

Secure Lightweight Entity Authentication with Strong PUFs: Mission Impossible?

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Contents

- Mission: Secure Lightweight Entity Authentication



- Physically Unclonable Functions (PUFs). How can they help?

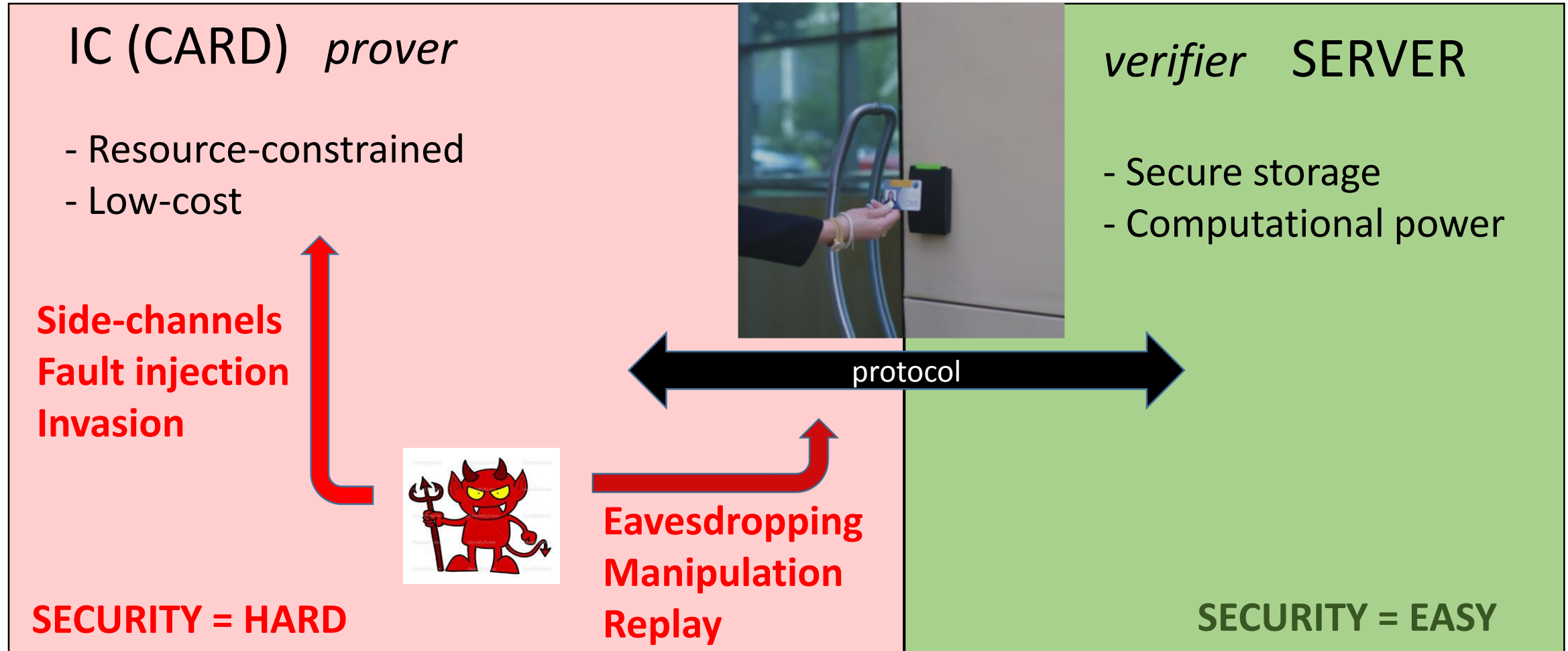


- 8 PUF Entity Authentication protocols
 - Security and practicality analysis
 - No protocol details here, only properties (*limited presentation time*)
- Conclusion

Lightweight Entity Authentication

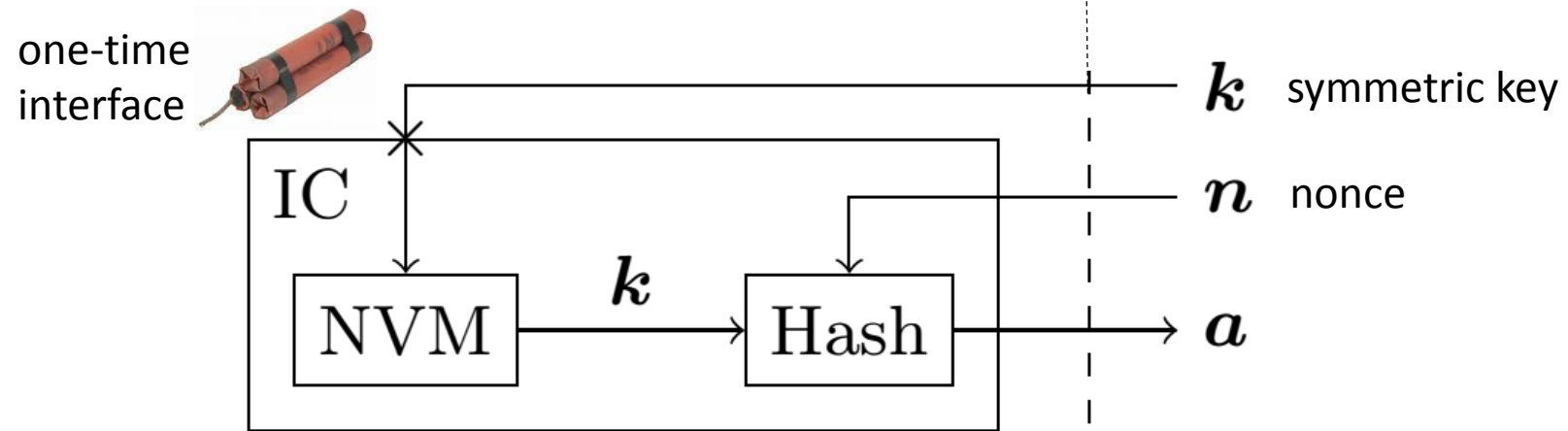
2 parties sharing
a secret

e.g. building access card



Lightweight Entity Authentication with PUFs?

Conventional approach: Non-volatile Memory (e.g. Flash)



NVM:

High manufacturing cost

Flash floating gate \neq CMOS compatible

PUFs = CMOS compatible

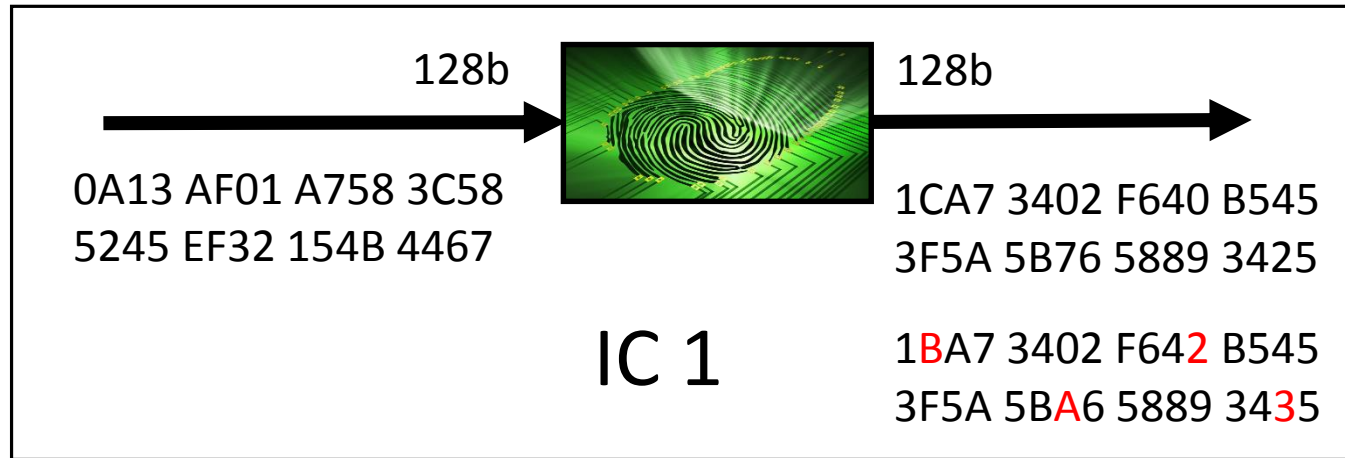
NVM:

Vulnerable to physical attacks

Robust electrical storage

PUFs = chemical storage

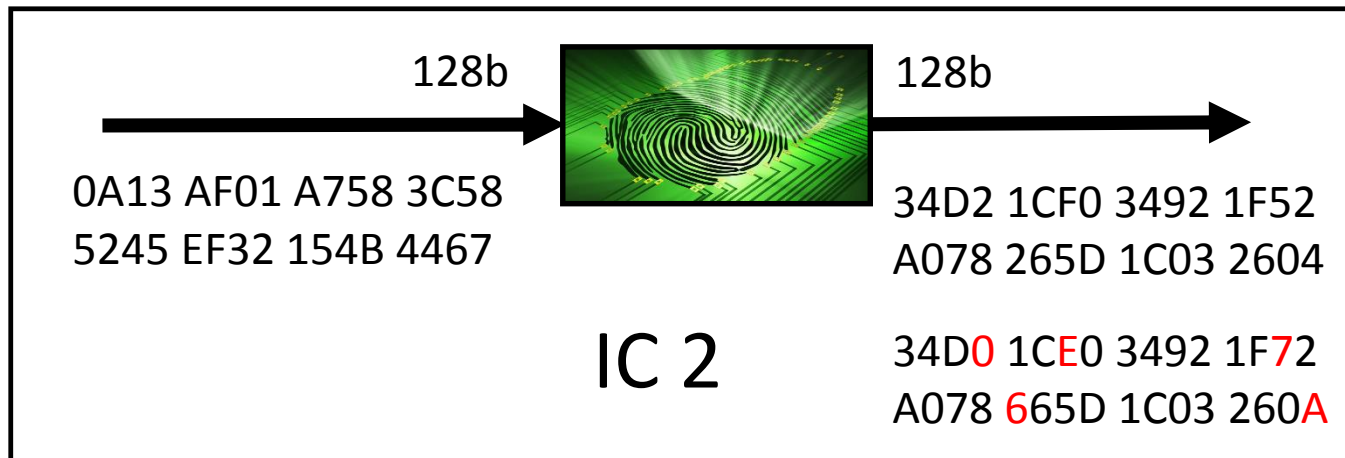
PUFs = Physically Unclonable Functions



Chip-dependent binary
function with noisy output

Evaluation 1

Evaluation 2 \approx 1-15% noise



Similar to biometrics, applied
to an IC rather than a human

Evaluation 1

Evaluation 2 \approx 1-15% noise

PUFs = Physically Unclonable Functions



Does not exist, but **all 8 protocols** need it to counteract brute-force and random guessing

The protocols need:

“Strong PUF”
(e.g. arbiter PUF)



+ lightweight solution for output expansion:
repeated evaluation e.g. $Out = PUF(PRNG(In))$

Note: For convenience, we define strong PUFs using the popular more recent notion of large input space rather than the original definition

The protocols cannot use:

“Weak PUF”
(e.g. SRAM PUF)



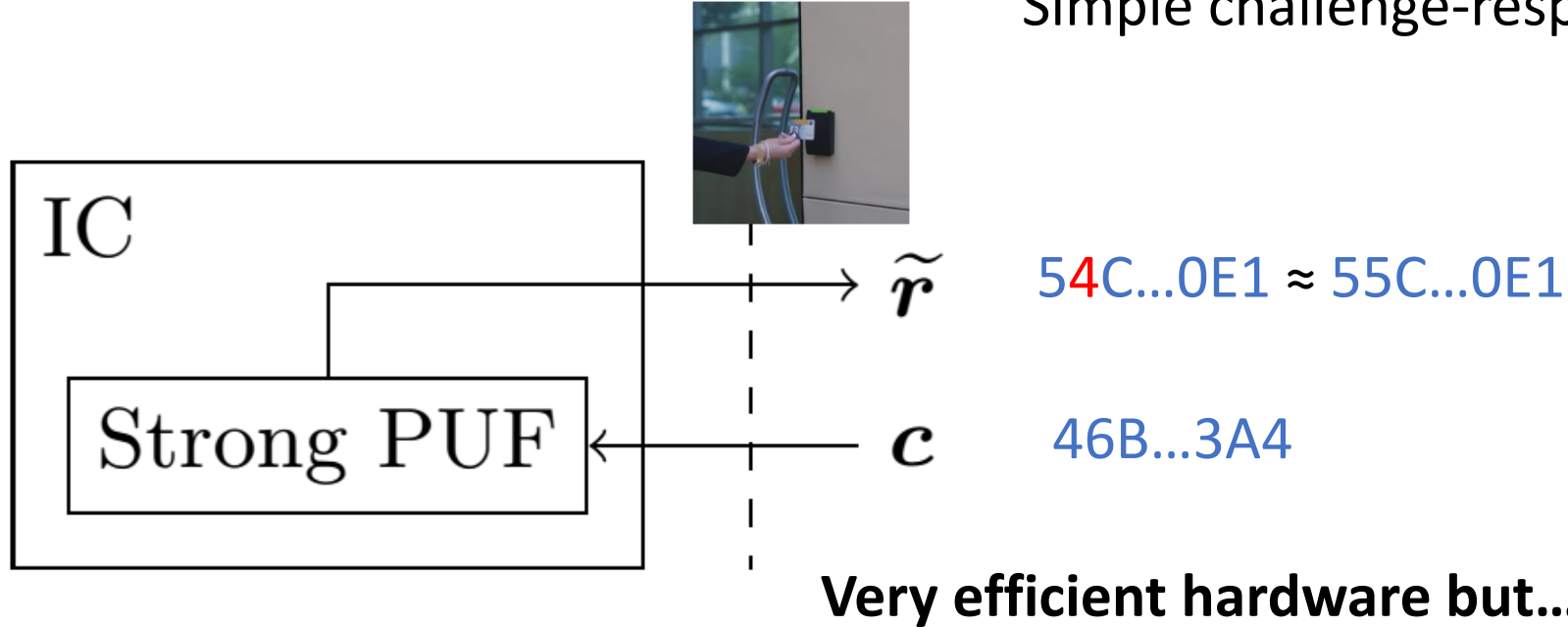
(This type of PUF is mainly used to generate a secret key)

Our reference protocol

Entity Authentication: Basic strong PUF protocol

07/17

Simple challenge-response



Prerecorded I/O pairs

1F2...51F	C73...467
46B...3A4	55C...0E1
E03...127	0A1...245

3 authentications, discard pair after use to avoid replay

No secure instantiation: PUF modeling attacks

- PUF I/O behavior is correlated
- Mathematical clone: learn full I/O behavior given a small training set (machine learning)
- No PUF has valid claim to be resistant and lightweight

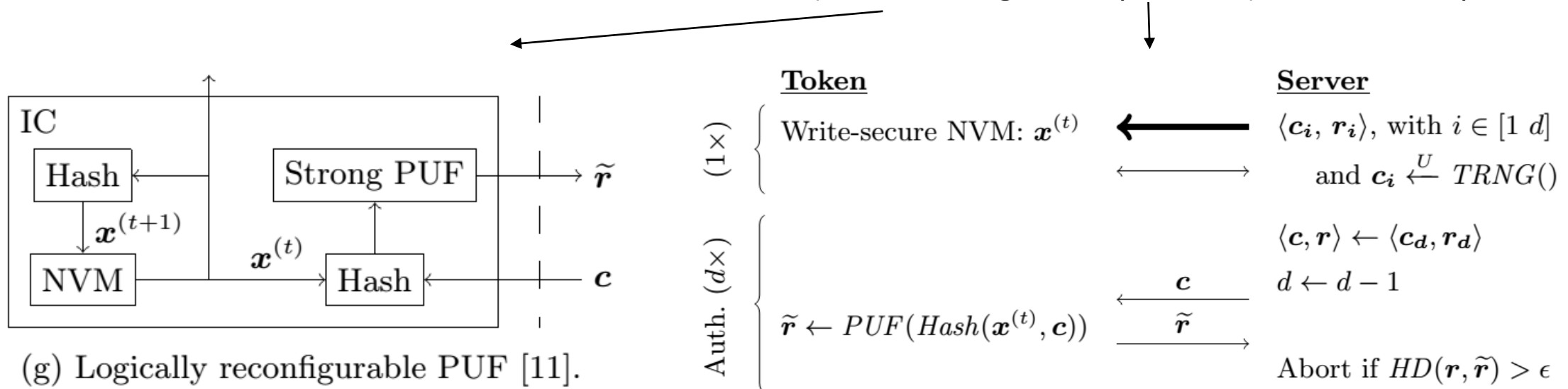
All 7 other protocols have additional building blocks: Hash, TRNG, Error-Correcting Code, ...

8 Strong PUF Entity Authentication Protocols

- Basic (2001, PhD MIT)
- Controlled (2002, ACM CCS)
- Öztürk et al. (2008, PerCom)
- *Hammouri et al.* (2008, journal)
- Logical Reconfiguration (2011, CHES & journal)
- Reverse Fuzzy Extractor (2012, FC)
- Slender (2012, SP & journal)
- Converse (2012, DATE & TRUST)

OUR FIRST CONTRIBUTION: FIRST OVERVIEW

Common framework with reference, same notation (IC block diagram & protocol), initiate comparison



8 Strong PUF Entity Authentication Protocols

OUR SECOND CONTRIBUTION: FIRST ANALYSIS

Security of the Protocol

Impersonation?

Denial-Of-Service?

Practicality of the Protocol

Are the PUF advantages (w.r.t. NVM) preserved?

- Low-cost manufacturing
- Improved physical security

PUF imperfections taken into account?

- Noisiness
- Prone to modeling

Are there PUF assumptions degrading the usability / generality ?

Efficiency & Scalability

The next 7 slides
Severe issues for
all-but-one protocol.
We do not
recommend their
usage






PUF noisiness taken into account?

Error Tolerance

Basic	
Reconfiguration	
Hammouri et al.	
Slender	

Design flaw: error amplification. Protocol does not function. No FPGA proof-of-concept.

Error Correction

Our reference protocol	
Controlled	
Öztürk et al.	
Reverse fuzzy extractor	
Converse	

Exhaustive search for error pattern by server. They assume noise < 1%. Might not be feasible for high noise. No FPGA Proof-of-concept.

Prevention of PUF modeling attacks?

Crypto (Hash)



Controlled

Reverse fuzzy extractor

Converse

Our reference protocol

Guarantee

Lightweight Protection



Öztürk et al.

Hammouri et al.

Slender

PUF needs high
resistance

No protection



Basic

Reconfiguration

Not usable, no
secure
instantiation

PUF advantages w.r.t. NVM preserved? Denial-Of-Service?

No NVM



Basic
Controlled
Reverse fuzzy extractor
Hammouri et al.
Slender
Converse
Our reference protocol

Write-Secure Reprogrammable NVM



Reconfiguration

Read/Write-Secure Reprogrammable NVM



Öztürk et al.

Undermines the advantages of PUF technology: low-cost manufacturing & physical security.

NVM = state vector (requires synchronization between PUF-IC and server). No user authentication for state update. Attacker can do it too: DoS.

PUF output expansion exploits?



We did not spot a problem

Basic Hammouri et al.

Controlled Reconfiguration

Öztürk et al. Converse



Token impersonation (for the chosen implementation only)

Reverse fuzzy extractor

PRNG = LFSR

Slender

PRNG = LFSR XOR LFSR

Not applicable (weak PUF)

Our reference protocol



PUF assumptions degrading the generality?

Every strong/weak PUF



Our reference protocol

Every strong PUF



Controlled

Converse

Every strong PUF satisfying the assumptions

Basic

Robust against modeling (does not exist)

Reconfiguration

Robust against modeling (does not exist)

Öztürk et al.

Both easy & hard to model (contradiction)

Hammouri et al.

Both easy & hard to model (contradiction)

Reverse fuzzy extractor

Noise lower limit (opposing stability)

Slender

Both easy & hard to model (contradiction)



Efficiency / Scalability (server storage)?

Constant w.r.t. # authentications



Our reference protocol

Öztürk et al.

Hammouri et al.

Reverse fuzzy extractor

Slender

Linear w.r.t. # authentications



Basic

Controlled

Reconfiguration

Special category



Converse

**Server and attacker face the
same brute-force workload.
Not usable.**

Comparison

- Resources
- Authenticity Type
- Server storage

	Weak PUF	Strong PUF ¹	NVM	TRNG	Gen	Rep	Hash	I× interface	Other	Token auth.	Server auth.	# Auth.	CRPs	Model	Key
Reference I	×	×	✓	×	×	×	✓	✓	×	✓	×	∞	×	×	✓
Reference II	✓	×	×	×	×	✓	✓	✓	×	✓	×	∞	×	×	✓
Naive	×	✓	×	×	×	×	×	×	×	✓	×	<i>d</i>	✓	×	×
Controlled	×	✓	×	×	×	✓	✓	✓	×	✓	×	<i>d</i>	✓	×	×
Öztürk et al.	×	✓ ²	✓ ⁶	×	×	×	×	✓	✓	✓	×	∞	×	✓	×
Hammouri et al.	×	✓ ²	×	✓	×	×	×	✓	✓	✓	×	∞	×	✓	×
Reconfiguration	×	✓ ³	✓ ⁶	×	×	×	✓	×	✓	✓	×	<i>d</i>	✓	×	×
Reverse FE	×	✓ ⁴	×	×	✓	×	✓	✓	✓	✓	✓	∞	✓	×	×
Slender	×	✓ ⁵	×	✓	×	×	×	✓	✓	✓	×	∞	×	✓	×
Converse	×	✓	×	✓	×	✓	✓	✓	✓	×	✓	∞	✓	×	×

¹ Including response expansion.

² Easy-to-model.

³ Robust against modeling.

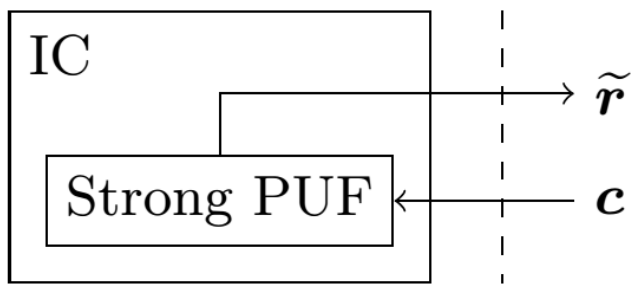
⁴ Non-determinism lower bound.

⁵ Both easy- and hard-to-model.

⁶ Reprogrammable.

Conclusion

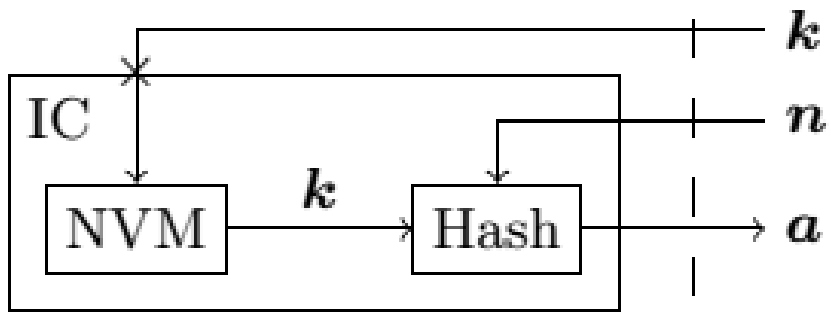
- Secure Lightweight Entity Authentication with Strong PUFs: **Mission Impossible?**
 - PUFs seem too brittle to be used without additional crypto (hash, ...).
 - PUF only one component in security architecture: still need TRNG, ...
 - Breakthrough: strong PUF with strong cryptographic properties (no machine learning)
- Not so very optimistic



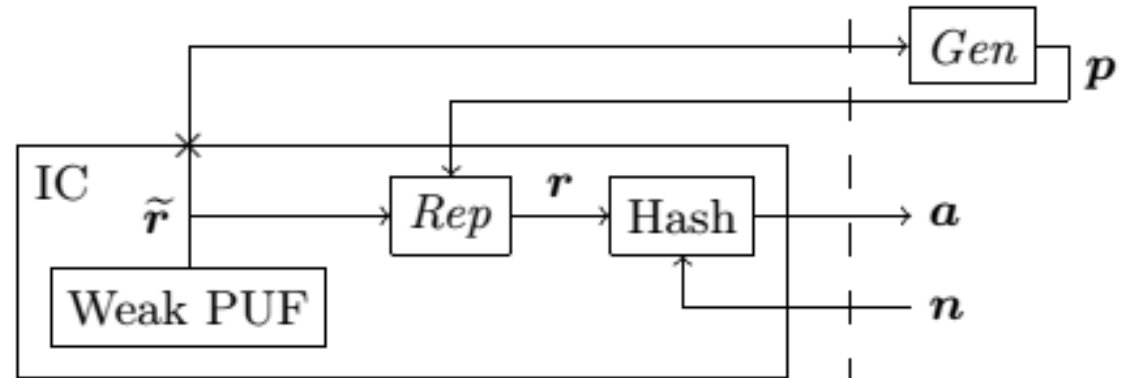
- Decade of research: no success. History of Machine Learning Attacks
- Infinity of output bits with limited number of circuit elements
- Most promising ideas (e.g. optics) are not lightweight
- Unavoidable trade-offs: security vs noise

- Coming soon: extended version on IACR Eprint (including 3 more protocols)

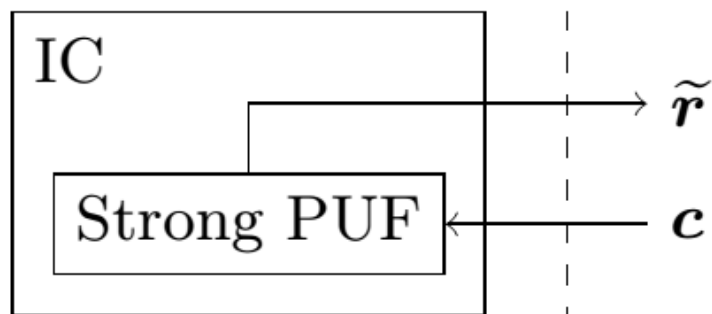
Appendix – Protocol Figures



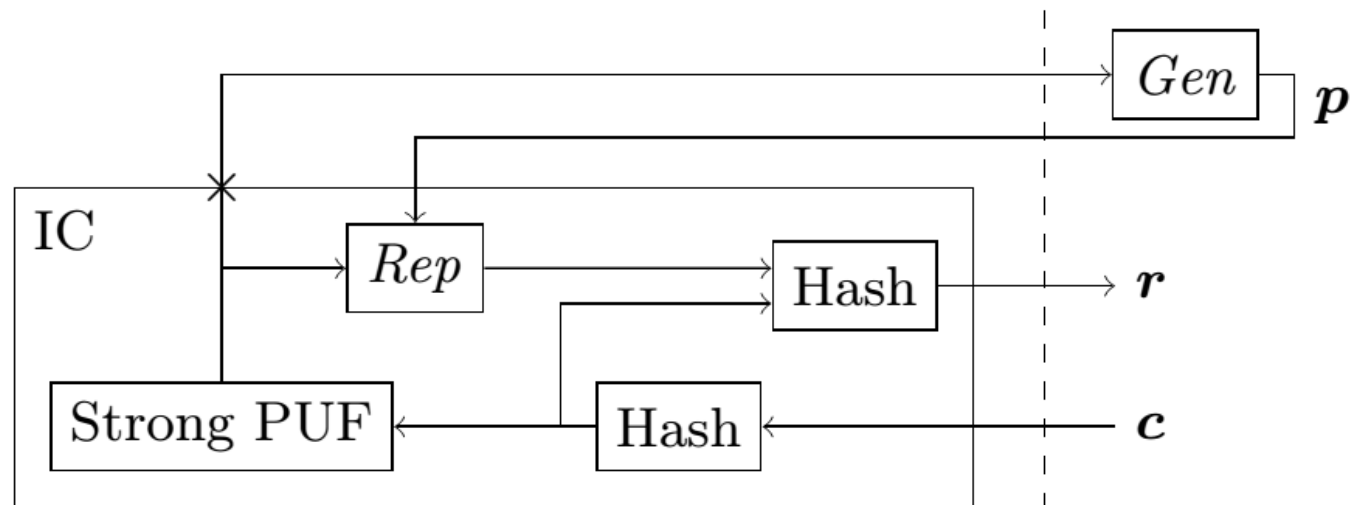
(a) Reference I: key in NVM.



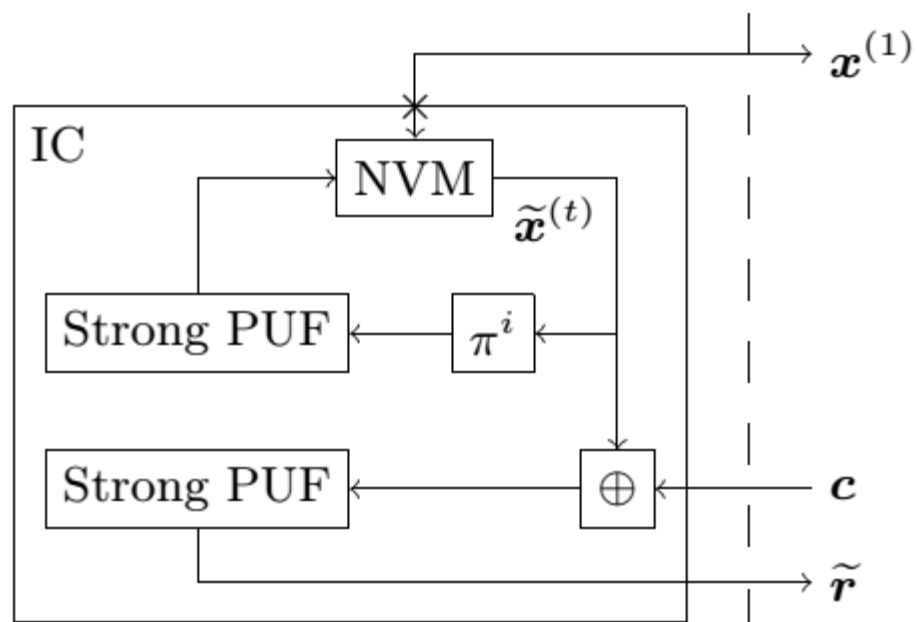
(b) Reference II: key via PUF and FE.



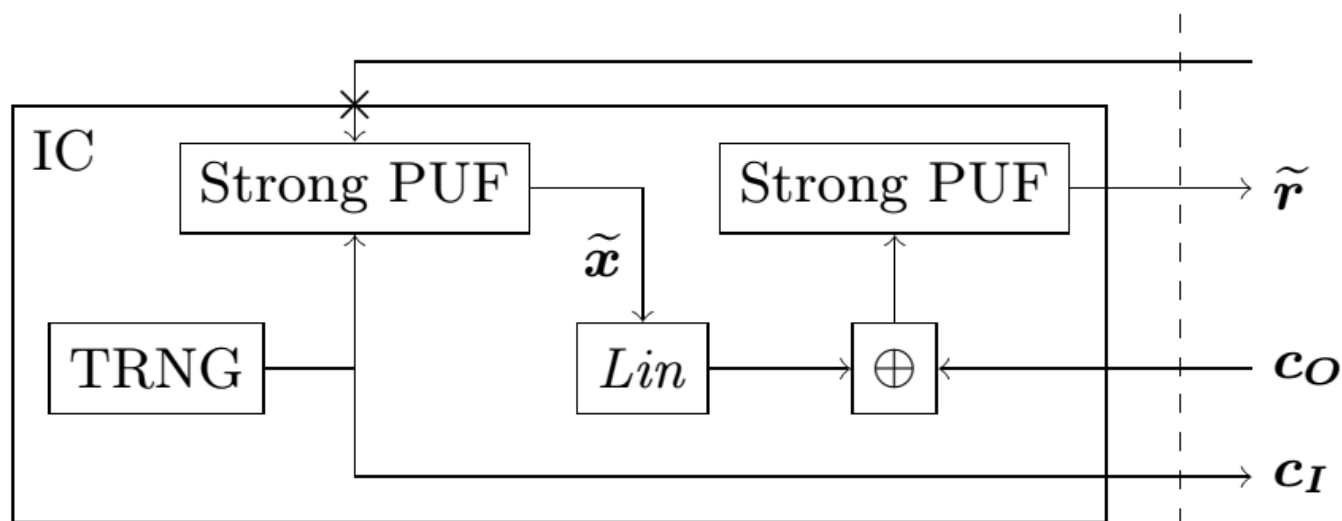
(c) Naive strong PUF [18].



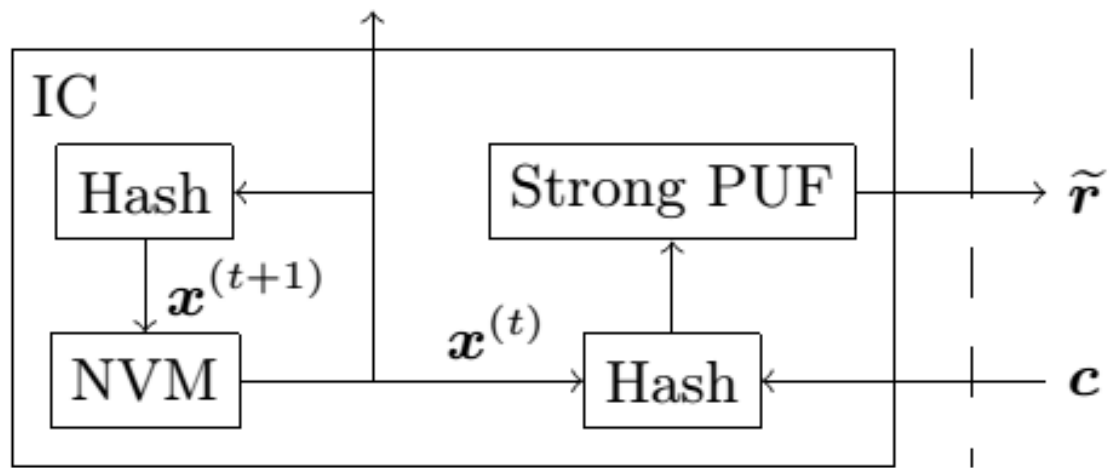
(d) Controlled PUF [4–6].



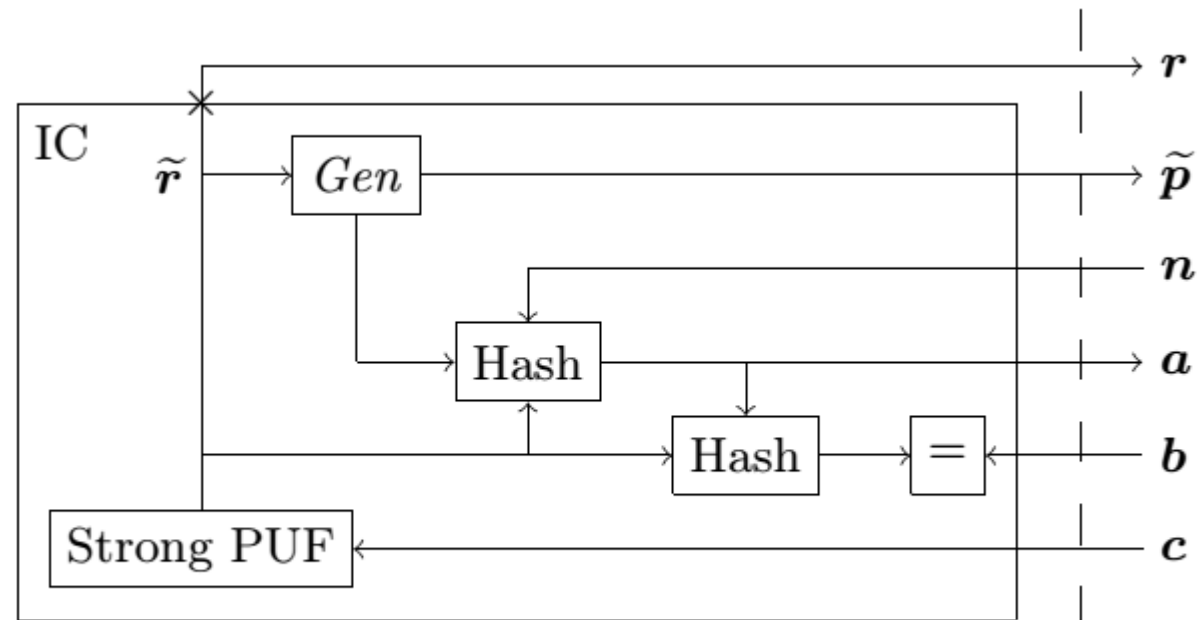
(e) Öztürk et al [17].



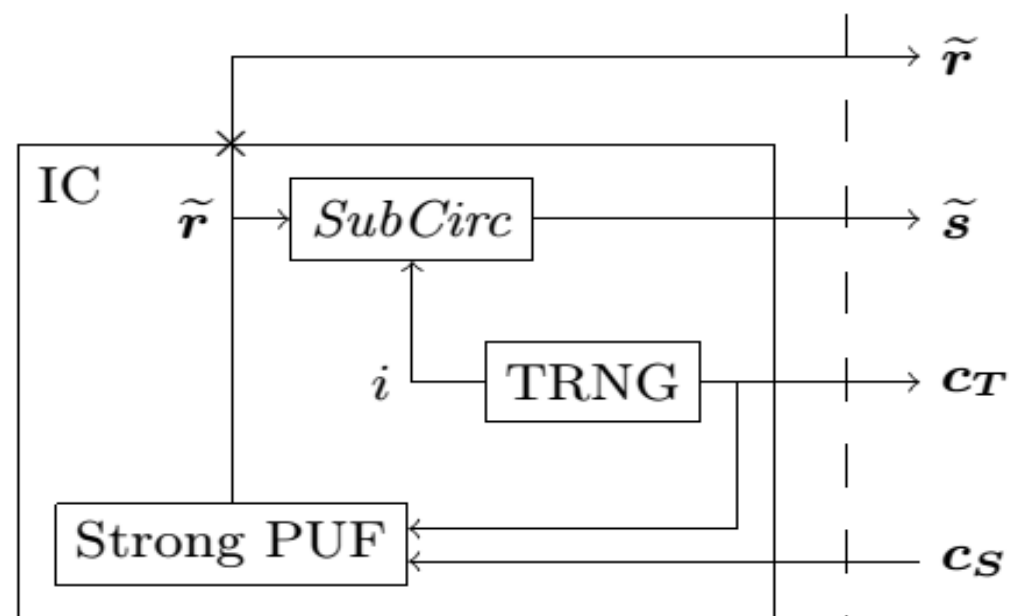
(f) Hammouri et al [8].



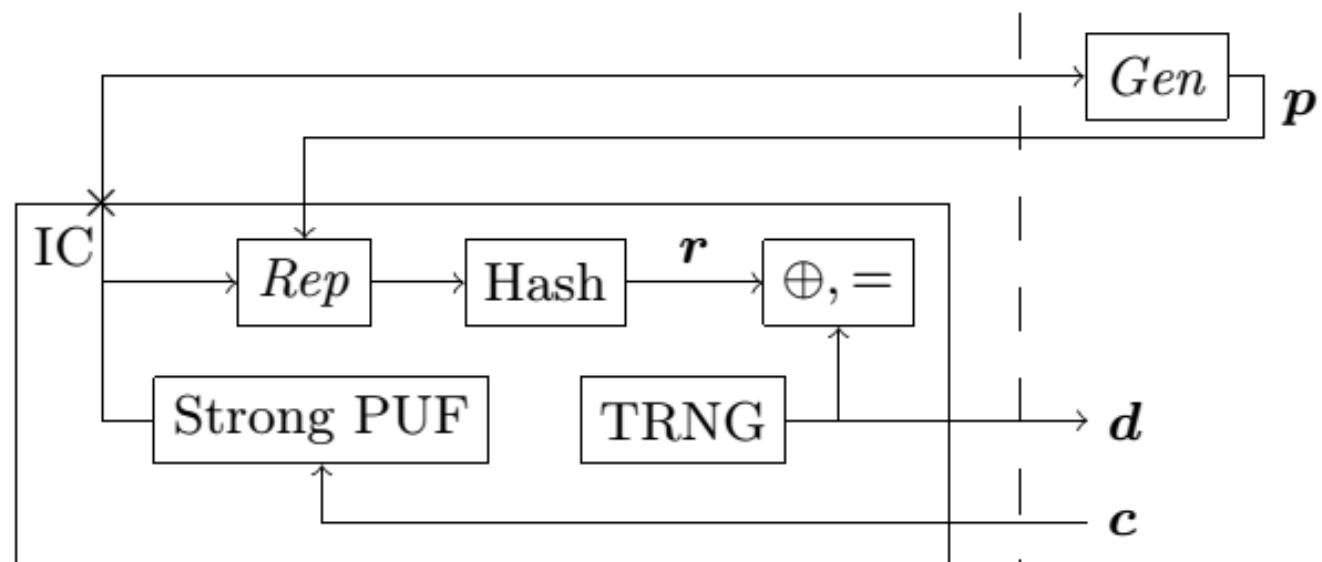
(g) Logically reconfigurable PUF [11].



(h) Reverse FE [24].



(i) Slender PUF [16].



(j) Converse [12].