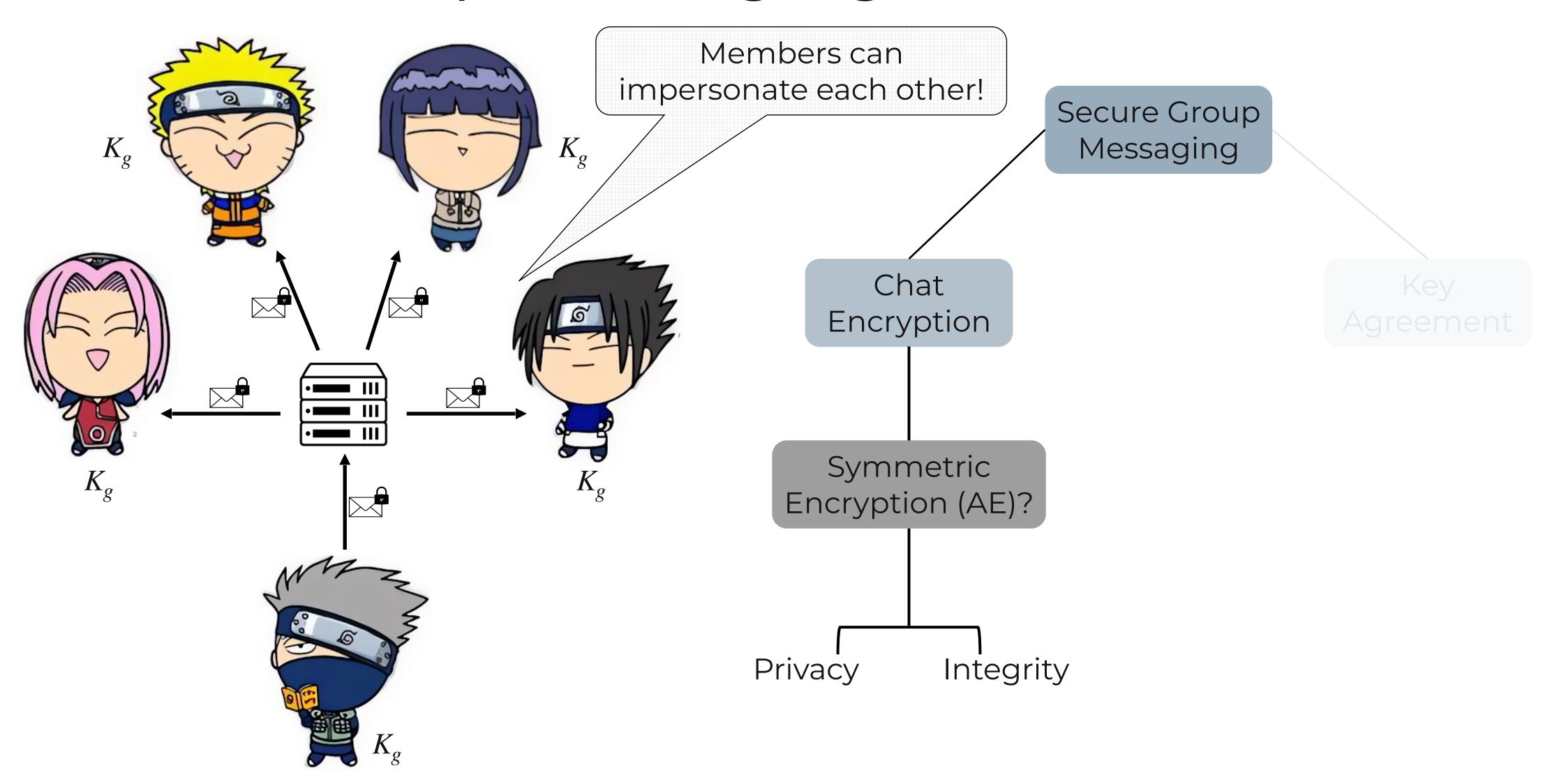
> Key Agreement

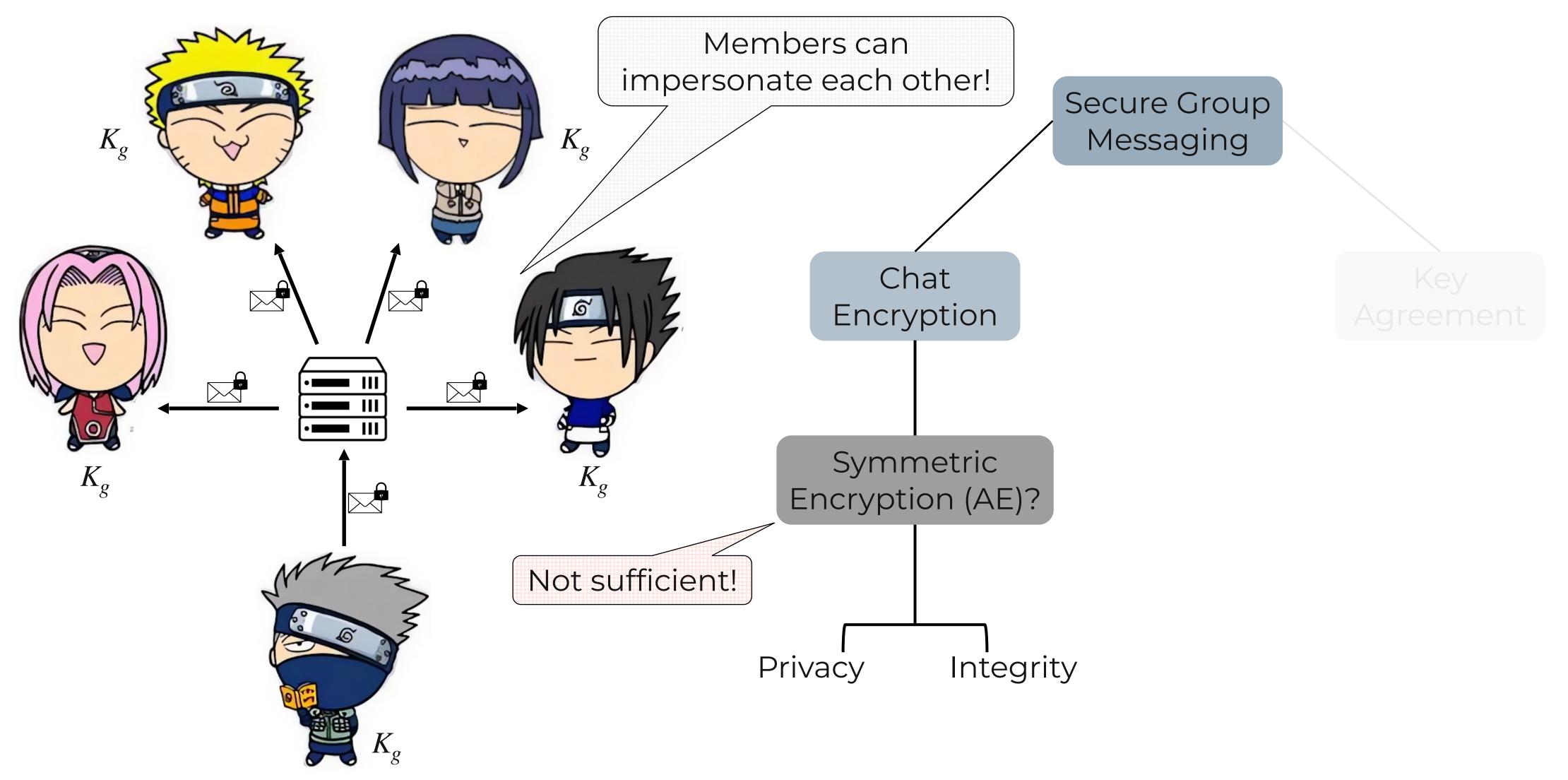


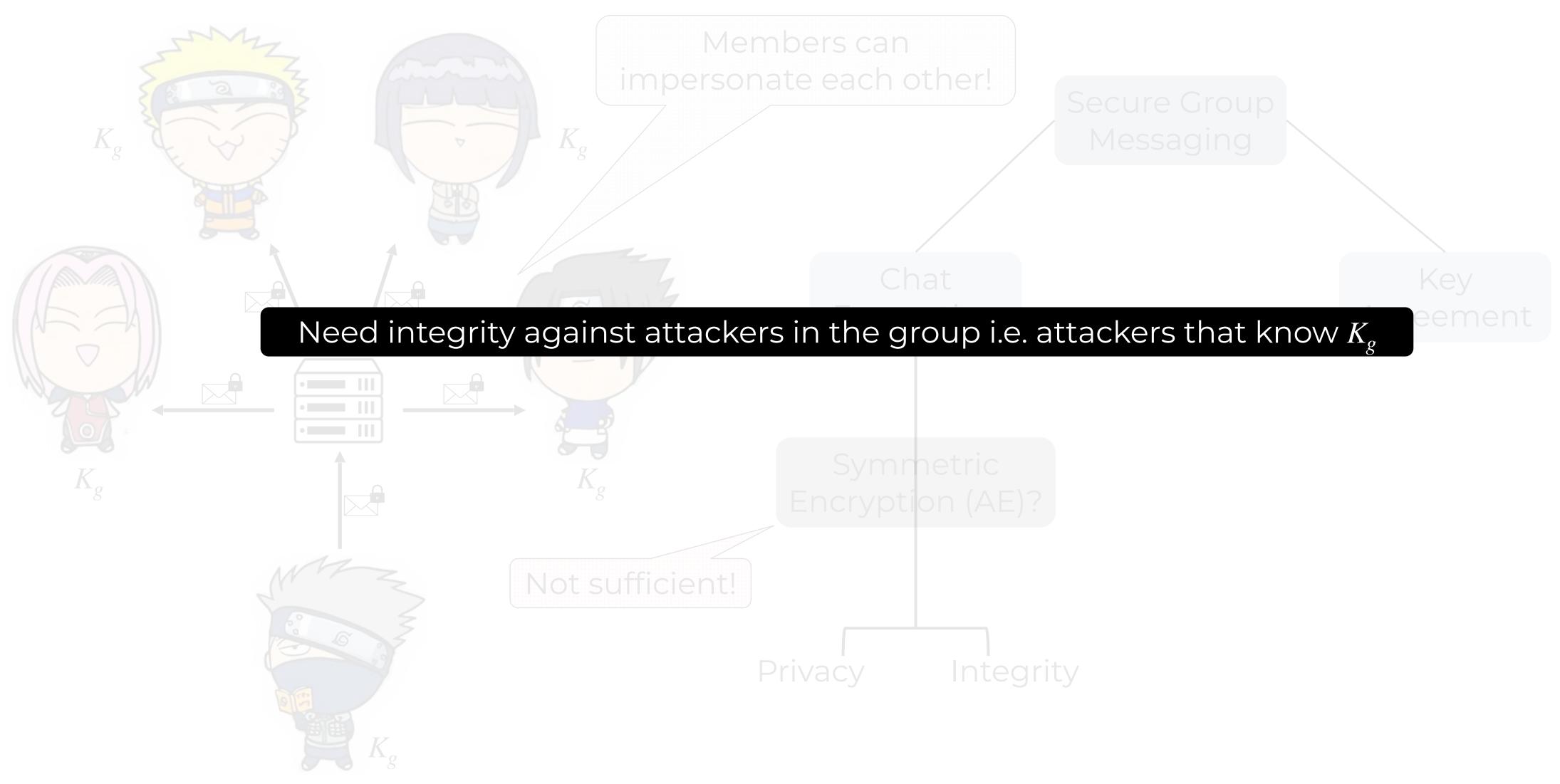


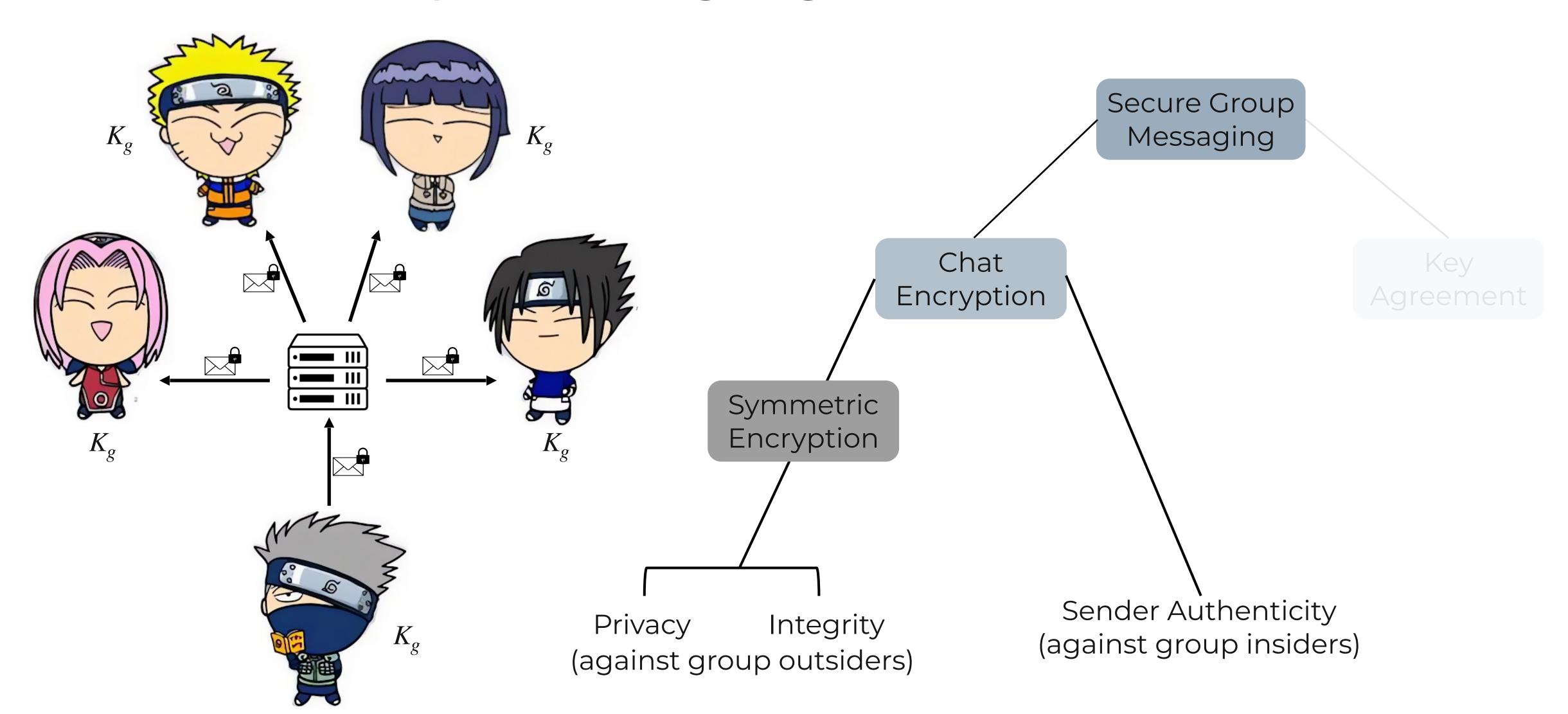


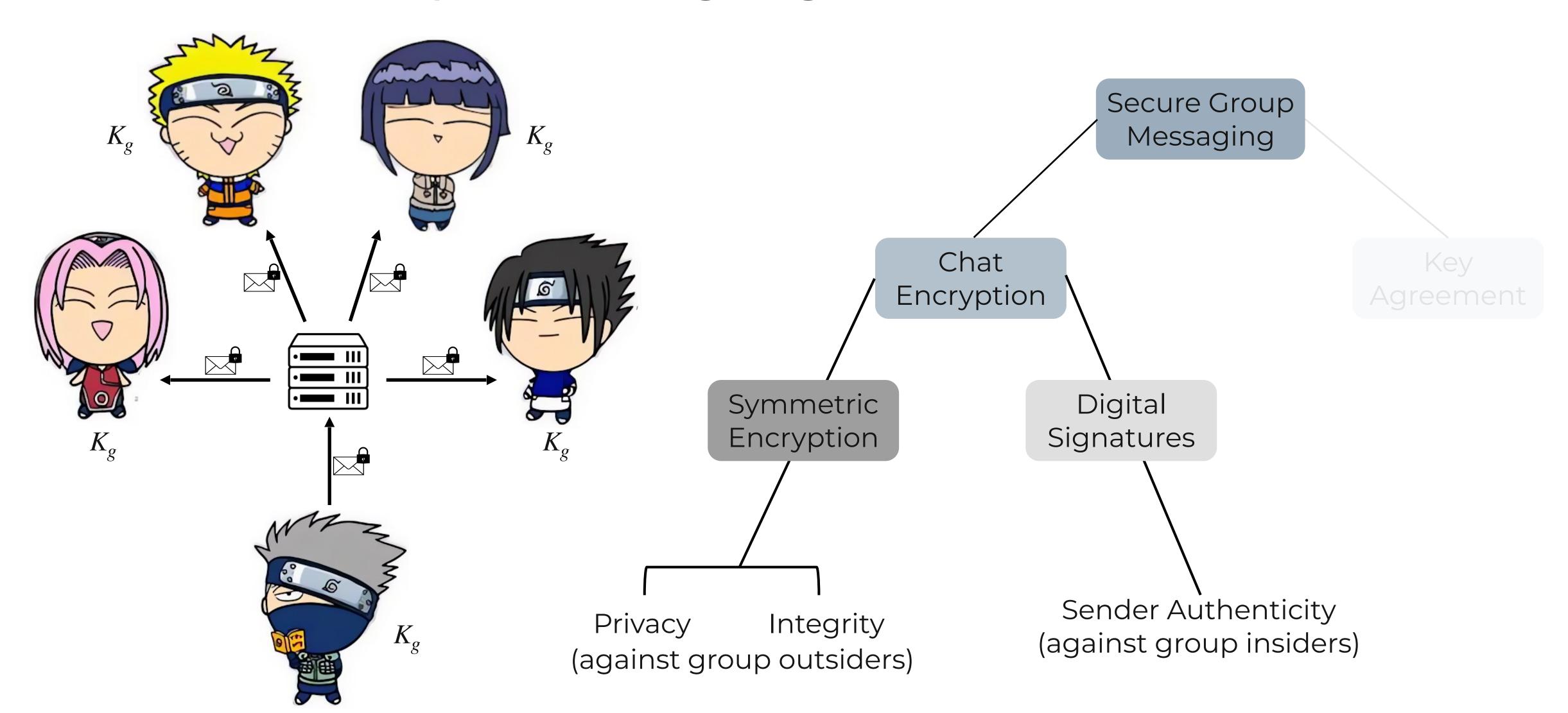


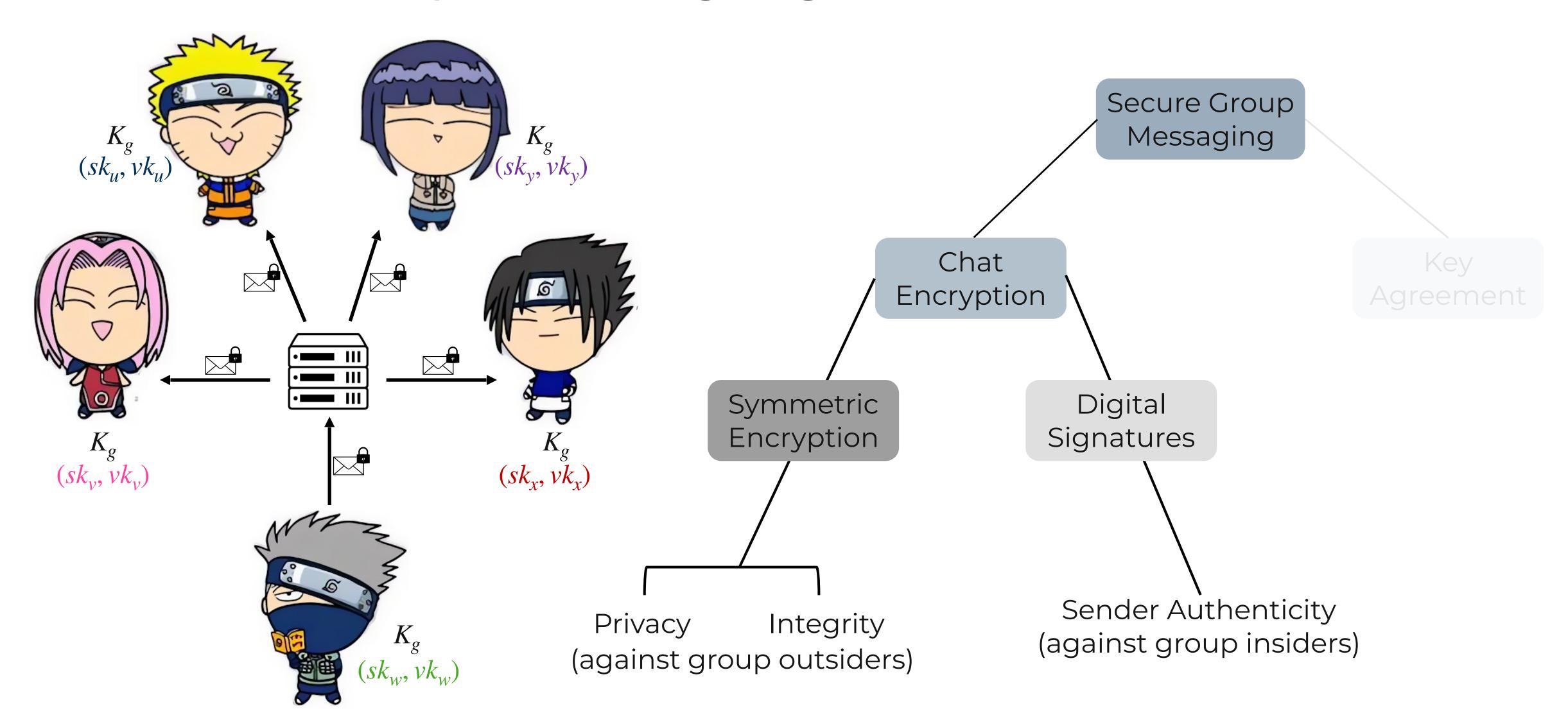


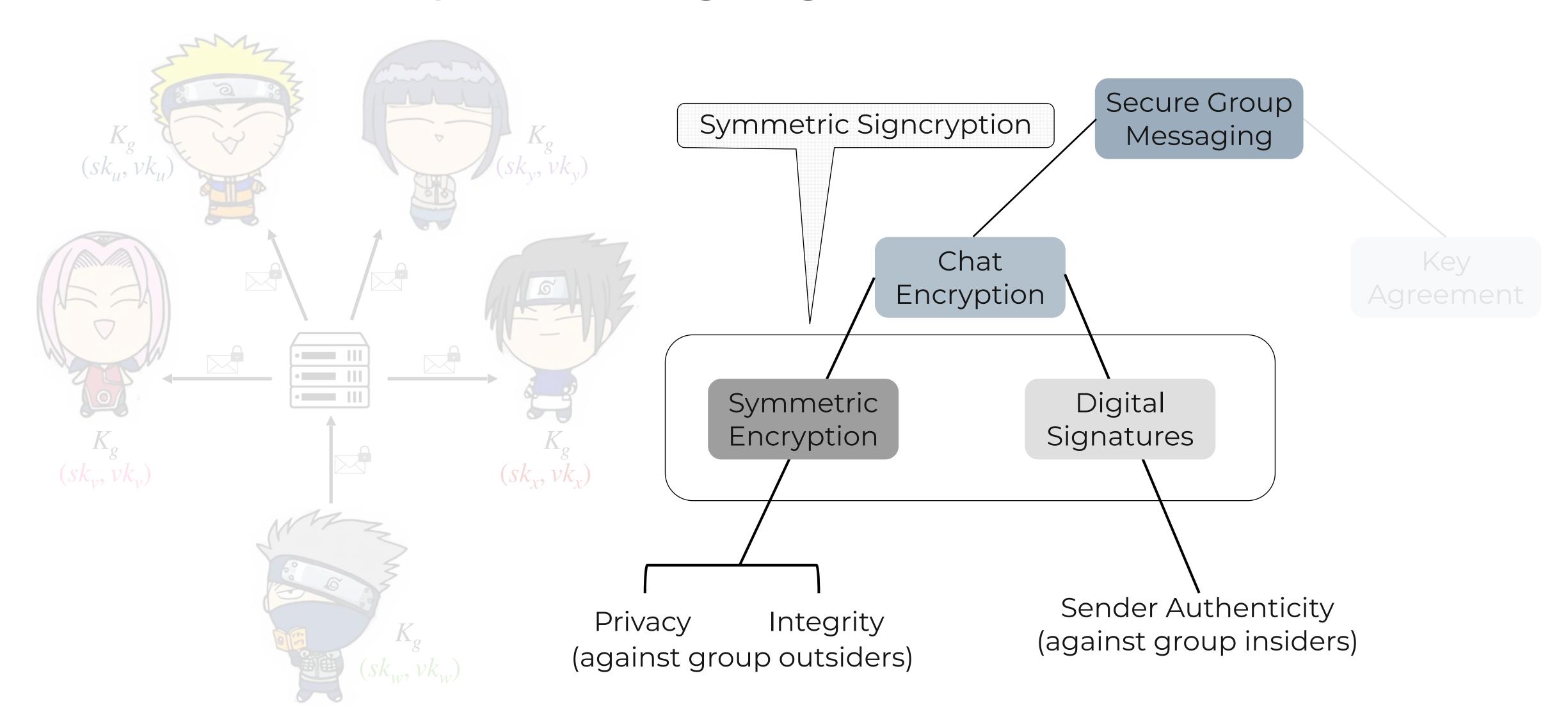












Symmetric Signcryption and E2EE Group Messaging in Keybase

Joseph Jaeger¹, Akshaya Kumar¹, and Igors Stepanovs²

Analyzing Group Chat Encryption in MLS, Session, Signal, and Matrix

Joseph Jaeger p and Akshaya Kumar

Symmetric Signcryption and E2EE Group Messaging in Keybase

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Symmetric Signcryption Model Analyzing Group Chat Encryption in MLS, Session, Signal, and Matrix

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Symmetric Signcryption and E2EE Group Messaging in Keybase

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Analyzing Group Chat Encryption in MLS, Session, Signal, and Matrix

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Symmetric Signcryption Model



Application







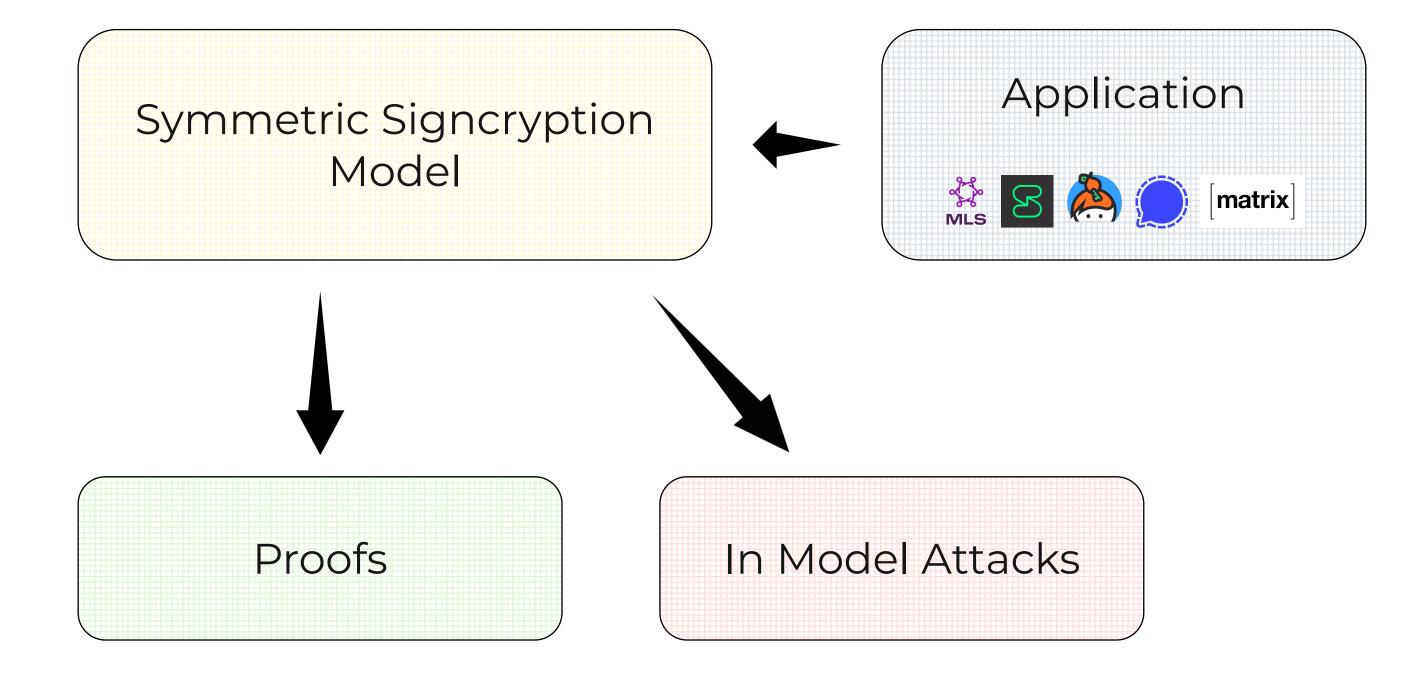


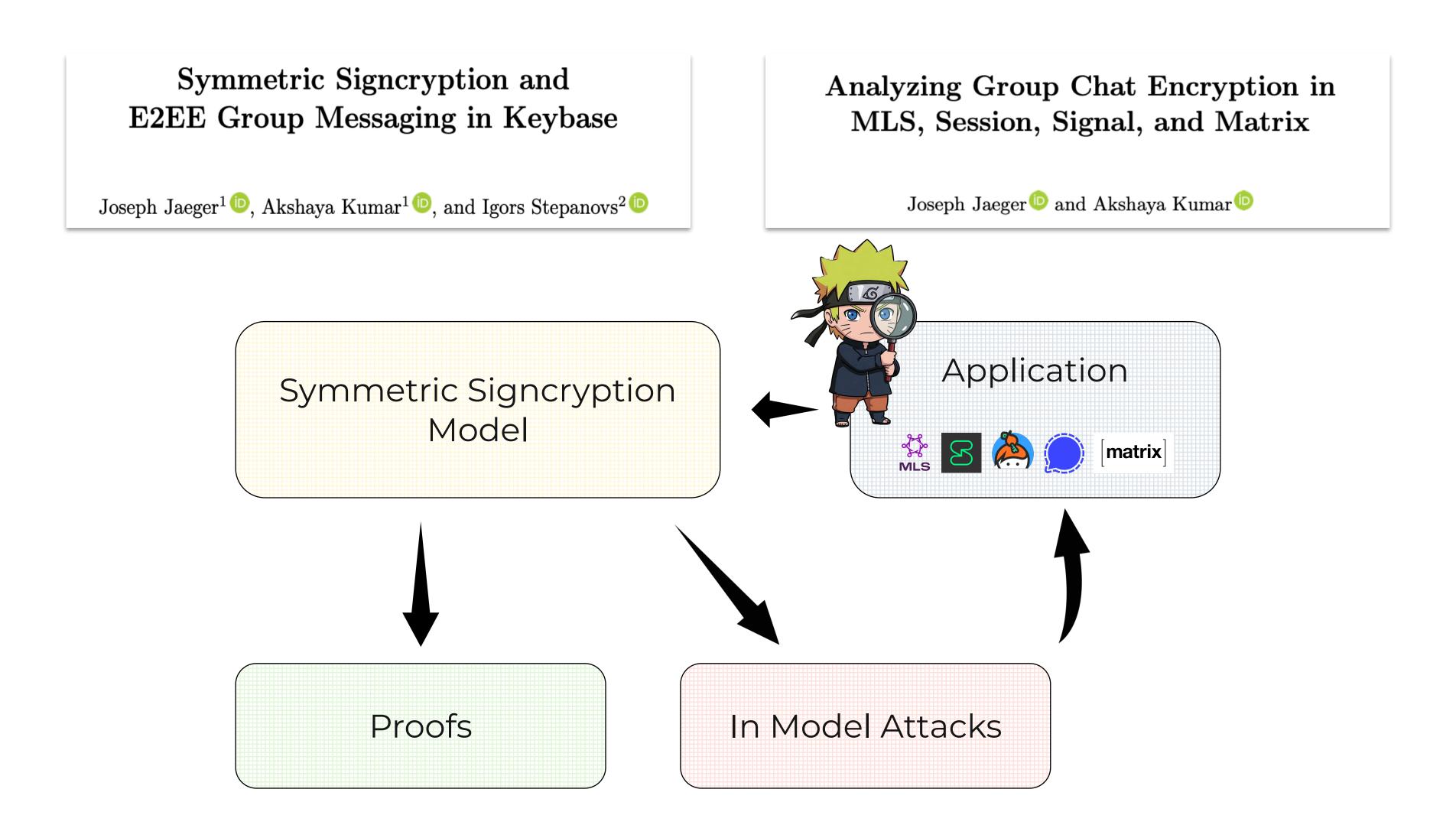
Symmetric Signcryption and E2EE Group Messaging in Keybase

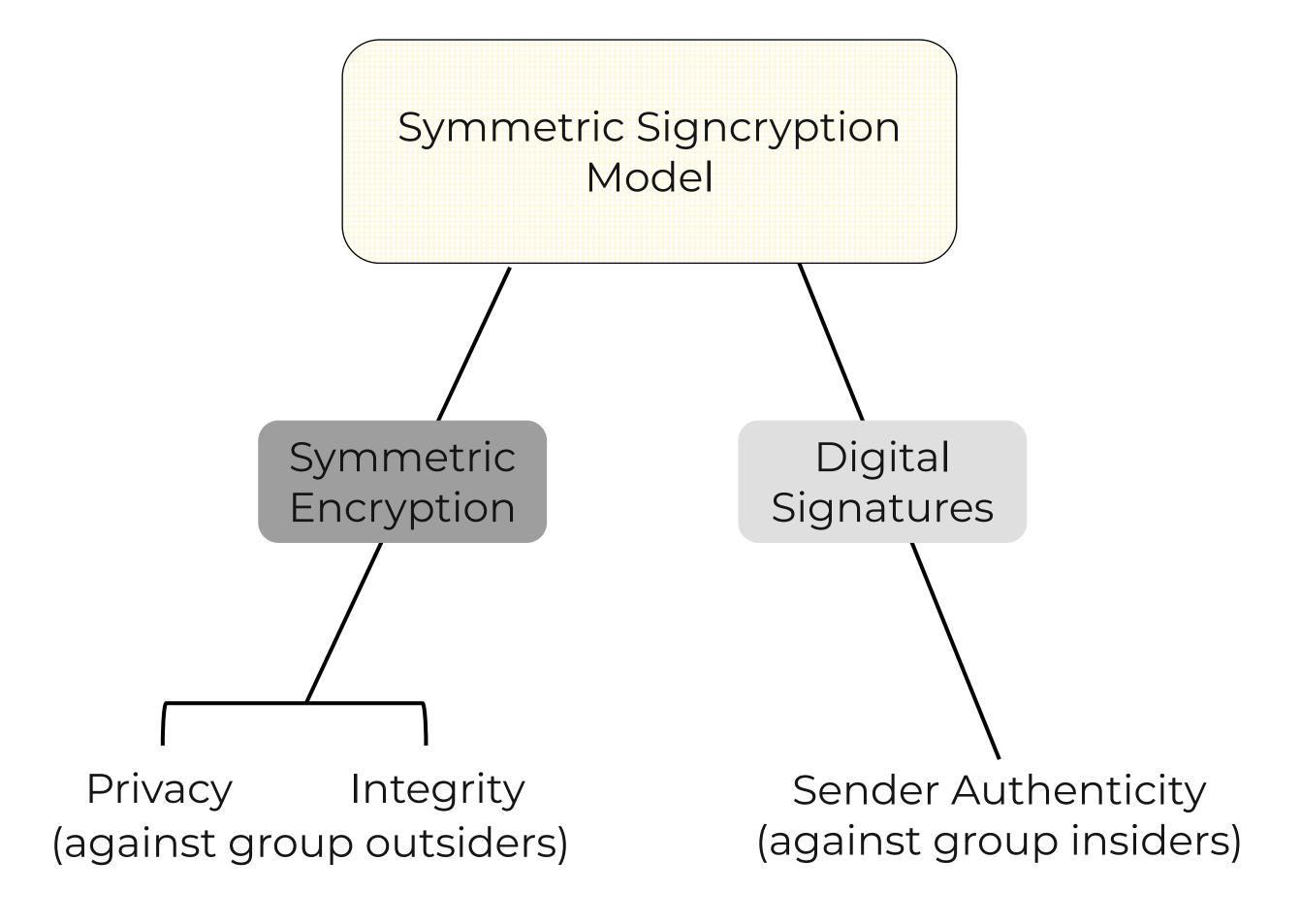
Joseph Jaeger¹, Akshaya Kumar¹, and Igors Stepanovs²

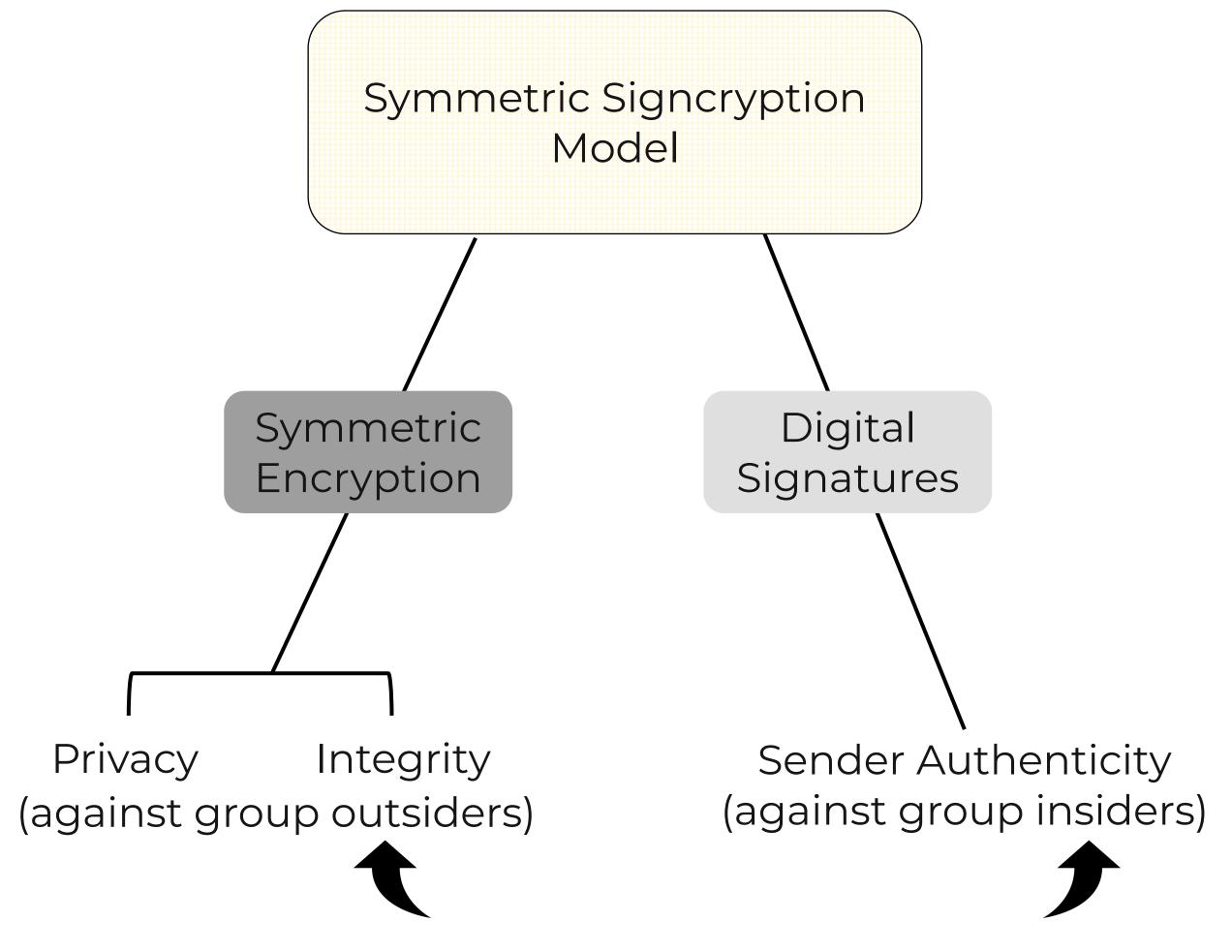
Analyzing Group Chat Encryption in MLS, Session, Signal, and Matrix

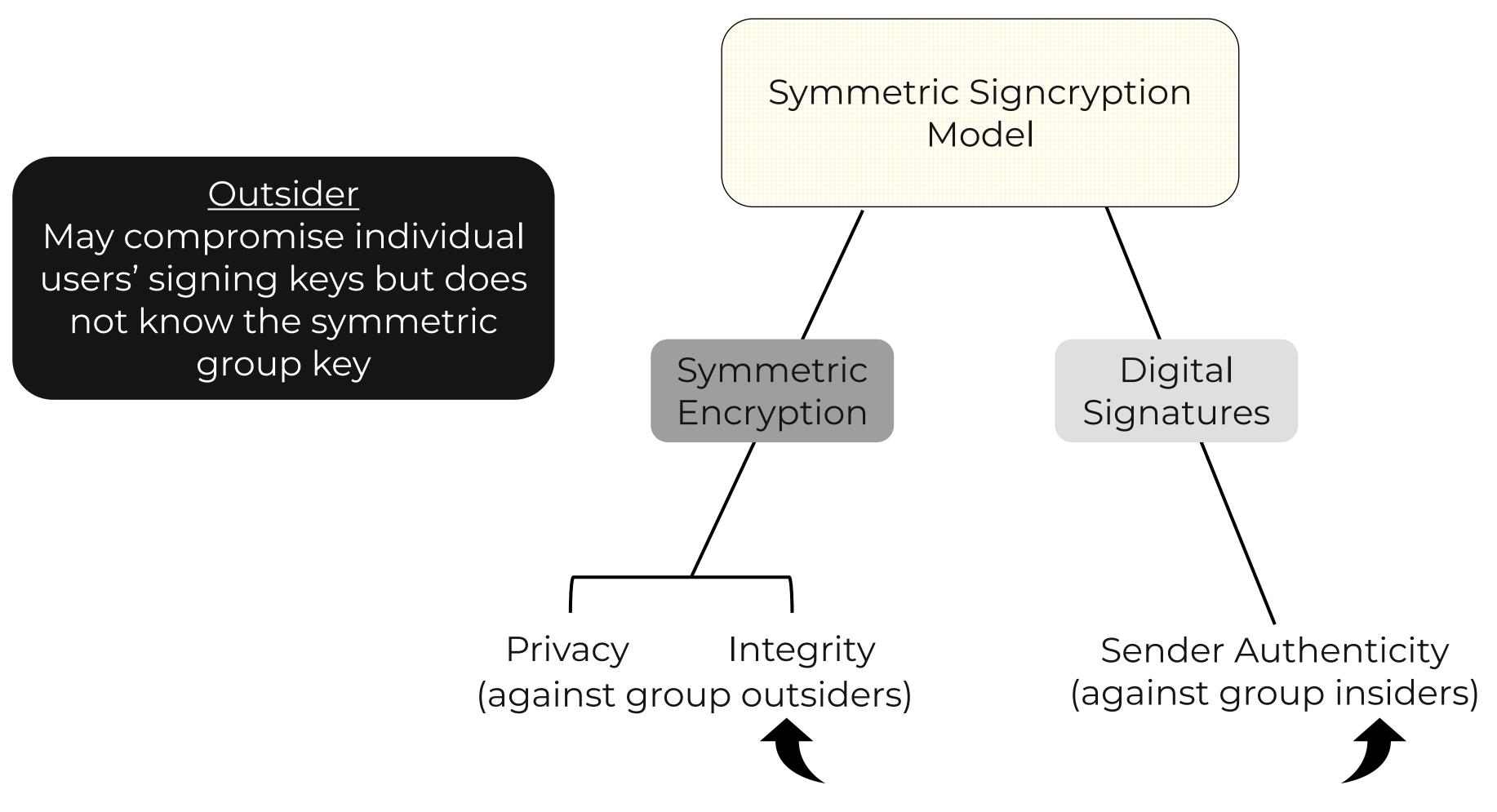
Joseph Jaeger p and Akshaya Kumar

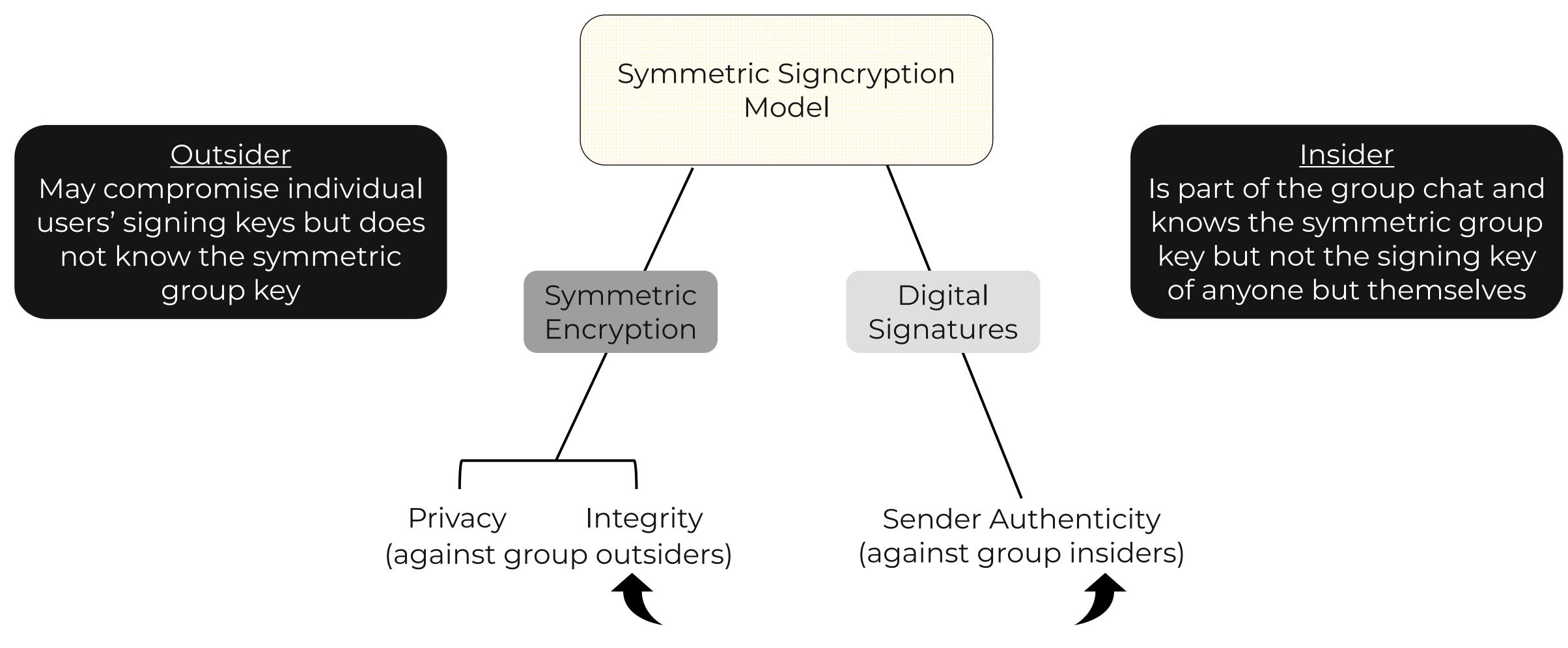


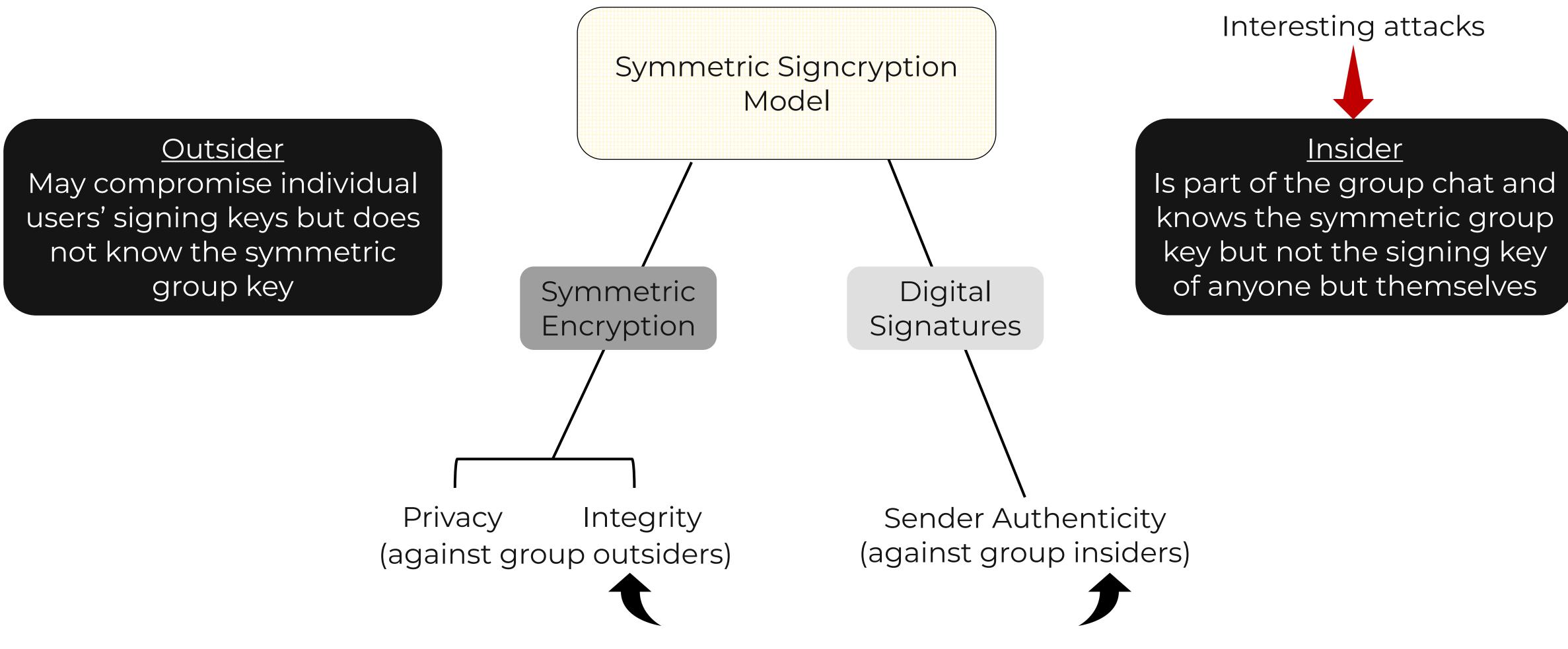










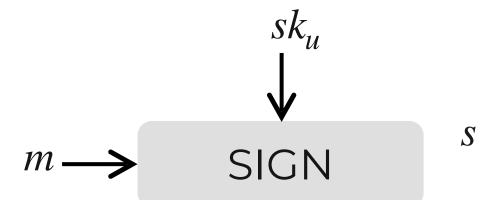


Sign-then-Encrypt (StE)

Sign-then-Encrypt (StE)

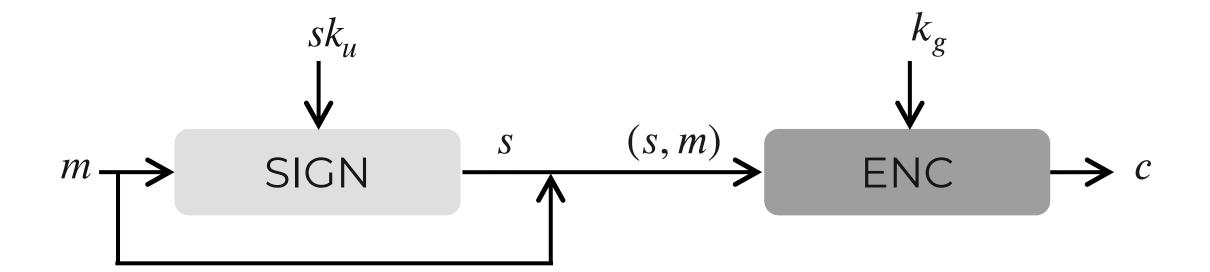
Encrypt-then-Sign (EtS)

Sign-then-Encrypt (StE)

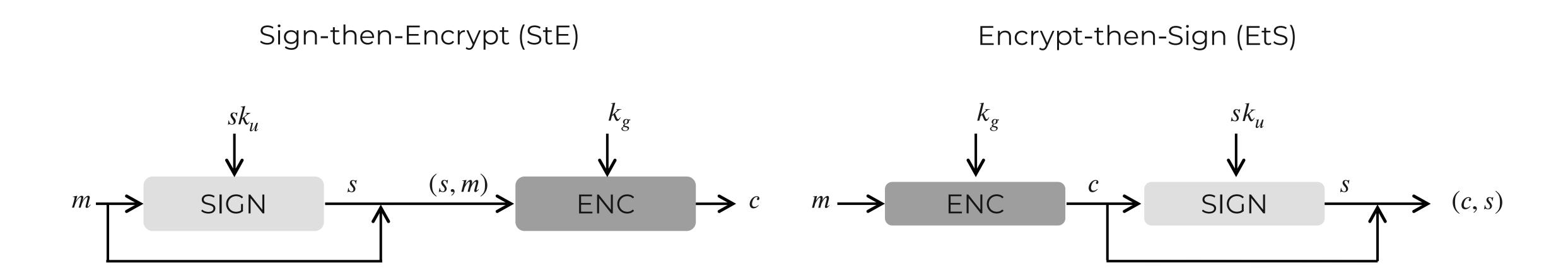


Encrypt-then-Sign (EtS)

Sign-then-Encrypt (StE)



Encrypt-then-Sign (EtS)



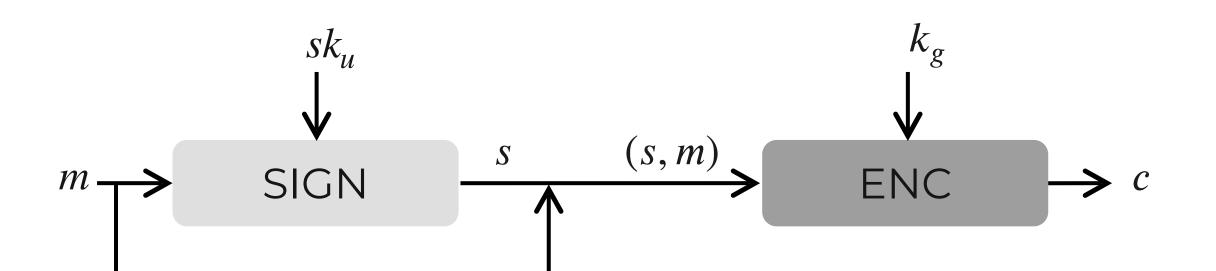
Sign-then-Encrypt (StE) sk_{u} $sk_$



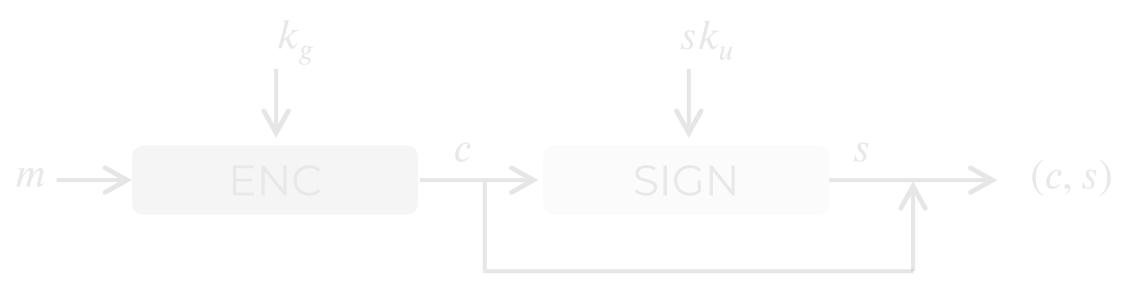
Sign-then-Encrypt (StE) sk_{u} $m \rightarrow SIGN$ Encrypt-then-Sign (EtS) k_{g} k_{g}

matrix

Sign-then-Encrypt (StE)



Encrypt-then-Sign (EtS)









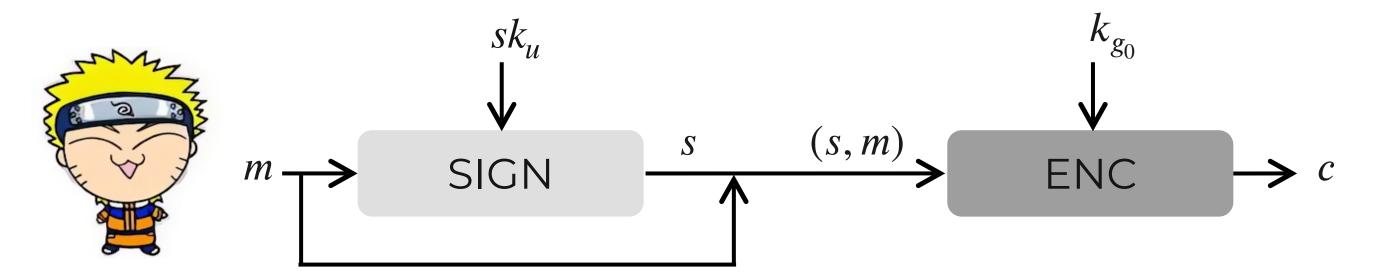




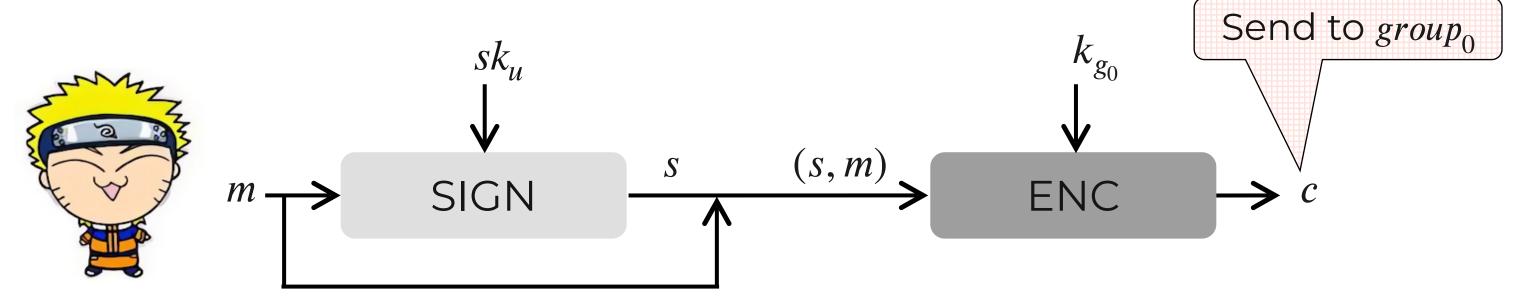




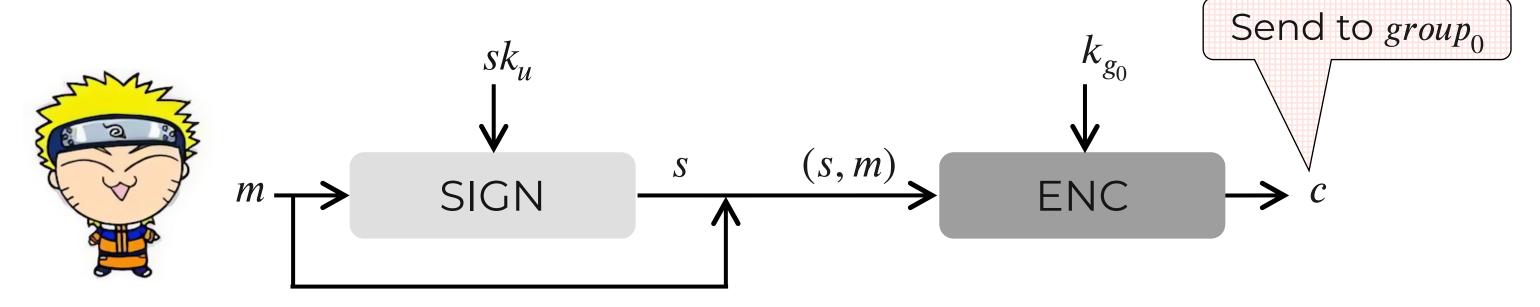


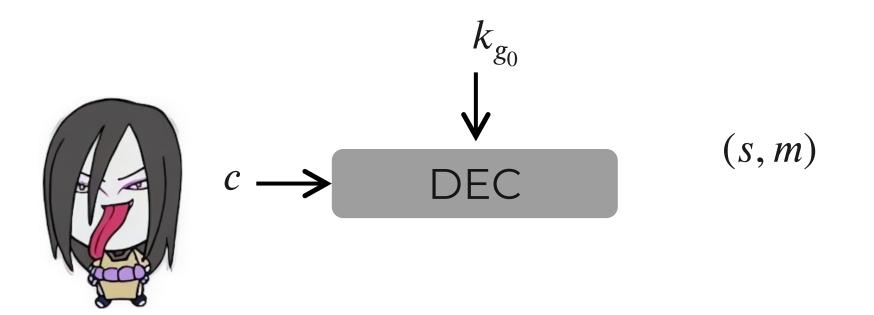




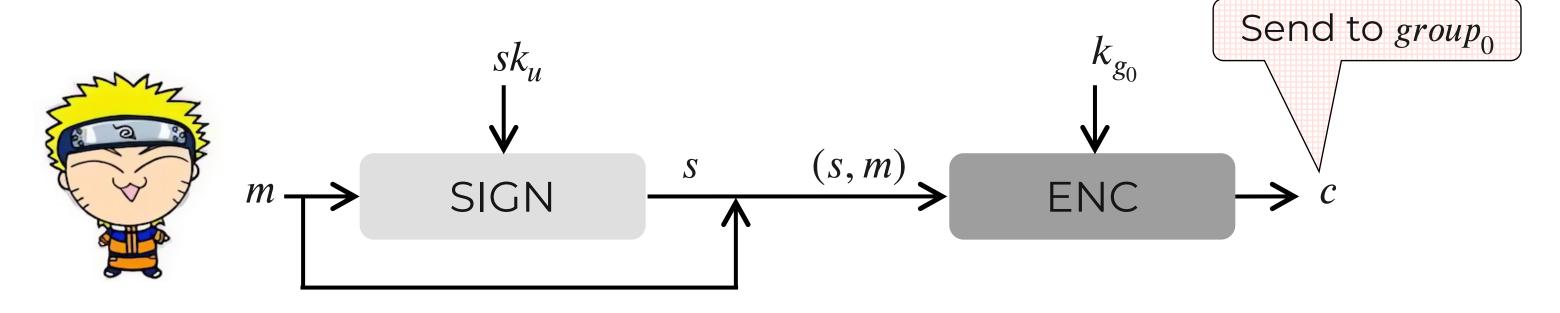


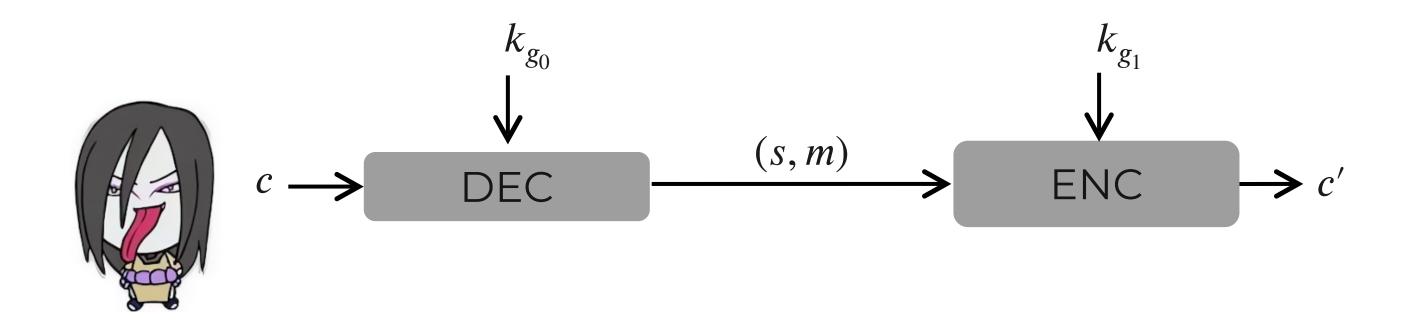




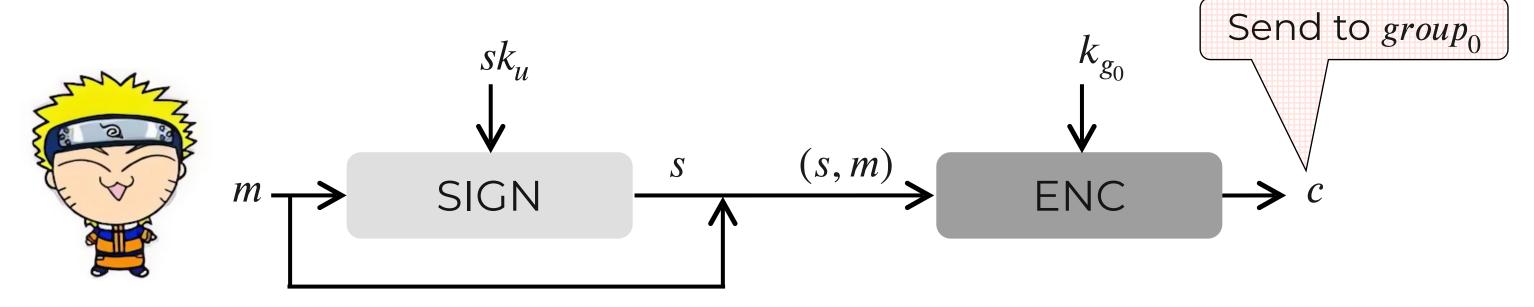


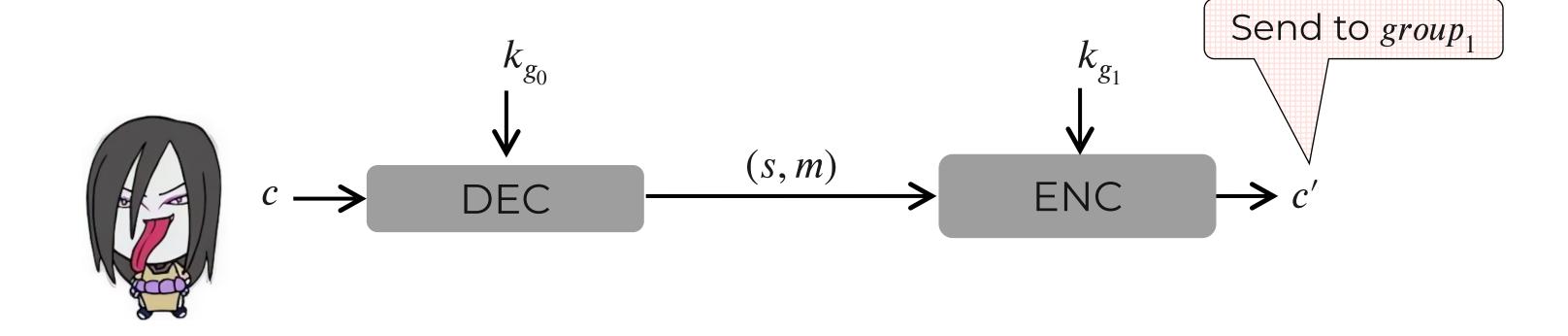


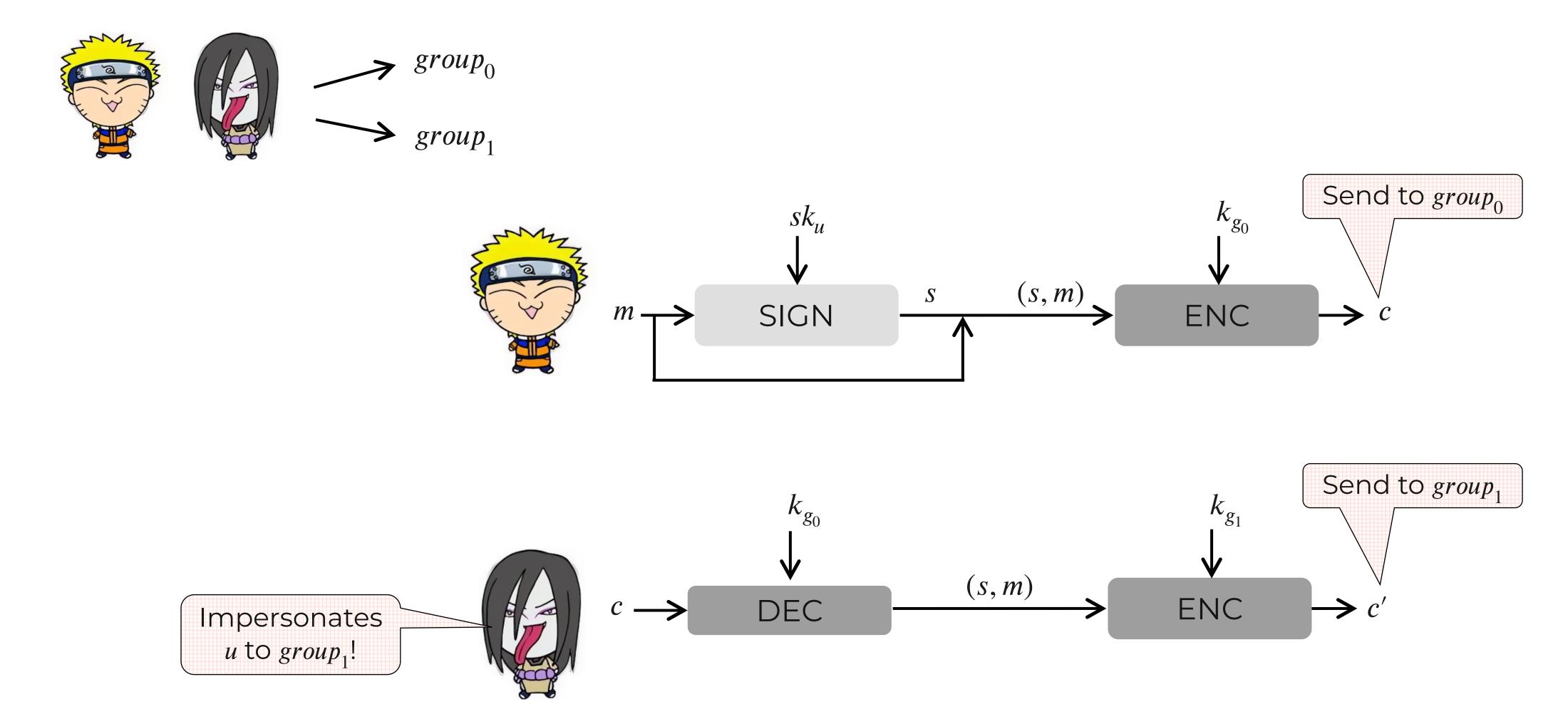


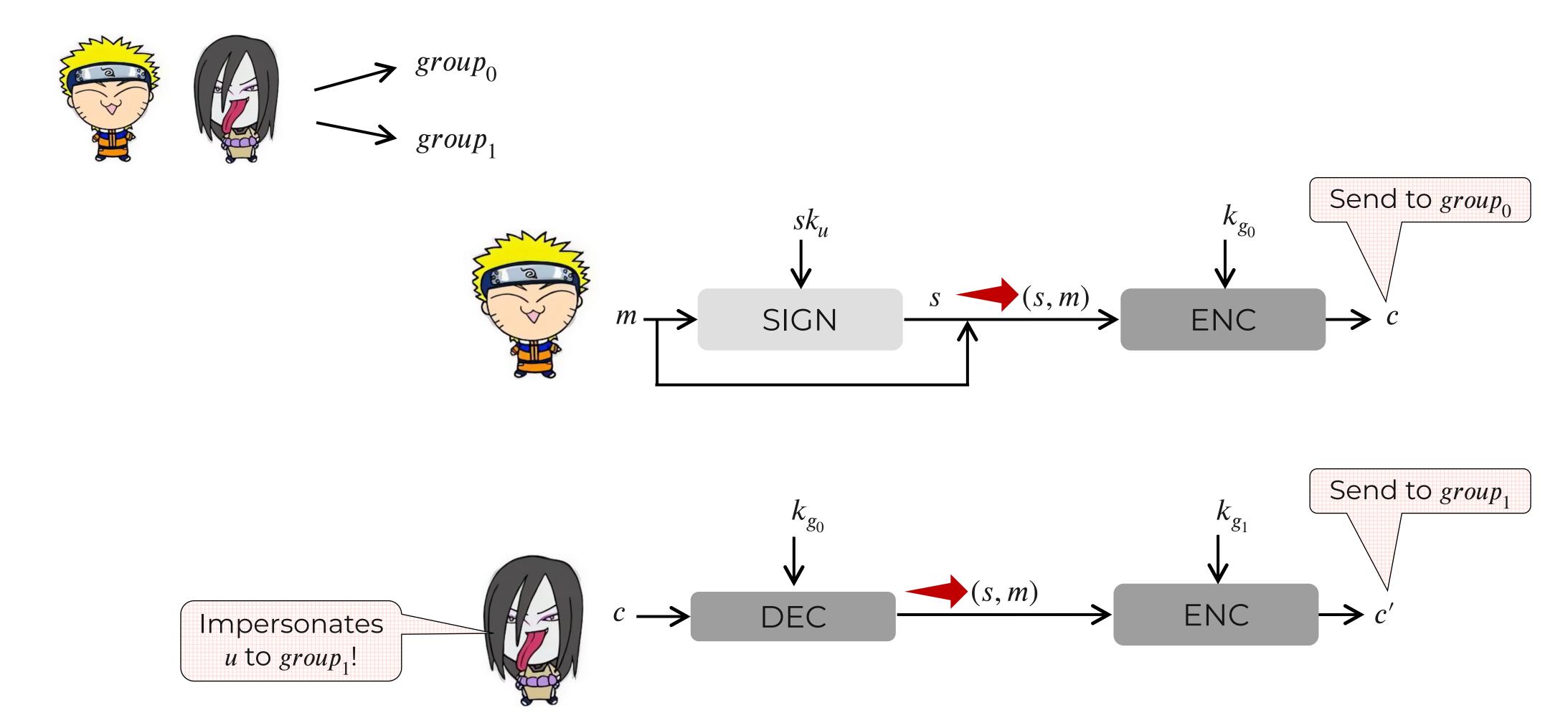


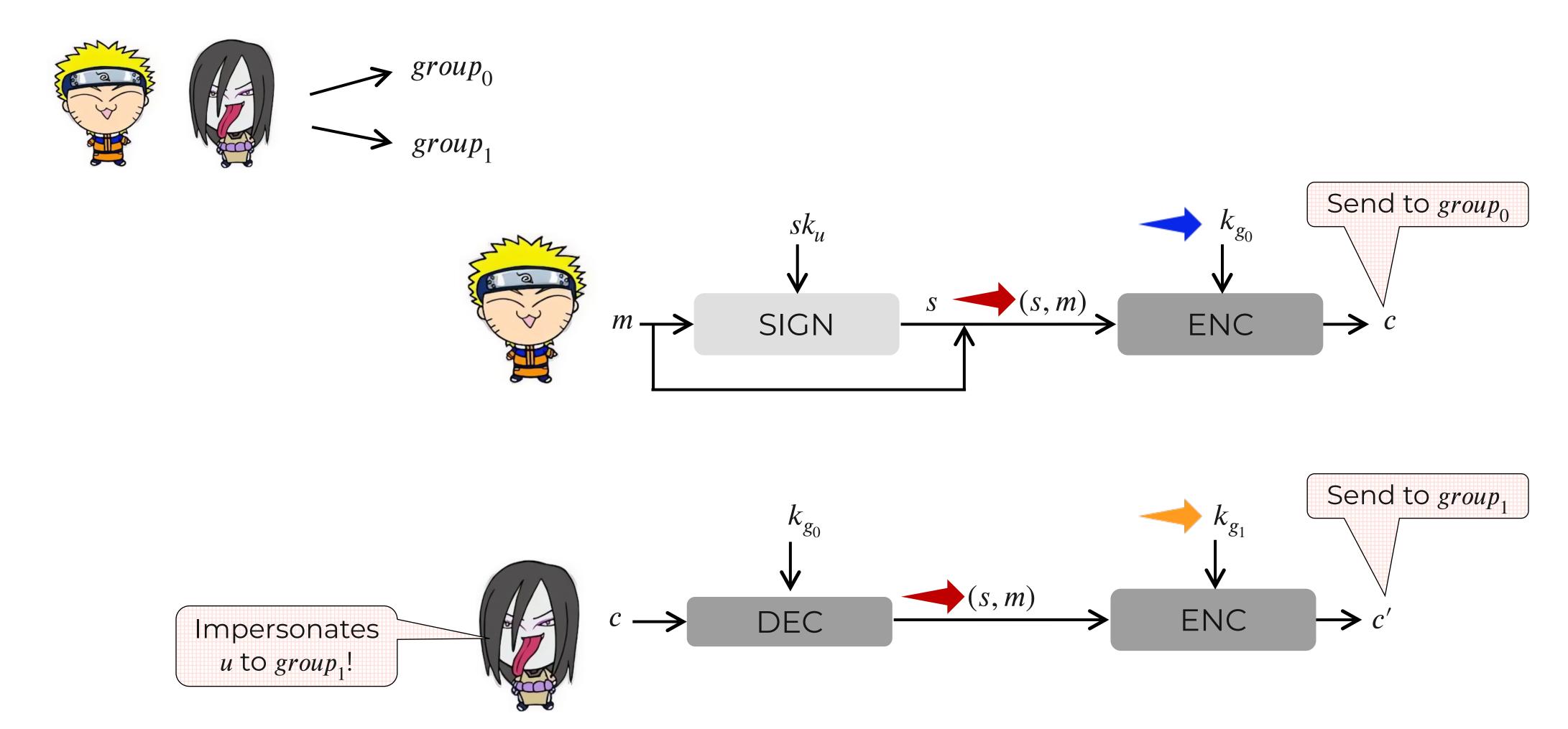


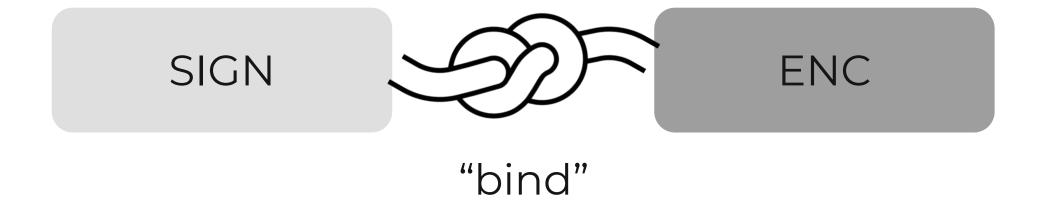


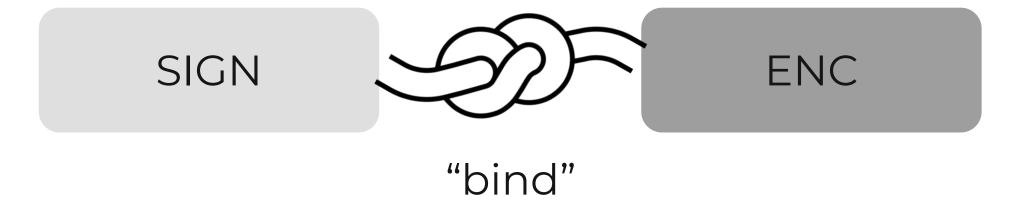




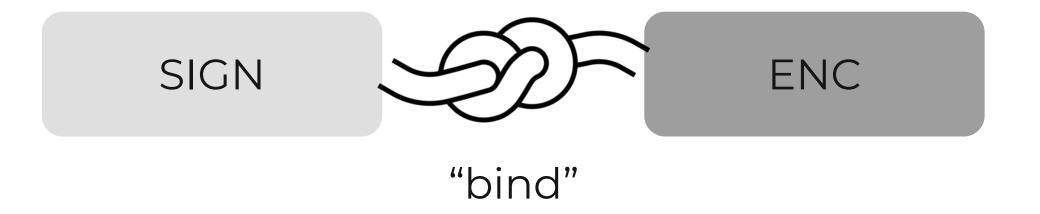


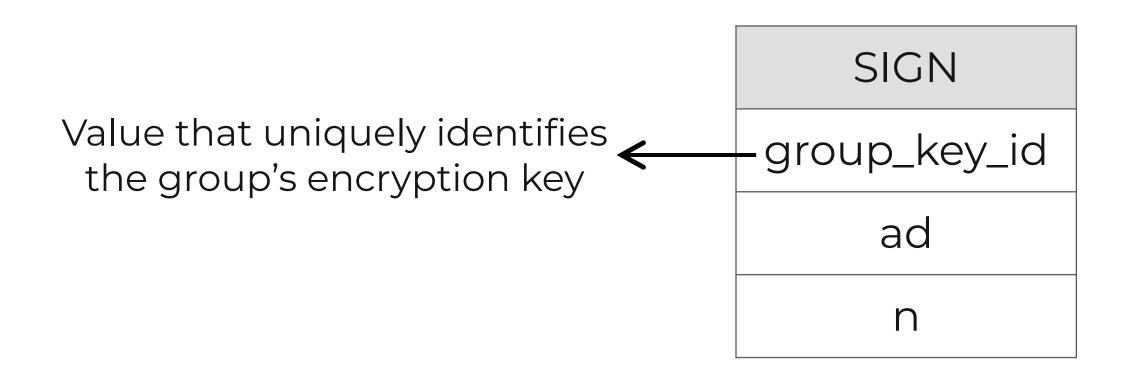


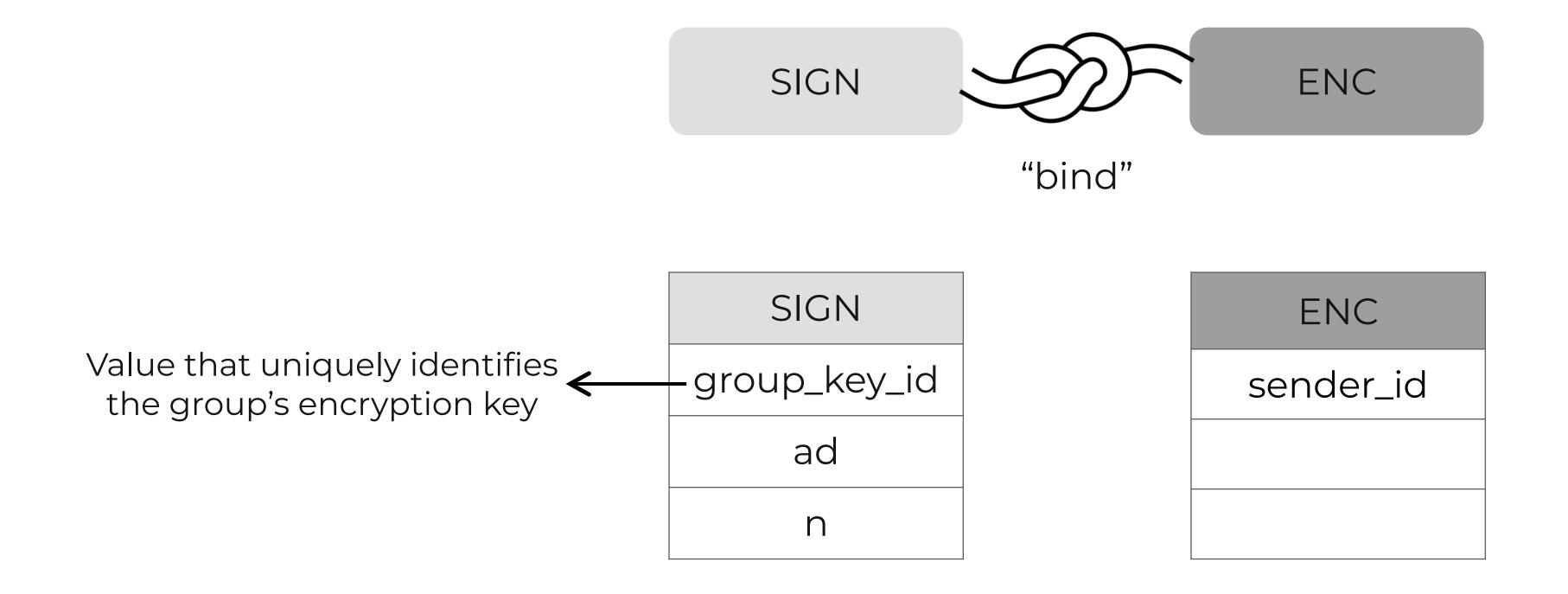




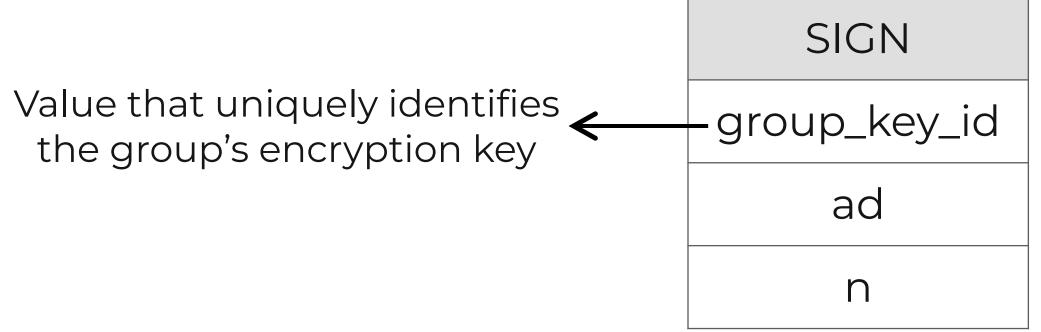
sign group_key_id ad n

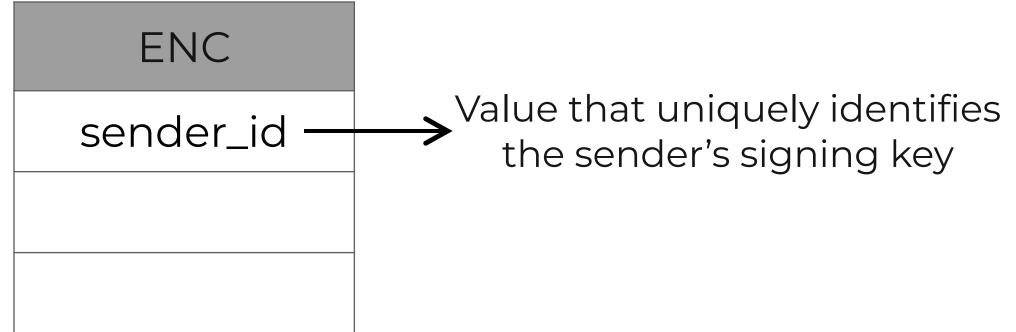








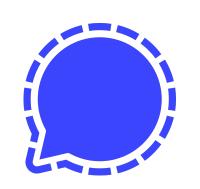










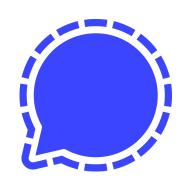


matrix









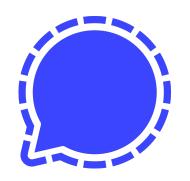














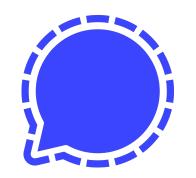




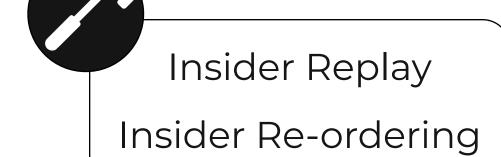














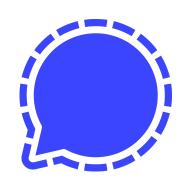


^{*} stolen signing key

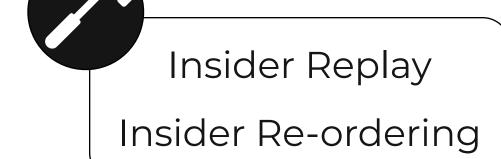
















No context binding



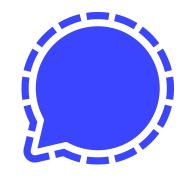
No context binding

^{*} stolen signing key

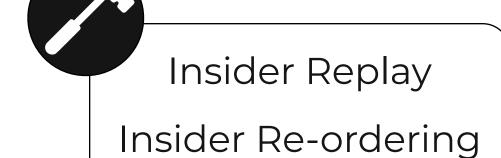




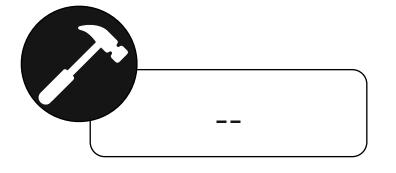




matrix









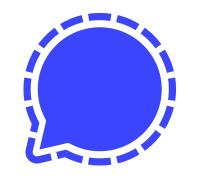


^{*} stolen signing key

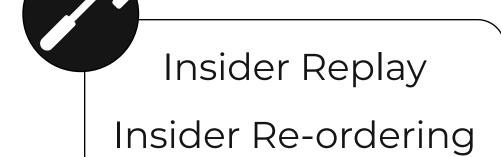




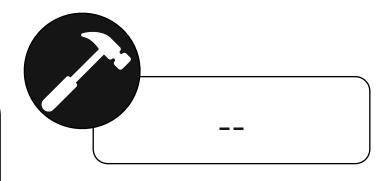














No context binding



No context binding



Key Reuse

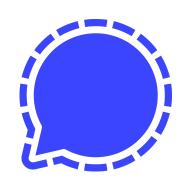
Key Cycle

^{*} stolen signing key





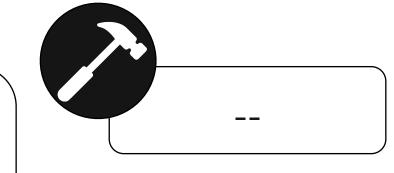




matrix











No context binding



No context binding



Key Reuse

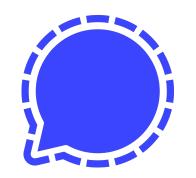
Key Cycle

^{*} stolen signing key * discovered by [BCG23]

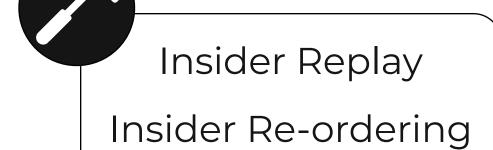




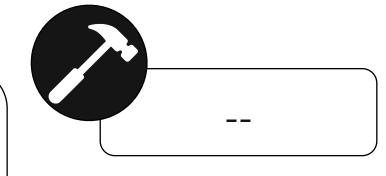




matrix











No context binding



No context binding



Key Reuse Key Cycle



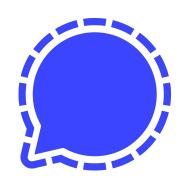
Unauthenticated
Symmetric
Encryption

^{*} stolen signing key * discovered by [BCG23]





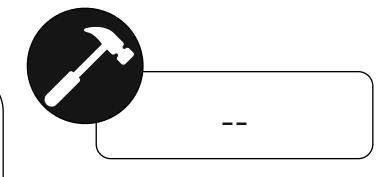




matrix

Insider Replay
Insider Re-ordering

Insider Replay
Outsider Replay
Outsider Forgery*









No context binding



No context binding



Key Reuse Key Cycle



Unauthenticated
Symmetric
Encryption



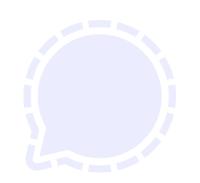
No context binding

^{*} stolen signing key * discovered by [BCG23]



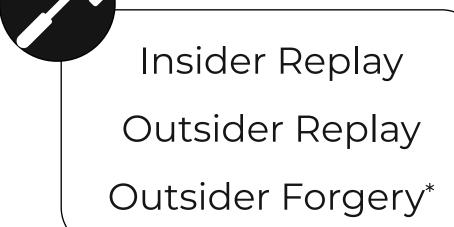


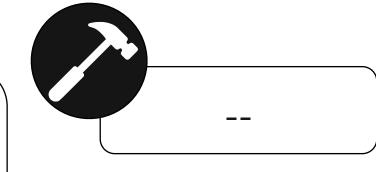




matrix

Insider Replay
Insider Re-ordering











No context binding



No context binding



Key Reuse Key Cycle



Unauthenticated
Symmetric
Encryption



No context binding

^{*} stolen signing key * discovered by [BCG23]

Case Study I: MLS



Encryption Key Derivation in MLS

Encryption Key Derivation in MLS

Ratchet tree

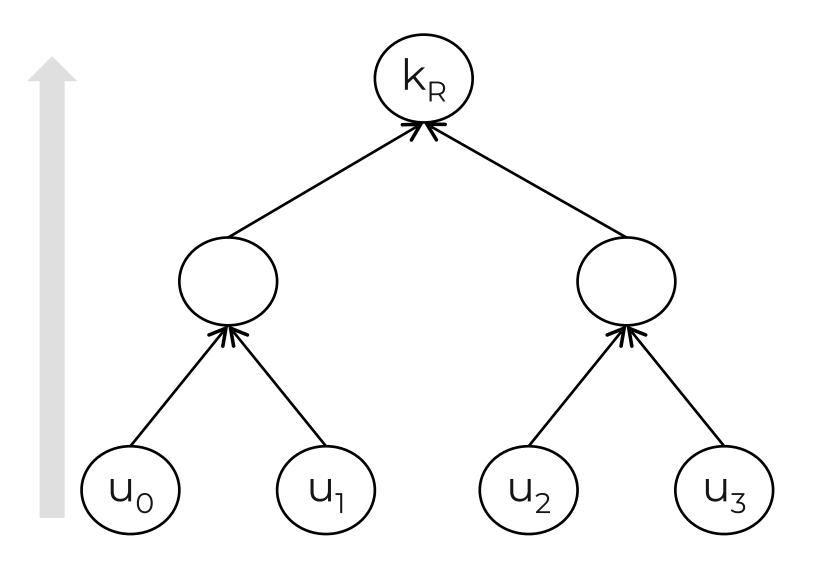




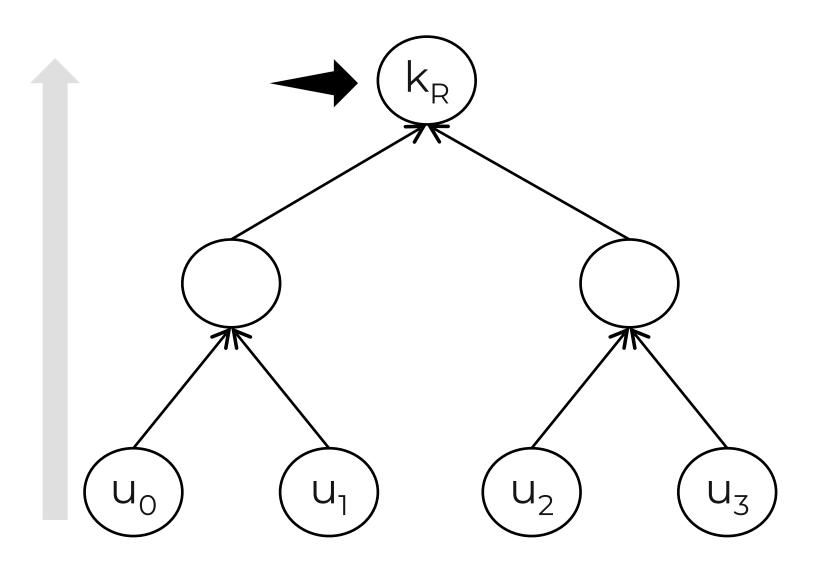


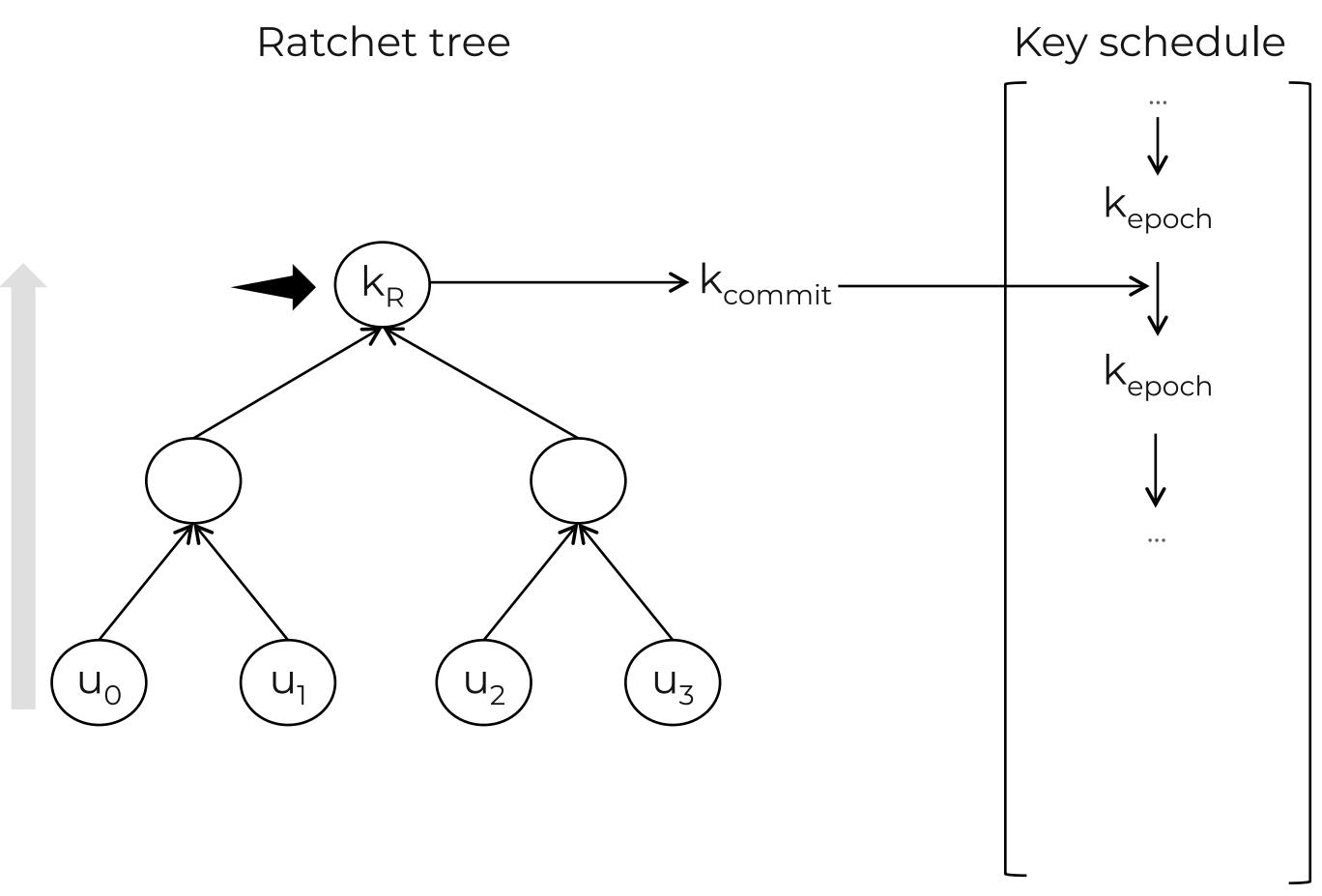


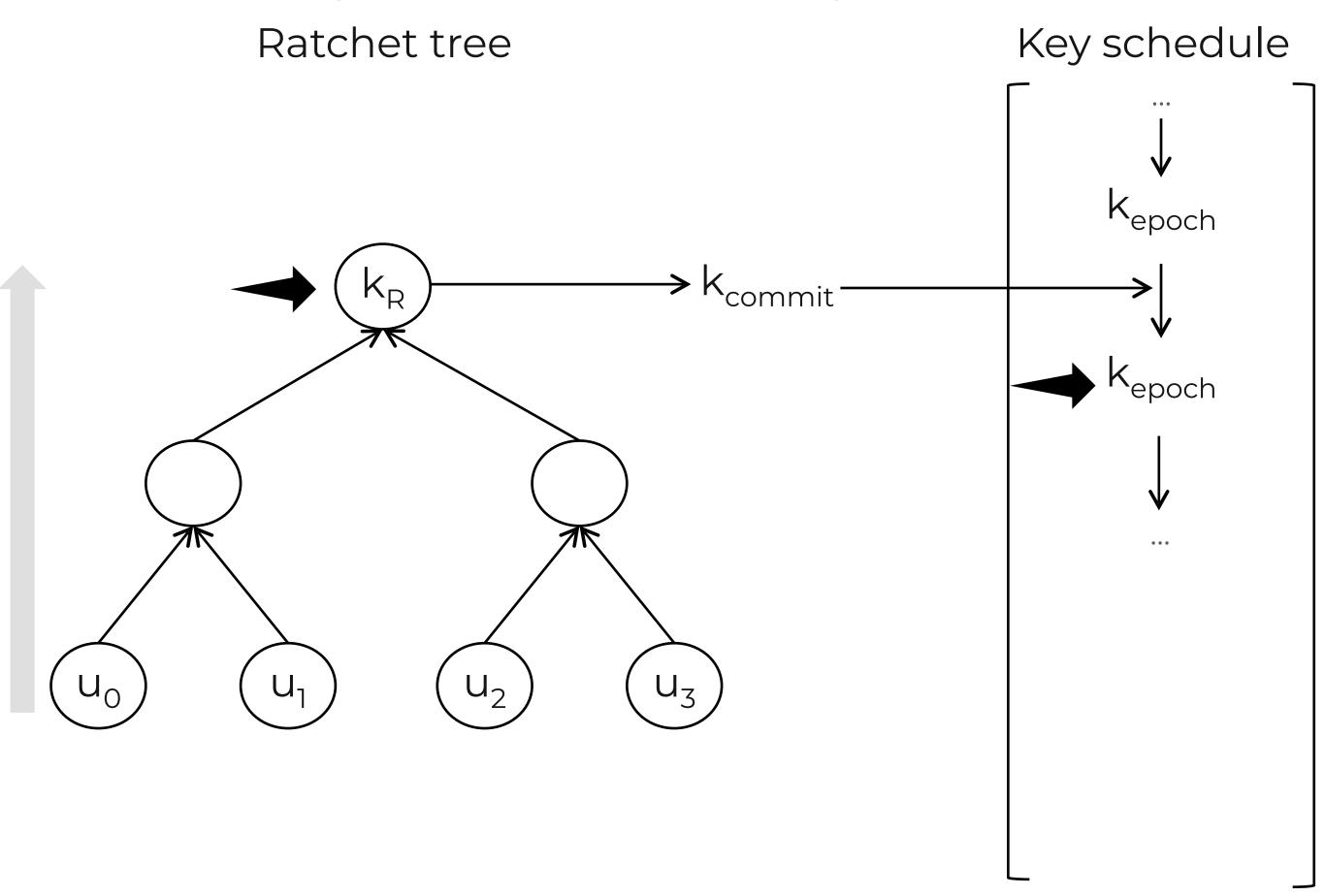
Ratchet tree

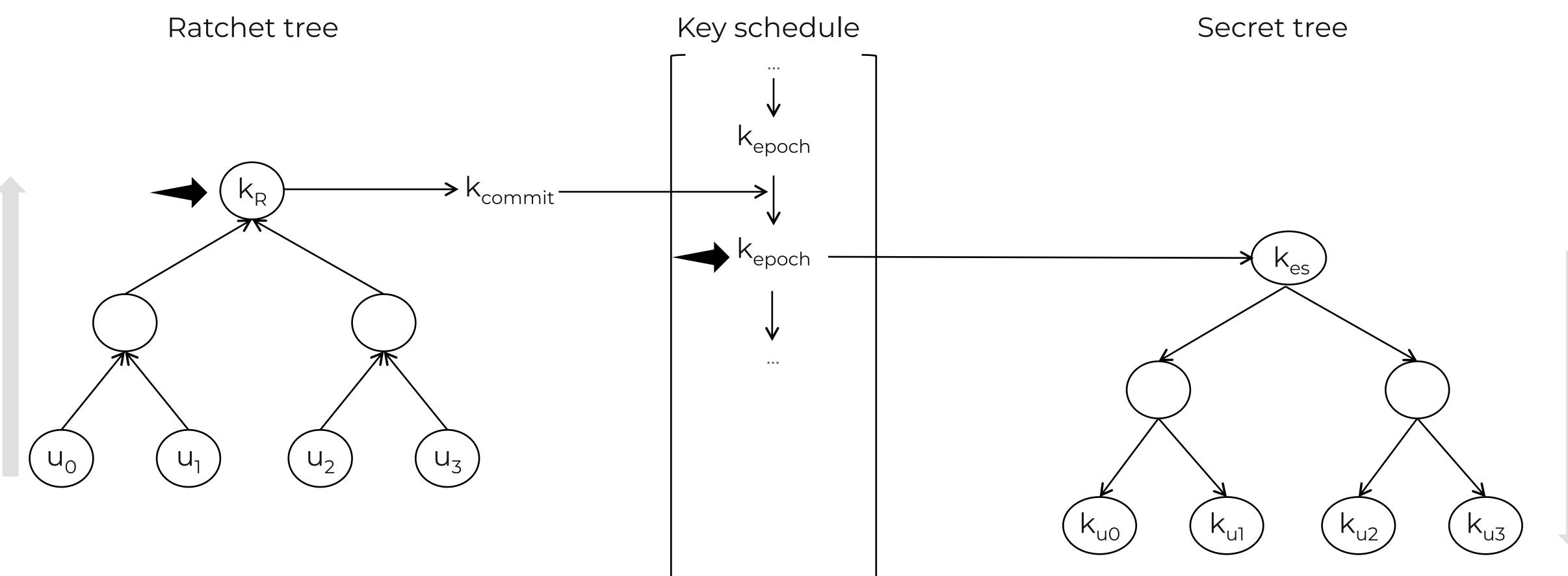


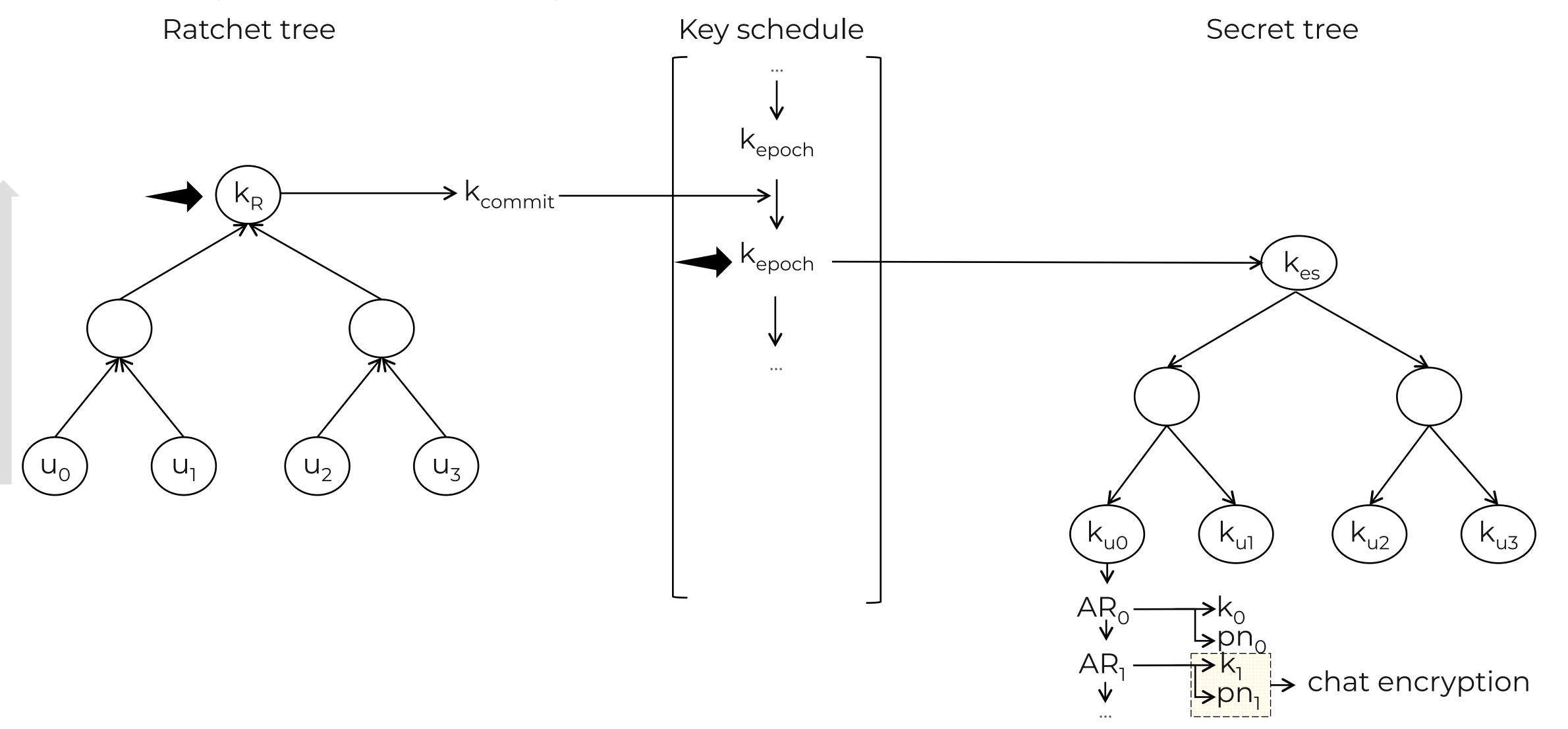
Ratchet tree

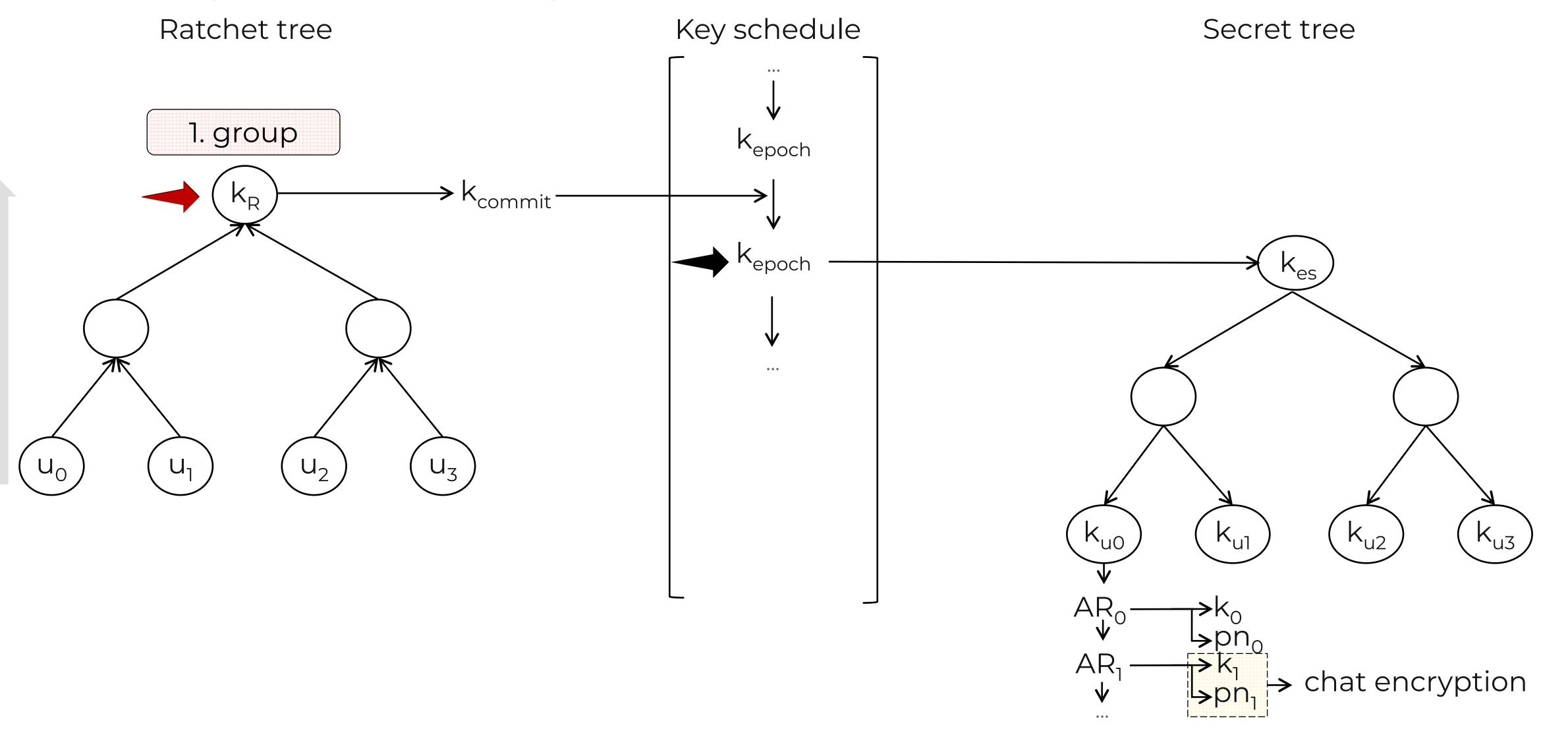


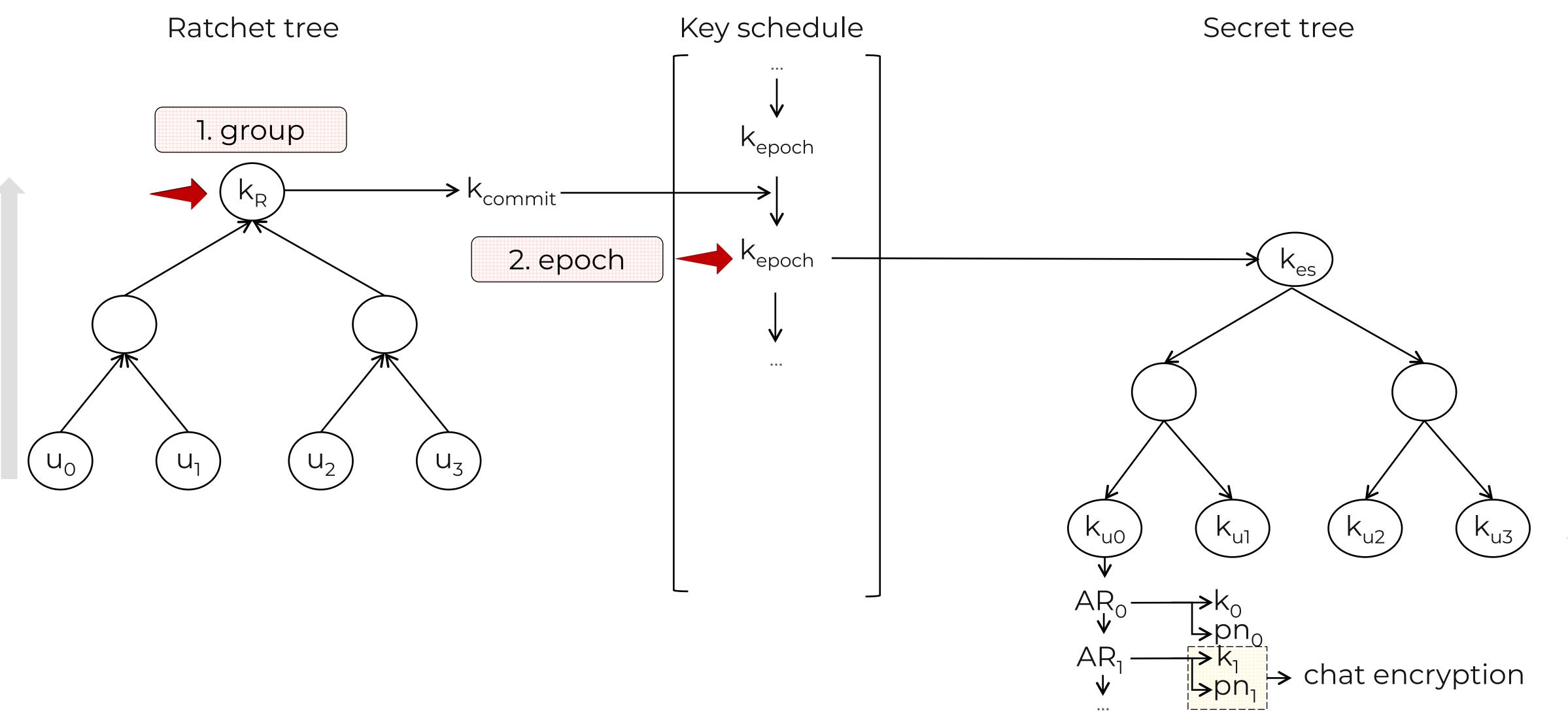


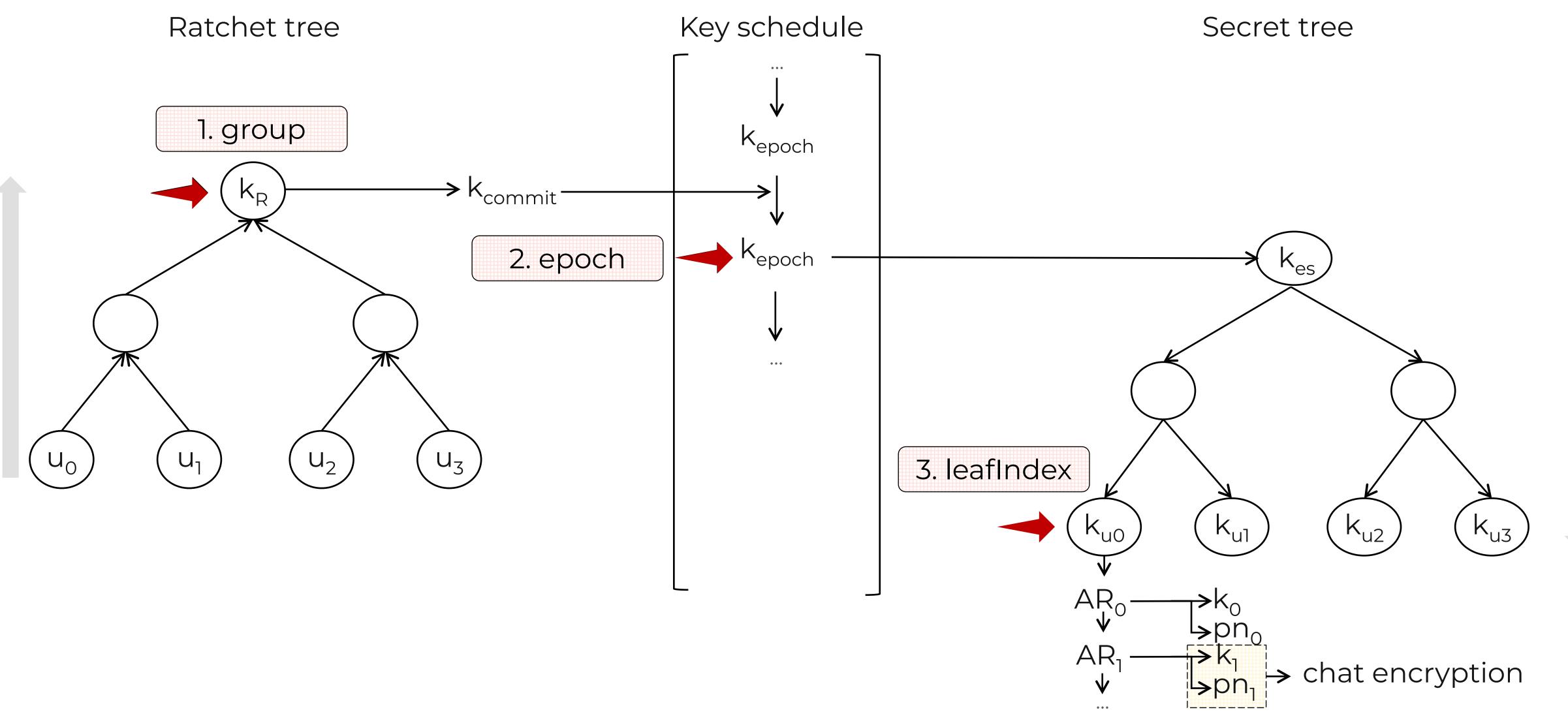


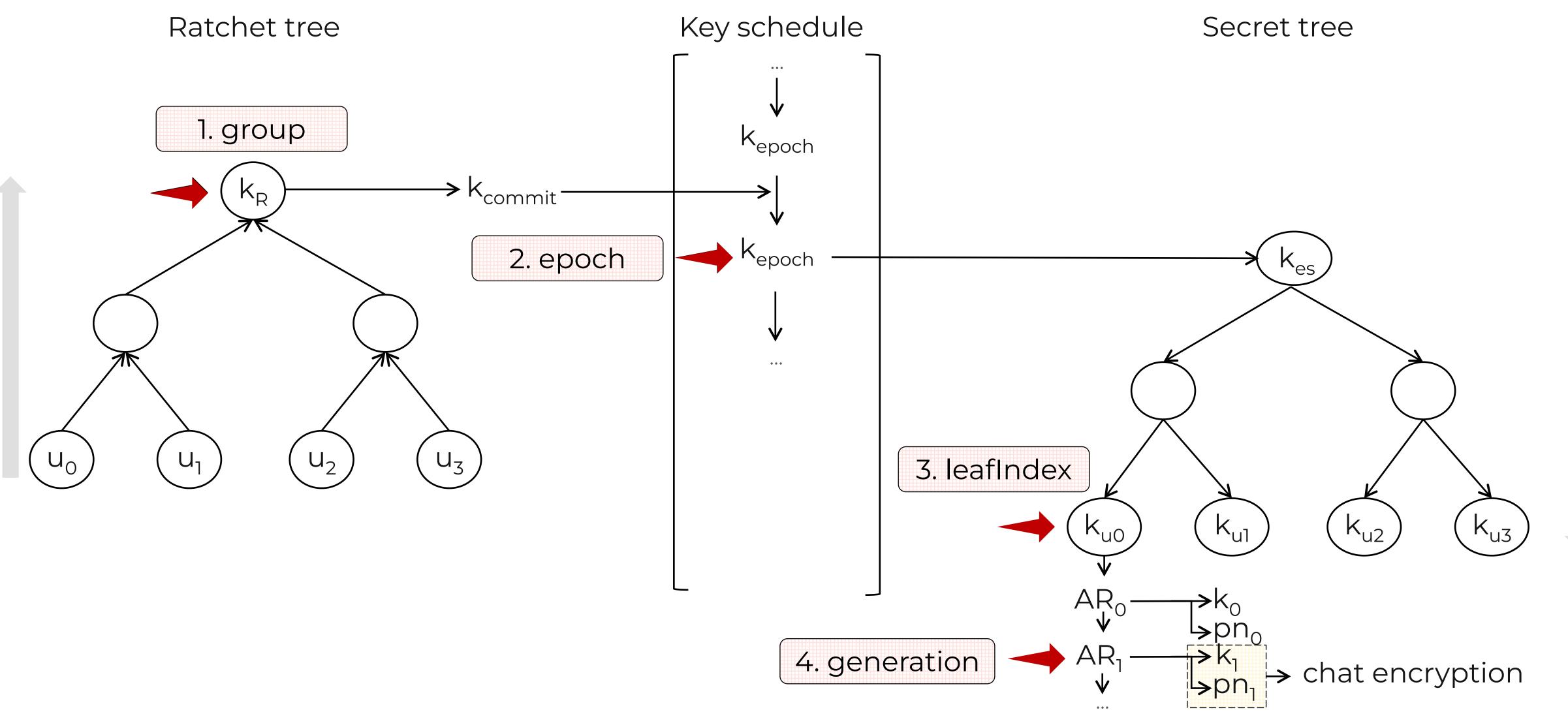


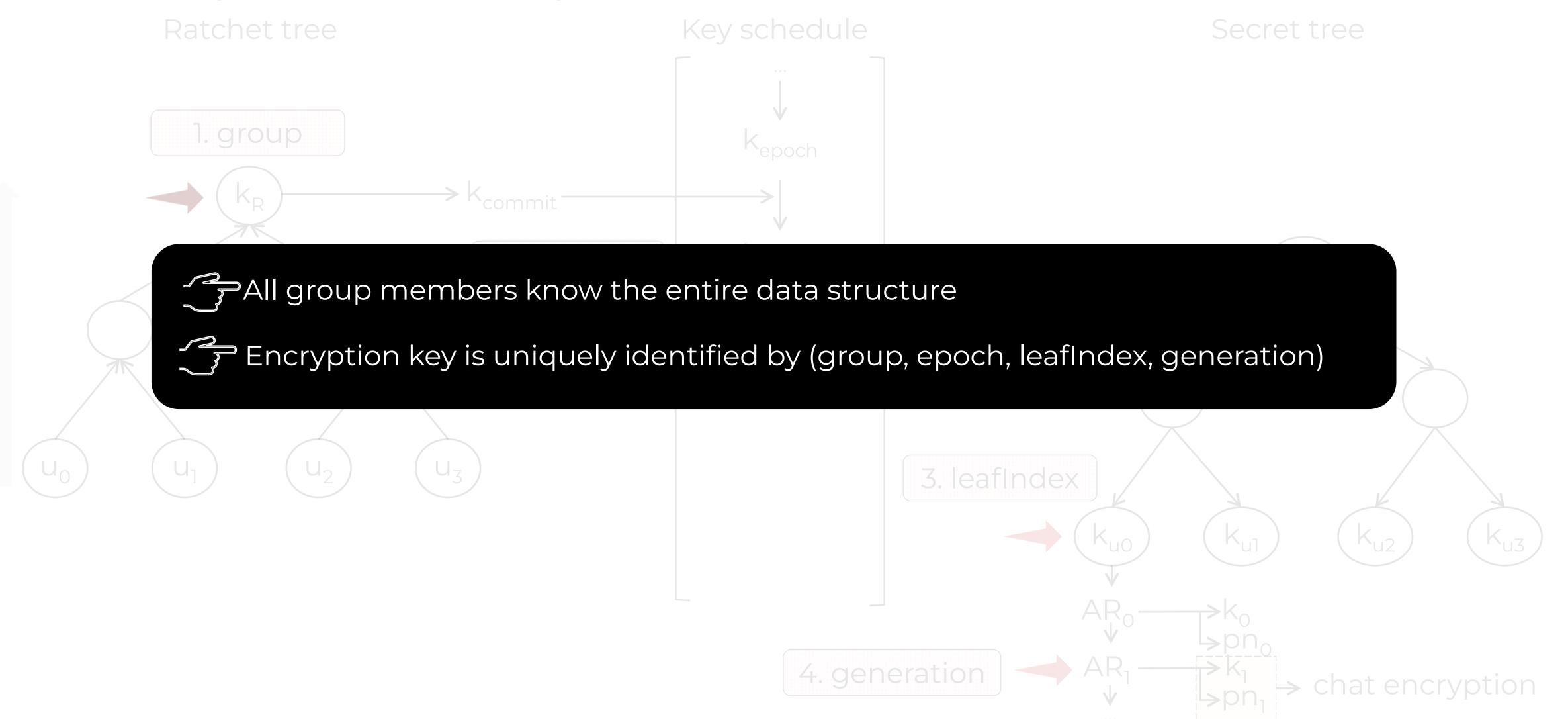










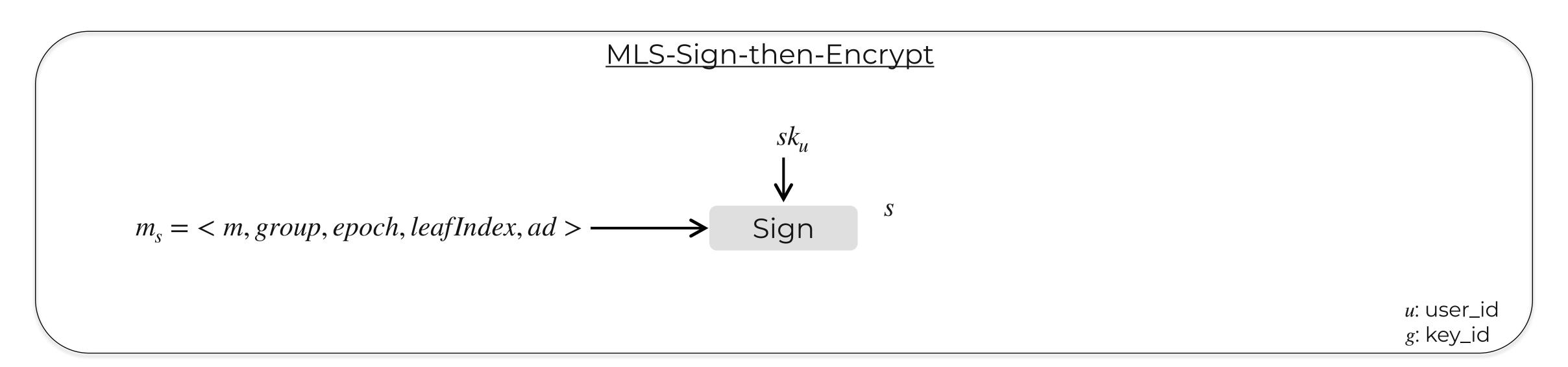


Chat encryption in MLS composes a digital signature scheme and a nonce-based encryption scheme in a Sign-then-Encrypt fashion

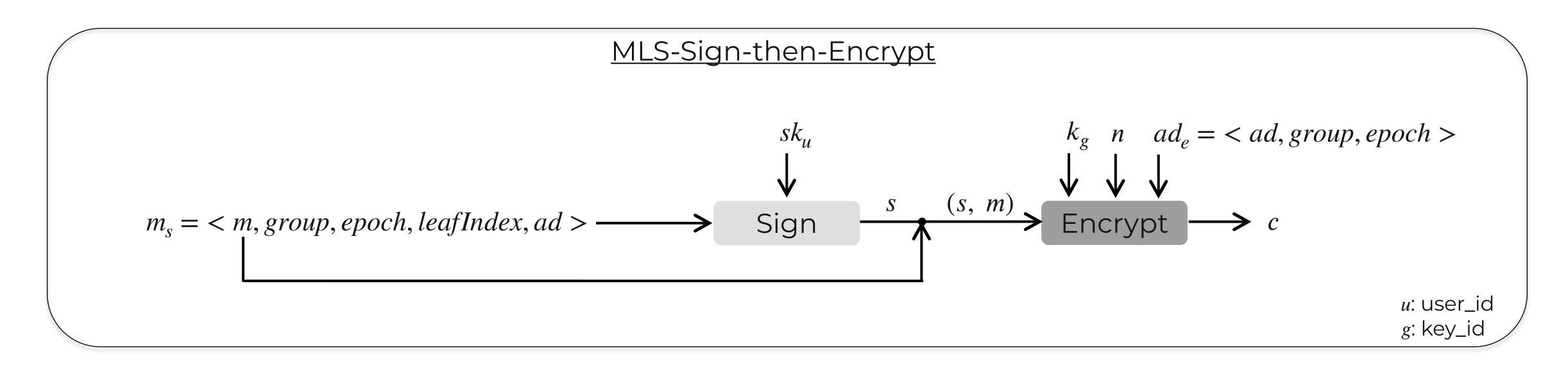


g: key_id

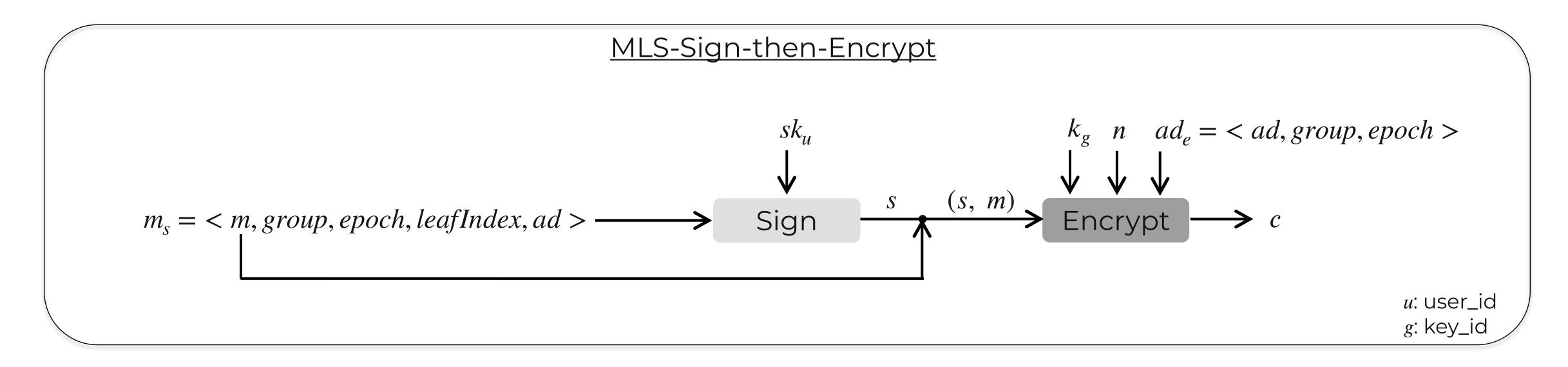
Chat encryption in MLS composes a digital signature scheme and a nonce-based encryption scheme in a Sign-then-Encrypt fashion



Chat encryption in MLS composes a digital signature scheme and a nonce-based encryption scheme in a Sign-then-Encrypt fashion

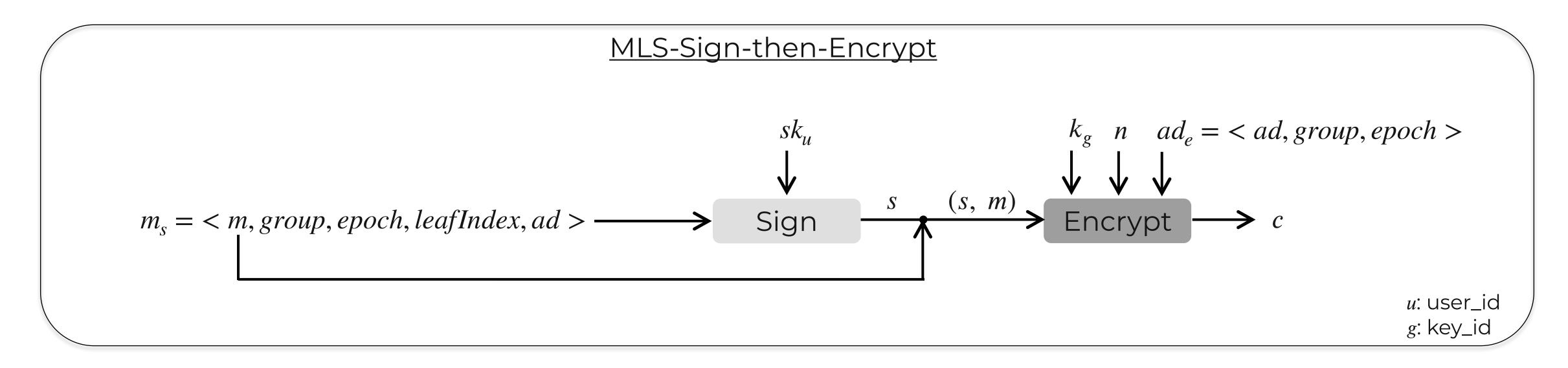


Chat encryption in MLS composes a digital signature scheme and a nonce-based encryption scheme in a Sign-then-Encrypt fashion



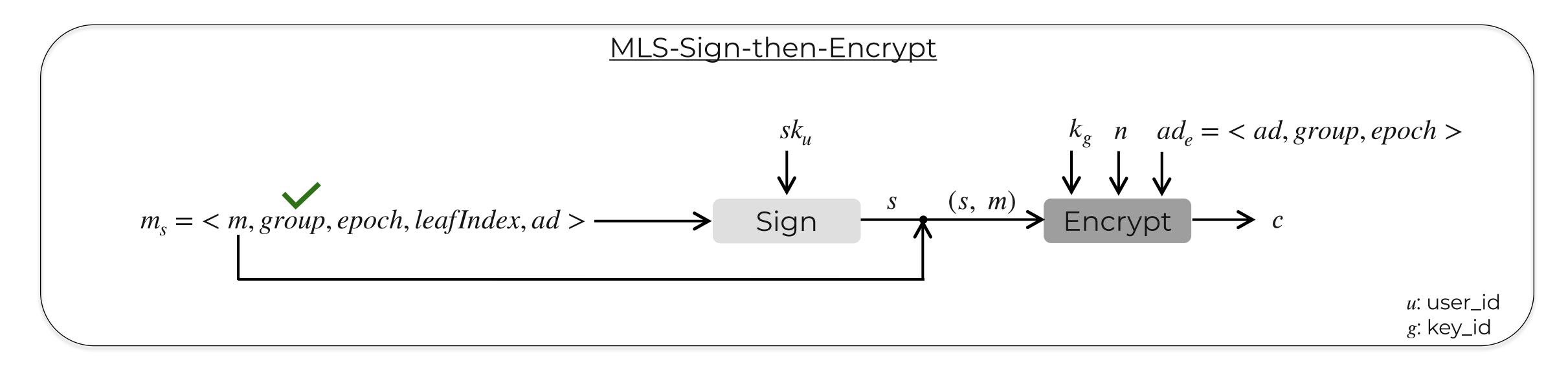
Intuition: s should authenticate the key identifier so that group insider cannot re-encrypt (s, m) using a different k and replay message to group

Chat encryption in MLS composes a digital signature scheme and a nonce-based encryption scheme in a Sign-then-Encrypt fashion



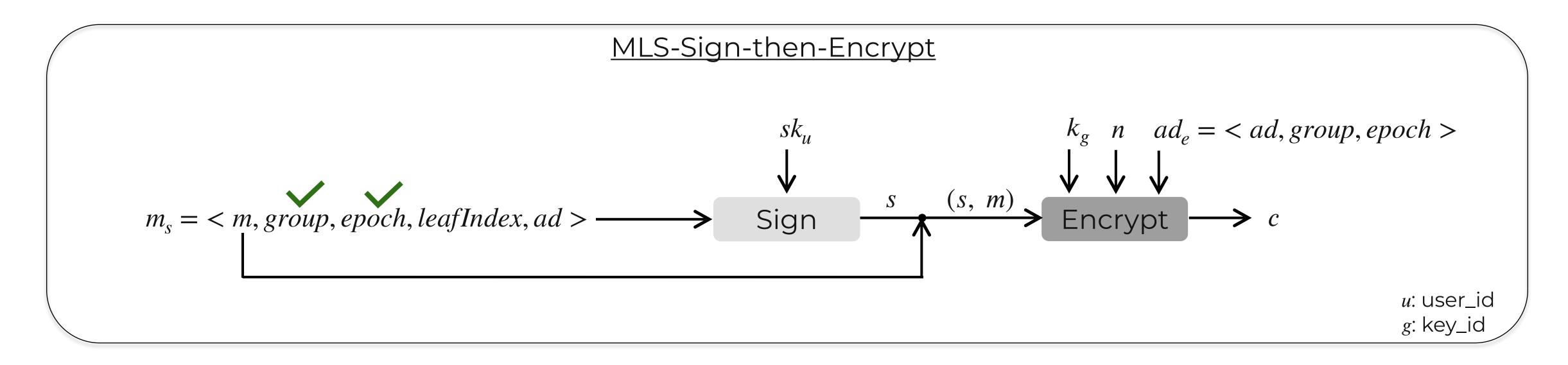
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Chat encryption in MLS composes a digital signature scheme and a nonce-based encryption scheme in a Sign-then-Encrypt fashion



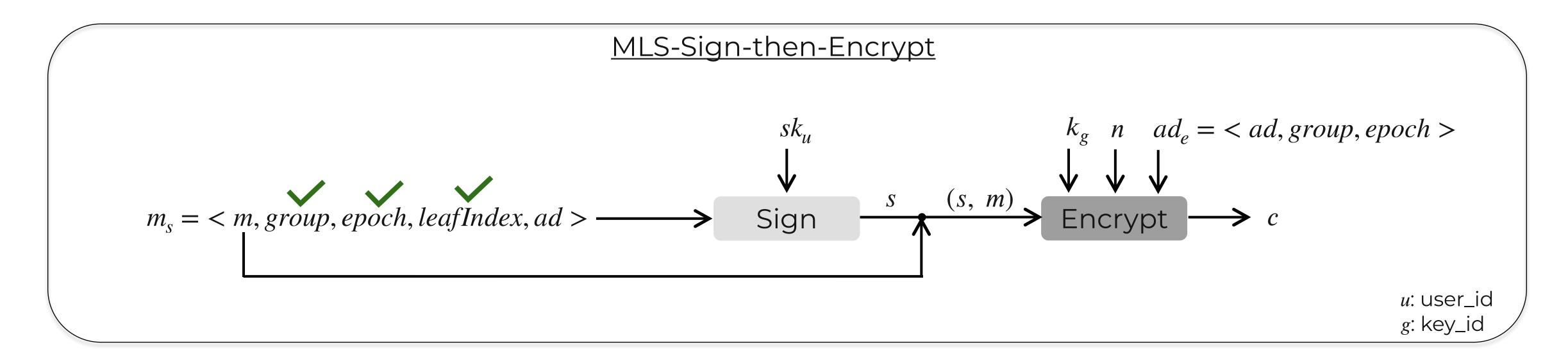
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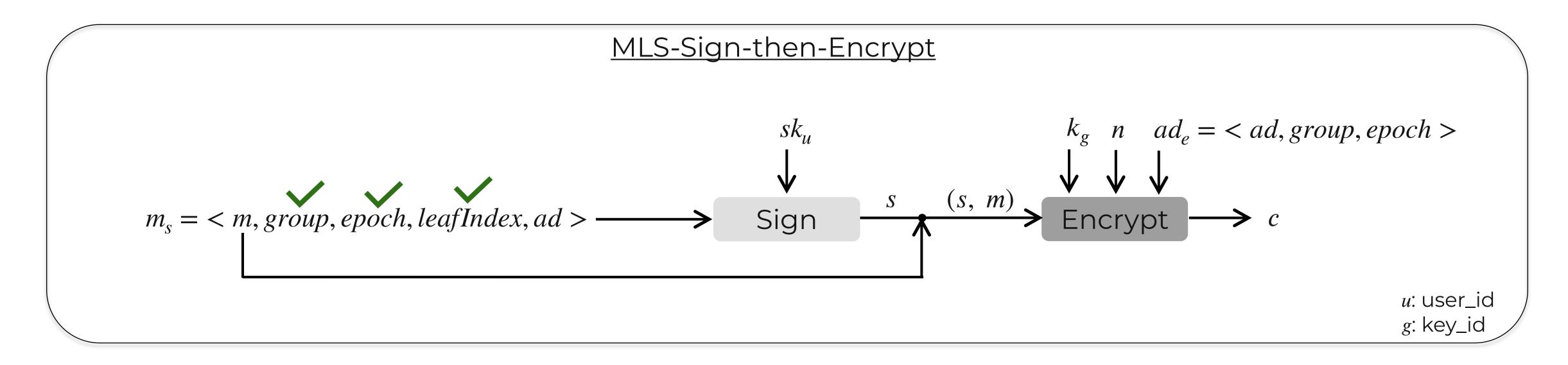
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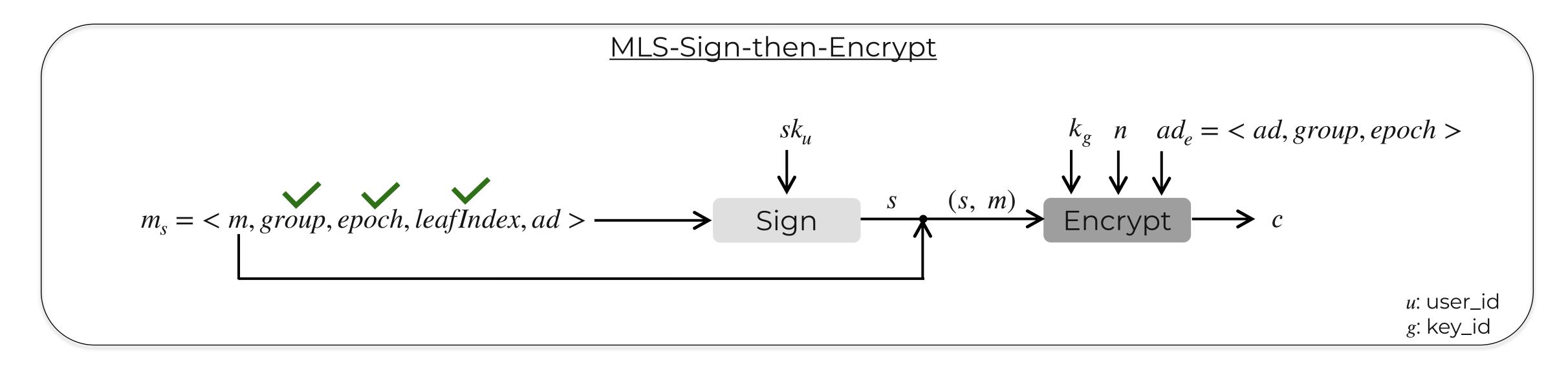
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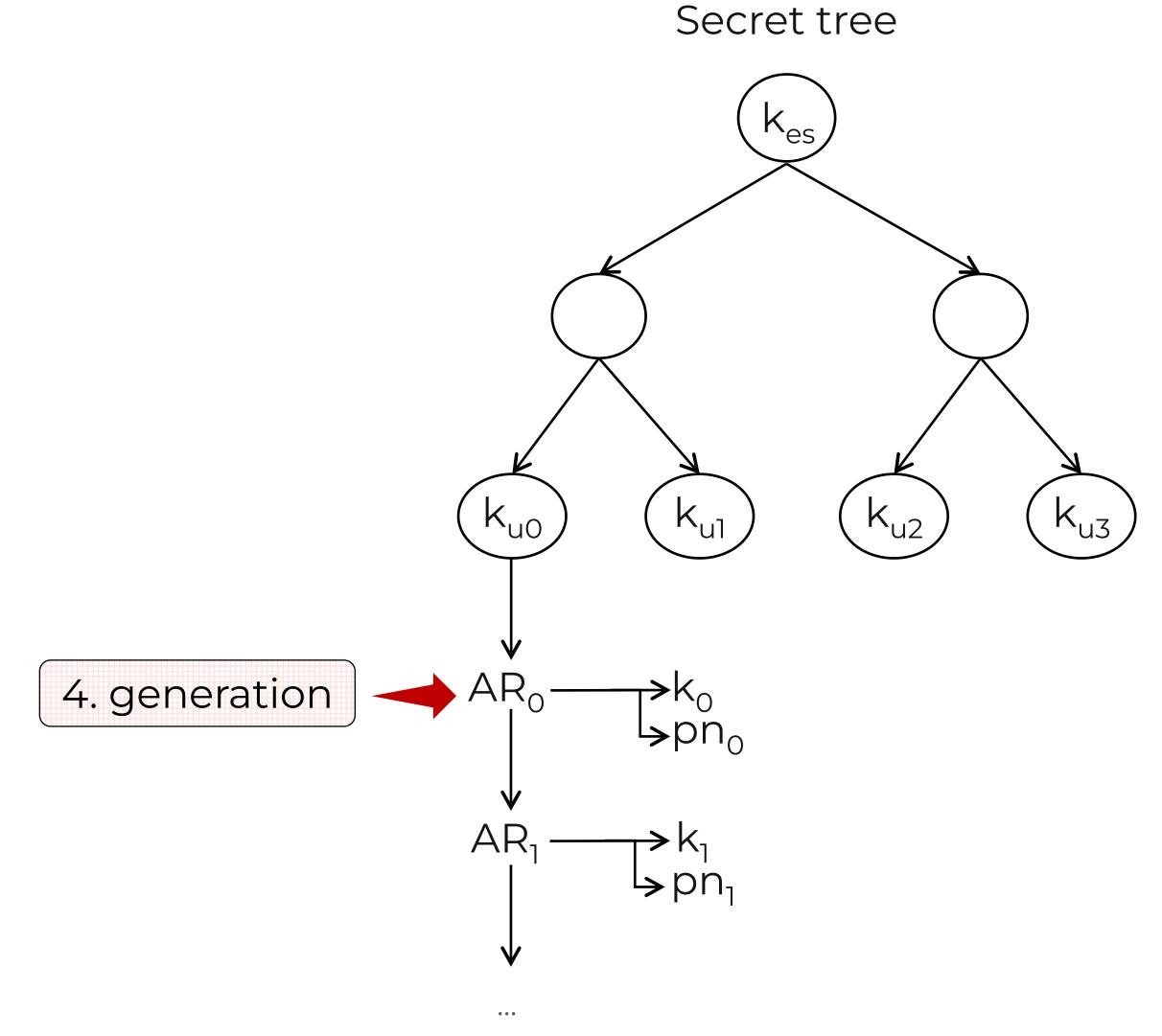
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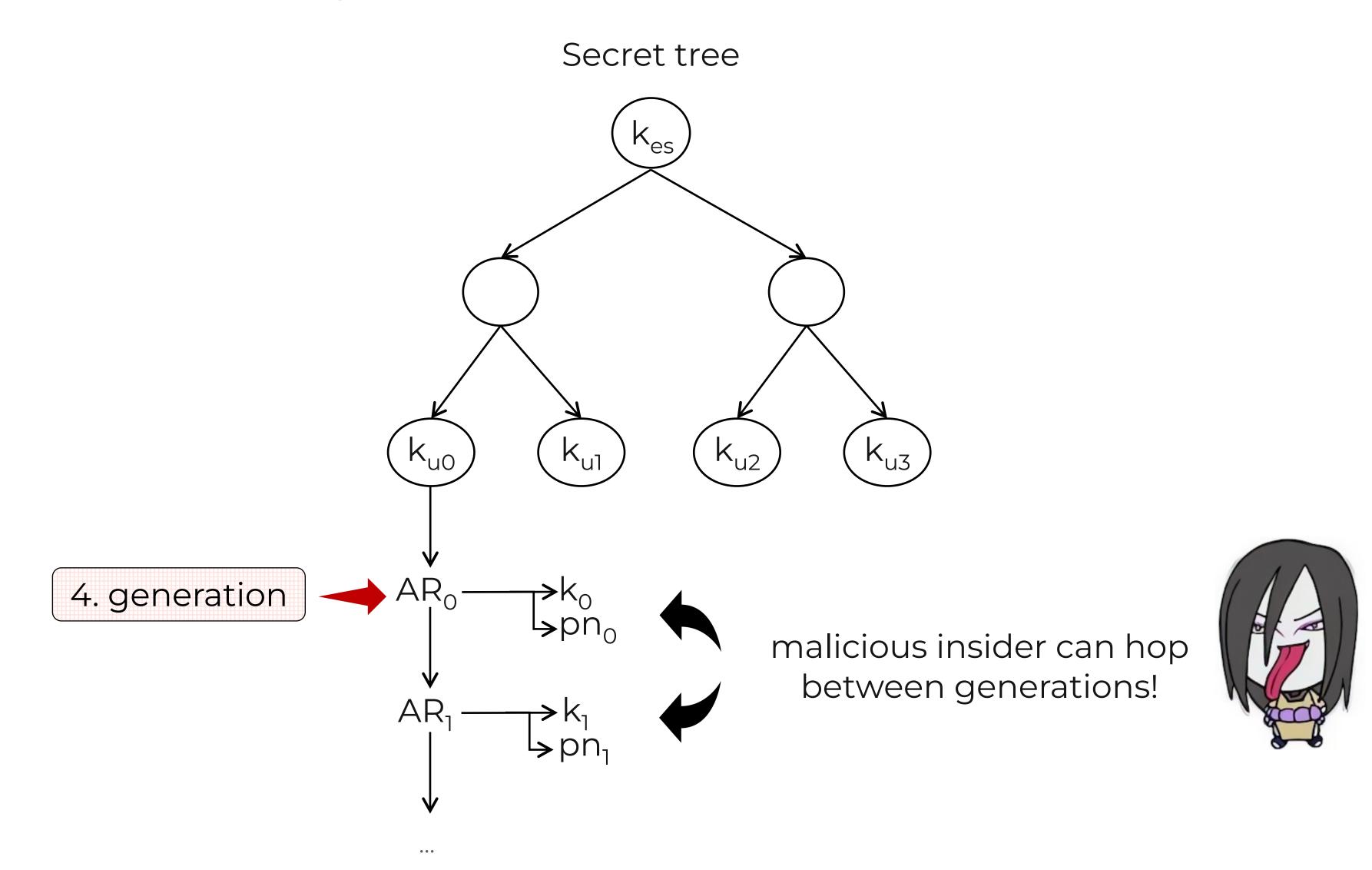


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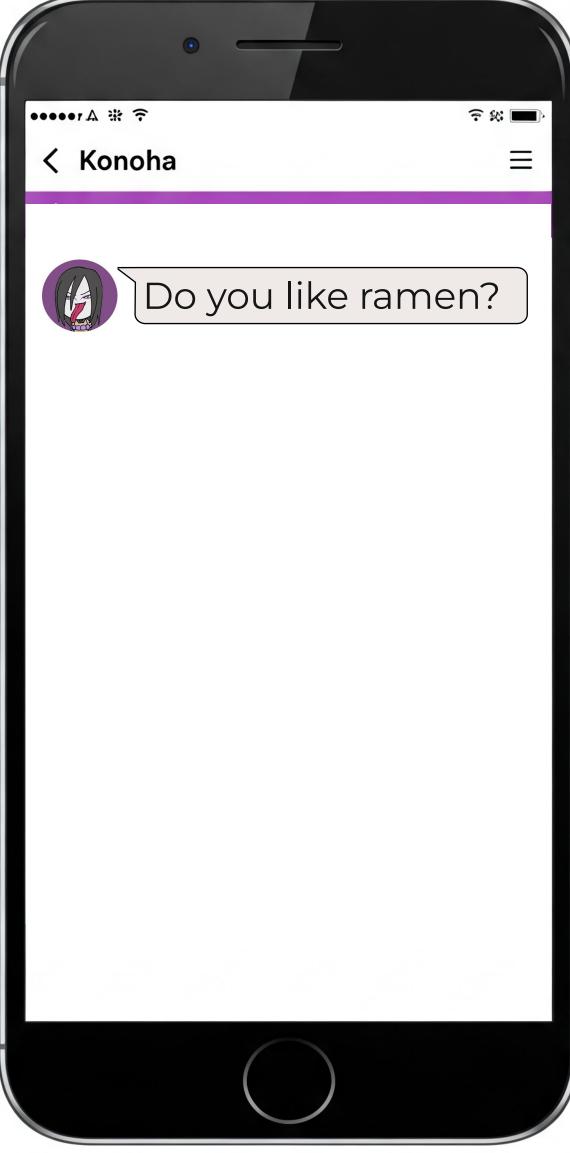
Recall: key identifier g = (group, epoch, leafIndex, generation)

group_key_id X





SIGN group_key_id 💢





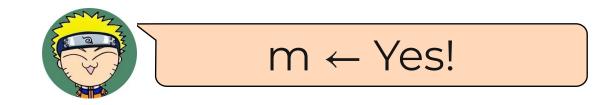


























m ← Yes!

MLS-Sign-then-Encrypt $s \leftarrow Sign(sk_u, m_s)$ $c \leftarrow Enc(k_{g_0}, n, s | | m, ad_e)$











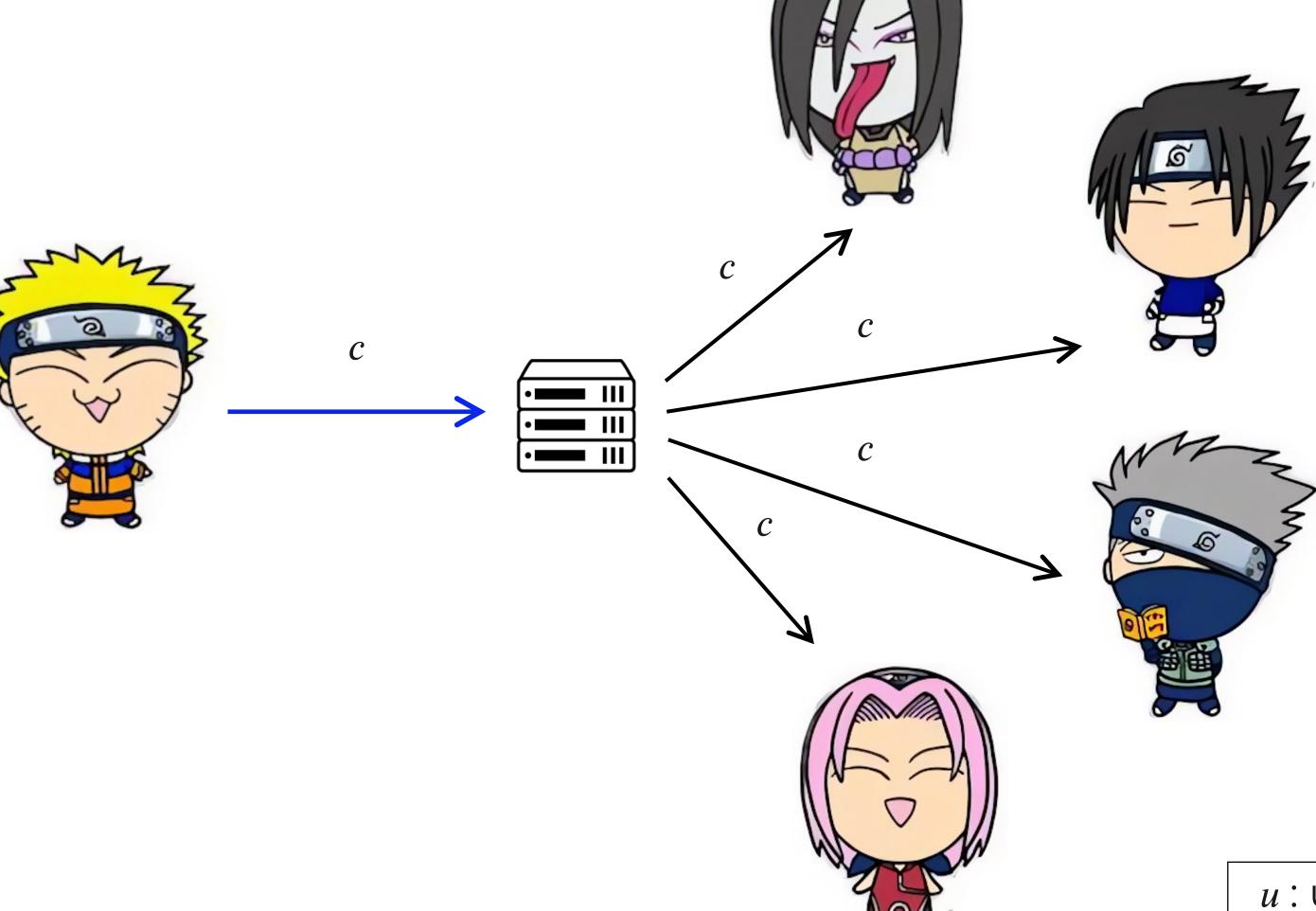


u: user_id g_0 : key_id

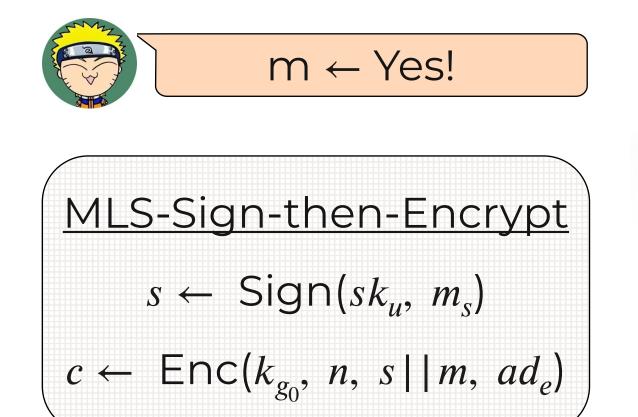


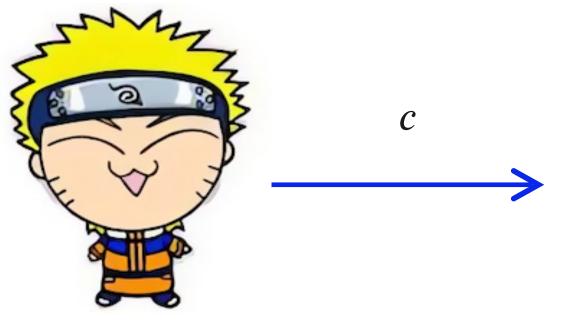
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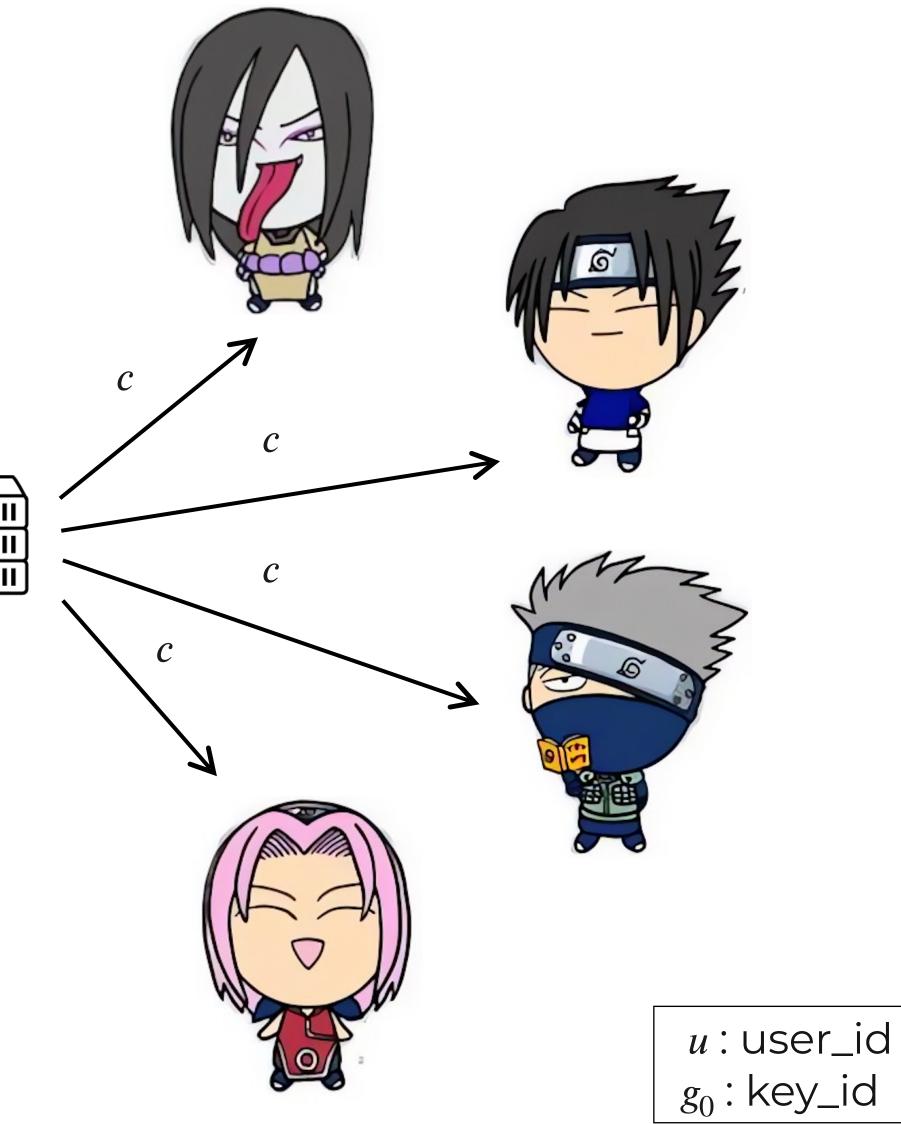
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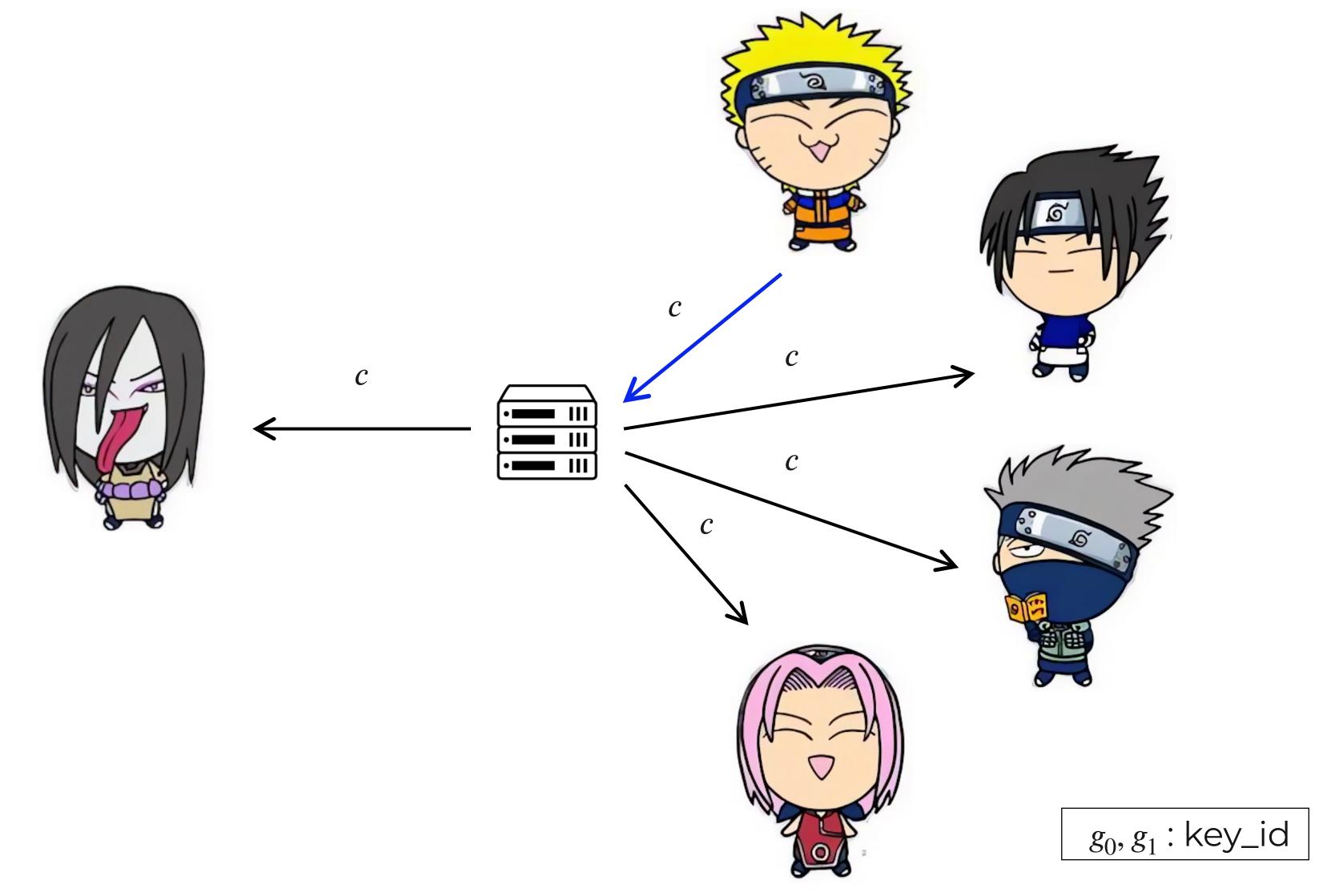
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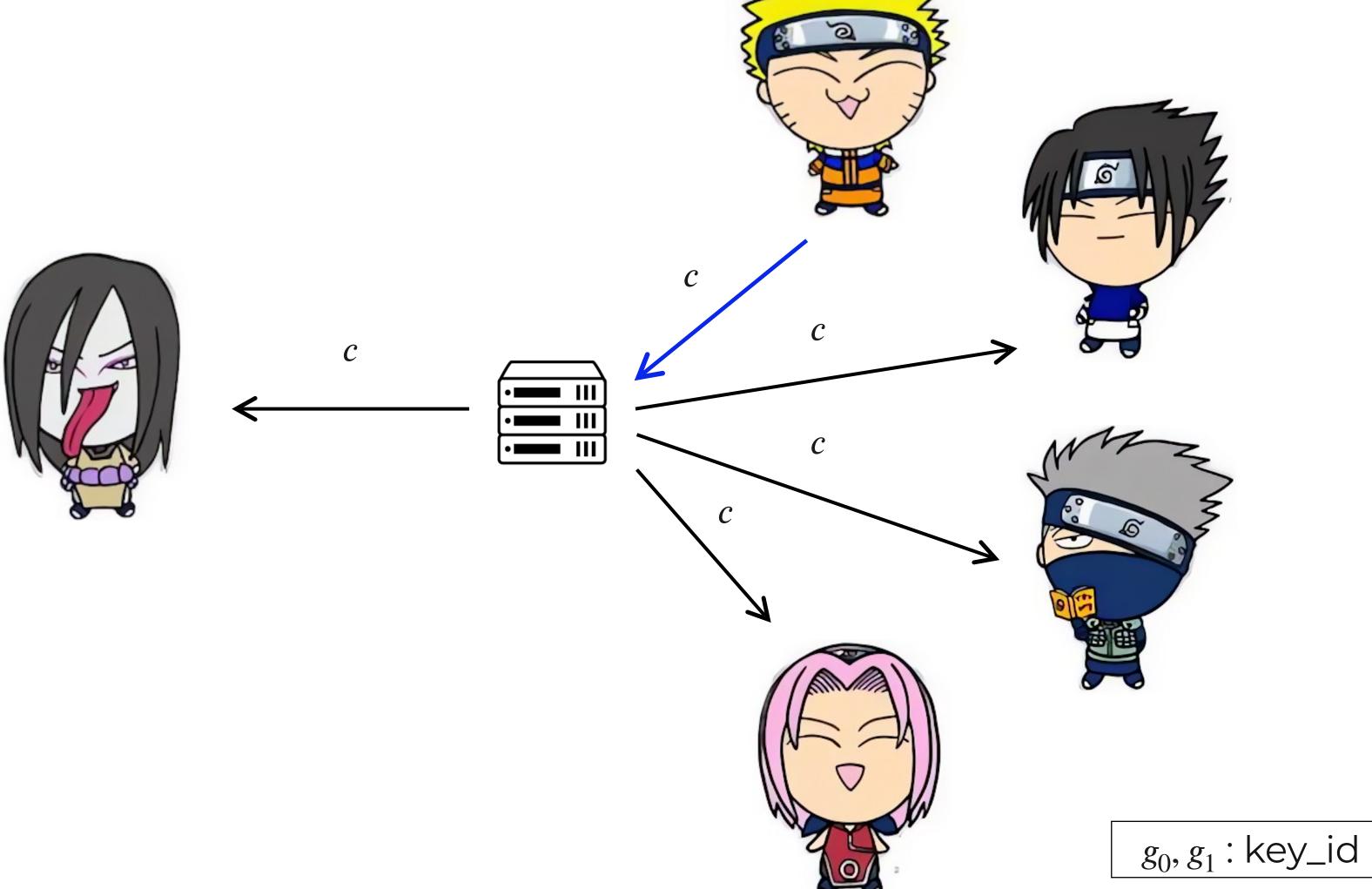




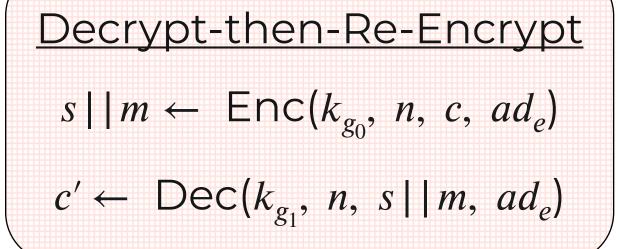




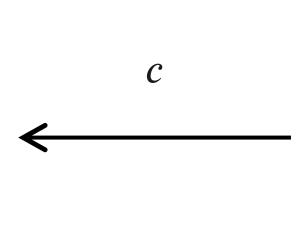
Recall: signature s does not authenticate the generation

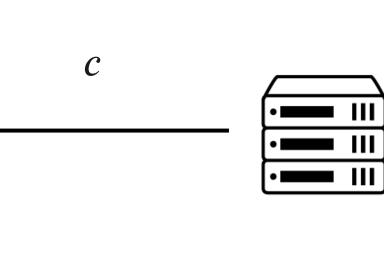


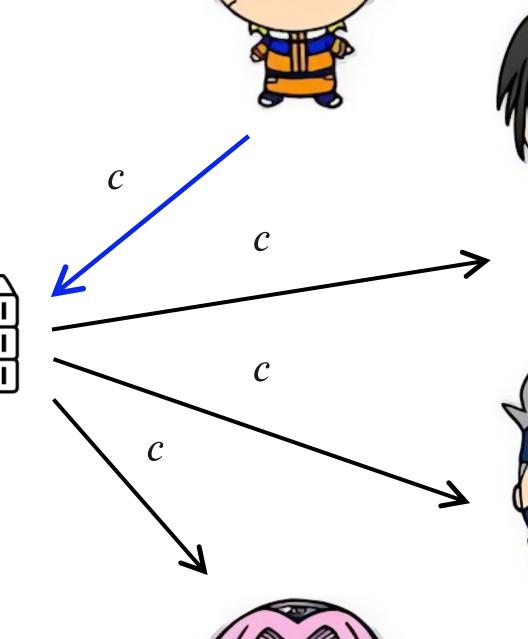
Recall: signature s does not authenticate the generation







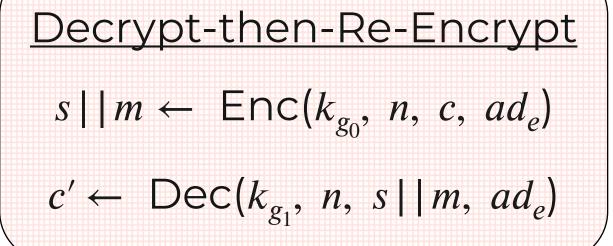




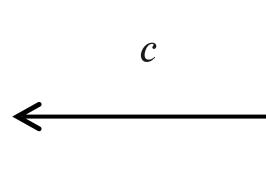


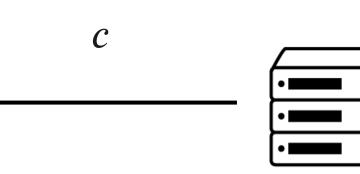
 g_0, g_1 : key_id

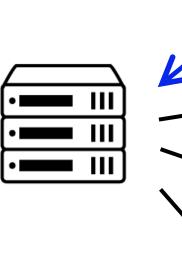
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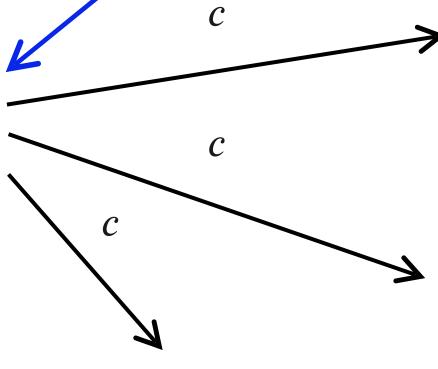




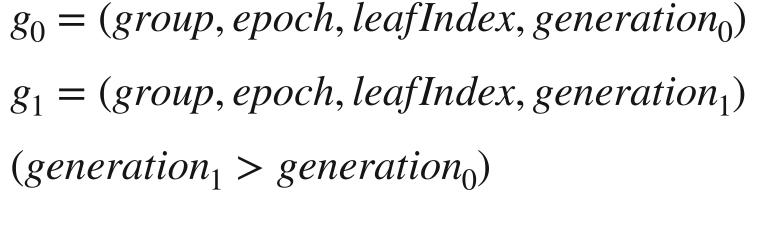






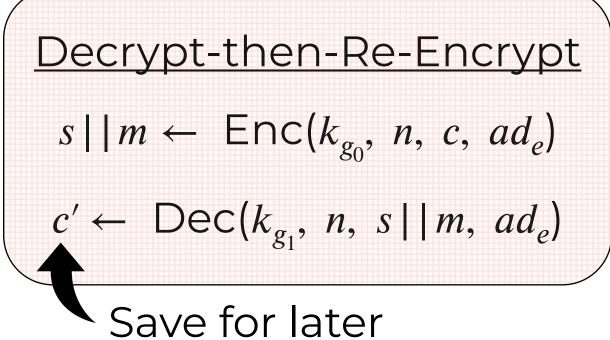




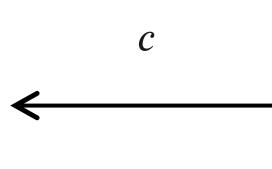


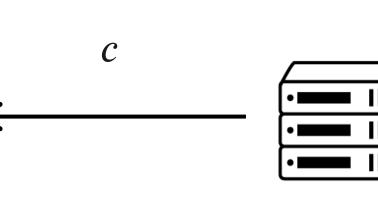


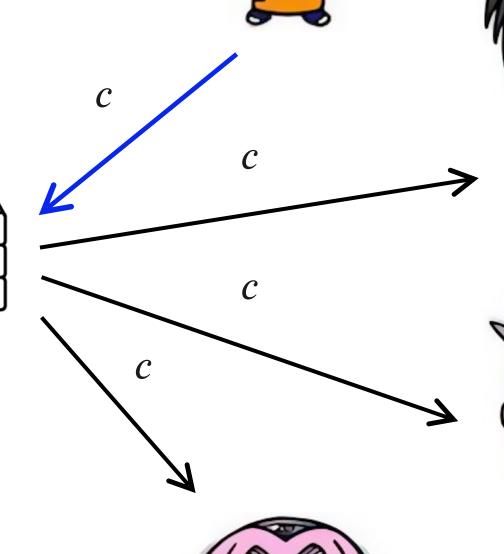
Recall: signature s does not authenticate the generation



















 $g_0 = (group, epoch, leafIndex, generation_0)$

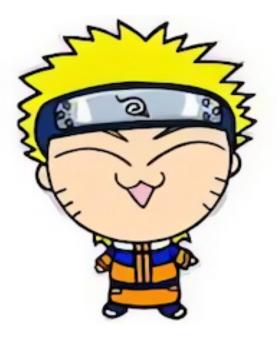
 $g_1 = (group, epoch, leafIndex, generation_1)$

 $(generation_1 > generation_0)$





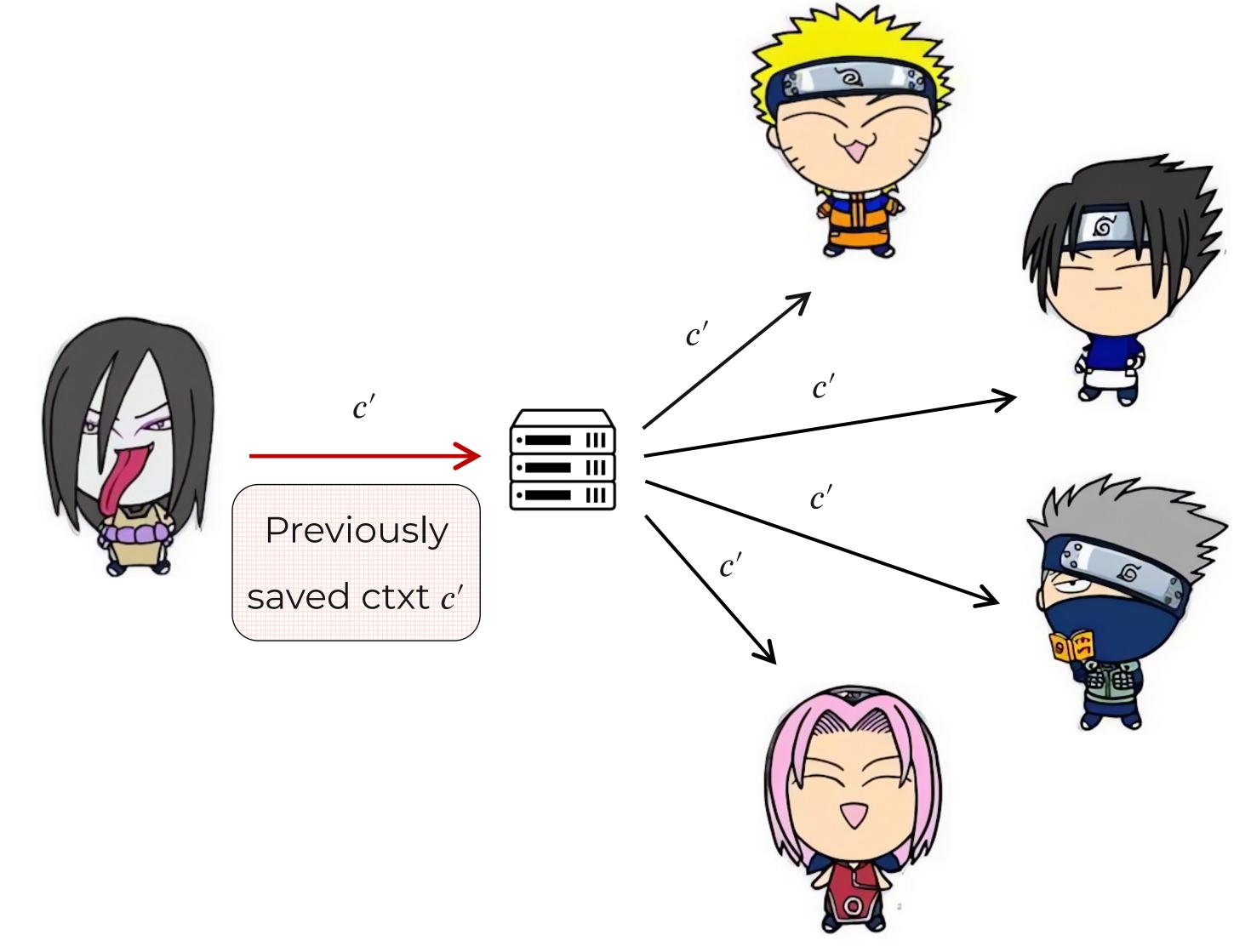


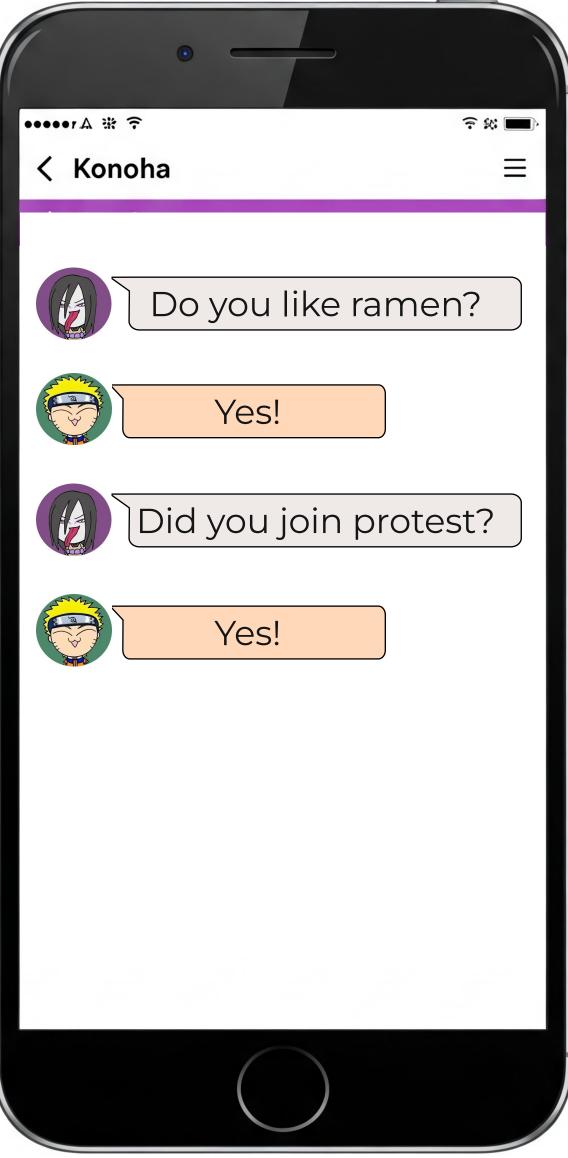


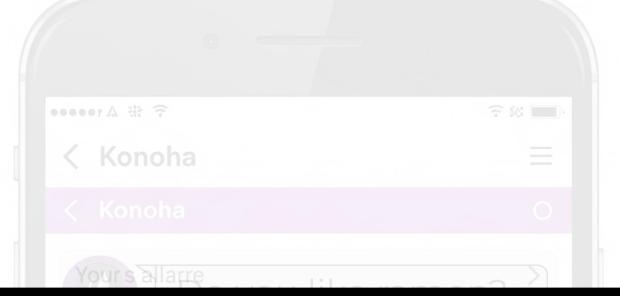












MLS aims to protect against forgeries by group members (aka insiders)

"[Knowledge] of the AEAD keys allows the attacker to send an encrypted message using that key, but cannot send a message to a group which appears to be from any valid client since they cannot forge the signature."

Attack results from lack of binding between signature and generation; mitigation is to bind them

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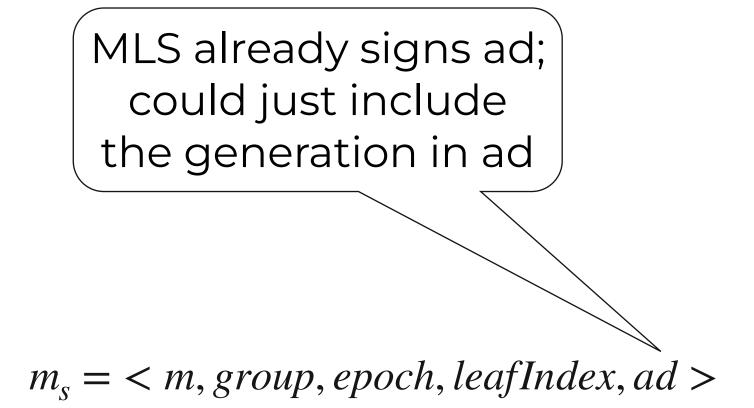
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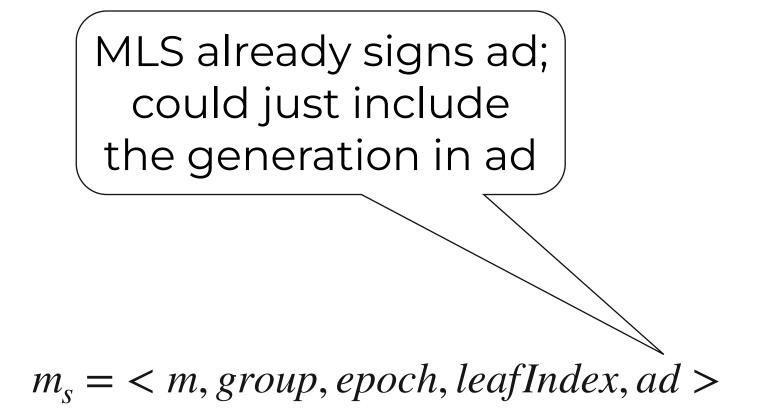
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Attack results from lack of binding between signature and generation; mitigation is to bind them

No Protection against Replay by Insiders

MLS does not protect against one group member replaying a PrivateMessage sent by another group member within the same epoch that the message was originally sent. Similarly, MLS does not protect against the replay (by a group member or otherwise) of a PublicMessage within the same epoch that the message was originally sent. Applications for whom replay is an important risk should apply mitigations at the application layer, as discussed below.

In addition to the risks discussed in {{symmetric-key-compromise}}, an attacker with access to the Ratchet Secrets for an endpoint can replay PrivateMessage objects sent by other members of the group by taking the signed content of the message and re-encrypting it with a new generation of the original sender's ratchet. If the other members of the group interpret a message with a new generation as a fresh message, then this message will appear fresh. (This is possible because the message signature does not cover the generation field of the message.) Messages sent as PublicMessage objects similarly lack replay protections. There is no message counter comparable to the generation field in PrivateMessage.



Applications can detect replay by including a unique identifier for the message (e.g., a counter) in either the message payload or the authenticated_data field, both of which are included in the signatures for PublicMessage and PrivateMessage.

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Case Study II: Session



MLS Insider Replay



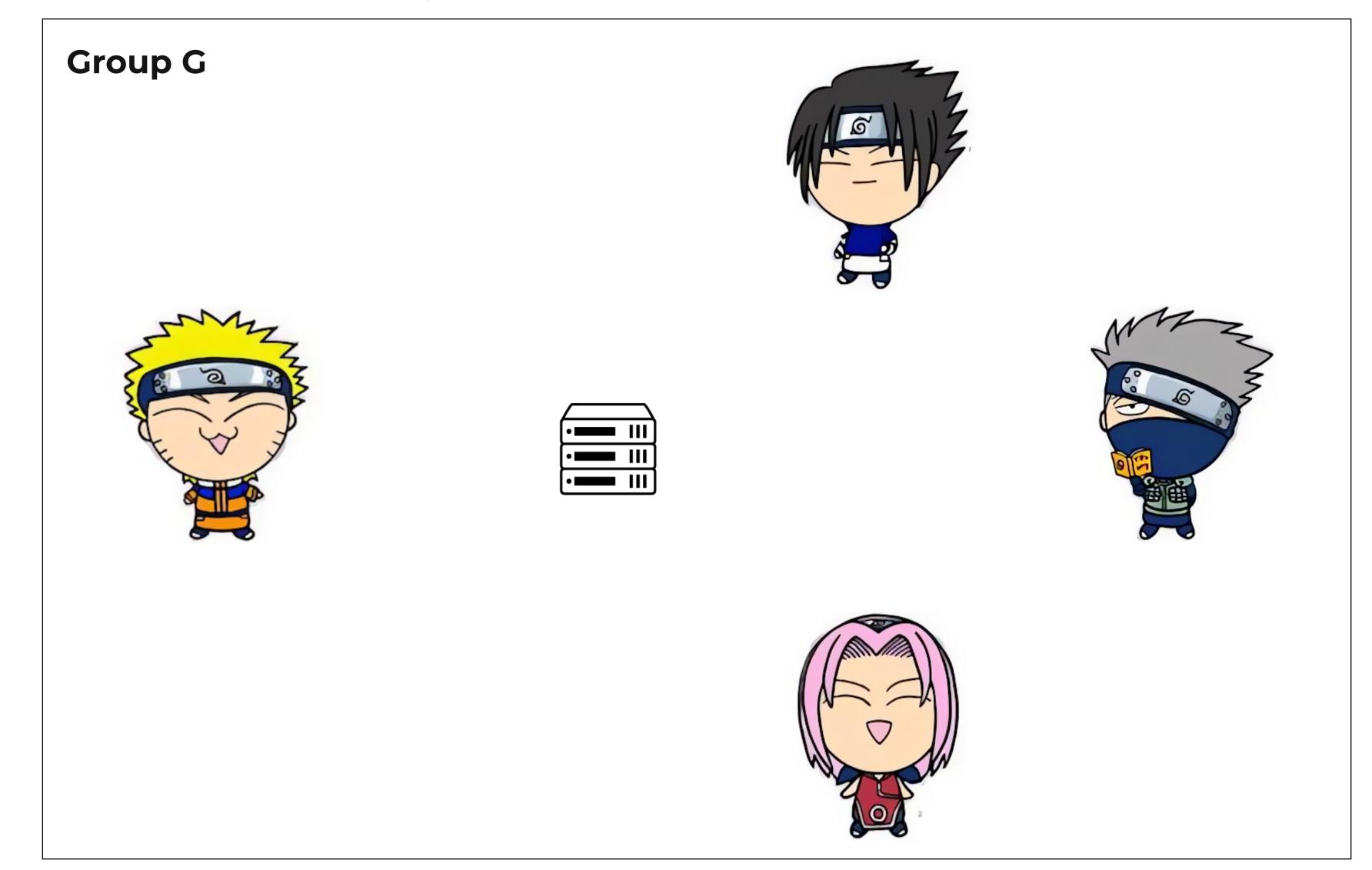
Session Insider Replay

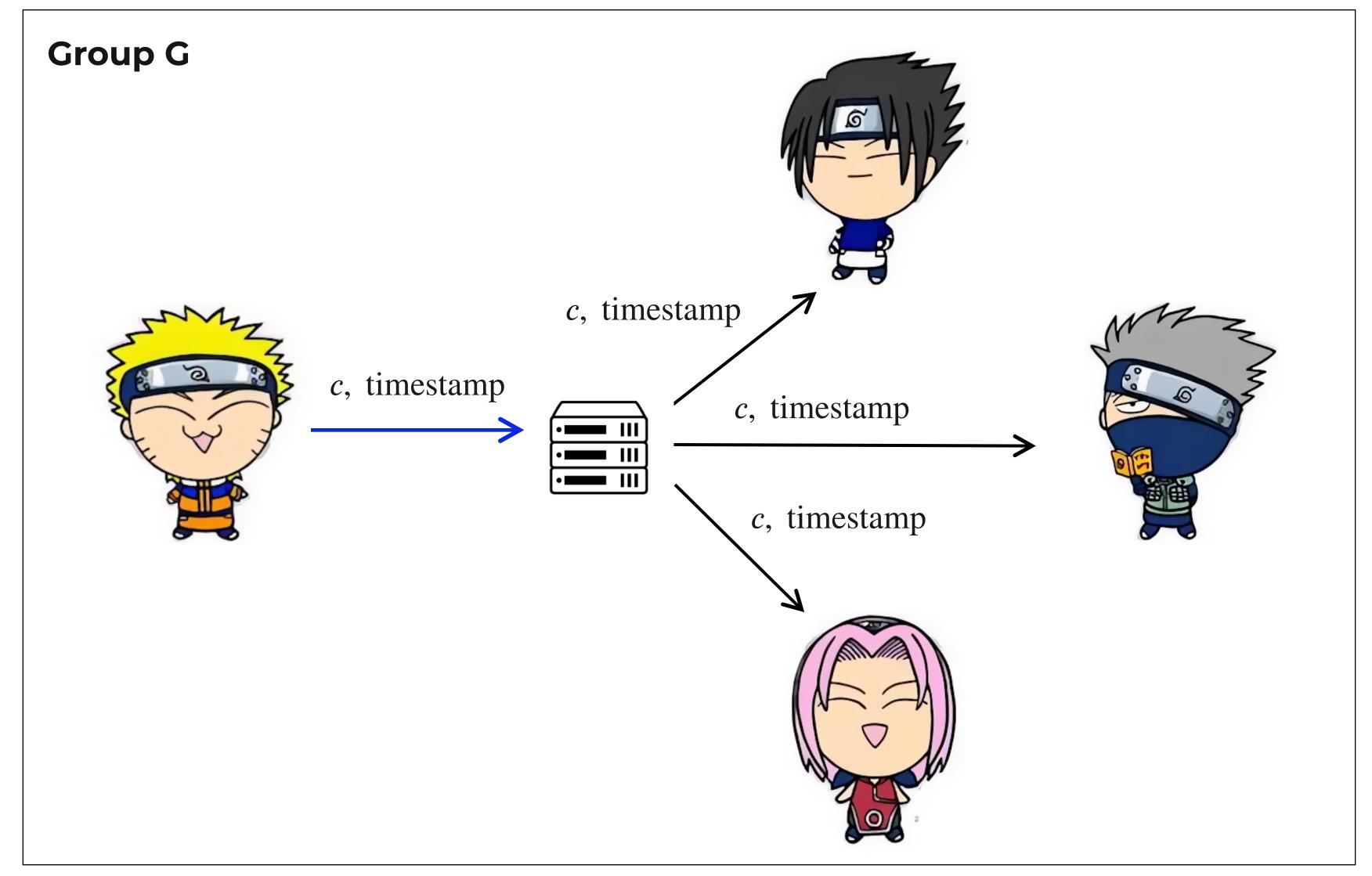
MLS Insider Replay

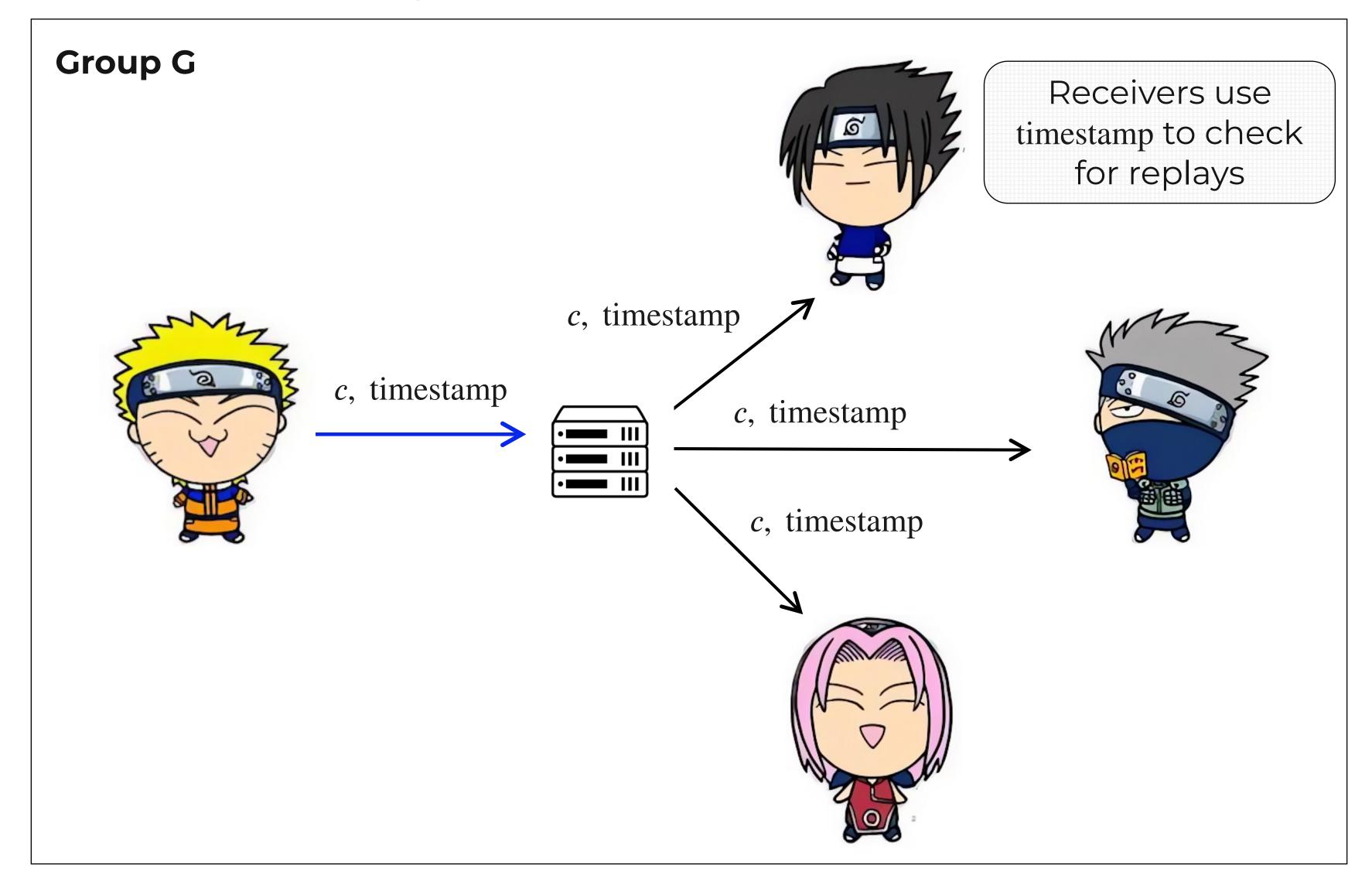


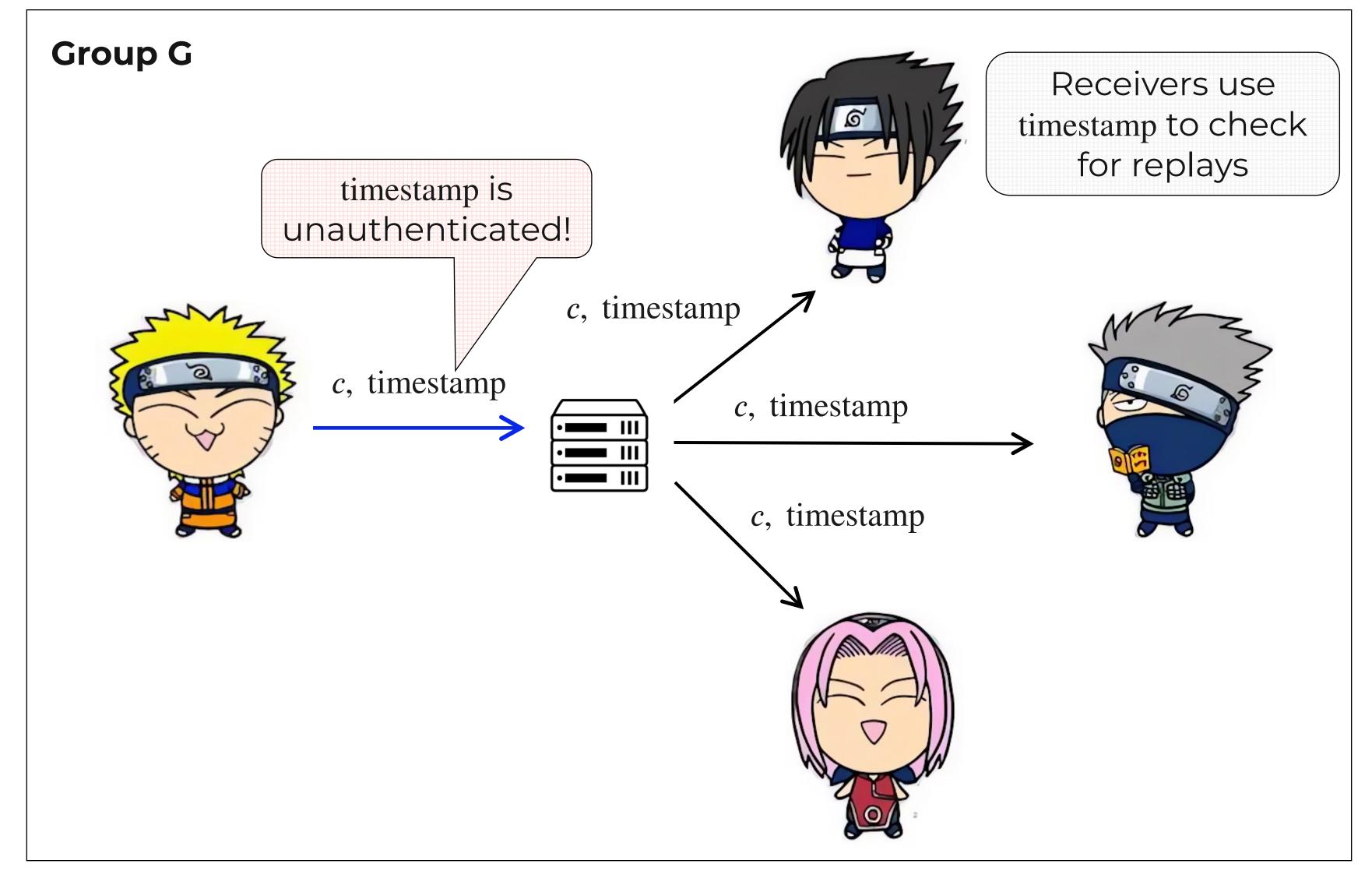
Session Insider Replay

SIGN group_key_id 💢









Group G



Receivers use timestamp to check for replays

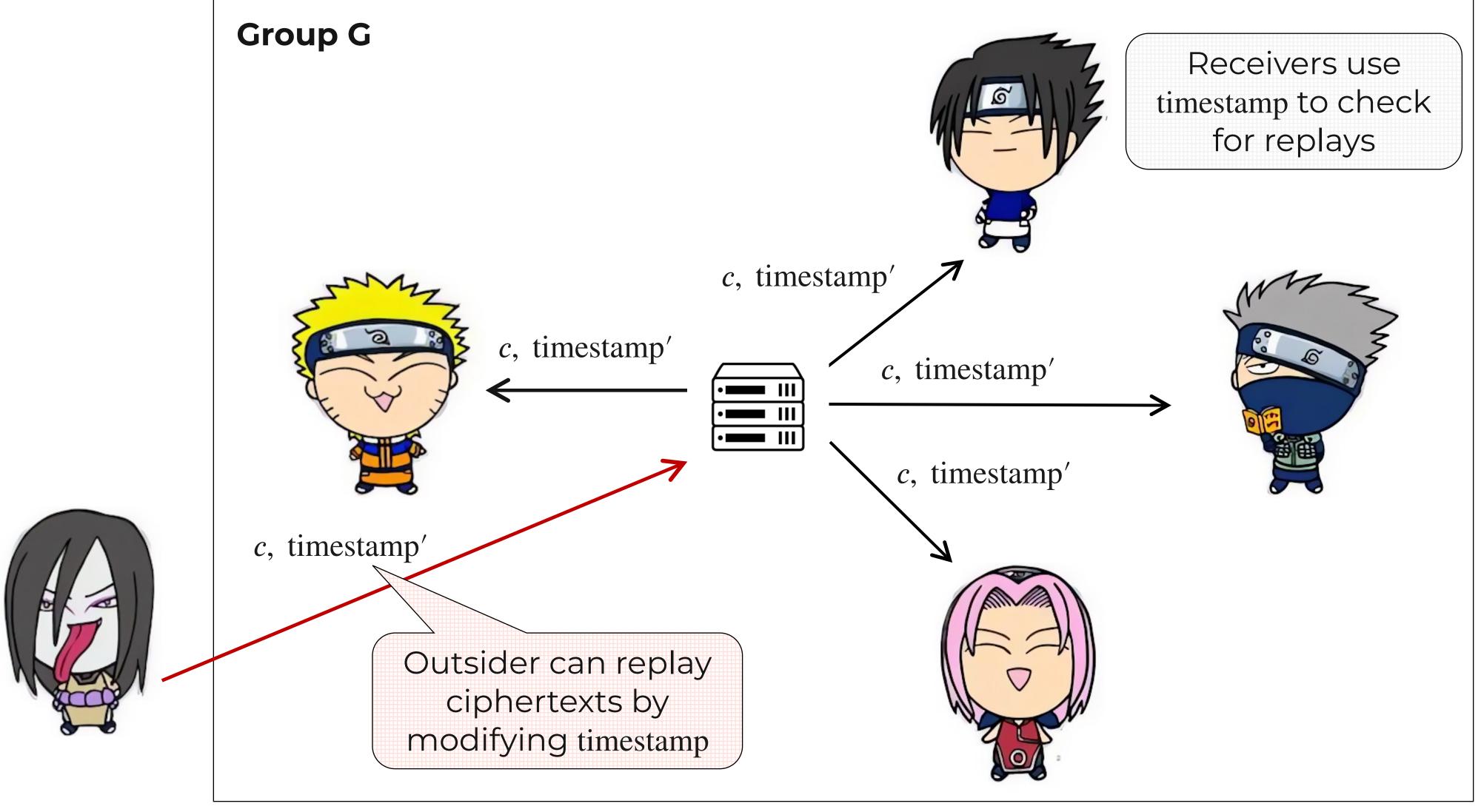


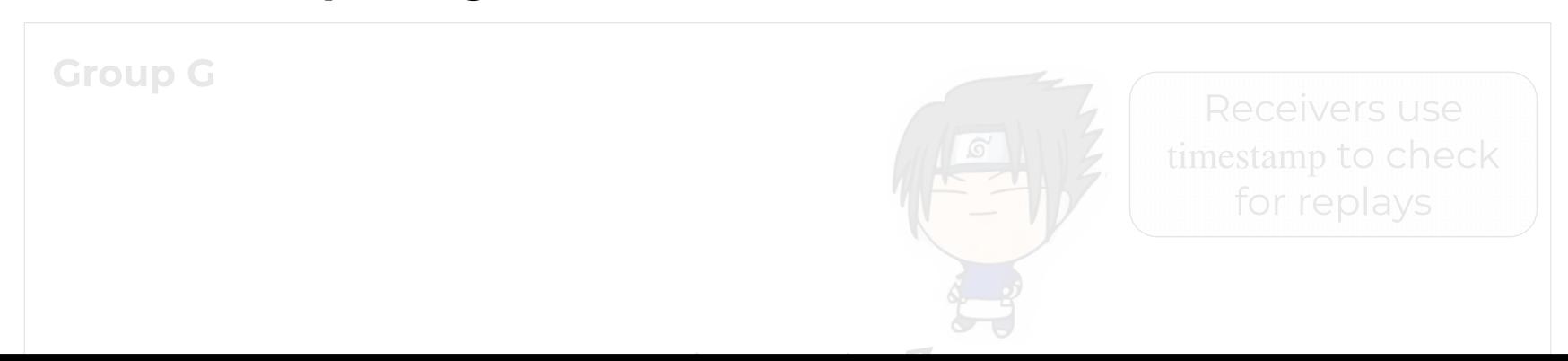




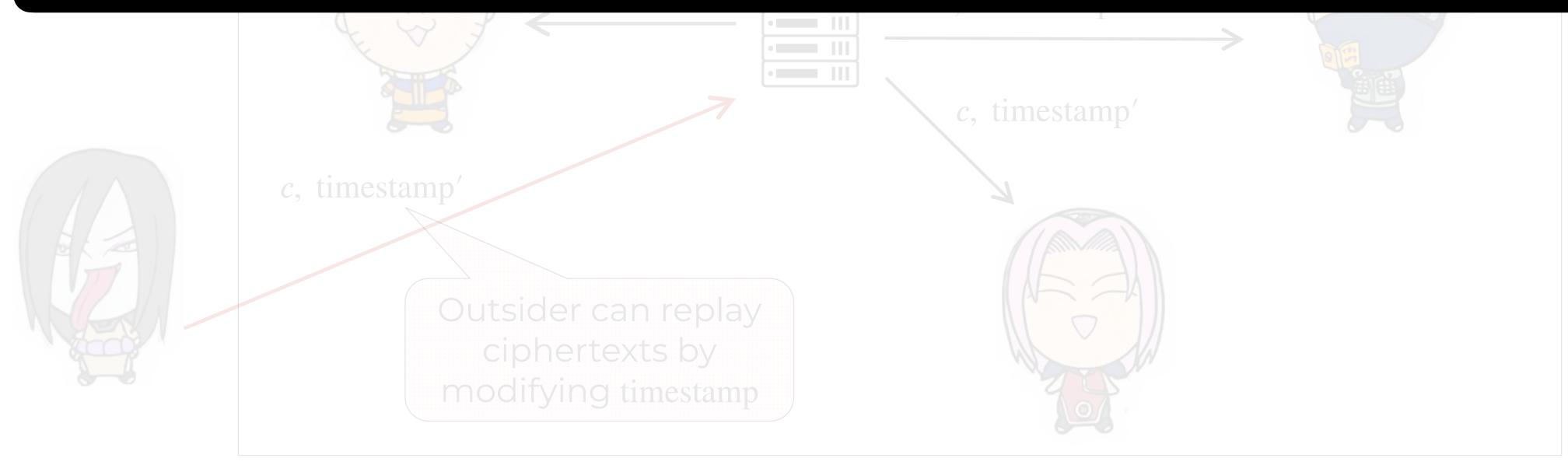








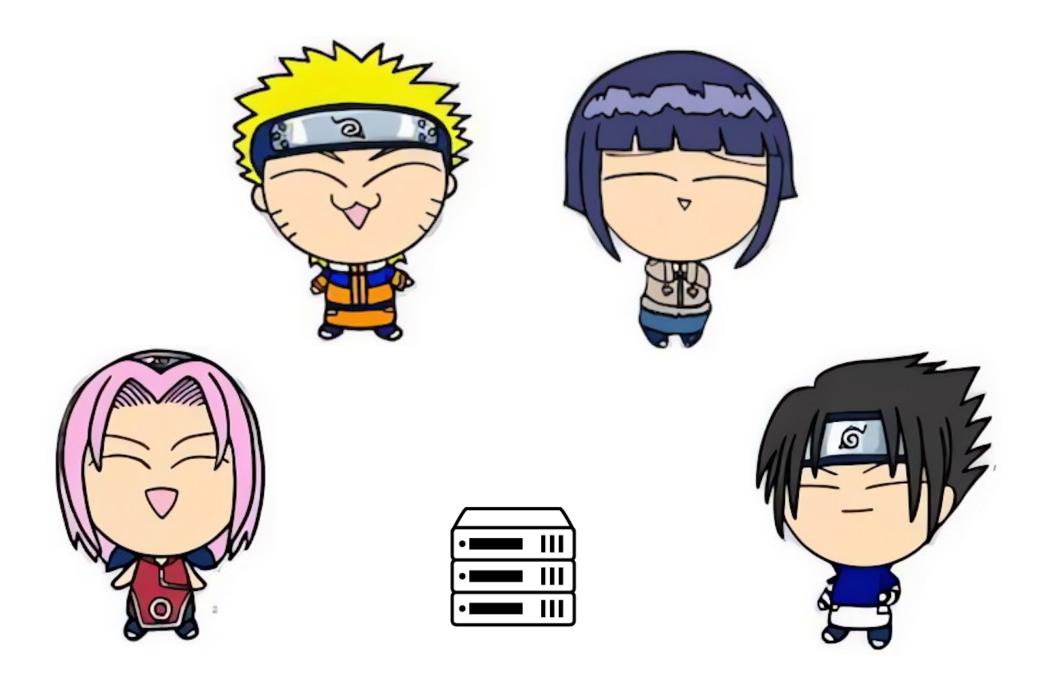
At the time of analysis, Session used the LegacyGroups protocol -- has since migrated to GroupsV2



Case Study III: Keybase



Keybase Encryption Key Derivation

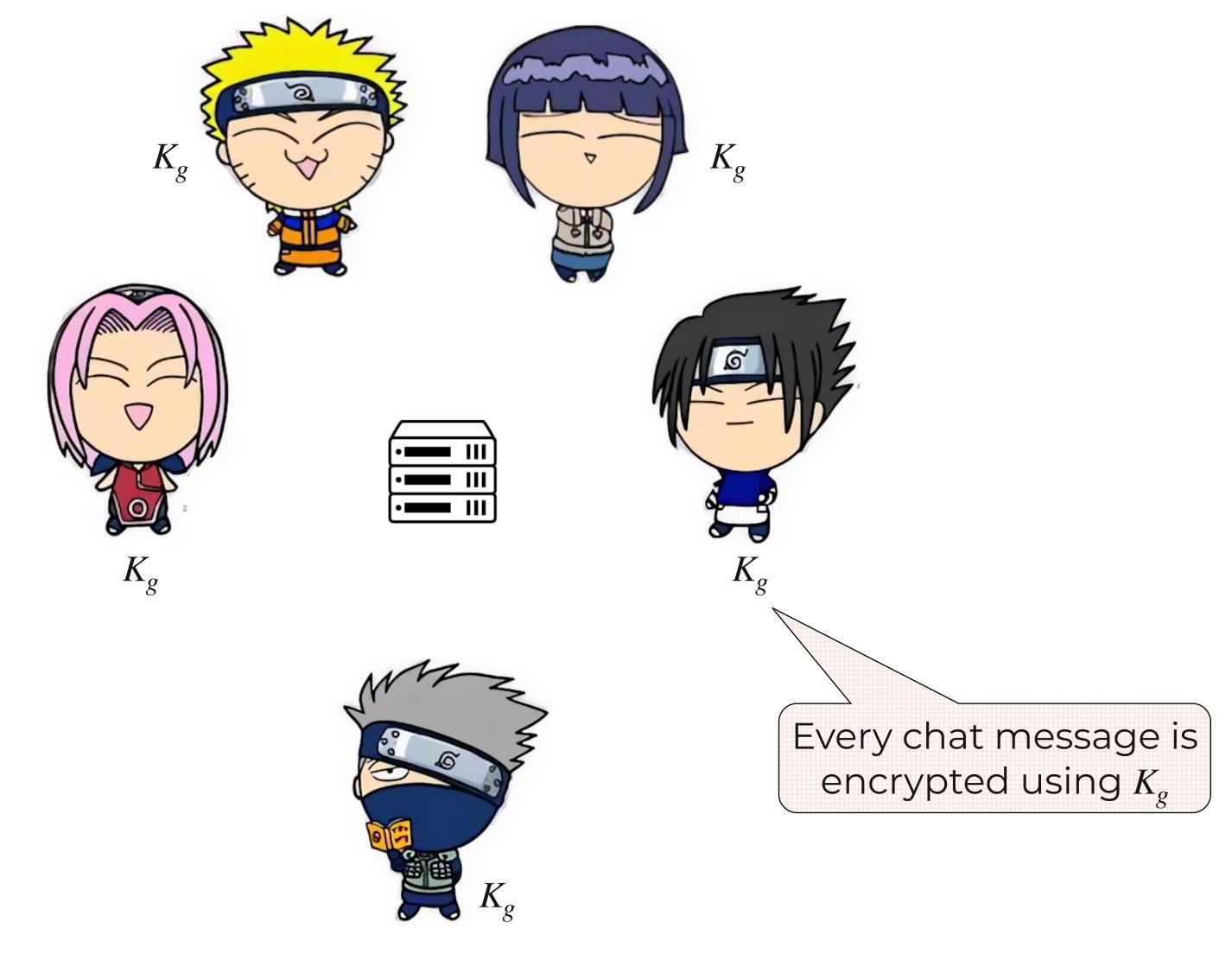




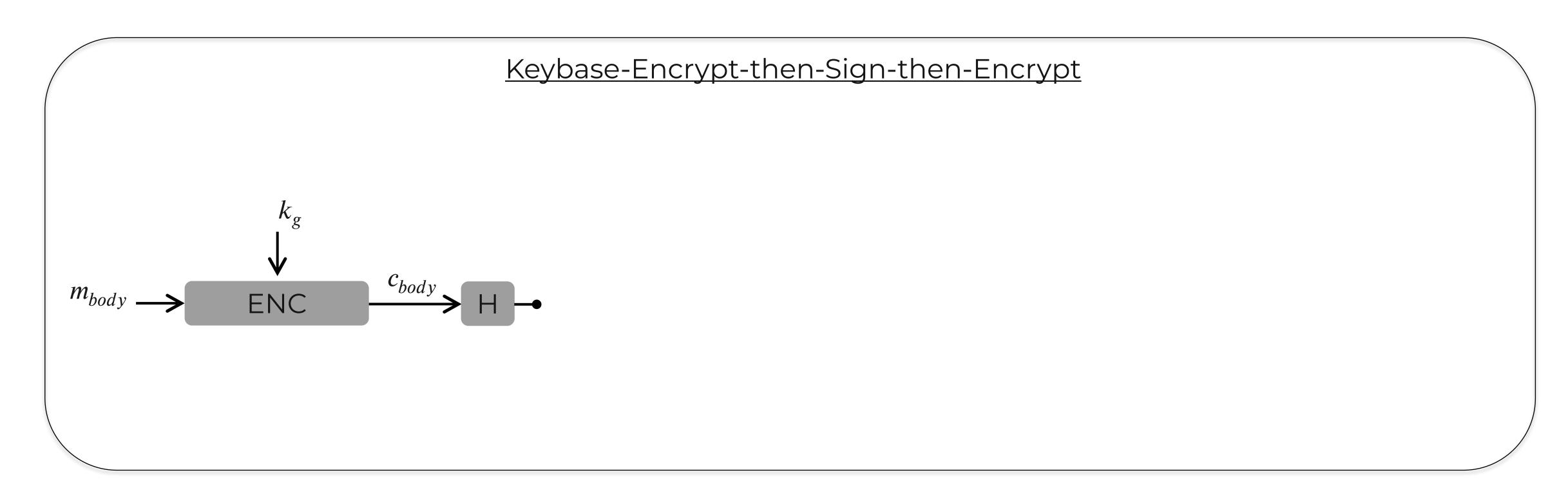
Keybase Encryption Key Derivation



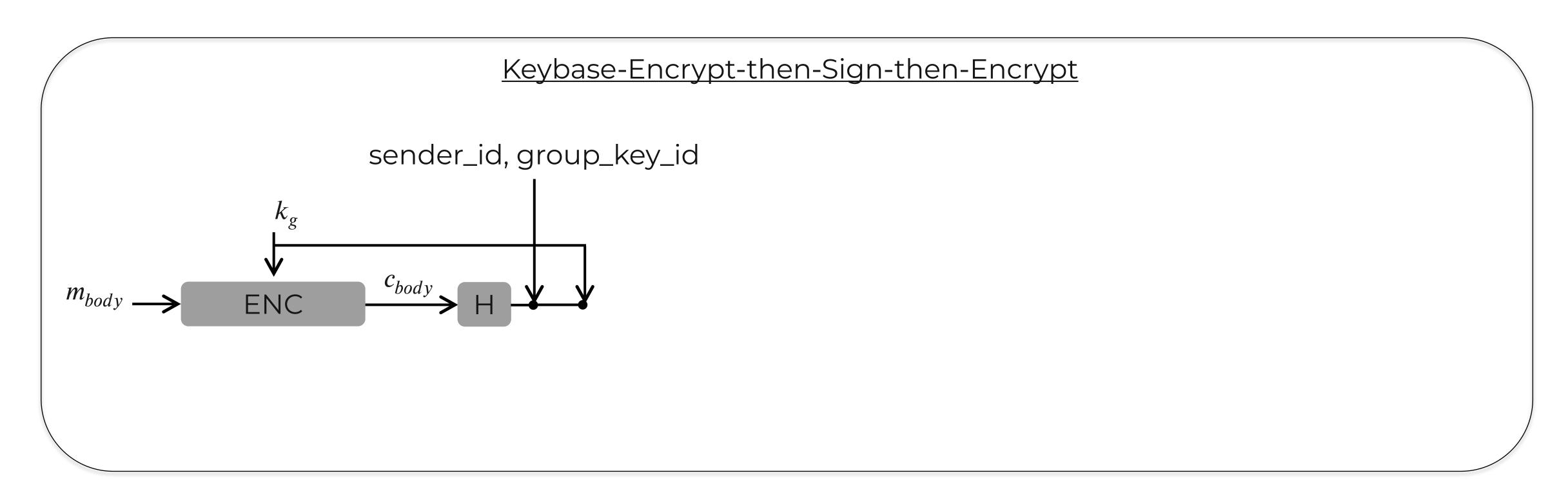
Keybase Encryption Key Derivation

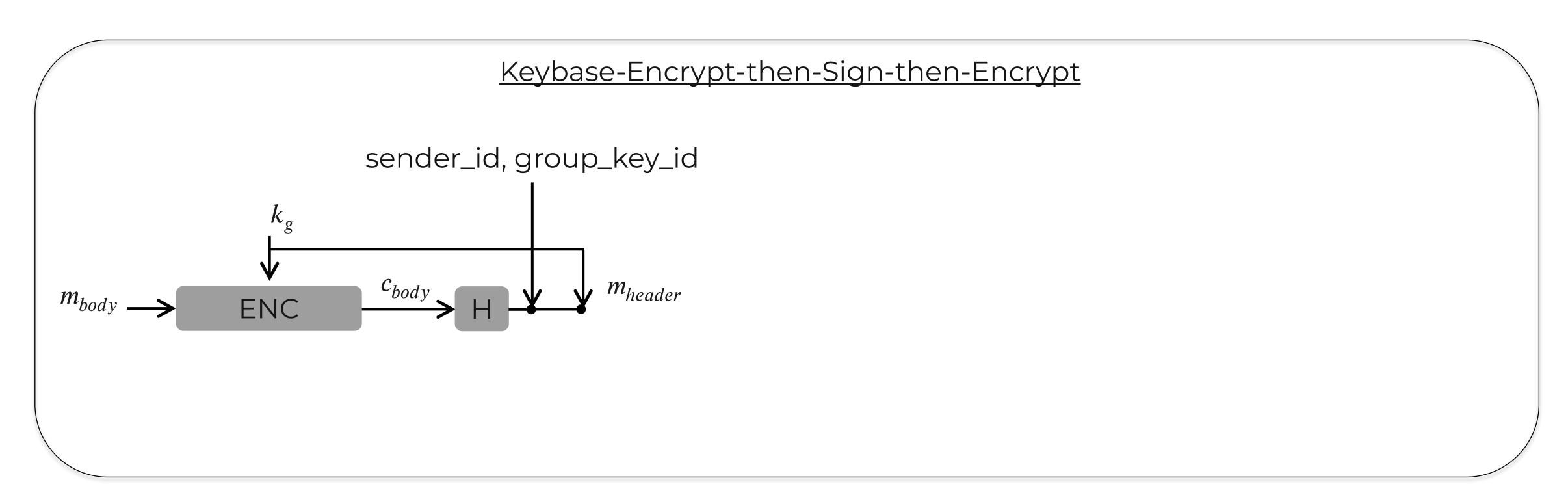


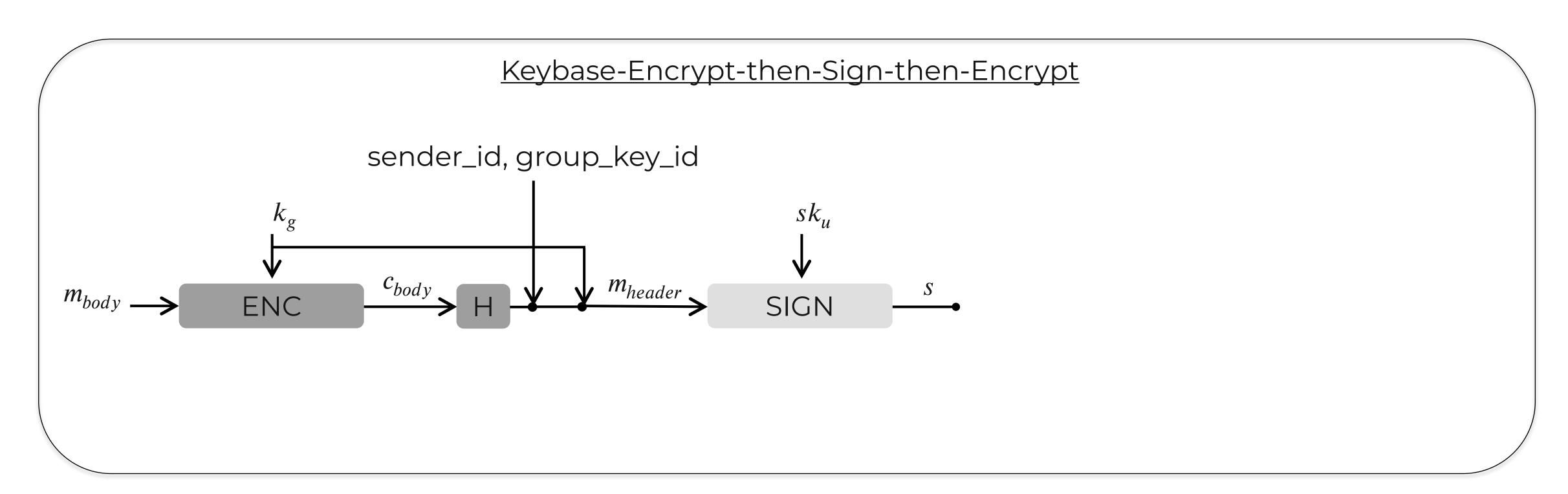
Keybase-Encrypt-then-Sign-then-Encrypt

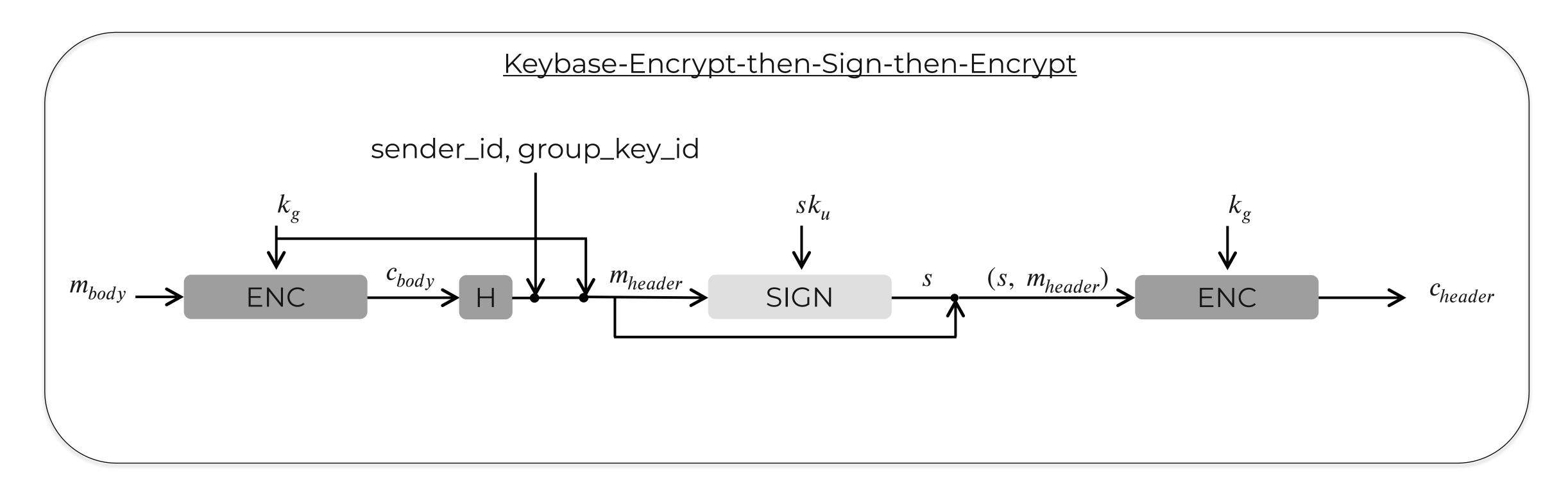


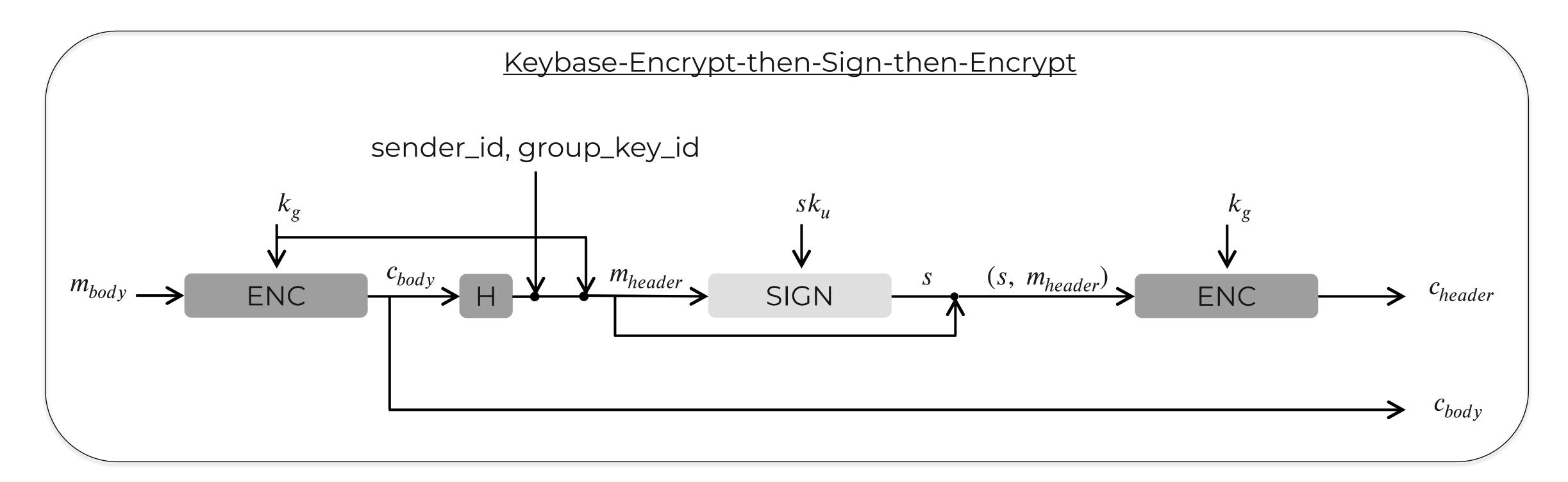


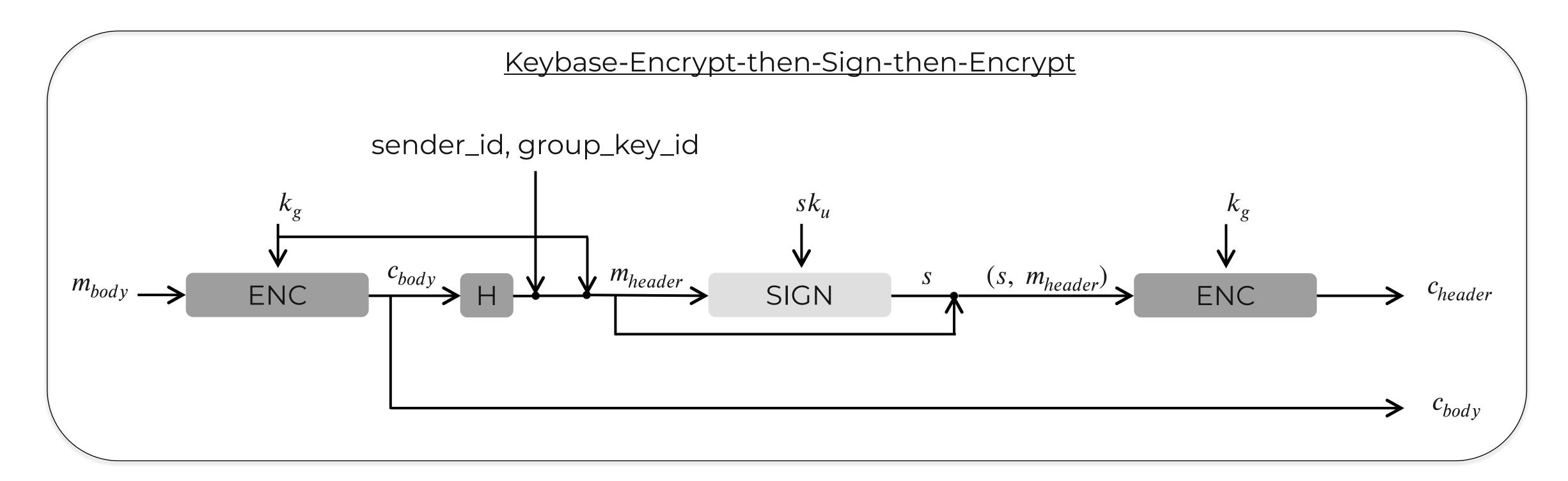


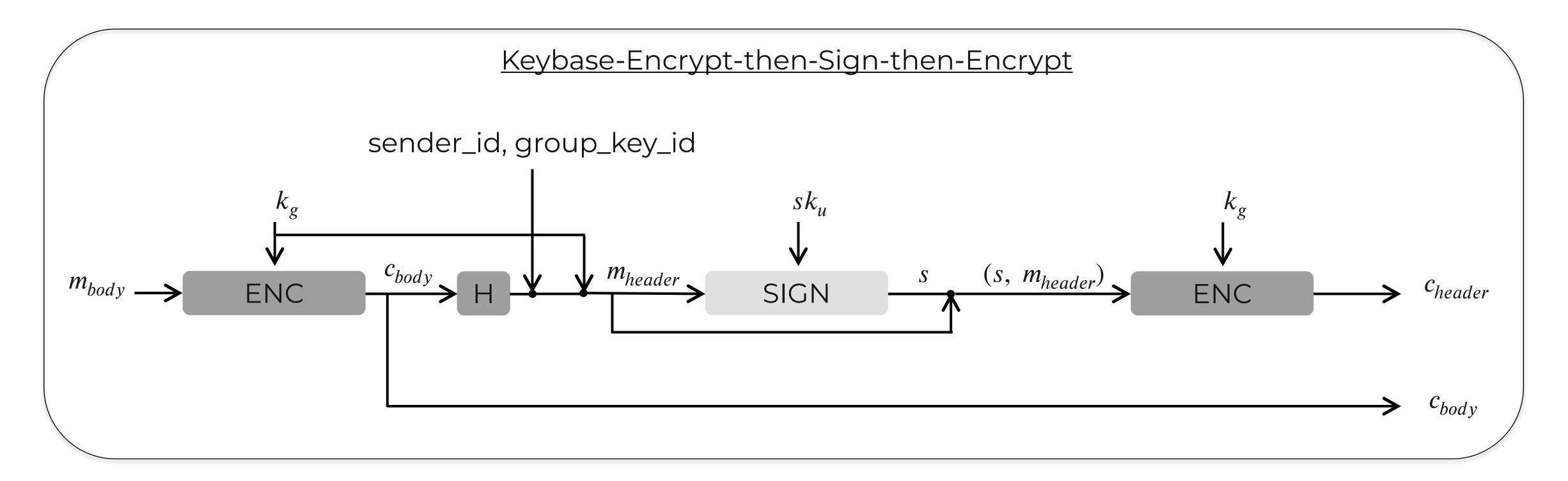








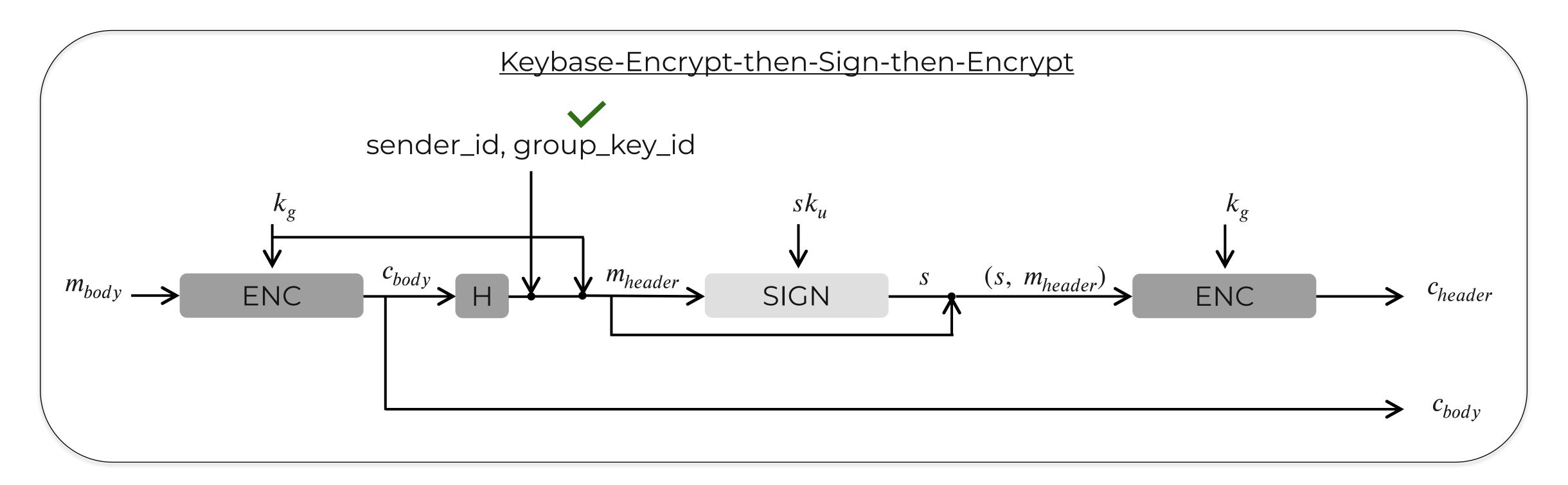


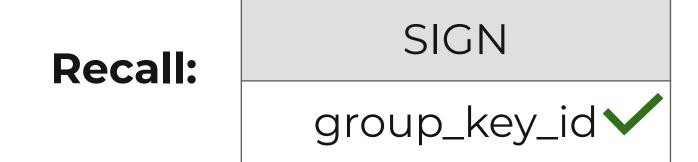


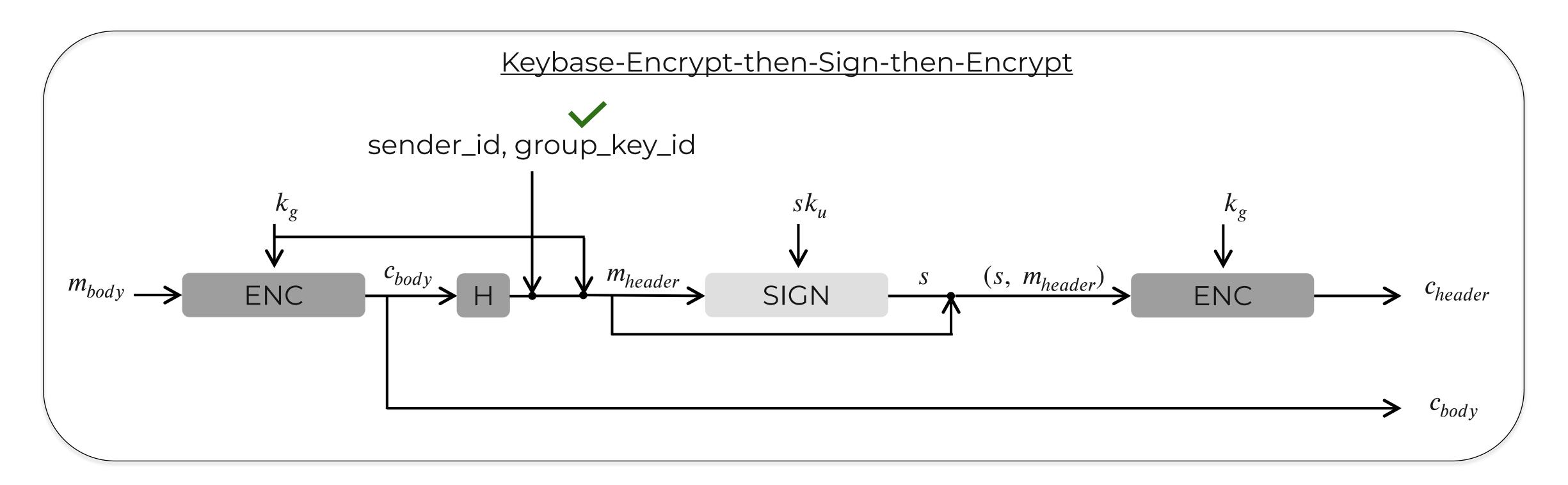
Recall:

SIGN

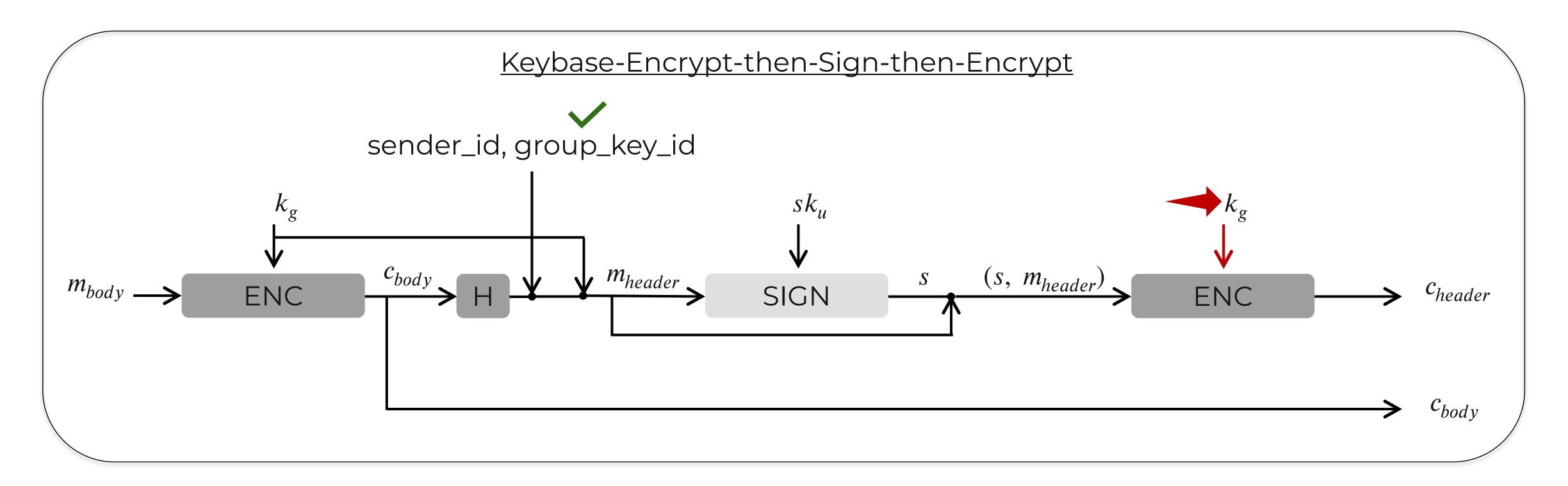
group_key_id



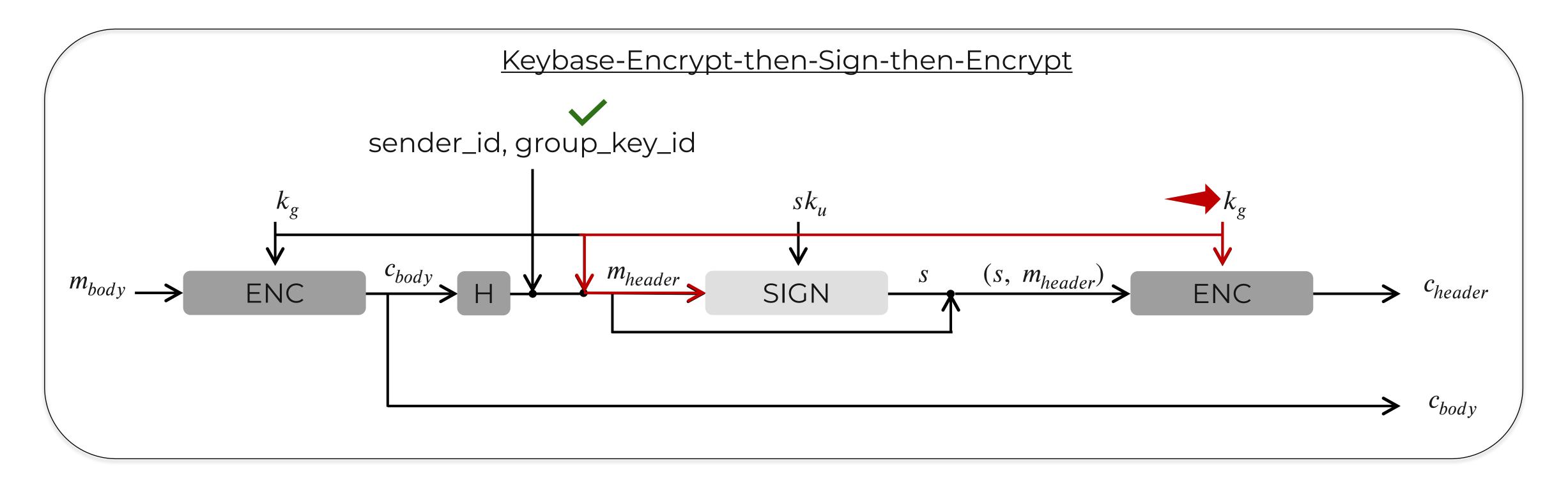




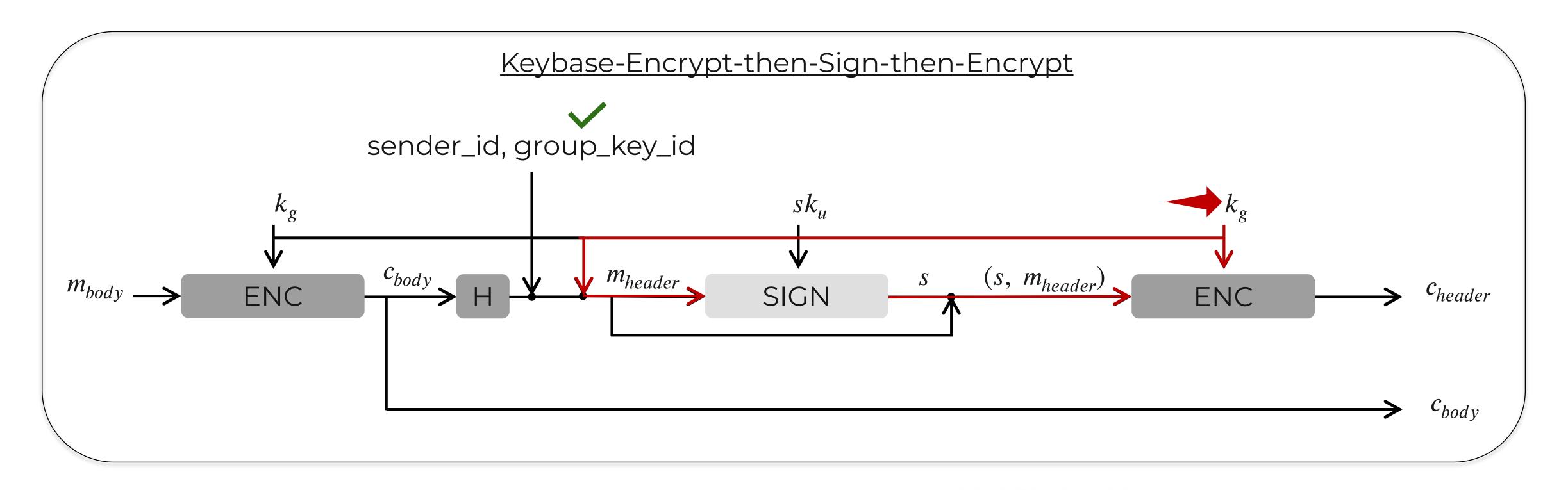




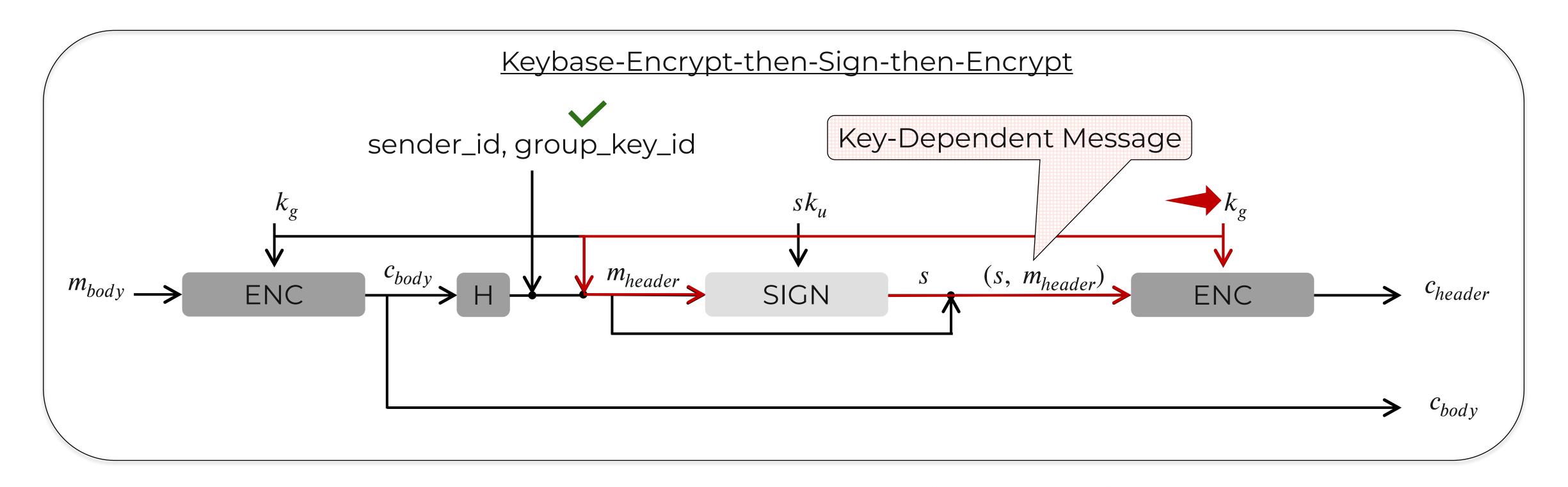




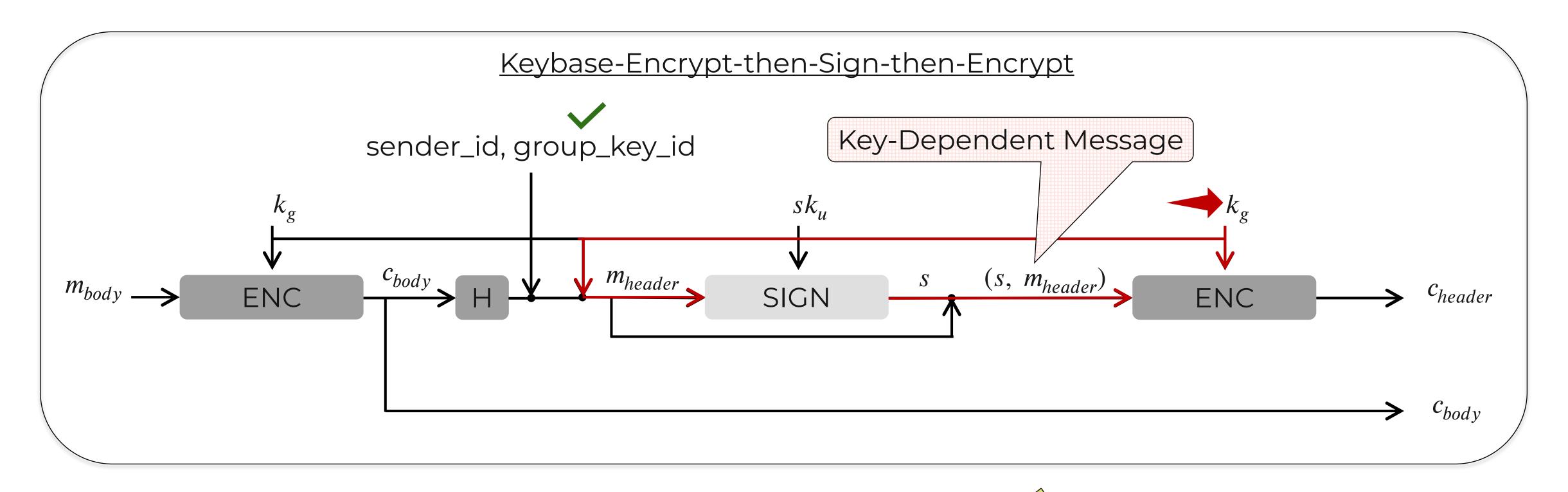




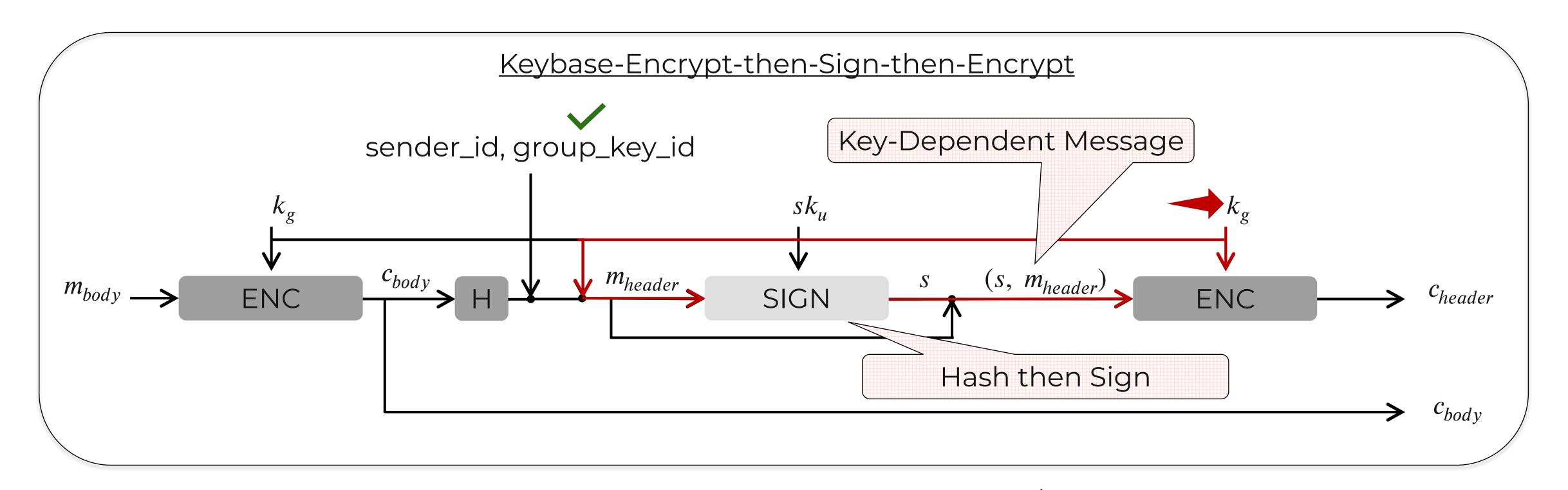


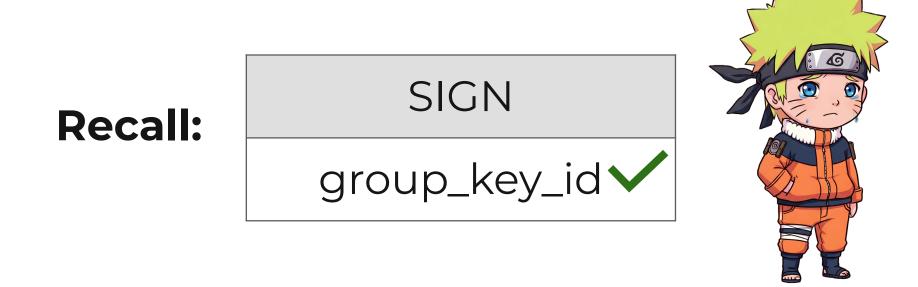


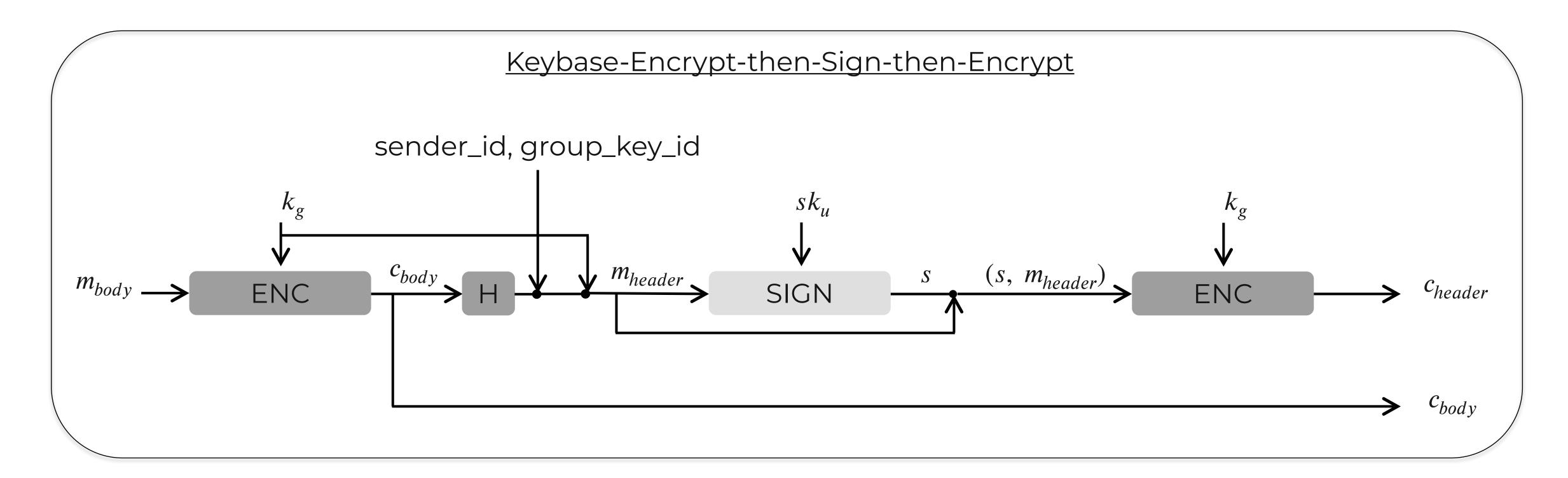


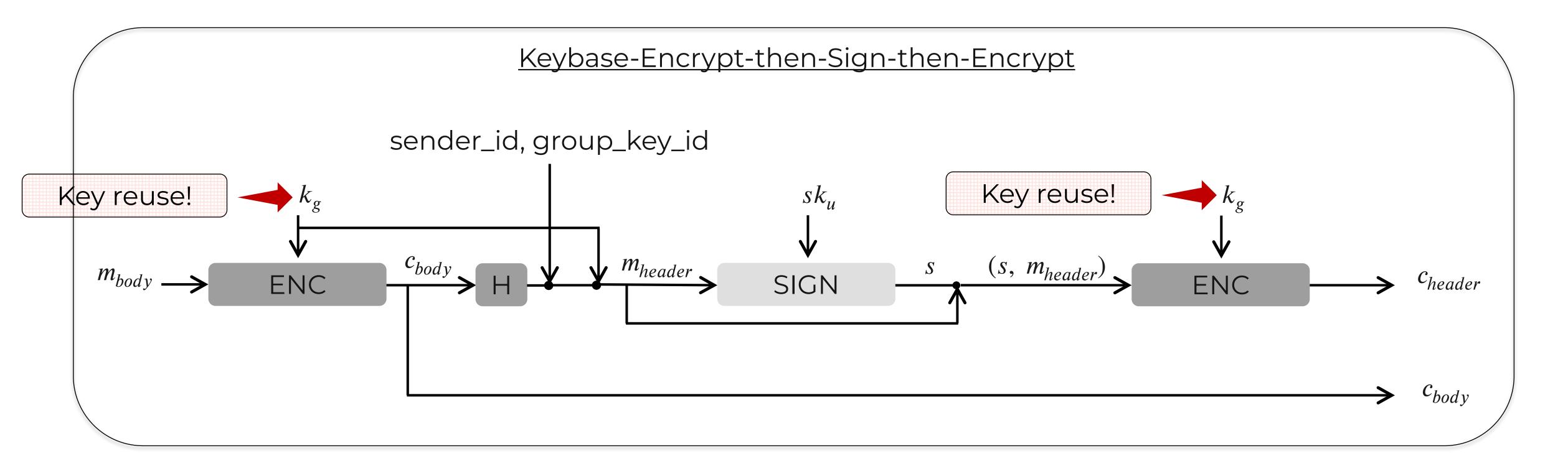


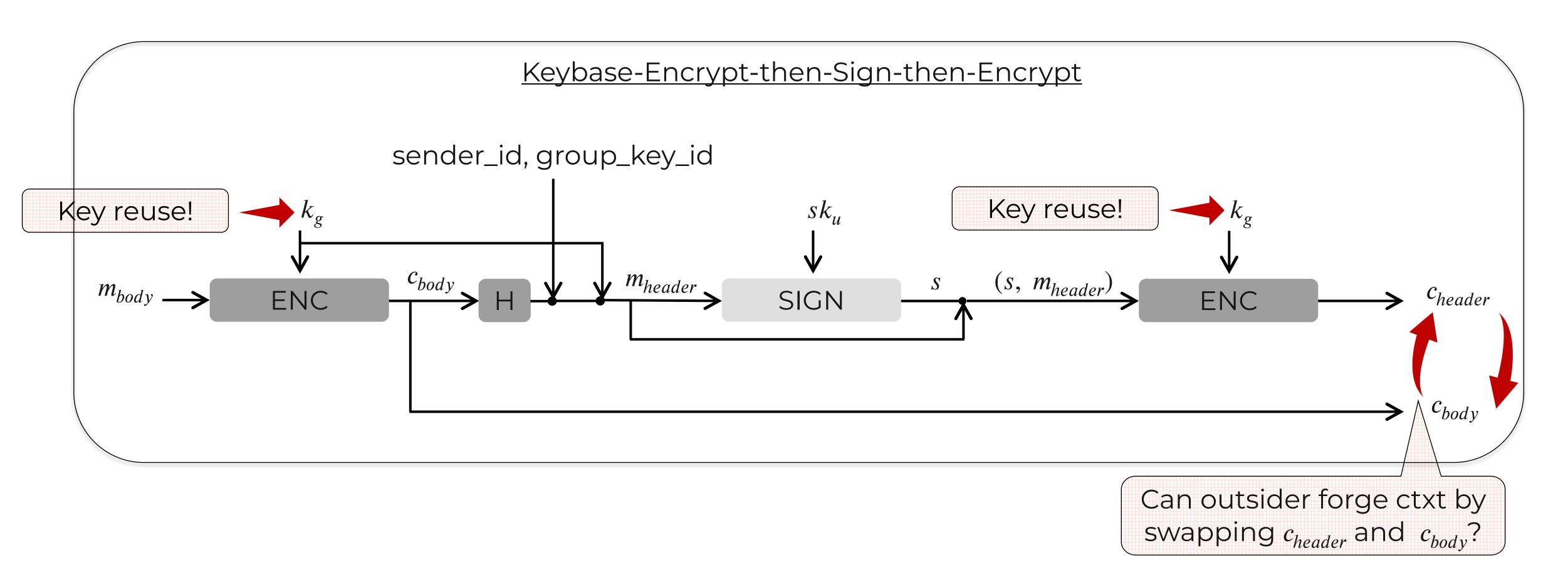


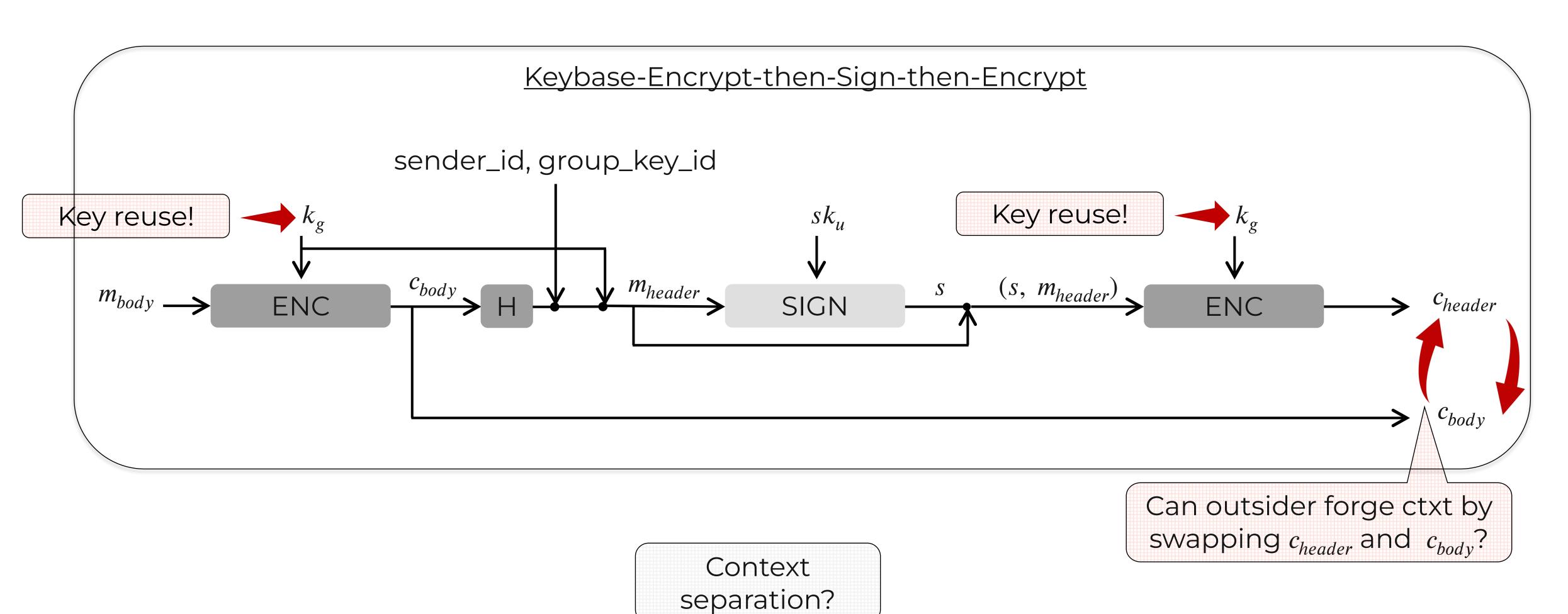


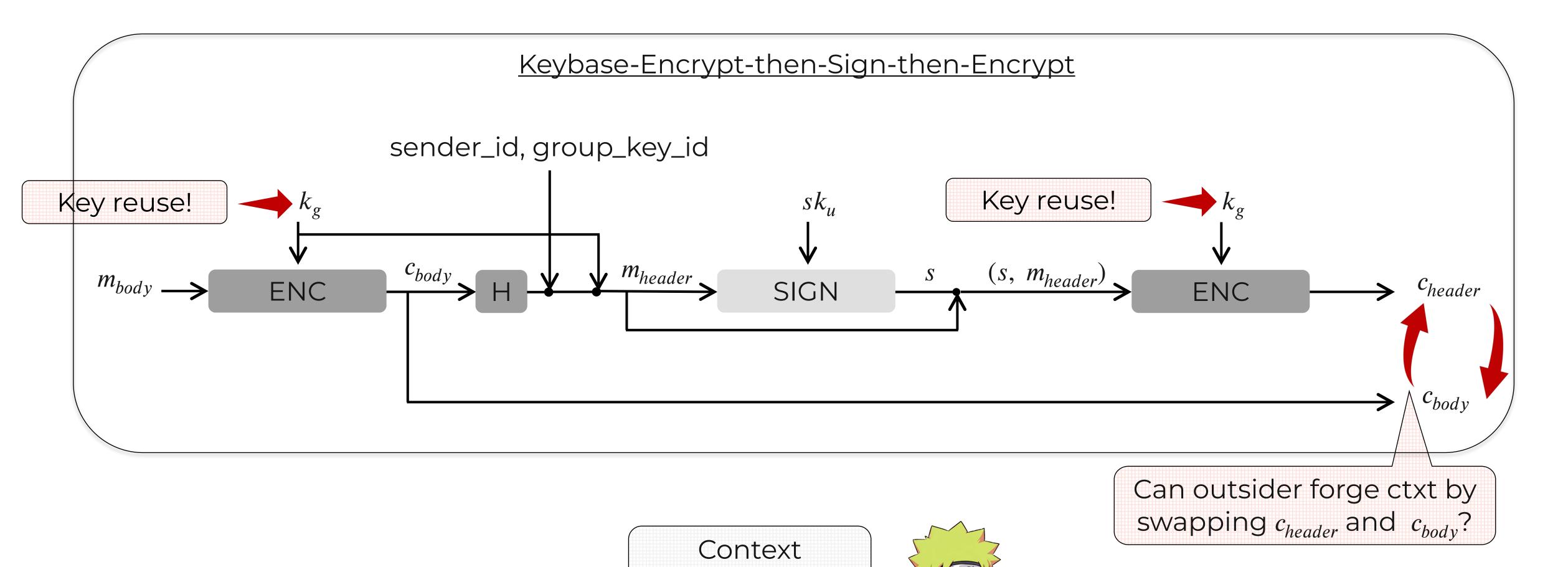




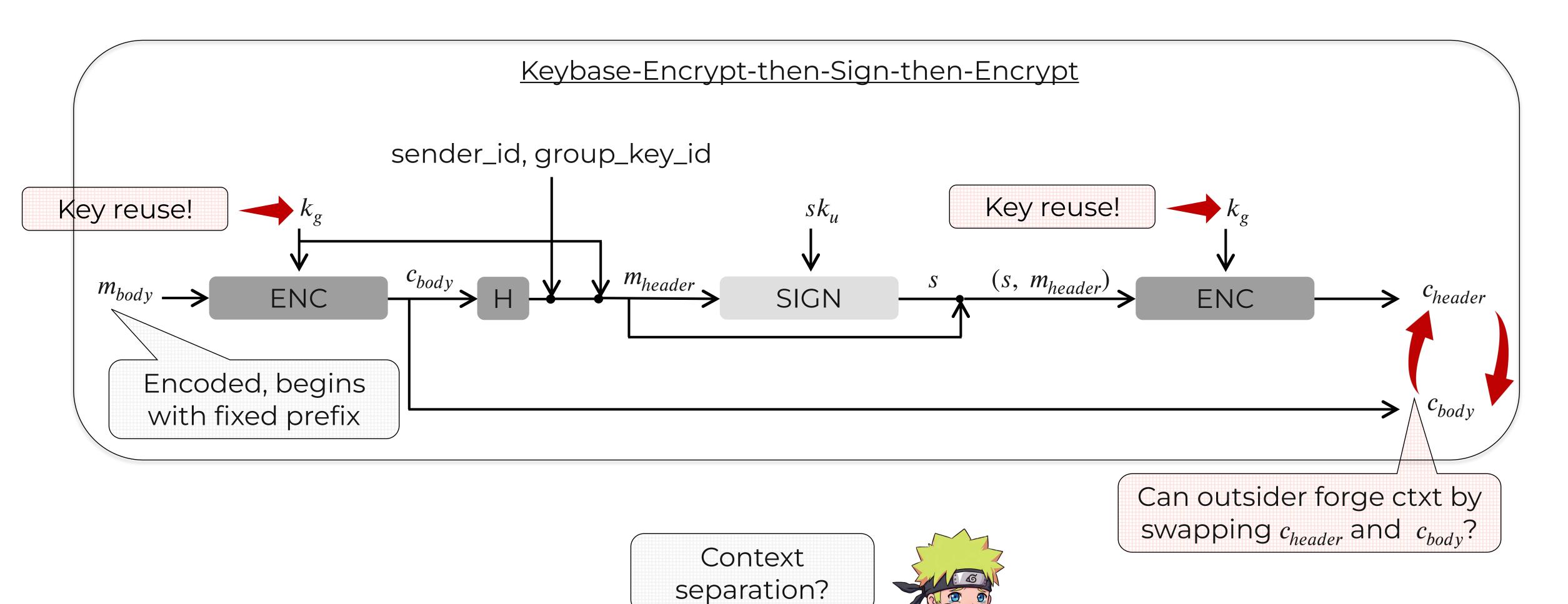




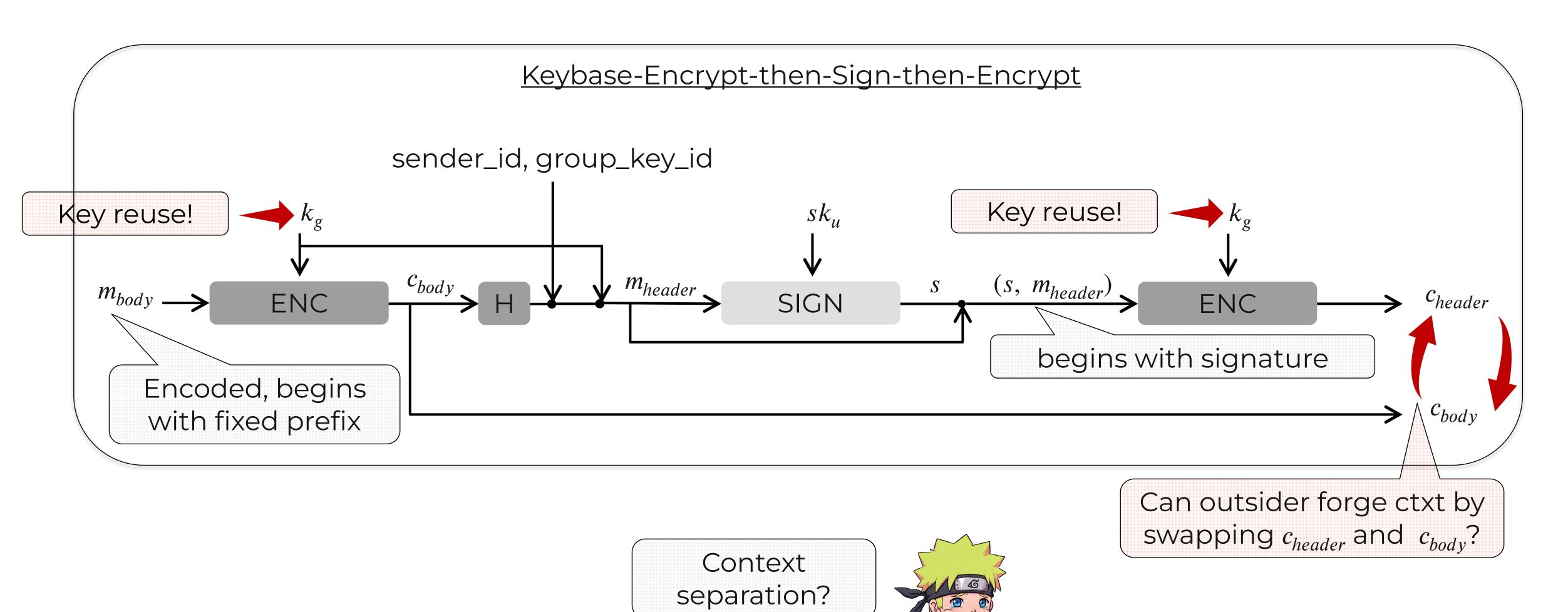


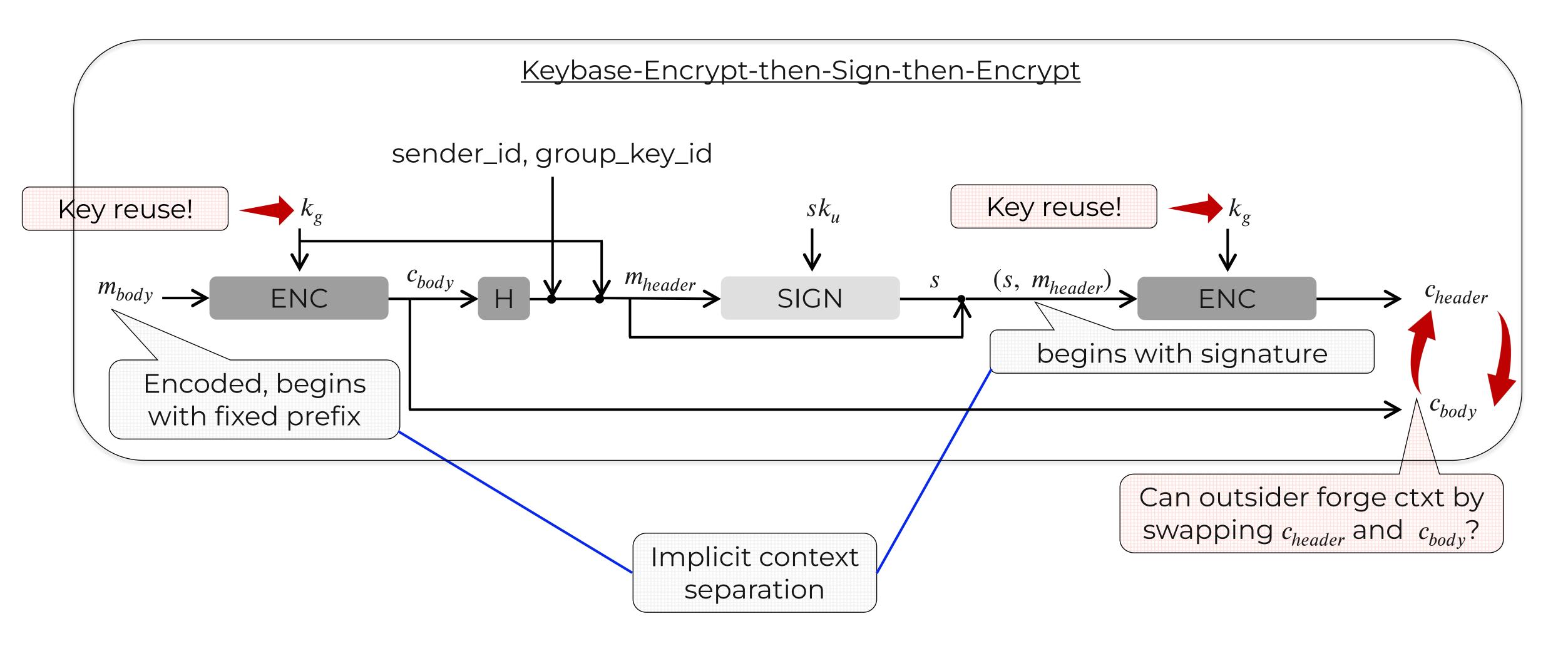


separation?



32





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Reusing keys across different contexts is bad



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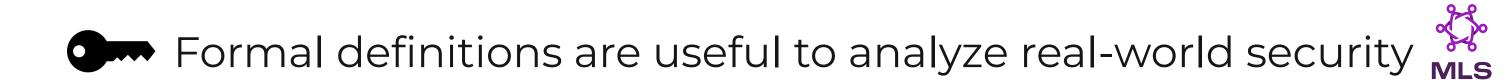


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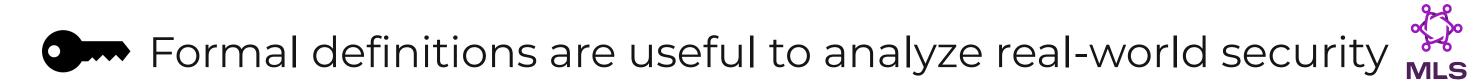


Combining encryption with MAC/signatures is non-trivial



















Details in the papers!

Coming Soon

https://eprint.iacr.org/2024/799.pdf