

Using Formally Verified Post-Quantum Algorithms at Scale

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~~CRYSPEN~~



A Collaboration is Born

Why PQC?

- Quantum Computers soon? Transition now!
 - Attack: store now, decrypt later
- Industry standards, government customers
 - By ~2026!
- 1st Priority: Key Exchange in SSH/TLS/...
- Next: digital signatures
- Many products (some OSS), industry-wide effort → Open Source

Why Verify Lattice-Crypto Implementations?

- Goal: **no implementation vulnerabilities** in **optimized** code
- Experience from Elliptic-Curve Cryptography
 - Auditing code is important but challenging
 - Subtle bugs missed in high-profile implementations
- Simpler than ECC? (No carry chains, standard representations)
 - Yes, in reference implementations
- Tricky optimizations: vectorization, deferred reductions, decoding

The Technical Details

FIPS 203

Federal Information Processing Standards Publication

Module-Lattice-Based Key-Encapsulation Mechanism Standard

Category: Computer Security

Subcategory: Cryptography

Information Technology Laboratory
National Institute of Standards and Technology
Gaithersburg, MD 20899-8900

This publication is available free of charge from:
<https://doi.org/10.6028/NIST.FIPS.203>



We implement and verify
this standard.

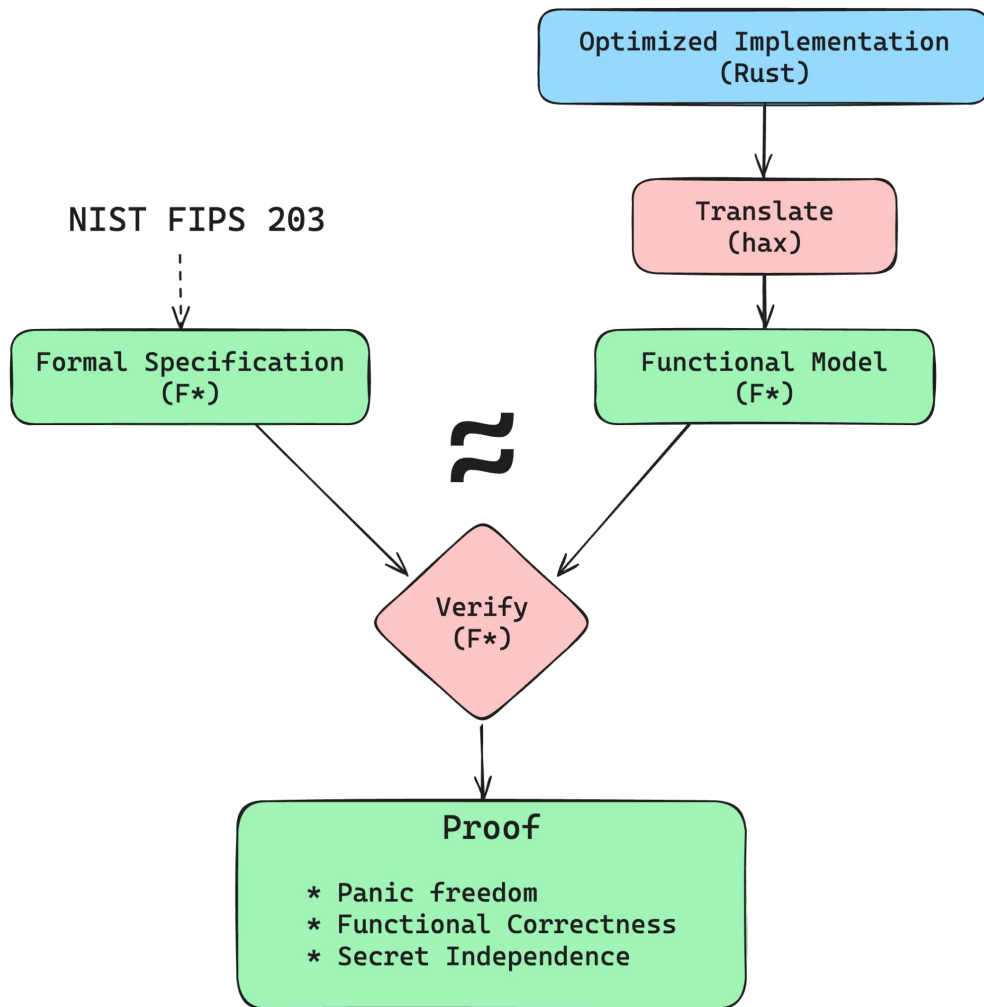
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Implementing ML-KEM in Rust

- Pure Rust code: 16 KLOC
- Optimized for multiple platforms
 - Portable + AVX2 + AArch64
 - 2 KLOC for SIMD optimizations (using intrinsics)
- Easy to integrate and deploy
 - Cargo crate: [libcrux-ml-kem](#)
 - PQCA's official Rust implementation

Mathematics	Low-Level Formats	Algorithms	High-Level APIs
Field, polynomial, matrix	(de)serialization	Sampling, IND-CPA, IND-CCA	ML-KEM 512/768/1024
3k lines	3k lines	6k lines	4k lines

Verifying crypto code written in Rust using `hax` and F^*



Writing Crypto Code in Rust

```
pub fn barrett_reduce(input: i32) -> i32 {  
    let t = (input as i64 * 20159) + (0x4_000_000 >> 1);  
    let quotient = (t >> 26) as i32;  
    let result = input - (quotient * 3329);  
    result  
}
```

Signed Barrett Reduction: with modulus **3329**
(in constant time, so cannot directly use **%**)

Specifying Correctness

```
#[requires(input <= 0x4_000_000 && input >= -0x4_000_000)]  
#[ensures(|result| result <= 3328 && result >= -3328 &&  
         modulo(result, 3329) == modulo(input, 3329))]  
pub fn barrett_reduce(input: i32) -> i32 {  
    ...  
}
```

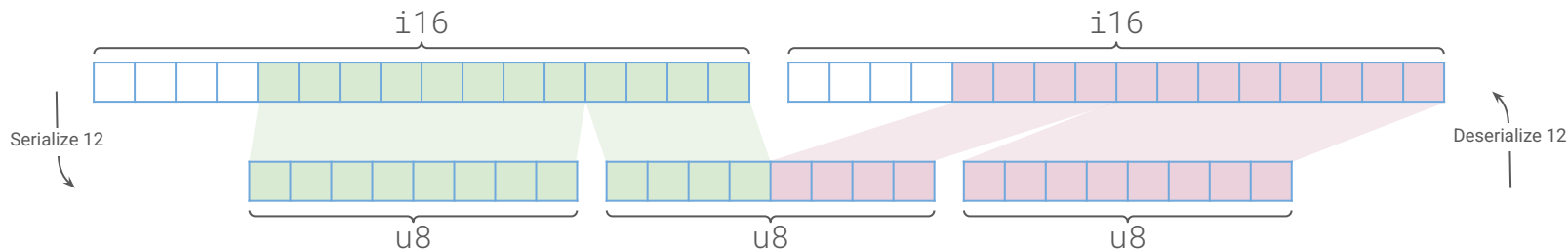
Expected behaviour: compute a signed representative of the input field element (modulo 3329)

Preventing Panics in Rust Code

```
#[requires(input <= 0x4_000_000 && input >= -0x4_000_000)]  
#[ensures(...)]  
pub fn barrett_reduce(input: i32) -> i32 {  
    let t = (input as i64 * 20159) + (0x4_000_000 >> 1);  
    let quotient = (t >> 26) as i32;  
    let result = input - (quotient * 3329);  
    result  
}
```

These arithmetic operations may overflow or underflow
causing the code to panic at run-time

Verifying (De-)Serialization Automatically



24 hand-optimized variants!

12 bits per integer 11 bits per integer 10 bits per integer 5 bits per integer 4 bits per integer 1 bits per integer

6 variants

serialize deserialize

2 variants

portable avx2

2 variants

A new F* tactic that can prove every variant automatically!

Enforcing Secret Independence

Type-based static analysis enforces a “constant-time” discipline

- **arithmetic operations** with input-dependent timing (e.g. division) over secret integers
- **comparison** over secret values
- **branching** over secret values
- array or vector **accesses** at secret indices

Prevents a large class of remote timing attacks (at source level).

Does not prevent compiler-induced leaks, micro-architectural attacks, ...

KyberSlash: a new timing vulnerability

```
void poly_tomsg(uint8_t msg[KYBER_INDCPA_MSGBYTES], const poly *a)
{
    unsigned int i,j;
    uint16_t t;

    for(i=0;i<KYBER_N/8;i++) {
        msg[i] = 0;
        for(j=0;j<8;j++) {
            t = a->coeffs[8*i+j];
            t += ((int16_t)t >> 15) & KYBER_Q;
            t = (((t << 1) + KYBER_Q/2)/KYBER_Q) & 1;
            msg[i] |= t << j;
        }
    }
}
```

Bug found in our
Rust code during
formal verification

Bug also present in
PQ-Crystals,
PQ-Clean, ...
(used in production)

KyberSlash: Exploiting secret-dependent division timings in Kyber Implementations.
IACR Transactions on Cryptographic Hardware and Embedded Systems, 2025(2),
209-234. Bernstein, D. J., Bhargavan, K., Bhasin, S., Chattopadhyay, A., Chia, T. K.,
Kannwischer, M. J., Kiefer, F., Paiva, T. B., Ravi, P., & Tamvada, G.

Scaling the Proof Effort

- Full formal verification of a large code-base
 - Source Rust code: 16 KLOC
 - Generated F* model: 28 KLOC (Portable + AVX2)
- Multiple automation strategies
 - SMT-based automation for low-level mathematics
 - Tactic-based automation for serialization
 - Type-based secret independence analysis
- Still needs many manual F* proofs + annotations for the full proof

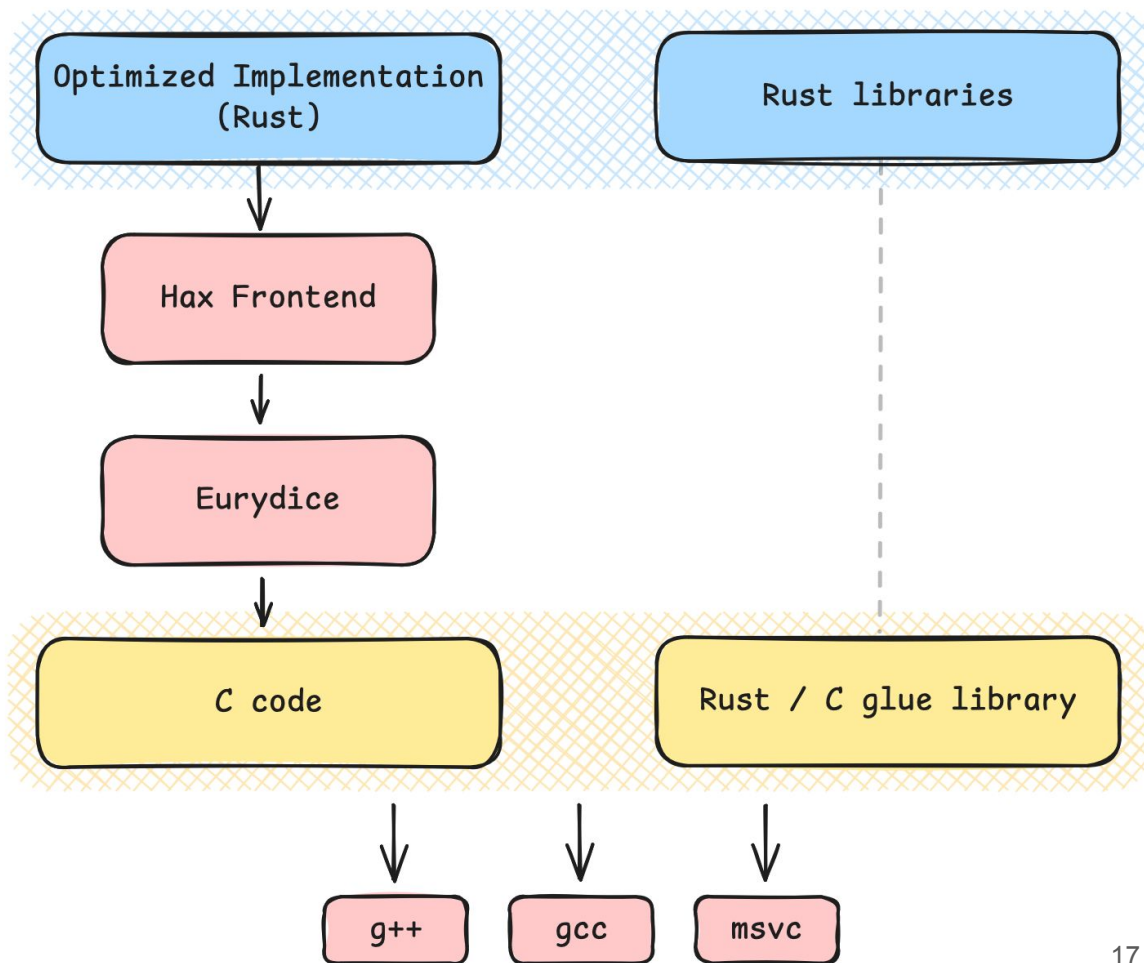
Mathematics	Low-Level Formats	Algorithms	High-Level APIs
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6k lines of F*	5k lines of F*	6k lines of F*	4k lines of F*

Verified PQC at Scale

Compiling verified Rust to C code

Many mainstream crypto
libraries need C code

NSS, BoringSSL,
OpenSSH, OpenSSL, ...



Integration Challenges

- C code size is larger than Rust
 - From monomorphizing ML-KEM variants from Rust
 - ~40KB optimized for speed
- Match existing APIs in the crypto library
 - Opaque secret keys, Alignment, Strict aliasing
- C++ toolchain compatibility (yes, even iOS, MSVC, ARM, bigendian...)
- Scale: ~100 build configurations, and evolving

Maintainability and Performance

- Establishing speciality tooling
 - Change workflow: modify Rust code, re-prove, re-generate C
 - Review specs, not code – but computer-check proofs!
 - Continuous integration for tools (ARM/Intel × Debian/MacOS)
 - Long-term support from Cryspen
- AVX2: ~2x faster than BoringSSL reference implementation
 - A great argument for at-scale deployment!

Takeaways

- PQC is coming, verification is important
 - Demonstrated with ML-KEM and KyberSlash
 - Deployments in OpenSSH, NSS, PQCA, Signal, Dropbear
- Many challenges need solving between
 - Formally verified fast code
 - At-scale deployment
- Next up: ML-DSA

Try our Rust or C code today!

