



Introduction to Cryptography

History

SHA2 implementation

- Usage
- Algorithm & optimizations
- Proposed implementation
- Results
- Conclusions





Cryptography - History and Usage

1900 BC is the 1st kwon usage of cryptography by the Egyptians – simple substitution scheme

1 AD Julius Cesar used the simple letter shift cipher in the Gallic Wars

300 AD a new mathematical cryptographic scheme was used by Sun Tzu

1412 AD an 14vol encyclopedia on Cryptanalysis is compiled by the Arabs

In the 2nd world war the enigma rotor machine is used (substitution scheme using a continuously changing alphabet)



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Msg: *OMNIA GALLIA EST DIVISA IN PARTES TRES* Enc: RPQLD JDOOLD HVW GLYLVD LQ SDUWHV WUHV

 $CRT(2,3,2)_{(357)} = ??(23)$







Cryptography - History and Usage

In the 70's with the development of complex electronic systems much more complex cryptographic systems have been created, such as Lucifer and DES.

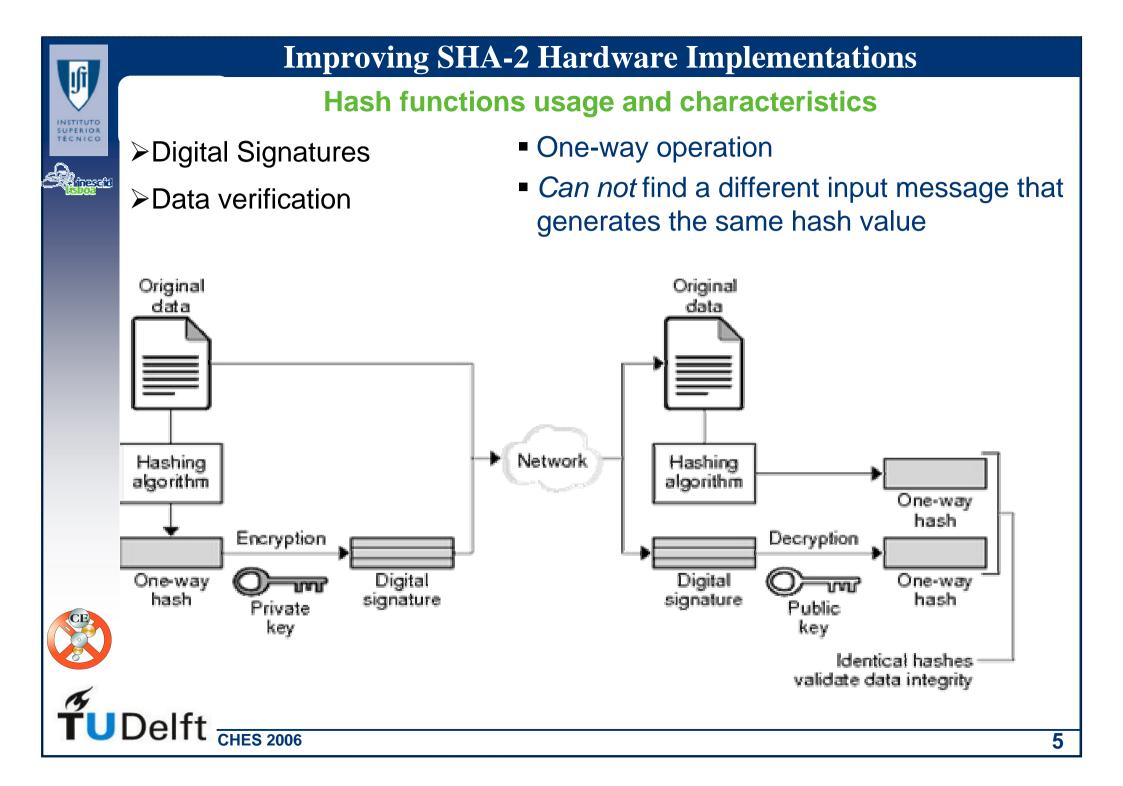
nowadays a variety of cryptographic algorithms exist an they are present in almost every day actions:

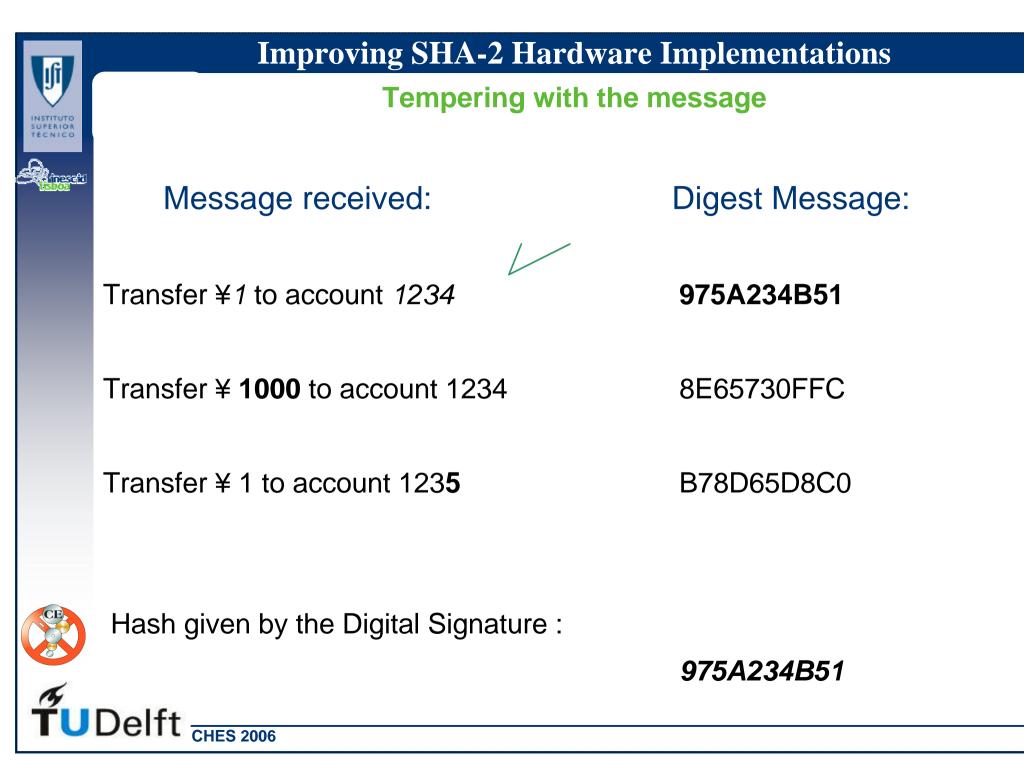
- Accessing the internet
- E-shopping
- ATM machines
- Emails
- Buildings access
- Pay TV
- Anti-car theft systems
- Private communications

• ...

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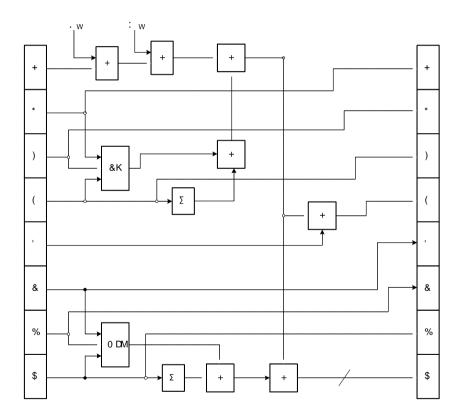




SHA2 Hash Function

Characteristics:

Based on simple logical and arithmetic operations



Additional operations:

- The value W_t is computed from the input data (data block expansion).

- After the 64 round (for SHA256) the resulting value has to be added to the intermediate Digest Message.

SHA512 requires 80 rounds using 64 bit values

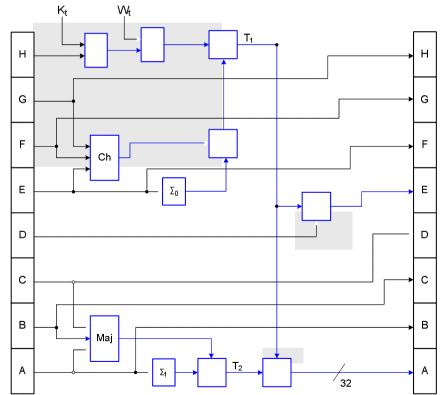




SHA2 Hash Function

Characteristics:

- Based on simple logical and arithmetic operations
- Only the values A and E requires computation



Optimization:

- Only the A and E computation depends on values computed in the previous round

- Values H_t, G_t, F_t and D_t, C_t, B_t do not depend on the values of round *t*.

- Part of the round *t* computation can be performed in round *t*-1





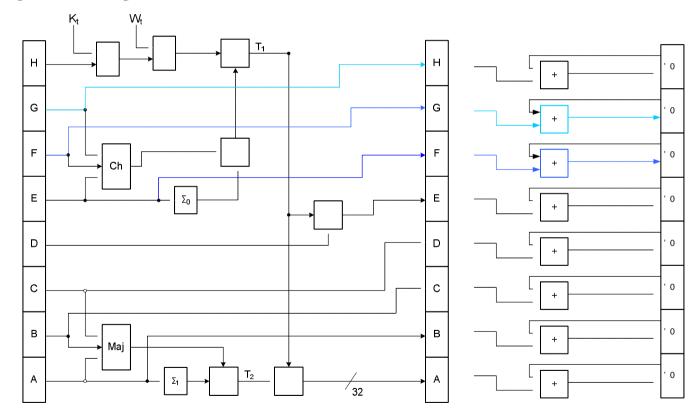
SHA2 Hash Function

Characteristics:

- Only the values A and E requires computation
- After the rounds the values have to be added to the Digest Message

-
$$H_t = G_{t-1} = F_{t-2}$$
; $D_t = C_{t-1} = B_{t-2}$

- DM is known in the 1st round







SHA2 Hash Function

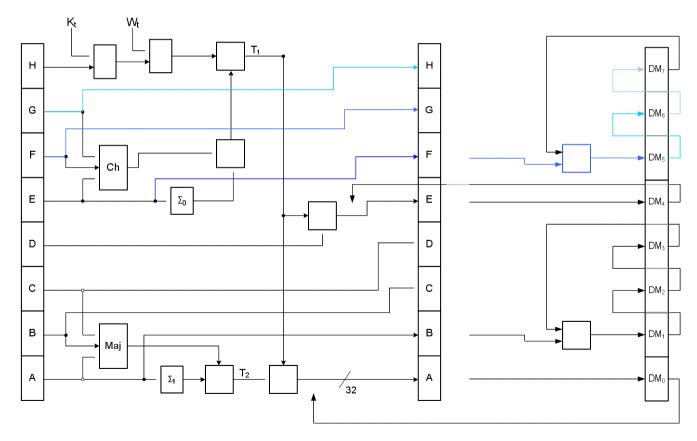
Characteristics:

Optimization:

- Only the values A and E requires computation
- After the rounds the values have to be added to the Digest Message



- DM is known in the 1st round





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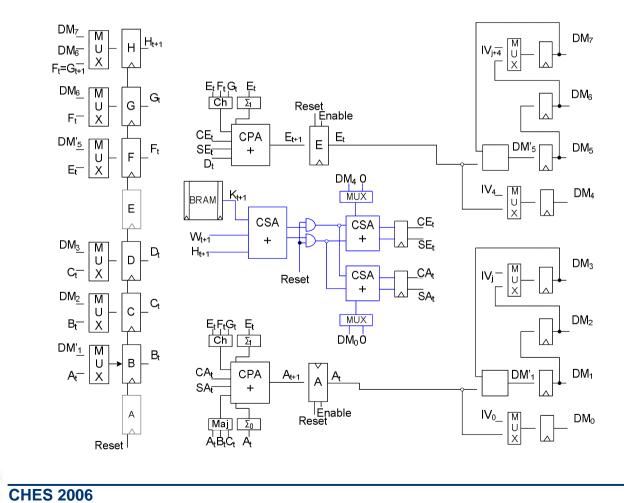
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Improving SHA-2 Hardware Implementations

SHA2 hardware implementation

SHA2 core:

- Variable IV initializations
- The pipeline has to be filled





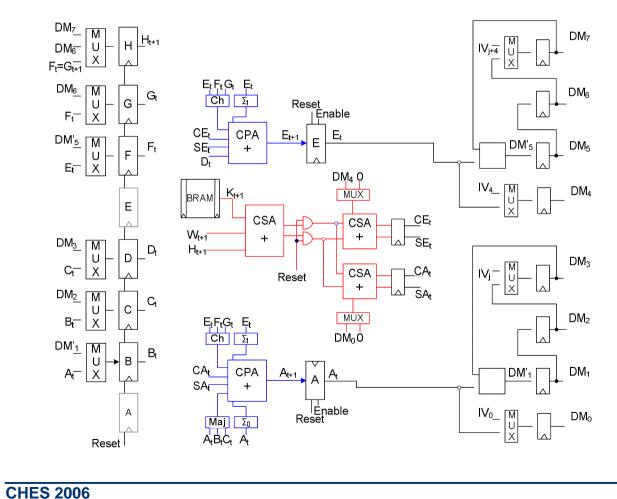
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Improving SHA-2 Hardware Implementations

SHA2 hardware implementation

SHA2 core:

- Round *t* is being calculated in blue while round *t*-1 is being calculated in red
- Critical path ~ 6 input adder

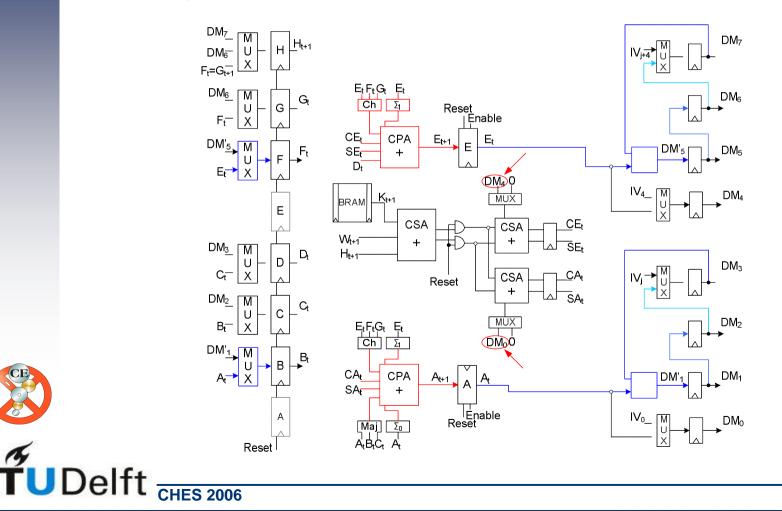




SHA2 hardware implementation

SHA2 core:

- The DM is added to the calculated values e.g. $DM_5 = DM_5 + F_{t+1} = DM_5 + E_t$
- DM₀ and DM₄ are calculated in the round hardware





Results for the standalone SHA2 core

SHA256 core:

- Variable IV initializations allowed
- 1.4 Gbits throughput
- Only 5% utilization of the available logic (Virtex II Pro 30)

Efficiency improved:

- 50% improvement of the Throughput/Slice metric regarding commercial cores
- 100% improvement of the Throughput/Slice metric regarding academia cores

SHA256	Sklav	Our	McEv.	Our	Helion	Our
Device	XCV	XCV	XC2V	XC2V	XC2PV-7	XC2PV-7
Slices	1060	764	1373	797	815	755
BRAMs	= 1	1	=1	1	1	1
Frequency	83	82	133	150	126	174
Throughput	326	646	1009	1184	977	1370
Throughput/Slice	0.31	0.84	0.74	1.49	1.2	1.83
Improved	171 %		101%		53%	



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Results for the standalone SHA2 core

SHA512 core:

- Variable IV initializations allowed
- 1.8 Gbits throughput
- Only *13%* utilization of the available logic (Virtex II Pro 30)

Efficiency improved:

• 77% improvement of the Throughput/Slice metric.

SHA512	Sklav	Lien	Our	McEv	Our	Our
Device	XCV	XCV	XCV	XC2V	XC2V	XC2VP-7
Slices	2237	2384	1680	2726	1666	1667
BRAMs	n.a.	n.a.	2	= 1	1	1
Frequency	75	56	70	109	121	141
Throughput	480	717	889	1329	1534	1780
Throughput/Slice	0.21	0.31	0.53	0.49	0.92	1.01
Improved	165%	77 %		88%		



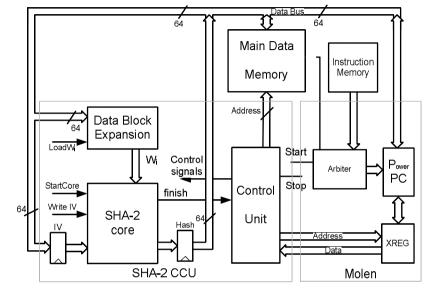


MOLEN with an SHA2 CCU

Main characteristics:

Low FPGA utilization.

- **7**% of a XCV2P30 for the SHA256
- **13** % of a XCV2P30 for the SHA512



Throughput of 785 Mbits/s @100MHz for SHA256

Throughput of **1200** *Mbits/s* @100MHz for SHA512

Minimal software integration costs

•A large range of encryption applications can be speedup just be being recompiled for the MOLEN processor with the SHA2 core.



Prototype @: **TUDelft** CHES 2006



SHA2 core in the MOLEN processor

Minimal Software integration costs:

Original Software:

Declaration: void SHA256(){ instructions...}

Usage : SHA256(Data,size, IV);

Modified for MOLEN:

Declaration:#pragma call_fpga encrypt
void SHA256(){ /*implemented in Hardware*/}Usage :SHA256(Data,size, IV);

Significant Speedup for a minimal area cost:

7% occupation of a Virtex II Pro 30 FPGA (994 Slices)

Throughput of 785Mbits (in MOLEN @100MHz) instead of 5 Mbits (in Software @300MHz), 153x Speedup.





Conclusions – SHA2

Efficiency gains to existing state-of-the-art

Higher Throughput/Slice ratio

- 50 % when compared with known SHA256 commercial cores
- 100% when compared with academia SHA256 related art
- 77 % when compared with academia SHA512 related art

Throughput of 1.37 Gbits for SHA-256(@ 174 MHz)5% Occupation (533 Slices)Throughput of 1.78 Gbits for SHA-512(@ 141 MHz)13% Occupation (1667 Slices)

MOLEN implementation

Low FPGA utilization: **7%-13%**

(Virtex II Pro 30)

(@ 100 MHz)

High throughput: **785** – **1200** Mbits/s

SHA256 throughput speedup of 153x

Minimal software integration costs

CB

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* Prototype running @100MHz

